Checkpoint4

1. Typescript for compilation

I chose to implement on the Windows operating system, so I cannot use the 'rm' command. Here, I opted to use the 'del' command to achieve the same functionality.

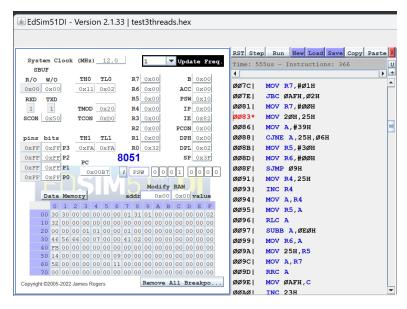
```
C:\清大資工\OS\112os\ppc4>make clean del *.hex *.ihx *.lnk *.lst *.map *.mem *.rel *.rst *.sym *.asm *.lk

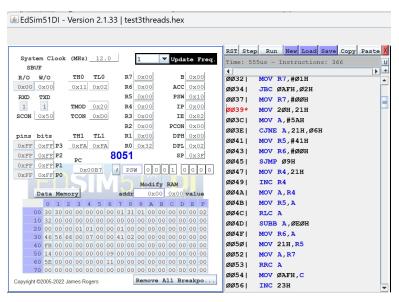
C:\清大資工\OS\112os\ppc4>make sdcc -c test3threads.c test3threads.c:57: warning 158: overflow in implicit constant conversion sdcc -c preemptive.c preemptive.c preemptive.c:96: warning 85: in function ThreadCreate unreferenced function argument : 'fp' sdcc -o test3threads.hex test3threads.rel preemptive.rel
```

2-1.Take screenshots when the Producer1 and Producer2 running and show semaphore changes.



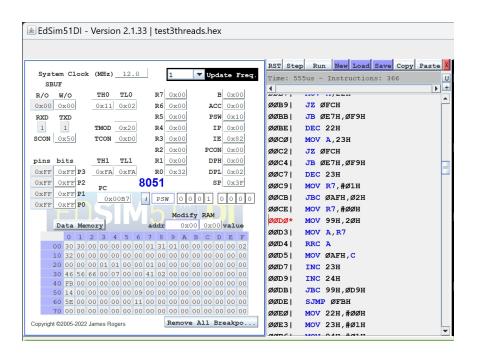






- 1.At 0083, executing MOV 20H,25H in the producer(buffer \leftarrow _input1), the values of semaphores(full,empty,mutex) (at 0x22, 0x24,0x23) is changed (0,0,0).
- 2.Because the producer is in the critical section, and hasn't produced an input. Also, it hasn't finished producing ,so the flag (26H) and flag1(27H) are 0. After it finish, it will set flag1(27H) to 1
- 1-1.At 0039, executing MOV 20H,21H in the producer1(buffer ← _input), the values of semaphores(full,empty,mutex) (at 0x22, 0x24,0x23) is changed (0,0,0).
- 2-1.Because the producer1 is in the critical section, and hasn't produced an input. Also, it hasn't finished producing ,so the flag (26H) and flag1(27H) are 0. After it finish, it will set flag(26H) to 1..

2-2. Take screenshots when the Consumer is running and show semaphore changes



- 1.At 00D0, executing MOV 99H, 20H in the consumer(SBUF ← buffer), the values of semaphores(full, empty, mutex)(at 0x22, 0x24,0x23) is changed (0,0,0).
- 2.Because the consumer is in the critical section, and hasn't consumed an input. Besides, one of the semaphore flag(26H) or flag1(27H) is changed from 0 to 1 and the other remains 0. The one whose flag becomes 1 will be the next producer to produce.

2-3. Show and explain UART output to show the unfair version, if any, and the fair version.

1.If using the Round-Robin scheduling policy, the producer executing after the consumer will always fill the buffer. Therefore, the other producer will never get the chance to produce.

2.Adding a semaphore flag to each producer and making them wait for their own flag, then signaling the other one after they finish, will result in a fair version.