

CMPE 352

Signal Processing & Algorithms

Spring 2019

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Review Questions (1)

- What are the two main constraints that must be taken into account for the choice of the sampling frequency f_s ?
 1. Constraint due to the signal spectrum (sampling theorem).
 2. Implementation also restricts the choice of f_s .
- How can the constraint due to the implementation be expressed?

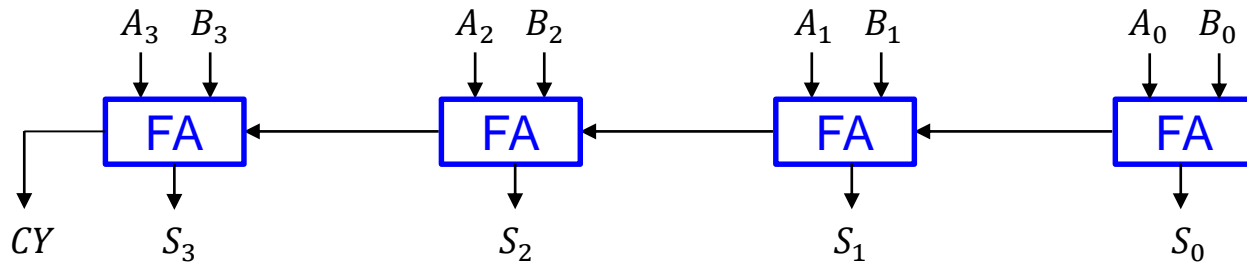
The sampling interval T_s must be larger than total processing or computation time required for each sample T_{proc} : $T_s > T_{proc}$

Note: *pipelining* can help to trade-off clock speed versus latency

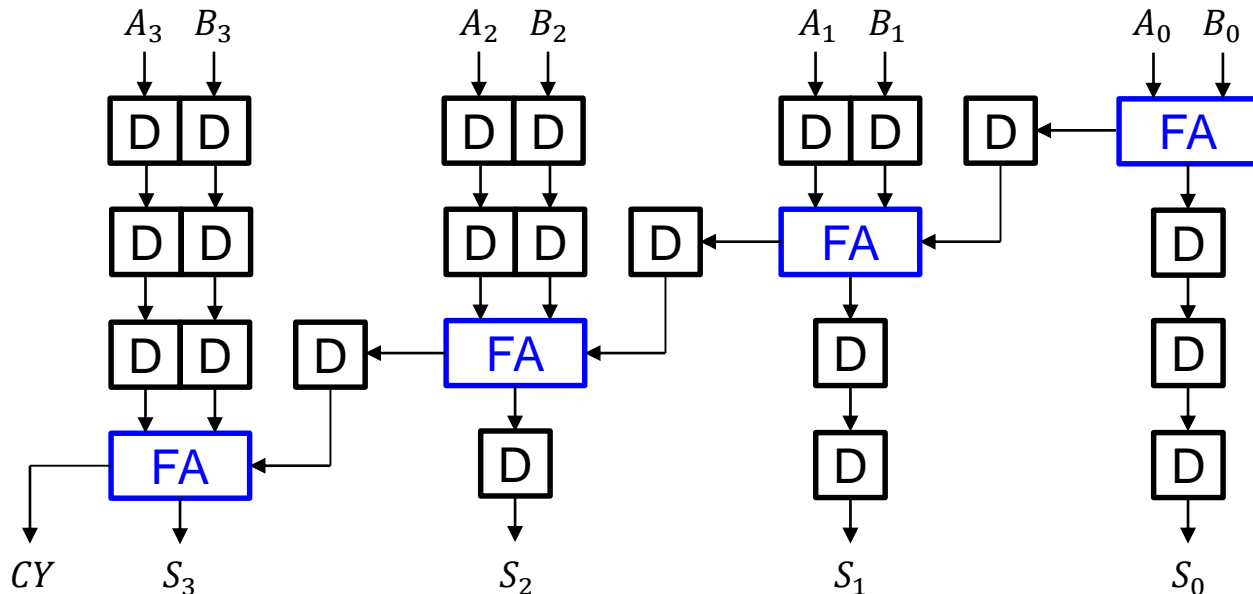
Note on Pipelining

Example: Adding two numbers.

Consider adding two 4-bit numbers: (A_3, A_2, A_1, A_0) and (B_3, B_2, B_1, B_0)



Direct Full-Adder
(FA) based
implementation



Implementation with
pipelining using
delay (D) elements
(D flip-flops)

Review Questions (2)

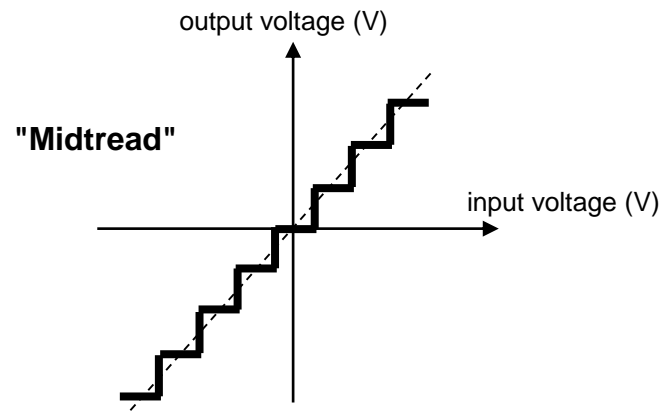
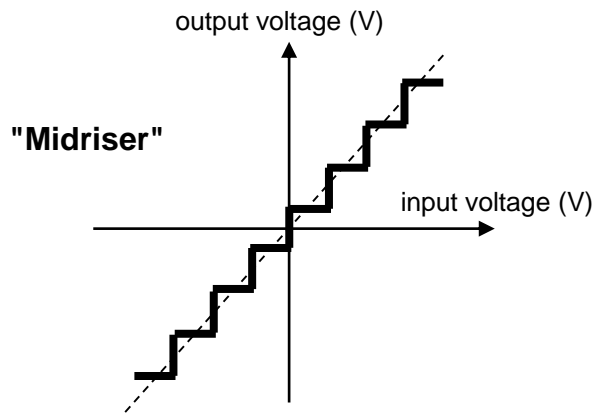
- What is the purpose of the quantization operation performed in association with sampling?

To be able to represent a signal sample with a finite number of bits or bytes.

- How is the quantization operation realized?

By dividing the signal variation range into a finite number of levels and rounding-off the current sample value to the nearest level.

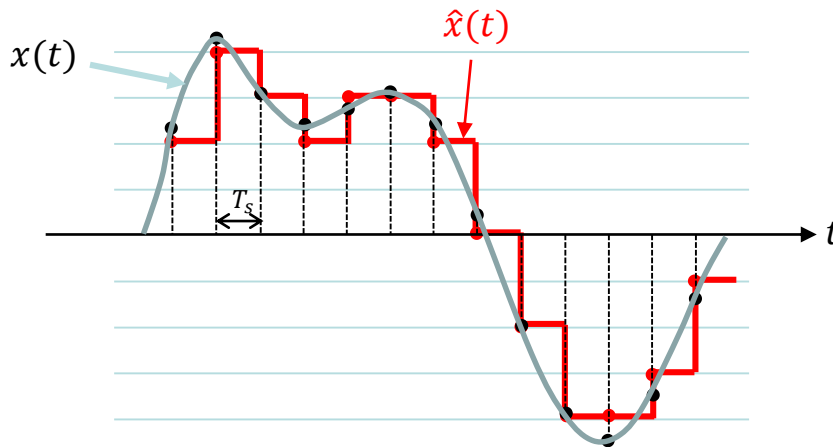
- What is the shape of the input-output characteristic of a quantizer?



Review Questions (3)

- What is a quantization-error signal?

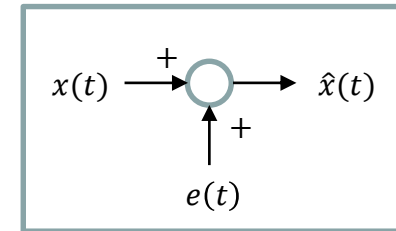
The difference between the quantized signal and the original unquantized analog signal.



$$e(t) = \hat{x}(t) - x(t)$$



$$\hat{x}(t) = x(t) + e(t)$$



Model of the Quantizer

- How did we define a quantizer model that expresses the quantization error?
- What is the power of the quantization error (= quantization noise)?
- What is the effect of increasing by 1 the number of quantizer bits?

$$N_q = \frac{\Delta v^2}{12} = \frac{x_p^2}{3L^2} \quad (\Delta v = \frac{2x_p}{L} : \text{quantization step})$$

To reduce N_q by 6 dB.

Review Questions (4)

- What is a PCM?

Pulse-code modulation (PCM) is a method used to digitally represent sampled analog signals.

- How is it realized?

The amplitude of the analog signal is sampled regularly at uniform intervals, and each sample is quantized to the nearest value, represented by N bits, within a range of digital steps.

- In PCM, what are the two basic parameters that determine the fidelity to the original analog signal?
 - the sampling rate
 - the bit depth N

Problem 1

A sinusoidal signal is transmitted using PCM. Find the minimum number of bits required to achieve a target signal-to-quantization-error ratio (SNR) of 25 dB.

Problem 2

Consider a sinusoidal signal given by $x(t) = 3 \sin(1000\pi t)$. Find the signal-to-quantization-error ratio (SNR) when the signal is quantized using a 9-bit PCM.

Problem 3

A sinusoidal TV signal with a frequency of 42 MHz is transmitted using binary PCM. The number of quantization levels is 1024. Calculate

- a) The word length N
- b) The average SNR
- c) The bit rate R

Problem 4

At what minimum frequency can the signal

$$x(t) = e^{-2|t|}$$

be sampled if we assume that the "essential bandwidth" B of the signal is the frequency at which $|X(j\omega)|$ drops to 1% of its peak value ?

The Discrete Fourier Transform (DFT) - 1

- The (continuous-time) Fourier transform and the inverse Fourier transform are given by

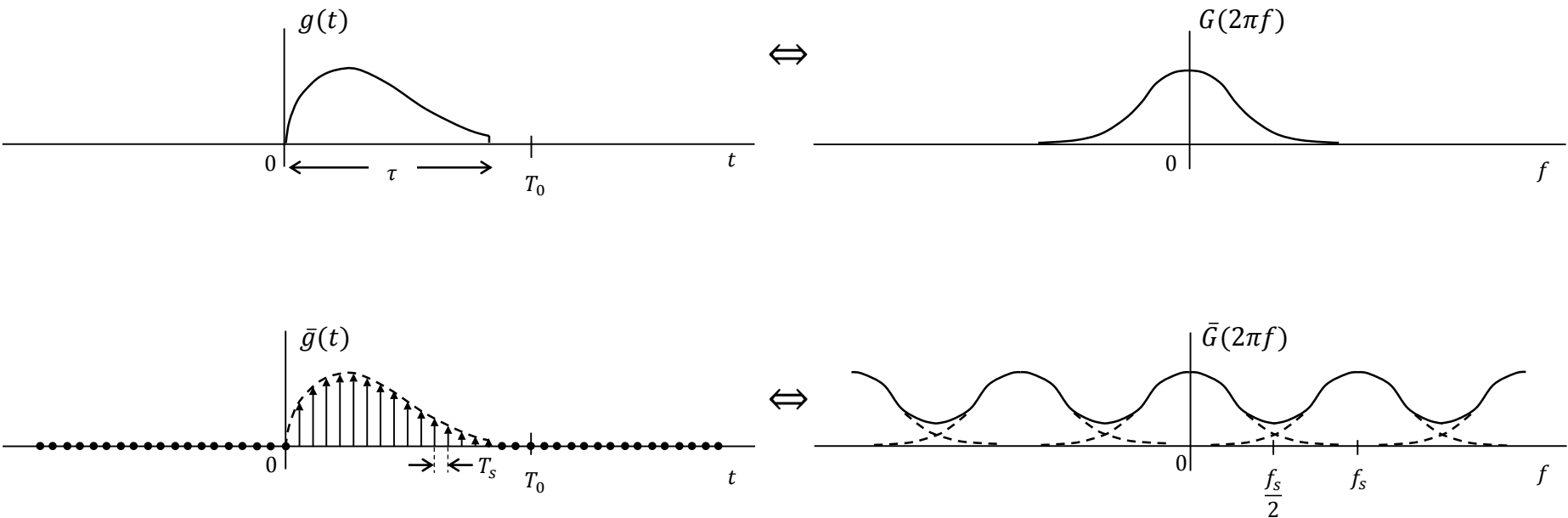
$$G(\omega) = \int_{-\infty}^{\infty} g(t) e^{-j\omega t} dt \qquad g(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} G(\omega) e^{j\omega t} d\omega$$

- We want to compute the (continuous-time) Fourier transform of $g(t)$ numerically.
- Hence we want to derive an **algorithm** to compute the Fourier transformation.
- This means that, practically, we have to use the samples of $g(t)$ and obtain the response (the Fourier transform) as samples of $G(\omega)$
 \Rightarrow we need to find the relationship between the samples of $g(t)$ and the samples of $G(\omega)$.
- Furthermore, for the algorithm to be practical, the data (i.e., the number of samples of $g(t)$ and $G(\omega)$) must be finite

The Discrete Fourier Transform (DFT) -2

- Let $g(t)$ be a signal starting at $t = 0$ and having duration τ . We will consider $g(t)$ on the interval $[0, T_0]$, where $T_0 \geq \tau$ (that is: $g(t) = 0$ for $\tau \leq t \leq T_0$, hence this makes no difference in the computation of $G(\omega)$).
- We sample $g(t)$ at intervals of T_s seconds: then there are a total of N samples:

$$N = \frac{T_0}{T_s}$$



The Discrete Fourier Transform (DFT) -3

- Let us make the approximation (for T_s sufficiently small)

$$G(\omega) = \int_0^{T_0} g(t) e^{-j\omega t} dt \cong \sum_{k=0}^{N-1} g(kT_s) e^{-j\omega kT_s} T_s$$

- and consider the samples of $G(\omega)$ at uniform intervals of ω_0 : $G_r = G(r\omega_0)$. We have therefore:

$$G_r = G(\omega)|_{\omega=r\omega_0} = \sum_{k=0}^{N-1} \underbrace{T_s g(kT_s)}_{g_k} e^{-jr\omega_0 T_s k} \Rightarrow \boxed{G_r = \sum_{k=0}^{N-1} g_k e^{-jr\Omega_0 k}}$$

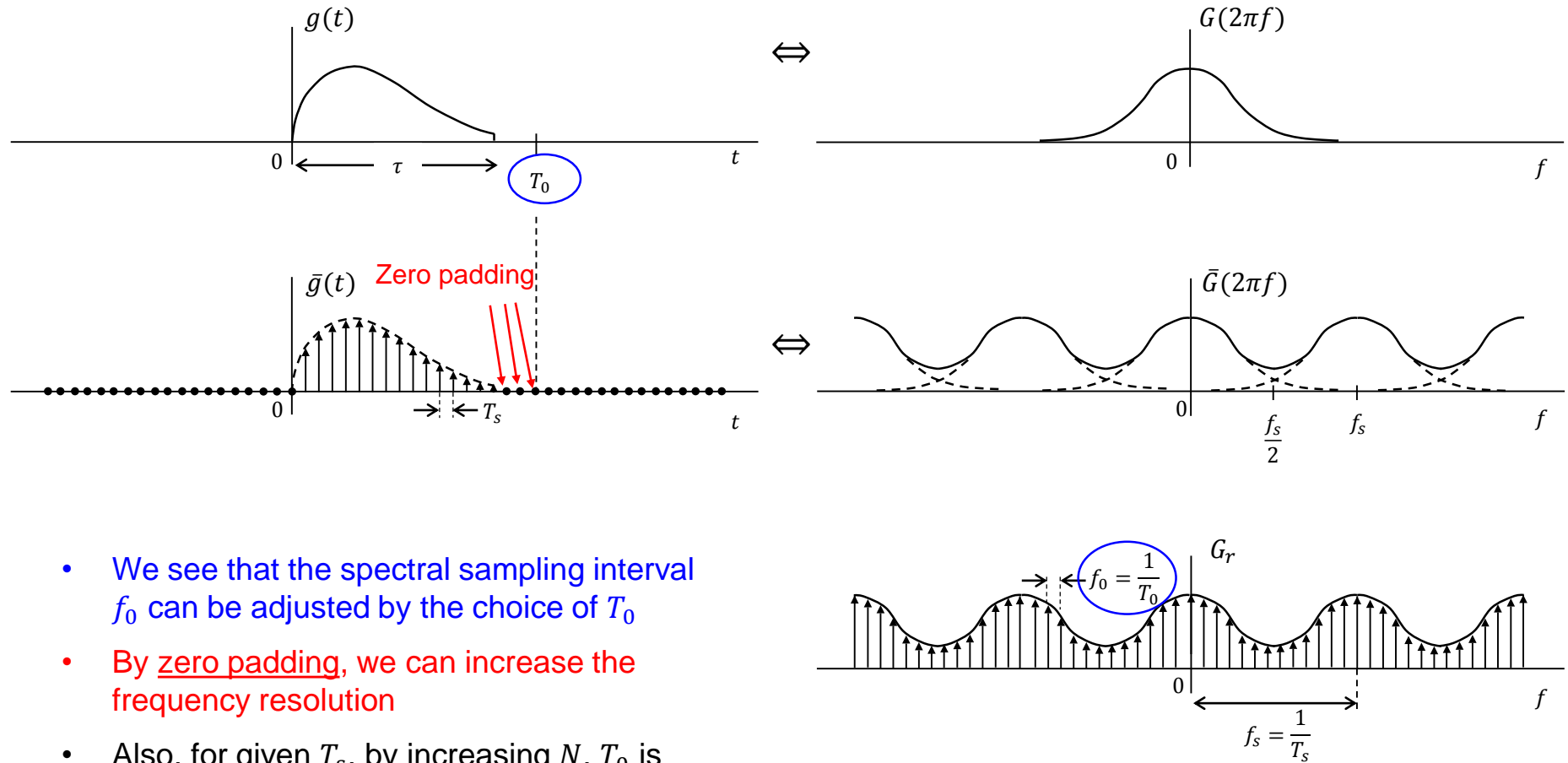
Relates g_k and G_r as desired

- G_r is periodic with period $\frac{2\pi}{\Omega_0}$: $G_{r+\frac{2\pi}{\Omega_0}} = G_r$
 \Rightarrow only $\frac{2\pi}{\Omega_0}$ samples G_r can be independent

- But G_r is also specified by N independent values of $g_k \Rightarrow$ hence $N = \frac{2\pi}{\Omega_0}$

$$N = \frac{2\pi}{\Omega_0} = \frac{2\pi}{\omega_0 T_s} = \frac{2\pi N}{\omega_0 T_0} \Rightarrow \boxed{\omega_0 = \frac{2\pi}{T_0} \text{ or: } f_0 = \frac{1}{T_0}}$$

Relationship Between Samples of $g(t)$ and $G(\omega)$



- We see that the spectral sampling interval f_0 can be adjusted by the choice of T_0
- By zero padding, we can increase the frequency resolution
- Also, for given T_s , by increasing N , T_0 is increased ($T_0 = NT_s$)

The Discrete Fourier Transform (DFT) - 4

- We now want to find the inverse relationship, that is, express g_k in terms of G_r

$$G_r = \sum_{k=0}^{N-1} g_k e^{-jr\Omega_0 k} \Rightarrow G_r e^{jm\Omega_0 r} = \sum_{k=0}^{N-1} g_k e^{-jr\Omega_0 k} e^{jm\Omega_0 r}$$

$$\sum_{r=0}^{N-1} G_r e^{jm\Omega_0 r} = \sum_{r=0}^{N-1} \sum_{k=0}^{N-1} g_k e^{-jr\Omega_0 k} e^{jm\Omega_0 r}$$

$$= \sum_{k=0}^{N-1} g_k \underbrace{\sum_{r=0}^{N-1} e^{j(m-k)\Omega_0 r}}_{(\Omega_0 = \frac{2\pi}{N})}$$

$$= \begin{cases} N & \text{for } m - k = 0, \pm N, \pm 2N, \dots \\ 0 & \text{otherwise} \end{cases}$$

$$= g_m N$$

$$\Rightarrow g_k = \frac{1}{N} \sum_{r=0}^{N-1} G_r e^{jk\Omega_0 r}$$

The Discrete Fourier Transform (DFT) - 5

$$g_k = \frac{1}{N} \sum_{r=0}^{N-1} G_r e^{jk\Omega_0 r}$$

- Note that $g_{k+N} = g_k$: hence g_k is periodic with period of N samples (because $N\Omega_0 = 2\pi$)
- Recall that G_r is also periodic with period of N samples

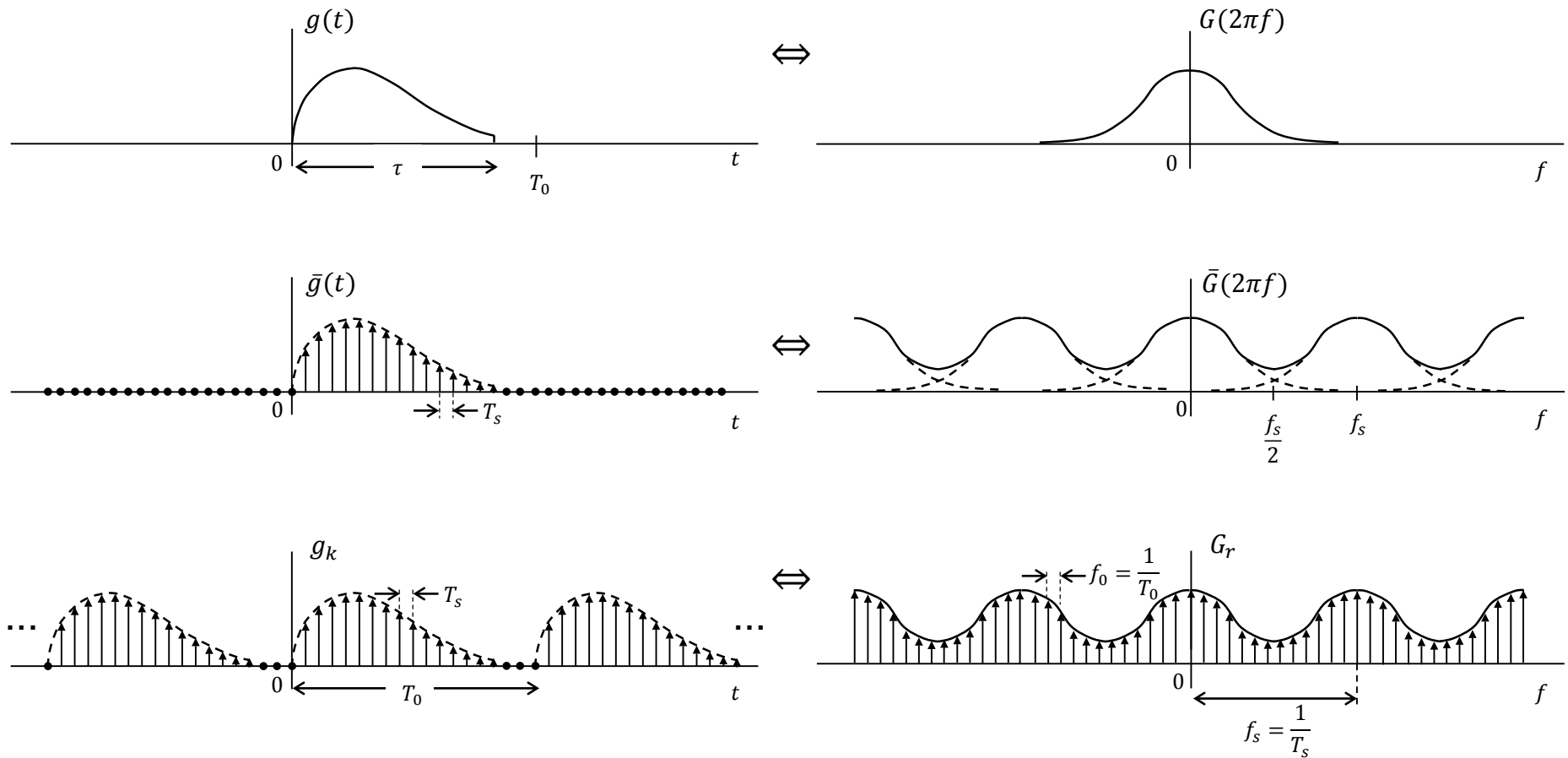


g_k ($k = 0, 1, \dots, N - 1$) and G_r ($r = 0, 1, \dots, N - 1$) are both periodic with period N

- G_r 's period N is equal to $\frac{1}{T_s}$:

$$N\omega_0 = \frac{T_0}{T_s} \frac{2\pi}{T_0} = \frac{2\pi}{T_s} = \omega_s \rightarrow \frac{1}{T_s} = f_s \text{ is } G_r \text{'s period}$$

Discrete Fourier Transform (DFT)



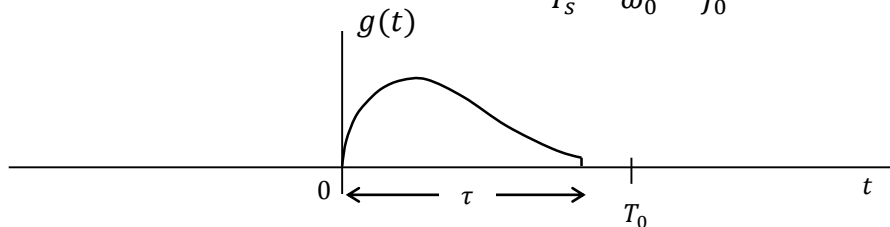
Discrete Fourier Transform (DFT) -- Summary

$$g_k = \frac{1}{N} \sum_{r=0}^{N-1} G_r e^{jk\Omega_0 r}$$

$$g_k = T_s g(kT_s)$$

$$\omega_0 = \frac{2\pi}{T_0} = 2\pi f_0$$

$$N = \frac{T_0}{T_s} = \frac{\omega_s}{\omega_0} = \frac{f_s}{f_0}$$

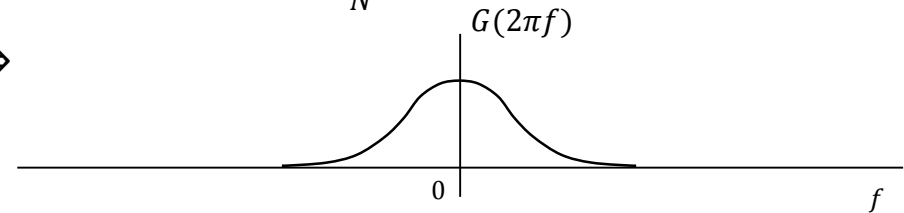
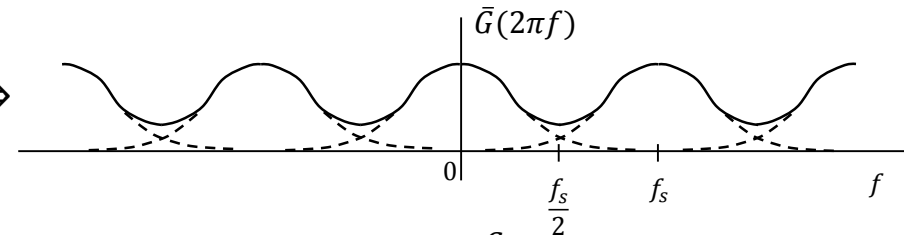
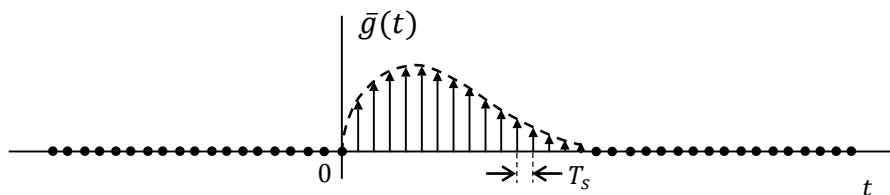

 \Leftrightarrow

$$G_r = \sum_{k=0}^{N-1} g_k e^{-jr\Omega_0 k}$$

$$G_r = G(r\omega_0)$$

$$\omega_s = \frac{2\pi}{T_s} = 2\pi f_s$$

$$\Omega_0 = \omega_0 T_s = \frac{2\pi}{N}$$


 \Leftrightarrow

 \Leftrightarrow
