

Generating LLVM IR using Template Meta Programming

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Accelerate

- Embedded array language
- Compiled online
- CUDA backend

LLVM Backend

- CPU/GPU
- no external process
- incomplete

Overview LLVM

- Compiler library
- Portable
- Haskell bindings (llvm-general)
- Auto-Vectorization
- Static Single Assignment (SSA)

Code Example LLVM

```
int abs(int x) {  
    int res;  
    if (x>0) {  
        res = x;  
    } else {  
        res = -x;  
    }  
    return res;  
}
```

```
define i32 @abs(i32 %x) {  
entry:  
    %1 = icmp sgt i32 %x, 0  
    br i1 %1, label %then, label %else  
  
then:                                ; preds = %entry  
    br label %5  
  
else:                                ; preds = %entry  
    %4 = sub nsw i32 0, %x  
    br label %5  
  
end:                                  ; preds = %else, %then  
    %res.0 = phi i32 [ %x, %then ],  
                  [ %4, %else ]  
    ret i32 %res.0  
}
```

llvm-general

- uses ADT to represent syntax tree
- translates from and to C++ Objects
- exposes Optimizer/JIT compiler

Quasiquoting

```
abs :: Global
abs = [llg|
  define i32 @abs(i32 %x) {
  entry:
    %1 = icmp sgt i32 %x, 0
    br i1 %1, label %end, label %else

  else:
    %4 = sub nsw i32 0, %x
    br label %5

  end:
    %res.0 = phi i32 [ %x, %entry ],
                  [ %4, %else ]
    ret i32 %res.0
  }
|]
```

Control Structures

```
i32 foo(i32 %m, i32 %n) {  
  entry:  
    br label %for  
  
  for:  
    for i32 %i in %m to %n with i32 [ 0, %entry ] as %j,  
                                label %end {  
      %k = add i32 %i, %j  
      ret i32 %k  
    }  
  
  end:  
    ret i32 %j  
}
```


Control Structures

```
i32 foo(i32 %m, i32 %n) {  
    entry:  
    br label %for  
  
    for:  
    for i32 %i in %m to %n with i32 [ 0, %entry ] as %j {  
        %k = add i32 %i, %j  
        ret i32 %k  
    }  
  
    end:  
    ret i32 %j  
}
```

Control Structures

```
define i64 @foo(i64 %start, i64 %end) {  
  entry:  
    %x = i64 0  
  
  for:  
    for i64 %i in %start to %end {  
      %x = add i64 %i, %x  
    }  
  
  exit:  
    ret i64 %x  
}
```

SSA



BRAUN, Matthias et al.: Simple and Efficient Construction of Static Single Assignment Form. In: JHALA, Ranjit ; BOSSCHERE, Koen (eds.): *Compiler Construction*. Vol. 7791. Lecture Notes in Computer Science. Springer Berlin Heidelberg, 2013, pp. 102–122

Local value numbering

```
writeVariable(variable, block, value):  
    currentDef[variable][block] <- value  
  
readVariable(variable, block):  
    if currentDef[variable] contains block:  
        # local value numbering  
        return currentDef[variable][block]  
    # global value numbering  
    return readVariableRecursive(variable, block)
```

Global value numbering

```
readVariableRecursive(variable, block):  
  if |block.preds| = 0:  
    # First block  
    val <- variable  
  else:  
    # Break potential cycles with operandless phi  
    val <- new Phi(block)  
    writeVariable(variable, block, val)  
  writeVariable(variable, block, val)  
  return val
```

```

define i64 @foo(i64 %start, i64 %end) {
entry:
    br label %for.head

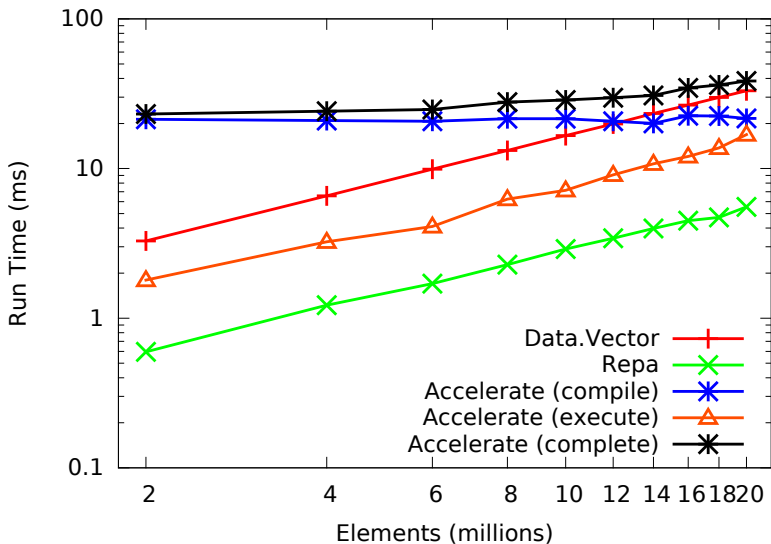
for.head:                                ; preds = %n0, %entry
    %x.12 = phi i64 [ 0, %entry ], [ %x.6, %n0 ]
    %i.4 = phi i64 [ %start, %entry ], [ %i.9, %n0 ]
    %for.cond.3 = icmp slt i64 %i.4, %end
    br i1 %for.cond.3, label %n0, label %for.end

n0:                                       ; preds = %for.head
    %x.6 = add i64 %i.4, %x.12
    %i.9 = add nuw nsw i64 %i.4, 1
    br label %for.head

for.end:                                ; preds = %for.head
    ret i64 %x.12
}

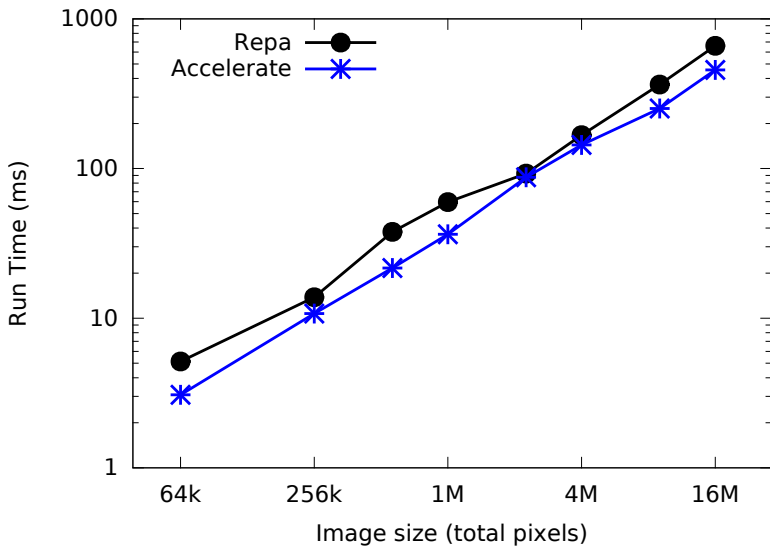
```

Dot product



Step	Time (ms)
getTarget	0.098
llvmOfAcc	0.0771
startFunction	0.0403
withContext	0.0879
withModuleFromAST	6.8032
withMachine	0.2997
libinfo	0.0801
optimizeModule	4.0379
withMCJIT	0.05
withModuleInEngine	0.2374
getGlobalFunctions	9.8508

Ray tracer



Ray tracer

