



# Assignment Project Exam Help

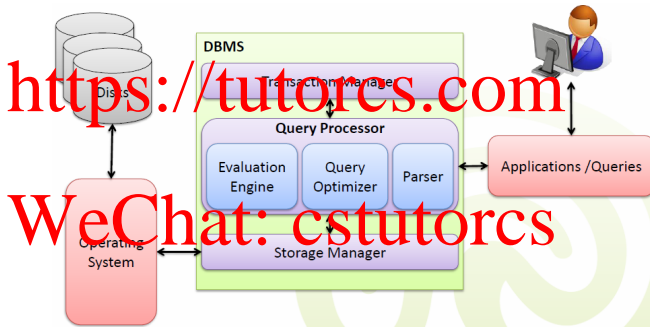
Query Processing

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## Query Processing – Overview

1. Users **submit** SQL queries to a DBMS.
2. The DBMS **processes and executes** them in a database.



- **Note:** SQL is a declarative language, so it is the task of DBMSs to decide how SQL queries should be executed.

## Query Processing – Example

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```
SELECT name FROM Person WHERE age<21;
```

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name
Rickon
Brann

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### • Questions:

- How does a relational DBMS process this?
- How can a relational DBMS process this efficiently?



## Query Processing – Example

SELECT name FROM Person WHERE age < 21

High-level language  
(SQL)



$\pi_{name}(\sigma_{age < 21}(\text{Person}))$

Low-level language  
(Relational Algebra)



$\pi_{name}$   
|  
 $\sigma_{age < 21}$   
|  
Person

Execution plan  
(Query tree)



name
Rickon
Bran

Query result

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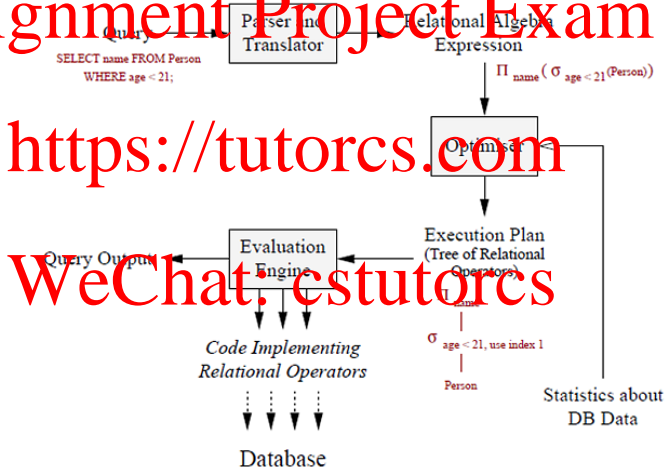
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## Query Processing – Example

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## Query Processing Steps

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- **Query parser and translator**

- 1 Check the syntax of SQL queries
- 2 Verify that the relations do exist
- 3 Transform into relational algebra expressions

- **Query optimiser**

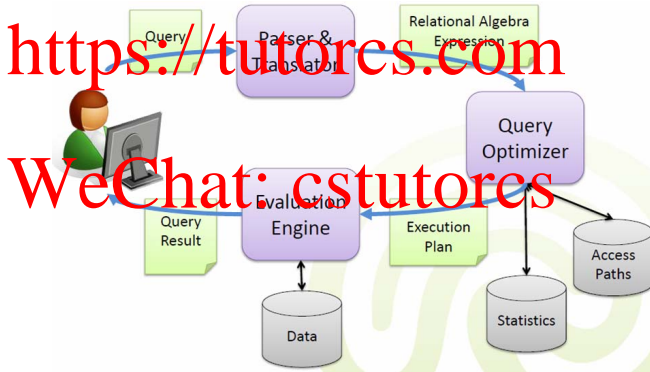
- 1 Transform into the best possible execution plan
- 2 Specify the implementation of each operator in the execution plan

- **Evaluation engine**

- 1 Evaluate the query execution plan
- 2 Return the result to the user

## Query Processing – Parser

- The **parser** checks the syntax of the query:
  - Validation of table names, attributes, data types, access permission ...;
  - Either the query is executable or an error message is generated.





## Query Processing – Parser

- Consider the relation schema:

Person(id:integer, name:string, age:integer, address:string)

- Note:** **System catalog** (also called **data dictionary**) is used at this stage, which contains the information about data managed by the DBMS.

Example:

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attr_name	rel_name	type	position
id	Person	integer	1
name	Person	string	2
age	Person	integer	3
address	Person	string	4
...	...	...	...

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- Question:** Can the following query be accepted by the parser?

```
SELECT fname, lname FROM Person WHERE address<21;
```





## Query Processing – Parser

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- Consider the relation schema.

Person(id:integer, name:string, age:integer, address:string)

- Question: Can the following query be accepted by the parser?

```
SELECT fname, lname FROM Person WHERE address<21;
```

- Answer: The query **would be rejected** because

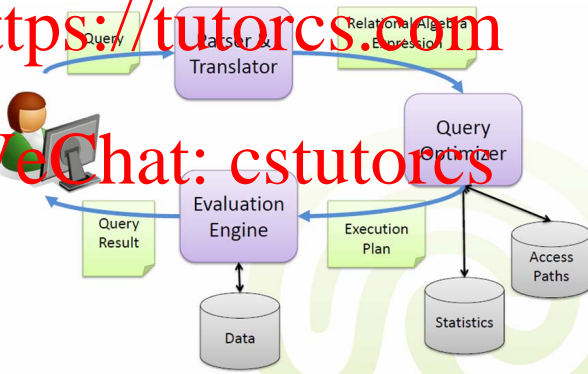
- 1 The attributes fname and lname are not defined;
- 2 The attribute address is not comparable with 21.

## Query Processing – Translator

- The **translator** translates queries into RA expressions (not necessarily equivalent due to duplicates):
  - A query is first decomposed into **query blocks**.
  - Each query block is translated into an RA expression.

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## Recall: RA and SQL Queries

### RA operators

- selection  $\sigma_\varphi$
- projection  $\pi_{A_1, \dots, A_n}$
- Cartesian product  $R_1 \times R_2$
- join  $R_1 \bowtie_\varphi R_2$  and  $R_1 \bowtie R_2$
- renaming  $\rho_{R(A_1, \dots, A_n)}$
- union  $R_1 \cup R_2$
- intersection  $R_1 \cap R_2$
- difference  $R_1 - R_2$

### SQL statement

```
SELECT attribute_list
  FROM table_list
  [WHERE condition]
  [GROUP BY attribute_list
  [HAVING group_condition]]
  [ORDER BY attribute_list];
```

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$\sigma_\varphi(R) \Leftrightarrow \text{SELECT } * \text{ FROM } R \text{ WHERE } \varphi;$

$\pi_{A_1, \dots, A_n}(R) \Leftrightarrow \text{SELECT DISTINCT } A_1, \dots, A_n \text{ FROM } R;$

$R_1 \times R_2 \Leftrightarrow \text{SELECT DISTINCT } * \text{ FROM } R_1, R_2;$

...

- Aggregate operations in SQL require extended RA expressions.



## Recall: RA and SQL Queries

- Nested subqueries are decomposed into separate query blocks.

- Example:

```
SELECT Lname, Fname
FROM EMPLOYEE
WHERE Salary > (SELECT Salary
FROM EMPLOYEE
WHERE ssn=5);
```

### Outer query block

```
SELECT Lname, Fname FROM EMPLOYEE WHERE
Salary > c
```

⇓ translated

$\pi_{Lname, Fname}(\sigma_{Salary > c}(EMPLOYEE))$

### Inner query block

```
(SELECT Salary FROM EMPLOYEE WHERE
ssn=5)
```

⇓ translated

$\pi_{Salary}(\sigma_{ssn=5}(EMPLOYEE))$

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## Query Processing – Query Optimiser

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- 1 Transform into the best possible execution plan

There are different possible relational algebra expressions for a single query

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(will be covered in this course)

- 2 Specify the implementation of each operator in the execution plan

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There are different possible implementations for a relational algebra operator!

(will not be covered in this course)



## Query Processing – Query Optimiser

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- SQL queries only specify **what data to be retrieved** and **not how to retrieve data**.
- There are **many possible execution plans** for a SQL query.
- Query optimiser is responsible for identifying **an efficient execution plan**:
  1. enumerating alternative plans (typically, a subset of all possible plans);
  2. choosing the one with the least estimated cost.
- Query optimisation is one of the most important tasks of a relational DBMS.  
**A good DBMS must have a good query optimiser!**

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## Equivalent RA Expressions

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- Suppose that we have:

Students(matNr, firstName, lastName, email)

Exams(matNr, crsNr, result, semester)

Courses(crsNr, title, unit)

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```
SELECT lastName, result, title
```

```
FROM STUDENTS, EXAMS, COURSES
```

```
WHERE STUDENTS.matNr=EXAMS.matNr AND
```

```
EXAMS.crsNr=COURSES.crsNr AND result $\leq$ 1.3;
```

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- **Question:**

How many equivalent RA expressions for this SQL query can you find?



## Equivalent RA Expressions

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```
Students(matNr, firstName, lastName, email)
Exams(matNr, crsNr, result, semester)
Courses(crsNr, title, unit)
```

```
SELECT lastName, result, title
```

```
FROM STUDENTS, EXAMS, COURSES
```

```
WHERE STUDENTS.matNr=EXAMS.matNr AND
```

```
EXAMS.crsNr=COURSES.crsNr AND result≤1.3;
```

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### Answer:

- 1  $\pi_{lastName, result, title}(\sigma_{result \leq 1.3}((\sigma_{Students.matNr=Exams.matNr} Exams) \bowtie_{Exams.crsNr=Courses.crsNr} Courses))$
- 2  $\pi_{lastName, result, title}(\sigma_{result \leq 1.3}(\sigma_{EXAMS.crsNr=Courses.crsNr} (\sigma_{Students.matNr=Exams.matNr} (Students \times Exams \times Courses))))$
- 3  $\pi_{lastName, result, title}((Students \bowtie_{Students.matNr=Exams.matNr} (\sigma_{result \leq 1.3} (Exams))) \bowtie_{Exams.crsNr=Courses.crsNr} Courses)$

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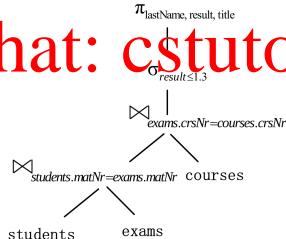
## Query Trees

- Each RA expression can be represented as a **query tree**:

- leaf nodes** represent the input relations;
- internal nodes** represent the intermediate result;
- the root node** represents the resulting relation.

- Example:**

$\pi_{\text{lastName}, \text{result}, \text{title}} (\sigma_{\text{result} \leq 1.3} ((\text{Students} \bowtie_{\text{Students.matNr}=\text{Exams.matNr}} \text{Exams}) \bowtie_{\sigma_{\text{Exams.crsNr}=\text{Courses.crsNr}}} \text{Courses}))$



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## Query Trees

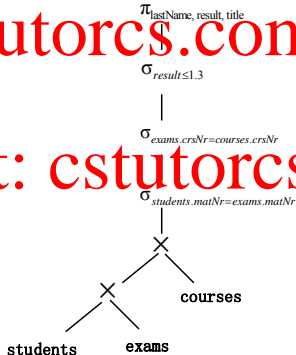
- **Exercise:** Can you draw the query tree for the following RA expression?

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$$\pi_{\text{lastName}, \text{result}, \text{title}}(\sigma_{\text{result} \leq 1.3}(\sigma_{\text{Exams.crsNr}=\text{Courses.crsNr}}(\sigma_{\text{Students.matNr}=\text{Exams.matNr}}(\text{Students} \times \text{Exams} \times \text{Courses}))))$$

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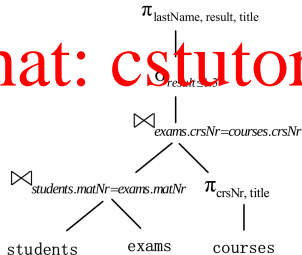
## Query Trees

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- For each query tree, computation proceeds **bottom-up**:
  - child nodes must be executed before their parent nodes;
  - but there can exist multiple methods of executing sibling nodes, e.g.,
    - process sequentially,
    - process in parallel.

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## Execution Plan

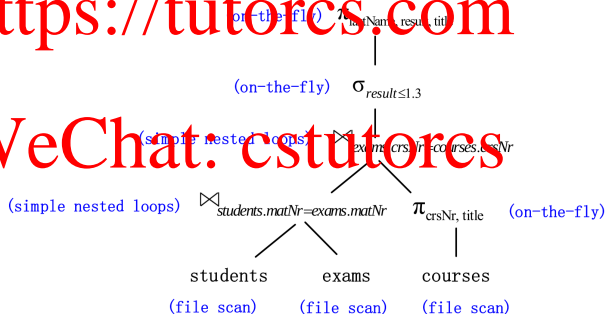
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- A query execution plan consists of an (extended) query tree with additional annotation at each node indicating:

- (1) the *access method* to use for each table, and
- (2) the *implementation method* for each RA operator.

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## Query Processing – Evaluation Engine

The evaluation engine executes an execution plan, and returns the query answer to the user.

