CS 160 Compilers

Lecture 15: Optimization

Yu Feng

Fall 2021

Slide designed by Prof. Alex Aiken, with modifications

Optimization

- Optimization is our last compiler phase
- Most complexity in modern compilers is in the optimizer

  - Also by far the largest phase

- First, we need to discuss intermediate languages

Why IR?

- When should we perform optimizations?
  - On AST
    - Pro: Machine independent
    - Con: Too high level
  - On assembly language
    - Pro: Exposes optimization opportunities
    - Con: Machine dependent
    - Con: Must reimplement optimizations when retargetting
  - On an intermediate language
    - Pro: Machine independent
    - Pro: Exposes optimization opportunities

Intermediate Languages

- Intermediate language = high-level assembly
  - Uses register names, but has an unlimited number
  - Uses control structures like assembly language
  - Uses opcodes but some are higher level
    - E.g., push translates to several assembly instructions
    - Most opcodes correspond directly to assembly opcodes

Three-Address IR

- Each instruction is of the form

$$x := y \text{ op } z$$

$$x := \text{op } y$$

- y and z are registers or constants

- Common form of intermediate code
- The expression x + y * z is translated

$$t_1 := y * z$$

$$t_2 := x + t_1$$

- Each subexpression has a name

Intermediate Code Generation

- Similar to assembly code generation
- But use any number of IL registers to hold intermediate results

Intermediate Code Generation

- You should be able to use intermediate code

  - At the level discussed in lecture

- You are not expected to know how to generate intermediate code
  - Because we won't discuss it
  - But really just a variation on code generation . . .

An Intermediate Language

$P \rightarrow S\ P \mid \varepsilon$

$S \rightarrow$ id := id op id
| id := op id
| id := id
| push id
| id := pop
| if id relop id goto L
| L:
| jump L

- id's are register names
- Constants can replace id's
- Typical operators: +, -, *

Basic Blocks

- A basic block is a maximal sequence of instructions with:
  - no labels (except at the first instruction), and
  - no jumps (except in the last instruction)

- Idea:
  - Cannot jump into a basic block (except at beginning)
  - Cannot jump out of a basic block (except at end)
  - A basic block is a single-entry, single-exit, straight-line code segment

Basic Block Example

- Consider the basic block

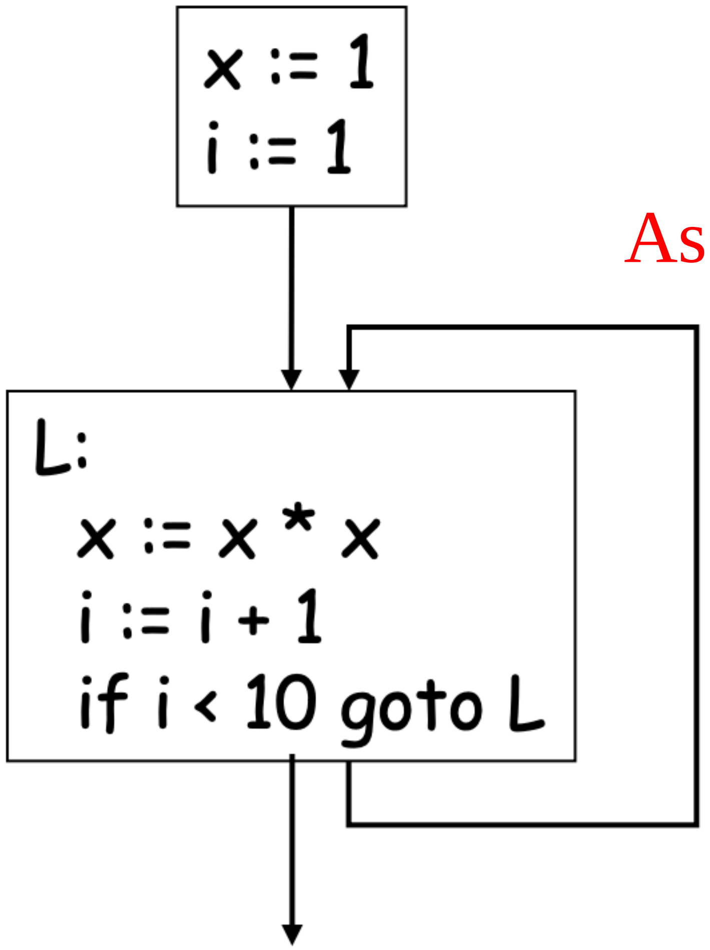
```
1.  L:
2.      t := 2 * x
3.      w := t + x
4.      if w > 0 goto L'
```


- (3) executes only after (2)
    - We can change (3) to w := 3 * x
    - Can we eliminate (2) as well?

Control-Flow Graphs

- A control-flow graph is a directed graph with
  - Basic blocks as nodes
  - An edge from block A to block B if the execution can pass from the last instruction in A to the first instruction in B
    - E.g., the last instruction in A is *jump* L_B
    - E.g., execution can fall-through from block A to block B

CFG Example



- The body of a method (or procedure) can be represented as a control-flow graph

- There is one initial node

- All "return" nodes are terminal

Optimization Overview

- Optimization seeks to improve a program's resource utilization
  - Execution time (most often)
  - Code size
  - Network messages sent, etc.
- Optimization should not alter what the program computes
  - The answer must still be the same

Classification of Optimization

- 1. ***Local*** optimizations: Apply to a basic block in isolation

- 2. ***Global*** optimizations: Apply to a control-flow graph (method body) in isolation

- 3. ***Inter-procedural*** optimizations: Apply across method boundaries

- Most compilers do (1), many do (2), few do (3)

Cost of Optimization

- In practice, a conscious decision is made not to implement the fanciest optimization known

- Why?

  - Some optimizations are hard to implement
  - Some optimizations are costly in compilation time
  - Some optimizations have low benefit
  - Many fancy optimizations are all three!
- Goal: Maximum benefit for minimum cost

Local Optimization

- The simplest form of optimizations
- No need to analyze the whole procedure body
    - Just the basic block in question
- Example: algebraic simplification

Algebraic Simplification

- Some statements can be deleted

x := x + 0

x := x * 1

- Some statements can be simplified

x := x * 0 $\Rightarrow$ x := 0

y := y ** 2 $\Rightarrow$ y := y * y

x := x * 8 $\Rightarrow$ x := x << 3

x := x * 15 $\Rightarrow$ t := x << 4; x := t - x

Constant Folding

- Operations on constants can be computed at compile time
  - If there is a statement $x := y\ op\ z$
  - And y and z are constants
  - Then $y\ op\ z$ can be computed at compile time
- Example: $x := 2 + 2$ => $x := 4$
- Example: $if\ 2 < 0\ jump\ L$ can be deleted
- When might constant folding be dangerous?

Control-flow Optimizations

- Eliminate unreachable basic blocks:
    - Code that is unreachable from the initial block
    - E.g., basic blocks that are not the target of any jump or "fall through" from a conditional
- Why would such basic blocks occur?
- Removing unreachable code makes the program smaller
    - And sometimes also faster
    - Due to memory cache effects (increased spatial locality)

Static Single Assignment (SSA)

- Some optimizations are simplified if each register occurs only once on the left-hand side of an assignment

- Rewrite intermediate code in *single assignment form*

x := z + y		b := z + y
a := x	$\Rightarrow$	a := b
x := 2 * x		x := 2 * b

(b is a fresh register)

Non-trivial due to loops and recursions

Common Subexpression Elimination

- ***If***
  - Basic block is in single assignment form
  - A definition x := is the first use of x in a block
- ***Then***
  - When two assignments have the same rhs, they compute the same value
- Example:

```
x := y + z                    x := y + z
...              ⇒            ...
w := y + z                    w := x
```

(the values of x, y, and z do not change in the … code)

Copy Propagation

- If $w := x$ appears in a block, replace subsequent uses of w with uses of x
    - Assumes single assignment form
    - Example:

b := z + y		b := z + y
a := b	⇒	a := b
x := 2 * a		x := 2 * b

- Only useful for enabling other optimizations
    - Constant folding
    - Dead code elimination

Applying Local Optimizations

- Each local optimization does little by itself
- Typically optimizations interact
  - Performing one optimization enables another
- Optimizing compilers repeat optimizations until no improvement is possible
  - The optimizer can also be stopped at any point to limit compilation time

Peephole Optimizations

- Write peephole optimizations as replacement rules where the rhs is the improved version of the lhs

$$i_1, \ldots, i_n \rightarrow j_1, \ldots, j_m$$

- The "peephole" is a short sequence of (usually contiguous) instructions

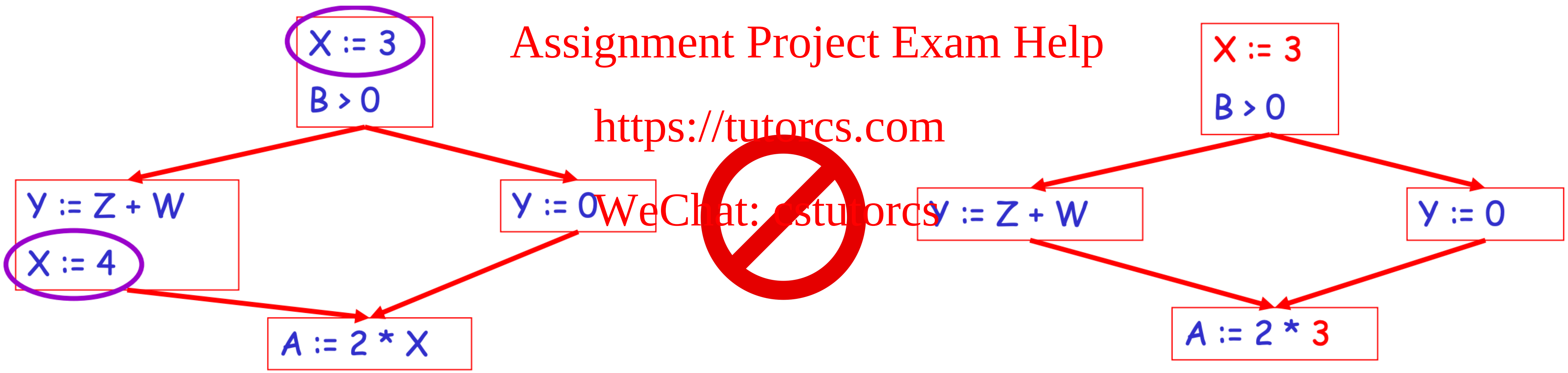
- The optimizer replaces the sequence with another equivalent one (but faster)

Global Optimizations

- Extend same optimizations to an entire control-flow graph

Global Optimizations

- Extend same optimizations to an entire control-flow graph



Correctness

- The correctness condition is not trivial to check
- “All paths” includes paths around loops and through branches of conditionals
- Checking the condition requires global analysis
- An analysis of the entire control-flow graph

An Example

- Initial code:

```
a := x ** 2
b := 3
c := x
d := c * c
e := b * 2
f := a + d
g := e * f
```

An Example

- Algebraic optimization:

```
a := x ** 2
b := 3
c := x
d := c * c
e := b * 2
f := a + d
g := e * f
```

An Example

- Algebraic optimization:

```
a := x * x
b := 3
c := x
d := c * c
e := b << 1
f := a + d
g := e * f
```

An Example

- Copy propagation:

```
a := x * x
b := 3
c := x
d := c * c
e := b << 1
f := a + d
g := e * f
```

An Example

- Copy propagation:

```
a := x * x
b := 3
c := x
d := x * x
e := 3 << 1
f := a + d
g := e * f
```

An Example

- Constant folding:

```
a := x * x
b := 3
c := x
d := x * x
e := 6
f := a + d
g := e * f
```

An Example

- Dead code elimination:

```
a := x * x
```



```
f := a + a
g := 6 * f
```

- This is the final form