

Relational AlgebraOperators

SELECT

symbol = σ (sigma)syntax = σ (R)

<selection condition>

Where: selection condition is a boolean expression on the attributes of R.

R = relational algebra expression

EMPLOYEE

attributes.

| e-id | e-name | e-surname | salary | → headers tuple |
|------|--------|-----------|--------|--------------------|
| e001 | Ece | Ece | 50 | |
| e050 | Oya | Ece | 20 | |
| e100 | Ali | Ece | 70 | |
| e2 | Ali | Omar | 20 | |

θ_1 = Use Relational Algebra to select all employees with salary > 30. out of EMPLOYEE relation.

$R \leftarrow \sigma_{\langle \text{salary} > 30 \rangle} (\text{EMPLOYEE})$

$R =$

| e-id | e-name | e-surname | salary |
|------|--------|-----------|--------|
| e001 | Ece | Ece | 50 |
| e100 | Ali | Ece | 70 |

θ_2 = Use R.A. to select all emp. having Name = Ece and salary > 20

$R \leftarrow \sigma_{\langle \text{e-name} = \text{Ece} \rangle \text{ AND } \langle \text{salary} > 20 \rangle} (\text{EMPLOYEE})$

$R \leftarrow \sigma_{\langle \text{e-name} = \text{ece} \rangle} \left(\sigma_{\langle \text{salary} > 20 \rangle} (\text{EMPLOYEE}) \right)$

Def. = Cardinality = number of Tuple
Def. = Degree = number of attribute

Example:

degree (EMPLOYEE) = 4

SELECT op. does not change the degree.

cord (EMPLOYEE) = 4 (tuples, rows)

cord (R₁) = 2

cord (R₂) = 1

PROJECT

symbol = π

syntax = π (R)
 | list attribute

Example: Use R.A. To select $\langle e_id, salary \rangle$
att. out of EMPLOYEE.

$R_3 \leftarrow \pi_{e_id, salary} (EMPLOYEE)$

degree = 2
cord = 4

$R_3 =$

| <u>e_id</u> | salary |
|-------------|--------|
| e001 | 50 |
| e050 | 20 |
| e100 | 70 |
| e2 | 20 |

Example: Use R.A. to project EMPLOYEE over salary.

$R_4 \leftarrow \pi_{\underline{salary}} (EMPLOYEE)$

degree = 1
cord = 3

$R_4 =$

| <u>salary</u> |
|---------------|
| 50 |
| 20 |
| 70 |

UNION, INTERSECTION and MINUS op.Syntax: $R \cup S, R \cap S, R - S$ R and S must be union compatible:

They must have same number of attribute

$$\left. \begin{array}{l} R(A_1, A_2, \dots, A_n) \\ S(B_1, B_2, \dots, B_n) \end{array} \right\} \begin{array}{l} n = m \\ \text{dom}(A_i) = \text{dom}(B_i) \\ \forall i = 1, \dots, n \end{array}$$

Example: Retrieve The SSN of all employees who either work in DNO5 OR directly supervise an employee from DNO5 (department number 5)

DEPT5-EMP \leftarrow 6 (EMPLOYEE)

degree (DEPT5-EMP) = 10

card (DEPT5-EMP) = 4

SSN-DEPT5-EMP \leftarrow $\pi_{\text{SSN}}(\text{DEPT5-EMP})$

degree = 5?

card = 4

SSN-SUPERVISE \leftarrow $\pi_{\text{SUPER-SSN}}(\text{DEPT5-EMP})$

degree = 1

card = 2

 $R = \text{SSN-DEPT5-EMP} \cup \text{SSN-SUPERVISE}$ • Cartesian ProductSyntax = $R \times S$

Semantic:

$$\theta = R(A_1, A_2, \dots, A_n) \times S(B_1, B_2, \dots, B_m)$$

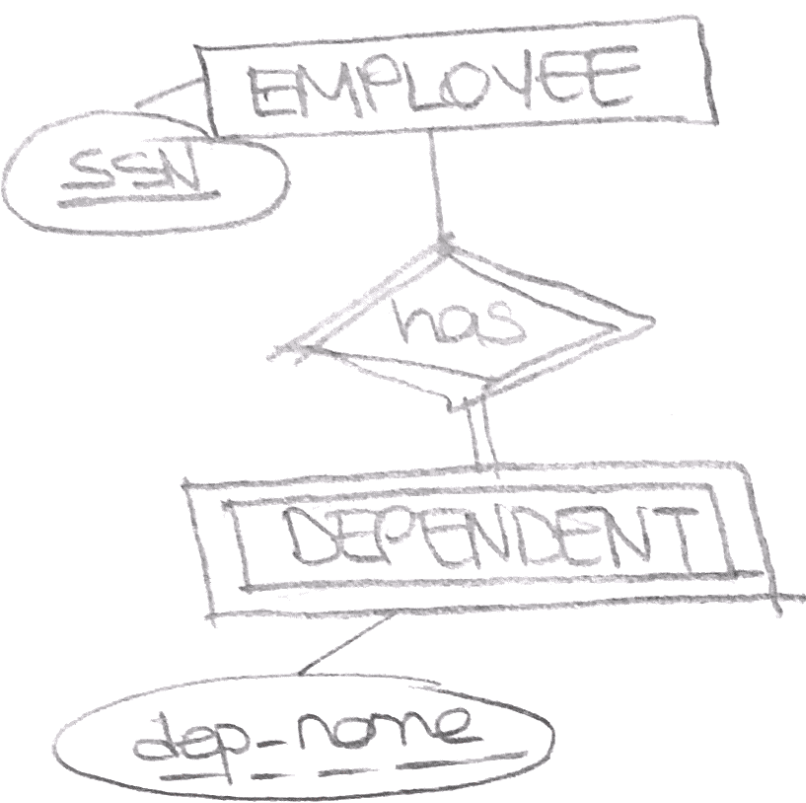
$$\theta = (A_1, A_2, \dots, A_n, B_1, B_2, \dots, B_m)$$

$$\text{degree}(\theta) = \text{degree}(R) + \text{degree}(S)$$

If $|R| = n_r$ and $|S| = n_s$ then

$$|\theta| = n_r \times n_s$$

Example: Use RA to retrieve the name of the dependent of each F employee together with "FNAME" and "LNAME" of employee.


$$FEM_DEP \leftarrow \pi(6(\text{DEPENDENT}))$$

ESSN, DEPENDENTNAME

degree = 2

$$\text{cond} = 4$$

FEM-DEP

| ESSN | DEP. Name |
|------|-----------|
| 333 | Alice |
| 333 | Joy |
| 123 | Alice |
| 123 | Elizabeth |

Temp ← FEM_DEP X EMPLOYEE

$$\text{degree} = 10 + 2$$
$$\text{count} = 8 \times 4$$

| | | |
|------|------|--|
| Temp | ESSN | |
| | 333 | |
| | | |