

## Spectral partitioning:

Given  $L = V \Lambda V^T$

Find a clustering into  $k$  groups:

① Take:

$$V_k = \begin{bmatrix} \vec{v}_1 & \dots & \vec{v}_k \end{bmatrix} \xrightarrow{\text{row}} \vec{w}_i \text{ (row)}$$

$N \times k$        $\vec{w}_i$   $k \times 1$   
" "

② Take ~~the~~ rows of  $V_k(i):)$  as descriptors of each node.

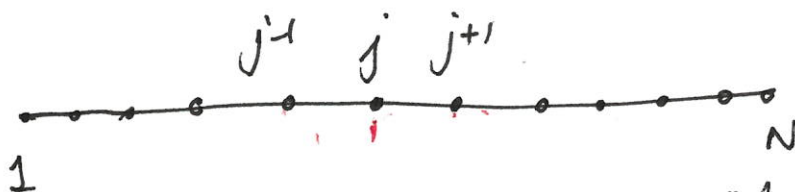
③ Carry out  $k$ -means on the vectors

$$\{\vec{w}_i\}_{i=1}^N \quad \vec{w}_0 \in \mathbb{R}^k$$

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# Connection of graphs with random walks.

Considers discretization of a 1D domain

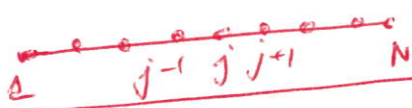


Discrete-time Markov chain

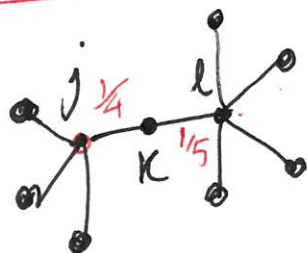
Defined in terms of:  $(\vec{p}_t)_{N \times 1} \equiv$  each coordinate containing the probability of the random walker at node  $i$  and time  $t$ .

$$\vec{1}^T \vec{p}_t = 1. \quad (\text{Normalized})$$

Random walk on:



$$p_{t+1}^{(j)} = \frac{1}{2} (p_t^{(j-1)} + p_t^{(j+1)})$$



$$p_{t+1}^{(k)} = \frac{1}{5} p_t^{(k)} + \frac{1}{4} p_t^{(j)}$$

This applies generally to any graph:

For a given graph with adjacency matrix  $A$  we have a Discrete-time Markov chain:

$$\vec{p}_{t+1} = \underbrace{(A\bar{D}^{-1})}_{\vec{M}} \vec{p}_t = \vec{M} \vec{p}_t$$

Transition matrix

Solution:  $\vec{p}_t = M^t \vec{p}_0$

Associated continuous-time process:

$$\frac{d\vec{p}(t)}{dt} = - \underbrace{(I - A\bar{D}^{-1})}_{\vec{L}_{RW}} \vec{p}(t)$$

$$L_{RW} D = L = D - A$$

$$\underline{L_{RW} = L D^{-1}}$$

And the isospectral (with  $L_{RW}$ )

Symmetrized normalized Laplacian:

$$\mathcal{L} = D^{-1/2} L_{RW} D^{1/2} = \bar{D}^{-1/2} L \bar{D}^{-1/2}$$

$$\underline{\mathcal{L} = \mathcal{L}^T} \quad \checkmark$$

See  
'Spectral  
clustering' section  
above.

$\Rightarrow$  Similar algorithms based on  
eigenvectors of  $\mathcal{L}$ .

(Ng & Jordan)

Note: They normalise by row too!

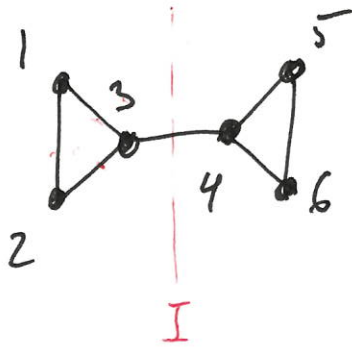
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② Alternative to spectral clustering:  
Modularity. (Newman)

Find a partition of the graph that  
will have maximally block diagonal  
structure in the adjacency matrix,  
and has more 'blocks' than  
expected at random.

Idea behind modularity: Count edges within blocks and between blocks

Example:



$$A = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 \end{bmatrix}$$

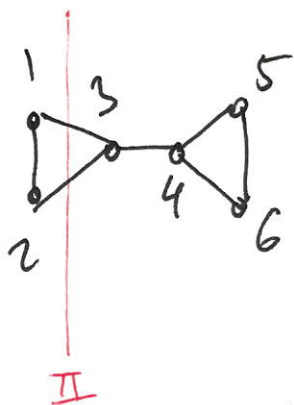
$$H_I = \begin{bmatrix} 1 & 0 \\ 1 & 0 \\ 1 & 0 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \end{bmatrix}_{N \times K}$$

Into  $\underline{c}$  groups

$$\underline{c=2}$$

$$\frac{1}{2} \left( \underset{I}{H}^T \underset{I}{A} \underset{c \times c}{H} \right) = \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \end{bmatrix} A \begin{bmatrix} 1 & 0 \\ 1 & 0 \\ 1 & 0 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \end{bmatrix} =$$

$$= \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 0 \\ 2 & 0 \\ 2 & 1 \\ 1 & 2 \\ 0 & 2 \\ 0 & 2 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 6 & 1 \\ 1 & 6 \end{bmatrix} = \begin{bmatrix} 3 & \frac{1}{2} \\ \frac{1}{2} & 3 \end{bmatrix}$$



Same A

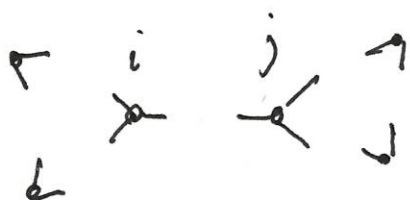
$$H_{II} = \begin{bmatrix} 1 & 0 \\ 1 & 0 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \\ 0 & 1 \end{bmatrix}$$

$$\frac{1}{2} (H_{II}^T A H_{II}) = \begin{bmatrix} 1 & 1 \\ 1 & 4 \end{bmatrix}$$

Maximise  
H       $\text{Tr} \left[ \frac{1}{2} (H^T A H) \right]$

Second ingredient : Null model.

Configuration model



$$\vec{d} = \begin{pmatrix} d_1 \\ \vdots \\ d_N \end{pmatrix} \quad 2E = \sum_{i=1}^N d_i$$

$$R_{ij} = \frac{d_i d_j}{2E}$$

Discount the edges that are expected to be present at random.



$$R = \frac{1}{2E} \vec{d} \vec{d}^T$$

(Expected edges if we rewired at random)

where  $\vec{d} = A\vec{1}$

Construct the modularity matrix:  $Z$

$$\frac{1}{2E} \text{Tr} \left[ H^T \left[ A - \frac{1}{2E} \vec{d} \vec{d}^T \right] H \right] = Q$$

and try to make it as modular as possible:

Modularity optimisation:

$$\max_H Q$$

How is  $Q$  optimised:

① Similarly to spectral clustering, it can be shown that we can use the leading eigenvector of  $Z$  to maximise  $Q$ .

Relaxation to  $\vec{S} \in \mathbb{R}^N$  etc

Then we can effect bipartitions in a recurrent manner until  $Q$  does not increase.

Stopping criterion:  $\Delta Q < 0$  in the iteration.

② In reality modularity is optimised through a greedy agglomerative algorithm that performs better than the spectral relaxation on the eigenvectors of  $Z$ .

This method is called Louvain optimisation and has become an industry standard.

It was proposed by Blondel et al in 2008.