# Medium Access Control

EE450: Introduction to Computer Networks

Professor A. Zahid

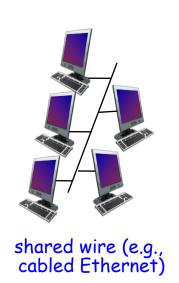
### Medium Access Control

- Single shared broadcast channel
- Two or more simultaneous transmissions by nodes: interference
  - collision if node receives two or more signals at the same time

#### Multiple Access Protocol

- Distributed algorithm that determines how nodes share channel, i.e., determine when node can transmit
- Communication about channel sharing must use channel itself!

# Multiple Access Links







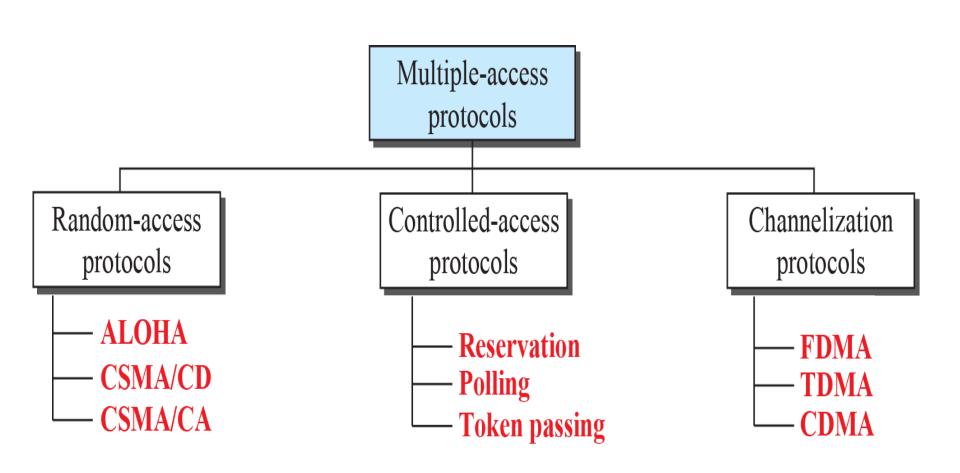


shared RF (satellite)

humans at a cocktail party (shared air, acoustical)

- Old-fashioned Ethernet
- Upstream HFC (In Cable Access Networks)
- · Wi-Fi: 802.11 wireless LAN

# Taxonomy of MAC Protocols

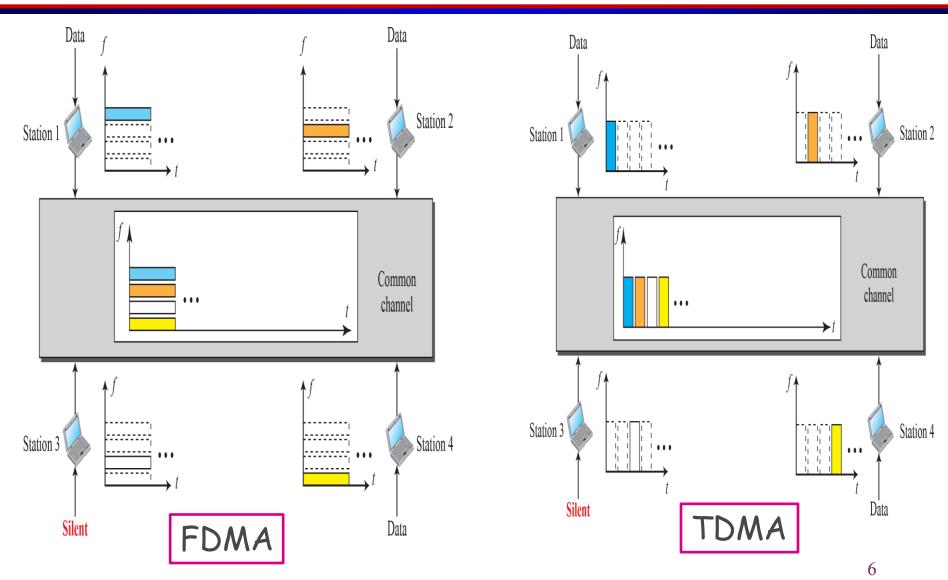


## Classifications of MAC Protocols

#### Three broad classes:

- Channel Partitioning
  - Divide channel into smaller "pieces" (time slots, frequency, code) for example TDMA, FDMA or CDMA
  - Allocate a piece to each node for exclusive use
- Random Access
  - Channel not divided, allow collisions. Examples: ALOHA, CSMA/CD, CSMA/CA
  - "Recover" from collisions for example via delayed retransmissions
- "Taking turns"
  - Nodes take turns, but nodes with more to send can take longer turns. Examples: Polling, Token Passing

# Channel Partitioning: FDMA/TDMA



## IEEE802 Standards for LANs

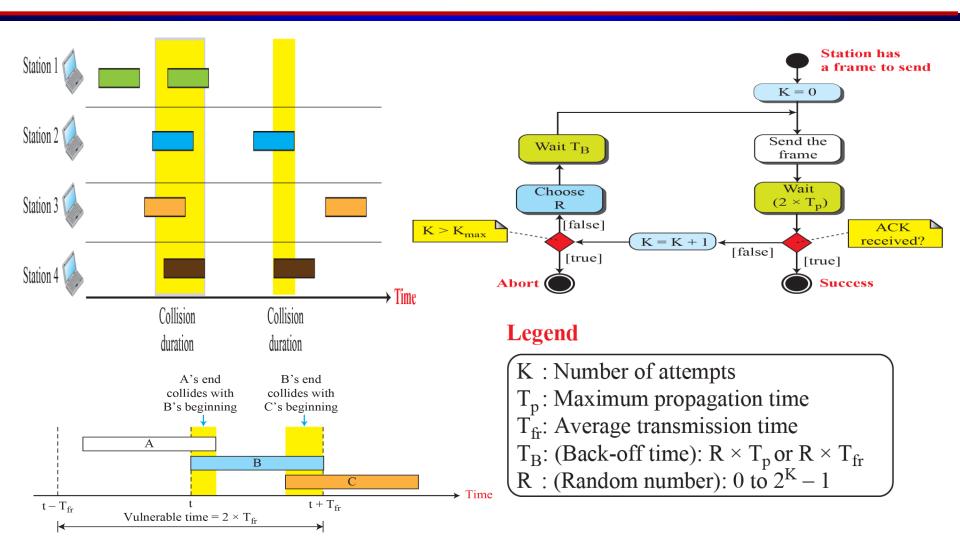
LLC: Logical link control MAC: Media access control

Upper layers		Upper layers			
		LLC			
Data link layer		Ethernet MAC	Token Ring MAC	Token Bus MAC	•••
Physical layer		Ethernet physical layers (several)	Token Ring physical layer	Token Bus physical layer	•••
ransmission medium	Transmission medium				

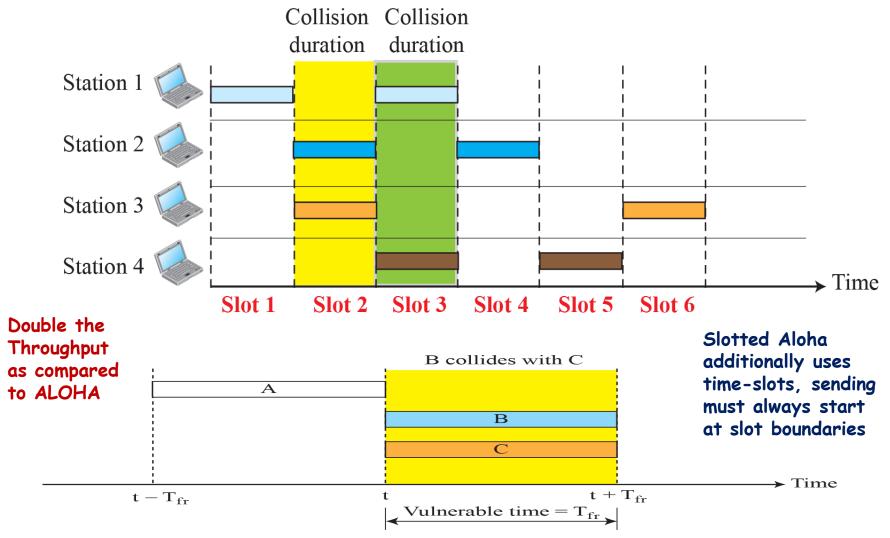
OSI or Internet model

**IEEE Standard** 

## Random Access Protocols: ALOHA



## Slotted ALOHA



# IEEE802.3 (Based on Ethernet) "Carrier Sense Multiple Access"

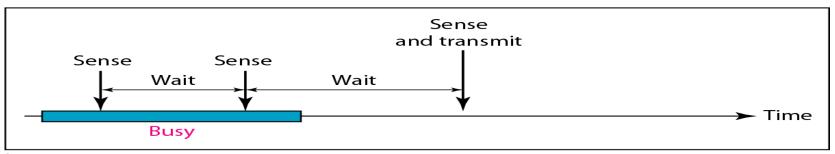
# Carrier Sense Multiple Access

- CSMA/CD: Carrier sense, multiple access with collision detection
  - collisions detected within short time
  - colliding transmissions aborted, reducing waste
  - Persistent, non-persistent and P-persistent retransmission
- Collision Detection:
  - On baseband bus, collision produces much higher signal voltage than transmitted signal
  - For twisted pair (Hub-topology) activity on more than one port is collision

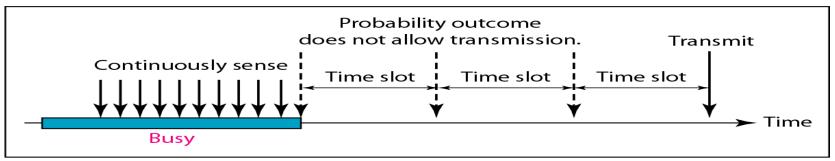
# Behavior of Three Persistent Scenarios



a. 1-persistent



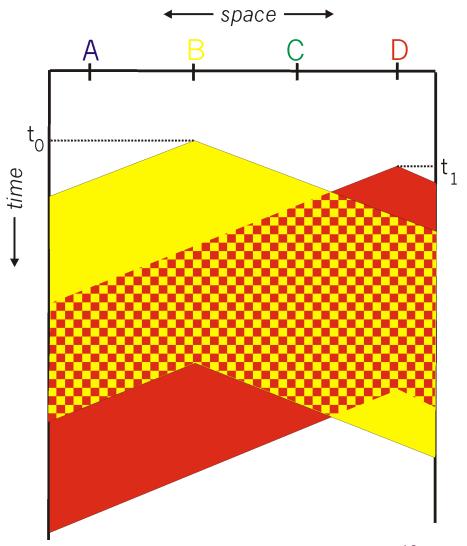
b. Nonpersistent



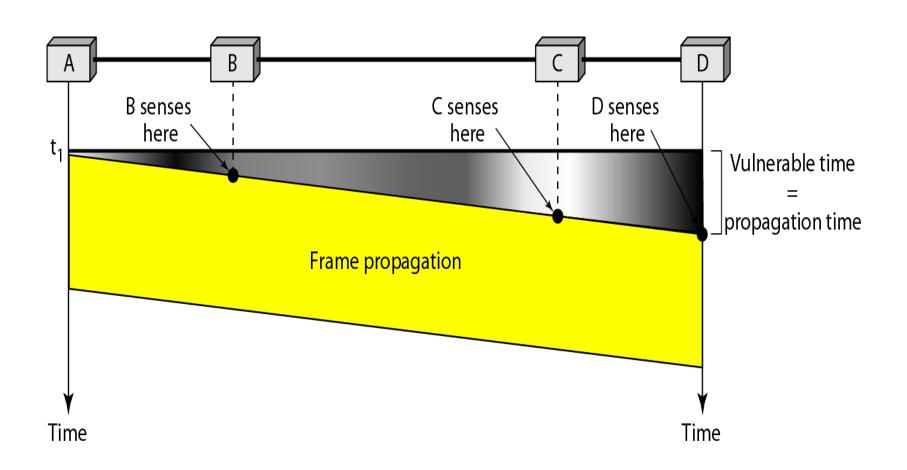
c. p-persistent

## Collisions in CSMA/CD

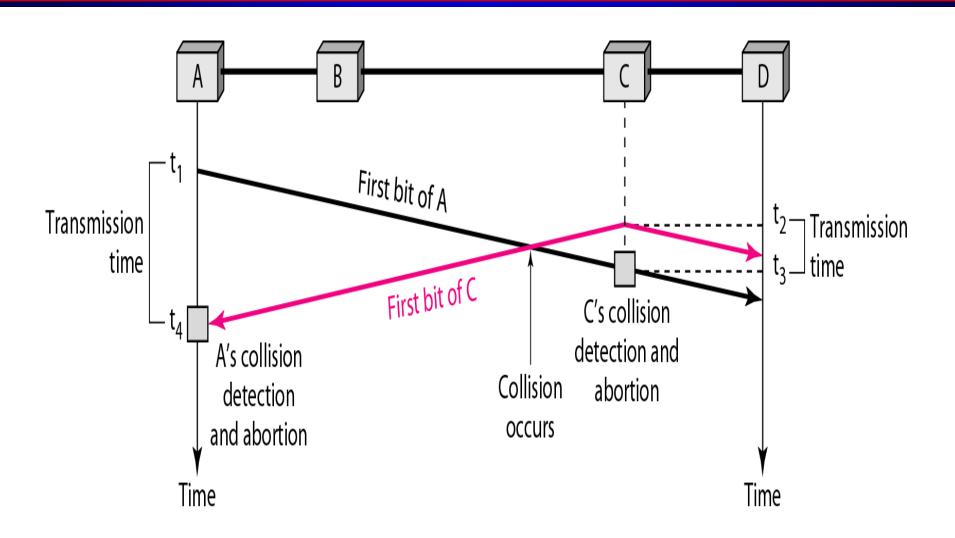
- Collisions can still occur: propagation delay means two nodes may not hear each other's transmission
- When collision occur, entire frame is wasted
- Collision is detected by comparing transmitted and received signal strengths (Hard to do in WLANs, TBD)



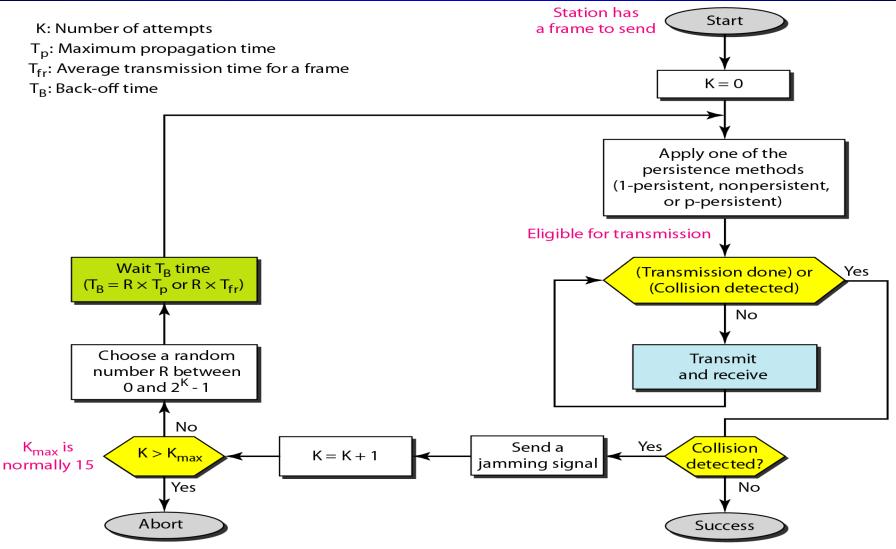
## Vulnerable Time in CSMA



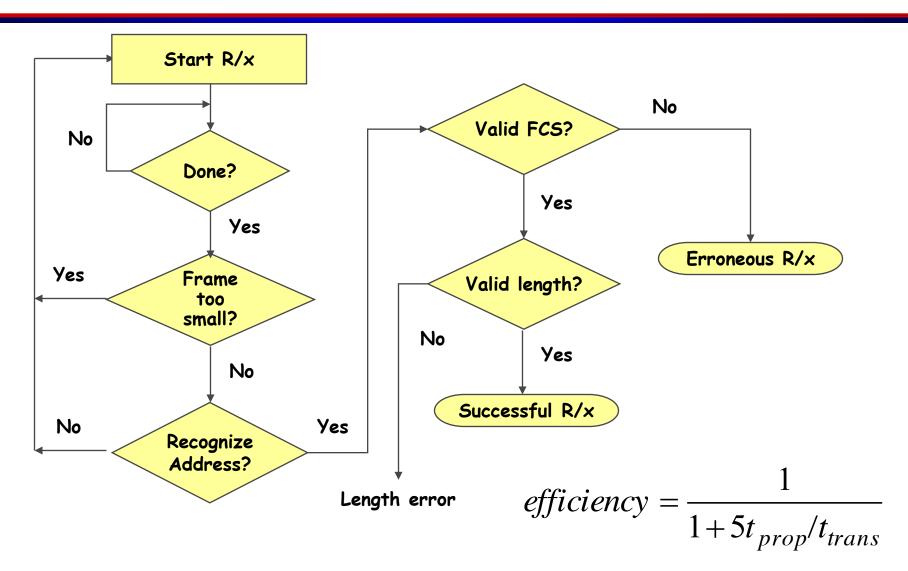
## Collision Detection



## Flow Chart of CSMA/CD



## Receive Process in IEEE802.3

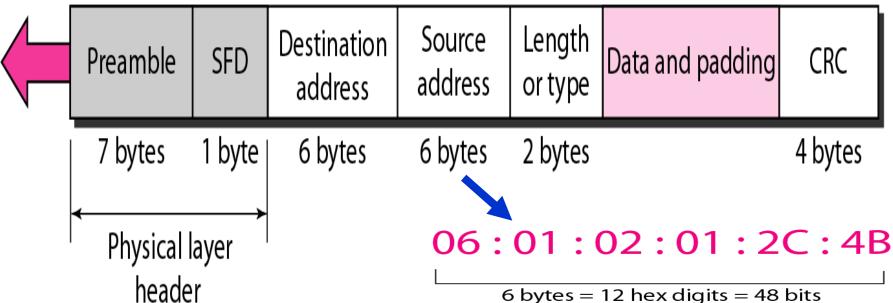


## IEEE802.3 MAC Frame

Preamble: 56 bits of alternating 1s and 0s.

SFD: Start frame delimiter, flag (10101011)

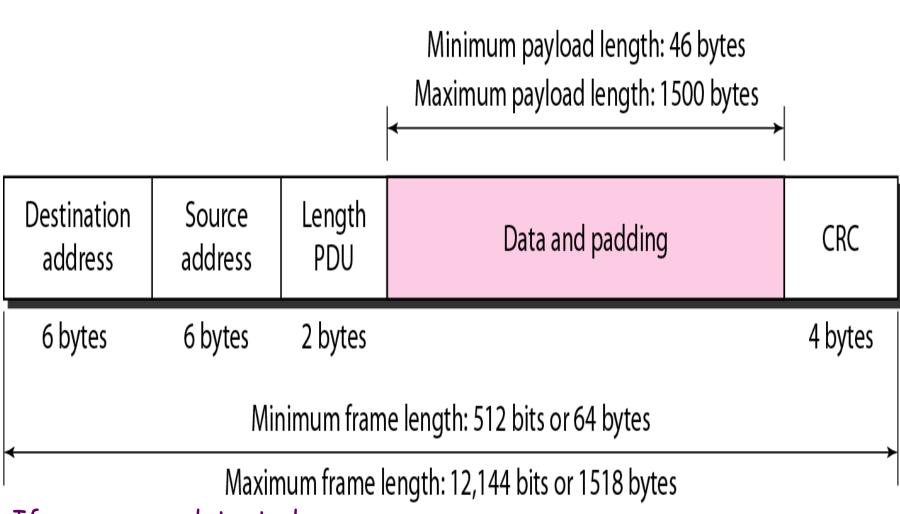




MAC address is burned in NIC ROM (sometimes software settable)

Type: Indicate Network Layer Protocol (mostly IP)

# IEEE802.3 Frame Length Limits



If errors are detected, Frame is dropped

#### MAC Addresses

Source and destination MAC addresses.
 These are the hardware addresses. They are 48-bits long each

```
Ethernet MAC Address

XX XX XX XX XX bytes

Vendor Part Vendor Assigned
24 bits 24 bits
```

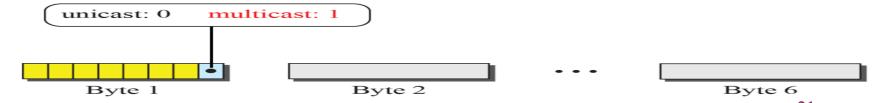
IEEE Organizationally Unique Identifier (OUI)

- allows vendor to build hardware with unique addresses

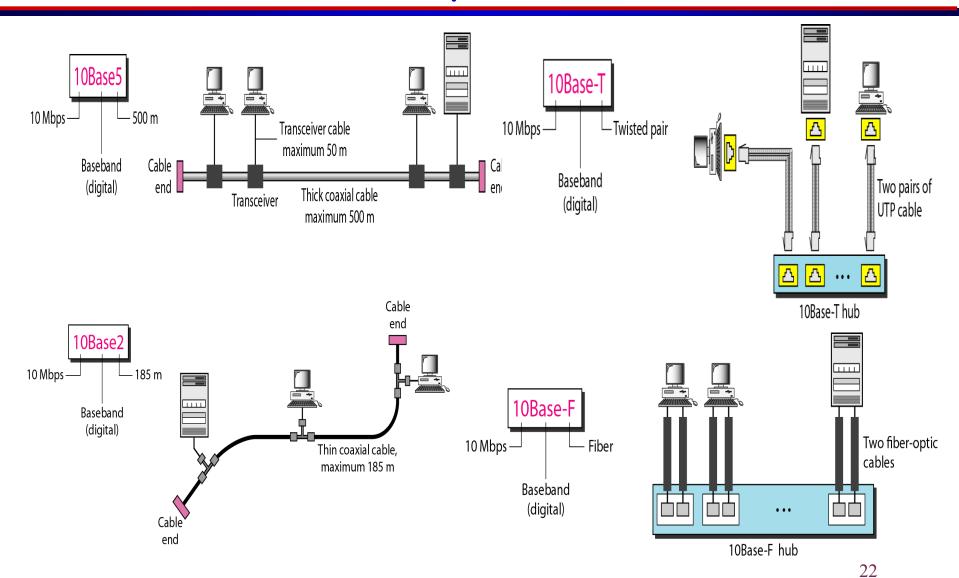
http://standards.ieee.org/regauth/oui/ http://www.cavebear.com/CaveBear/Ethernet/

# Types of MAC Addresses

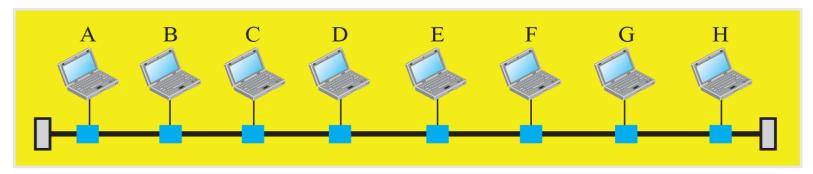
- Unicast: one interface to one interface
- Broadcast: all 1's destination address
  means that every attached interface to a
  LAN should read the frame.
  - MAC Address: FF:FF:FF:FF:FF
- Multicast: an interface can be configured to read frames sent to one or more multicast addresses.



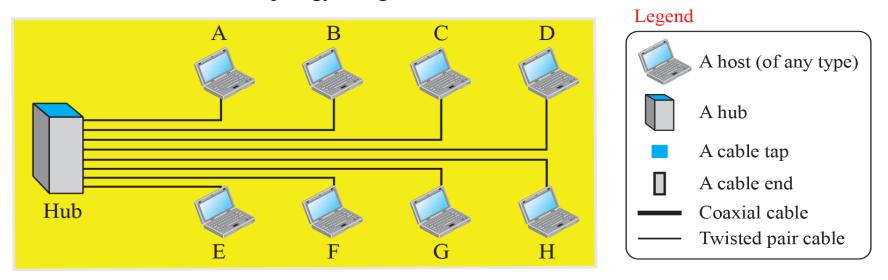
# 10Base? Implementations



# Shared Ethernet Implementations

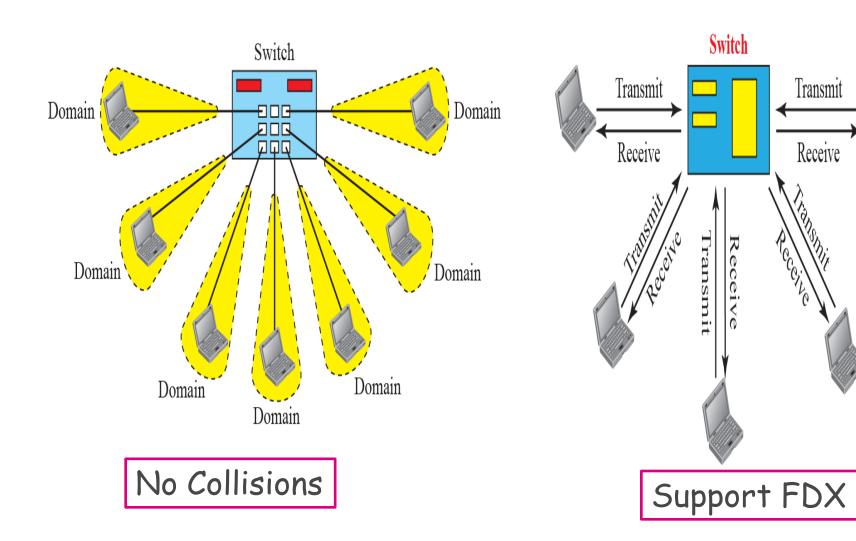


a. A LAN with a bus topology using a coaxial cable



b. A LAN with a star topology using a hub

# Switched Ethernet



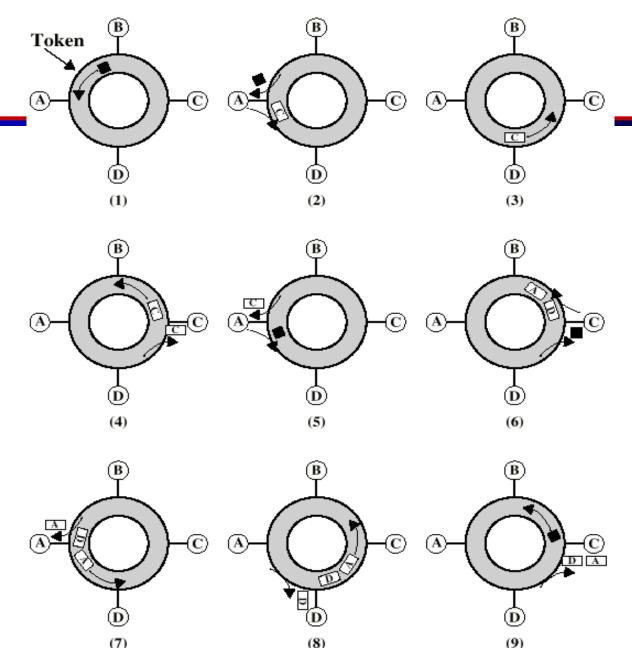
# IEEE802.5 (Based on IBM) "Token-Passing Rings"

# IEEE802.5 Token-Passing Rings

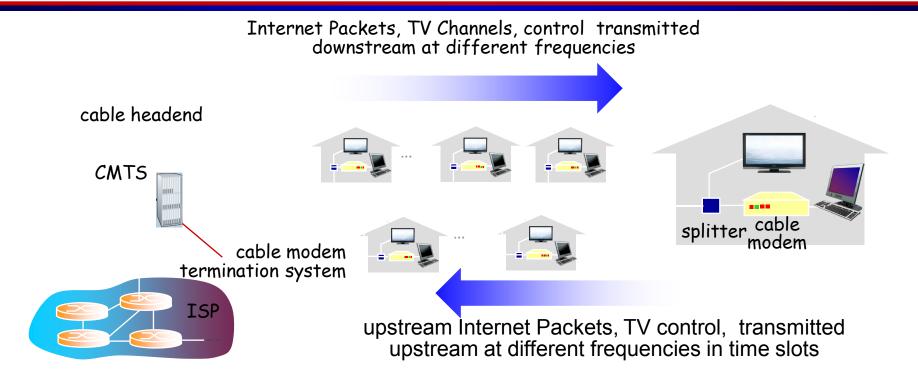
- Frames flow in one direction
- Special bit pattern (token) rotates around ring. The token is 24-bit long
- Node having a frame to transmit must capture token first
- Node must release token after done transmitting
- Node remove frame when it comes back around
- Stations get round-robin service

# Token Ring

- Supports 4Mbps (UTP), 16 Mbps (STP) and 100 Mbps (Fiber)
- Token Holding Timer
   ~ 10 msec which
   limits the frame
   length ~ 4500 Bytes
   for the 4Mbps Rings
   and ~ 18,000 Bytes
   for the faster Rings



#### Cable Access Network



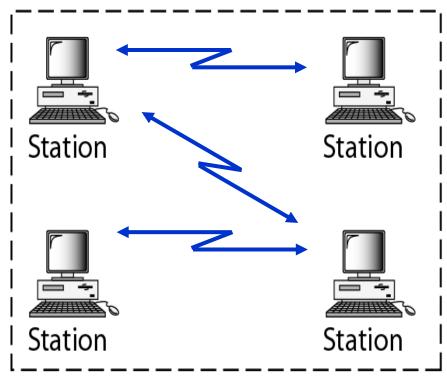
- multiple 40Mbps downstream (broadcast) channels
  - single CMTS transmits into channels
- multiple 30 Mbps upstream channels
  - multiple access: all users contend for certain upstream channel time slots (others assigned)

# IEEE 802.11 Wireless LANs (Wi-Fi)

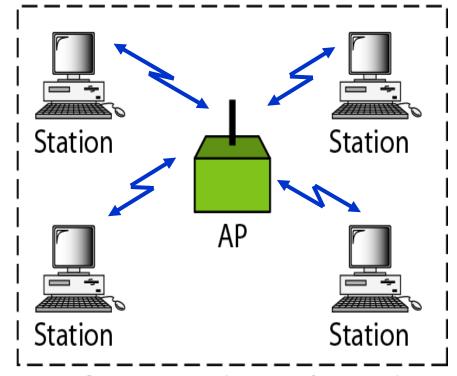
## Ad-hoc vs. Infrastructure WLANs

**BSS**: Basic service set

AP: Access point

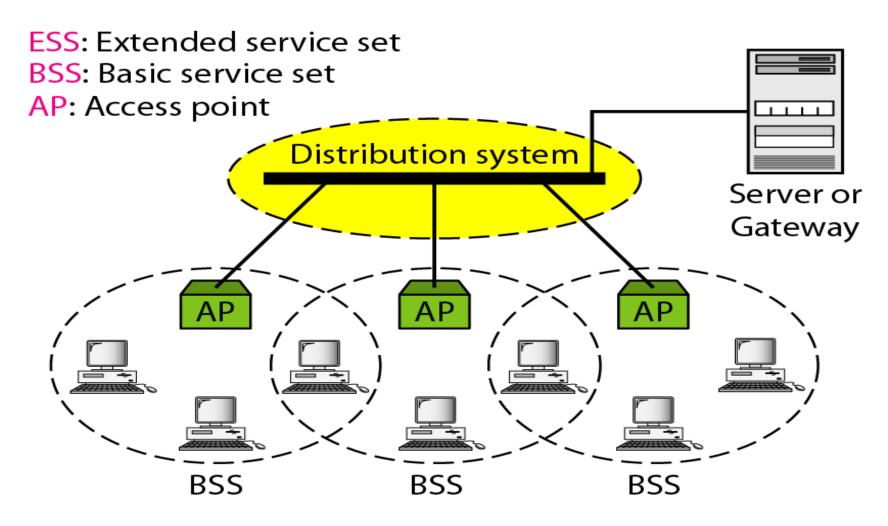


Ad hoc network (BSS without an AP)

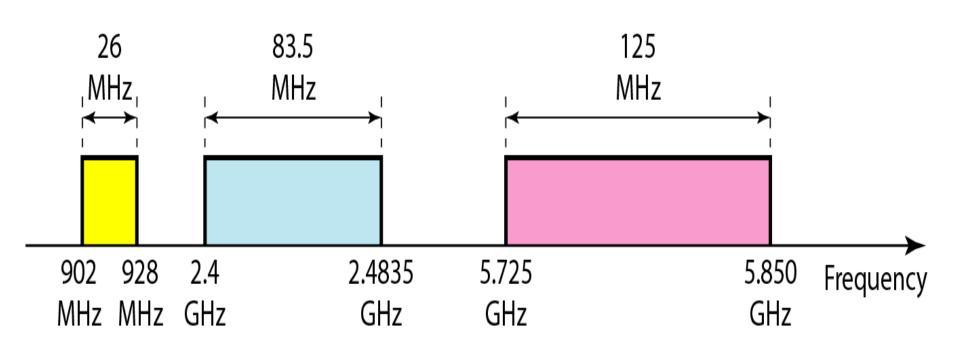


Infrastructure (BSS with an AP)

## Extended Service Sets



# Unregulated Band (ISM)



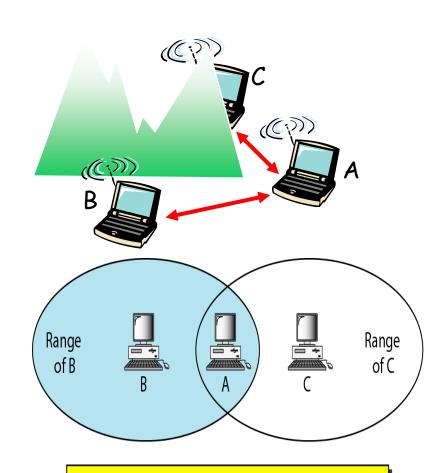
ISM: Industrial, Scientific and Medical band

## Wireless Link Characteristics

- Differences from wired link ....
  - Decreased Signal Strength: Radio signal attenuates as it propagates through matter (path loss)
  - Interference from other sources: standardized wireless network frequencies (e.g., 2.4 GHz) shared by other devices (e.g., phone); devices (motors) interfere as well
  - Multipath propagation: Radio signal reflects off objects ground, arriving ad destination at slightly different times

## Hidden Terminal Problem

- Hidden terminal problem
  - B, A hear each other
  - C, A hear each other
  - B, C can not hear each other
  - Means B, C unaware of their interference at A

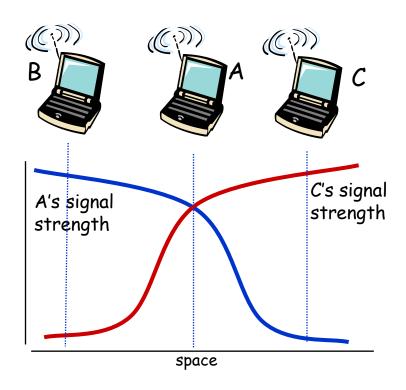


B and C are hidden from each other with respect to A.

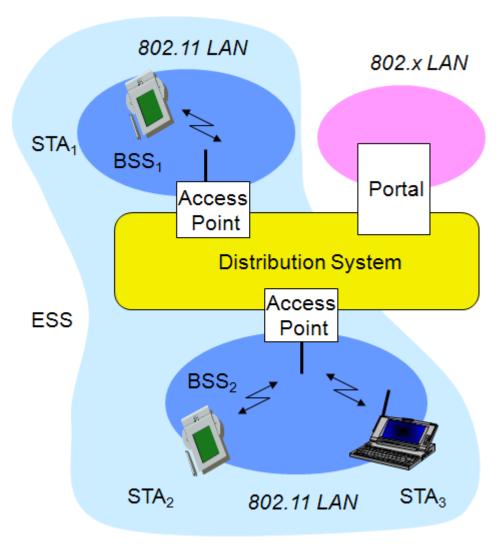
# Signal Fading

## Signal fading:

- A, B can hear each other
- A, C can hear each other
- B, C can not hear each other interfering at A
- Signal losses its strength as distance increases

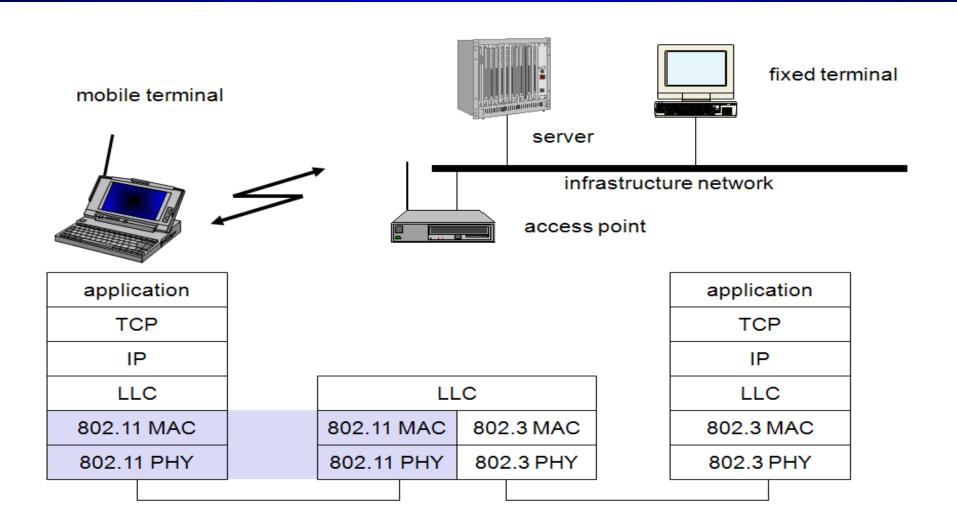


## 802.11 Infrastructure Network



- Station (STA)
  - terminal with access mechanisms to the wireless medium and radio contact to the access point
- Basic Service Set (BSS)
  - group of stations using the same radio frequency
- Access Point
  - station integrated into the wireless LAN and the distribution system
- •Portal (Bridge/Router)
  - to other (wired) networks
- Distribution System
  - interconnection network to form one logical network (EES: Extended Service Set) based on several BSS

## 802.11 in the TCP/IP Stack



## IEEE 802.11 Wireless LAN

### 802.11b

- 2.4-5 GHz unlicensed spectrum
- up to 11 Mbps
- direct sequence spread spectrum (DSSS) in physical layer
  - all hosts use same chipping code

### 802.11a

- 5-6 GHz range
- up to 54 Mbps

### 802.11g

- 2.4-5 GHz range
- up to 54 Mbps

### 802.11n: multiple

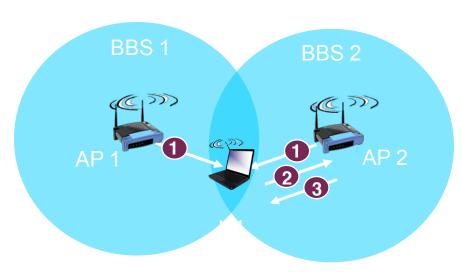
#### antennae

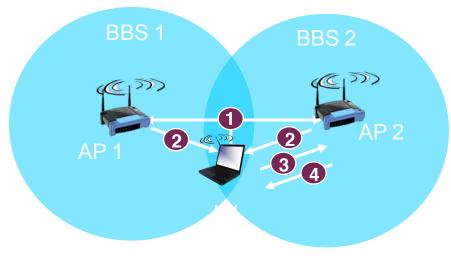
- 2.4-5 GHz range
- up to 200 Mbps
- all use CSMA/CA for multiple access
- \* all have base-station and ad-hoc network versions

## Channel Association

- 802.11b: 2.4GHz-2.485GHz spectrum divided into
   11 channels at different frequencies
  - AP admin chooses frequency for AP
  - Interference possible: channel can be same as that chosen by neighboring AP!
- Host: must associate with an AP
  - Scans channels, listening for Beacon frames containing AP's name (SSID) and MAC address
  - Selects AP to associate with
  - May perform authentication
  - Run DHCP to get IP address in AP's subnet

## 802.11: Passive/Active scanning





### passive scanning:

- (1) beacon frames sent from APs
- (2)association Request frame sent: H1 to selected AP
- (3)association Response frame sent from selected AP to H1

#### active scanning:

- (1) Probe Request frame broadcast from H1
- (2) Probe Response frames sent from APs
- (3) Association Request frame sent: H1 to selected AP
- (4) Association Response frame sent from selected AP to H1

### IEEE802.11 MAC Protocol

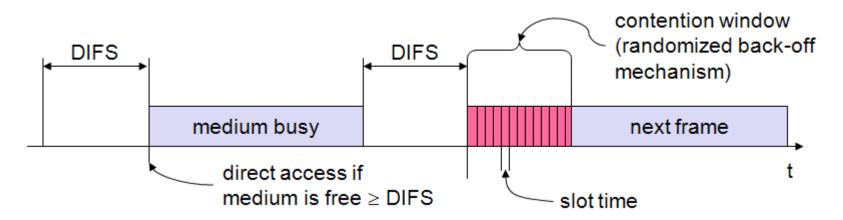
- Avoid collisions: 2+ nodes transmitting at same time
- 802.11: CSMA sense before transmitting
  - Don't collide with other transmissions
- 802.11: No collision detection!
  - Difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
  - Can't sense all collisions in any case: hidden terminal, fading
  - Goal: avoid collisions: CSMA/C(ollision)A(voidance)

## 802.11 MAC Procedures

- Traffic services
  - Asynchronous Data Service (mandatory) DCF
  - Time-Bounded Service (optional) PCF
- Access methods
  - DCF CSMA/CA (mandatory)
    - collision avoidance via randomized back-off mechanism
    - ACKs for data frames (not for broadcasts)
  - DCF w/ RTS/CTS (optional)
    - avoids hidden terminal problem
  - PCF (optional)
    - access point polls terminals according to a list

#### Distributed Coordination Function

## DCF: CSMA/CA



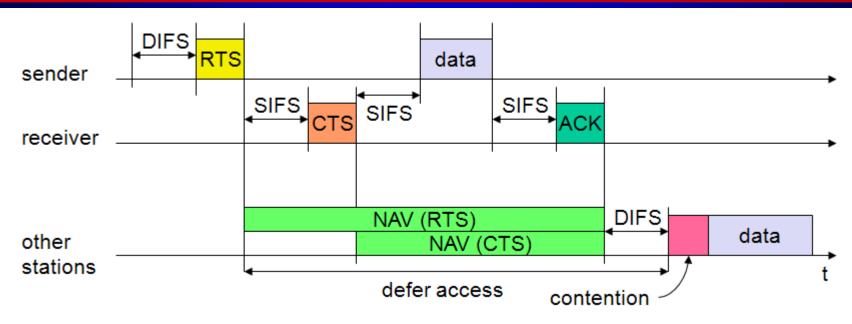
- station ready to send starts sensing the medium (Carrier Sense based on CCA, Clear Channel Assessment)
- if the medium is free for the duration of an Inter-Frame Space (IFS), the station can start sending (IFS depends on service type)
- if the medium is busy, the station has to wait for a free IFS, then the station must additionally wait a random back-off time (collision avoidance, multiple of slot-time)
- if another station occupies the medium during the back-off time of the station, the back-off timer stops (fairness)

## Avoiding Collisions

- Idea: allow sender to "reserve" channel rather than random access of data frames: avoid collisions of long data frames
- Sender first transmits small request-to-send (RTS) frames to BS using CSMA
  - RTSs may still collide with each other (but they're short)
- BS broadcasts clear-to-send CTS in response to RTS
- CTS heard by all nodes
  - Sender transmits data frame
  - Other stations defer transmissions

avoid data frame collisions completely using small reservation packets!

## DCF w/RTS & CTS



- Station send RTS with reservation parameter (amount of time the data frame needs the medium) after waiting for DIFS
- Acknowledgement via CTS after SIFS by receiver (if ready to receive)
- Sender can now send data at once, acknowledgement via ACK
- Other stations store medium reservations distributed via RTS and CTS

# Collision Avoidance using RTS/CTS

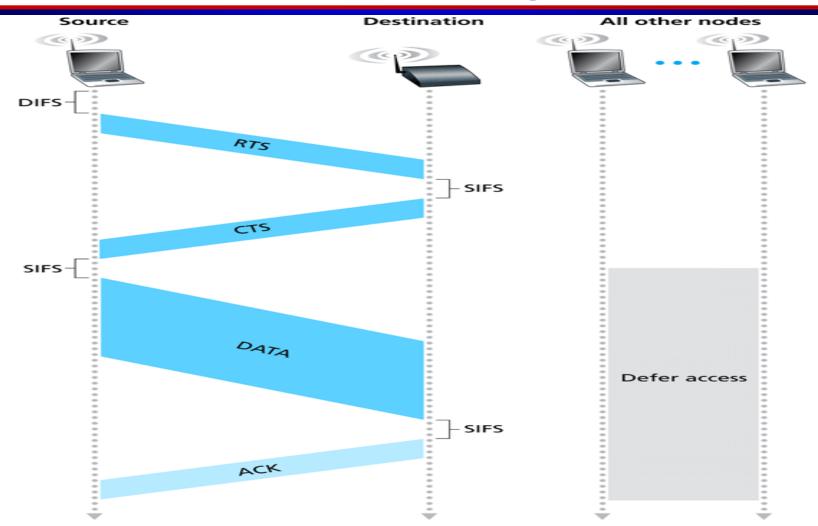
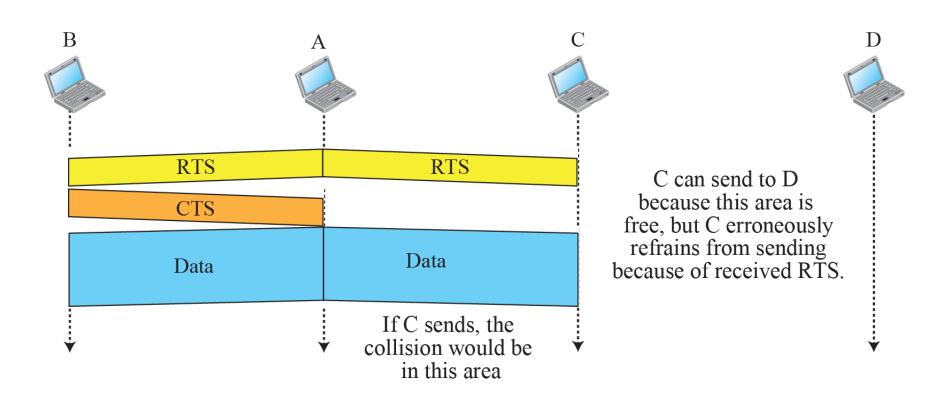
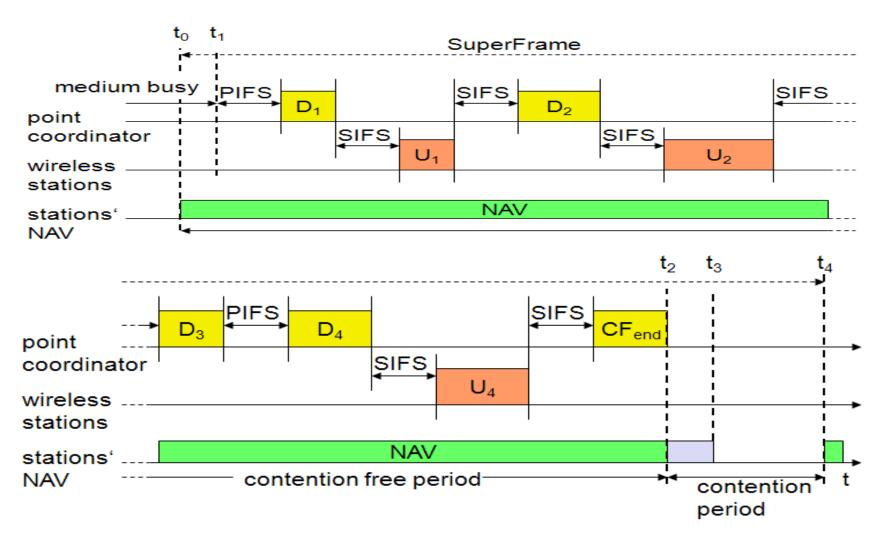


Figure 6.12 ◆ Collision avoidance using the RTS and CTS frames

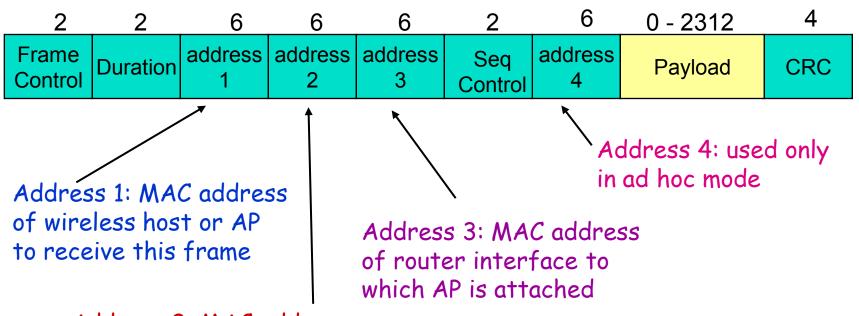
# Exposed Terminal Problem



## 802.11 PCF (Point Coordination Function)

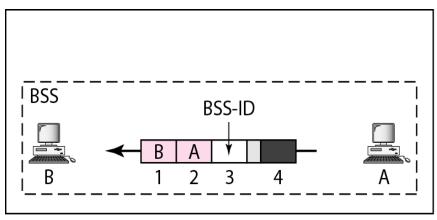


### IEEE802.11 Frame Structure

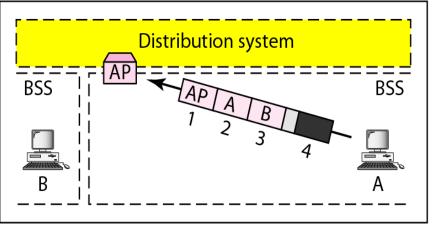


Address 2: MAC address of wireless host or AP transmitting this frame

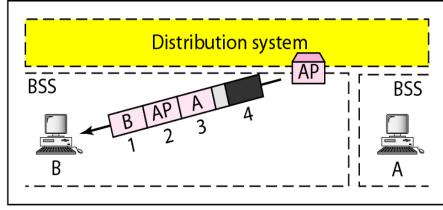
## Addressing Mechanisms



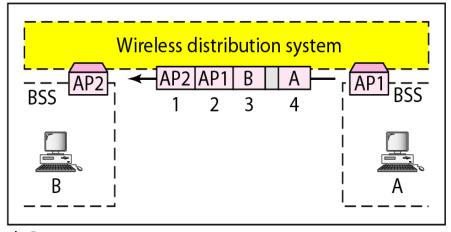
a. Case 1



c. Case 3



b. Case 2



d. Case 4

## IEEE802.11 Frame Addressing

