

Homework #1**MATH4710J — Numerical Methods**

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Problem 1.

1. Prove that \mathbb{N} , \mathbb{Z} , and \mathbb{Q} have the same number of elements.
2. Prove that $[0, 1]$ has as many elements as \mathbb{R} .
3. Prove that $[0, 1]$ has more elements than \mathbb{N} . Hint: understand Cantor's diagonal argument.

Solution.

1. **\mathbb{N} , \mathbb{Z} , and \mathbb{Q} have the same number of elements**

We want to show that the sets of natural numbers \mathbb{N} , integers \mathbb{Z} , and rational numbers \mathbb{Q} are all countably infinite, i.e., there exists a bijection between each pair.

Step 1: \mathbb{N} and \mathbb{Z} have the same cardinality.

Define a function $f : \mathbb{N} \rightarrow \mathbb{Z}$ by

$$f(n) = \begin{cases} \frac{n}{2} & \text{if } n \text{ is even} \\ -\frac{n-1}{2} & \text{if } n \text{ is odd} \end{cases}$$

This function lists the integers in the order $0, 1, -1, 2, -2, 3, -3, \dots$ as n increases. It is a bijection, so $|\mathbb{N}| = |\mathbb{Z}|$. **Step 2: \mathbb{N} and \mathbb{Q} have the same cardinality.**

Every rational number can be written as p/q with $p \in \mathbb{Z}$, $q \in \mathbb{N}$, and $\gcd(p, q) = 1$. Arrange all such pairs (p, q) in a grid and enumerate them diagonally (Cantor's diagonal argument), skipping duplicates and zero denominators. This process lists all rationals in a sequence, showing that \mathbb{Q} is countable.

Therefore, there exist bijections $\mathbb{N} \leftrightarrow \mathbb{Z} \leftrightarrow \mathbb{Q}$, so all three sets have the same number of elements.

2. **$[0, 1]$ has as many elements as \mathbb{R}**

To show that the interval $[0, 1]$ has the same cardinality as \mathbb{R} , we can construct a bijection between these two sets. Consider the function $f : (0, 1) \rightarrow \mathbb{R}$ defined by

$$f(x) = \tan\left(\pi x - \frac{\pi}{2}\right)$$

This function maps the open interval $(0, 1)$ onto the entire real line \mathbb{R} . To include the endpoints 0 and 1, we can extend this function by defining:

$$f(0) = -\infty, \quad f(1) = +\infty$$

This shows that there is a bijection between $[0, 1]$ and \mathbb{R} , hence they have the same cardinality.

3. $[0, 1]$ has more elements than \mathbb{N}

To show that the interval $[0, 1]$ has more elements than \mathbb{N} , we can use Cantor's diagonal argument. Assume, for the sake of contradiction, that there is a bijection between \mathbb{N} and $[0, 1]$. This means we can list all numbers in $[0, 1]$ as a sequence:

$$x_1, x_2, x_3, \dots$$

where each x_i is represented in its decimal expansion. Now, we construct a new number y in $[0, 1]$ by changing the i -th digit of x_i to a different digit (for example, if the i -th digit is 5, change it to 6; if it's 9, change it to 0, etc.). This new number y differs from each x_i in at least one decimal place, meaning y cannot be in the list. This contradiction shows that there is no bijection between \mathbb{N} and $[0, 1]$, hence $[0, 1]$ has more elements than \mathbb{N} . □

Problem 2.

1. Prove the Cauchy-Schwarz inequality (1.20|1.39) over the complex numbers.
2. Show that a distance is always positive.

Solution.

1. Cauchy-Schwarz Inequality:

$$|\langle x, y \rangle| \leq \|x\| \|y\|, \quad x, y \in V,$$

Proof:

$$0 \leq \langle x - \lambda y, x - \lambda y \rangle \tag{1}$$

$$= \langle x, x \rangle - \bar{\lambda} \langle y, x \rangle - \lambda \langle x, y \rangle + |\lambda|^2 \langle y, y \rangle. \tag{2}$$

Take $\lambda = \frac{\langle y, x \rangle}{\langle y, y \rangle}$ (when $y \neq 0$), we have

$$0 \leq \langle x, x \rangle - \frac{|\langle x, y \rangle|^2}{\langle y, y \rangle},$$

from which it follows that

$$|\langle x, y \rangle|^2 \leq \langle x, x \rangle \langle y, y \rangle,$$

and hence

$$|\langle x, y \rangle| \leq \|x\| \|y\|.$$

which is the Cauchy-Schwarz inequality.

2. Distance is Always Positive:

A distance function $d : X \times X \rightarrow \mathbb{R}$ on a set X must satisfy the following properties for all $x, y, z \in X$:

1. $d(x, y) = 0$ if and only if $x = y$.
2. Symmetry: $d(x, y) = d(y, x)$.
3. Triangle inequality: $d(x, z) \leq d(x, y) + d(y, z)$.

Proof

$$\forall x, y \in X, d(x, y) + d(y, x) = 2d(x, y) \geq d(x, x) = 0$$

□

Problem 3.

1. Let f be a linear map from a vector space V_1 into a vector space V_2 . Show that the dimension of V_1 is the sum of the dimensions of the kernel and of the image of f . This result is called the **rank-nullity theorem**.
2. Prove that the composition of two linear maps is a linear map.
3. Prove that the inverse of a linear map is a linear map.

Solution.

1. Let $f : V_1 \rightarrow V_2$ be a linear map. We want to show that

$$\dim(V_1) = \dim(\ker(f)) + \dim(\operatorname{im}(f)).$$

Let $\{v_1, v_2, \dots, v_k\}$ be a basis for $\ker(f)$. We can extend this basis to a basis for V_1 by adding vectors $\{v_{k+1}, v_{k+2}, \dots, v_n\}$ such that $\{v_1, v_2, \dots, v_n\}$ is a basis for V_1 . The images of the added vectors under f , $\{f(v_{k+1}), f(v_{k+2}), \dots, f(v_n)\}$, form a basis for $\operatorname{im}(f)$ (Since $\{f(v_{k+1}), f(v_{k+2}), \dots, f(v_n)\}$ can't be 0 unless $v_{k+1}, v_{k+2}, \dots, v_n$ are in $\ker(f)$). Therefore, we have

$$\dim(V_1) = n = k + (n - k) = \dim(\ker(f)) + \dim(\operatorname{im}(f)).$$

2. Let $f : V_1 \rightarrow V_2$ and $g : V_2 \rightarrow V_3$ be linear maps. We want to show that the composition $g \circ f : V_1 \rightarrow V_3$ is linear. For any $u, v \in V_1$ and scalar c , we have

$$(g \circ f)(u + v) = g(f(u + v)) = g(f(u) + f(v)) = g(f(u)) + g(f(v)) = (g \circ f)(u) + (g \circ f)(v),$$

and

$$(g \circ f)(cu) = g(f(cu)) = g(cf(u)) = cg(f(u)) = c(g \circ f)(u).$$

Thus, $g \circ f$ is linear.

3. Let $f : V_1 \rightarrow V_2$ be a bijective linear map. We want to show that the inverse map $f^{-1} : V_2 \rightarrow V_1$ is linear. For any $f(u), f(v) \in V_2$ and scalar c , we have

$$f^{-1}(f(u) + f(v)) = f^{-1}(f(u + v)) = u + v = f^{-1}(f(u)) + f^{-1}(f(v)),$$

and

$$f^{-1}(c(f(u))) = f^{-1}(f(cu)) = cu = cf^{-1}(f(u)).$$

And since f is bijective, every element in V_2 can be written as $f(u)$ for some $u \in V_1$. Thus, f^{-1} is linear.

□

Problem 4.

Intuitively a complete space has “no point missing” anywhere. In particular it means that any Cauchy sequence converges inside the space. In this exercise we show that e is not rational while we can find a Cauchy sequence of rationals converging to e .

1. Show that e is irrational.
2. Show that the sequence $(u_n)_{n \in \mathbb{N}}$ defined by $u_n = (1 + \frac{1}{n})^n$ is a Cauchy sequence converging to e .

3. Is \mathbb{Q} complete? Explain.

Solution.

1. To show that e is irrational, we can use a proof by contradiction. Assume that e is rational, i.e., $e = \frac{p}{q}$ for some integers p and q . Consider the series expansion of e :

$$e = \sum_{n=0}^{\infty} \frac{1}{n!} = 1 + 1 + \frac{1}{2!} + \frac{1}{3!} + \cdots$$

Multiplying both sides by $q!$, we get

$$q!e = q! \left(1 + 1 + \frac{1}{2!} + \frac{1}{3!} + \cdots \right) = q! + q! + \frac{q!}{2!} + \frac{q!}{3!} + \cdots$$

The left side is an integer since $e = \frac{p}{q}$ and $q!e = p(q-1)!$. However, the right side is not an integer because the terms $\frac{q!}{n!}$ for $n > q$ are not integers. This leads to a contradiction, hence e is irrational.

2. Let

$$a_n = \left(1 + \frac{1}{n}\right)^n, \quad n \geq 1,$$

and define the function

$$f(x) = x \ln\left(1 + \frac{1}{x}\right), \quad x \geq 1.$$

Notice that for integer n we have $f(n) = \ln a_n$. To show that $\{a_n\}$ is increasing, it suffices to prove that f is increasing on $[1, \infty)$.

Let $u = 1/x > 0$. Then

$$f'(x) = \ln(1+u) + x \cdot \frac{1}{1+u} \cdot \frac{du}{dx}.$$

Since $u = 1/x$, we have $du/dx = -1/x^2 = -u^2$. Substitution yields

$$f'(x) = \ln(1+u) - \frac{u}{1+u}.$$

Define $g(u) = \ln(1+u) - \frac{u}{1+u}$. Then

$$g'(u) = \frac{1}{1+u} - \frac{1}{(1+u)^2} = \frac{u}{(1+u)^2} > 0 \quad (u > 0),$$

and $g(0) = 0$. Hence $g(u) > 0$ for all $u > 0$, so $f'(x) > 0$ for $x > 0$. Therefore f is increasing, and thus $\ln a_n = f(n)$ is increasing in n , i.e. $\{a_n\}$ is monotone increasing.

Next, we provide an upper bound. Using the inequality $\ln(1+t) \leq t$ valid for $t > -1$, we obtain for $t = 1/n$:

$$\ln a_n = n \ln\left(1 + \frac{1}{n}\right) \leq n \cdot \frac{1}{n} = 1.$$

Therefore $a_n \leq e^1 = e$ for all n . Consequently, the sequence $\{a_n\}$ is monotone increasing and bounded above by e . By the monotone convergence theorem for real numbers, $\{a_n\}$ converges in \mathbb{R} .

Finally, in the metric space $(\mathbb{R}, |\cdot|)$, every convergent sequence is Cauchy. Hence $\{a_n\}$ is a Cauchy sequence.

3. The set of rational numbers \mathbb{Q} is not complete. A metric space is said to be complete if every Cauchy sequence in that space converges to a limit that is also within the same space. However, there exist Cauchy sequences of rational numbers that converge to irrational numbers, which are not in \mathbb{Q} . For example, the sequence defined by $a_n = (1 + \frac{1}{n})^n$ converges to e , which is irrational. Since $e \notin \mathbb{Q}$, this shows that \mathbb{Q} is not complete. For simpler example, consider the sequence defined by $b_n = 3.1, 3.14, 3.141, 3.1415, \dots$, which converges to π . Since π is irrational and not in \mathbb{Q} , this is another example of a Cauchy sequence in \mathbb{Q} that does not converge to a limit in \mathbb{Q} . Thus, \mathbb{Q} is not a complete metric space.

□

Problem 5.

1. Write the pseudocode for at least one the following strategy to approximate π .
 - (a) The polygons method;
 - (b) Machin's formula $\frac{\pi}{4} = 4 \arctan \frac{1}{5} - \arctan \frac{1}{239}$ and Taylor series;
2. Implement at least one of the previous algorithms in MATLAB.

Solution.

1. a) For the unit circle (radius 1), the side length of an inscribed regular n -gon is

$$c_n = 2 \sin \frac{\pi}{n}, \quad \text{so} \quad p_n = nc_n = 2n \sin \frac{\pi}{n}$$

is the inscribed polygon perimeter. The perimeter of the circumscribed regular n -gon is

$$P_n = 2n \tan \frac{\pi}{n}.$$

Thus

$$n \sin \frac{\pi}{n} \leq \pi \leq n \tan \frac{\pi}{n}.$$

If we double the number of sides, the inscribed side lengths satisfy

$$c_{2n} = 2 \sin \frac{\pi}{2n} = \sqrt{2 - 2\sqrt{1 - \frac{c_n^2}{4}}}.$$

Consequently

$$p_{2n} = 2n c_{2n}.$$

Repeating the doubling step yields arbitrarily tight bounds for π .

Algorithm 1: Polygon method to approximate π

Input : number of sides n (initially 6), number of iterations k **Output:** approximation of π

```

1 Function AlgoHw( $n, k$ ):
2    $c_n \leftarrow 2 \sin(\pi/n)$ ;
3    $p_n \leftarrow n \cdot c_n$ ;
4   for  $i \leftarrow 1$  to  $k$  do
5      $c_n \leftarrow \sqrt{2 - 2\sqrt{1 - c_n^2/4}}$ ;
6      $p_n \leftarrow 2n \cdot c_n$ ;
7      $n \leftarrow 2n$ ;
8   end for
9   return  $p_n$ 
10 end

```

b) Using Machin's formula and Taylor series to approximate π :

$$\frac{\pi}{4} = 4 \arctan \frac{1}{5} - \arctan \frac{1}{239}$$

The Taylor series expansion for $\arctan(x)$ is given by:

$$\arctan(x) = \sum_{n=0}^{\infty} \frac{(-1)^n x^{2n+1}}{2n+1}$$

To approximate π , we can truncate the series after a finite number of terms. Here is the pseudocode for the algorithm:

Algorithm 2: Machin's formula to approximate π

Input : number of terms N **Output:** approximation of π

```

1 Function AlgoHw( $N$ ):
2    $\arctan(\frac{1}{5}) \leftarrow 0$ ;
3    $\arctan(\frac{1}{239}) \leftarrow 0$ ;
4   for  $n \leftarrow 0$  to  $N - 1$  do
5      $\arctan(\frac{1}{5}) \leftarrow \arctan(\frac{1}{5}) + \frac{(-1)^n (1/5)^{2n+1}}{2n+1}$ ;
6      $\arctan(\frac{1}{239}) \leftarrow \arctan(\frac{1}{239}) + \frac{(-1)^n (1/239)^{2n+1}}{2n+1}$ ;
7   end for
8    $\pi \leftarrow 4 \cdot (4 \cdot \arctan(\frac{1}{5}) - \arctan(\frac{1}{239}))$ ;
9   return  $\pi$ 
10 end

```

2. Here is a MATLAB implementation of the polygon method to approximate π :

Listing 1: Polygon Method to Approximate Pi

```

function pi_approx = approximate_pi_polygon(n, k)
% n: initial number of sides (e.g., 6)
% k: number of iterations (doublings)
c_n = 2 * sin(pi / n); % side length of inscribed polygon
p_n = n * c_n;          % perimeter of inscribed polygon
for i = 1:k
    c_n = sqrt(2 - 2 * sqrt(1 - c_n^2 / 4));
    p_n = 2 * n * c_n;

```

```
        n = 2 * n;  
    end  
    pi_approx = p_n;  
end
```

□