

Homework #2 - [Chapter 3]
(CS 372)

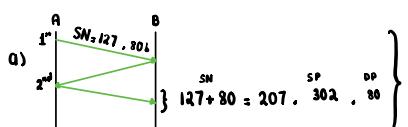
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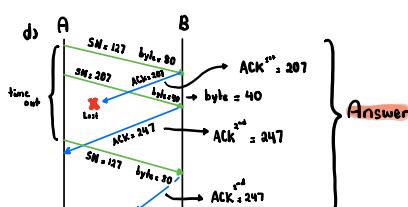
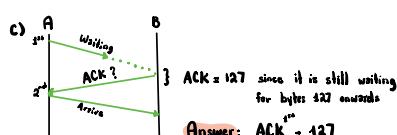
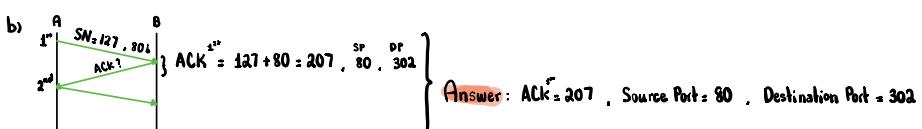
1. a) * Window size (N) = 4
* SN range = 1,024
* k = next in-order sn } Assuming k is the receiver acknowledging all $k-1$ pkts, then the sender's window can be in range of $[k, k+N-1]$.
② Assuming that sender's window contain $k-1$ & N pkts w/ include $k-1$, then we have sender's window be $[k-N, k-1]$

Then with these two cases and knowing window size = 4 we can conclude that at t time, set of possible SN is : Answer: $t = [k-N, k]$

- b) If no N ACKs have been received by the sender, the ACK is still propagating back so window will be $[k-N, k-1]$. This mean the sender does send all the pkt of $k-N$ where ACK could never be less than $k-N-1$ to $k-1$. So the possible value of ACK field in all possible messages is:

Answer: $k-N-1$ to $k-1$

2. a) (Send acknowledgement once received from A)
Host B received up to 126 bytes from A (A send back-to-back)
1st seg = 80 bytes & 2nd seg = 40 bytes data
SN = 127
SP = 302 DP = 80
A)  } $127 + 80 = 207$, 302 , 80 } Answer: Sequence N = 207, Source Port = 302, Destination Port = 80



3. * $\alpha = 0.1$
 * SampleRTT₁ = most recent & SampleRTT₂ = next most recent

$$a) \text{EstimatedRTT}_4 = (1-\alpha)^3 \text{SampleRTT}_4 + (1-\alpha)^2 \text{SampleRTT}_3 + (1-\alpha) \text{SampleRTT}_2 + \alpha \text{SampleRTT}_1$$

$$\text{Answer: } \text{ERTT}_4 = 0.9^3 \text{SRTT}_4 + 0.9^2 \text{SRTT}_3 + 0.9 \text{SRTT}_2 + 0.1 \text{SRTT}_1$$

b) Knowing from part A, we can determine the generalize formula to be:

$$\text{Answer: } (1-\alpha)^{n-1} \text{SampleRTT}_n + \frac{\alpha}{1-\alpha} \sum_{j=1}^{n-1} (1-\alpha)^j \text{SampleRTT}_j$$

c) If we move to ∞ from B, we get:

$$\text{Answer: } \lim_{n \rightarrow \infty} \text{EstimatedRTT}'' = \frac{\alpha}{1-\alpha} \sum_{j=1}^{\infty} (1-\alpha)^j \text{SRTT}_j \rightarrow \frac{1}{9} \sum_{j=1}^{\infty} (0.9)^j \text{SRTT}_j$$

* It is called for exp. moving average as the EstimatedRTT depend on previous recent SampleRTT to get the next EstimatedRTT which will produce the exp. moving average.

4. a) The interval where TCP have a slow start is:

$$\text{Answer: } t = [1, 6] \text{ & } [23, 26]$$

- b) The interval where TCP have a congestion avoidance is:

$$\text{Answer: } t = [6, 16] \text{ & } [17, 22]$$

- c) Looking at the event loss at transmission round at $t=15$, we see CongWin ≈ 42 and take $\frac{42}{2} = 21$ where $t=16$ start at the threshold, then:

Answer: This is a triple duplicate ACK

- d) Looking at the event loss at transmission round at $t=22$, we see CongWin ≈ 26 and take $\frac{26}{2} = 13$ where $t=23$ start at 1 where set CongWin to 1, then:

Answer: This is a timeout event

- e) The ssthreshold at the first transmission round is:

$$\text{Answer: } \text{threshold} \approx 32$$

- f) The ssthreshold at the 18th transmission round is:

$$\text{Answer: } \text{CongWin} = \frac{42}{2} = 21$$

- g) The ssthreshold at the 24th transmission round is:

$$\text{Answer: } \text{CongWin} = \frac{26}{2} = 13$$

- h) The 70th segment would be around 7th transmission round.

* We can see this as 1st trans. = 1 segs, 2nd trans = 2-3 segs, 3rd trans = 4-7 segs, ..., 7th trans = 64-96 segs.

Answer: 7th transmission round

- i) Assuming pkt loss after the 26th round by triple ACK, we can see
the 26th round = CongWin at 8 and we can take $\frac{8}{2} = 4$
Answer: CongWin & ssthresh = 4
- j) TCP Tahoe use instead TCP Reno & assume triple ACKs received at round 16th,
the ssthreshold & congWin at 19th \Rightarrow CongWin set to 1 at 19th & ss thresh = 21 $\Rightarrow \frac{21}{2}$
Answer: CongWin = 1 & ssthreshold = 21
- h) If Tahoe is used, timeout at 22nd round, the pkts send from 17-22th round
* 17th = 1 pkt * 21 = 16 pkts
* 18 = 2 pkts * 22 = 21 pkts - threshold } Add to get 52 pkts, increase exponentially
* 19 = 4 pkts
* 20 = 8 pkts } Answer: 52 pkts

5. a) * Assume cwnd increase by 1 MSS
* TCP uses AIMD for its congestion control w/o slow start
* Constant RTT } * Take 1RTT to increase CongWin = 7 MSS follow that:
→ 2RTT = 8 MSS
3RTT = 9 MSS
4RTT = 10 MSS
5RTT = 11 MSS
6RTT = 12 MSS } Answer: Take 6RTT from cwnd 6 \rightarrow 12 MSS
- b) Avg. throughput in time = 6RTT $\Rightarrow \frac{(6+7+8+9+10+11)}{6} \frac{\text{MSS}}{\text{RTT}}$
Answer: Avg. throughput $8.5 \frac{\text{MSS}}{\text{RTT}}$
7. a) * 10 Mbps link $\rightarrow 1.0 \times 10^7$ bps
* TCP segment size 1,500 bytes $\rightarrow 12,000$ bits
* $s = 150$ msec, ignore slow start
 0.15 sec
 $\frac{(1.0 \times 10^7)(0.15)}{12,000} = 125$ segs. \rightarrow Answer: Max window size = 125 segments
- b) Avg. window size? $\frac{3W}{4} = \frac{3 \cdot 125}{4} = 94$ segs.
Avg. throughput? $\frac{3W \cdot \text{MSS}}{4 \cdot \text{RTT}} = \frac{94 \cdot (12,000)}{0.15} = 7.52 \times 10^6$ bps } Answer: Avg. window size = 94 & Avg. throughput = 7.52×10^6 bps
- c) We would take the $\frac{W \cdot \text{RTT}}{2} = \frac{125 \cdot 0.15}{2} = 9.4$ sec } Answer