

Computer Architecture

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Sequential Logic

Combinatorial Logic

- So far we ignored the issue of time
- The inputs were just “there” – fixed and unchanging
- The output was just a function of the input
 - Not of anything that happened “previously”
- The output was computed “instantaneously”
- This is sometimes called “Combinatorial Logic”

Sequential Logic

Hello, Time

- Use the same hardware over time
 - Inputs change and outputs should follow

□ E.g.

```
For i = 1 ... 100:  
  a[i]=b[i]+c[i]
```

- Remember “State”

- Memory
- Counters

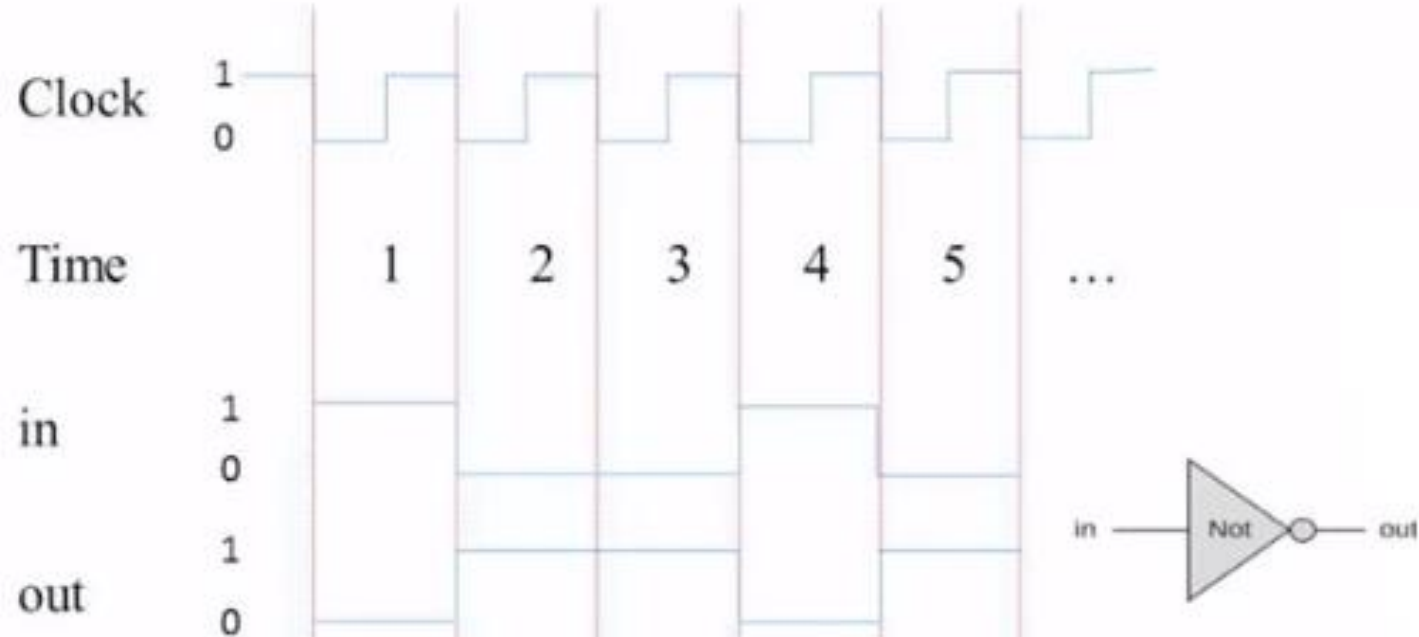
□ E.g.

```
For i = 1 ... 100:  
  sum = sum + i
```

- Deal with speed

Sequential Logic

The Clock



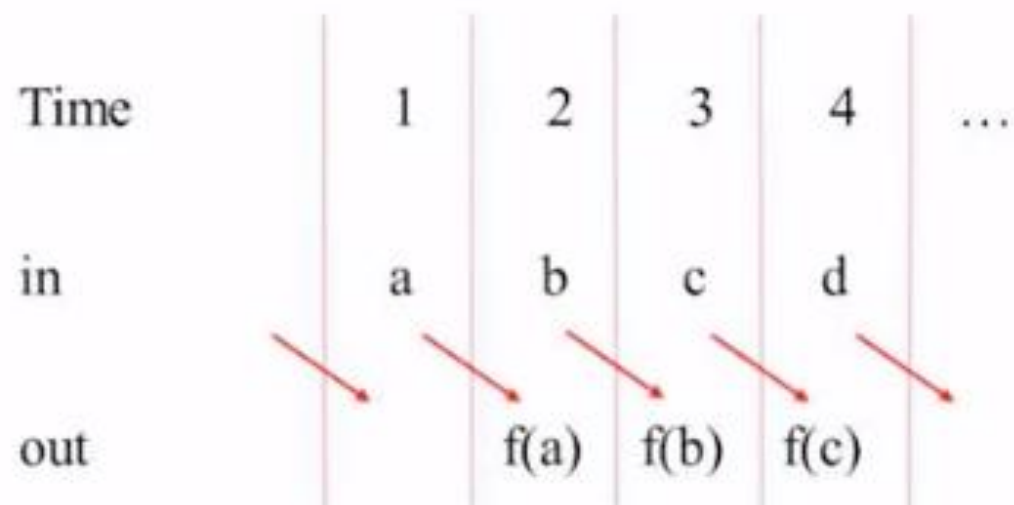
Combinatorial Logic vs. Sequential Logic

- Combinatorial: $out[t] = function(in[t])$
- Sequential: $out[t] = function(in[t-1])$

Time	1	2	3	4	...
in	a	b	c	d	
	↓	↓	↓	↓	
out	f(a)	f(b)	f(c)	f(d)	

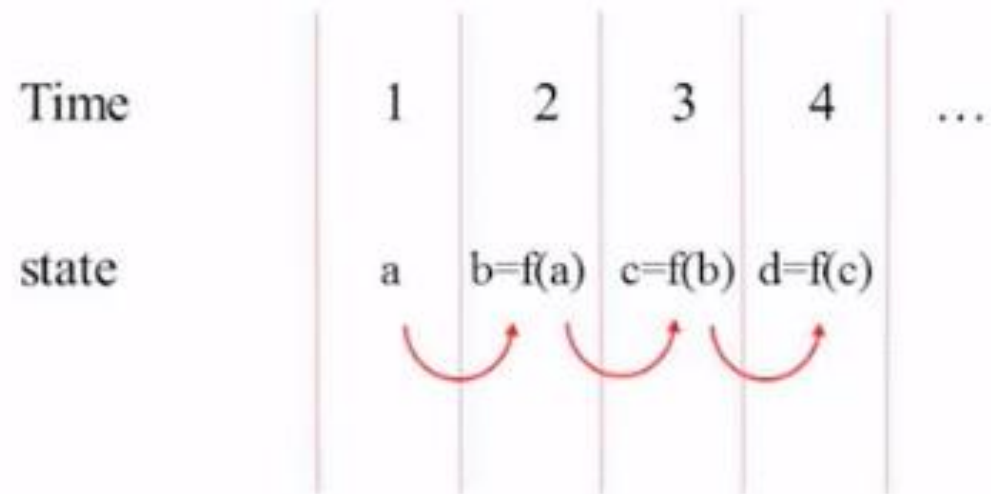
Combinatorial Logic vs. Sequential Logic

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- Sequential: $out[t] = function(in[t-1])$



Combinatorial Logic vs. Sequential Logic

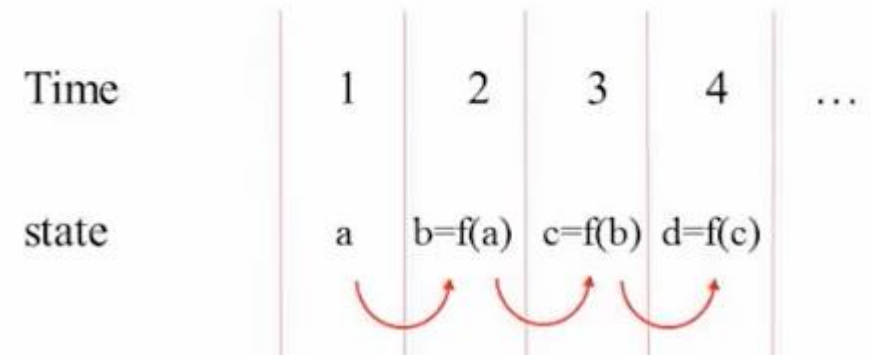
- Combinatorial: $out[t] = function(in[t])$
- Sequential: $state[t] = function(state[t-1])$



Remembering State

- Missing ingredient: remember one bit of information from time $t-1$ so it can be used at time t .
- At the “end of time” $t-1$, such an ingredient can be at either of two states: “remembering 0” or “remembering 1”.
- This ingredient remembers by “flipping” between these possible states.
- Gates that can flip between two states are called Flip-Flops.

$$state[t] = function(state[t-1])$$



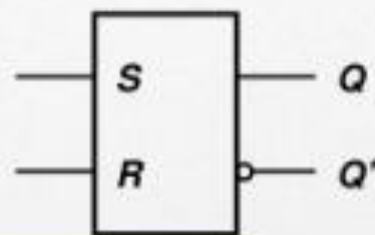
Flip-Flop

- "Flip-flop" is the common name given to two-state devices **which offer basic memory for sequential logic operations.**
- Flip-flops are heavily used for digital data storage and transfer and are commonly used in banks called "registers" for the **storage of binary numerical data.**

Flip-Flop

- Flip-flop are basic storage/memory elements.
- Flip-flop are essentially 1-bit storage devices.
- Types of flip-flops are:
 - 1. SR Flip-flop
 - 2. JK Flip-flop
 - 3. D Flip-flop
 - 4. T Flip-flop
- Application of flip-flop:
 - 1. Counter
 - 2. Register
 - 3. Memory
 - 4. Logic controller
 - 5. Frequency Divider

SR Flip-Flop



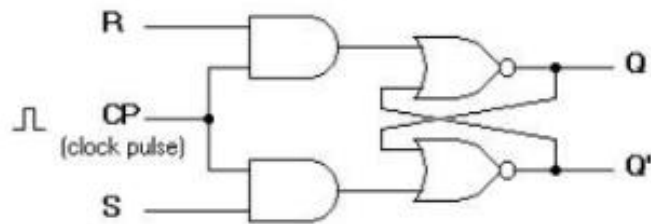
- The **simplest** binary storage device.
- SR Flip-flop have **2 inputs (SET & RESET)** and **2 outputs (Q & Q')**.

NOTE: Q & Q' are compliments of each other

- The SR flip flop is sometimes referred to as an **SR latch**. The Term latch refers to its use as a **temporary memory storage device**.

Flip Flops- SR Flip flop (set reset)

- If S and R is 0, $Q_{next}=Q$; $Q(t+1)=Q(t)$



(a) Logic diagram

Q	S	R	Q(t+1)
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	indeterminate
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	indeterminate

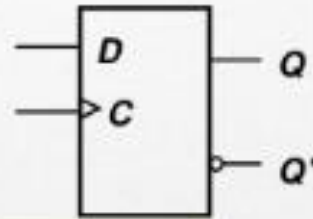
(b) Truth table

Clocked SR flip-flop

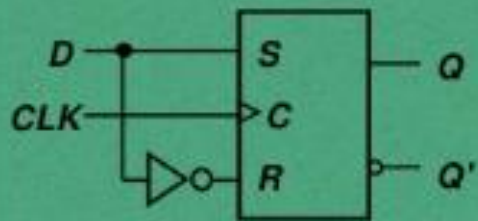
S	R	Q	Q_{next}	Q'_{next}
0	0	0	0	1
0	0	1	1	0
0	1	x	0	1
1	0	x	1	0
1	1	x	x	x

D-flip flop

D Flip-Flop Truth Table:



- **D flip-flop:** single input D (data)
- D=HIGH a **SET** state
- D=LOW a **RESET** state
- Q follows D at the clock edge.
- D flip-flop formed by add **NOT** gate between SR input.

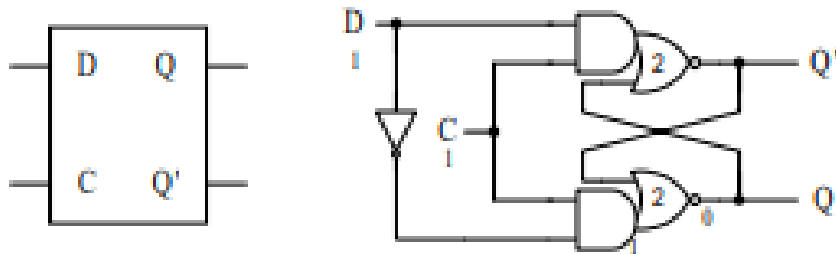


D	CLK	Q(t+1)	Comments
1	↑	1	Set
0	↑	0	Reset

↑ = clock transition LOW to HIGH

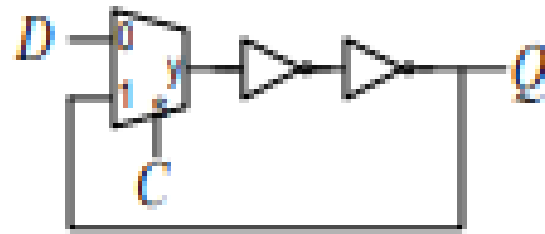
Gated D- flip flop

- when the C input is 0, the Q output remains same.
- The latch "closes", the Q output retains 'ts last value and no longer change 'n response to D
- when the C input is 1, the Q output follows the D input.
- At this stage, the latch is said to be "open" and the path from D input to Q output is "transparent"



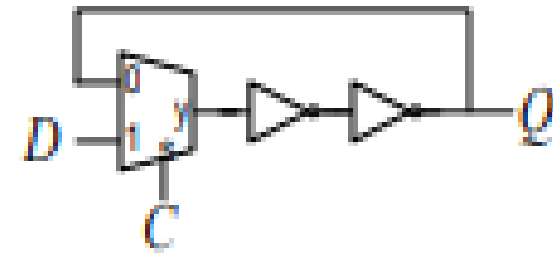
C	D	Q	Q _{next}
0	x	0	0
0	x	1	1
1	0	x	0
1	1	x	1

Gated D latch can also be implemented using a multiplexer.



negative latch

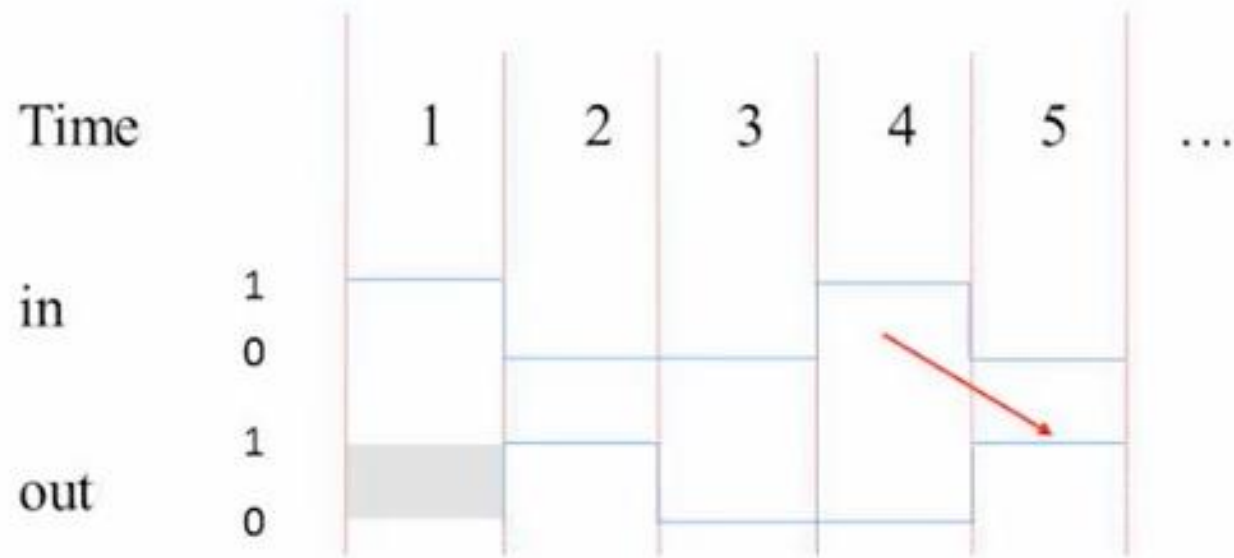
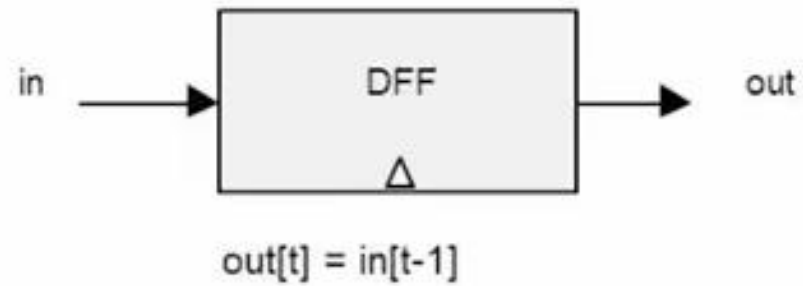
D passes to Q when $C=0$



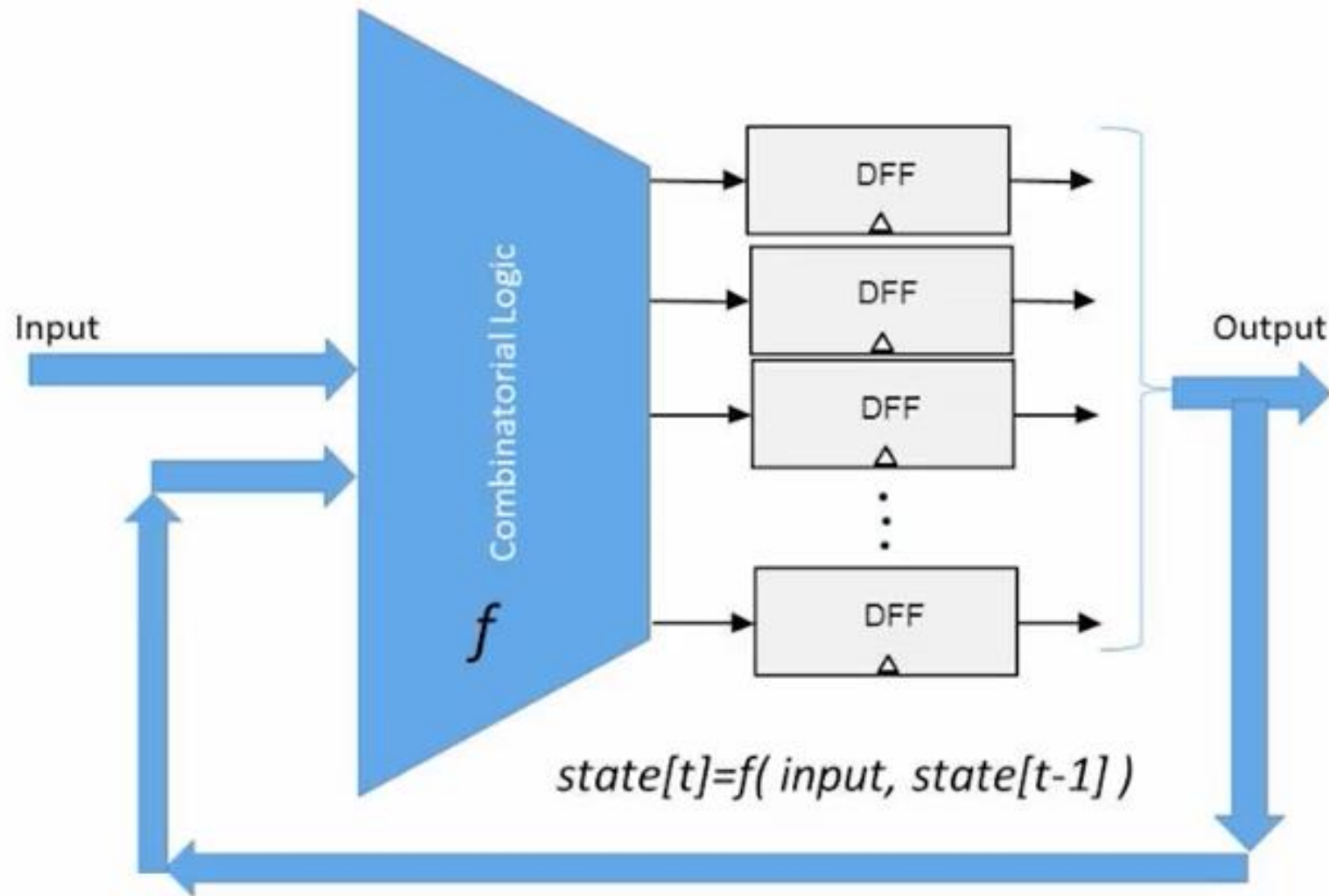
positive latch

D passes to Q when $C=1$

Flip Flops

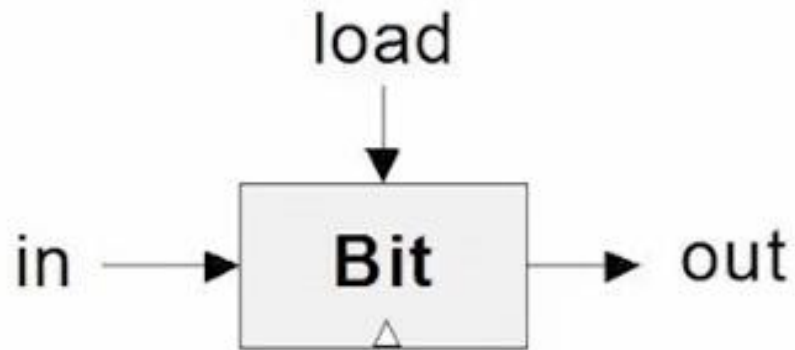


Sequential Logic Implementation



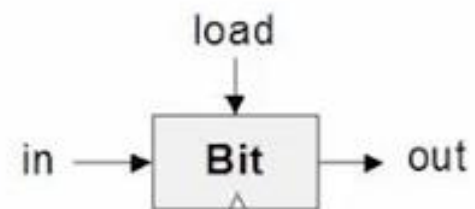
Remembering For Ever: 1-bit Register

- Goal: remember an input bit “forever”: until requested to load a new value.

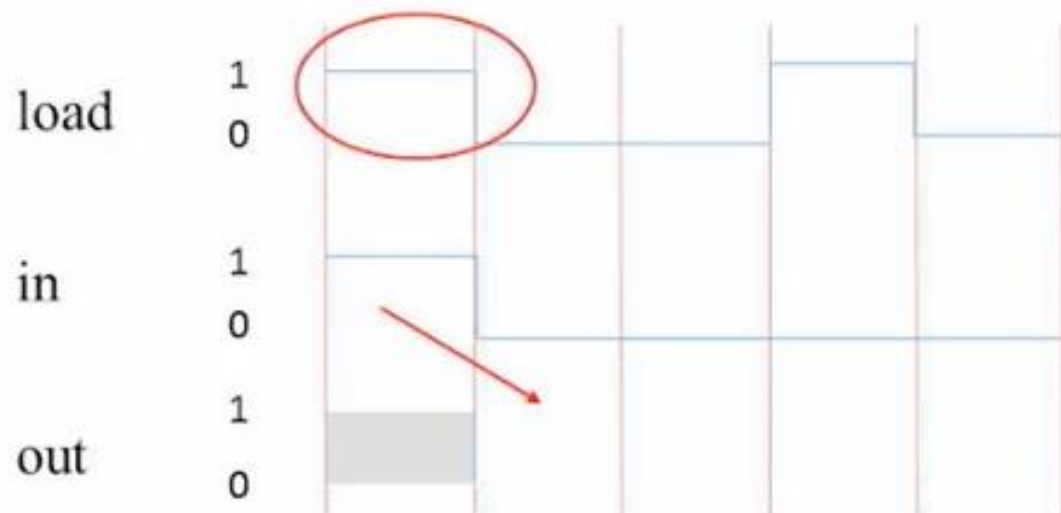


if $\text{load}(t-1)$ then $\text{out}(t) = \text{in}(t-1)$
else $\text{out}(t) = \text{out}(t-1)$

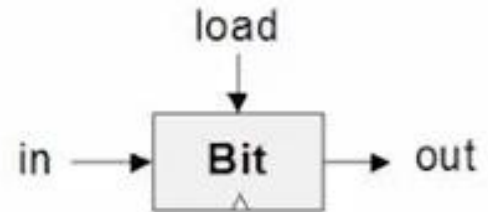
1-Bit Register



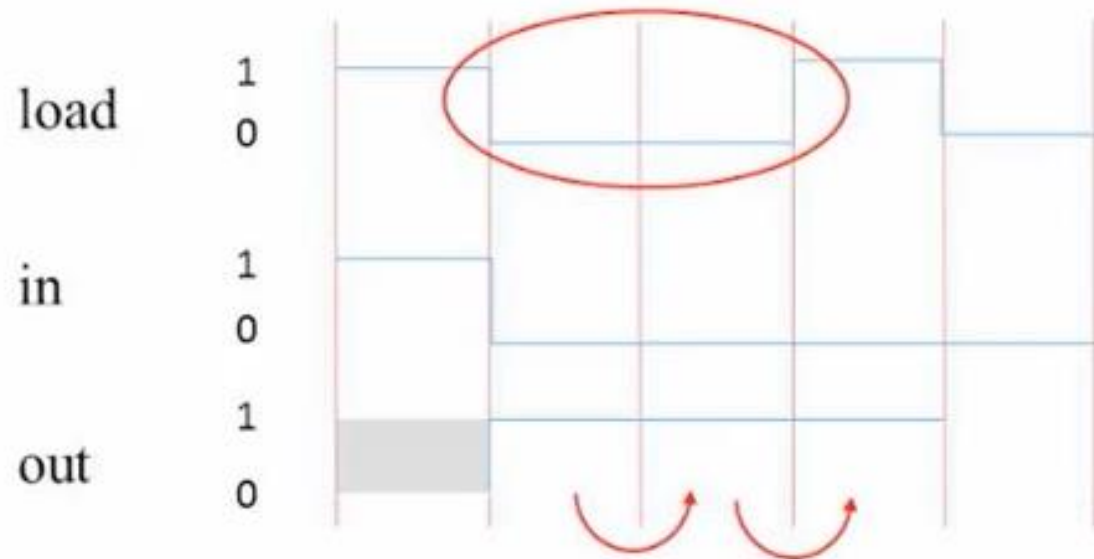
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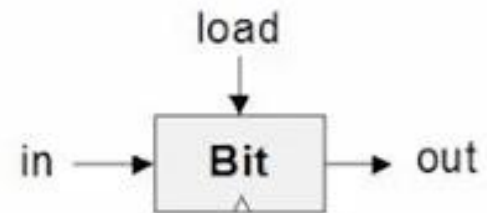
1-Bit Register



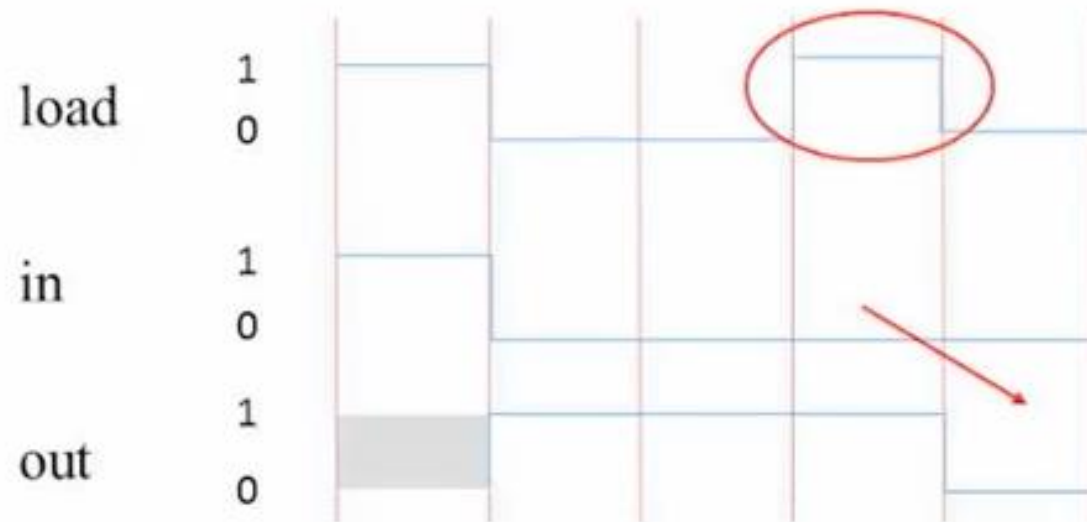
if load(t-1) then out(t)=in(t-1)
else out(t)=out(t-1)

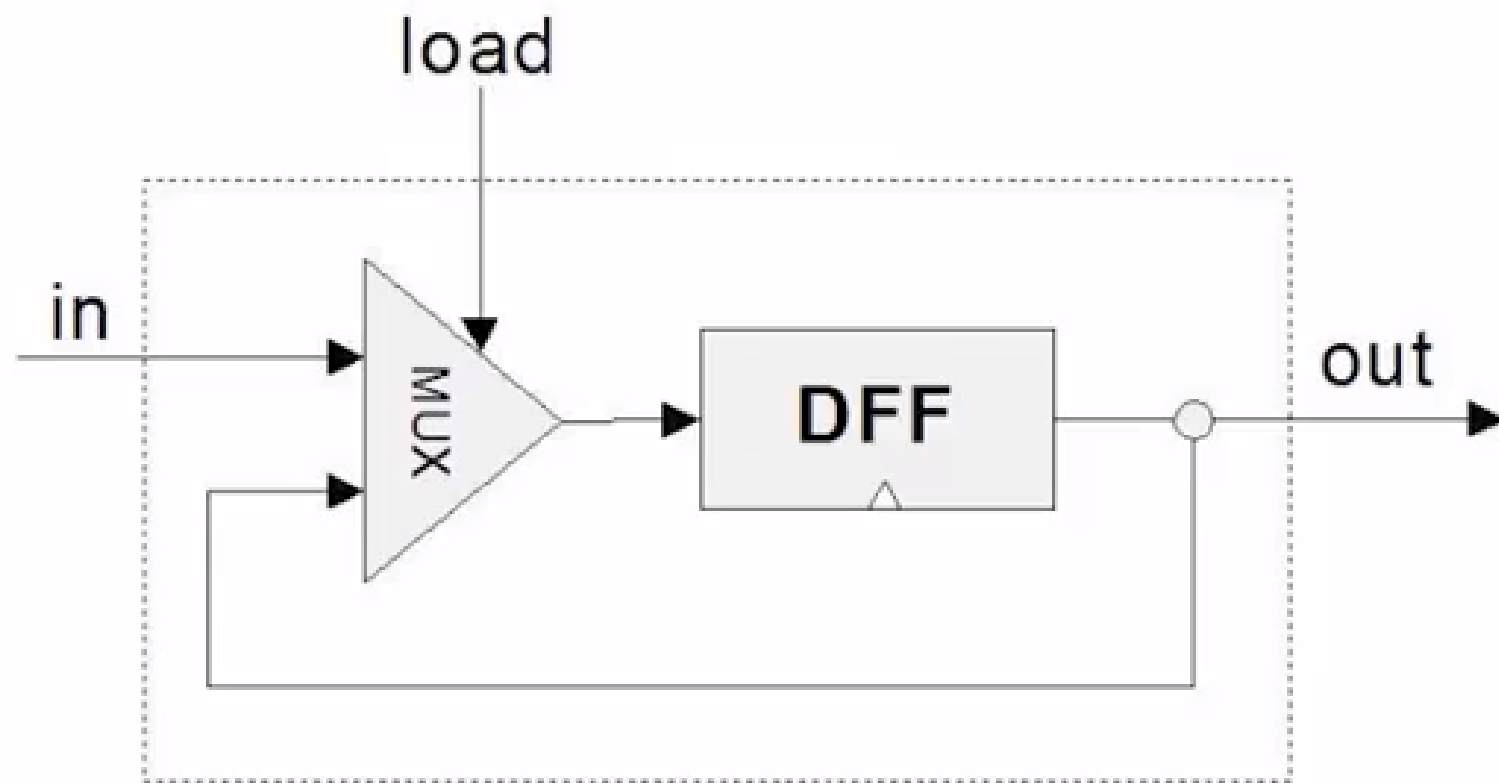


1-Bit Register

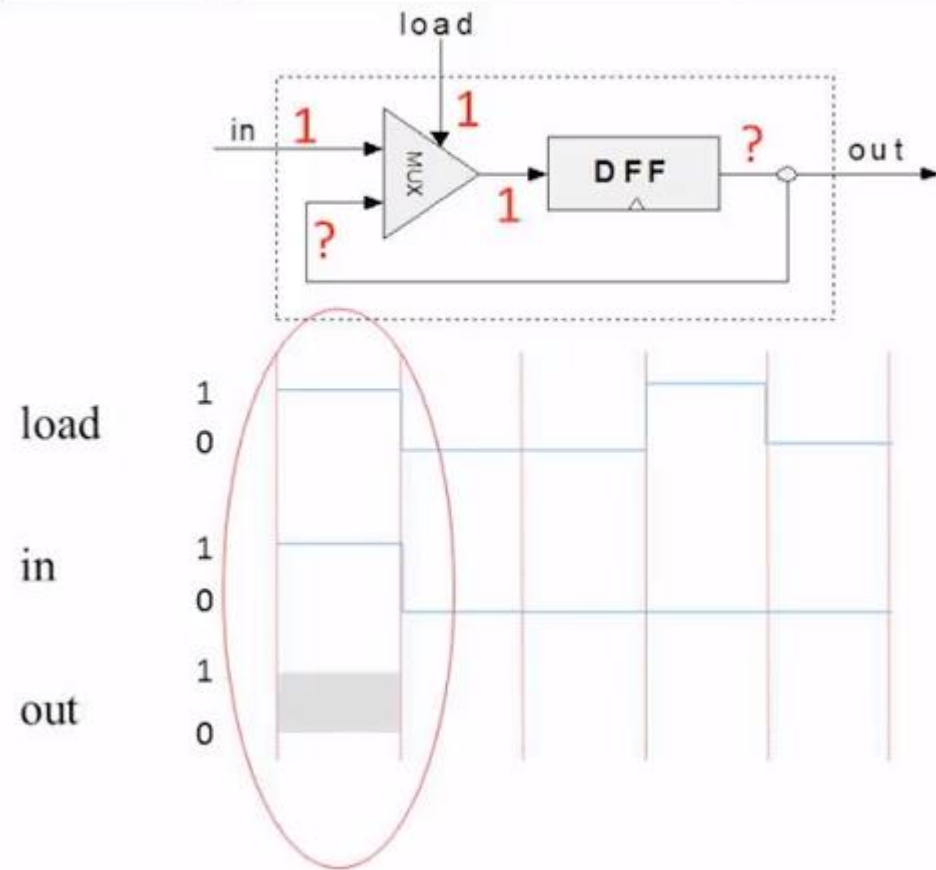


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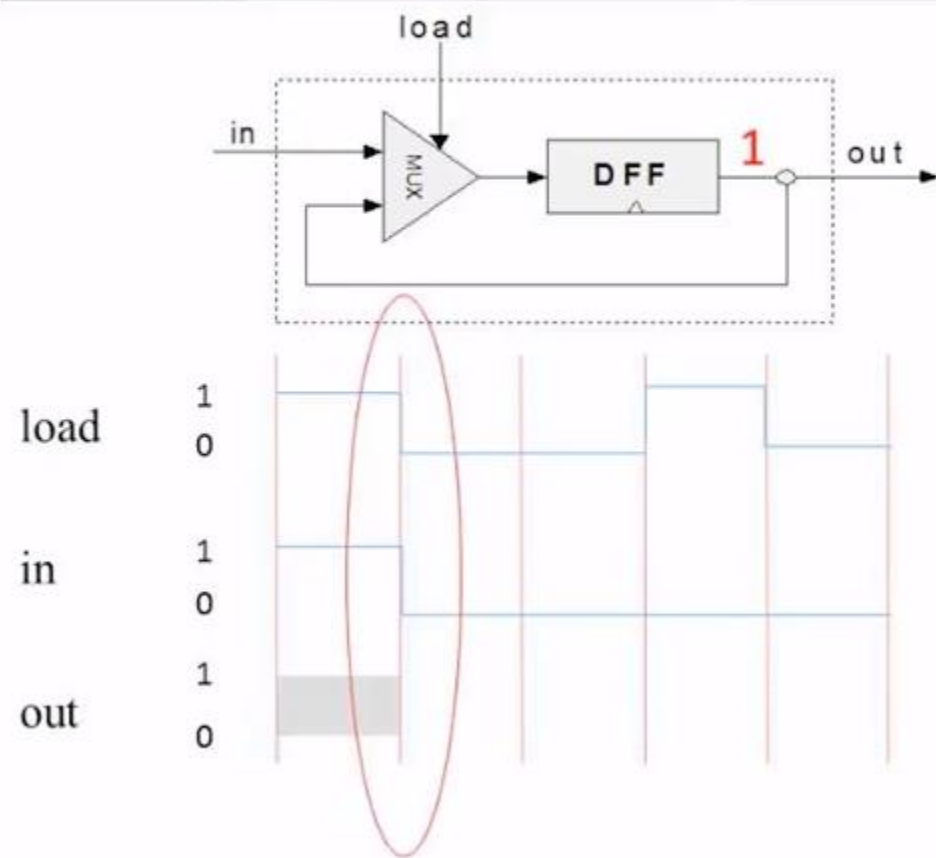




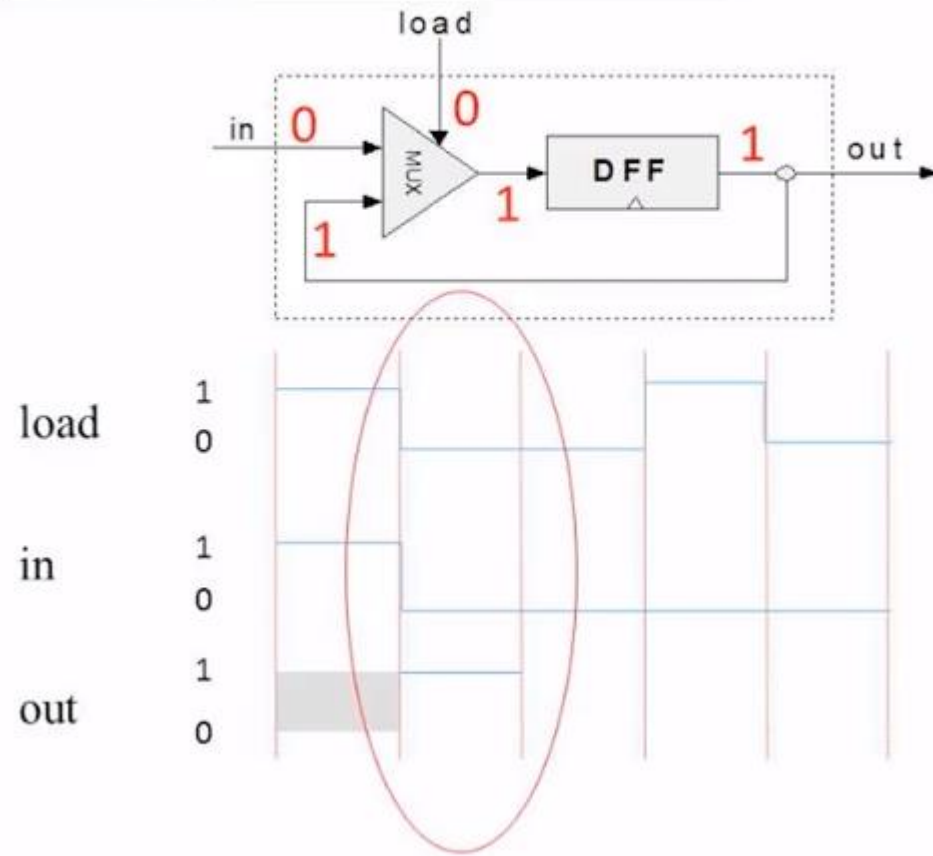
1-Bit Register

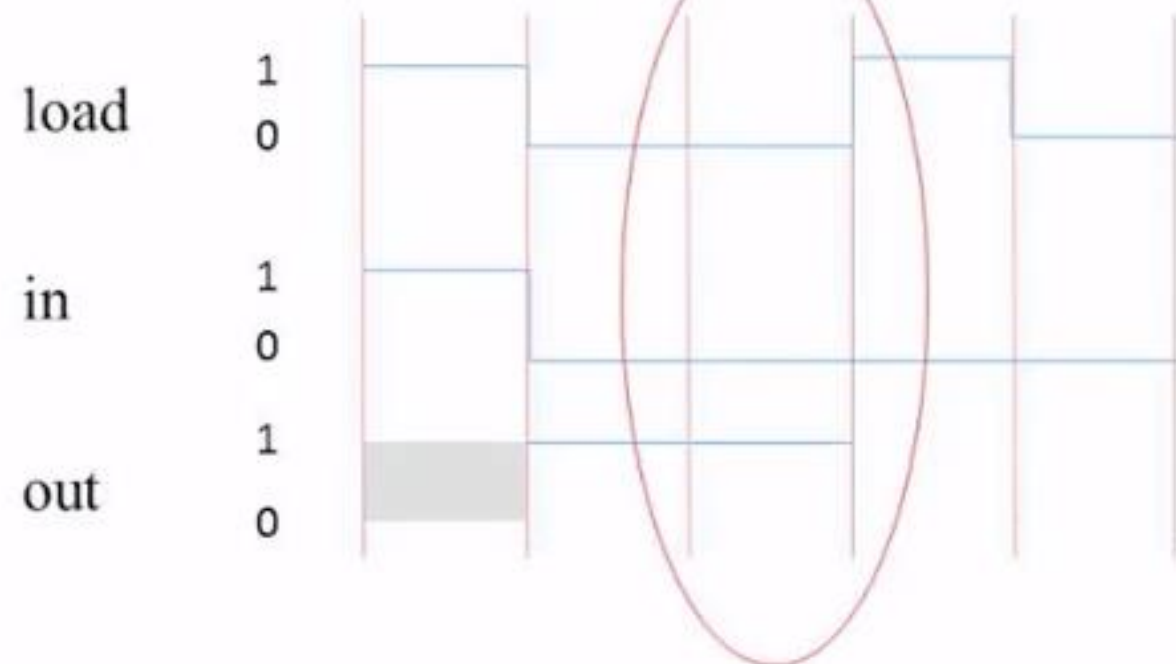
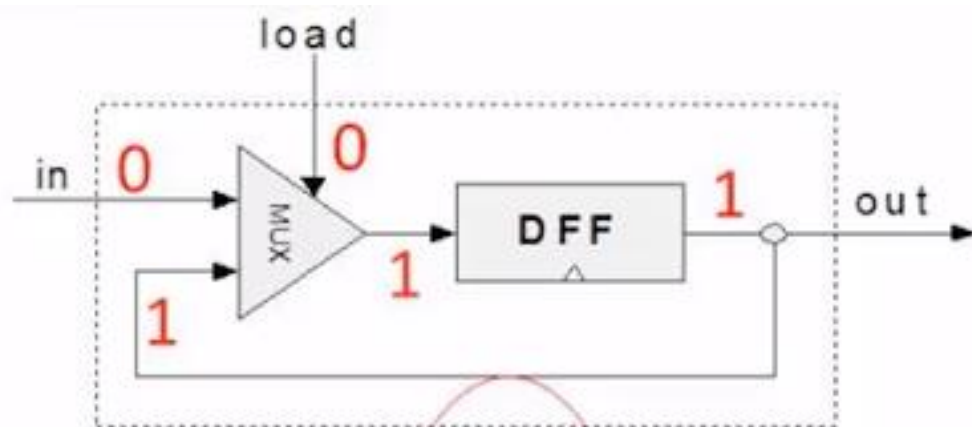


1-Bit Register

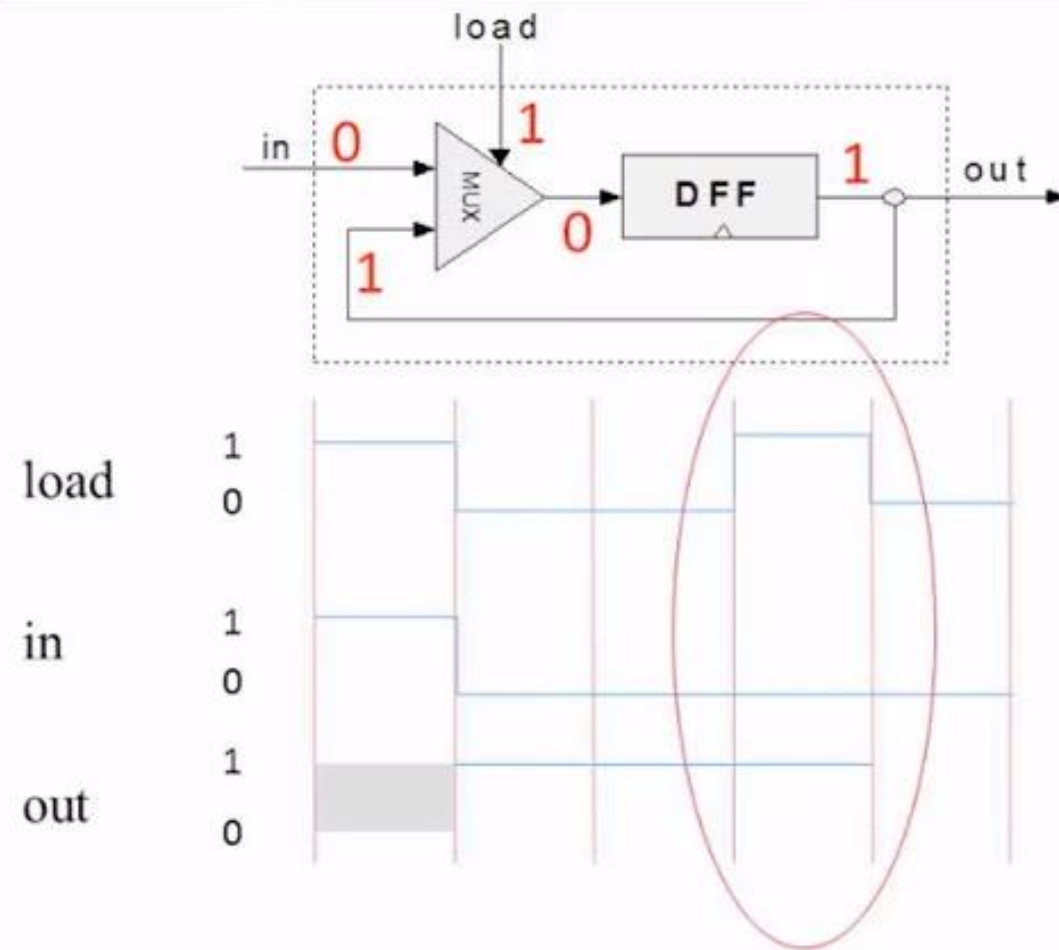


1-Bit Register

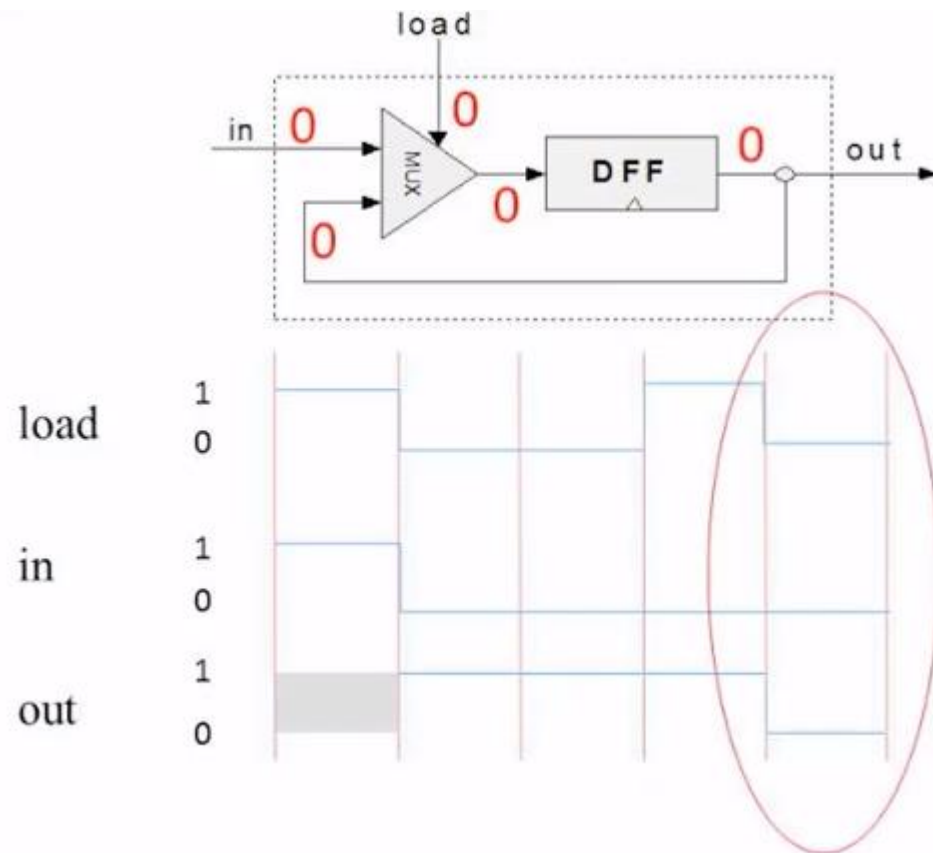




1-Bit Register



If load is 1, it takes the input value; if load is 0 it save the previous state..



- **MEMORY**

Memory

Memory:

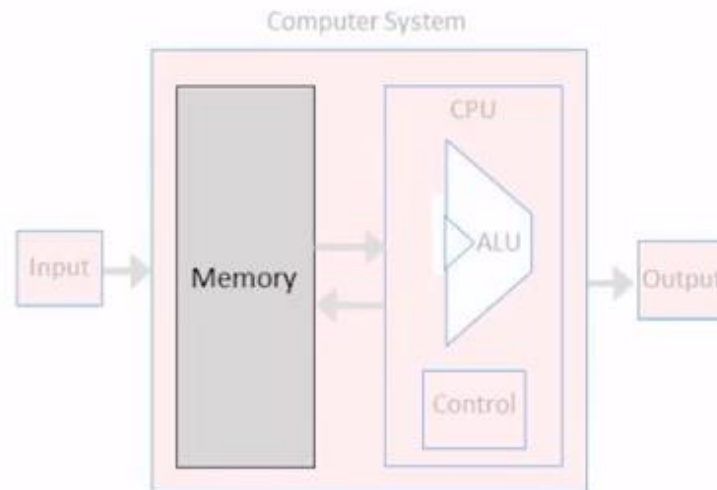
- Main memory: RAM, ...
- Secondary memory: disks, ...
- Volatile / non-volatile

RAM:

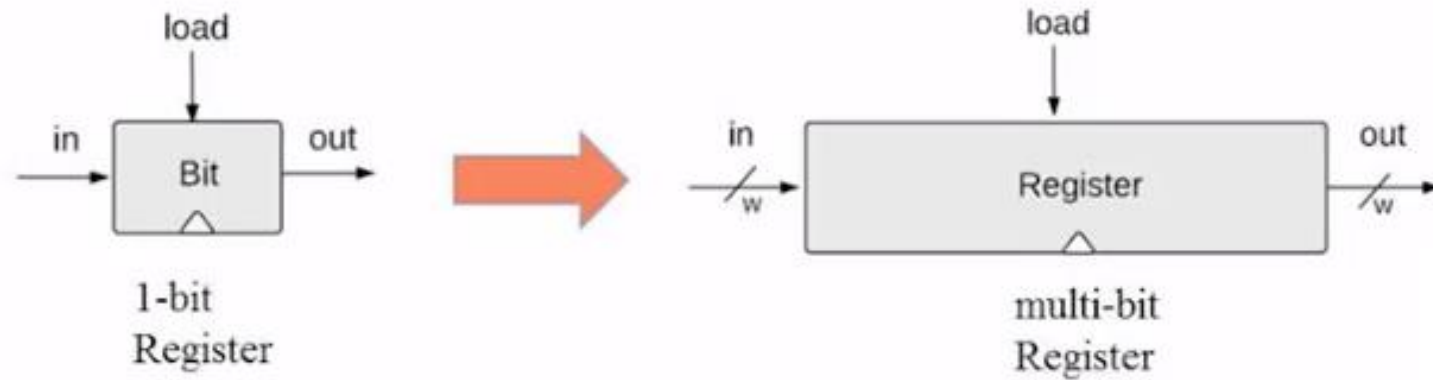
- Data
- Instructions

Perspective:

- Physical
- Logical

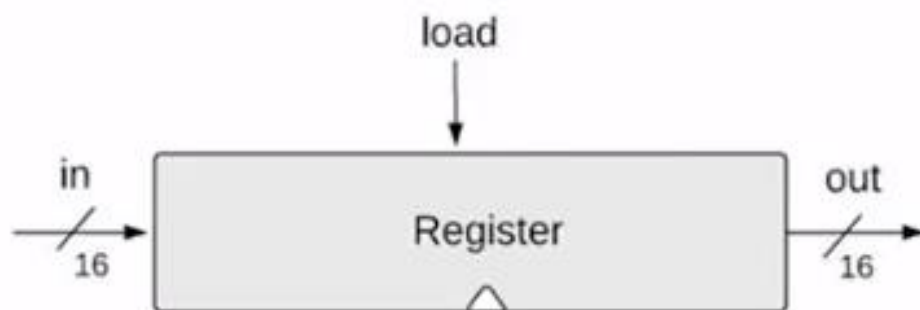


The most basic memory element: Register



- w (word width): 16-bit, 32-bit, 64-bit, ...
(from now on we will talk about 16-bit registers,
without loss of generality)
- Register's state: the value which is currently stored
inside the register

Register / read logic



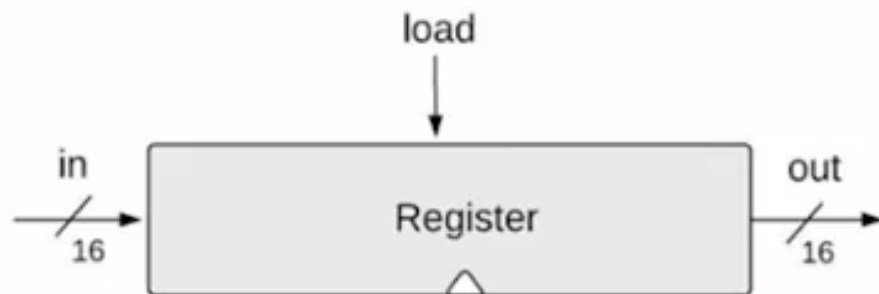
To read the Register:

probe out

Result:

out emits the Register's state

Register / write logic



To set Register = v

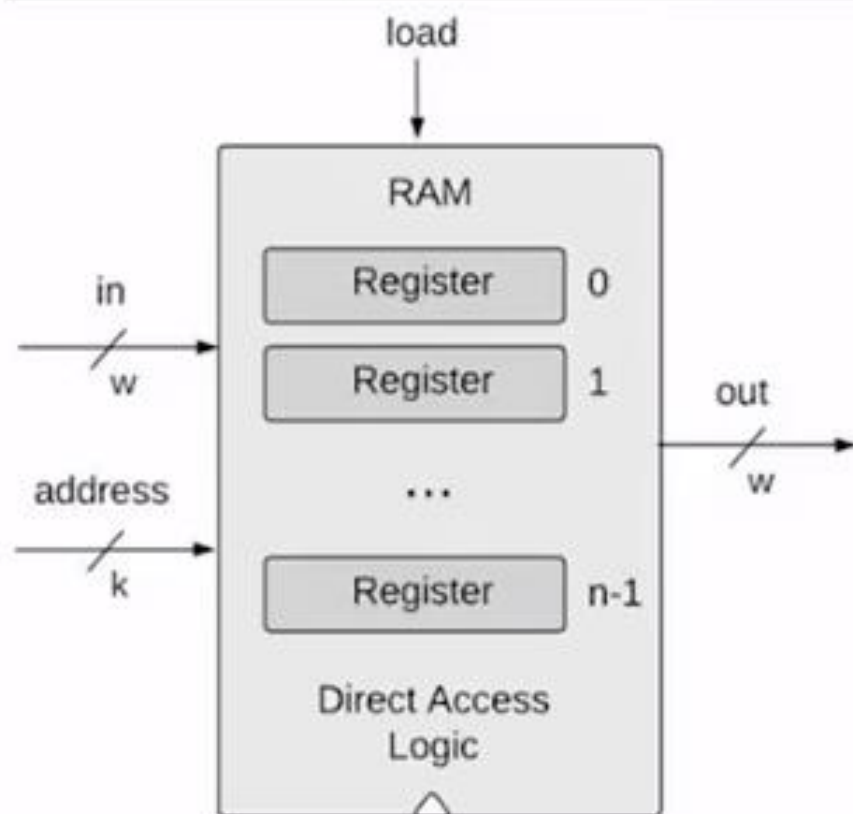
set in = v

set load = 1

Result:

- The Register's state becomes v
- From the next cycle onward,
out emits v

RAM unit



RAM abstraction:

A sequence of n addressable registers, with addresses 0 to $n-1$

At any given point of time, only *one* register in the RAM is selected

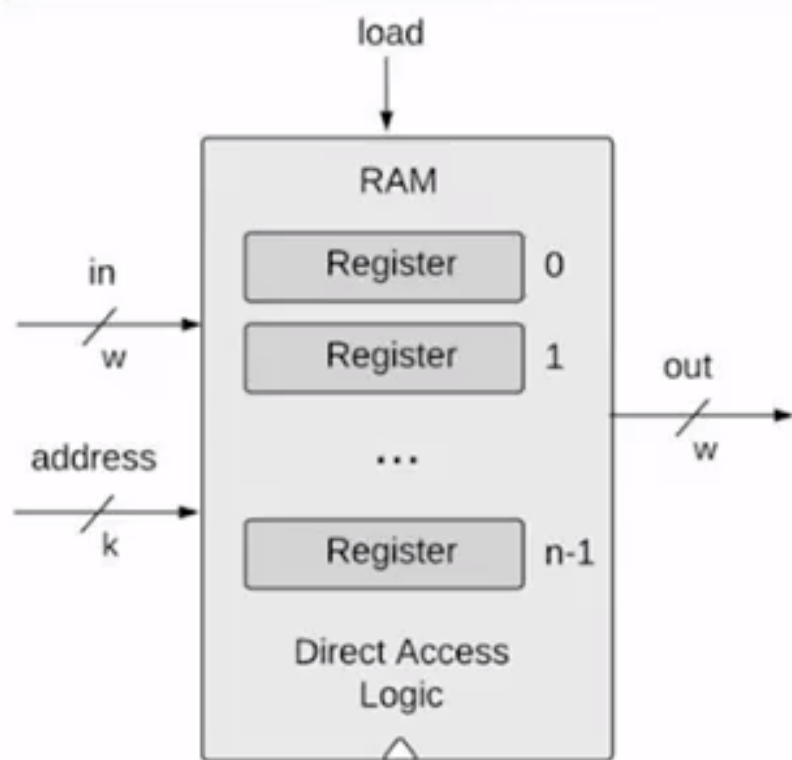
k (width of address input):

$$k = \log_2 n$$

w (word width):

No impact on the RAM logic
(Hack computer: $w=16$)

RAM / Read Logic



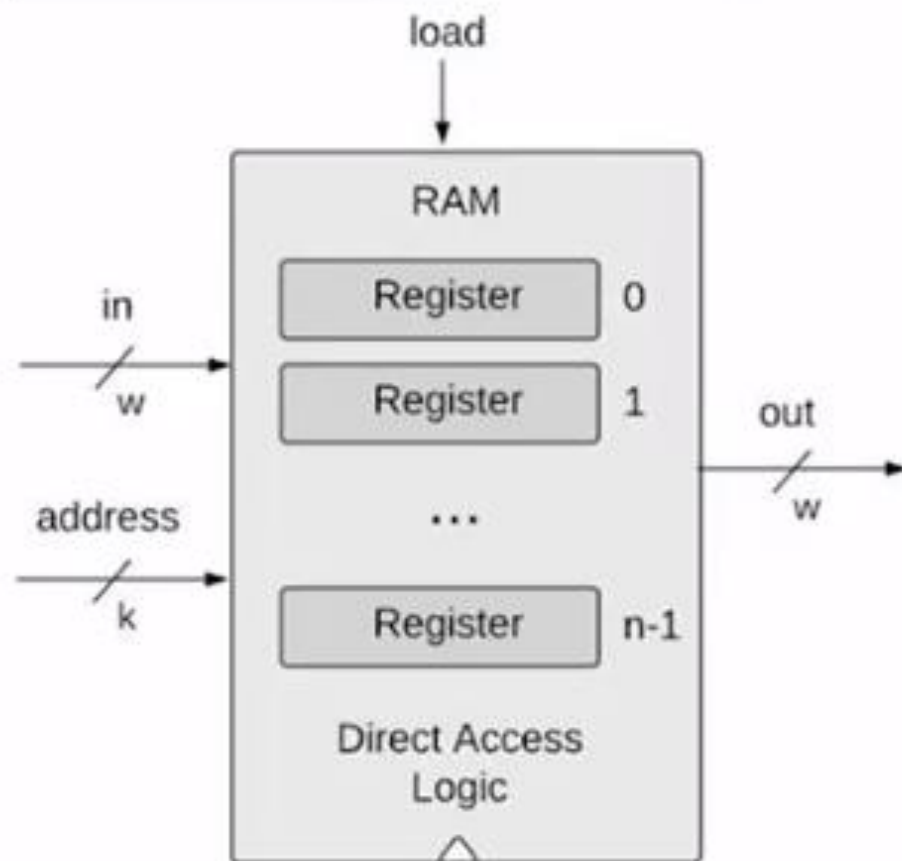
To read Register i :

set address = i

Result:

out emits the state of Register i

RAM / Write Logic



To set Register i to v :

set address = i

set in = v

set load = 1

Result:

- The state of Register i becomes v
- From the next cycle onward, out emits v