NJS: Database Protection Algorithm

Non-deterministic and post-quantum cryptography

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Abstract: NJS is a cryptographic protection algorithm for relational databases with non-deterministic symmetric encryption, making it possible to search data with almost the same speed as a clear text search (depending on the parameterization). The algorithm has the characteristic of performing a fast encryption on the data and a slightly slower decryption that is only performed on the client workstation. The entire process of searching, changing, adding and deleting data is performed on the server with the encrypted data. The NJS cipher is not a form of homomorphic encryption, but it can replace it with some search limitations. One advantage is the fact that noise added to the message does not interfere with its decryption, regardless of the number of operations performed on each record in a database table.

Keywords: symmetric cryptography, database protection, non-deterministic cryptography.

1 Introduction

In this paper we propose a new symmetric non-deterministic algorithm: NJS¹. The algorithm attempts to realize an encryption with security level equivalent to RSA from a simple structure. Such a goal is achieved without using the factorization problem or the discrete logarithm problem. Therefore, without compromising the algorithm in the scenario where quantum computers are a reality. The RSA algorithm has been attacked in several ways [2][3][4] but most attacks start from the knowledge of the modulus N and the exponent used in encryption which is usually fixed.

In section 2 we present the structure of the NJS algorithm briefly. Section 3 presents the rationale of the algorithm, the proposed encryption of the data, the structure of the cipher key, and the interaction between numerical operations of different algebraic groups in arbitrary block sizes without the need for padding. Section 4 presents operational details such as the application of noise levels, examples of non-deterministic encryption, efficient use of hash functions to prevent decryption errors, search techniques and

 $^{1~\}mathrm{NJS}$ is the acronym for Noêmia Josefina da Silva.

security in the use of the SQL language [5]. Section 5 makes an interesting comparison (respecting the difference between symmetric and asymmetric encryption) between the RSA system and the NJS algorithm. In this section we show that the use of the fixed exponent 2¹⁶+1 in RSA reduces the exponentiation of long integers to a single modular multiplication. Without delving into the question of a possible vulnerability of RSA, we have reasoned about the NJS algorithm, showing the use of at least 9 algebraic operations against a single operation in RSA. Can we at least conjecture that the level of security is equivalent? In section 6 we present the structure of the source code files that are available in annex III. Annex I presents an example of a system key and annex II shows a table of 30 records with dummy data in its native encrypted form. Such records are intended to give an objective idea of the increase in data volume in the encrypted records. In the conclusion we present some data thoughtfully.

The remainder of the introduction deals with the limitations of this work and our contributions.

1.1 Limitations of this paper

The NJS algorithm cannot fully replace a homomorphic encryption [6]. However, it can allow a group of users sharing the same key to access a cloud-hosted server [7] in encrypted form. The algorithm also allows searches on encrypted data, changes, deletions and additions of new records (at a rather high but feasible cost).

NJS is a highly parameterizable algorithm. In our example code we work with blocks of 480 bits and modules of at least 512 bits. These values can be adjusted almost without limit, and this limit only comes up against the speed at which the data can be processed. It has not yet been possible to establish a formal balance between security and execution speed for the NJS algorithm.

In order to make queries and data changes feasible, it is necessary to employ simpler primary keys, which can minimize the cost of accessing the encrypted data. In general, it is not possible to search for information that is smaller than the proposed block (480 bits), unless the block matches the entire field of a given record.

Another disadvantage of the proposed scheme is the use of "fake queries" to protect existing data relationships from attackers who might monitor SQL operations on the database. The use of "fake queries" has the advantage of confusing those who are monitoring the SQL operations, however the server must be able to handle a much larger number of requests to perform tasks that in a decrypted database would be much less costly.

The prototype executable that can be obtained from the source code should be further explored in the future. The ciphering and decryption results were encouraging, both when considering the performance of the algorithm and the non-deterministic characteristic of the cipher. Despite this, it was not possible to formalize these observations through mathematical proofs. No major cryptoanalytic studies have been performed to date. However, in Section 5 we present a comparison of the NJS algorithm with the RSA algorithm and make a credible argument for the feasibility of this proposal.

Another negative point, but not a deterrent to using the NJS algorithm, is the increase in data volume. Using parameters of the same order of magnitude as those illustrated in the source code, it is possible to see an increase in the database of about 6.7 times compared to the same database unencrypted.

1.2 Our contribution

The NJS algorithm can encrypt a relational type database [8] with numerous tables using a non-deterministic cipher that is supposed to have a security level similar to RSA. But unlike RSA, the NJS algorithm will not suffer from quantum attacks because it does not use factorization or the discrete logarithm problem to implement its security. Its symmetric structure allows it to hide from the attacker the main pillars that he could use to carry out his attacks.

The algorithm clearly tries to benefit from employing numerical operations involving different algebraic groups, namely addition, multiplication, and the XOR operation. Such a strategy was successfully employed in the symmetric IDEA algorithm [1] (International Data Encryption Algorithm). NJS innovates this proposal by allowing the application of such operations with data blocks of 480 bits (proposed size) or even larger, without any limitations. The IDEA algorithm employs 64-bit blocks, with 16-bit words (using the prime 2¹⁶+1 as a modulus for multiplication). This block size is now obsolete.

With the proposed parameterization in our source code the NJS algorithm can generate up to 2056 combinations of cryptograms for the same block of text using the same key. The proposed text block size is 480 bits (60 characters). This value seems reasonable to us to avoid too many repetitions in the data while protecting, within a certain limit, the primary key of the table.

The NJS algorithm can also be used to encrypt simple data (not in a relational database structure). In this case it is possible to employ noise levels [9] of the magnitude of 256 bits (or larger) thus avoiding any kind of pattern repetition in the data, even if it is a concatenation of a single byte repeated billions of times. This is a huge advantage from the operational point of view of a non-deterministic cipher. The fact is further justified because noise removal is almost instantaneous in a legitimate decryption and very difficult in the absence of the correct cipher key. However such a practice prevents searching encrypted data due to the impossibility of mapping all possible search variants. If the encryption is of a simple file it can be advantageous from an operational point of view.

Finally, the data compression algorithm used in NJS generates a result that contains only printable characters, easy to save in files with a simple structure such as TXT or CSV.

2 Algorithm Structure

The NJS algorithm behaves like a non-deterministic cipher, i.e., using a fixed key (K) it is always possible to obtain a different cipher (C) using the same cleartext (P).

The algorithm works with random data mixed with the clear message and the cipher key. The random data used in the encryption process can come from any source, it just has to be concentrated in a parameterized range of values to allow correct decryption.

The cost of encryption is minimal but the cost of decrypting a message is somewhat high, which translates into one more security parameter for the algorithm. If to legitimately decipher a cryptogram it is necessary to process more information then it is also very likely that an attacker, trying to break the code without access to the keys, will have bigger problems than the legitimate decryptor.

The structure of the algorithm is very simple and works by using different algebraic groups to encrypt a message. This strategy has precedents in cryptography and was successfully employed in the IDEA algorithm. A key Sbox is used before processing to make it difficult to identify the starting point of the data. This Sbox is used with a function based on its inverse Sbox, which makes it possible to differentiate the keys used for each field of a database record.

The use of the NJS algorithm is advantageous when we have an encrypted database that needs to be accessed by several legitimate clients. The database can be hosted on an insecure server but the decryption is only performed on the client computer. Of course, the data traffic between the client and the server can be strengthened from a security point of view by using a post-quantum asymmetric system, but the algorithm seems to be enabled to work without this extra security barrier if necessary.

3 Logical Rationale

Let \mathbf{M} be a message (clear text) represented by the vector $\mathbf{Y} = [b_1, b_2, b_3, b_4, b_5, ..., b_n]$, where each $\mathbf{b_x}$ value is contained in the range [0, 255]. The vector \mathbf{Y} represents a field of a database record or a fraction of such a field. Thus a record can have \mathbf{X} number of fields. A table \mathbf{T} may contain millions of \mathbf{X} records, and a database \mathbf{B} may contain tens, hundreds, or even thousands of \mathbf{T} tables.

Table fields:

$$\begin{array}{l} X_1 = Y_{11}[b_1,\,b_2,\,b_3,\,...,\,b_n] \,+\,Y_{12}[b_1,\,b_2,\,b_3,\,...,\,b_n] \,+\,...\,+\,Y_{1n}[b_1,\,b_2,\,b_3,\,...,\,b_n] \\ X_2 = Y_{21}[b_1,\,b_2,\,b_3,\,...,\,b_n] \,+\,...\,+\,Y_{2n}[b_1,\,b_2,\,b_3,\,...,\,b_n] \\ X_3 = Y_{31}[b_1,\,b_2,\,b_3,\,...,\,b_n] \end{array}$$

Table records:

$$R_1 = X_1 + X_2 + ... + X_n$$

Full table:

$$T_1 = R_1 + R_2 + R_3 + ... + R_n$$

Database:

$$B = T_1 + T_2 + T_3 + ... + T_n$$

A field X of an R-record of a T-table is separated into blocks of size Y (512 bits or a value close to it). Each block Y is encoded as a long unsigned integer (BigUint). This long integer is modified by at least 3 algebraic operations: addition, multiplication and XOR. The order of execution of these operations does not depend on the cipher key and can be defined by a pseudo-random number generator (or even a truly random generator). The operations are modular for some prime P of size greater than 512 bits.

$$Y_1 = (257^3 * (b_3+1)) + (257^2 * (b_2+1)) + (257^1 * (b_1+1)) + (257^0 * (b_0+1))$$

This is an example of a 32-bit block. To avoid filling null bytes to complete a block (which for a database record would be hardly applicable in practice) we operate in base 257 instead of base 256. So we don't have any message represented by the value 0, in any of its bytes. This way we can work with incomplete blocks. Note also that it is necessary to add one unit to the value of each byte in the encryption and then subtract it in the decryption.

The word limits will be extended. In the case of 32-bit words we would have a limit of 4294967295 ($256^4 - 1$) and in the case of the NJS algorithm we would have 4362470400 ($257^4 - 1$). This implies that if we use blocks of 512 bits for processing the fields ($256^{64} - 1$) we will have to have a prime modulus **P** that is larger than ($257^{64} - 1$). This implies an increase in data volume of about 1.28 times.

The system key is represented by a vector with 48 values of at least 512 bits each, plus an Sbox (a vector that represents a permutation between all 256 8-bit values) and of course the Sbox⁻¹ (inverse Sbox).

If a field in a record has the structure $R_1 = Y_1[b_1, b_2, b_3, ..., b_n]$, where b_n represents a value in the range [1, 256], then the corresponding cryptogram has the following structure:

$$R_{[0]} = [(r \ *257^{59}) + ((S[b_1] \ xor \ F_i)^* \ 257^{58}) + (S[b_2] \ xor \ F_i)^* \ 257^{57}) + ... + \ (S[b_n] \ xor \ F_i)^* \ 257^0)]$$

 $R_{[1]} = (R_0 \text{ ADD } K_{11}) \text{ mod } M_1$

 $R_{[2]} = \ (R_1 \ MUL \ K_{12}) \ mod \ M_1$

 $R_{[3]} = R_2 XOR K_{13}$

Here we see the 4 states of the register $\mathbf{R}_{[\mathbf{r}]}$, where $\mathbf{R}_{[0]}$ is its initialization, which is done by passing through the Sbox \mathbf{S} and adding a noise value \mathbf{r} . Besides this a \mathbf{F}_i value (field index) is added, which is directly linked to the field number of the register associated with the Sbox \mathbf{S}^{-1} (inverse Sbox).

Then an addition (ADD) and a multiplication (MUL) are performed in module M (Keys K_{11} and K_{12}). Finally an XOR operation is performed with key K_{13} . This procedure is executed 3, 6, 9 or 12 times. The variable that defines these values comes from a pseudo-random number generator and has no relation to the cipher key.

The complete cipher looks like this:

$$R_{[0]} = [(r \ *257^{59}) + ((S[b_1] \ xor \ F_i)^* \ 257^{58}) + (S[b_2] \ xor \ F_i)^* \ 257^{57}) \\ + ... + \ (S[b_n] \ xor \ F_i)^* \ 257^0)]$$

$\overline{R_{[1]} = (R_0 \text{ ADD } K_{XA}) \text{ mod } M_N}$		$R_{[1]} = \ (R_0 \ MUL \ K_{XA}) \ mod \ M_N$	Number of
$R_{[2]} = (R_1 \text{ MUL } K_{XB}) \text{ mod } M_N$	or	$R_{[2]} = (R_1 \text{ ADD } K_{XB}) \text{ mod } M_N$	rounds:
$R_{[3]} = R_2 XOR K_{XC}$		$R_{[3]} = R_2 XOR K_{XC}$	3, 6, 9 or 12.

The presented structure allows you to choose from 8 possible variations to encrypt each database field. The number of repetitions can represent a key. However, for operational ease, fixed values were set for the number of rounds (3, 6, 9 or 12). Such values were chosen to thwart side-channel attacks [10]. These attacks analyze the algorithm's processing time. As an example we could say that it would not be possible to easily distinguish two encrypted records with 6 rounds from another encrypted record with 12 rounds, or even 4 encrypted records with 3 rounds each.

The set of keys in the algorithm can be defined thus:

Table 1: NJS system keys

Round	Encryption / Decryption Operations	Prime module
	$\mathrm{K}_{11} \; (\mathrm{MUL} \; \mathrm{or} \; \mathrm{ADD})$	M_1
1	$\mathrm{K}_{12} \; (\mathrm{MUL} \; \mathrm{or} \; \mathrm{ADD})$	M_1
	$\mathrm{K}_{13} \; \mathrm{(XOR)}$	
	$\mathrm{K}_{21} \; \mathrm{(MUL \; or \; ADD)}$	M_2
2	$\mathrm{K}_{22} \; \mathrm{(MUL \; or \; ADD)}$	$ m M_2$
	$\mathrm{K}_{23} \; (\mathrm{XOR})$	
	K ₃₁ (MUL or ADD)	M_3
3	$\mathrm{K}_{32} \; \mathrm{(MUL \; or \; ADD)}$	M_3
	$\mathrm{K}_{33} \; (\mathrm{XOR})$	
	K ₄₁ (MUL or ADD)	M_4
4	$\mathrm{K}_{42} \; \mathrm{(MUL \; or \; ADD)}$	M_4
	$\mathrm{K}_{43} \; (\mathrm{XOR})$	

	K ₅₁ (MUL or ADD)	M_5
5	$\mathrm{K}_{52} \; (\mathrm{MUL} \; \mathrm{or} \; \mathrm{ADD})$	M_5
	K ₅₃ (XOR)	
	K ₆₁ (MUL or ADD)	M_{6}
6	$\mathrm{K}_{62} \; (\mathrm{MUL} \; \mathrm{or} \; \mathrm{ADD})$	M_{6}
	K ₆₃ (XOR)	
	K ₇₁ (MUL or ADD)	M_{7}
7	K ₇₂ (MUL or ADD)	M_{7}
	K_{73} (XOR)	
	K ₈₁ (MUL or ADD)	M_8
8	K ₈₂ (MUL or ADD)	M_8
	K_{83} (XOR)	
	K ₉₁ (MUL or ADD)	M_9
9	K ₉₂ (MUL or ADD)	M_{9}
	K_{93} (XOR)	
	K _{A1} (MUL or ADD)	$ m M_A$
10	K _{A2} (MUL or ADD)	$ m M_A$
	K _{A3} (XOR)	
	K _{B1} (MUL or ADD)	$ m M_B$
11	K _{B2} (MUL or ADD)	$ m M_B$
	K_{B3} (XOR)	
	K _{C1} (MUL or ADD)	$ m M_{C}$
12	K _{C2} (MUL or ADD)	$ m M_{C}$
	K_{C3} (XOR)	

An important note is to establish the relationship between the prime modulus M_x , and the keys used K_{x1} , K_{x1} and K_{x3} .

1. M_x must be a prime number of at least 512 bits: It is important that this value cannot be attacked by exhaustive search. There are about $1.88 * 10^{151}$ prime numbers of 512 bits.

$$\frac{2^{512}}{\log(2^{512})} - \frac{2^{511}}{\log(2^{511})} = 1.88 \times 10^{151}$$

2. $K_{\rm x1}$ and $K_{\rm x2}$ can be any number of 512 bits smaller than $M_{\rm x}.$ Since $M_{\rm x}$ is a prime number there will always be the multiplicative inverse K_{x1}^{-1} or K_{x2}^{-1} that will allow the data to be deciphered. The inverse of the sum is given by M_x-K_{x1} or $M_x-K_{x2},$ all these values being unsigned. $\space{1mm}{7}$

3. The key K_{x3} is used in an XOR operation. Thus, in decryption the value of K_{x3} matches the value of K_{x3} . An important remark should be made here: K_{x3} need not necessarily be smaller than M_x since in XOR operations we do not use a modular operation. However, the value M_{x+1} must be larger in number of bits than M_x and K_{x3} for decryption to be possible. The IDEA system takes advantage of the fact that $2^{16} + 1$ is a prime number. In the case of the NJS algorithm we do not have a limit of the cipher block by the number of bits. So for decryption to be possible, if M_x has 512 bits and K_{x3} has 512 bits then M_{x+1} must contain at least 513 bits. This is true for the complete string of M_x and K_{x3} values.

The full key can be understood as a set of 48 values of approximately 512 bits plus one Sbox^2 .

Table 2: Total system key size

Keys	Size
K_x	3 * 12 * 512 (bits)
$ m M_x$	512 + 513 + 514 + 515 + 516 + 517 + 518 + 519 + 520 + 521 + 522 + 523 (bits)
Sbox	$256! \approx 2^{1684} \text{ bits}$
Total	26326 bits

A brute force attack is very complicated because none of these values are exposed (K_x, M_x) and Sbox and all the keys are used only in the client environment. All operations performed on the database are done with the encrypted values. To have a higher level of protection we recommend the use of a different Sbox for each database table. One last remark is to point out that it is possible to work with larger key numbers (1024 bits or more) but this will increase the size of the database considerably.

4 Operational details

In order to operate the NJS algorithm it is necessary to add an extra field to each database table. This field must have as content a hash value. For the NJS algorithm we adopted SHA3-224 but any strong hash function could have been used. In this case we opted for a smaller hash so that the impact on the database volume would be as small as possible. The hash value for this field is given by concatenating the contents of all fields in the record.

A record in a database could be represented as follows:

² A complete example of a key can be seen in Annex I.

Table 3: Example of a record

Field	Value	
Name	Tairis Geissler Behrends Delazare Groskopf	
Document	180651975-23	
Phone	(64)98073-4980	
Date of Sale	11/12/2017	
Sale Value	1377.23	
Hash	dcc6efc903fbd40ae78c31ace6e39b39ea0d8e19e15673822eb27e8c	

The SHA3-224 hash value is required to allow decryption of the data. Since the cipher is non-deterministic in nature then a parameter is needed to differentiate a "correct decryption" from an "incorrect decryption". Examples of possible encryptions with the same key:

Table 4: Ciphered record (version 1)

Field	Value	
Name)TI20v/40Bt+uCkDg/ t}<*800=+0Jryb60OX0SU`zmYk40cU60R:80cH&ijn0vY?! 70KM0MuWd!0Duv#mX20ZN0EJ09	
Document	$+ w10@hkY;*(-F[0W]zL70f~s0*tNbNaJH}00=)0CmPgV0psYUNFzA90f^m~]400~'<50G\{+00050Pywz20[vh402$	
Phone	$1090 MI0c5010 Tt QuTkXu < 80b30') frm: W\&Z < J\{40\$n0a0\$f0a\}B60w + tb\$01`j00k D\&X\#IS0zRJd80 THd10r < 100 > 0 > 7$	
Date of Sale	$ \begin{array}{l} Nq_U-3060-vkD)C0defm0*80z0/)_40F0MR0T!M`,00qK?\\ A\{30N20BgJ\#Nl0N0XQ]jT10!Rjq_zFFT)R10,0Nr80i?^9 \end{array} $	
Sale Value	_#p30SHJ_(`*~?dPCRj0050?ovH#0/[N60? cWct00rofZ0do70#f;}p0r900/iaM=-os*_fQ'i0>LJ}DF@-pHzV	
Hash	_Yr+XADIJv0x]*UwYg40Y90dN@(0=60600zBYrAyey0y_700`<80byl0i YP*C}<;60/wZ0h\$ 0J,e*&)040@tD4040^b	

Table 5: Ciphered record (version 2)

Field	Value	
Name	$\begin{array}{l} M70?90kP0l200O{>}KCTOBJ0B500t{+}O:VD70nif0c70UUca?Cvf?\\ 0)NjEv/TNShvQ0@(090P{+}kwsBMA50q~L_Y)Kv!90`6 \end{array}$	
Document	vC0zOM70/S=bhk800?_qE&Qxt/>40Jb20i&K0Zk0o<0wER- 0e0s])30LGZF`0y;40/06090nQY0hI,0[0=S0MMNSM~0M203	
Phone	^uBT0yw>DB)0!D020\$hz70i0!0'L0{>(`0! 0N=700XOx90nWj00k&0KSU&}/ hoO[40800z^jO[S,zwjUdyz0^iXKCQM0	
Date of Sale	70rI00:XB^m{40/!<40T60+DY0N0AYY*W,[Y qL^+0J,! #J0n7060m+Dvhbz_o601040Dr`eCTNwnmt>G30_Jf0,A20V)	

Sale Value	~m`>W t^50'gdQD2050jgm^MpV{v70R20`020viz kbMEb*S[yB]q90,vC0dj)<40fVgZT^cX'40,=W#0i klf}c
Hash	BBLbUR,0@_Qi&MY20Q ^K+q90k>E-;0q'fWTrXUs10kmA30^Wqx[T60QMD;b0=z{Yk0a,i? Rex0gFY00O70E0DB9

From the table plus the hash field it is possible to safely encrypt and decrypt the records. The noise r has 257 variations and the encryption functions are 8. This way we have a total of 2056 encryption variations for the same data with the same key. The noise level can easily be increased without compromising the performance of the algorithm, but by adding too high a noise level the speed of data search is compromised.

An apparently effective noise level is 16, thus allowing 128 combinations for each cipher block. This value seems reasonable to us in terms of security and effective in not hindering database searches. Thus, since a field can be encrypted in 128 different ways, it is necessary to search all of them to obtain the data for decryption.

But in the search it is necessary to include random values to avoid cryptoanalysis by search frequency, indicating that those data were produced by the same cleartext. So it is highly recommended that when referring an SQL query to the database more queries are done with random values that the legitimate client can easily isolate in decrypting the data but the attacker will be stuck in them, not knowing exactly which cryptograms are true and which are false.

Example of an SQL query:

```
Select Code From Table
Where (Code = x_1) or (Code = x_2) ... or (Code = x_{256})
```

In this case the legitimate client of the system, who has access to use the keys, creates a search vector $\mathbf{Y} = [\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3, \mathbf{x}_4, ..., \mathbf{x}_{254}, \mathbf{x}_{255}, \mathbf{x}_{256}]$. The contents of \mathbf{x}_1 through \mathbf{x}_{128} represent all possible encryptions for a given search key and the values \mathbf{x}_{129} through \mathbf{x}_{256} represent random values, however with characteristics similar to the true data with respect to their size.

The legitimate user applies a random permutation to the data in vector Y. This vector could look like this: $\mathbf{Y} = [\mathbf{x}_{159}, \, \mathbf{x}_{126}, \, \mathbf{x}_{13}, \, \mathbf{x}_{16}, \, \mathbf{x}_{149}, \, ..., \, \mathbf{x}_5, \, \mathbf{x}_{249}]$. Then he does an SQL query on the database with the data in this random order, and even if the database manager changes this order, the legitimate user of the system will be able to separate the real data from the fake data because he has the decryption key. The SHA3-224 hash allows the data to be verified by concatenating all the fields except the last one, which is the hash (this hash is encrypted and can only be accessed by the user who has access to the system keys).

It is important to note that while the fake data returns no search results, the real data may also return nothing, since a particular field in a record may not have been encrypted in all possible ways by the key in use.

So the legitimate user can get rid of the useless data, but the attacker will most likely have great difficulty knowing which data is "correct" and which is "random data". Of course, for each database query there will be a cost of 256 queries. But it is a worthwhile effort if the server is powerful enough to handle the requests and the response time is satisfactory. In this case, we pay the price of security with a slower performance of searching, changing, adding and deleting data.

5 Cryptoanalysis

The NJS algorithm is symmetric and non-deterministic as we have already shown in this paper. RSA is perhaps the world's best known asymmetric system that is still used in many scenarios and has not yet been completely broken. It will only be effectively broken with the advent of quantum computers.

But why mention RSA in this context if we are talking about a symmetric cryptosystem? The answer is simple: the RSA system and the NJS algorithm share a very close level of security, at least theoretically. RSA is a deterministic cryptosystem, i.e. a message, encrypted with the same key, always produces the same ciphertext. In the case of the NJS algorithm we have a non-deterministic system that can produce different ciphertexts from the same key, using a pseudo-random (or totally random) number generator.

Employing the techniques of fast powers we have that an RSA cipher corresponds to only 2 simple modular multiplications (being even more precise only one). This is because there are only two "1" bits in the number 65537. And what does this have to do with the NJS algorithm? In the NJS algorithm the clear message is encrypted by at least three multiplications, three sums and three XOR operations, and to perform these operations three different modules are used. In the RSA system the module is public, in NJS there is no public information other than the cryptogram.

In an RSA encryption we clearly see what has been explained in this section³:

 $^{3~{\}rm Let's}$ use small numbers to make our reasoning easier.

Table 6:	RSA Key
Prime Number p	246497
Prime Number q	8909161
N = p * q	2196081459017
D = (p-1) * (q-1)	2196072303360
e	65537
e ⁻¹	1845334723073

Consider a message M, whose content is 1234567890. To encrypt this message we can consider exponentiation:

 $C = M^e \mod N = 1234567890^{65537} \mod 2196081459017 = 1033311264004$

The operation of modular exponentiation with giant numbers (as is mandatory in a case of real encryption with RSA) is very costly, and can only be accomplished by employing fast exponentiation techniques. To apply this technique in the cited example we would have the following situation:

Table 7: Speed Exponentiation Technique

Message	Exponent	e = 65537 (binary)	Result
	1	1	1234567890
	2	0	479610188505
	4	0	991099202795
	8	0	1515777153364
	16	0	1521376824852
	32	0	1515437109397
	64	0	693892638319
	128	0	920413295621
1234567890	256	0	1401497722674
	512	0	601437657220
	1024	0	1027144228206
	2048	0	1173322269509
	4096	0	1762446373179
	8192	0	499695743722
	16384	0	1441246836953
	32768	0	1665320632773
	65536	1	222544470823

Then we have that:

$$C = M^e \mod N = 1234567890^{65537} \mod 2196081459017 = 1033311264004$$

Which is equivalent to:

$$C = C * C^{65536} \mod N = 1234567890 * 222544470823 \mod N = 1033311264004$$

In fact, the multiplicative inverse of 222544470823 to modulo N (2196081459017) exists (we will name this value K) and can be represented by L (1977535799413). And we conclude our reasoning by saying that:

$$M = C * L \mod N = 1033311264004 * 1977535799413 \mod N = 1234567890$$

With this exposition we want to show that an RSA encryption with the fixed exponent 65537 $(2^{16} + 1)$ is equivalent to a single multiplication of any message M by a value smaller than the modulus N (which we call K). It is clear that for each M there will be a different K value, and consequently for each K value there will be a corresponding multiplicative inverse L.

Therefore, it is plausible to conjecture that, given the longevity of the RSA algorithm, the codebook ${\bf R}$ formed by all possible messages generates a codebook ${\bf S}$ formed by all possible cryptograms and the numerical ratio between the codebook ${\bf R}$ and the codebook ${\bf S}$ is sufficient to provide the necessary security against all known cryptoanalytical attacks applied against RSA. If such a fact were not true the RSA would already have been broken without needing to factor the N-module, which, to date has not been performed, or at least not disclosed.

The \mathbf{R} codebook and \mathbf{S} codebook mentioned in the previous paragraph can be understood (to remain using the RSA key we presented) as a permutation between the elements of the clear message and the cryptogram.

Table 8: Codebook

Position of the clear text	Plain Text R-code book	Position of the cipher text	Cryptogram S-Codebook
1	1234567890	4	1033311264004
2	1234567891	8	1933023821024
3	1234567892	10	2018464103304
4	1234567893	1	556086417688
5	1234567894	7	1590524243055
6	1234567895	9	1547302878353
7	1234567896	6	1480606826760
8	1234567897	5	701131536877
9	1234567898	2	1522501699385
10	1234567899	3	863040139835

In the case of fixed exponent RSA $(2^{16}+1)$ the order of the codebook S is determined only by the value of N (modulus). When we refer to the order of the S-codebook we are referring to the values [4, 8, 10, 1, 7, 9, 6, 5, 2, 3]. Of course, by increasing our R-codebook and calculating our S-codebook we will obtain a list with N-1 elements arranged in a seemingly random order (although it is not). As we clearly show in this section the modular exponentiation of RSA, in this specific case, can be replaced by a simple modular multiplication on N, although for each message M this multiplier is distinct.

The NJS algorithm attempts to replace this single modular RSA multiplication with a set of at least 9 operations using distinct algebraic groups. The operations are contained in vector $A = [MUL_1, ADD_1, XOR_1, MUL_2, ADD_2, XOR_2, MUL_3, ADD_3, XOR_3]$ or in vector $B = [ADD_1, MUL_1, XOR_1, ADD_2, MUL_2, XOR_2, ADD_3, MUL_3, XOR_3]$ or in vector C (which can be defined as A^*2 , A^*3 or A^*4) or finally in vector D (which can be defined as B^*2 , B^*3 or B^*4). The total number of operations can vary between 9, 18, 27 or 36 changes in the clear message M, and the number of modules used can vary between 3, 6, 9 or 12. These parameters can be easily adjusted and adapted to promote a better level of security.

It seems reasonable to us to believe that this change made in the NJS algorithm relative to the RSA algorithm should not compromise the security of the algorithm. However, we point out here that we have no mathematical proof for such a statement. But considering the fact that the NJS algorithm is not deterministic like RSA the number of distinct S-codebooks relative to a fixed R-codebook (plaintext) increases exponentially (without the key being changed).

Disregarding the initial noise that is added to the plain text we can establish the following relationship between an R codebook and its corresponding codebooks:

Table 9: Different encryption possibilities

R code book Number of elements	S Codebook Possible permutations	
1	8	
10	1073741824	
100	2.037×10^{90}	
1000	1.230×10^{903}	
10000	7.940×10^{9030}	
100000	9.970×10^{90308}	
1000000	9.704×10^{903089}	

The number of elements in the codebook refers to the modular limit N. In the NJS algorithm the size of the N module we recommend is at least 512 bits. It is also important to note that in the case of the RSA system the N module is exposed (the system attacker has access to it) and in the case of the NJS algorithm the various N modules are not accessible to a potential attacker, given that the NJS algorithm is symmetric. The modules used in the NJS cryptosystem must be prime numbers.

6 Source-code

To facilitate the understanding of the algorithm we have made available the complete source code written in the Rust language. The code is divided into 8 files with the following names:

- [1] main.rs: This file contains the call to all functions of the NJS algorithm.
- [2] compress.rs: File with functions to compress and decompress numeric data using only printable characters and easy to save even in TXT files.
- [3] show_message.rs: Functions to show messages on the screen.
- [4] math.rs: Functions to calculate SHA3-224 hashes, create long integers, and potentiation of long integers.
- [5] intermediate_file.rs: Processing functions, hash calculation, encryption, data pattern transformations, and more.
- [6] cripto.rs: Data encryption and decryption.
- [7] read_write.rs: Reading and writing files.
- [8] cargo.toml: Rust language configuration file.
- [9] database_min.txt: Small example database with dummy values.

The source code is available in Annex III of this paper.

Conclusion

The NJS algorithm presents a non-deterministic encryption model that can be used to encrypt relational databases. The security level of the algorithm seems to be close to the security level of RSA with fixed exponent $(2^{16}+1)$ although this is not proven in this work. The encryption technique (although non-deterministic) allows to have an exact decryption (at least for secure hash functions).

The NJS system can protect information on a cloud-hosted database server by performing direct lookups on the encrypted data using a larger number of requests for this. Although there is a stress that the server must endure due to a significant increase in the number of requests, the data remains encrypted at all times and is decrypted only on the client computer. To encrypt a simulated database in TXT format with 100,000 records took 43 seconds (using a personal computer with i5, 3210M, 2.5 GHz, 6 GB RAM processor)

but decryption was much more costly spending 302 seconds. Importantly, the database was increased 6.72-fold. Encrypting a large database can take a long time but searching and changing records can be done quickly on the client workstation with bearable processing stress. On the other side, the database server is forced to make fast and frequent responses, requiring the machine to have sufficient hardware to support this demand.

Although we have presented some analysis points in section 5 of this paper, we know that the algorithm is not properly validated in this respect, and this is a long road for future work. Given the rigor necessarily indispensable in the validation of a cryptographic algorithm, further research is needed.

We finally conclude this work by making available the source code written in the Rust language (annex III). It proves that the NJS algorithm works flawlessly and will certainly be a useful tool to better understand its structure, either in its highlights or weaknesses that can be pointed out in cryptoanalysis.

Acknowledgments

We thank all the supporters of this work. The NJS algorithm is entirely dedicated to Noêmia Josefina da Silva, a woman who knew how to face life's challenges and overcome them with an unshakable faith. A unique example of maximum overcoming portrayed indelibly in her honorable and extremely ethical life.

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Annex I: Example of system key

Encryption key:

\mathbf{k}_{11}	$\frac{79859942866320310995615816391444895386935266232877036021557362979267032608024305957988557}{5172159941160996858791492426951050125434098259893785644324473094}$
k_{12}	$\frac{38949712688808959188135718653498207769865355640446944689388230064769283614335319722839038}{4194426695893807917767545151499602887246947718593727448997248475}$
k ₁₃	$\frac{1652692441615594879688938721164988030021629470487536871620129120090230908908754745799804672338033800526737708514965500441016077256195537282122354800594924}{22354800594924}$
M_1	$\frac{10722467218978268926258696226747927491676137674641045614263412653289622125863632210496158644295701991444380932249987824423920733594443161811202603577008773}{20100000000000000000000000000000000000$
\mathbf{k}_{21}	$\frac{3941381251296432391240385372518378389889668192544602147367559773318242289362065146375652446705743394905107240300826602999576955133001725533228366416498}{2000000000000000000000000000000000000$
\mathbf{k}_{22}	$\frac{679968032814501248768476391390032756738339217273517137049704455723925545469640575828725844859598672641936331271979289600118870937100720954066294927776550}{2}$
k_{23}	3835248109069771657151213531800158263562575766525562242988648811472403732399692957653608768504173863214872485435240209474647530755370512731260367140136452
M_2	$\begin{bmatrix} 14131977366805436785561604598983193438661364711763400248572879753402250305836809006390647865165179066112442451653522369947632069276109866770003652323187791 \end{bmatrix}$
k ₃₁	8701184376278655751331556964155239257951186269523030826634992031106613114893203081022028234991995476452470920166304139172438412582464232396044419717239034
\mathbf{k}_{32}	$\frac{13401271165779039677132077429203632328263144769874312970068148876314128392349501785954083}{025823662536353054832141092803043131351501829397794476306697797200}$
\mathbf{k}_{33}	$73650558202136522578961847987037021097186454439664812513820381124745401730583432481268162\\01085689521196154809391056247962583625831519782811009777420268488$
M_3	$\frac{62751567957182121908589996037575700747066488676349468457385537869782914085904630220317018}{343273624072066926036275655625788047574496799302455205246748565509}$
\mathbf{k}_{41}	$\frac{106988375097127211870602749384275294570332059231887791278038521312263089527954377285336409952542023814717712254977581950219094673003828603055290712969021149}{2}$
k_{42}	
k_{43}	$\frac{69788261443455392943670574880623061062023346236205339840506429701634279754909208766209825}{809871772721626107228624907791963563980572861272445177425858810129}$
M_4	$\frac{56318219381744816544439511826752610591319102858264549263126651054674431571364048378997632}{7078295689508583931189713378020461035111061571711358175709862669281}$
\mathbf{k}_{51}	$\frac{317367178828690618632118206497405695612253257366419375002247340737375277156645777092645969343443984728120217084563821153744606396051632248214402428190943820}{2}$
\mathbf{k}_{52}	$\frac{33855909614234757713825352827420391274856621863677939168223138347061081927775922134325113}{3863422154101312082161426639930879444864114808225291457998801536749}$
\mathbf{k}_{53}	89200174180240769003124910785659935975568174525955772008586876752967346230835068674249612931710530071211780004404035322416526103821775522237486401422533964
M_5	$33329743085315113737010539623297528881942999390631349022533616076173198088719102556131655\\56452639960948605869910003763383801375060793785465554200995461907467$

\mathbf{k}_{62}	$\begin{bmatrix} 62078941727844029018562430691641600148906615539760974904152320361983552898878086338072462\\ 5315105746738160711515857826425880862823431747203180962260225557075 \end{bmatrix}$
\mathbf{k}_{63}	$\frac{82761951118548026171665188689023929724371871054693335351606328165846865807392013931608453}{2288748260709746842094866590537398443302680752591017630247877018164}$
M_6	$\frac{99573602313013214890895192703276583876676791337236341796997645988639569160753269275188196}{00196215681241225228051832571702783982090495383748302721565963337207}$
k_{71}	$\frac{35003842897595275488593485236584007234296697263908474131710746107275869188057085136303883}{10382189102763575543001892976610026095416646397685085920654636174580}$
\mathbf{k}_{72}	$\frac{48038991922995433230632041599925761424532221305292720226724449796326861499282267885657359}{45251474418239196023030971051538936883008977042636655287059018166508}$
k ₇₃	$\frac{58987590331259835327523673564078773249192890112121417041530501183387851831438843038086181}{79712495623483198034552686468407356268370134766461896934886729989200}$
M_7	$\frac{39336567463872650496661187719967295324131671510484908235533132078268965879985258102876153}{189975637397803255967592003990034868715526061633664076597821547293153}$
\mathbf{k}_{81}	$\frac{58322910032767642749125845422490023323794838826454813770832496247564439479337065191241799}{04353686527271447612249232616222101378568238038791702869116715412384}$
\mathbf{k}_{82}	$\frac{20667359221089305799651070270390355608903182165393345873196101035759543152442880133277661}{55639207932505253195089353084835441222656328647980638370435195223679}$
k_{83}	8517096333202442940332690116185642973007465316314727214817751099953687312168758973588885647323336295639280105466124593940423282734707279076090573637085090889
$ m M_8$	$\begin{bmatrix} 17642302649280537945253519667519774118884596216709714632191493770377727497144907906504078\\ 0065967098297333212276392743191552230738505519831810235757400980386471 \end{bmatrix}$
k_{91}	$\frac{13507438403455255706724055352895150509589959270452191297861659327605429880306341007002634}{57212457699324987213515820194852225349473294173470236326449624180564}$
\mathbf{k}_{92}	$\frac{46708533305859862246342212395656626817956263986528451684539220950284947582850968288996469}{102829137996028889939932253615710652507912419720255383696419754095008}$
\mathbf{k}_{93}	$\frac{89863313163967525687555563825704682515897308079485684094045147873611610871146150758815673}{077156715279412146409320241309980819503775179165765310459810392097430}$
M_9	$\frac{26515645151195228084453144221862383016684567652219022973054189202845637336141316449287721}{4978703953815327028686632303249577196727133760024471154069296166709617}$
k_{A1}	$\frac{17761668501285473553082503814022867024281423756331360535753353092030137327693262410648282}{3869389442271099009008253407617780462477030070460389023891643549384758}$
k_{A2}	$\frac{12307934745841656314236506693947789169770726545251550232413974790307816067968572643640196}{8264126233552740948506939798378342033491735883412403247875785573802278}$
k_{A3}	$\frac{19085468430034164508690178824501640446773070236234221397112424814629212573529505878863918}{1172579016001234921786707930677227420135596224864681373291165079328504}$
$ m M_A$	$\frac{33472743437802238628232913416469647259031995665490214381969900993319961831052072424776937}{61428579390816977510197372565370535689628837057251024553191256985937601}$
$k_{\rm B1}$	$\frac{16291413379806191998310257157307071695782234925453539712760587087772823923687488532281396}{25142330463138009979138144450660871209198834994992155138819752752139914}$
k_{B2}	$\frac{96611745787633898896747381304569411984203601392418969801696773350477520351284762022106693}{8931175651215660345952187457885434446403919701747723740130023475656217}$
k_{B3}	$\frac{17055168194510082814534981625915296297672355393748241841663982835024337034726302312298674}{53581482908661022261923489950309410526603388975337157910467857787639168}$
$ m M_{B}$	$\frac{71906169247722049318659987421117273548263388179005807314756475373808924882203457851093233}{88049519310522590087131429105564260801558422560148681900427878042427877}$

$k_{\rm C1}$	
k_{C2}	$\frac{23824064693359259554861463289975866469588906139761712641988273604471077261508001292222741}{08165663496353116745052155104878171493827799341793008673752655633285596}$
k_{C3}	$\frac{42505579115717305913948872663460264689492004477578030477039025278694294274175075394391187}{27832879733822553991227675505750080624780759023309103784838956065341027}$
$ m M_{C}$	$\frac{54501625714709153513679587851807791930939770596103501005387258885974417197042133106478994}{971207681668450887172994046333566119814339928055166169672264611027110417}$

Decryption key:

k ₁₁ -1	
k ₁₂ -1	$\frac{46909210154436142854196599732595358204466507570120817780852957029638088852277813967765120}{34061141177653033055237065260332117677781772365708940493776907720}$
k ₁₃ ⁻¹	$\frac{16526924416155948796889387211649880300216294704875368716201291200902309089087547457998046}{72338033800526737708514965500441016077256195537282122354800594924}$
M_1	$\frac{10722467218978268926258696226747927491676137674641045614263412653289622125863632210496158}{644295701991444380932249987824423920733594443161811202603577008773}$
k ₂₁ -1	$\frac{10190596115509004394321219226464815048771696519218798101205319980084008016474743860014995}{422378473322717537344413221543344632492320976865044470423956771293}$
k ₂₂ -1	$\frac{74625408217079413041868889062488099372953616766749096609017492086647883538364688139851566}{1065048489419473944580260803556822874234435848619335362737271967}$
k ₂₃ -1	$\frac{38352481090697716571512135318001582635625757665255622429886488114724037323996929576536087}{68504173863214872485435240209474647530755370512731260367140136452}$
M_2	$\frac{14131977366805436785561604598983193438661364711763400248572879753402250305836809006390647}{865165179066112442451653522369947632069276109866770003652323187791}$
k ₃₁ -1	$\begin{bmatrix} 54050383580903466157258439073420461489115302406826437630750545838676300971011427139294990\\ 108281628595614455116109351486615609161914335070059160827031326475 \end{bmatrix}$
k ₃₂ -1	$\frac{55166610868712856244865043441162380233950704810569034232218842807868814522154318307410670}{775636551710115116263161270196905021500080324071491294262481378139}$
k ₃₃ -1	$\frac{73650558202136522578961847987037021097186454439664812513820381124745401730583432481268162}{01085689521196154809391056247962583625831519782811009777420268488}$
M_3	$\frac{62751567957182121908589996037575700747066488676349468457385537869782914085904630220317018}{343273624072066926036275655625788047574496799302455205246748565509}$
k ₄₁ -1	$\frac{45619381872032095357379236888325081134285896935075770135322798923448122618568610650463991}{7125753665693866218934735796070241940438057743108302884996893648132}$
k ₄₂ -1	$\frac{12662894560411496715599214033734503578697658366294305692346068347071341491024804555663575}{158889760211902844728485199053395544938089443755939917310618134632}$
k ₄₃ -1	
M_4	$\frac{56318219381744816544439511826752610591319102858264549263126651054674431571364048378997632}{7078295689508583931189713378020461035111061571711358175709862669281}$
k ₅₁ -1	$\frac{30156071297028207550689357558323471925820466816967155272511142668799445317152644785205195}{87109195976220485652825439942230056768664742153217339798567270963647}$
k_{52}^{-1}	$83607754988894773204195816239129479131097760255346625357992622653089092977547310998772792\\2380128680656188384225302151012056729182016324868737300420198651949$

${ m k}_{53}^{-1}$	$89200174180240769003124910785659935975568174525955772008586876752967346230835068674249612\\931710530071211780004404035322416526103821775522237486401422533964$
M_5	$33329743085315113737010539623297528881942999390631349022533616076173198088719102556131655\\56452639960948605869910003763383801375060793785465554200995461907467$
k_{61}^{-1}	$96879749985301351437454124584019090448940593986593509972204701879508929903320512336282788\\20742340746957358448413443487317874892963994939057101291976907528412$
k_{62}^{-1}	$\frac{12992796215483682498372734500070520323481541289985898822149177506249028320068754134865794}{99511852506225864082227915489122117831327292593168904954837841337349}$
k ₆₃ -1	$\frac{82761951118548026171665188689023929724371871054693335351606328165846865807392013931608453}{2288748260709746842094866590537398443302680752591017630247877018164}$
M_{6}	$99573602313013214890895192703276583876676791337236341796997645988639569160753269275188196\\00196215681241225228051832571702783982090495383748302721565963337207$
k_{71}^{-1}	$\frac{35836183174113122947801839196308894600702001784094060822362057467541378961179549589245764}{879593448295039680424590111013424842620109415235978990677166911118573}$
k ₇₂ -1	$\frac{27022493229487031671447974820306537179927272650829693984524487079073864903469408172704027}{770681993450029548163333926030892816967336219810621844602569321637611}$
k ₇₃ -1	$58987590331259835327523673564078773249192890112121417041530501183387851831438843038086181\\79712495623483198034552686468407356268370134766461896934886729989200$
M_7	$39336567463872650496661187719967295324131671510484908235533132078268965879985258102876153\\189975637397803255967592003990034868715526061633664076597821547293153$
k ₈₁ -1	$\frac{17059073548952861517762261213294873885646647828445166494483168807902083102351537254591660}{0161613411770061764664143510575330129359937281793018532888284264974087}$
k_{82}^{-1}	$\frac{17550227059375372237214690178144965381622493307463874312388996443782860877532794833333310}{4244349172997529785232911718598220394449346773574207938573270075999885}$
k ₈₃ -1	8517096333202442940332690116185642973007465316314727214817751099953687312168758973588885647323336295639280105466124593940423282734707279076090573637085090889
M_8	$\frac{17642302649280537945253519667519774118884596216709714632191493770377727497144907906504078}{0065967098297333212276392743191552230738505519831810235757400980386471}$
k_{91} -1	$\frac{26380570767160675527385903668333431511588668059514501060075572609569583037338253039217695}{1521491496116002041473116483054724971377660465851000917742846542529053}$
k ₉₂ -1	$\frac{19192497812325753715947538795646491437789916569660953385218804216891469175065075746327080}{9251983718727573571024549840505407353605970963598335105504580222137717}$
k_{93}^{-1}	$89863313163967525687555563825704682515897308079485684094045147873611610871146150758815673\\077156715279412146409320241309980819503775179165765310459810392097430$
M_9	$\frac{26515645151195228084453144221862383016684567652219022973054189202845637336141316449287721}{4978703953815327028686632303249577196727133760024471154069296166709617}$
k _{A1} -1	$\frac{31696576587673691272924663035067360556603853289857078328394565684116948098282746183712109}{37559189948545878501189119157752755227151806986790635529299613436552843}$
k _{A2} -1	$\frac{16777729356905918783318414669976457264212568580897743863523967696199069151478253310557824}{44672971075779064328651566999867139042234917358055068772216323896006193}$
k _{A3} -1	$\frac{19085468430034164508690178824501640446773070236234221397112424814629212573529505878863918}{1172579016001234921786707930677227420135596224864681373291165079328504}$
$ m M_A$	$33472743437802238628232913416469647259031995665490214381969900993319961831052072424776937\\61428579390816977510197372565370535689628837057251024553191256985937601$
$k_{\rm B1}^{-1}$	$\frac{55614755867915857320349730263810201852481153253552267601995888286036100958515969318811837}{62907188847384580107993284654903389592359587565156526761608125290287963}$

$k_{\rm B2}^{-1}$	$\frac{59546631245252356888398909394113473475509524159676893203146481744278826237800484284387416}{34375738396937361308961230100995076218153351169712180851384167372565182}$
$k_{\rm B3}^{-1}$	$\begin{bmatrix} 17055168194510082814534981625915296297672355393748241841663982835024337034726302312298674\\ 53581482908661022261923489950309410526603388975337157910467857787639168 \end{bmatrix}$
$ m M_{B}$	
k _{C1} -1	$\frac{52409478096216568524397493211846751752519800961934140401372733705269441354140477818840133}{127240500453063148828399832711538534305825282824928359995257071194619526}$
k _{C2} ⁻¹	$\frac{44600100520712772545029384541829422187201732511866482786918855121853130557404655702518681}{485631023417772288774363854198398244970249941064254119317535006939329779}$
k _{C3} ⁻¹	$\frac{42505579115717305913948872663460264689492004477578030477039025278694294274175075394391187}{27832879733822553991227675505750080624780759023309103784838956065341027}$
$ m M_{C}$	$\frac{54501625714709153513679587851807791930939770596103501005387258885974417197042133106478994}{971207681668450887172994046333566119814339928055166169672264611027110417}$

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```

Annex II: Example of a table

Table with fictitious data:

1	['Airtom Fregapani Carpeta Cunha Merraco', '382248368-07', '(41)98136-7571', '23/03/2019', '21837.88']
2	["Layne Bellani Dallarmellina", "257253618-57", "(16)95336-9701", "08/11/2007", "21022.53"]
3	['Yuri Buchmann', '179790017-93', '(45)97040-6314', '23/11/2011', '21288.60']
4	["Sibeli Ceolotto Grosseli", "991276898-05", "(74)94410-3803", "09/04/2022", "11690.95"]
5	["Soila Giacomel", "634814762-00", "(48)91308-7529", "21/04/2022", "5269.11"]
6	['Anelito Broccardo Bourguignon Balbinotti', '906757754-34', '(62)99754-9447', '13/07/2010', '10543.24']
7	$\label{eq:control_energy} \begin{tabular}{ l l l l l l l l l l l l l l l l l l l$
8	['Jonatham Caruso', "338338942-18", "(96)99154-0933", "06/12/2018", "9825.37"]
9	['Cipriano Gentina', "860563543-65', "(67)92805-7573', "22/02/2015', "8507.21']
10	["Tailaine Bozzoni Anesia", "286895488-10", "(89)90742-9484", "28/03/2020", "11678.59"]
11	['Edimarcos Celleghin Rigone', "246561540-73", "(89)92624-4731", "15/08/2002", "134.58"]
12	['Ivano Dobger Green Claudete', "002409746-39", "(97)95326-8142", "03/09/2017", "4601.07"]
13	['Laynara Bacelar Bentele Cattani', '138044932-14', '(32)92012-0287', '27/09/2021', '9088.54']
14	['Walner Mangieri Crass', "892722759-00', "(32)91775-6893', "20/03/2016', "9674.10']
15	['Alyson de Zan', '958168983-24', '(41)94519-7954', '20/08/2010', '3426.03']
16	['Noilson Buhnert', "565435673-27", "(17)95770-0423", "14/12/2018", "3532.94"]
17	["Rose Bizarra", "834993628-07", "(87)95527-1914", "15/03/2012", "14800.47"]
18	['Welder Guaresi', "764225432-77', "(45)92426-6847', "09/02/2013', "6279.69']
19	["Erique Brunkmer Maiolino Divino", "598989914-28", "(65)95271-0645", "09/01/2006", "861.20"]
20	["Bernadino Friedrich Brantano Fried ani Clelia Deolinda", "808203951-03", "(5)95105-9519", "09/03/2018", "21192.92"]
21	["Aldivan Aguzzi", "485410475-10", "(54)94240-0122", "09/02/2013", "1535.84"]
22	["Emerson Moura Filho Dias Bierbrauer Ferraz", "697960271-56", "(89)99572-4141", "23/08/2017", "21738.87"]
23	["Sheilla Faller Farinon Barneche", "756795415-03", "(12)95491-3910", "13/07/2003", "12639.07"]
24	['Blenda Favatto', '194238594-70', '(84)95599-8669', '19/11/2005', '16011.41']
25	["Kelita Edilaine Barboza", "784825094-63", "(74)92803-6674", "04/06/2017", "5262.86"]
26	["Verenice Drugg", "793829161-39", "(19)97986-9204", "18/05/2007", "2967.25"]
27	["Jocenildo Aquilio Winters", "464520485-70", "(31)95886-0344", "21/08/2016", "5846.81"]
28	["Mairia Etcheverry Cuomo Zaccardi", "321857950-29", "(46)95030-2306", "02/05/2013", "1017.65"]
29	["Jefter Gasquez Legnaghi", "210287841-20", "(95)91085-3949", "07/06/2016", "6406.82"]
30	['Ibrain Eisenhut', '037333705-85', '(86)90453-8840', '21/09/2003', '8821.55']

Encrypted table:

1	$ wnU90L@00^LZK0C6", "U60FgljH90ZFVDl,p;!w<100=:EqgzGb^ab$
2	$ \begin{tabular}{ll} $ ["e\sim \{FN100!; !vF,gzh[\sim xQ^Nf' 10St0@Umlwzh50J@-e6020(Ots\{b0ihud30z[I@Umh>=q\&bO0FEL:0g=20wx\{7", "p=60VqfnjxYcyX(/6090UO30,0k{ } {)'o(CnZ60zte] 30k$EQ#Df@d,0! [qt&Qw=u70D[tNbsbCa#", "60x00}\}%=n]=;/0pYKn_:60F80#d::nodrp,@ %30E^0V10]BL00800w0$l50Ymj30a50M0300-fUx0@0p+lk]O;ug0a1010_6", "10-aR*KQ!;]UUs`$70QzAPO<\sim M0yv} {40R$Gs{duy#90^VvpFhw0!FW30[0't/0T}S imCB-k[770]w40)#[kdqe0 %[:@p70b-iqHYWll'K0m`N]Cwn0BCx$100X0o050T;2000C40HL! Y0Yn 70sfsO\simk 70?yglxtAbKJ", "b`10900M50:)O0NnVC{Vvy} 30O,)'0aq'}K0i(OO30nDz>Vm10`(&60s%%0u$Sy10~MgEzy#0L010!H;[(70HJG?Ce"]$
3	$["0m5010aDmytq']?=P*M20F90'y!/J:jAQS>w,\$N\sim090'0:,?\\ [U\#mnGa0d0y20c'0m_VC+uAsM30l\$P0H10q;?szlx9",\\ "[R30[x@S{:w10}]t40AL\#M>00t00{Y30q0@m;y]c'(FWL{m}]c70r01040-%*MZ&00VO\\ \%lL0:50)V0f+t50t0eobg?]=/3", ")neX`uQObG_`nBP[!00[*0+0Qgkc0u400300pd2000 $30MQ0t60804", "ZU~T^Wu)R ^A&<{u90Dj} sP{%0\#jVe@*F00:0f{% ieD", "20{V[&} =G{Cuc90l/B0@J20W} +{/*U00,/Rq*d70~+b!+*snCN80pPm(y0Q=_dU`)fCPuGJrY0Ky90e$0#i`!9"]}$
4	$["Uo]00][!\{t(+mdt\sim<70C@bjG_aCN_900?\ 30QKB0dR+@y60HDwu)=GeR0novSrm5","\&0=90DT$r0,w0BI+;e90bv600\{\#XD0JRw0dH 0g':\}\\ [30\$W)T100mouW80cOW0:?\\ 90m0DofP+D,K]*10SiN50abOpq0ELnEu600woeUf@Wh^i+s' dH?%Z< x<","10$90z-? x0mEWe,10mSzQ!PUK@Xr0=E<30j40*>uXO10?NHI@T0mhl50DT30b50$
5	$ \begin{tabular}{l} & & & & & & & & & & & & & & & & & & $
6	$ \begin{tabular}{ll} $ ["] A@800600_vq50>O+\#020h/O_? v2040^1\%eq]'0EY]CDz60E70;m !KsbR040 \\ D0'z0*)fp0Q>80030Yq1002080D\{050?m]qefB90~]H?Mc 90oc'G0$y90C-10MrG0/V? \\ XfyO50xB900$80_yi$yC!kOGK[IeB0]a0Ns_fmod3000'0;yc0]50$$`2", \\ "M/10*0&bN0,hsC];;pjSMU60(=P?tJKnO-ilEtV^#$4:90+0@-{})", \\ "#70uXpPUa0l70(R#070R[/`dA$@%&N~{}nj$*s}]IooZC0Gkmj>U)80J#}-js}+jC0ob#?] \\ [trAQO30:Ai]u0q", "_WnU50@0]50sEY*m>30b$110iW{^4030Jj;20LxBGET$0,E>aRdJC80n00UaH b500]EkN20'0K&rO?50g[*20@70]`P7"] \\ \end{tabular} $

7	$ \begin{tabular}{ll} & ["urFq?{wuk/k^;Sl\&u0?0x_[`t+_f z^Qo10b50hyJzjQ(B50AV0Ht/X^Gwy}n@F>'u/v[s^*jtsWhPJE<`5", "A0qg[_#! z^Qo10b50hyJzjQ(B50AV0Ht/X^Gwy}n@F>'u/v[s^*jtsWhPJE<`5", "A0qg[_#! z^Qo10b50hyJzjQ(B50AV0Ht/X^Gwy}n@F>'u/v[s^*jtsWhPJE<`5", "A0qg[_#! z^Qo10b50hyJzjQ(B50AV0H0; z^2vVWM-P YpMJ60! z^2vVWM-P YpMJ60! yO070H?;dcQ50KB$0v0AZwj3", "#0V00Qpxh0`buTls_}0_Swld;_$k~G'T0@~Hz'X})%=AIw^N^=o %00G20?!Y`ShP90_];S070rN/`0700<0%/{/0cD", "20V:z20b;O~`)hf^f^e/Xh+_V10060Mu80q0nG$uJ^030Q10a$iZBQ80~50^20Zs00+X]}F^m>wu*th(0 E40d0YSvhD0N6", "=-600vcm[)80g) W@cS/0f0v50m70-d2010Mn070aklonQ0C0scUssb fmh0i0a`030CF80u80,/v60 @kw>g{0m0Rg^qlC~", "10~0;A0+t0w]0L0$DMx60)/*V;ujz~T{}AMI,60xl)(G50elOAQ70Umkg'f/_Jn<20rWg70i^70 Ui0[lMGk~ta80-)*] } \end{tabular}$
8	$ ["+[\&m0\$IDd[RE(\$(f0u\$rDq_0O)]90^=x'_w10T0j10r^*'0' h90ecL \\ +0x/\#X20/Y00Tm60BYv070+0jsKSZ50v(40v0", "+ovXK#HL@_%080Iqwd30eM~=0N-hu60:x900T\{Mp!*1qv:60\&m,bt0AV*UZf0r\$M@0K[0f0p\{00_TxpMX)S \\ \%u0=b[>", "s0gUz70mR-((\%k0W);900HC\$C_ya:r20_0nX]w -V^0b0bm!EOXI^-< ("[0Kg+0if0\%Up]uQ0Gv}C90e\{jw)^*q=0^*", "@040\&;Hz!0teW~{}^xhEE y \\ 0+p?]:l\&&E0[)F]:Osy=(`z;00TQ?=Pt@/E`80OVa=0000t\{QC'\#bb?50X10Ftw~O90X6', "=Rv X)\&cRw=V40[+LYc800^0V?=RAl@q90D:-*E`h 0:Wz)}!]nAr60%^YhD/vxq+p%`{}V0B30f %;dd0Z_r(m80S"] $
9	$ ["F^vzTPVZ20b60wGol0_vI0en]i]30M>`(WdH060^e?K0cR;]GD\#/F_k0V+ezl0*(@-+0t20VA/kD$
	$30R; (lse^L0L\{", "mC0, (L+iurz; 'c080AW0~0X80ri`*y^DD`o0*kV, X?ah'0(0070FGK!0`v') 100L0\{C0$Ceh>qE$
10	$ \begin{tabular}{ll} [",q)mfdY-400-U0IPi`Li]03070QiZT$0uNtF+LX-s-}0\#70KFd=0YAqsctCB10QE)\&d?jc90A]) \\ (0z,IMVL00Mqip70", "KN/[$
11	$ ["]QO ; *oCRR:D(@\{wkCM0X=050[x60,):h(BEUb\}ymGvU> *\#aG\}`IGMh\$N<70\%090\&U,h\#Ec@f\$00\%L-sR0n300R", ""?`a[O\{Ij_gIX['-HV=a-(c0gszh\%_k0qaeBv-L?\{O20IJC;mtyb", "%(lpI,j~G'x20qE0uL'0j30myU %VBfC'i0U=]70C:z\}L]DlTcYO0m00d900JUiy}wo0LNCzF?YuUXybMU')M0**30"] } $
12	$ \begin{tabular}{ll} ["\&PPo%`soZ'90_i,EYP60Y]GM=0>cmry!60_W-Cr90yp>Th400:0b }EF)m\$sYu)/aX! \\ S`\#=c`P\{0VW0\#kq10:_{0-"}, "n-S4090U\{0(_{90}*wUX0\$-10AH\{0!+<\\i40f\#>30NT9040>^{0}-zIOchs902030 G?>W00O_N-00C\#?`W>eJ70OzbA30Ie\{w0", "?'/0*600Z_0g4040@0o/A=:dK70n*0=AYap0l=\{0X)fO/;giq T ;O)0GT/&", "OLN90Bc@baT0]:r=Y j^$ $240_$`TO0WID0@b-lbx*r\#qD@#:r^{vN,:80}_{u}'YOeM0w70Iu>b103000T00HJgc40", "#TQLFNBrouQWs0O@=n90/EBpP/90>8040~IN00#)00RJJay 40RON+S]Iw00[]!nU`KWRwrqI? ff%LH40nl sO5010U", "30N]aL`h0Jt{KM`$_70^}Z0g0W0gT^,x+TgS*20,10x=V\#GfdHuQ60XG=QuC10Ljmk]S^Qh-ZW0 N200Gd80d0C=@60d07"] $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$$

13	$ \begin{tabular}{ll} ["I/hzx40QG\$(QINyS40800CaNCjp\&00P0s\#El;Q0iqVpPxV<(X)EL0a20Q\\ \%q0z=I\\ \%G/\#I-0RRu40ddkJ040k40=0O`I`b0h>vn40Tp%oxYtbxoL/YR", "?\\ beG040r(BbA*o500jo\{040D3060o$31100i0goj;palz\{0Trc080\{=ba?W\#06080J0)Tz\$80\}!\\ 10`;>}-DhG\#rNpa'\$@90A7", "S5060Z0,70\}is*60i;50)0ix*20xq]0C\&zXB>QD80t)Pnqk-w90Z30t,]>-n=r[}Un50'0z300-T30[{Q608010*h85}%\\ (D;v'>0/0F}\}0n40jINWw-z90e'{})@+dA{\{Vh0`EBv&0iy0)0L~90e!?t8", "XR00yY[p>!JI*ZSN?}\\ _R04050zWAh=yG0DIKeXNt70smpkz{p0j?0g}=v9040vj;Q$
14	$ \begin{tabular}{ll} $ A80=0-90'$ &h30=60>J'DPMKD[00Aj10p70rF,20^mXF!60'wA0_%k0x0-? \\ Zw[[40MI'G', "?=YY*iPKI<60wRL*(kndx! nQE90O20^80MEn<60fL6000?;lJ0<0^"V20%Kh40lq$j=80 @w] $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $$
15	$ \begin{tabular}{ll} $ ["+q0]g0800wX+iR0gb10X0lE*XP\#kN''\%u>fvE*eFjp70[]LpZMSida()BD]=+pAGH90EZk^HY? \\ z40@O:E[Y0\{p3", $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$$
16	$ \begin{tabular}{lllllllllllllllllllllllllllllllllll$
17	$ \begin{tabular}{l} [")@020fC'SGg]< v] $f0i010Y70+$W?T20$)NjUexd70$'lK?z60-Qb60gE<-'Q20bc? Pp&i'0rJdmaP$0@tnA-802060X0", "F]0: :0/T'u-80O10vsx>IH Mx0,R>SzPF`M^1060si>XRUKn0@#+60\{0v\{A)]liVOzf-&TJV_$0PRS0d0>(40000yi", "Da: [KKmq*g8BkJj`ly)g)Z70$Vk'80e[a/40[[WA_%yNxJ_Jp0f;oDzMNrVfDPq$20KK0&'d*i&:t}e0! *CU", "BApIq=ie'_mK[xT-00!*0`(30c_V0]\{SN\{\&::t2030Uw0-40v-*ARh0QV0X $0w?N %0i80SwT;sa'wbBF0i='mA-B-x', "TxL(,VqZXJlB-;ygq!w,aTDn'Jw30+ML, [P&f700RkGujAg200%20bFPpEE50aL%'<8060?+T;30$=_Jw010Ef00/hPq", "Jnu=;eKU#WSDMoC-s,S[^SP0VUkD0j}v'MwxqT`u40]QZ>a^)`_>^0ij_0bY70TeIY)n)RJSi'P_+FG70YG7"] \end{tabular} $
18	$ \begin{tabular}{l} ["80\{10wSHB $$ \%FE\{10mX90jd0QZ=n0\#@OzDO)\}B\&KSsy0\{Cir^*\}Pks'w >u'W^^0g)w[m'x)r^>-%d0'HB1", "$sh&70xL060@]s#dM00LeU0]rARmB_Jiqh10r&'wkE! $$0Cz0uu[wm00Tn^*-{MO-30", "}:abZXUm0u'40gWDy$gV@[Gk90e+50v0D#>80'T O90'/YaSGE70[a80?s`PqeH70V0qMliq,0lh=lbx70U[!]dA]3", "@1010-0`T0U{vtil@YgY,-#y$10wVu0>s))/0b$;050)M0%)0hAKwj<:80qf90`ShXO00rKOKG:r10<}0:ln`o(ML 3"] \end{tabular} $

19	$ \begin{tabular}{ll} $ ["H0'Q)m0fU0/[rSArmFATtuEd/;vD7000YYZKi80cJd;iSW+pZ:P]r@50-030vURH-60-L60 \\ C{_j<0M}08050ec0=;",">DZ!:I*00Xf0,UO0{&&XS20(I0fDj>)N^40<=R700+0Nm!dG_00'H03", "^g=TQ0k,qh(KT$/)AYf&+070#j0V_y0[b10l_uH %b0jcJ0A]OPL!J090EPE80/10'0Rve/-&$TLO^{`F0eT06020!{0VaS0,$&'H0>0'sng^{?}%iJL//l0-D,Ye0DI%-0{Aw/RTb0Np0p}l%:;/f", "y0+FH00=R50Rf40<]r[60pS$zioL#v*zOti)-QeSVCii_Uv]60>YJjB=ao*0-*G %B{Z0+*KOoB-V}30IM_DE4", "UF_^)y-001000nfc70vcF10kEbK}T0YRr4040f*QWz)0ynj=B30kZlUf_x>-d50X[200q70+10'%L600f N,10T'rt#E#Z2"] \\ \end{tabular}$
20	$ \begin{tabular}{l} $ "s\#040CqyGIM8", "5020lf]W-;C*:=-x0yuTU,@n/CGu0zMyBG80E70t\}0e0oivn4010Fu? \\ 3080I*Ae50MV30s\{,[!Wmm%10+lI0X80tu^?;]", "!0\#R60/Z_0xOFR^8000u0X40Fg0=^XO]gs20 \{P0r0x pURrAk0xGz0r050c=2000U0X40I(80-G[>)0y(Ft)qE=3", "/zzC}UAtR00shH#QA{?!! Pu70wi:0<;+O%R80 KJe:SW900!g;A30pO0z0S0vHAe>m90-30ezQt0100TWVyIuaOrf"] \\ \end{tabular}$
21	$ \begin{tabular}{ll} ["500=R90Di*Tq30>(\%F80T;Ov>olPvN[e{Pz~K20}60U*+000uPXj0g0k[jtw30+jhh_&I40*ew:0)400 \\ 900[\#0pr0(d>0/60C5", $$$"FX_/0RSe;WL80?/0SMfym500``af_f>#~j0:RyeGNX0B3020[`s020~40h*`ugbsAX60L`NFG10ueG PNM:pyStuCm0", "(zg`>IV=xg90uB}qd:;600iF0}HevVGKO`0a00XJQ=20-auX70c*0;T@Rk0~;0,}0B?LOMDc+Si\#'U/2", "40OQ:G-S80?pC(P,0&$CGbH+;g? $$$$U{=Wb$\#WV00^xVXiO80CHr@0[x]10)0EUTYK\{00S7040b0L000vZM*>`muU400t800oZ8'] $$$} \end{tabular}$
22	$\label{eq:continuous} $$ ["K-Mj0TP]80H=@e0\#[@wx+-rpY0!0>=Kj:T10T000;S0+-/VuUt/$$ $$ CPp]O0yLLtfh00t80clX]oE,g30ux'+0mHW07", "40z0s+u60 aAH]k30d!$ $$ Dy@J0vnJFfdn*dOAju]0^z800IvxF?s:r050ZKIV0/h$#90INTdh80_u0/%yu])Tk'Y(H)Gy", "H0iitU50600[Y#%wiV?uGj00EZ300'*BIV70{!} $$ QfNK:RF400vD90Y0-ba$
23	$\label{eq:control_eq} $$ ["x0Sh0h tYE?0EG]au00h-0'L)u0=(MC\$=+j+10y600@\%y20ZT70/?010!80\&picp'0_eqG WyQ]Z40yT:0z^@IIA{vk}", "Ultg@nC[0ltR0'\}=A?F,)\$]_OA;W $$ -kO2050+m80>=0igyN{#0(0Pkym=^0i,GS?rK703040*0c+70KF+mU90Z+a0+4", ")0'0:g0*eW090.jbYFJ?fPM`070#20_?+Q0x=!L0GU/0{v+);;iQo60)OP$ $$ cWg0:0A-QkW='+cV?-,H)10%pOY*300@nkP!/w;0ECk@ia8", "h#tLzxl'Jlqj40j-40050ul_,Gq1060fS0-20;0f70902050nPPX$Y@Uc)r>yAV]Gn0NDE(V@PPtQ}U5060=!^y!070^Dm", "?!>f0:p90X''?CnZyfm>0dU'H(/PRc*M0EsW]>);50Vz70**h0;Gx30S $$ GO0PxP30e>Y0Q-/JxV)]&$
24	['gj <jl40y fp*g,g-c0e;<jgv0x=")#g~WKRsESE%Y?YnGZjd/X30!ShP~">0PNza\$ 0fiGDupL^Tv;? pse)60NU', "q;XQF>l10eO0n0Nje0U%50r(&!0< Pfwzk? G20cR<{Mx,JA40d}b*a)G`ITm^tWfBuzie20a/;[oWv_00CVF}jJZ70', "0Z-/ZK70{&050^ue^TG+>lp/{X;0Y #oo[L<0T\$10^&iy0<r0scfhcj0n80dynv0a!} ohyw="">50-0jI'l0&qoVgtp<,", "2lH'd; 80>v/)H[=010 QvML*DCEq+'10g70r40) _yS;ApOi;I40f*0&vu30-0+]F60U30(C:dSD!NjOV[r<0Lax q', "C#{100X<yy}] ")ag100va]sbk0="" 090d80kv}60qh)*le0vva+klg40m'20y<&%?30v00lo5000f='20JVC0rdP901"]</td' 90?vso00jgecor30q`hjy`ra70-080~tky'c:la,a60tg0\$6',="" :ug0o0nor@="Q60#,/Sa0?" [0xgki5060l;0f80ak@jg^;@30<l\$~100:=""></yy}]></r0scfhcj0n80dynv0a!}></jl40y>

25	$\label{eq:control_control_control} $$ ["\%:,mc)0UEnXq?(70[0E\sim]WhHyT,l:900WsH+*5040+0'_b-20@-dr0D = `E70XDbeLy:3030TV$
26	$ \begin{tabular}{ll} $ \begin{tabular}{ll}$
27	$\label{eq:control_control_control} $$ [^*B'oMQIeHf\}=? fk:frA0U=Aco004", $$ $$ [R:frA0U=Aco004", $$ "H_o>r,:N!:0MsXBMt-'cSwM50es40! $$ 90\}aXIw20$0sENJmH10'r;2", $$ [CW>mhynRS20f0\}0[0G0qX\}h w09019040N $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ 900:YsydZj{[q30nd;Enlsy!y0U+eH$zVY)-C#=IyXP0iJ^Umi}uu^2", $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$ $$$
28	$ \begin{tabular}{ll} $ ["mNL,o30T0iiY@40V\%0p?TZMEs0LJ0Zz, {20v80? vq$WiN,WF10{^MY20lc50y0sL$-T{RxY#uer&D4", "@0e0z20R]F!00i! o-U(j$I&u20)m0W{durMRDE0b+Sp=/HU?JW80f0W80Z+00Nah=? w+kIEC0-0_8", "#s0^*_][fvwYbk&+,=@WEP]zkfdK{PKUW@rCGp+p0u;qLCqni0l=W[&=0B(0?-nx^*;10vm+L0^*,&10l}-QIj{", "%L)xa0J$
29	$ \begin{tabular}{lllllllllllllllllllllllllllllllllll$
30	$ \begin{tabular}{ll} $ \begin{tabular}{ll}$

Annex III: Source code

```
[1] main.rs
______
NJS: Database Protection Algorithm
______
The NJS algorithm is of symmetric and non-deterministic type. It can be used to protect information in a relational database.
The acronym NJS is the acronym of Noêmia Josefina da Silva, the great honoree of this work.
Designer: Edimar Veríssimo (Yugi).
______
* /
use std::time::Instant;
extern crate base64;
extern crate num bigint as bigint;
use num bigint::BigUint;
// Use another source code file [show_mwssage.rs]
mod show message;
use crate::show message::*;
// Use another source code file [read write.rs]
mod read write;
use crate::read write::*;
// Use another source code file [math.rs]
mod math;
// Use another source code file [intermediate file.rs]
mod intermediate file;
use crate::intermediate file::*;
// Use another source code file [cripto.rs]
mod cripto;
use crate::cripto::*;
// Use another source code file [cripto.rs]
mod compress;
use crate::math::*;
extern crate rand;
extern crate num;
```

// Configurable variables and constants

```
const NUM COL : usize = 5; // number of columns in table:
const LEN BLOCK : i32 = 60; // size block in bytes - max 512
const NOISE: u16 = 257; // range noise [1,257]:
static FILE PATH : &str = "/home/yuqi/criptografia/NJS"; // path of database
const TOT COL : usize = NUM COL+1; // don't change this value
const EMPTY STRING: String = String::new();
// do no change this value!!!!
static COMPRESSAO : &str = " !#$%&()*+,-/:;<=>?@ABCDEFGHIJKLMNOPQRSTUVWXYZ[]^ `abcdefghijklmnopqrstuvwxyz{}~\";
fn main() {
   message header("NJS: Database Protection Algorithm");
   let processing time total = Instant::now();
   let info data = read data();
   msg("\n\nShow Database Record - plaintext:\n ");
   show record(&info data,5);
   let encrypted = cypher data(&info data);
   msg("\n\nShow Database Record - Compress File:\n ");
   show record2 (&encrypted, 5);
   let plaintext = decypher data(&encrypted);
   msg("\n\nShow Database Record - plaintext:\n ");
   show record(&plaintext,5);
   message header ("END OF THE PROCESS!");
   ptime(processing time total);
fn read data() -> Vec<[String; NUM COL]> {
   msg("[01] Reading Database: ");
   let processing time = Instant::now();
   let info data = read database(DATABASE);
   ptime(processing time);
   return info data;
fn cypher_data(database : &Vec<[String;NUM_COL]>) -> Vec<[String;TOT_COL]> {
   msg("[01] Encrypted Database: ");
   let processing time = Instant::now();
   let cipher data = transform database(&database);
```

```
msg("\nWrite encrypted text: ");
    let result = write database encrypted("encrypted database.txt", &cipher data);
    ptime(processing time);
    return cipher data;
fn decypher_data(database : &Vec<[String;TOT_COL]>) -> Vec<[String;NUM_COL]> {
    msq("[02] Decrypted Database: ");
    let processing time = Instant::now();
    msq("[03] uncompress: ");
    // key tuples
    let ichave = controle chave().1;
    msq("[04] decrypted: ");
    let database encrypted = inv cipher database3(&database, &ichave);
    let dados = database encrypted;
    msg("[06] Write: ");
    let plaintext = verify database(&dados);
   msq("\nWrite plain text: ");
    let result = write database("texto claro.txt", &plaintext);
    ptime (processing time);
    return plaintext;
pub fn controle chave() -> (Vec<BigUint>, Vec<BigUint> , [u8;256], [u8;256]) {
    // Chave do sistema cifragem
    let chave = vec![
big integer("7985994286632031099561581639144489538693526623287703602155736297926703260802430595798855751721599411609968587914924269510
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3920733594443161811202603577008773", false),
```

```
999576955133001725533228366416498", false),
big integer("6799680328145012487684763913900327567383392172735171370497044557239255454696405758287258448595986726419363312719792896001
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```

```
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```

```
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```

```
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1;
```

```
// Sbox
   let sbox =
       64, 124, 108, 28, 188, 254, 134, 154, 187, 235, 199, 240, 251, 2, 70, 156,
       3, 133, 66, 107, 128, 150, 79, 77, 130, 83, 58, 121, 49, 246, 94, 101,
     176, 60, 57, 238, 202, 36, 126, 198, 158, 86, 194, 47, 174, 18, 144, 26,
     151, 120, 175, 205, 201, 42, 82, 95, 179, 164, 112, 186, 217, 104, 45, 223,
     200, 99, 195, 247, 46, 78, 102, 93, 34, 65, 97, 105, 92, 115, 139, 31,
      85, 145, 234, 110, 33, 123, 137, 173, 72, 40, 140, 252, 80, 185, 160, 215,
     227, 177, 218, 21, 10, 161,
                                  7, 71, 245, 109, 43, 116, 131, 142, 20, 113,
     232, 129, 16, 210, 253, 30, 56, 167, 216, 63, 208, 125, 189, 25,
     106, 163, 152, 182, 19, 69,
                                   9, 29, 27, 193, 178, 51, 228, 90, 122,
                             4, 8, 118, 157, 165, 96, 206, 135, 127, 147,
     203, 12, 41, 35, 225,
     219, 244, 192, 59, 204, 184, 153, 14, 75, 111, 15, 168, 214, 138, 114, 249,
     169, 207, 141, 32, 166, 143, 53, 149, 226, 22, 38, 146, 74, 76, 183, 224,
       6, 222, 233, 117, 48, 91, 236,
                                      5, 170, 162, 243, 55, 213, 197, 248, 61,
      37, 11, 24, 100, 242, 87, 239, 103, 136, 89, 62, 54, 211, 237, 181, 88,
     159, 23, 231, 196, 155, 220, 132, 13, 50, 67, 44, 221, 148, 17, 171, 39,
     180, 68, 81, 209, 98, 119, 191, 230, 250, 255, 190, 84, 212, 229, 241, 172
   ];
   let isbox = [
       159, 126, 13, 16, 149, 199, 192, 102, 150, 134, 100, 209, 145, 231, 167, 170,
114, 237, 45, 132, 110, 99, 185, 225, 210, 125, 47, 136, 3, 135, 117, 79,
179, 84, 72, 147, 37, 208, 186, 239, 89, 146, 53, 106, 234, 62, 68, 43,
196, 28, 232, 139, 127, 182, 219, 203, 118, 34, 26, 163, 33, 207, 218, 121,
 0, 73, 18, 233, 241, 133, 14, 103, 88, 143, 188, 168, 189, 23, 69,
92, 242, 54, 25, 251, 80, 41, 213, 223, 217, 141, 197, 76, 71,
154, 74, 244, 65, 211, 31, 70, 215, 61, 75, 128, 19,
                                                         2, 105,
58, 111, 174, 77, 107, 195, 151, 245, 49, 27, 142, 85,
                                                         1, 123, 38, 157,
                                                        90, 178, 109, 181,
20, 113, 24, 108, 230, 17,
                            6, 156, 216, 86, 173, 78,
46, 81, 187, 158, 236, 183, 21, 48, 130, 166, 7, 228, 15, 152, 40, 224,
94, 101, 201, 129, 57, 153, 180, 119, 171, 176, 200, 238, 255, 87, 44, 50,
32, 97, 138, 56, 240, 222, 131, 190, 165, 93, 59, 8,
                                                         4, 124, 250, 246,
162, 137, 42, 66, 227, 205, 39, 10, 64, 52, 36, 144, 164, 51, 155, 177,
122, 243, 115, 220, 252, 204, 172, 95, 120, 60, 98, 160, 229, 235, 193, 63,
191, 148, 184, 96, 140, 253, 247, 226, 112, 194, 82, 9, 198, 221, 35, 214,
11, 254, 212, 202, 161, 104, 29, 67, 206, 175, 248, 12, 91, 116,
   ];
   return (chave, ichave, sbox, isbox);
```

[2] compress.rs

```
DATABASE COMPRESSION ROUTINES
______
Function: transform vetor()
Compress the database - auxiliary function
Parameters:
   vetor: binary vector in byte format : u8
   mapa : binary vector in byte format : u8
Return:
   ret: String compressed : String
______
* /
pub fn transform vetor(vetor : Vec<u8>, mapa : Vec<u8>) -> String {
   let mut ret = String::new();
   let mut controle : [i8;2] = [-1, -1];
   let mut posicao controle : usize = 0;
   // Transforming the field contents into compressed format
   for ct in vetor{
       if ct == 46 {
          if posicao controle == 0 {
              ret.push str(".");
           } else {
              let valor2 : Vec<u8> = [(controle[0]+49) as u8].to vec();
              let s = String::from utf8 lossy(&valor2);
              ret.push str(&s);
              ret.push str(".");
              posicao controle = 0;
           continue;
       if ct == 48 {
           if posicao controle == 0{
              ret.push str("0");
          } else {
              let valor2 : Vec<u8> = [(controle[0]+49) as u8].to vec();
              let s = String::from utf8 lossy(&valor2);
              ret.push str(&s);
              ret.push str("0");
              posicao controle = 0;
```

```
continue;
       // Makes the control 0 to 8
       controle[posicao controle] = (ct as i8)-49;
       posicao controle = posicao controle + 1;
       if posicao controle > 1 {
           let valor : usize = ((controle[0]*9) + controle[1]) as usize;
           let valor2 : Vec<u8> = [mapa[valor]].to vec();
           let s = String::from utf8 lossy(&valor2);
           ret.push str(&s);
           posicao controle = 0;
   // Reading possible final residue
   if posicao controle == 1{
       let valor2 : Vec<u8> = [(controle[0]+49) as u8].to vec();
       let s = String::from utf8 lossy(&valor2);
       ret.push str(&s);
   return ret;
Function: transform vetor dec()
Uncompress the database - auxiliary function
   vetor: binary vector in byte format : u8
   mapa : binary vector in byte format : u8
   ret: String compressed : String
______
pub fn transform vetor dec(vetor : Vec<u8>, mapa : Vec<u8>) -> String {
   let mut ret = String::new();
   for ct in vetor{
       if ct == 46 {
           ret.push str(".");
           continue;
```

```
if ct == 48 {
        ret.push str("0");
        continue;
   // read the position element in the map
   let index element = mapa.iter().position(|&valor| valor == ct as u8).unwrap or(100);
    // Transcribing uncompressed characters
   if index element == 100{
        let \overline{\text{num}} = \text{ct} - 48;
        let r = num.to_string();
        ret.push_str(&r);
        continue;
    // separate 2 elements found
   let c1 : usize = (index_element / 9) + 1;
   let c2 : usize = (index_element % 9) + 1;
    // descompress elements found
   let mut s = c1.to string();
   s.push str(&c2.to string());
   ret.push str(&s);
}
return ret;
```

[3] show_message.rs

```
DISPLAYS MESSAGES ON THE SCREEN
use crate::NUM COL;
use crate::TOT COL;
Function: message_header()
Show a message on the screen (header)
    message 01: Message to show on the screen : &str
Return:
pub fn message header(message 01 : &str){
    let mut size separator = 80;
    let separator = &vec!["="; size separator];
    let mut s = String::from("");
    let mut t = String::from("");
    size separator = size separator - message 01.len();
    size separator = size separator / 2;
    for ct in 0..size separator {
        t.push str(" ");
    for ct in separator {
        s.push str(ct);
    println!("");
    println!("{}", s);
    println!("{}{}",t,message 01);
    println!("{}", s);
   println!("");
```

```
Function: msq()
Show a message on the screen
    message 01: Message to show on the screen : &str
Return:
pub fn msg(message 01 : &str) {
    print!("{}\n", message_01);
Function: ptime()
Show a message on the screen - Processing time % \left( \frac{1}{2}\right) =0
   message 01: Start of processing : std::time::Instant
Return:
pub fn ptime(pro time : std::time::Instant) {
    let elapsed time = pro time.elapsed();
    let final time = elapsed time.as secs();
    println!("The Processing Time is {} seconds.", final time);
Function: show record()
presents a sample of the records of a matrix
Parameters:
    regdatabase: vector with data : &Vec<[String;5]>
    n: number of registers to show: usize
Return:
______
pub fn show record(regdatabase : &Vec<[String;NUM COL]>, n : usize) {
    // Show the firsts register
   for ct in 0..n {
```

```
print!("[ {} ] ", ct+1);
        for ct2 in 0..NUM COL {
            if ct2 != (NUM COL-1)  {
                print!("{} | ",regdatabase[ct][ct2]);
            } else {
                print!("{}",regdatabase[ct][ct2]);
        println!("");
Function: show record()
presents a sample of the records of a matrix
    regdatabase: vector with data : &Vec<[String;5]>
    n: number of registers to show : usize
Return:
    No
pub fn show record2(regdatabase : &Vec<[String;TOT COL]>, n : usize) {
    // Show the firsts register
    for ct in 0..n {
        print!("[ {} ] ", ct+1);
        for ct2 in 0..TOT COL {
            if ct2 != (TOT COL-1) {
                print!("{} | ",regdatabase[ct][ct2]);
            } else {
                print!("{}",regdatabase[ct][ct2]);
        println!("");
```

[4] math.rs

```
MATHEMATICAL AND RELATED FUNCTIONS
use sha3::{Digest, Sha3 224};
extern crate num bigint as bigint;
use num bigint::BigUint;
Function: h224()
Extract hash with 224 bits of String
Return 56 characters hexadecimals
   msg: Message for extract hash: String
Return:
   ret: hash : String
______
pub fn h224(msq : String) -> String {
   let mut hasher = Sha3 224::new();
   hasher.update(msq);
   let result = hasher.finalize();
   let mut ret = String::from("");
   for ct in result {
       let lbyte = format!("{:x}", ct);
       if ct < 16 {
           ret.push str("0");
       ret.push str(&lbyte);
   return ret;
Function: big integer()
Lets you create a BigUint number
```

```
Parameters:
   value: str for convert to BigUint : &str
   check: enables or disables error checking : bool
Return:
   number 1: BigUint
______
* /
pub fn big integer(value: &str, check: bool) -> BigUint {
   return value.parse::<BigUint>().unwrap();
   let number 1 = &value.to string();
   // Security check - optional
   if check {
      let check value = number_1.as_bytes();
       for ct in check value {
          if ct < &48u8 || ct > &57u8 {
             println!("* * * * * There was a error in the conversion for BigUint number * * * * *");
             return BigUint::parse_bytes(b"0", 10).unwrap();
      }
   return number 1.parse::<BigUint>().unwrap();
   */
______
Function: pow big min()
Raises a bigUint to a small exponent
   num: large integer without sign to raise to an exponent: &BigUint
   exponent: exponent to raise the base : &BigUint
Return:
   result: BigUint
______
pub fn pow big min(num : &BigUint , exponent : &BigUint) -> BigUint {
   let mut result = BigUint::parse bytes(b"1", 10).unwrap();
   let control = BigUint::parse bytes(b"1", 10).unwrap();
   let limit = BigUint::parse bytes(b"0", 10).unwrap();
   let mut exp = exponent.clone();
```

```
// if exp = 0, return 1
if exp == limit {
    return control;
}

loop {
    result = result * num;
    exp = exp - &control;

    if exp == limit {
        break;
    }
}

return result;
```

[5] intermediate_file.rs

```
INTERMEDIATE FILE - FUNCTIONS
extern crate rand;
use rand::thread rng;
use rand::Rng;
extern crate num bigint as bigint;
use num bigint::BigUint;
extern crate bigdecimal;
// Use another source code file [math.rs]
use crate::math::*;
use crate::NUM COL; // Number of collumns of table (cleartext)
use crate::TOT COL; // Number of collumns of encrypted table - do not change this value!!!
use crate::NOISE;
use crate::LEN BLOCK;
use crate::EMPTY STRING;
use crate::cripto::*;
use crate::controle chave;
use crate::compress::*;
use crate::COMPRESSAO;
Function: transform database()
creates the intermediate file (ready to be encrypted)
    v: matrix with the data in plain text: &Vec<[String; NUM COL]>
Return:
    x: numerically encoded file : Vec<[String;TOT COL]</pre>
pub fn transform database(v : &Vec<[String;NUM COL]>) -> Vec<[String;TOT COL]> {
    // Prepare return this function:
    let mut x = Vec::new();
    // size of database
    let sizevet = v.len();
    // System key:
```

```
let chave = controle chave().0;
    let sbox = controle chave().2;
    let isbox = controle chave().3;
    // Compression vector
    let mapa = COMPRESSAO.as bytes();
    // how to populate a String vector or declare the vector with default element
    let mut vet: [String; TOT COL] = [EMPTY STRING; TOT COL];
    for ct in 0..sizevet {
        // convert each field of database
        for ct2 in 0..NUM COL{
            // Differentiating the SBOX for each record field
            let campo : usize = isbox[(ct2.clone() % 256) as usize] as usize;
            vet[ct2] = convert text(&v[ct][ct2].as bytes().to vec(),&LEN BLOCK,&NOISE,&sbox,&campo) ;
        // Add Hash Value (SHA3-224) - new field (see that vet contains TOT COL elements)
        let campo : usize = isbox[NUM COL % 256] as usize;
        vet[NUM COL] = convert text(&concatenate texthash(&v[ct]).as bytes().to vec(), &LEN BLOCK, &1, &sbox, &campo);
        // encrypting record
        let ret = cipher database3(&vet, &chave);
        // Applying data compression
        for ct2 in 0..TOT COL{
            let tmp = (ret[ct2]).as bytes().to vec();
            vet[ct2] = transform vetor(tmp, mapa.to vec());
        // include the record in the return vector
        x.push(vet.clone());
    return x;
Function: convert text()
Function convert text into number-text
Parameters:
    texto: text for convert number: &[u8]>
    sizeblock: size of the block that will be generated (in bytes) : i32
    noise: noise for encryption - value from 1 to 256 : u16;
```

```
sbox: sbox system key: u8
   campo: variable to distinguish each field of the table
Return:
   ret: converted number string : String
______
fn convert text(texto: &[u8], sizeblock: &i32, noise: &u16, sbox: &[u8], campo:&usize) -> String {
   let mut ret = String::from("");
   let mut contador : i32 = sizeblock-2; // 6
   let block len = sizeblock -2;
   // size block:
   let big exp = big integer(&(sizeblock-1).to string(), false);
   // base 256 informat bigUint
   let base = big integer("257", false);
   // generate a random number (noise):
   let mut rng = thread rng();
   let number rand: u16 = rng.gen range(0, noise).try into().unwrap();
   // variable for beginning process
   let mut resultado : BigUint = number rand * pow big min(&base , &big exp);
   // Controlling the size of the information
   let tamanho real : i32 = texto.len() as i32;
   let limite : i32 = tamanho real % (sizeblock-1);
   let mut tamanho : i32 = tamanho real - limite;
   // Very important this code for splitting short fields (sizeblock or less)
   if tamanho real <= block len {</pre>
       tamanho = tamanho real;
       contador = tamanho -1;
   let campo cifra = campo.clone() as u8;
   // Main processing
    for ct in 0..tamanho{
       // Adding the information in the variable with the noise
       // passing the SBOX - differential for each field using SBOX
       let tmp = sbox[(texto[ct as usize]) as usize] ^ campo cifra;
       resultado += ((tmp+1) as u16) * pow big min(&base , &big integer(&contador.to string(),false));
       contador = contador - 1;
       // Checking the end of each block
       if contador < 0 {
```

```
// Reset blocksize
            contador = block len;
            // Adding the value to the cumulative string
            let linha = &resultado.to string();
            ret.push str(linha);
            if ct != (tamanho-1) {
                ret.push str(".");
            // Resetting the noise vector
            let number rand: u16 = rng.gen range(0, noise).try into().unwrap();
            resultado = number rand * pow big min(&base , &big exp);
    // Code to check if there is an incomplete piece of block left to be processed:
    if tamanho != tamanho real {
        contador = tamanho real - tamanho -1;
        let number rand: u16 = rng.gen range(0, noise).try into().unwrap();
        resultado = number rand * pow big min(&base , &big exp);
        for ct in tamanho..tamanho real {
            // passing the SBOX - differential for each field using SBOX
            let tmp = sbox[(texto[ct as usize]) as usize] ^ campo cifra;
            resultado += ((tmp+1) as u16) * pow big min(&base , &big integer(&contador.to string(),false));
            contador = contador - 1;
            // create the final string
            let linha = &resultado.to string();
            ret.push str(".");
            ret.push str(linha);
    return ret;
Function: concatenate texthash()
Return hash SHA3-224 for concatenated fields
Parameters:
```

```
vet: vect with fields of table: [String; NUM COL]>
Return:
   ret: hash of concatened fields: String (Hexadecimal)
______
fn concatenate texthash(vet : &[String; NUM COL]) -> String {
   let mut hash text = String::from("");
   for ct2 in 0..NUM COL{
       hash text.push str(&vet[ct2]);
   return h224 (hash text);
/*
______
Function: convert text dec()
Function convert text into number-text
   texto: text for convert number: &String
   sizeblock: size of the block that will be generated (in bytes): i32
   others: auxiliary variables to speed up processing
Return:
   ret: converted number string : String
       ______
fn convert text dec(texto: &String, sizeblock: i32, zero: &BigUint, bigu: &[BigUint], bigu pot: &[BigUint], guarda: &BigUint,
isbox: &[u8], campo:&usize) -> String {
   let mut ret = String::from("");
   let mat = texto.split(".");
   let mut valor : BigUint = big integer("0", false);
   // main loop of this function
   for texto parte in mat {
       let mut num : BigUint = texto parte.parse::<BigUint>().unwrap();
       // Alternative code for removing noise
       let num2 = num.clone();
       num = num - (guarda * (num2 / guarda));
       let campo cifra = campo.clone();
       // converting to string
       let mut limite bloco2 = sizeblock.clone();
```

```
let mut contador : usize = 0;
            for controle in (0..limite bloco2).rev(){
                valor = &bigu_pot[controle as usize] * &bigu[256];
                if valor <= num{</pre>
                    contador = controle as usize + 1;
                    break;
            // reducing the threshold value to increase the speed
            limite bloco2 = (contador as i32).clone();
            // if the values are equal we end the loop
            if num == valor {
                let vet : Vec<u8> = [isbox[(255^campo cifra)as usize]].to vec();
                let s = String::from utf8 lossy(&vet);
                ret.push str(&s);
                break;
            // alternative decryption code
            let prop = &num / &bigu_pot[contador];
            let m = prop.to string().parse::<u16>().unwrap();
            num = num - (&bigu pot[contador] * &bigu[m as usize]);
            let tmp = ((m-1) as u8) ^ campo cifra as u8;
            let vet : Vec<u8> = [isbox[tmp as usize]].to vec();
            let s = String::from utf8 lossy(&vet);
            ret.push str(&s);
            if &num == zero{
                break;
    return ret;
/*
Function: verify hash()
data checking function
```

for ct3 in (0..sizeblock).rev(){

```
Parameters:
   vet: vect with fields of table: [String;TOT COL]>
Return:
   ret: verify hash : bool
______
pub fn verify hash(vet : &[String;TOT COL]) -> bool {
   let mut hash text = String::from("");
   for ct2 in 0..NUM COL{
       hash text.push str(&vet[ct2]);
   let hash1 = h224 (hash text);
   let hash2 = vet[NUM COL].clone();
   if hash1 == hash2 {
       return true;
   } else {
       return false;
______
Function: verify database()
Does the final check of the data decryption
Parameters:
   v: matrix with the data in plain text in decryption: &Vec<[String; NUM TOT]>
   x: matrix plain text : Vec<[String; NUM COL]
pub fn verify database(v : &Vec<[String;TOT COL]>) -> Vec<[String;NUM COL]> {
   // Prepare return this function:
   let mut x = Vec::new();
   // size of database
   let sizevet = v.len();
   // how to populate a String vector or declare the vector with default element
   let mut vet: [String; NUM COL] = [EMPTY STRING; NUM COL];
   let mut contador = 0;
```

```
for ct in 0..sizevet {
        // data verification
        let ret = verify hash(&v[ct]);
        if ret == false {
            println!("There was an error in decrypting the data!!!!");
            contador = contador + 1;
        // Copy data for return
        for ct2 in 0..NUM COL{
            vet[ct2] = v[ct][ct2].clone();
        // include the record in the return vector
        x.push(vet.clone());
    // checking if there were any errors in the decryption
    println!("\nNumber of errors: {}",contador);
    return x;
Function: transform database dec()
inverse function transform database()
Parameters:
    v: matrix with the data in plain text: &Vec<[String;TOT COL]>
    others: auxiliary variables to speed up processing
Return:
    x: numerically encoded file : Vec<[String; NUM COL]</pre>
* /
                                                                             transform database dec2(v
[String; TOT COL], zero: &BigUint, bigu: &Vec < BigUint >, bigu pot: &Vec < BigUint >, guarda: &BigUint, isbox: &[u8]) -> [String; TOT COL] {
    // how to populate a String vector or declare the vector with default element
    let mut vet: [String; TOT_COL] = [EMPTY_STRING; TOT_COL];
    // convert each field of database
    for ct in 0..TOT COL{
        let campo : usize = isbox[ct.clone() % 256] as usize;
        vet[ct] = convert text dec(&v[ct], LEN BLOCK, &zero, &bigu, &bigu pot, &guarda, &isbox, &campo);
    return vet;
```

[6] cripto.rs

```
FILE ENCRYPTION FUNCTION
extern crate num bigint as bigint;
use num bigint::BigUint;
use crate::TOT COL;
// Use another source code file [math.rs]
use crate::math::big integer;
use crate:: EMPTY STRING;
use crate::compress::*;
use crate::intermediate file::*;
use crate::LEN BLOCK;
use crate::math::*;
use crate::COMPRESSAO;
use crate::controle chave;
extern crate rand;
use rand::thread rng;
use rand::Rng;
/*
function cipher database()
Function for cipher the intermediate database
The key is passed as parameter
Parameters:
    database: database intermediate format : Vec<[String;TOT COL]>
    chave: vector with the keys and modules
    modulo: module for operations
    x : return database encrypted: Vec<[String;TOT COL]>
______
pub fn cipher database3(database : &[String;TOT COL], chave : &Vec<BigUint>) -> [String;TOT COL] {
    // how to populate a String vector or declare the vector with default element
    let mut cifra: [String; TOT_COL] = [EMPTY_STRING; TOT_COL];
       // random data insertion point
       let mut rng = thread rng();
       let number_rand: u8 = rng.gen_range(0, 8).try_into().unwrap();
```

```
for campo in 0..TOT COL {
           // create the string for reading data
           let mut resultado = "".to string();
           // Table field to be encrypted
           let registro = database[campo].split(".");
           // temporary vector
           let mut y: Vec<BigUint> = Vec::new();
           // transform String Vector in type BigUint
           for ct3 in registro{
               y.push(ct3.parse::<BigUint>().unwrap());
           // Encryption
           let tam2 = y.len();
           for ct4 in 0..tam2{
               // Non-deterministic encryption
               y[ct4] = cifrar(&y[ct4], chave, &number rand);
               resultado.push str(&y[ct4].to string());
               if ct4 != (tam2-1)  {
                   resultado.push str(".");
           // storing field information in the accumulator
           cifra[campo] = resultado;
   return cifra; // Function return
/*
Function to encrypt the data
Receives the BigUint value and returns the encrypted BigUint value with the NJS pattern
______
fn cifrar(dados : &BigUint, chave : &Vec<BigUint>, controle: &u8) -> BigUint{
   // Choose a random operation type to encrypt
   let mut ret : BigUint = dados.clone();
   match controle {
       0 =>
       // ADD MUL XOR
```

```
let mut fator = 0;
    loop {
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 44 {
            break;
},
1 =>
// MUL ADD XOR
    let mut fator = 0;
    loop {
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 44 {
            break;
},
2 =>
// ADD MUL XOR
    let mut fator = 0;
    loop {
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 32 {
            break;
},
3 =>
// MUL ADD XOR
```

```
let mut fator = 0;
    loop {
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 32 {
           break;
},
4 =>
// ADD MUL XOR
    let mut fator = 0;
   loop {
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 20 {
           break;
},
5 =>
// MUL ADD XOR
   let mut fator = 0;
    loop {
        ret = (ret * (&chave[1+fator])) % &chave[3+fator];
        ret = (ret + &chave[0+fator]) % &chave[3+fator];
        ret = ret ^ (&chave[2+fator]);
        fator = fator + 4;
        if fator > 20 {
           break;
},
6 =>
// ADD MUL XOR
```

```
let mut fator = 0;
           loop {
               ret = (ret + &chave[0+fator]) % &chave[3+fator];
               ret = (ret * (&chave[1+fator])) % &chave[3+fator];
               ret = ret ^ (&chave[2+fator]);
               fator = fator + 4;
               if fator > 8 {
                  break;
       },
       7..=255 =>
       // MUL ADD XOR
           let mut fator = 0;
           loop {
               ret = (ret * (&chave[1+fator])) % &chave[3+fator];
               ret = (ret + &chave[0+fator]) % &chave[3+fator];
               ret = ret ^ (&chave[2+fator]);
               fator = fator + 4;
               if fator > 8 {
                  break;
   };
   return ret;
Function to decrypt the data
Takes the BigUint value and returns the decrypted BigUint value with the NJS pattern
______
fn decifrar(dados : &BigUint, ichave : &Vec<BigUint>, tipo: &u8) -> BigUint{
   let controle : u8 = tipo.clone();
   let mut ret : BigUint = dados.clone();
   match controle {
```

```
0 =>
   // XOR MUL ADD
       let mut fator : i32 = 44;
       loop {
       ret = ret ^ &ichave[(2+fator) as usize];
       ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
        ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
       fator = fator - 4;
           if fator < 0 {
               break;
    },
   1 =>
    // XOR ADD MUL
       let mut fator : i32 = 44;
       loop {
       ret = ret ^ &ichave[(2+fator) as usize];
        ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
       ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
       fator = fator - 4;
           if fator < 0 {
               break;
   },
2 =>
// XOR MUL ADD
   let mut fator : i32 = 32;
   loop {
   ret = ret ^ &ichave[(2+fator) as usize];
   ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
   ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
   fator = fator - 4;
       if fator < 0 {
           break;
},
3 =>
// XOR ADD MUL
```

```
let mut fator : i32 = 32;
   1000 {
   ret = ret ^ &ichave[(2+fator) as usize];
   ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
   ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
   fator = fator - 4;
       if fator < 0 {
           break;
},
4 =>
// XOR MUL ADD
   let mut fator : i32 = 20;
   loop {
   ret = ret ^ &ichave[(2+fator) as usize];
   ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
   ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
   fator = fator - 4;
       if fator < 0 {
           break;
},
5 =>
// XOR ADD MUL
   let mut fator : i32 = 20;
   loop {
   ret = ret ^ &ichave[(2+fator) as usize];
   ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
   ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
   fator = fator - 4;
       if fator < 0 {
           break;
6 =>
// XOR MUL ADD
   let mut fator : i32 = 8;
   loop {
   ret = ret ^ &ichave[(2+fator) as usize];
```

```
ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
       ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
       fator = fator - 4;
          if fator < 0 {
              break;
   },
   7...=255 =>
   // XOR ADD MUL
       let mut fator : i32 = 8;
       loop {
       ret = ret ^ &ichave[(2+fator) as usize];
       ret = (ret + &ichave[(0+fator) as usize]) % &ichave[(3+fator) as usize];
       ret = (ret * (&ichave[(1+fator) as usize])) % &ichave[(3+fator) as usize];
       fator = fator - 4;
          if fator < 0 {
              break;
} ;
   return ret;
______
function cipher database()
Function for cipher the intermediate database
The key is passed as parameter
Parameters:
   database2: database intermediate format : Vec<[String;TOT COL]>
   ichave: vector with the keys and modules
Return:
   x : return database encrypted: Vec<[String; TOT COL]>
_____
pub fn inv cipher database3(database2: &Vec<[String;TOT COL]>, ichave: &Vec<BigUint>) -> Vec<[String;TOT COL]> {
   // Sbox for decryption
   let isbox = controle chave().3;
```

```
// attempt to speed up the code
let mut bigu : Vec<BigUint> = vec![big integer("0",false);512];
for ct in 0..512{
   bigu[ct] = big integer(&(ct+0).to string(), false);
// size block - exponent:
let big exp = big integer(&(&LEN BLOCK-1).to string(), false);
// base 256 in format bigUint
let base = big integer("257", false);
let zero = big integer("0", false);
let guarda = pow big min(&base, &big exp);
// attempt to speed up the code
let mut biqu pot : Vec<BiqUint> = vec![biq integer("0",false);512];
for ct in 0..512{
   bigu pot[ct] = pow big min(&base, &big integer(&(ct+0).to string(), false));
let mut ret = Vec::new();
// how to populate a String vector or declare the vector with default element
let mut cifra: [String; TOT COL] = [EMPTY STRING; TOT COL];
// Map to decompress
let mapa = COMPRESSAO.as bytes();
// Here we start the actual decryption
for database in database2 {
    let mut database3 = database.clone();
    // Decompressing the data before decrypting
    for ct2 in 0..TOT COL{
        let tmp = (database[ct2]).as bytes().to vec();
        database3[ct2] = transform vetor dec(tmp, mapa.to vec());
    for tipo in 0..8 { // decryption variations
        for campo in 0..TOT COL {
            // create the string for reading data
            let mut resultado = "".to string();
            // Table field to be encrypted
            let registro = database3[campo].split(".");
            // temporary vector
            let mut y: Vec<BigUint> = Vec::new();
            // transform String Vector in type BigUint
            for ct3 in registro{
```

```
y.push(ct3.parse::<BigUint>().unwrap());
            // decryption
            let tam2 = y.len();
            for ct4 in 0..tam2{
                y[ct4] = decifrar(&y[ct4], ichave, &tipo);
                resultado.push str(&y[ct4].to string());
                if ct4 != (tam\overline{2}-1) {
                    resultado.push str(".");
            \//\ storing field information in the accumulator
            cifra[campo] = resultado;
        // Checking the decryption control:
        let tmp = &transform database dec2(cifra.clone(),&zero,&bigu,&bigu pot,&guarda,&isbox);
        if verify hash(tmp) == true {
            break;
    // Determining the correct decryption
   ret.push(transform database dec2(cifra.clone(), &zero, &bigu, &bigu pot, &guarda, &isbox));
return ret;
```

[7] read_write.rs

```
use std::fs::File;
use std::io::{BufReader, BufRead};
use std::io::Write;
use crate::FILE PATH;
use crate::NUM COL;
use crate::TOT COL;
use crate:: EMPTY STRING;
/*
______
Function: read database()
Read a text file and put the memory
Parameters:
   filename: filename: &str
Return:
   array with fields read from text files : Vec<[String;5]>
______
pub fn read_database(lfilename : &str) -> Vec<[String;NUM COL]> {
   // Preparing to read the file path
   let fpath = FILE PATH.to string();
   let mut filename = String::from("");
   filename.push str(&fpath);
   filename.push str("/");
   filename.push str(&lfilename);
   // Create the vector for processing
   let mut read vector = Vec::new();
   // Create a return vector:
   let mut ret database = Vec::new();
   // Open the file in read-only mode (ignoring errors).
   let file = File::open(filename).expect("Error reading file. Set the FILE PATH variable");
   let reader = BufReader::new(file);
```

```
// Read lines from file
    for line in reader.lines() {
        let line = line.unwrap(); // Ignore errors.
        read vector.push(line.to string());
    // how to populate a String vector or declare the vector with default element
    let mut y: [String; NUM COL] = [EMPTY STRING; NUM COL];
    // Main processing
    for line in read vector {
        // transform a string in a vector
        let word 01 = line.split("|");
        let mut campo = 0; // variable to control the divisions in each field
        for ct in word 01 {
            y[campo] = ct.trim().to string();
            campo = campo + 1;
        // put new element in matrix
        ret database.push(y.clone());
    // Return matrix with data
    return ret database;
Function: write database()
write a text file in disk
Parameters:
    filename: filename: &str
    database : table with fields : Vec<[String; NUM COL]>
Return:
    std::io::Result<()>
pub fn write database(lfilename: &str, database : & Vec<[String;NUM COL]>) -> std::io::Result<()> {
    // Preparing to write the file path
    let fpath = FILE PATH.to string();
    let mut filename = String::from("");
    filename.push str(&fpath);
    filename.push str("/");
    filename.push str(&lfilename);
```

```
let mut file = File::create(filename)?;
   let tam = database.len();
   for ct in 0..tam {
       let mut registro = String::from("");
       for ct2 in 0..NUM COL{
           registro.push str(&database[ct][ct2]);
           if ct2 != (NUM COL-1)  {
               registro.push str(" | ");
       }
       registro.push str("\n");
       file.write_all(registro.as_bytes())?;
   Ok(())
______
Function: write database encrypted()
write a text file in disk
Parameters:
   filename: filename: &str
   database : table with fields : Vec<[String; NUM COL]>
   std::io::Result<()>
pub fn write database encrypted(lfilename: &str, database : & Vec<[String;TOT COL]>) -> std::io::Result<()> {
   // Preparing to write the file path
   let fpath = FILE PATH.to string();
   let mut filename = String::from("");
   filename.push str(&fpath);
   filename.push str("/");
   filename.push str(&lfilename);
   let mut file = File::create(filename)?;
   let tam = database.len();
   for ct in 0..tam {
       let mut registro = String::from("");
       for ct2 in 0..TOT COL{
           registro.push str(&database[ct][ct2]);
```

```
if ct2 != (TOT_COL-1) {
            registro.push_str("|");
        }
      registro.push_str("\n");
      file.write_all(registro.as_bytes())?;
}
Ok(())
```

[8] cargo.toml

```
[package]
name = "projeto"
version = "0.1.0"
edition = "2021"

# See more keys and their definitions at https://doc.rust-lang.org/cargo/reference/manifest.html
[dependencies]
rand = "0.7.3"
base64 = "0.13.0"
num-bigint = "0.2"
sha3 = "0.10.1"
bigdecimal = "0.1"
num = "0.1"
```

[9] database_min.txt

```
Helda Fabrizio Corto | 576837751-29 | (27)94640-5409 | 23/03/2013 | 2886.22
Charlles Dockhorn Cossio Campara | 712237047-62 | (35)98394-9535 | 26/03/2022 | 26772.52
Guimar Acauan Azmus Brinke Caurrinhos Gabani | 222770563-79 | (17)92477-7699 | 25/06/2020 | 242.59
Iliana Fiorese Azzalini | 438695869-02 | (48)98827-9027 | 13/06/2010 | 2115.71
Diulia Albeck Grassetto Tessaro Frota Goodwin | 067880700-09 | (68)96658-8199 | 17/12/2012 | 3873.97
Nei Gottschalk | 053327513-38 | (65)92373-4258 | 07/04/2007 | 540.83
Denilda Ancinelo Efrem | 957259222-55 | (97)92088-8319 | 14/11/2015 | 19850.56
Weldson Biancardi Andersson | 252363691-66 | (93)90986-1043 | 04/06/2018 | 1982.85
Andreisa Bettger Beghi | 933807975-77 | (42)98156-2205 | 13/09/2014 | 15600.99
Manuela Brostolon Annunziato | 044044698-51 | (68)97933-6839 | 08/09/2015 | 809.54
Waderson Corpas | 559673060-30 | (69)94085-9162 | 29/05/2016 | 2102.14
Genisson Fisher Benassatto | 385677859-21 | (75)98790-8784 | 29/02/2020 | 833.54
Kenya de Albernas | 124195273-65 | (79)90834-0095 | 14/02/2009 | 8454.10
Daira Grossu Alexandre Chioca Giuduce | 810065369-15 | (69)92933-3889 | 26/05/2018 | 6452.43
Joventino Goffad Eichendorf | 673812302-46 | (71)98164-5349 | 28/07/2012 | 5515.55
Eliosmar Bartici | 010206101-57 | (14)92937-7096 | 23/05/2014 | 14363.99
Cleyde Bonsegno | 658708865-32 | (17)90338-9135 | 12/11/2009 | 19779.87
Mariliane Agnetti Hooper Mamiro | 903954380-09 | (61)97100-6801 | 30/01/2009 | 9827.69
Dalessandro Fischborn | 084775090-57 | (2)90918-5500 | 09/03/2005 | 11463.12
Milana Gallagher Graeber Brunel | 307368986-02 | (18)94207-4931 | 03/07/2013 | 7474.54
Thaiza Bizzo | 433528700-09 | (6)93230-0841 | 14/02/2007 | 5281.97
Edso Adornes | 522440324-82 | (18) 98719-2635 | 02/11/2012 | 11772.34
Dinara Bordina | 666403069-08 | (19)98194-5912 | 14/09/2010 | 9015.75
Veleda Deleon | 489682721-70 | (4)97794-6577 | 25/02/2017 | 11497.68
Alessandra Fruhan | 026879622-84 | (64)91386-5413 | 07/03/2021 | 29204.55
Kevilin Cicarini | 675339496-82 | (53)94512-1129 | 10/05/2010 | 4400.94
Teuma Gagetti | 577271378-60 | (24)94922-2718 | 16/11/2017 | 2591.46
Tertuliano Noqueira | 542293619-52 | (47)96606-6691 | 12/05/2015 | 1128.66
Edleni Borella | 573910000-00 | (6)96148-7154 | 06/04/2006 | 3172.58
Marcondes Goncalves Invernizzi | 814251795-11 | (93)96488-6412 | 03/09/2019 | 9892.98
Oracio Dellonardo Benfatti Burkhard Rassato Groehs | 619610284-78 | (38)96902-3369 | 19/01/2017 | 7130.25
Arival Romero Guglielmi | 277339799-00 | (51)91348-1062 | 15/04/2022 | 1922.99
Kele Amando Calimam Ferdinandi Lamha | 106388769-40 | (53)92471-4046 | 08/09/2004 | 9829.03
Dineide Canes Caldana | 530842160-46 | (13)98514-9811 | 09/10/2010 | 160.82
Milena Rush | 276606579-01 | (28)91459-2818 | 24/03/2003 | 13064.85
Vicentina Filippini | 249851544-80 | (9)94903-2774 | 28/02/2021 | 1587.54
Guibson Haarstick Medrado Galucci | 325921292-03 | (87)90098-5893 | 29/06/2015 | 14882.66
Mercedes Eickhof Torezani | 796662050-51 | (86)99325-1241 | 13/01/2018 | 7370.43
Olimpo Brackmann Bonese | 576690026-00 | (19)93264-6139 | 06/07/2003 | 4478.52
Joseph Hansen Hooper | 332949425-92 | (98)93336-4199 | 13/03/2019 | 26269.86
Dielen Basseto Baumy | 137307536-10 | (38)90641-6453 | 06/02/2019 | 8465.11
Ubelina Henrichsen | 094461172-01 | (5)98934-1424 | 14/04/2003 | 29672.59
Raylson Cotrim | 507793620-53 | (45)99607-1257 | 15/07/2016 | 3082.02
Valguer Durival | 754068785-00 | (14)92171-5559 | 06/04/2015 | 10903.54
Wellisson Faccipieri Brigati | 661419240-18 | (3)95307-1899 | 13/03/2017 | 1253.10
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Junia Covalesci Bacello | 821422627-06 | (6)99994-8329 | 13/10/2012 | 3989.39
Anthony Ferezini | 193213661-49 | (3)99370-3567 | 28/12/2004 | 4035.06
Maxel Dothes Davrison Collevati Concari | 197281777-80 | (7)92845-8966 | 22/07/2011 | 8420.54
Evelaine Diuliana Appol Sperotto Altreiter Cunshnir | 441433682-93 | (98)90803-1733 | 27/11/2013 | 27632.18
Gilda Dondi Altenetter | 926027157-27 | (93)95397-8657 | 07/05/2010 | 15480.82
Claudelina Bindchen | 393694677-38 | (79)95852-0543 | 26/02/2006 | 9936.42
Hermelindo Carnas Burkle | 389464377-00 | (14)90595-5222 | 13/06/2017 | 5896.58
Jerfesom Boettner Carturan | 639160956-10 | (42)93681-0888 | 22/06/2021 | 3514.91
Priciane Finkler | 738359137-80 | (54)91060-6514 | 23/07/2014 | 1888.22
Augusto Gasparotto Guzzi | 644584713-01 | (32)94002-1741 | 17/02/2012 | 6157.24
Maria Eduarda Mayer | 253775982-40 | (75)99151-8798 | 24/03/2008 | 4070.19
Gilcemar Basenova Pais Adamo Cappellaro Price | 220589100-08 | (82)94961-1207 | 09/06/2013 | 6167.68
Domiciano Avona Campisano | 267052015-62 | (65)92130-5471 | 15/10/2013 | 4900.23
Zenobio Favari Gonzalez | 150921090-57 | (12)90768-0145 | 19/10/2007 | 2173.97
Angelice Zerbone Goncales Guitarra Pozinni | 820656852-01 | (81)91228-6841 | 10/06/2017 | 25591.18
Semirames Buriani Bazilevitz | 515970813-97 | (38)99523-3468 | 21/11/2018 | 7002.86
Francelino Manfio | 306758165-94 | (6)92511-9672 | 12/08/2003 | 752.88
Elizelda Negri | 530334571-99 | (51)96020-7853 | 02/02/2003 | 4729.22
Leoberto Ferreira Collina | 002488228-09 | (24)93590-8240 | 01/11/2020 | 14179.04
Analice Lunaidi Dompieri Adamska Barbabela de Mattos Eichholz | 270491236-29 | (79)96628-2192 | 14/10/2014 | 6056.81
Diniz Gularte | 851958482-21 | (15)90821-4967 | 25/12/2010 | 28197.80
Glauce Cuadro Gay | 069647173-19 | (49)90547-1777 | 26/08/2020 | 9827.14
Josenita Praba | 498921301-35 | (4)94805-2797 | 01/09/2011 | 7103.62
Maronita Frias | 226708101-32 | (53)91155-2100 | 08/07/2014 | 26293.21
Valmiro Mccormick Borri Cagnin Saccomano | 770116621-93 | (16)91286-6260 | 18/05/2021 | 7361.57
Divanei Jungllut Geanine | 148619691-98 | (65)97231-5457 | 13/08/2015 | 3556.45
Helenir Dechsler Good Deutschle | 262414126-07 | (67) 98576-6817 | 30/10/2017 | 2954.49
Gorge Lupino | 838556605-52 | (94)91727-5409 | 08/06/2008 | 2095.84
Jaelson Haynes Kluge | 868852555-02 | (42)96432-3574 | 02/02/2020 | 105.45
Raimuda Bailey Cauzzi | 350563621-41 | (42)92151-9150 | 04/12/2008 | 25666.28
Oliverio Flocke | 069221046-60 | (12)99694-3983 | 20/09/2007 | 4382.77
Melke Gherele Braschi Fiorenza Entringer Dixen | 097443834-68 | (62)95688-7918 | 25/03/2011 | 9457.95
Lisangela Bernar | 292645790-55 | (42)94942-2105 | 27/05/2019 | 9016.15
Naile Aleci Barabani | 002037831-72 | (87)90956-2350 | 12/07/2014 | 1884.80
Cleidir Valani | 815054587-47 | (98) 98285-7133 | 14/06/2003 | 18234.37
Claudimara Carpegiani | 412974882-39 | (31)92616-0470 | 01/07/2014 | 3513.99
Lucca Lorenzetto | 125347476-90 | (28)92250-0784 | 09/04/2010 | 8851.75
Kesia Cesarini | 218791950-05 | (45)92171-3962 | 12/05/2015 | 6603.49
Albino Librenti Battilana Hafner Gloeden Girolami | 144451013-12 | (73)91332-5386 | 27/05/2017 | 19945.15
Iracena Gilande | 079017033-94 | (82)99576-8375 | 19/09/2011 | 17562.33
Marcon Camos Cohen | 479097814-49 | (37)92889-1684 | 15/08/2006 | 1758.15
Maria Sophia Ewing Dickhut | 071088715-37 | (12)94968-0890 | 13/12/2010 | 1566.77
Romenique Strong | 330054630-59 | (8)91860-4120 | 25/08/2016 | 28308.64
Zefira Aloisio | 668150351-50 | (68)98958-9075 | 10/02/2020 | 5651.68
Christina Fresh | 247400546-03 | (84)91423-2110 | 11/08/2003 | 6949.80
Walice Astolfo Braccer | 584819180-93 | (12)98373-1613 | 02/04/2006 | 26481.65
Sumiko Boarato | 397672799-41 | (42)98357-0431 | 03/05/2008 | 521.79
Kayane Frimm | 345470664-00 | (96)98225-2581 | 04/07/2019 | 4840.16
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Tanara Fabriz | 133997017-35 | (41)96469-8159 | 27/09/2014 | 11436.58

Clailson Cossari Fanto Bombel Brugalli | 452399361-52 | (83)90077-8258 | 05/03/2008 | 383.57

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