

Online Learning Without Prior Information

Speaker: Yi-Shan Wu

Institute of Information Science

Academia Sinica

Taiwan

November 19, 2017

- 1 Online Convex Optimization
- 2 Online Convex Optimization with Unconstrained Domains and Losses (NIPS2016)
- 3 Online Learning Without Prior Information (COLT2017)

Online Linear Optimization

1 Setting:

- Sample at time t : (\mathbf{x}_t, y_t)
- Hypothesis : \mathbf{w}
- Let $f_t(\mathbf{w}) = \langle \mathbf{w}, \mathbf{g}_t \rangle$, where $\mathbf{g}_t \in \nabla f_t(\mathbf{w}_t)$
- Let $R(\mathbf{w}) = \frac{1}{2\eta} \|\mathbf{w}\|_2^2$

Online Linear Optimization

1 Setting:

- Sample at time t : (\mathbf{x}_t, y_t)
- Hypothesis : \mathbf{w}
- Let $f_t(\mathbf{w}) = \langle \mathbf{w}, \mathbf{g}_t \rangle$, where $\mathbf{g}_t \in \nabla f_t(\mathbf{w}_t)$
- Let $R(\mathbf{w}) = \frac{1}{2\eta} \|\mathbf{w}\|_2^2$

2 Regret:

$$\text{Regret}(\mathbf{u}) = \sum_{t=1}^T f_t(\mathbf{w}_t) - f_t(\mathbf{u})$$

Online Linear Optimization

1 Setting:

- Sample at time t : (\mathbf{x}_t, y_t)
- Hypothesis : \mathbf{w}
- Let $f_t(\mathbf{w}) = \langle \mathbf{w}, \mathbf{g}_t \rangle$, where $\mathbf{g}_t \in \nabla f_t(\mathbf{w}_t)$
- Let $R(\mathbf{w}) = \frac{1}{2\eta} \|\mathbf{w}\|_2^2$

2 Regret:

$$\text{Regret}(\mathbf{u}) = \sum_{t=1}^T f_t(\mathbf{w}_t) - f_t(\mathbf{u})$$

3 FTRL Algorithm:

$$\forall t, \mathbf{w}_t = \underset{\mathbf{w}}{\operatorname{argmin}} \sum_{i=1}^{t-1} f_i(\mathbf{w}) + R(\mathbf{w})$$

$$\text{Thus, } \mathbf{w}_{t+1} = \mathbf{w}_t - \eta \mathbf{g}_t = \mathbf{w}_t - \eta \nabla f_t(\mathbf{w}_t)$$

lemma

Let w_1, w_2, \dots be the sequence of vectors produced by this OLO algorithm, we have

$$\begin{aligned} \text{Regret}(\mathbf{u}) &= \sum_{t=1}^T f_t(\mathbf{w}_t) - f_t(\mathbf{u}) \\ &\leq R(\mathbf{u}) - R(\mathbf{w}_1) + \sum_{t=1}^T (f_t(\mathbf{w}_t) - f_t(\mathbf{w}_{t+1})) \\ &\leq \frac{1}{2\eta} \|\mathbf{u}\|_2^2 + \sum_{t=1}^T \langle \mathbf{w}_t - \mathbf{w}_{t+1}, \mathbf{g}_t \rangle \\ &= \frac{1}{2\eta} \|\mathbf{u}\|_2^2 + \sum_{t=1}^T \eta \|\mathbf{g}_t\|_2^2 \end{aligned}$$

$$\text{Regret}(\mathbf{u}) \leq \frac{1}{2\eta} \|\mathbf{u}\|_2^2 + \sum_{t=1}^T \eta \|\mathbf{g}_t\|_2^2$$

Usually, we consider $U = \{\mathbf{u} : \|\mathbf{u}\| \leq B\}$ and let L be such that

$\frac{1}{T} \sum_{t=1}^T \|\mathbf{g}_t\|_2^2 \leq L^2$, then by setting $\eta = \frac{B}{L\sqrt{2T}}$, we have

$$\text{Regret}_T(U) \leq BL\sqrt{2T}$$

Previous works and problem now

$B = \max_{\mathbf{u} \in U} \ \mathbf{u}\ $	$L_{\max} = \max_t \ z_t\ $	Regret bound
V	V	$O(BL_{\max}\sqrt{T})$
V	X	$O(BL_{\max}\sqrt{T})$
X	V	$O(\ \mathbf{u}\ \log(\ \mathbf{u}\) L_{\max} \sqrt{T})$
X	X	??

Previous works and problem now

$B = \max_{\mathbf{u} \in U} \ \mathbf{u}\ $	$L_{\max} = \max_t \ z_t\ $	Regret bound
V	V	$O(BL_{\max}\sqrt{T})$
V	X	$O(BL_{\max}\sqrt{T})$
X	V	$O(\ \mathbf{u}\ \log(\ \mathbf{u}\) L_{\max} \sqrt{T})$
X	X	??

Lower bound

For any $\epsilon > 0$,

$$\text{Regret}_T(U) \geq O(\|\mathbf{u}\| \log(\|\mathbf{u}\|) L_{\max} \sqrt{T}) + L_{\max} \exp[(\max_t \frac{\|g_t\|}{L(t)})^{1/2-\epsilon}]$$

where $L(t) = \max_{t' < t} \|g_{t'}\|$ and L_{\max} are unknown in advance.

Lower Bound with unknown L_{\max} and unbounded U

Theorem 1

For any $c, k, \epsilon > 0$, there exists a T and an adversarial strategy picking $g_t \in \mathbb{R}$ in response to $w_t \in \mathbb{R}$ such that

$$\begin{aligned} R_T(u) &= \sum_{t=1}^T g_t w_t - g_t u \\ &\geq (k + c\|u\| \log\|u\|) L_{\max} \sqrt{T} \log(L_{\max} + 1) + k L_{\max} e^{((2T)^{1/2-\epsilon})} \\ &\geq (k + c\|u\| \log\|u\|) L_{\max} \sqrt{T} \log(L_{\max} + 1) + k L_{\max} e^{[(\max_t \frac{\|g_t\|}{L(t)})^{1/2-\epsilon}]} \end{aligned}$$

for some $u \in \mathbb{R}$ where $L_{\max} = \max_{t \leq T} \|g_t\|$ and $L(t) = \max_{t' < t} \|g_{t'}\|$

Proof of lower bound

$$R_T(u) = \sum_{t=1}^T g_t w_t - g_t u$$

High level concept: for adversary, the goal is to pick a sequence of g_t s such that the algorithm suffer high regret.

- Case 1: when w_t is small ($w_t < \frac{1}{2} \exp\{T^{1/2}/4\log(2)c\}$), there is a large u

$$\Rightarrow \text{choose } g_t = -1$$

$$\Rightarrow L_{\max} = 1, \quad \max_t \frac{\|g_t\|}{L(t)} = 1$$

- Case 2: when $w_t < \frac{1}{2} \exp\{T^{1/2}/4\log(2)c\}$, let $u = 0$

$$\Rightarrow \text{choose } g_t = 2T$$

$$\Rightarrow L_{\max} = 2T, \quad \max_t \frac{\|g_t\|}{L(t)} = 2T$$

The algorithm RESCALED EXP

High level concept: by "guess-and-double" strategy

- 1 Initialize a guess L for L_{max} to $\|g_1\|$, then we can run a "known- L_{max} " algorithm.
- 2 The "known- L_{max} algorithm uses the Follow-the-Regularized-Leader(FTRL) framework.
- 3 If $\|g_t\| > 2L$, then update the guess to $\|g_t\|$.
The periods during which L is constant is called "epochs".
- 4 Need to prove that the "known- L_{max} " algorithm does not suffer too much regret when seeing a g_t that violate the assumed bound L .

The algorithm RESCALEDEXP

RESCALEDEXP

Initialize: $k \leftarrow \sqrt{2}$, $M_0 \leftarrow 0$, $w_1 \leftarrow 0$, $t_\star \leftarrow 1$ // t_\star is the start-time of the current epoch.
for $t = 1$ **to** T **do**
 Play w_t , receive subgradient $g_t \in \partial \ell_t(w_t)$.
 if $t = 1$ **then**
 $L_1 \leftarrow \|g_1\|$
 $p \leftarrow 1/L_1$
 end if
 $M_t \leftarrow \max(M_{t-1}, \|g_{t_\star:t}\|/p - \|g\|_{t_\star:t}^2)$.
 $\eta_t \leftarrow \frac{1}{k\sqrt{2(M_t + \|g\|_{t_\star:t}^2)}}$
 //Set w_{t+1} using FTRL update
 $w_{t+1} \leftarrow -\frac{g_{t_\star:t}}{\|g_{t_\star:t}\|} [\exp(\eta_t \|g_{t_\star:t}\|) - 1] = \operatorname{argmin}_w \left[\frac{\psi(w)}{\eta_t} + g_{t_\star:t} w \right]$
 if $\|g_t\| > 2L_t$ **then**
 //Begin a new epoch: update L and restart FTRL
 $L_{t+1} \leftarrow \|g_t\|$
 $p \leftarrow 1/L_{t+1}$
 $t_\star \leftarrow t + 1$
 $M_t \leftarrow 0$
 $w_{t+1} \leftarrow 0$
 else
 $L_{t+1} \leftarrow L_t$

The algorithm RESCALEDEXP

Theorem 2

Let $M_{\max} = \max_t M_t$. Then if $L_{\max} = \max_t \|g_t\|$ and $L(t) = \max_{t' < t} \|g_{t'}\|$, RESCALEDEXP achieves regret:

$$R_T(\mathbf{u}) \leq O\left(L_{\max} \log\left(\frac{L_{\max}}{L_1}\right) [(\|\mathbf{u}\| \log(\|\mathbf{u}\|) + 2)\sqrt{T} + \exp(8 \max_t \frac{\|g_t\|^2}{L(t)^2})]\right)$$

Improved version-lower bound

For any $\gamma \in (1/2, 1]$, $k > 0$, $T_0 > 0$, and any online optimization algorithm picking $w_t \in \mathbb{R}$, there exists a $T > T_0$, a $u \in \mathbb{R}$, and a sequence $g_1, \dots, g_T \in \mathbb{R}$ with $\|g_t\| \leq \max(1, 18\gamma(4k)^{1/\gamma}(t-1)^{1-1/2\gamma})$ on which the regret is:

$$\begin{aligned} R_T(\mathbf{u}) &= \sum_{t=1}^T g_t w_t - g_t \mathbf{u} \\ &\geq k \|\mathbf{u}\| L_{\max} \log^\gamma(T \|\mathbf{u}\| + 1) \sqrt{T} \\ &\quad + \max_{t \leq T} L_{\max} \frac{L_{t-1}^2}{\|g\|_{1:t-1}^2} \exp\left[\frac{1}{2} \left(\frac{L_t/L_{t-1}}{288\gamma k^2}\right)^{1/2\gamma-1}\right] \end{aligned}$$

Improved version: lower bound

First dimension tradeoff

Fix $\gamma = 1$, for $k > 0$ there is tradeoff between

$$k\|\mathbf{u}\|L_{\max}\log(T\|\mathbf{u}\|)\sqrt{T}$$

and

$$\exp\left[\left(\frac{L_t/L_{t-1}}{k^2}\right)\right]$$

Second dimension tradeoff

Fix k , for $\gamma \in (1/2, 1]$ there is tradeoff between

$$\|\mathbf{u}\|L_{\max}\log^{\gamma}(T\|\mathbf{u}\|)\sqrt{T}$$

and

$$\exp\left[(L_t/\gamma L_{t-1})^{1/(2\gamma-1)}\right]$$

Improved version: design algorithm

Still uses the FTRL framework with different regularizer.

$$\mathbf{w}_{t+1} = \operatorname{argmin}_{\mathbf{w} \in W} \psi_t(\mathbf{w}) + \sum_{t'=1}^t f_{t'}(\mathbf{w})$$

where

$$\psi_t(\mathbf{w}) = \frac{k}{a_t \eta_t} \psi(a_t \mathbf{w})$$

$$\frac{1}{\eta_0^2} = 0, \quad \frac{1}{\eta_t^2} = \max\left(\frac{1}{\eta_{t-1}^2} + 2\|g_t\|_*^2, L_t\|g_{1:t}\|_*\right)$$

$$a_1 = \frac{1}{(L_1 \eta_1)^2}, \quad a_t = \max\left(a_{t-1}, \frac{1}{(L_t \eta_t)^2}\right)$$

a_t and η_t are carefully chosen functions of observed gradients g_1, \dots, g_t that guarantee the desired asymptotics in the regret bound

Improved version: adaptive regularizers

Definition 1

A convex function $f : W \rightarrow \mathbb{R}$ is σ -strongly convex with respect to a norm $\|\cdot\|$ if for all $x, y \in W$ and $g \in \partial f(x)$ we have

$$f(y) \geq f(x) + g \cdot (y - x) + \frac{\min(\sigma(x), \sigma(y))}{2} \|x - y\|^2$$

Definition 2

Let W be a closed convex subset of a vector space s.t. $0 \in W$. Any differentiable function $\psi : W \rightarrow \mathbb{R}$ that satisfies the following conditions:

- ① $\psi(0) = 0$
- ② $\psi(x)$ is σ -strongly-convex with respect to some norm $\|\cdot\|$ for some $\sigma : W \rightarrow \mathbb{R}$ s.t. $\|x\| \geq \|y\|$ implies $\sigma(x) \leq \sigma(y)$
- ③ For any C , there exists a B s.t. $\psi(x)\sigma(x) \geq C$ for all $\|x\| \geq B$

is called a $(\sigma, \|\cdot\|)$ -adaptive regularizer.

Improved version: adaptive regularizers

Define

$$h(\mathbf{w}) = \psi(\mathbf{w})\sigma(\mathbf{w}), \quad h^{-1}(x) = \max_{h(\mathbf{w}) \leq x} \|\mathbf{w}\|$$

$$D = \max_t \frac{L_{t-1}^2}{(\|g\|_*^2)_{1:t-1}} h^{-1}\left(\frac{5L_t}{k^2 L_{t-1}}\right)$$

Theorem 3

Suppose ψ is a $(\sigma, \|\cdot\|)$ -adaptive regularizer and g_1, \dots, g_T is some arbitrary sequence of subgradients, then FTRL with regularizer ψ_t achieves regret

$$R_T(u) \leq kL_{\max} \frac{\psi(2uT)}{\sqrt{2T}} + 2L_{\max} D + \frac{45L_{\max}}{\sigma_{\min}}$$

Goal: Choose a $h(x) = \psi(x)\sigma(x)$ s.t. $h^{-1}(x) \approx \exp(x)$
and $\psi(2uT)/\sqrt{2T} = O(\|u\|\sqrt{T}\log(T\|u\| + 1))$

Improved version: γ -optimal

Theorem 4

If ψ is an $(\sigma, \|\cdot\|)$ -adaptive regularizer s.t.

$$\psi(x)\sigma(x) \geq \Omega(\gamma \log^{2\gamma-1}(\|x\|)) \quad (1)$$

$$\psi(x) \leq O(\|x\| \log^\gamma(\|x\| + 1)) \quad (2)$$

Then for any $k \geq 1$, FTRL with regularizers $\psi_t(\mathbf{w}) = \frac{k}{a_t \eta_t} \psi(a_t \mathbf{w})$ yields regret

$$\begin{aligned} R_T(u) &\leq O[k L_{\max} \sqrt{T} \|u\| \log^\gamma(T \|u\| + 1)] \\ &\quad + \max_t \frac{L_{\max} L_{t-1}^2}{\|g\|_{1:t-1}^2} \exp[O((\frac{L_t}{k^2 \gamma L_{t-1}})^{1/2\gamma-1})] \end{aligned}$$

We call regularizers that satisfy these conditions γ -optimal.

Note that this actually match the lower bound for all $\gamma \in (1/2, 1]$.

Improved version: choose adaptive regularizers

Proposition 1

Let $\phi : \mathbb{R}^+ \rightarrow \mathbb{R}$ be a three-times differentiable function that satisfies

- ① $\phi(0) = 0$
- ② $\phi'(x) \geq 0$
- ③ $\phi''(x) \geq 0$
- ④ $\phi'''(x) \leq 0$
- ⑤ $\lim_{x \rightarrow \infty} \phi(x)\phi''(x) = \infty$

Then $\psi(\mathbf{w}) = \phi(\|\mathbf{w}\|)$ is a $(\phi''(\|\cdot\|), \|\cdot\|)$ -adaptive regularizer.

Improved version: 1-optimal adaptive regularizer

Proposition 2

Let $\phi(x) = (x + 1)\log(x + 1) - x$. Then $\psi(w) = \phi(\|w\|)$ is a 1-optimal, $(\phi''(\|\cdot\|), \|\cdot\|)$ -adaptive regularizer (pf):

- ① $\phi(0) = 0$
- ② $\phi'(x) = \log(x + 1)$
- ③ $\phi''(x) = \frac{1}{x+1}$
- ④ $\phi'''(x) = -\frac{1}{(x+1)^2}$
- ⑤ $\phi(x)\phi''(x) = \log(x + 1) - \frac{x}{x+1}$

Thus, FTRL with regularizers

$\psi_t(w) = \frac{k}{\eta_t a_t} ((\|w\| + 1)\log(\|w\| + 1) - \|w\|)$ achieves the regret:

$$k\|\mathbf{u}\|L_{\max}\log(T\|\mathbf{u}\| + 1)\sqrt{T} + L_{\max}\max_{t \leq T} \frac{L_{t-1}^2}{\|\mathbf{g}\|_{1:t-1}^2} \exp\left[\left(\frac{5L_t/L_{t-1}}{k^2}\right)\right]$$

Improved version: γ -optimal adaptive regularizer

Proposition 3

Given $\gamma \in (1/2, 1]$, set $\phi(x) = \int_0^x \log^\gamma(z+1)dz$. Then $\psi(w) = \phi(\|w\|)$ is a γ -optimal, $(\phi''(\|\cdot\|), \|\cdot\|)$ -adaptive regularizer

Thus, FTRL with regularizers $\psi_t(w) = \int_0^{\|w\|} \log^\gamma(z+1)dz$ achieves the regret:

$$R_T(u) \leq O[kL_{\max}\sqrt{T}\|u\|\log^\gamma(T\|u\|+1)] \\ + \max_t \frac{L_{\max}L_{t-1}^2}{\|g\|_{1:t-1}^2} \exp[O((\frac{L_t}{k^2\gamma L_{t-1}})^{1/2\gamma-1})]$$

and the update rule:

$$w_{t+1} = -\frac{g_{1:t}}{a_t\|g_{1:t}\|} [\exp((\eta_t\|g_{1:t}\|/k)^{1/\gamma}) - 1]$$

The algorithm FREEREX

FREEREX, for $\gamma = 1$

Algorithm 1 FREEREX

Input: k .

Initialize: $\frac{1}{\eta_0^2} \leftarrow 0, a_0 \leftarrow 0, w_1 \leftarrow 0, L_0 \leftarrow 0, \psi(w) = (\|w\| + 1) \log(\|w\| + 1) - \|w\|$.

for $t = 1$ **to** T **do**

 Play w_t , receive subgradient $g_t \in \partial \ell_t(w_t)$.

$L_t \leftarrow \max(L_{t-1}, \|g_t\|)$.

$\frac{1}{\eta_t^2} \leftarrow \max\left(\frac{1}{\eta_{t-1}^2} + 2\|g_t\|^2, L_t\|g_{1:t}\|\right)$.

$a_t \leftarrow \max(a_{t-1}, 1/(L_t\eta_t)^2)$.

 //Set w_{t+1} using FTRL update

$w_{t+1} \leftarrow -\frac{g_{1:t}}{a_t\|g_{1:t}\|} \left[\exp\left(\frac{\eta_t\|g_{1:t}\|}{k}\right) - 1 \right] // = \operatorname{argmin}_w \left[\frac{k\psi(a_t w)}{a_t \eta_t} + g_{1:t} w \right]$

end for