

Algorithms: The Dynamic Dining Philosophers

Problem Statement. We need to develop and implement a protocol that allows nodes to join or leave a running dining philosophers network at runtime in an orderly fashion while maintaining an acyclic priority graph, ensuring safety, and preventing deadlock.

Allowed Assumption

External controllers

If a philosopher p informs an external controller that it is *eating*, the external controller will direct p to become either *thinking* or *leaving* within a bounded time. External controllers will not send duplicate or invalid control signals.

Solution.

```

1 infinite_loop(Ref, Nodel, Neighbors) ->
2   spelling, Node} ! {self(), Ref, become_hungry},
3   % {spelling, Node} ! {self(), Ref, stop_eating},
4   % {spelling, Node} ! {self(), Ref, leave},
5   receive
6     %{Ref, eating} ->
7     %print("~p is eating.\n", [Ref]);
8     {Ref, gone} ->
9     print("~p is gone.\n", [Ref]);
10    Reply ->
11    print("Got unexpected message: ~p~n", [Reply])
12    after ?TIMEOUT -> print("Timed out waiting for reply!")
13  end,
14  infinite_loop(Ref, Nodel, Neighbors)
15 end.
```

```

1 philosophize(Ref, joining, Node, ForksNeighborsList)->
2   print("joining"),
3   %philosophize(Ref, thinking, Node, Neighbors);
4   requestJoin(Ref, Node, ForksNeighborsList),
5   philosophize(Ref, thinking, ForksNeighborsList).
6
7 %requests each neighbor to join the network, one at a time,
8 %when joining there shouldn't be any other requests for forks or leaving going on
9 requestJoin(-, -, []) -> ok;
10 requestJoin(Ref, Node, ForksNeighborsList)->
11   try
12     io:format("Process ~p at node ~p sending request to ~n~s",
13       [self(), Node, hd(Neighbors)]),
14     io:format("After ~n"),
15     {list_to_atom(hd(Neighbors)), Node} ! {self(), Ref, requestJoin},
16     receive
17       {Ref, ok} ->
18         io:format("Got reply (from ~p): ok!",
19           [Ref]);
20       Reply ->
21         io:format("Got unexpected message: ~p~n", [Reply])
22     after ?TIMEOUT -> io:format("Timed out waiting for reply!")
23   end,
24   requestJoin(Ref, Node, tl(Neighbors))
```

```

25  catch
26  _:_ -> io:format("Error getting joining permission.\n")
27  end.

```

```

1  philosophize(Ref, thinking, Node, ForksNeighborsList)->
2      receive
3          {self(), NewRef, leave} ->
4              print("leaving"),
5              philosophize(NewRef, leaving, ForksNeighborsList);
6          {self(), NewRef, become_hungry} ->
7              print("becoming hungry"),
8              philosophize(NewRef, hungry, ForksNeighborsList)
9  after ?TIMEOUT -> print("Timed out waiting for reply!")
10 end;

```

```

1  philosophize(Ref, hungry, Node, ForksNeighborsList)->
2      receive
3          {self(), NewRef, leave} ->
4              print("leaving"),
5              philosophize(NewRef, leaving, Node, ForksNeighborsList);
6      % {self(), NewRef, Fork} -> AND/OR CHECK IF ALL NEIGHBORS ARE NOT EATING?
7          %print("got fork"),
8          % check if has all forks
9          % continue with philosophize(NewRef,
10
11      %end;
12
13      % want to receive all forks and then start eating
14      {controller, Node} ! {NewRef, eating},
15      print("eating"),
16      philosophize(Ref, eating, Node, ForksNeighborsList)
17 end;

```

```

1  philosophize(Ref, eating, Node, ForksNeighborsList)->
2      receive
3          {self(), NewRef, stop_eating} ->
4              % handle forks and hygenity if doing that
5              print("stopping eating"),
6              philosophize(NewRef, thinking, Node, ForksNeighborsList);
7          {self(), NewRef, leave} ->
8              print("stopping eating and leaving"),
9              %get rid of forks
10             philosophize(NewRef, leaving, Node, ForksNeighborsList)
11  after ?TIMEOUT -> print("Timed out waiting for reply!")
12 end;

```

```

1  philosophize(Ref, leaving, Node, ForksNeighborsList)->
2      %need to gather forks and then leave with them
3      {controller, Node} ! {Ref, gone}.

```

Progress.

$p.joining \rightsquigarrow p.thinking$

After philosopher p has started up, it has given itself the *joining* state. Assuming that the neighbors in which p knows about are running correctly and the network runs as expected, all other philosophers are bound to hear p 's request to join eventually and thus p is guaranteed to be given the state *thinking*.

$p.hungry \rightsquigarrow p.eating$

When *hungry*, a philosopher p must be passed the forks once all neighbor philosophers are not eating

$p.leaving \rightsquigarrow p.gone$

Safety.

In order to show that the algorithm satisfies the Safety requirement, we must show that the relevant stages satisfy the Safety properties:

initially $p.joining$

p is given the state of *joining* in which the philosopher p is requesting to join the group and cannot possibly gain any other state until granted acceptance.

$p.joining$ **next** $(p.joining \vee p.thinking)$

This requirement is satisfied because p can be constantly trying to join the party but may be waiting infinitely or he/she can be granted the state of thinking (which is the only initial state in the party).

$p.thinking$ **next** $(p.thinking \vee p.hungry \vee p.leaving)$

When thinking, the philosopher p can continue thinking, the external controller can issue the order to become hungry, or the controller may tell the philosopher to leave. No other "state" transitions are available to the philosopher at the *thinking* state.

$p.hungry$ **next** $(p.hungry \vee p.eating \vee p.leaving)$

If a philosopher p is told by the external controller to become hungry, then it may be told to leave, it may become eating due to its hungry nature, or p may remain hungry (which would likely imply a failure in the system).

$p.eating$ **next** $(p.eating \vee p.thinking \vee p.leaving)$

$p.leaving$ **next** $(p.leaving \vee p.gone)$

$p.gone$ **next** $(p.gone)$

$p.eating \rightarrow \langle \forall q | q \in p.neighbors \triangleright /q.eating \rangle$ (when a philosopher p is *eating*, none of its neighbors is *eating*)

$(p.thinking \vee p.hungry \vee p.eating) \langle \rightarrow \forall q | q \in p.neighbors \triangleright p \in q.neighbors \rangle$ (when a philosopher p is *thinking*, *hungry*, or *eating*, each of p 's neighbors knows that p is one of its neighbors)

$p.gone \rightarrow \langle \forall q \triangleright q \notin q.neighbors \rangle$ (when a philosopher p is *gone*, it is not in any other philosopher's set of neighbors)

□