

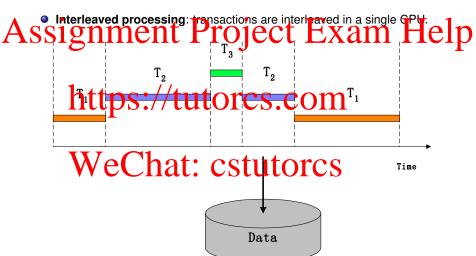
## Assignment Project Exam Help

https://ctutenransactions

WeChat: cstutorcs



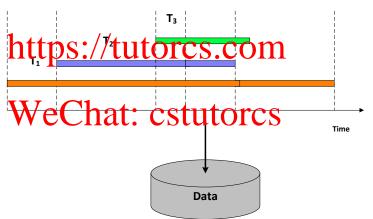
#### **Concurrent Transactions**





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## Assignment Project Exam Help





#### **Concurrent Transactions**

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- → Increase throughput (average number of completed transactions)
- For example, while one transaction is waiting for an object to be transaction from district the CPIP carchrocess prother transaction (because I/O activity can be done in parallel with CPU activity).
- --> Reduce latency (average time to complete a transaction)
- For example, interleave execution of a short transaction with a cong transaction usually allows the short one to be completed more quickly.
- But the DBMS has to guarantee that the interleaving of transactions does not lead to inconsistencies, i.e., concurrency control.



## Why is Concurrency Control Needed?

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Concurrency control is needed for preventing the following problems:

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The dirty read problem

When the peaked that probatutores

The phantom read problem



## (1) - The Lost Update Problem

## Assignification of the property of the second of the secon

```
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
T2: SELECT balance FROM ACCOUNT WHERE name='Bob';
T1: UPDATE ACCOUNT SET balance=balance-100 WHERE name='Bob';
T2: UPDATE ACCOUNT SET balance=balance+500 WHERE name='Bob';
T2: COMMIT;
```

Steps	T	<i>T</i> <sub>2</sub>					
1111	rad(E)	CS	M	or		Steps	B(Bob)
2		read(	B)			before 1	\$200
3	write(B) (B:=B-100)					after 2	\$200
4	commit					after 4	\$100
5				=B+500)		after 6	\$700
6		comr	nit		_		



## (1) - The Lost Update Problem

## Assignification of the property of the second of the secon

```
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
T2: SELECT balance FROM ACCOUNT WHERE name='Bob';
T1: UPDATE ACCOUNT SET balance=balance-100 WHERE name='Bob';
T2: UPDATE ACCOUNT SET balance=balance+boo WHERE name='Bob';
T2: COMMIT;
```

Steps	T	T <sub>2</sub>			
	Prad(II) 2	cstut	orc	Steps	B(Bob)
2''		read(B)		before 1	\$200
3	write(B) (B:=B-100)			after 2	\$200
4	commit			after 4	\$100
5		write(B) (B:=B	8+500)	after 6	\$700
6		commit			

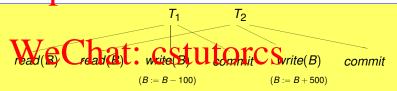
• Answer: Bob's balance should be \$600. The update by  $T_1$  is lost!



## (1) - The Lost Update Problem

A SS occurs vite nive transactions to date the same object, and one transaction position of the object which has already been updated by another transaction (write-write conflicts).

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• write(B) by  $T_2$  overwrites B, and the update by  $T_1$  is lost.



#### (2) - The Dirty Read Problem

## Assignituble actions (E). To ject Exam Help

```
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
T1: UPDATE ACCOUNT SET balance=balance-100 WHERE name='Bob';
T2: SELECT balance FROM ACCOUNT WHERE name='Bob';
T3: BDRT; S / LILLOT CS COUNT NHERE name='Bob';
T4: DPDATE ACCOUNT SET balance=balance+boo WHERE name='Bob';
T2: COMMIT:
```

Steps	T	T <sub>2</sub>
1 <b>VV</b> (	Pread(II)	<b>estutore</b>
2	write(B) (B:=B-100)	
3		read(B)
4	abort	
5		write(B) (B:=B+500)
6		commit

Steps	B(Bob)
Defore 1	\$200
after 1	\$200
after 2	\$100
after 4	\$200
after 6	\$600



#### (2) - The Dirty Read Problem

## Assignification of the property of the second of the secon

```
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
T1: UPDATE ACCOUNT SET balance=balance-100 WHERE name='Bob';
T2: SELECT balance FROM ACCOUNT WHERE name='Bob';
T1 AFORT
T2: UPDATE ACCOUNT SET balance=balance+500 CHEIEIname='Bob';
T2: COMMIT:
```

Steps	<i>T</i> <sub>1</sub>	<i>T</i> <sub>2</sub>	Steps
TTT	read(B)	4 4	<u> </u>
2 <b>W</b> (	Write(B)(13B100)	cstutoro	<b>Defore</b>
3		read(B)	after 1
4	abort	1000(2)	after 2
5	abort	write(B) (B:=B+500)	after 4
6		commit	after 6
		Commit	

Steps	B(Bob)
<b>Operation</b>	\$200
after 1	\$200
after 2	\$100
after 4	\$200
after 6	\$600

Answer: Bob's balance should be \$700 since T<sub>1</sub> was not completed.



#### (2) - The Dirty Read Problem

A S Secure when one transaction pould read the value of an object that his a positive with the secure of an object that his a positive with the secure conflicts).

Example:



•  $T_1$  fails and must change the value of B back to \$200; but  $T_2$  has read the uncommitted ( $\cong dirty$ ) value of B (\$100).



## (3) - The Unrepeatable Read Problem

## A SS to with the v (100) while A ce with a raws (500) frame (500) frame (500) and (500)

```
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
T2: SELECT balance FROM ACCOUNT WHERE name='Bob';
T3: WPDATE ACCOUNT SET balance=balance-500 WHERE name='Bob';
T1: SELECT balance FROM ACCOUNT WHERE name='Bob';
```

Steps	$T_{1}$	$T_2$			
	ad(Ib)	ี 21	•	$CS^1$	m
2		re	ad(B)		u
3		ıw	rite(B)	) (B:=B-5	500)
4		CC	mmit		
5	read(B)				

Steps	B(Bob)
hafqre C	\$500
after 2	\$500
after 3	\$0
after 4	\$0
after 5	\$0



## (3) - The Unrepeatable Read Problem

# Example: Bob checks his account (T<sub>1</sub>) twice (takes time to decide whether St. with Have \$100) while Alce withdraws \$500 fram Poble account T<sub>2</sub>. C property of the second with the second wit with the second with the second with the second with the second

## her since the correspond

Steps	<i>T</i> <sub>1</sub>	$T_2$	
2W	read(B)	aetad(B)CStU	1
3		write(B) (B:=B-500)	·
4		commit	
5	read(B)		

Steps	B(Bob)
before 1	\$500
af e(2 S	\$500
after 3	\$0
after 4	\$0
after 5	\$0

 Answer: Bob received two different account balances \$500 and \$0, even though he hasn't withdrawn any money yet.



(B = 500)

(B = 500)

## (3) - The Unrepeatable Read Problem

# Assignment Project Exam Help another transaction but is still in progress (could issue two read for the object, or a write after reading the object) (read-write conflicts).

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# WeChat: Cstutorcs read(B) read(B) write(B) commit read(B)

(B=0)



#### (4) - The Phantom Read Problem

**Example:** A query is subrigited for finding all customers whose actor into the balance see that \$3.00 ( $I_1$ ) while Alice is opening a new account with the balance \$200 ( $I_2$ ).

Assume that only Bob (B) has an account whose balance is less than \$300 before Alice (A) opens his new account.

T1 SILICE have FROM ACCOUNT HERE has an Common Scott

T2: INSERT INTO ACCOUNT(id, name, balance) VALUES(99, 'Alice', 250);

 $T_2$ : COMMIT;

T<sub>1</sub>: SELECT name FROM ACCOUNT WHERE balance<300;

Sees		at: C
1	read(R)	
2		write(R)
3		commit
4	read(R)	

Sep C	Query result
before 1	$R = \{B\}$
after 1	$R = \{B\}$
after 2	$R = \{A, B\}$
after 4	$R = \{A, B\}$



#### (4) - The Phantom Read Problem

Example: A query is subjurged for finding all customers whose actor in S balances are less than \$310 ( $I_1$ ) by the balance \$200 ( $I_2$ ).

Assume that only Bob (B) has an account whose balance is less than \$300 before Alice (A) opens his new account.

```
T_1 SILCT have FROM Account (id, name, balance) VALUES (99, 'Alice', 250); T_2: INSERT INTO Account (id, name, balance) VALUES (99, 'Alice', 250); T_2: COMMIT;
```

 $T_1$ : SELECT name FROM ACCOUNT WHERE balance<300;

Seps	e( 'h	at· c	stii	Stepp	Query result
1	read(R)	<del></del>	blu	before 1	$R = \{B\}$
2		write(R)		after 1	$R = \{B\}$
3		commit		after 2	$R = \{A, B\}$
4	read(R)			after 4	$R = \{A, B\}$

 Answer: T<sub>1</sub> reads Account based on the condition balance<300 twice but gets two different results {B} and {A, B}.



## (4) - The Phantom Read Problem

## Assignment Project Exame Help conditions of another transaction so that, by the same search condition, the

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transaction obtains different results at different times.

