

A k-flip local search algorithm for SAT and MAX SAT

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Abstract

Local search can be applied to SAT by determining whether it is possible to increase the number of satisfied clauses for a given truth assignment by flipping at most k variables. However, for a problem instance with v variables, the search space is of order v^k . A naive approach that enumerates every combination is impractical for all but the smallest of problems. This paper outlines a hybrid approach that plays to the strength of modern SAT solvers to search this space more efficiently. We describe an encoding of SAT to a related problem – k-Flip MAX SAT – and show how through repeated application, it can be used to solve SAT and MAX SAT problems. Finally, we test the algorithm on a benchmark set with different values of k to see how it performs.

1 Introduction

- sat problems have hundreds or thousands of variables, doesn't scale
 - explain k-flip max sat
 - explain ipasir and justify it for this problem

2 The encoding

At a high-level, the encoding works by introducing a set of variables A that represents a hypothetical SAT solver's current truth assignment of variables within some formula F . A corresponding set of variables A' is introduced that is allowed to differ by at most k truth assignments from A . We use a counter circuit and a less-than comparator to enforce this constraint.

For each clause in F , we introduce a variable whose intended meaning is that the related clause has not been satisfied by A' . Collectively, we call this set U . We enforce that the number of true literals in U is less than the SAT solver's current number of unsatisfied clauses for F . We once again use a counter circuit and less-than comparator to enforce this constraint.

2.1 Flipped variables

2.2 Unsatisfied clauses

2.3 Parallel counter

2.4 Less-than comparator

3 Repeated application

4 Empirical results

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
430	0	0	0	1	4	4	8	10	17	17	17	22	21	28	24	32	42	31	41	37
429	0	0	1	10	16	15	20	22	17	25	26	21	27	26	30	22	14	21	18	22
428	0	3	11	13	14	14	16	17	14	12	13	11	10	6	6	6	4	7	1	1
427	0	7	11	18	14	23	8	9	10	5	4	6	2	0	0	0	0	1	0	0
426	1	6	15	9	9	3	7	2	1	1	0	0	0	0	0	0	0	0	0	0
425	3	9	10	7	2	1	0	0	1	0	0	0	0	0	0	0	0	0	0	0
424	4	13	5	2	1	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0
423	5	10	5	3	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
422	7	7	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
421	3	4	4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
420	9	4	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
419	6	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
418	8	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
417	11	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
416	4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
415	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
414	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
413	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Table 1: caption