

Northeastern University

**Department of Electrical and Computer
Engineering**

EECE2323: Digital Systems Design Lab

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**Lab # 1-3: 8-Bit Adder, Partial Arithmetic Logic
Unit, and Arithmetic Logic Unit**

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1 Introduction

The objective of the series of Lab 1, 2, and 3 were to become familiarized with Verilog hardware descriptions (Xilinx) and use Verilog code to implement small features on Zedboard hardware. The goals of the following experiments were to implement a simple 8-bit adder, a partial arithmetic logic unit (ALU), and a full ALU by extending partial ALU. The Verilog code was first written and tested in the Xilinx platform, and then a bitstream was generated for the hardware testing on the Zedboard, and functionality was verified on the board.

2 Design Approach

2.1 Logic Design in Verilog Code

The purpose of an 8-bit adder was to have two inputs of signed 8-bit and produce a signed 8-bit output and an overflow. Overflow was calculated using:

$$ovf = (f[7]! = a[7]) \quad \&\& \quad (f[7]! = b[7]) \quad (1)$$

Appendix A shows details of the code used.

The partial ALU (PALU) had a 2-bit selector and two 8-bit inputs, and the 2-bit selector determined how the output is determined. If the selector is 00, the result would be the sum of two 8-bit inputs. If it is 01, the output will be a complement of input b. If it is 10, AND operator logic will apply to the inputs. Finally if it is 11, then OR operator logic will apply. Appendix B shows details of the code used.

Finally the full ALU had a 3-bit selector. Full ALU is an extension of partial ALU, as the selector now has 8 different states. In addition, the full ALU has one more output called *take_branch*. The selector from 000 to 011 is same as PALU. If the selector is 100, the output is an arithmetic shift right of an input a. If it is 101, the output is a logic shift left of an input a. If it is 110, the output f is 0, and the *take_branch* will be the result of $a == b$. If the selector is 111, the output f will also be 0, but the *take_branch* will be the result of $a \neq b$. The full code is shown in Appendix C.

2.2 Software Used

- Xilinx Vivado - Xilinx Vivado was used to perform all model simulations, including testbenches. Additionally, Xilinx Vivado provided the means with which our models were uploaded onto the ZedBoard FPGAs.
- Icarus Verilog - Icarus Verilog was used for rapid development and verification of several of the modules used in our experiments.

3 Results and Analysis

3.1 Design Simulation

All simulation tests were done using Xilinx Vivado's simulation environment. All Verilog modules that were tested were shown to satisfy their respective requirements.

3.1.1 8-Bit Adder

Shown here is the software simulation of the 8-Bit Adder module testbench. A 200ns delay is observed, followed by two trivial test cases where inputs a and b share no common set bits. Subsequent tests show that the adder module successfully handles cases where there are common set bits in both a and b . The final case shows the overflow condition being obeyed.



Figure 1: Simulation results for lab 1, **8-Bit Adder**.

3.1.2 Partial Arithmetic and Logic Unit

Shown here is the software simulation of the 8-Bit Partial ALU module testbench. The simulation begins with a series of test cases which demonstrate that the ADD ($s = 0$) test case behaves correctly, followed by tests for other operations. Subsequent tests show that the partial ALU module successfully handles other operations. The final case shows the overflow condition being obeyed for the ADD ($s = 0$) instruction.

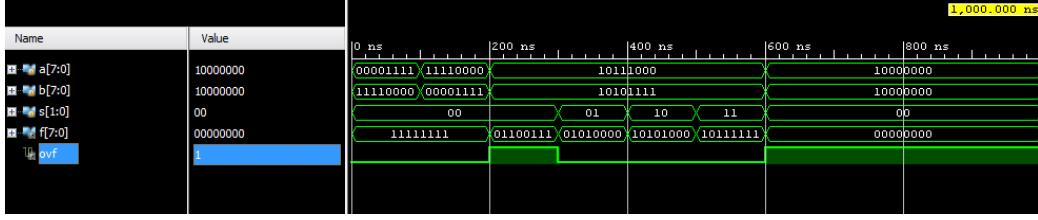


Figure 2: Simulation results for lab 2, **Partial Arithmetic and Logic Unit**.

3.1.3 Arithmetic and Logic Unit

Shown here is the software simulation of the 8-Bit ALU module testbench. The simulation begins with a series of test cases which demonstrate that the ADD ($s = 0$) test case behaves correctly, including a test for the ovf bit. Subsequent tests show that the ALU module successfully handles operations 1 through 5. The final four cases show the behavior of the BEQ and BNE instructions. In order, the four cases show that:

- `take_branch` is not set when the instruction is BEQ and $a \neq b$.
- `take_branch` is set when the instruction is BNE and $a \neq b$.
- `take_branch` is set when the instruction is BEQ and $a = b$.
- `take_branch` is not set when the instruction is BNE and $a = b$.

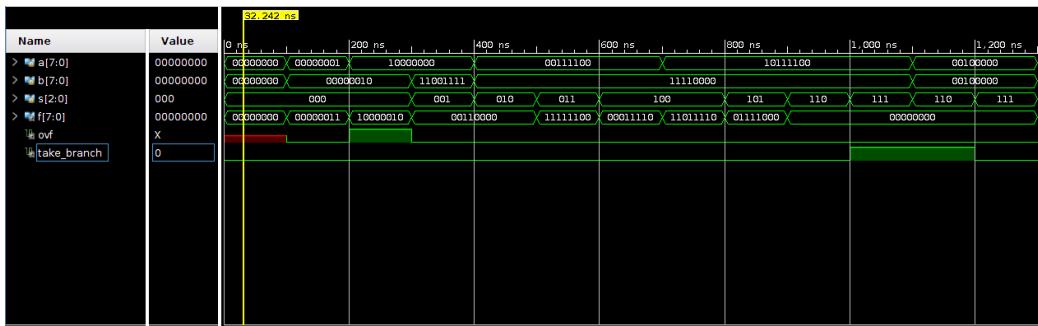


Figure 3: Simulation results for lab 3, **Arithmetic Logic Unit**.

3.2 Hardware Testing

Here we discuss our methods for testing our models in hardware. All hardware testing was done on a ZedBoard with an integrated Field-Programmable Gate Array (FPGA)¹.

3.2.1 8-Bit Adder

The eight-bit adder was tested using the ZedBoard's on-board switches to partially manipulate the inputs. Switches 0 through 3 control bits 0, 5, 6, and 7 of the first operand while switches 4 through 7 control bits 0, 5, 6, and 7 of the second operand. The hardware testing results can be seen in the [8-Bit Adder section](#) of the appendix.

3.2.2 Partial Arithmetic and Logic Unit

The eight-bit partial ALU was tested using the ZedBoard's on-board switches to partially manipulate the inputs. Switches 0 and 1 control bits bits 0 and 1 of the instruction (*s*) input. Switches 2 through 4 control bits 0, 6, and 7 of the first operand while 5 through 7 control bits 0, 6, and 7 of the second operand. The hardware testing results can be seen in the [8-Bit Partial ALU section](#) of the appendix.

3.2.3 Arithmetic and Logic Unit

The eight-bit partial ALU was tested using the Virtual Input/Ouput (VIO) functionality provided by Xilinx Vivado. This proved to be a more efficient way of testing the hardware behavior of our modules. A comprehensive suite of test cases was performed to ensure that all input and output bits were either set or unset during at least one test case. The hardware testing results can be seen in the [8-Bit ALU section](#) of the appendix.

¹The particular FPGA used is the Xilinx XC7Z020CLG484-1.

4 Conclusions

The first three Lab sessions introduced the basics of Verilog and the functionalities of the hardware Zedboard. Implementation of logics with Verilog and verification of the logics through the FPGA on the ZedBoard were explored. The results came out successful. In future lab, the knowledge of thorough testing and effectively utilizing the software obtained from this lab would be useful.

Appendices

A 8-Bit Adder Module: `eightbit_adder.v`

```
module eightbit_palu(
    input signed [7:0] a,
    input signed [7:0] b,
    output reg [7:0] f,
    output reg ovf
);
    always @(a, b) begin
        f = a + b;
        // overflow check for addition
        ovf = (f[7] != a[7]) && (f[7] != b[7]);
    end
endmodule
```

Figure 4: `eightbit_adder.v` implementing an 8-bit adder in Verilog.

B Partual ALU Verilog Module: `eightbit_palu.v`

```
module eightbit_palu(
    input signed [7:0] a,
    input signed [7:0] b,
    input [1:0] s,
    output [7:0] f,
    output ovf
);
    always @(a, b, s) begin
        // operations on 'f'
        case(s)
            0: f = a + b;
            1: f = ~b;
            2: f = a & b;
            3: f = a | b;
        endcase
        // overflow check for addition
        ovf = (s == 0) && (f[7] != a[7]) && (f[7] != b[7]);
    end
endmodule
```

Figure 5: `eightbit_palu.v` implementing an 8-bit partial ALU in Verilog.

C ALU Verilog Module: `eightbit_alu.v`

```
module eightbit_alu(
    input signed [7:0] a,
    input signed [7:0] b,
    input [2:0] s,
    output reg [7:0] f,
    output reg ovf,
    output reg take_branch
);
    always @(a, b, s) begin
        // operations on 'f'
        case(s)
            0: f = a + b;
            1: f = ~b;
            2: f = a & b;
            3: f = a | b;
            4: f = a >>> 1;
            5: f = a << 1;
            6: f = 0;
            7: f = 0;
        endcase
        // overflow check for addition
        ovf = (s == 0) && (f[7] != a[7]) && (f[7] != b[7]);
        // 'take_branch' special cases
        take_branch = ((s == 6) && (a == b)) || ((s == 7) && (a != b));
    end
endmodule
```

Figure 6: `eightbit_alu.v` implementing an 8-bit ALU in Verilog.

D Hardware Test Results for Lab 1: 8-Bit Adder

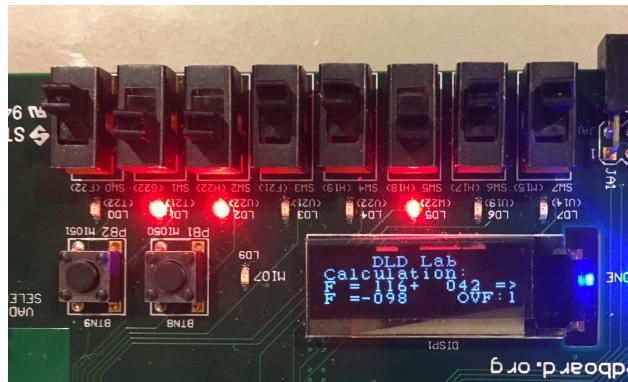


Figure 7: First test case.

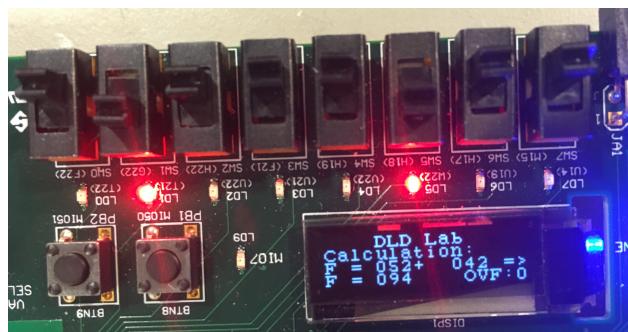


Figure 8: Second test case.

E Hardware Test Results for Lab 2: Partial Arithmetic Logic Unit

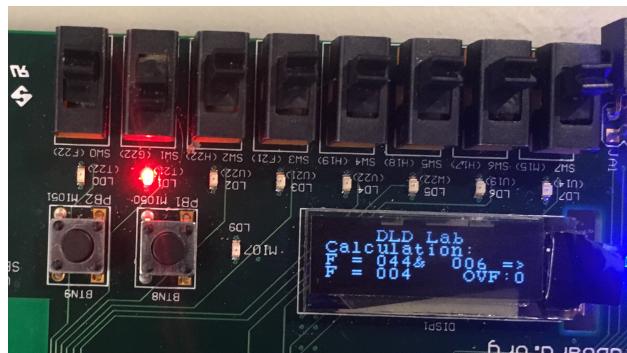


Figure 9: First test case.

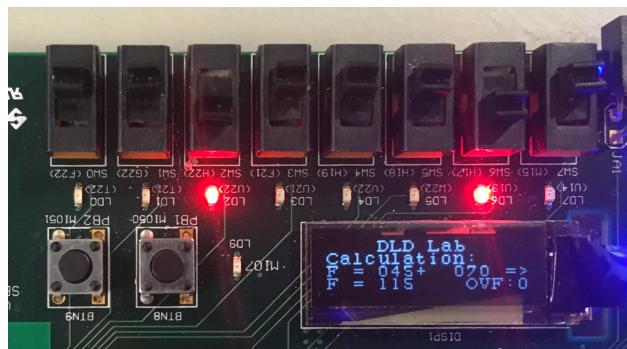


Figure 10: Second test case.



Figure 11: Third test case.

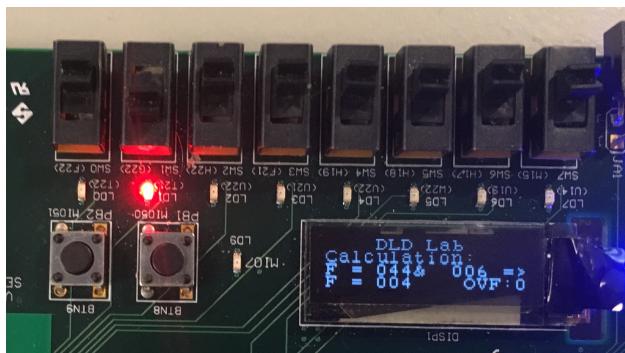


Figure 12: Fourth test case.

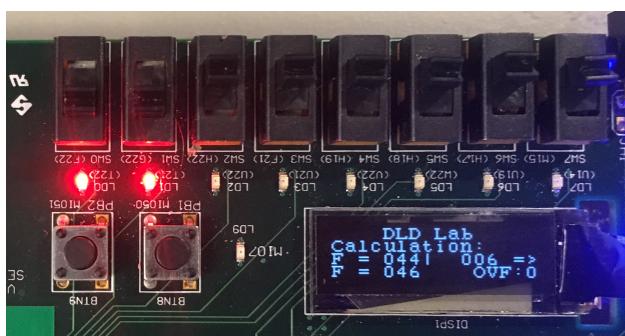


Figure 13: Fifth test case.

F Hardware Test Results for Lab 3: Arithmetic Logic Unit

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 000		Output	hw_vio_1
alu_output[7:0]	[B] 0110_1000		Input	hw_vio_1
alu_ovf_flag	[B] 1		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 14: First test case – $a + b$.

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 001		Output	hw_vio_1
alu_output[7:0]	[B] 0100_1011		Input	hw_vio_1
alu_ovf_flag	[B] 0		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 15: Second test case – $\sim b$.

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1000_0000		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 010		Output	hw_vio_1
alu_output[7:0]	[B] 1000_0000		Input	hw_vio_1
alu_ovf_flag	[B] 0		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 16: Third test case – $a \cdot b$.

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1000_0000		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 011		Output	hw_vio_1
alu_output[7:0]	[B] 1011_0100		Input	hw_vio_1
alu_ovf_flag	[B] 0		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 17: Fourth test case – $a|b$.

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1000_0000		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 100		Output	hw_vio_1
alu_output[7:0]	[B] 1101_1010		Input	hw_vio_1
alu_ovf_flag	[B] 0		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 18: Fifth test case – $a >>> 1$.

Name	Value	Activity	Direction	VIO
alu_1st_input[7:0]	[B] 1011_0100		Output	hw_vio_1
alu_2nd_input[7:0]	[B] 1000_0000		Output	hw_vio_1
alu_operation_sel[2:0]	[B] 101		Output	hw_vio_1
alu_output[7:0]	[B] 0110_1000		Input	hw_vio_1
alu_ovf_flag	[B] 0		Input	hw_vio_1
alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 19: Sixth test case – $a << 1$.

hw_vio_1					
	Name	Value	Activity	Direction	VIO
[+]	alu_1st_input[7:0]	[B] 1011_0100	▼	Output	hw_vio_1
[+]	alu_2nd_input[7:0]	[B] 1000_0000	▼	Output	hw_vio_1
[+]	alu_operation_sel[2:0]	[B] 110	▼	Output	hw_vio_1
[+]	alu_output[7:0]	[B] 0000_0000	▼	Input	hw_vio_1
-	alu_ovf_flag	[B] 0		Input	hw_vio_1
-	alu_take_branch_output	[B] 0		Input	hw_vio_1

Figure 20: Seventh test case – $beq(a, b)$ ($a \neq b$).

hw_vio_1					
	Name	Value	Activity	Direction	VIO
[+]	alu_1st_input[7:0]	[B] 1011_0100	▼	Output	hw_vio_1
[+]	alu_2nd_input[7:0]	[B] 1011_0100	▼	Output	hw_vio_1
[+]	alu_operation_sel[2:0]	[B] 110	▼	Output	hw_vio_1
[+]	alu_output[7:0]	[B] 0000_0000		Input	hw_vio_1
-	alu_ovf_flag	[B] 0		Input	hw_vio_1
-	alu_take_branch_output	[B] 1	↑	Input	hw_vio_1

Figure 21: Eighth test case – $beq(a, b)$ ($a = b$).

hw_vio_1					
	Name	Value	Activity	Direction	VIO
[+]	alu_1st_input[7:0]	[B] 1011_0100	▼	Output	hw_vio_1
[+]	alu_2nd_input[7:0]	[B] 1011_0100	▼	Output	hw_vio_1
[+]	alu_operation_sel[2:0]	[B] 111	▼	Output	hw_vio_1
[+]	alu_output[7:0]	[B] 0000_0000		Input	hw_vio_1
-	alu_ovf_flag	[B] 0		Input	hw_vio_1
-	alu_take_branch_output	[B] 0	▼	Input	hw_vio_1

Figure 22: Ninth test case – $bne(a, b)$ ($a = b$).