

metavariable, x, xs, y, uf, f, d

term, t, u	$::=$ $ $ x $ $ z $ $ $t\ u$ $ $ $t\ ;\ u$ $ $ $\text{case } t \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\}$ $ $ $\text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}$ $ $ $\text{case } t \text{ of } \{U_r\ x \mapsto u\}$ $ $ $\text{case } t \text{ of } \{\text{roll}_R\ x \mapsto u\}$ $ $ $\text{alloc}_A\ d.\ t$ $ $ $t \triangleleft u$ $ $ $t \triangleleft 1.d'.u$ $ $ $t \triangleleft 2.d'.u$ $ $ $t \triangleleft \langle d_1, d_2 \rangle .u$ $ $ $\star l$ $ $ (t) $ $ $\backslash n\ \text{sp}\ t\ \backslash n\ \text{spe}$ $ $ $\backslash n\ \text{sp}\ t$ $ $ $t[\text{var_subs}]$	term variable value application pattern-matching on unit pattern-matching on sum pattern-matching on product pattern-matching on exponentiated value unroll for recursive types allocate data fill terminal-type destination fill sum-type destination with variant 1 fill sum-type destination with variant 2 fill product-type destination note: $\star l$ is not a part of the user syntax S S S M
var_sub, vs	$::=$ $ $ $x := t$	variable substitution
var_subs	$::=$ $ $ vs $ $ $vs, \text{var_subs}$	variable substitutions
data_val, v	$::=$ $ $ $()$ $ $ $\lambda x:A.\ t$ $ $ $U_r\ t$ $ $ $\text{roll}_R\ t$	data value unit lambda abstraction exponential roll for recursive types
val, z	$::=$ $ $ v $ $ $[A]_l$ $ $ $1.l$ $ $ $2.l$ $ $ $\langle l_1, l_2 \rangle$ $ $ $C\bar{l}$	unreducible value note: l is not a part of the user syntax M
<i>label</i> , l	$::=$	label
label_stmt, s	$::=$ $ $ $l \triangleleft v$ $ $ $l \triangleleft 1.l'$ $ $ $l \triangleleft 2.l'$ $ $ $l \triangleleft \langle l_1, l_2 \rangle$ $ $ $l \triangleleft \emptyset$	label statement

	$ \begin{array}{ l} l \triangleleft C\bar{l} \\ l \triangleleft \dots \end{array} $	$ \begin{array}{ l} M \\ M \end{array} $	<p>TODO: hide. $l \triangleleft C\bar{l}$ is an alias for any heap cons</p>
label_stmts	$ \begin{array}{ l} s \\ s, \text{label_stmts} \end{array} $		label statements
heap_context, \mathbb{H}	$ \begin{array}{ l} \emptyset \\ \{\text{label_stmts}\} \\ \mathbb{H}_1 \sqcup \mathbb{H}_2 \end{array} $		label statements
type, A, B	$ \begin{array}{ l} 1 \\ R \\ A \otimes B \\ A \oplus B \\ A \multimap B \\ [A] \\ !A \\ (A) \\ W[r := A] \end{array} $	$ \begin{array}{ l} S \\ M \end{array} $	<p>unit type</p> <p>recursive type bound to a name</p> <p>product type</p> <p>sum type</p> <p>linear function type</p> <p>destination type</p> <p>exponential</p>
type_with_hole, W	$ \begin{array}{ l} r \\ 1 \\ R \\ W_1 \otimes W_2 \\ W_1 \oplus W_2 \\ W_1 \multimap W_2 \\ [W] \\ !W \\ (W) \end{array} $	$ \begin{array}{ l} S \end{array} $	<p>type hole in recursive definition</p> <p>unit type</p> <p>recursive type bound to a name</p> <p>product type</p> <p>sum type</p> <p>linear function type</p> <p>destination type</p> <p>exponential</p>
rec_type_bound, R	$ \begin{array}{ l} \mu r. W \end{array} $		recursive type bound to a name
rec_type_def	$ \begin{array}{ l} \mu r. W \end{array} $		
type_affect, ta	$ \begin{array}{ l} x : A \\ l : A \\ \bar{l} : \bar{A} \end{array} $		<p>type affectation</p> <p>var</p> <p>label</p> <p>labels</p>
type_affects	$ \begin{array}{ l} \text{ta} \\ \text{ta}, \text{type_affects} \end{array} $		type affectations
typing_context, $\Gamma, \Delta, \mathcal{U}, \Phi$	$ \begin{array}{ l} \end{array} $		typing context

	\emptyset $\{\text{type_affects}\}$ $\Gamma \sqcup \Delta$	
types, \bar{A}	$::=$ \bullet A $A \text{ types}$	empty type list
heap_constructor, C	$::=$ $(1.)$ $(2.)$ (\langle, \rangle)	
judg	$::=$ $\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H}$ $\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A$ $\Phi ; \mathcal{U} ; \Gamma \vdash t : A$ $C : \bar{A} \xrightarrow{c} A$ $A = B$ $t = u$ $\Gamma = \Delta$ $\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'$ $\text{type_affect} \in \Gamma$ $\text{label_stmt} \in \mathbb{H}$ $\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset$	
terminals	$::=$ \mapsto \star \otimes \oplus \dashv $:=$ \vdash \sqcup $;$ \cap \emptyset \longrightarrow \triangleright \neq \in \notin $\backslash n$ \langle \rangle $1.$ $2.$ Ur	

	\triangleleft $ $ \bigcirc $\underline{\hookrightarrow}$ $=$ \Downarrow \dots $\underline{\text{fix}}$	
formula	$::=$ $ $ judgement	
Ctx	$::=$ $ $ $\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset$ $ $ $\text{type_affect} \in \Gamma$	Γ and Δ are disjoint typing contexts with no clashing
Heap	$::=$ $ $ $\text{label_stmt} \in \mathbb{H}$	
Eq	$::=$ $ $ $A = B$ $ $ $t = u$ $ $ $\Gamma = \Delta$	
Ty	$::=$ $ $ $R \stackrel{\text{fix}}{=} \text{rec_type_def}$ $ $ $\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H}$ $ $ $\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A$ $ $ $\Phi ; \mathcal{U} ; \Gamma \vdash t : A$ $ $ $C : \bar{A} \hookrightarrow A$	\mathbb{H} is a well-typed heap given heap typing context Φ , u t is a well-typed term of type A given heap typing con Heap constructor C builds a value of type A given arg
Sem	$::=$ $ $ $\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'$ $ $ $\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'$	t reduces to t' , with heap changing from \mathbb{H} to \mathbb{H}' t reduces to t' , with heap growing from \mathbb{H} to \mathbb{H}'
judgement	$::=$ $ $ Ctx $ $ Heap $ $ Eq $ $ Ty $ $ Sem	
user_syntax	$::=$ $ $ metavariable $ $ term $ $ var_sub $ $ var_subs $ $ data_val $ $ val	

- | *label*
- | label_stmt
- | label_stmts
- | heap_context
- | type
- | type_with_hole
- | rec_type_bound
- | rec_type_def
- | type_affect
- | type_affects
- | typing_context
- | types
- | heap_constructor
- | judg
- | terminals

$\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset$

Γ and Δ are disjoint typing contexts with no clashing variable names or labels

$\text{type_affect} \in \Gamma$

$\text{label_stmt} \in \mathbb{H}$

$A = B$

$t = u$

$\Gamma = \Delta$

$R \stackrel{\text{fix}}{=} \text{rec_type_def}$

$\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H}$

\mathbb{H} is a well-typed heap given heap typing context Φ , unrestricted typing context \mathcal{U} and linear

$$\frac{\begin{array}{l} \Phi_1 ; \mathcal{U} ; \Gamma \vdash \mathbb{H}_1 \\ \Phi_2 ; \mathcal{U} ; \Delta \vdash \mathbb{H}_2 \\ \text{names}(\Phi_1) \cap \text{names}(\Phi_2) = \emptyset \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi_1 \sqcup \Phi_2 ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \mathbb{H}_1 \sqcup \mathbb{H}_2} \text{TYHEAP_UNION}$$

$$\frac{\begin{array}{l} C : \bar{A} \hookrightarrow A \\ \Phi \sqcup \{\bar{l} : \bar{A}\} ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \end{array}}{\Phi \sqcup \{\bar{l} : \bar{A}, l : A\} ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \sqcup \{l \triangleleft C\bar{l}\}} \text{TYHEAP_CTOR}$$

$$\overline{\{l : A\} ; \mathcal{U} ; \emptyset \vdash \{l \triangleleft \emptyset\}} \text{TYHEAP_NULL}$$

$$\frac{\Phi ; \mathcal{U} ; \Gamma \vdash v : A}{\Phi \sqcup \{l : A\} ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \sqcup \{l \triangleleft v\}} \text{TYHEAP_VAL}$$

$\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \\ \Phi ; \mathcal{U} ; \Delta \vdash t : A \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A} \text{TYCOMMAND_DEF}$$

$\Phi ; \mathcal{U} ; \Gamma \vdash t : A$

t is a well-typed term of type A given heap typing context Φ , unrestricted typing context

$$\overline{\Phi ; \mathcal{U} ; \{x : A\} \vdash x : A} \text{TYTERM_ID}$$

$$\begin{array}{c}
\frac{}{\Phi ; \mathcal{U} \sqcup \{x : A\} ; \emptyset \vdash x : A} \text{TYTERM_ID} \\
\\
\frac{}{\Phi ; \mathcal{U} ; \emptyset \vdash () : 1} \text{TYTERM_UNIT} \\
\\
\frac{\Phi ; \mathcal{U} ; \emptyset \vdash t : A}{\Phi ; \mathcal{U} ; \emptyset \vdash \text{Ur } t : !A} \text{TYTERM_EXP} \\
\\
\frac{}{\Phi ; \mathcal{U} ; \emptyset \vdash \textcolor{blue}{[A]} : \textcolor{blue}{[A]}} \text{TYTERM_LABELASDEST} \\
\\
\frac{}{\Phi \sqcup \{l : A\} ; \mathcal{U} ; \emptyset \vdash \star l : A} \text{TYTERM_DEREF} \\
\\
\frac{C : \bar{A} \xrightarrow{c} A}{\Phi \sqcup \{\bar{l} : \bar{A}\} ; \mathcal{U} ; \emptyset \vdash C \bar{l} : A} \text{TYTERM_CTOR} \\
\\
\frac{R \stackrel{\text{fix}}{=} \mu r . W}{\Phi ; \mathcal{U} ; \Gamma \vdash t : W[r := R]} \text{TYTERM_ROLL} \\
\frac{}{\Phi ; \mathcal{U} ; \Gamma \vdash \text{roll } t : R} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \sqcup \{x : A\} \vdash t : B}{\Phi ; \mathcal{U} ; \Gamma \vdash \lambda x : A . t : A \multimap B} \text{TYTERM_LAM} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \vdash t : A \multimap B \quad \Phi ; \mathcal{U} ; \Delta \vdash u : A}{\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset} \text{TYTERM_APP} \\
\frac{}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t u : B} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \vdash t : 1 \quad \Phi ; \mathcal{U} ; \Delta \vdash u : A \quad \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t ; u : A} \text{TYTERM_PATU} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \vdash t : A_1 \otimes A_2 \quad \Phi ; \mathcal{U} ; \Delta \sqcup \{x_1 : A_1, x_2 : A_2\} \vdash u : B \quad \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : B} \text{TYTERM_PATP} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \vdash t : A_1 \oplus A_2 \quad \Phi ; \mathcal{U} ; \Delta \sqcup \{x_1 : A_1\} \vdash u_1 : B \quad \Phi ; \mathcal{U} ; \Delta \sqcup \{x_2 : A_2\} \vdash u_2 : B \quad \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\} : B} \text{TYTERM_PATs} \\
\\
\frac{\Phi ; \mathcal{U} ; \Gamma \vdash t : !A \quad \Phi ; \mathcal{U} \sqcup \{x : A\} ; \Delta \vdash u : B \quad \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} : B} \text{TYTERM_PATE} \\
\\
\frac{R \stackrel{\text{fix}}{=} \mu r . W \quad \Phi ; \mathcal{U} ; \Gamma \vdash t : R \quad \Phi ; \mathcal{U} ; \Delta \sqcup \{x : W[r := R]\} \vdash u : B \quad \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{\text{roll } x \mapsto u\} : B} \text{TYTERM_PATR}
\end{array}$$

$$\begin{array}{c}
\frac{\Phi ; \mathcal{U} ; \Gamma \sqcup \{\mathbf{d} : \mathbf{[A]}\} \vdash t : \mathbf{1}}{\Phi ; \mathcal{U} ; \Gamma \vdash \text{alloc } \mathbf{d} . t : \mathbf{A}} \quad \text{TYTERM_ALLOC} \\
\\
\frac{\begin{array}{c} \Phi ; \mathcal{U} ; \Gamma \vdash t : \mathbf{[A]} \\ \Phi ; \mathcal{U} ; \Delta \vdash u : \mathbf{A} \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft u : \mathbf{1}} \quad \text{TYTERM_FILL} \\
\\
\frac{\begin{array}{c} \Phi ; \mathcal{U} ; \Gamma \vdash t : \mathbf{[A_1 \oplus A_2]} \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{\mathbf{d}' : \mathbf{[A_1]}\} \vdash u : \mathbf{B} \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft \mathbf{1.d' . u} : \mathbf{B}} \quad \text{TYTERM_FILLV1} \\
\\
\frac{\begin{array}{c} \Phi ; \mathcal{U} ; \Gamma \vdash t : \mathbf{[A_1 \oplus A_2]} \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{\mathbf{d}' : \mathbf{[A_2]}\} \vdash u : \mathbf{B} \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \vdash t \triangleleft \mathbf{2.d' . u} : \mathbf{B}} \quad \text{TYTERM_FILLV2} \\
\\
\frac{\begin{array}{c} \Phi ; \mathcal{U} ; \Gamma \vdash t : \mathbf{[A_1 \otimes A_2]} \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{\mathbf{d}_1 : \mathbf{[A_1]}, \mathbf{d}_2 : \mathbf{[A_2]}\} \vdash u : \mathbf{B} \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft \langle \mathbf{d}_1, \mathbf{d}_2 \rangle . u : \mathbf{B}} \quad \text{TYTERM_FILLP}
\end{array}$$

$\boxed{\mathbf{C} : \bar{\mathbf{A}} \hookrightarrow \mathbf{A}}$ Heap constructor \mathbf{C} builds a value of type \mathbf{A} given arguments of type $\bar{\mathbf{A}}$

$$\begin{array}{c}
\frac{}{(1.) : \mathbf{A} \hookrightarrow \mathbf{A \oplus B}} \quad \text{TYCTOR_V1} \\
\\
\frac{}{(2.) : \mathbf{B} \hookrightarrow \mathbf{A \oplus B}} \quad \text{TYCTOR_V2} \\
\\
\frac{}{(\langle, \rangle) : \mathbf{A} \ \mathbf{B} \hookrightarrow \mathbf{A \otimes B}} \quad \text{TYCTOR_PAIR}
\end{array}$$

$\boxed{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}$ t reduces to t' , with heap changing from \mathbb{H} to \mathbb{H}'

$$\begin{array}{c}
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t u \longrightarrow \mathbb{H}' \mid t' u} \quad \text{SEMMUT_UAPP} \\
\\
\frac{}{\mathbb{H} \mid (\lambda \mathbf{x} : \mathbf{A} . t) u \longrightarrow \mathbb{H} \mid t[\mathbf{x} := u]} \quad \text{SEMMUT_APP} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t ; u \longrightarrow \mathbb{H}' \mid t' ; u} \quad \text{SEMMUT_UPATU} \\
\\
\frac{}{\mathbb{H} \mid () ; u \longrightarrow \mathbb{H} \mid u} \quad \text{SEMMUT_PATU} \\
\\
\frac{}{\mathbb{H} \sqcup \{\mathbf{l} \triangleleft \mathbf{v}\} \mid \star \mathbf{l} \longrightarrow \mathbb{H} \sqcup \{\mathbf{l} \triangleleft \mathbf{v}\} \mid \mathbf{v}} \quad \text{SEMMUT_DEREFVAL} \\
\\
\frac{}{\mathbb{H} \sqcup \{\mathbf{l} \triangleleft \mathbf{C}\bar{\mathbf{l}}\} \mid \star \mathbf{l} \longrightarrow \mathbb{H} \sqcup \{\mathbf{l} \triangleleft \mathbf{C}\bar{\mathbf{l}}\} \mid \mathbf{C}\bar{\mathbf{l}}} \quad \text{SEMMUT_DEREFCOR} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{1.\mathbf{x}_1 \mapsto u_1, 2.\mathbf{x}_2 \mapsto u_2\} \longrightarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{1.\mathbf{x}_1 \mapsto u_1, 2.\mathbf{x}_2 \mapsto u_2\}} \quad \text{SEMMUT_UPATS} \\
\\
\frac{}{\mathbb{H} \mid \text{case } 1.\mathbf{l} \text{ of } \{1.\mathbf{x}_1 \mapsto u_1, 2.\mathbf{x}_2 \mapsto u_2\} \longrightarrow \mathbb{H} \mid u_1[\mathbf{x}_1 := \star \mathbf{l}]} \quad \text{SEMMUT_PATSV1} \\
\\
\frac{}{\mathbb{H} \mid \text{case } 2.\mathbf{l} \text{ of } \{1.\mathbf{x}_1 \mapsto u_1, 2.\mathbf{x}_2 \mapsto u_2\} \longrightarrow \mathbb{H} \mid u_2[\mathbf{x}_2 := \star \mathbf{l}]} \quad \text{SEMMUT_PATSV2}
\end{array}$$

$$\begin{array}{c}
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} \longrightarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}} \quad \text{SEMMUT_UPATP} \\
\\
\frac{\mathbb{H} \mid \text{case } \langle l_1, l_2 \rangle \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} \longrightarrow \mathbb{H} \mid u[x_1 := \star l_1, x_2 := \star l_2]}{\mathbb{H} \mid \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} \longrightarrow \mathbb{H} \mid u[x_1 := \star l_1, x_2 := \star l_2]} \quad \text{SEMMUT_PATP} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} \longrightarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{\text{Ur } x \mapsto u\}} \quad \text{SEMMUT_UPATE} \\
\\
\frac{\mathbb{H} \mid \text{case Ur } t \text{ of } \{\text{Ur } x \mapsto u\} \longrightarrow \mathbb{H} \mid u[x := t]}{\mathbb{H} \mid \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} \longrightarrow \mathbb{H} \mid u[x := t]} \quad \text{SEMMUT_PATE} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{\text{roll}_R x \mapsto u\} \longrightarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{\text{roll}_R x \mapsto u\}} \quad \text{SEMMUT_UPATR} \\
\\
\frac{\mathbb{H} \mid \text{case roll}_R t \text{ of } \{\text{roll}_R x \mapsto u\} \longrightarrow \mathbb{H} \mid u[x := t]}{\mathbb{H} \mid \text{case } t \text{ of } \{\text{roll}_R x \mapsto u\} \longrightarrow \mathbb{H} \mid u[x := t]} \quad \text{SEMMUT_PATR} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft u \longrightarrow \mathbb{H}' \mid t' \triangleleft u} \quad \text{SEMMUT_UFILLL} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \underline{[A]} \triangleleft t \longrightarrow \mathbb{H}' \mid \underline{[A]} \triangleleft t'} \quad \text{SEMMUT_UFILLL'} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft 1.d'.u \longrightarrow \mathbb{H}' \mid t' \triangleleft 1.d'.u} \quad \text{SEMMUT_UFILLV1} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft 2.d'.u \longrightarrow \mathbb{H}' \mid t' \triangleleft 2.d'.u} \quad \text{SEMMUT_UFILLV2} \\
\\
\frac{\mathbb{H} \mid t \longrightarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft \langle d_1, d_2 \rangle . u \longrightarrow \mathbb{H}' \mid t' \triangleleft \langle d_1, d_2 \rangle . u} \quad \text{SEMMUT_UFILLP} \\
\\
\frac{\mathbb{H} \mid \text{alloc}_A d . t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid t[d := \underline{[A]}] ; \star l}{\mathbb{H} \mid \text{alloc}_A d . t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid t[d := \underline{[A]}] ; \star l} \quad \text{SEMMUT_ALLOC} \\
\\
\frac{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A]} \triangleleft v \longrightarrow \mathbb{H} \sqcup \{l \triangleleft v\} \mid ()}{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A]} \triangleleft v \longrightarrow \mathbb{H} \sqcup \{l \triangleleft v\} \mid ()} \quad \text{SEMMUT_FILLV} \\
\\
\frac{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A]} \triangleleft C\bar{l} \longrightarrow \mathbb{H} \sqcup \{l \triangleleft C\bar{l}\} \mid ()}{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A]} \triangleleft C\bar{l} \longrightarrow \mathbb{H} \sqcup \{l \triangleleft C\bar{l}\} \mid ()} \quad \text{SEMMUT_FILLCTOR} \\
\\
\frac{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \oplus A_2]} \triangleleft 1.d'.t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft 1.l', l' \triangleleft \emptyset\} \mid t[d' := \underline{[A_1]}]}{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \oplus A_2]} \triangleleft 1.d'.t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft 1.l', l' \triangleleft \emptyset\} \mid t[d' := \underline{[A_1]}]} \quad \text{SEMMUT_FILLV1} \\
\\
\frac{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \oplus A_2]} \triangleleft 2.d'.t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft 2.l', l' \triangleleft \emptyset\} \mid t[d' := \underline{[A_2]}]}{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \oplus A_2]} \triangleleft 2.d'.t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft 2.l', l' \triangleleft \emptyset\} \mid t[d' := \underline{[A_2]}]} \quad \text{SEMMUT_FILLV2} \\
\\
\frac{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \otimes A_2]} \triangleleft \langle d_1, d_2 \rangle . t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft \langle l_1, l_2 \rangle, l_1 \triangleleft \emptyset, l_2 \triangleleft \emptyset\} \mid t[d_1 := \underline{[A_1]}, d_2 := \underline{[A_2]}]}{\mathbb{H} \sqcup \{l \triangleleft \emptyset\} \mid \underline{[A_1 \otimes A_2]} \triangleleft \langle d_1, d_2 \rangle . t \longrightarrow \mathbb{H} \sqcup \{l \triangleleft \langle l_1, l_2 \rangle, l_1 \triangleleft \emptyset, l_2 \triangleleft \emptyset\} \mid t[d_1 := \underline{[A_1]}, d_2 := \underline{[A_2]}]} \quad \text{SEMMUT_FILLP} \\
\\
\boxed{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'} \quad t \text{ reduces to } t', \text{ with heap growing from } \mathbb{H} \text{ to } \mathbb{H}'
\end{array}$$

$$\begin{array}{c}
\frac{\mathbb{H} \mid \text{roll}_R t \Downarrow \mathbb{H} \mid t}{\mathbb{H} \mid \text{roll}_R t \Downarrow \mathbb{H} \mid t} \quad \text{SEMMUT_ROLL} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t u \Downarrow \mathbb{H}' \mid t' u} \quad \text{SEMMUT_UAPP} \\
\\
\frac{\mathbb{H} \mid (\lambda x:A. t) u \Downarrow \mathbb{H} \mid t[x := u]}{\mathbb{H} \mid (\lambda x:A. t) u \Downarrow \mathbb{H} \mid t[x := u]} \quad \text{SEMMUT_APP}
\end{array}$$

$$\begin{array}{c}
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t ; u \Downarrow \mathbb{H}' \mid t' ; u} \text{SEMIMM_UPATU} \\
\\
\frac{}{\mathbb{H} \mid () ; u \Downarrow \mathbb{H} \mid u} \text{SEMIMM_PATU} \\
\\
\frac{}{\mathbb{H} \sqcup \{l \triangleleft v\} \mid \star l \Downarrow \mathbb{H} \sqcup \{l \triangleleft v\} \mid v} \text{SEMIMM_DEREFVAL} \\
\\
\frac{}{\mathbb{H} \sqcup \{l \triangleleft C\bar{l}\} \mid \star l \Downarrow \mathbb{H} \sqcup \{l \triangleleft C\bar{l}\} \mid C\bar{l}} \text{SEMIMM_DEREFTOR} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\} \Downarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\}} \text{SEMIMM_UPATS} \\
\\
\frac{}{\mathbb{H} \mid \text{case } 1.l \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\} \Downarrow \mathbb{H} \mid u_1[x_1 := \star l]} \text{SEMIMM_PATSV1} \\
\\
\frac{}{\mathbb{H} \mid \text{case } 2.l \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\} \Downarrow \mathbb{H} \mid u_2[x_2 := \star l]} \text{SEMIMM_PATSV2} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} \Downarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}} \text{SEMIMM_UPATP} \\
\\
\frac{}{\mathbb{H} \mid \text{case } \langle l_1, l_2 \rangle \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} \Downarrow \mathbb{H} \mid u[x_1 := \star l_1, x_2 := \star l_2]} \text{SEMIMM_PATP} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{Ur \ x \mapsto u\} \Downarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{Ur \ x \mapsto u\}} \text{SEMIMM_UPATE} \\
\\
\frac{}{\mathbb{H} \mid \text{case } Ur \ t \text{ of } \{Ur \ x \mapsto u\} \Downarrow \mathbb{H} \mid u[x := t]} \text{SEMIMM_PATE} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \text{case } t \text{ of } \{\text{roll}_R \ x \mapsto u\} \Downarrow \mathbb{H}' \mid \text{case } t' \text{ of } \{\text{roll}_R \ x \mapsto u\}} \text{SEMIMM_UPATR} \\
\\
\frac{}{\mathbb{H} \mid \text{case } \text{roll}_R \ t \text{ of } \{\text{roll}_R \ x \mapsto u\} \Downarrow \mathbb{H} \mid u[x := t]} \text{SEMIMM_PATR} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft u \Downarrow \mathbb{H}' \mid t' \triangleleft u} \text{SEMIMM_UFILLL} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid \overset{l}{[A]} \triangleleft t \Downarrow \mathbb{H}' \mid \overset{l}{[A]} \triangleleft t'} \text{SEMIMM_UFILLL'} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft 1.d' . u \Downarrow \mathbb{H}' \mid t' \triangleleft 1.d' . u} \text{SEMIMM_UFILLV1} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft 2.d' . u \Downarrow \mathbb{H}' \mid t' \triangleleft 2.d' . u} \text{SEMIMM_UFILLV2} \\
\\
\frac{\mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t'}{\mathbb{H} \mid t \triangleleft \langle d_1, d_2 \rangle . u \Downarrow \mathbb{H}' \mid t' \triangleleft \langle d_1, d_2 \rangle . u} \text{SEMIMM_UFILLP} \\
\\
\frac{\mathbb{H} \mid t[d := \overset{l}{[A]}] \Downarrow \mathbb{H}' \sqcup \{l \triangleleft \dots\} \mid ()}{\mathbb{H} \mid \text{alloc}_A \ d . t \Downarrow \mathbb{H}' \sqcup \{l \triangleleft \dots\} \mid \star l} \text{SEMIMM_ALLOC} \quad \text{--- Changes start from there ---} \\
\\
\frac{}{\mathbb{H} \mid \overset{l}{[A]} \triangleleft v \Downarrow \mathbb{H} \sqcup \{l \triangleleft v\} \mid ()} \text{SEMIMM_FILLV} \\
\\
\frac{}{\mathbb{H} \mid \overset{l}{[A]} \triangleleft C\bar{l} \Downarrow \mathbb{H} \sqcup \{l \triangleleft C\bar{l}\} \mid ()} \text{SEMIMM_FILLLCTOR}
\end{array}$$

$$\begin{array}{c}
\frac{\mathbb{H} \mid \mathfrak{t}[\mathbf{d}' := \underline{A_1}] \Downarrow \mathbb{H}' \sqcup \{l' \triangleleft \dots\} \mid ()}{\mathbb{H} \mid \underline{A_1 \oplus A_2} \triangleleft \mathbf{1.d'.t} \Downarrow \mathbb{H}' \sqcup \{l' \triangleleft \dots, l \triangleleft \mathbf{1.l'}\} \mid ()} \text{SEMIMM_FILLV1} \\
\\
\frac{\mathbb{H} \mid \mathfrak{t}[\mathbf{d}' := \underline{A_2}] \Downarrow \mathbb{H}' \sqcup \{l' \triangleleft \dots\} \mid ()}{\mathbb{H} \mid \underline{A_1 \oplus A_2} \triangleleft \mathbf{2.d'.t} \Downarrow \mathbb{H}' \sqcup \{l' \triangleleft \dots, l \triangleleft \mathbf{2.l'}\} \mid ()} \text{SEMIMM_FILLV2} \\
\\
\frac{\mathbb{H} \mid \mathfrak{t}[\mathbf{d_1} := \underline{A_1}, \mathbf{d_2} := \underline{A_2}] \Downarrow \mathbb{H}' \sqcup \{l_1 \triangleleft \dots, l_2 \triangleleft \dots\} \mid ()}{\mathbb{H} \mid \underline{A_1 \otimes A_2} \triangleleft \langle \mathbf{d_1}, \mathbf{d_2} \rangle . \mathfrak{t} \Downarrow \mathbb{H}' \sqcup \{l_1 \triangleleft \dots, l_2 \triangleleft \dots, l \triangleleft \langle l_1, l_2 \rangle\} \mid ()} \text{SEMIMM_FILLP}
\end{array}$$

Definition rules: 81 good 0 bad
 Definition rule clauses: 155 good 0 bad