

Sparse Matrix Factorization from an Optimization Point of View

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1 Introduction

2 NP-hardness

3 Existence of optimal solutions

4 A polynomial algorithm for easy instances

5 Multiple factors matrix factorization

6 Back to two factors: Optimization landscape

Sparse matrix factorization

Objective: Given A , find *multiple* factors $S^{(1)}, S^{(2)}, \dots, S^{(J)}$ such that:

$$A \approx S^{(1)}S^{(2)} \dots S^{(J)}$$

where $S^{(i)}$ are *sparse* matrices.

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Motivations

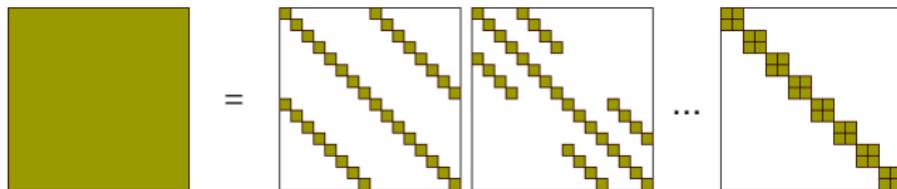
- *Fast* matrix vector products:

$$\underbrace{A}_{\text{dense}} \approx \underbrace{S^{(1)}S^{(2)} \dots S^{(J)}}_{\text{sparse}} \Rightarrow Ax \approx S^{(1)}(S^{(2)}(\dots(S^{(J)}x)))$$

- Reduce time + memory complexity

Applications

- Fast Fourier Transform, Fast Hadamard Transform, etc.



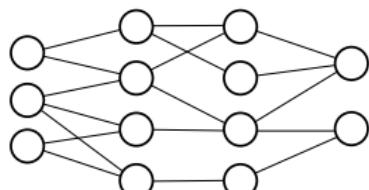
- Dictionary learning

- $A = XY^\top$, A data, X a base (words in a dictionary), Y representation of each sample using the dictionary.

[S. Foucart, H. Rauhut, A mathematical introduction to compressive sensing, ANHA, 2013]

- Sparse (linear) neural networks:

- Interpretability.
- Energy efficiency.



[T. Dao & all. Learning fast algorithms for linear transforms using butterfly factorizations, PMLR, 2019]

[B. Chen & all. Pixelated butterfly: Simple and efficient sparse training for neural network models, PMLR, 2022]

Sparse Matrix Factorization Problem

Given A and \mathcal{E}_j some sets of sparse matrices, solve:

$$\min_{S^{(1)}, \dots, S^{(J)}} \|A - \prod_{j=1}^J S^{(j)}\|_F^2 \text{ subject to: } S^{(j)} \in \mathcal{E}_j, \forall j \in \{1, \dots, J\}$$

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- Example of *sparsity patterns* \mathcal{E}_j : set of matrices with at most k nonzero entries per **row**, per **column** or in **total**.
- Known to be **NP-hard** (covers sparse PCA, sparse dictionary learning)

[Malik, NP-hardness and inapproximability of sparse PCA, IPL, 2017]

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A general formulation for sparse matrix factorization

Sparse Matrix Factorization Problem

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→ A challenging problem, how to deal with it?

Our approach: from two to multiple factors

- **Two** factors matrix factorization: the simplest nontrivial case.

Given A , minimize $\underset{X,Y}{\|A - XY^T\|_F^2}$ subject to: X, Y sparse matrices

- **Multiple** factors matrix factorization: butterfly factorization.

Two sub-problems of **two** factors matrix factorization

$$\underset{X, Y}{\text{Minimize}} \quad \|A - XY^T\|_F^2 \quad \text{subject to: } X, Y \text{ sparse matrices}$$

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Find two sets S_X, S_Y - the supports of two factors X, Y

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Find two sets S_X, S_Y - the supports of two factors X, Y

2) Optimize coefficients inside support

$$\underset{X \in \mathbb{R}^{m \times r}, Y \in \mathbb{R}^{n \times r}}{\text{Minimize}} \quad L(X, Y) = \|A - XY^\top\|_F^2$$

Subject to: $\text{supp}(X) \subseteq S_X$

$\text{supp}(Y) \subseteq S_Y$

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→ **Second problem:** Fixed support matrix factorization (FSMF).

Examples of FSMF

FSMF covers:

- Low rank matrix decomposition

$$A = \begin{matrix} & r \\ \text{---} & \text{---} \\ \text{---} & \text{---} \end{matrix} \times \begin{matrix} & r \\ \text{---} & \text{---} \\ \text{---} & \text{---} \end{matrix}$$

- LU decomposition

$$A = \begin{matrix} \text{yellow} & & & \\ & \text{yellow} & & \\ & & \text{white} & \\ & & & \text{yellow} \end{matrix} \times \begin{matrix} \text{white} & \text{yellow} & & \\ \text{yellow} & \text{white} & \text{yellow} & \\ & \text{yellow} & \text{white} & \text{yellow} \\ & & \text{yellow} & \text{white} \end{matrix}$$

- Hierarchical \mathcal{H} matrices

$$A = X \times Y^\top$$

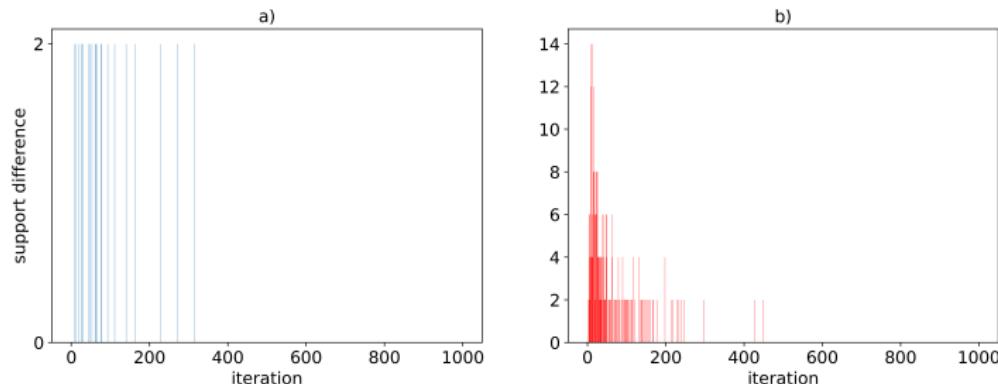
- Butterfly factorization (multiple factors)

$$\begin{matrix} \text{red} & \text{blue} \\ \text{blue} & \text{red} \end{matrix} = \begin{matrix} \text{red} & & & \\ & \text{red} & & \\ & & \text{red} & \\ & & & \text{red} \end{matrix} \times \begin{matrix} & & & \\ & \text{red} & & \\ & & \text{red} & \\ & & & \text{red} \end{matrix} \times \begin{matrix} & & & \\ & \text{red} & & \\ & & \text{red} & \\ & & & \text{red} \end{matrix} \times \begin{matrix} & & & \\ & \text{red} & & \\ & & \text{red} & \\ & & & \text{red} \end{matrix}$$

FSMF: further motivation

Understand the *asymptotic behaviour* of other existing heuristics:

- Supports asymptotically **don't change**
- Spurious local valley in the landscape of $L(X, Y)$: certain heuristic can lead to iterates **diverging** to infinity



Measure of the difference between two consecutive supports

Main contributions on FSMF

- ① Is the problem **polynomially tractable**?
→ We prove its *NP-hardness*.
- ② Does the problem have a **solution**?
→ We show an instance where optimal solutions do *not* exist.
- ③ Are there **easy instances**?
→ We individuate a family of *polynomially solvable* instances and proposed an *efficient algorithm*.
- ④ How well does **gradient descent** tackle the problem of FSMF?
→ We study the properties of the *landscape* of the function $L(X, Y) = \|A - XY^\top\|^2$ under the support constraints.

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NP-hardness

Theorem (1)

Sparse matrix factorization is NP-hard even with fixed support

Proof.

Rank-one matrix completion with noise is reducible to FISMF. □

- In line with recent results on matrix factorization:
 - non-negative matrix factorization (NMF)
 - weighted low rank
 - matrix completion

[N. Gillis, F. Glineur, Low-rank matrix approximation with weights or missing data is NP-hard. SIAM JMAA, 2010]

[S. A. Vavasis, On the complexity of nonnegative matrix factorization, SIOPT, 2010]

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LU decomposition and non-closedness

$$A = \begin{matrix} & \begin{matrix} \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \end{matrix} \\ A & = & \begin{matrix} \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \end{matrix} \end{matrix} \times \begin{matrix} & \begin{matrix} \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \end{matrix} \\ & \times & \begin{matrix} \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} & \text{---} & \text{---} \end{matrix} \end{matrix}$$

The diagram illustrates the LU decomposition of a 5x5 matrix A. Matrix A is shown as a grid where yellow cells represent non-zero elements and white cells represent zero elements. It is decomposed into two matrices: L (Lower triangular) and U (Upper triangular). Matrix L has ones on its main diagonal and zeros below it. Matrix U has ones on its main diagonal and zeros above it. The multiplication of L and U results in the original matrix A.

LU decomposition and non-closedness

- Fact: Any square matrix is the limit of a sequence of matrices having an LU decomposition.
 - But there exist square matrices that **do not have** an exact LU decomposition.

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- The set of matrices having LU decomposition is **not closed**
- For certain support constraints (S_X, S_Y) and matrices A , **(FSMF)** does not have an optimal solution.

LU decomposition and non-closedness

$$A = \begin{array}{c|c|c|c} \textcolor{yellow}{\blacksquare} & \square & \square & \square \\ \hline \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \square & \square \\ \hline \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \square \\ \hline \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} \\ \hline \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} \\ \hline \end{array} \times \begin{array}{c|c|c|c} \textcolor{yellow}{\blacksquare} & \square & \square & \square \\ \hline \square & \textcolor{yellow}{\blacksquare} & \square & \square \\ \hline \square & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \square \\ \hline \square & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} \\ \hline \square & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} & \textcolor{yellow}{\blacksquare} \\ \hline \end{array}$$

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- The set of matrices having LU decomposition is **not closed**
- For certain support constraints (S_X, S_Y) and matrices A , **(FSMF)** does not have an optimal solution.

Question : Given (S_X, S_Y) , do their FSMF instances always have optimal solutions? → This problem is *decidable* (not known to be *tractable* yet).

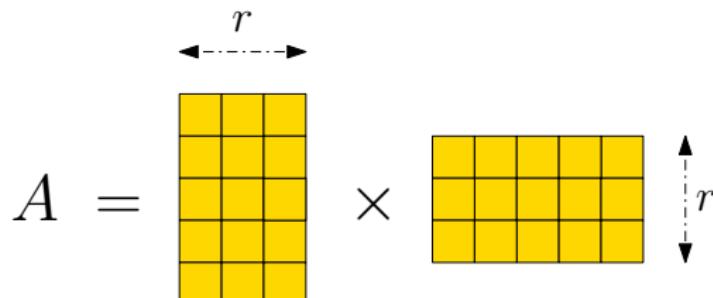
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Polynomially solvable instances

Example (Unconstrained matrix factorization)

When there is **no constraints** on the support of X and Y :

$$\underset{X \in \mathbb{R}^{m \times r}, Y \in \mathbb{R}^{n \times r}}{\text{Minimize}} L(X, Y) = \|A - XY^\top\|_F^2$$

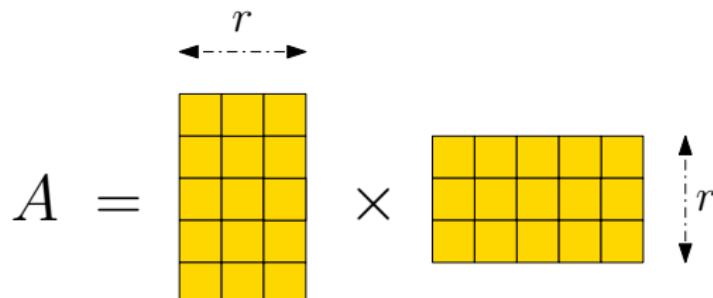


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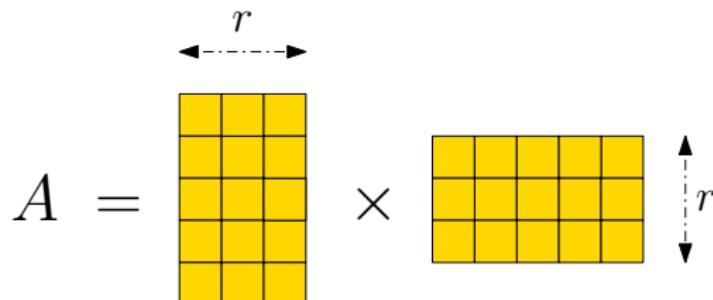
→ Solution: Use Singular Value Decomposition (SVD).

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- Solution: Use Singular Value Decomposition (SVD).
- Question: Can Singular Value Decomposition (SVD) still work in the *constrained case*?

SVD as a greedy algorithm

- 1) Decompose the problem:

$$A - XY^\top = A - \sum_{i=1}^r x_i y_i^\top = A - \sum_{i=1}^r \underbrace{M_i}_{\text{rank-one}} \quad (M_i := x_i y_i^\top)$$

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- 2) Finding the SVD:

$$\text{bestRankOneApprox}(A) \rightarrow M_1$$

$$\text{bestRankOneApprox}(A - M_1) \rightarrow M_2$$

...

$$\text{bestRankOneApprox}(A - M_1 \dots - M_{r-1}) \rightarrow M_r$$

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→ SVD is a greedy algorithm in disguise

Algorithm 1 Algorithm for *unconstrained* matrix factorization

```
1: for  $i \in \{1, \dots, r\}$  do
2:    $M_i :=$  best rank-one approximation of  $A - \sum_{k=1}^{i-1} M_k$ .
3: end for
```

SVD as a greedy algorithm: the *constrained* case

- How to generalize the greedy algorithm?

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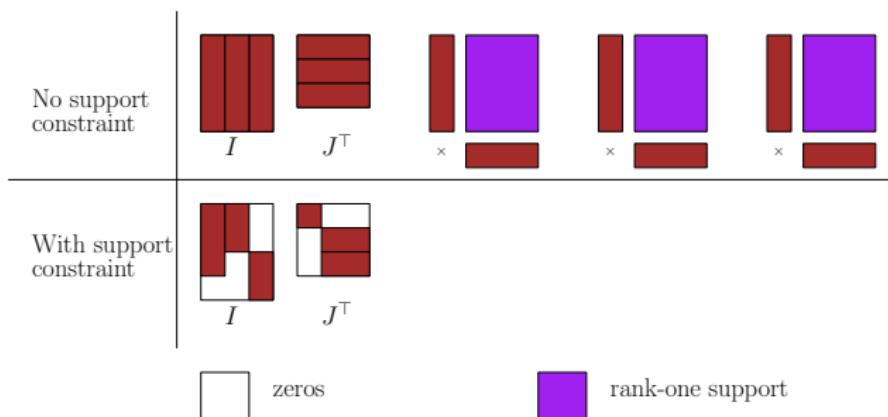
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SVD as a greedy algorithm: the *constrained* case

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$$XY^\top = M_1 + M_2 + M_3$$

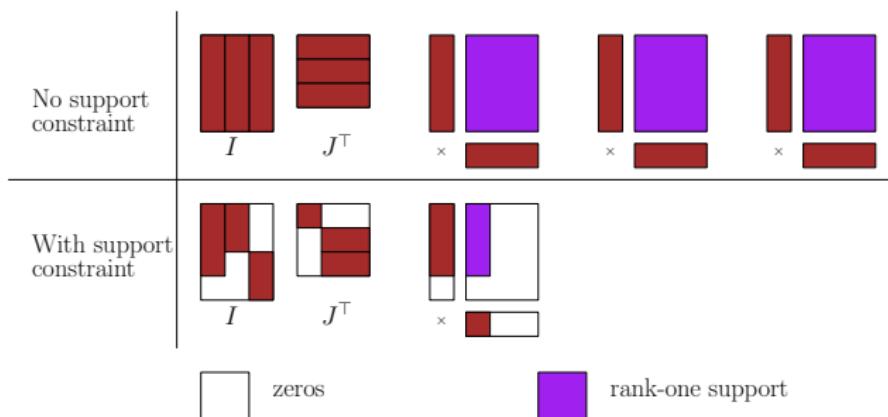


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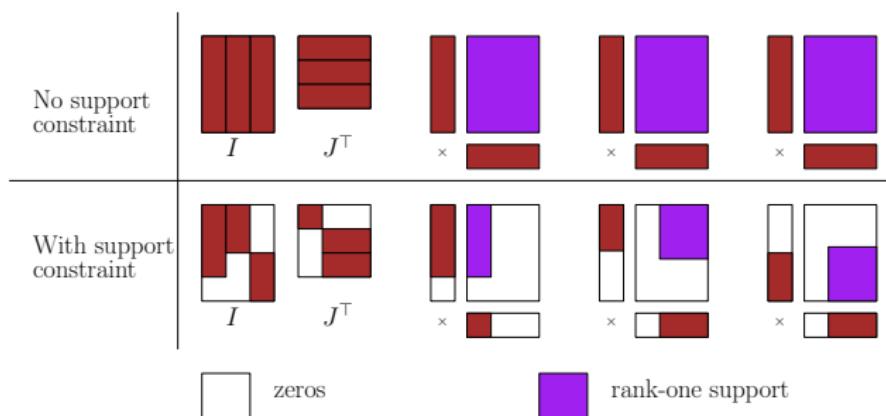


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- Finding optimal solution $(X, Y) \Leftrightarrow$ Finding optimal the rank-one constrained supports.

Algorithm 2 Algorithm for fixed-support matrix factorization

```
1: for  $i \in \{1, \dots, r\}$  do
2:    $S_i \leftarrow i\text{-th rank-one support}$ 
3:    $M_i := \text{best rank-one approximation of } (A - \sum_{k=1}^{i-1} M_k) \odot S_i$ 
4: end for
```

Example:

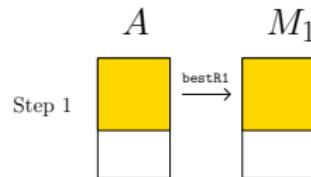
$$A \approx M_1 + M_2 + M_3$$

The diagram shows a tall white rectangle labeled A on the left, followed by a small gray double-headed arrow symbol. To its right is the expression \approx . After the \approx symbol are three rectangles stacked vertically. The top rectangle is yellow and labeled M_1 . The middle rectangle is blue and labeled M_2 . The bottom rectangle is teal and labeled M_3 . Between the M_1 and M_2 rectangles is a plus sign, and between the M_2 and M_3 rectangles is another plus sign.

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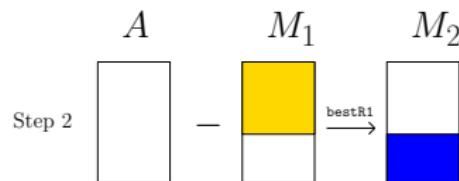
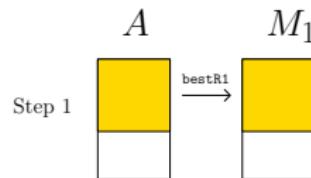
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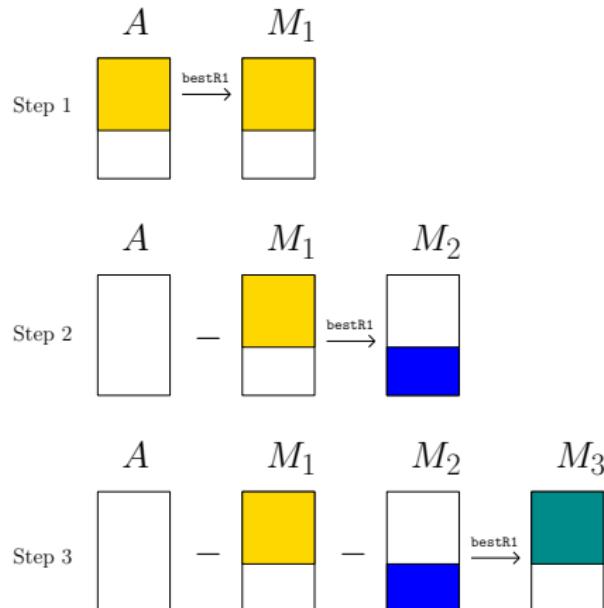
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```

Example:



Polynomial solvability characterized by rank-one supports

- The output of the greedy algorithm will always *satisfy the constraints*
- Is the output an optimal solution? Not always

Theorem (2)

If the rank-one supports are pairwise disjoint or identical, then the greedy algorithm gives an optimal solution.

For the same result with a weaker assumption:

[QT. Le, E. Riccietti, R. Gribonval, Spurious Valleys, NP-hardness, and Tractability of Sparse Matrix Factorization With Fixed Support, SIAM Journal of Matrix Analysis and Applications, In press, 2022.]

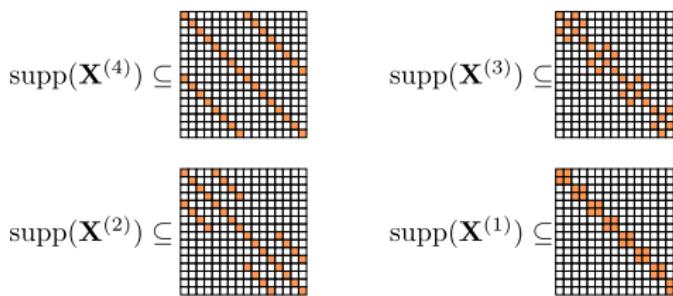
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Multiple-factors cases: the butterfly factorization

A special case: the butterfly factorization

Approximate **any** matrix by a product of $J \geq 2$ **butterfly** factors

Let $\mathbf{A} := \mathbf{X}^{(4)}\mathbf{X}^{(3)}\mathbf{X}^{(2)}\mathbf{X}^{(1)}$ such that:



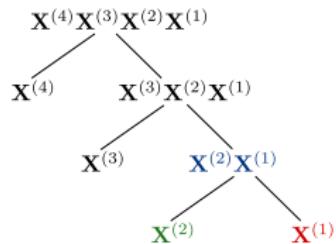
- It is **expressive**: the composition of matrices with a butterfly structure can accurately approximate any given matrix
- In neural networks faster training and inference time without harming the performance

[T. Dao & all. Kaleidoscope: An efficient, learnable representation for all structured linear maps, ICLR, 2020]

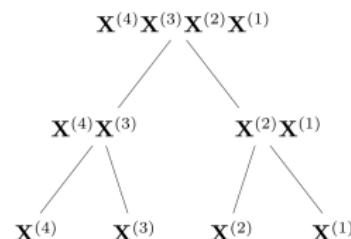
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Butterfly factorization: theoretical guarantees

Hierarchical approach: Use our algorithm to recover the partial factors: solve a sequence of **two factors** problems, if the supports are known



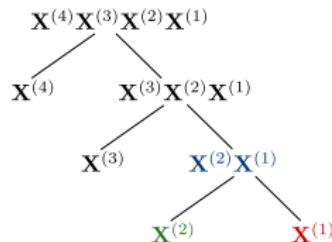
Left-to-right tree



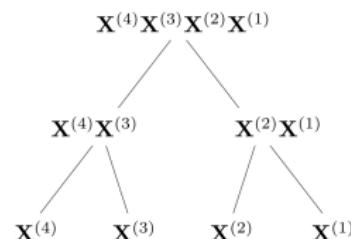
Balanced tree

Butterfly factorization: theoretical guarantees

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Left-to-right tree



Balanced tree

Theorem (3)

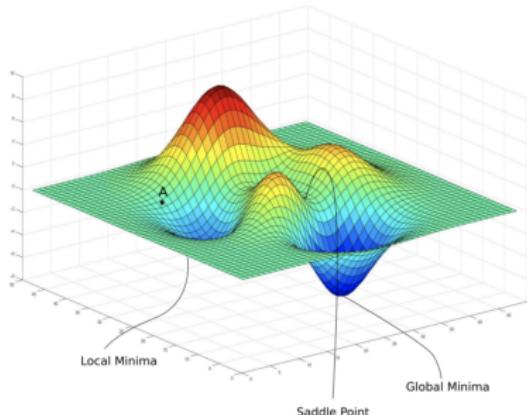
If E^* is best error approximation of $A \in \mathbb{R}^{N \times N}$, the solution of the hierarchical algorithm yields a distance/error E such that:

- Any tree: $E \leq (N/2 - 1) \times E^*$.
- Left-to-right (or right-to-left) tree: $E \leq N^c \times E^*$ where $c = \log_4 3 < 1$.

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Study of the landscape of the loss function

$$L(X, Y) = \|A - XY^T\|_F^2$$



Has been studied for:

- linear and shallow neural networks
- matrix sensing, phase retrieval, matrix completion ...

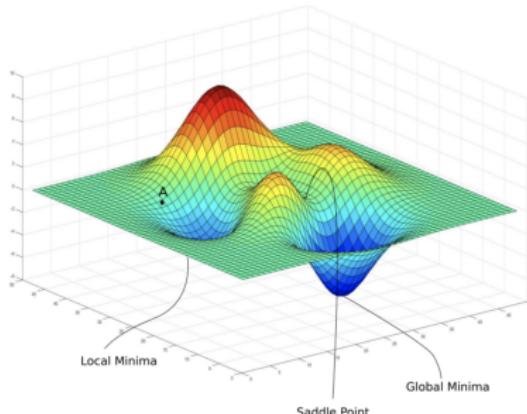
[Q. Li, Z. Zhu, G. Tang, The non-convex geometry of low-rank matrix optimization, Information and Inference, 2018]

[Z. Zhu & all. The global optimization geometry of shallow linear neural networks, JMIV, 2019]

[L. Venturi, A. S. Bandeira, J. Bruna, Spurious valleys in one-hidden-layer neural network optimization landscapes, JMLR, 2019]

Study of the landscape of the loss function

$$L(X, Y) = \|A - XY^T\|_F^2$$



Has been studied for:

- linear and shallow neural networks
- matrix sensing, phase retrieval, matrix completion ...

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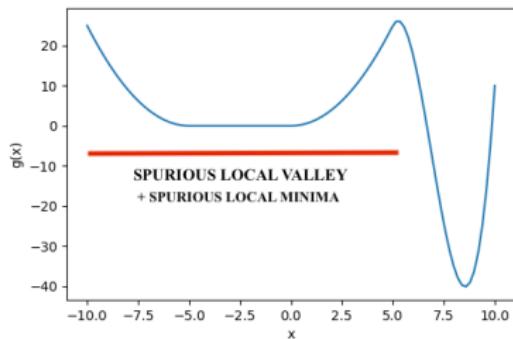
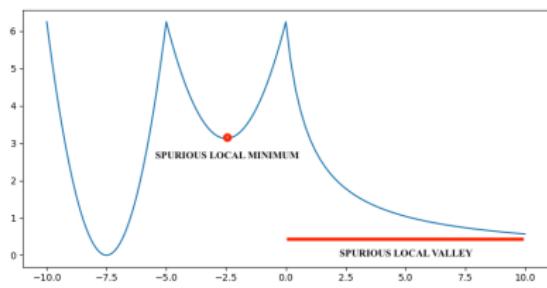
Never with **support constraints!**

Undesirable spurious objects

Example of undesirable spurious objects :

- Spurious local minima: local minima (but not global minima)
- Spurious local valleys: less known but equally troublesome

[L. Venturi, A. S. Bandeira, J. Bruna, Spurious valleys in one-hidden-layer neural network optimization landscapes, JMLR, 2019]



Example of spurious local minimum and spurious local valley. Two **undesirable objects**: may make the convergence of iterative methods difficult

Landscape of full support matrix factorization

- With unconstrained (full support) matrix factorization:

$$\underset{X \in \mathbb{R}^{m \times r}, Y \in \mathbb{R}^{n \times r}}{\text{Minimize}} \quad L(X, Y) = \|A - XY^\top\|^2$$

- The landscape of $L(X, Y)$ is *benign*:
 - No spurious local minima.¹
 - No spurious local valleys²

¹ [Z. Zhu & all. The global optimization geometry of shallow linear neural networks, JMIV, 2019]

² [L. Venturi, A. S. Bandeira, J. Bruna, Spurious valleys in one-hidden-layer neural network optimization landscapes, JMLR, 2019]

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→ Are there other instances similar to full support matrix factorization ?

¹ [Z. Zhu & all. The global optimization geometry of shallow linear neural networks, JMIV, 2019]

² [L. Venturi, A. S. Bandeira, J. Bruna, Spurious valleys in one-hidden-layer neural network optimization landscapes, JMLR, 2019]

Landscape of $L(X, Y)$ under sparsity constraints

Reminder : Fixed support matrix factorization

$$\underset{X \in \mathbb{R}^{m \times r}, Y \in \mathbb{R}^{n \times r}}{\text{Minimize}} \quad L(X, Y) = \|A - XY^\top\|^2$$

Subject to:

$$\begin{aligned} \text{supp}(X) &\subseteq S_X \\ \text{supp}(Y) &\subseteq S_Y \end{aligned}$$

Theorem (4)

If (S_X, S_Y) satisfy the condition of polynomial solvability in Theorem (2), then for all A , the landscape of $L(X, Y)$ does not contain **any** spurious local minimum and spurious local valley.

Conclusions

Take home message

For Fixed support matrix factorization (FSMF), we have:

- 1) It is **NP-hard** to solve
- 2) Easy instances with effective **direct algorithm** exists, competitive with gradient descent.
- 3) Those easy instances have **benign landscape**
- 4) Multiple factors can be dealt with by a hierarchical approach.

On-going works/perspectives

- Study the **closedness** of the set of solutions to FSMF.
- Can we enlarge the family of tractable instances of FFSMF?
- Can we improve the approximate factor in the case of *butterfly factorization*?



Available: an implementation of the algorithm in C++ via Python and Matlab wrappers **FA μ ST toolbox** at <https://faust.inria.fr/>.

To know more:

-  Q.-T. Le, E. Riccietti, and R. Gribonval (2022), Spurious Valleys, Spurious Minima and NP-hardness of Sparse Matrix Factorization With Fixed Support, *arXiv preprint*, arXiv:2112.00386.
-  L. Zheng, E. Riccietti, and R. Gribonval (2022), Efficient Identification of Butterfly Sparse Matrix Factorizations, *arXiv preprint*, arXiv:2110.01235.
-  Q.-T. Le, L. Zheng, E. Riccietti, and R. Gribonval (2022), Fast learning of fast transforms, with guarantees, *ICASSP 2022*

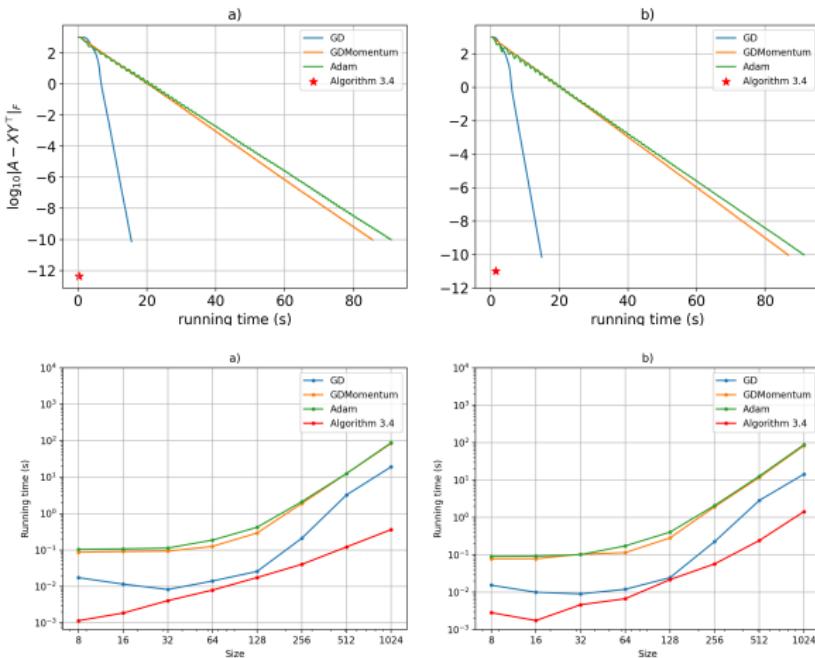
Definition: spurious local valleys

Definition (Spurious local valley - Informal)

$S \in \mathbb{R}^d$ is a spurious local valley if for all $x \in S$, there does not exist any *continuous path* connecting x and a global minimum x^* *without increasing* the loss function f .

Numerical results: 2 factors

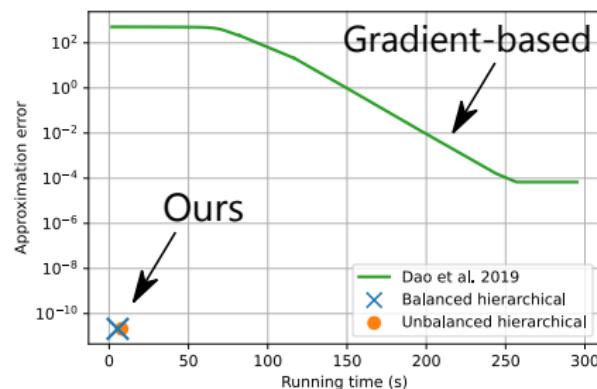
A the Hadamard matrix of size $2^J \times 2^J$, $J = 10$, two different supports



Numerical results: J factors

Approximation of the DFT matrix by a product of $J = 9$ butterfly factors.

Faster and more accurate in the
noiseless setting



Also more robust in the
noisy setting

