

Destination λ -calculus

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1 Term and value syntax

termvar, x, y, d Term-level variable
holevar, h Hole

term_value, v ::=

- $\langle v_1, \bar{v}_2 \rangle_H$
- $@h$
- $()$
- $\text{Inl } v$
- $\text{Inr } v$
- (v_1, v_2)
- $\rangle^m v$
- $\lambda x. t$

Term value

- Ampar
- Destination
- Unit
- Left variant for sum
- Right variant for sum
- Product
- Exponential
- Linear function

$\overline{\text{extended_value}}, \bar{v}$::=

- v
- h
- $\text{Inl } \bar{v}$
- $\text{Inr } \bar{v}$
- (\bar{v}_1, \bar{v}_2)
- $\rangle^m \bar{v}$

Pseudo-value that may contain holes

- Term value
- Hole
- Left variant with val or hole
- Right variant with val or hole
- Product with val or hole
- Exponential with val or hole

term, t, u ::=

- v
- x
- $t \succ u$
- $t \succ \text{case } () \mapsto u$
- $t \succ \text{case } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$
- $t \succ \text{case } (x_1, x_2) \mapsto u$
- $t \succ \text{case } \rangle^m x \mapsto u$
- $t \succ \text{mapL } x \mapsto u$
- $\text{to}_x t$
- $\text{from}_x t$
- alloc_A
- $t \triangleleft ()$
- $t \triangleleft \text{Inl}$
- $t \triangleleft \text{Inr}$
- $t \triangleleft (,)$
- $t \triangleleft \rangle^m$
- $t \triangleleft \bullet u$

Term

- Term value
- Variable
- Application
- Pattern-match on unit
- Pattern-match on sum
- Pattern-match on product
- Pattern-match on exponential
- Map over the left side of the ampar
- Wrap t into a trivial ampar
- Extract value from trivial ampar
- Return a fresh "identity" ampar object
- Fill destination with unit
- Fill destination with left variant
- Fill destination with right variant
- Fill destination with product constructor
- Fill destination with exponential constructor
- Fill destination with root of ampar u

2 Type system

type, A, B	$::=$	1 $A_1 \oplus A_2$ $A_1 \otimes A_2$ $!^m A$ $A_1 \ltimes A_2$ $A_1 \xrightarrow{m} A_2$ $^m[A]$	Type Unit Sum Product Exponential Ampar type (consuming A_1 yields A_2) Linear function Destination
multiplicity, m, n	$::=$	ν \uparrow ∞ $m_1 \cdot m_2$	Multiplicity (Semiring with product \cdot) Born now. Identity of the product One scope older Infinitely old / static. Absorbing for product Semiring product
typing_context, Δ	$::=$	Γ H $\Gamma \sqcup H$ $m \cdot \Delta$	Typing context Increase age of bindings by m
pos_context, Γ	$::=$	$\{\text{pos_assign}^*\}$ $\Gamma_1 \sqcup \Gamma_2$ $m \cdot \Gamma$	Positive typing context Increase age of bindings by m
pos_assign	$::=$	$x :_m A$ $@h :_m ^n[A]$	Positive type assignment Variable Destination (m is its own age; n is the age of values it accepts)
neg_context, H	$::=$	$\{\text{neg_assign}^*\}$ $H_1 \sqcup H_2$ $@^{-1} \Gamma$ $m \cdot H$	Negative typing context Inverse the sign of the context Increase age of bindings by m
neg_assign	$::=$	$h :^n A$	Negative type assignment Hole (n is the age of values it accepts, its own age is undefined)

$$\boxed{H_1 = H_2}$$

($@^{-1}$: "Inverse sign of context" operation)

ATAPP-EMPTY

$$\frac{}{@^{-1} \emptyset = \emptyset}$$

ATAPP-REC

$$\frac{}{@^{-1}(\{ @h :_m ^n[A] \} \sqcup \Gamma) = \{ h :^{m \cdot n} A \} \sqcup @^{-1} \Gamma}$$

$$\boxed{\Delta \Vdash e}$$

(Typing of effects (require both positive and negative contexts))

TYEFF-NOEFF

$$\frac{}{\emptyset \sqcup \emptyset \Vdash \varepsilon}$$

TYEFF-SINGLE

$$\frac{\Gamma \sqcup H \Vdash \bar{v} : A \quad h \notin \text{names}(\Gamma)}{m \cdot ((n \cdot \uparrow) \cdot \Gamma \sqcup \{ @h :_\nu ^n[A] \} \sqcup ^n \cdot H) \Vdash h := \bar{v}}$$

TYEFF-UNION

$$\frac{\begin{array}{c} \Gamma_1 \sqcup H_1 \sqcup @^{-1} \Gamma_{22} \Vdash e_1 \\ \Gamma_{21} \sqcup \Gamma_{22} \sqcup H_2 \Vdash e_2 \\ \text{names}(\Gamma_1 \sqcup H_1) \cap \text{names}(\Gamma_{21} \sqcup H_2) = \emptyset \end{array}}{\Gamma_1 \sqcup \Gamma_{21} \sqcup H_1 \sqcup H_2 \Vdash e_1 \cdot e_2}$$

$$\boxed{\Gamma \vdash v \mid e : A}$$

(Typing of commands (only a positive context is needed))

TYCMD-CMD

$$\frac{\begin{array}{c} \Gamma_{11} \sqcup \Gamma_{12} \vdash v : A \\ \Gamma_2 \sqcup @^{-1} \Gamma_{12} \Vdash e \\ \text{names}(\Gamma_{11}) \cap \text{names}(\Gamma_2) = \emptyset \end{array}}{\Gamma_{11} \sqcup \Gamma_2 \vdash v \mid e : A}$$

$$\Delta \Vdash \bar{v} : A$$

(Typing of extended values (require both positive and negative contexts))

$$\text{TYVALEXT-HOLE} \quad \frac{}{\emptyset \sqcup \{\mathbf{h} : \nu A\} \Vdash \mathbf{h} : A}$$

$$\text{TYVALEXT-DEST} \quad \frac{}{\{\mathbf{@h} : \nu \bar{n}[A]\} \sqcup \emptyset \Vdash \mathbf{@h} : \bar{n}[A]}$$

$$\text{TYVALEXT-UNIT} \quad \frac{}{\emptyset \sqcup \emptyset \Vdash () : \mathbf{1}}$$

$$\text{TYVALEXT-INL} \quad \frac{\Gamma \sqcup H \Vdash \bar{v} : A_1}{\Gamma \sqcup H \Vdash \text{Inl } \bar{v} : A_1 \oplus A_2}$$

$$\text{TYVALEXT-INR} \quad \frac{\Gamma \sqcup H \Vdash \bar{v} : A_2}{\Gamma \sqcup H \Vdash \text{Inr } \bar{v} : A_1 \oplus A_2}$$

$$\text{TYVALEXT-PROD} \quad \frac{\begin{array}{l} \Gamma_1 \sqcup H_1 \Vdash \bar{v}_1 : A_1 \\ \Gamma_2 \sqcup H_2 \Vdash \bar{v}_2 : A_2 \\ \text{names}(\Gamma_1 \sqcup H_1) \cap \text{names}(\Gamma_2 \sqcup H_2) = \emptyset \end{array}}{\Gamma_1 \sqcup \Gamma_2 \sqcup H_1 \sqcup H_2 \Vdash (\bar{v}_1, \bar{v}_2) : A_1 \otimes A_2}$$

$$\text{TYVALEXT-EXP} \quad \frac{\Gamma \sqcup H \Vdash \bar{v} : A}{m \cdot \Gamma \sqcup m \cdot H \Vdash \bar{v} : !^m A}$$

$$\text{TYVALEXT-AMPAR} \quad \frac{\begin{array}{l} \Gamma_1 \sqcup \emptyset \Vdash v_1 : A_1 \\ \Gamma_2 \sqcup @^{-1}\Gamma_1 \Vdash \bar{v}_2 : A_2 \end{array}}{\Gamma_2 \sqcup \emptyset \Vdash \langle v_1, \bar{v}_2 \rangle_H : A_1 \ltimes A_2}$$

$$\text{TYVALEXT-LAMBDA} \quad \frac{\Gamma \sqcup \{x : m A_1\} \vdash t : A_2}{\Gamma \sqcup \emptyset \Vdash \lambda x. t : A_1 \multimap A_2}$$

$$\Gamma \vdash t : A$$

(Typing of terms (only a positive context is needed))

$$\text{TYTERM-VAL} \quad \frac{\Gamma \sqcup \emptyset \Vdash v : A}{\Gamma \vdash v : A}$$

$$\text{TYTERM-VARNOW} \quad \frac{}{\{x : \nu A\} \vdash x : A}$$

$$\text{TYTERM-VARINF} \quad \frac{}{\{x : \infty A\} \vdash x : A}$$

$$\text{TYTERM-APP} \quad \frac{\Gamma_1 \vdash t : A_1 \quad \Gamma_2 \vdash u : A_1 \multimap A_2 \quad \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset}{m \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ u : A_2}$$

$$\text{TYTERM-PATUNIT} \quad \frac{\Gamma_1 \vdash t : \mathbf{1} \quad \Gamma_2 \vdash u : B \quad \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset}{\Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case } () \mapsto u : B}$$

$$\text{TYTERM-PATSUM} \quad \frac{\begin{array}{l} \Gamma_1 \vdash t : A_1 \oplus A_2 \\ \Gamma_2 \sqcup \{x_1 : m A_1\} \vdash u_1 : B \\ \Gamma_2 \sqcup \{x_2 : m A_2\} \vdash u_2 : B \\ \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset \end{array}}{m \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : B}$$

$$\text{TYTERM-PATPROD} \quad \frac{\begin{array}{l} \Gamma_1 \vdash t : A_1 \otimes A_2 \\ \Gamma_2 \sqcup \{x_1 : m A_1, x_2 : m A_2\} \vdash u : B \\ \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset \end{array}}{m \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case } (x_1, x_2) \mapsto u : B}$$

$$\text{TYTERM-PATEXP} \quad \frac{\begin{array}{l} \Gamma_1 \vdash t : !^{m'} A \\ \Gamma_2 \sqcup \{x : m \cdot m' A_1\} \vdash u : B \\ \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset \end{array}}{m \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case } !^{m'} x \mapsto u : B}$$

$$\text{TYTERM-MAPAMPAR} \quad \frac{\begin{array}{l} \Gamma_1 \vdash t : A_1 \ltimes A_2 \\ \uparrow \Gamma_2 \sqcup \{x : \nu A_1\} \vdash u : B \\ \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset \end{array}}{\Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{mapL } x \mapsto u : B \ltimes A_2}$$

$$\text{TYTERM-FILLCOMP} \quad \frac{\Gamma_1 \vdash t : \bar{n}[A_2] \quad \Gamma_2 \vdash u : A_1 \ltimes A_2 \quad \text{names}(\Gamma_1) \cap \text{names}(\Gamma_2) = \emptyset}{\Gamma_1 \sqcup (\bar{n} \cdot \uparrow) \cdot \Gamma_2 \vdash t \triangleleft \bullet u : A_1}$$

$$\text{TYTERM-FILLUNIT} \quad \frac{}{\Gamma \vdash t \triangleleft () : \mathbf{1}}$$

$$\text{TYTERM-FILLINL} \quad \frac{}{\Gamma \vdash t \triangleleft \text{Inl} : \bar{n}[A_1]}$$

$$\text{TYTERM-FILLINR} \quad \frac{}{\Gamma \vdash t \triangleleft \text{Inr} : \bar{n}[A_2]}$$

$$\text{TYTERM-FILLPROD} \quad \frac{\Gamma \vdash t : \bar{n}[A_1 \otimes A_2]}{\Gamma \vdash t \triangleleft (,) : \bar{n}[A_1] \otimes \bar{n}[A_2]}$$

$$\text{TYTERM-FILLEXP} \quad \frac{\Gamma \vdash t : \bar{n}[!^{n'} A]}{\Gamma \vdash t \triangleleft !^{n'} : \bar{n} \cdot n' [A]}$$

$$\text{TYTERM-ALLOC} \quad \frac{}{\emptyset \vdash \text{alloc}_A : \nu[A] \ltimes A}$$

$$\text{TYTERM-TOAMPAR} \quad \frac{\Gamma \vdash t : A}{\Gamma \vdash \text{to}_{\ltimes} t : \mathbf{1} \ltimes A}$$

$$\text{TYTERM-FROMAMPAR} \quad \frac{\Gamma \vdash t : \mathbf{1} \ltimes A}{\Gamma \vdash \text{from}_{\ltimes} t : A}$$

3 Effects and big-step semantics

effect, e	$::=$	Effect
	ε	No effect
	$\mathbf{h} := \bar{v}$	
	$e_1 \cdot e_2$	

$\text{eff_app}_1 = \text{eff_app}_2$

(**apply**: how effects are applied locally or winded up (we assume effect lists are ε -terminated))

$$\begin{array}{c}
 \text{EFFAPP-NOEFF} \\
 \hline
 \text{apply}(\varepsilon, \bar{v}_H) = \varepsilon, \bar{v}_H \\
 \\
 \text{EFFAPP-WINDUP} \\
 \hline
 \mathbf{h} \notin \text{names}(H) \\
 \hline
 \text{apply}(\mathbf{h} := \bar{v}_2 \cdot e, \bar{v}_1 H) = \mathbf{h} := \bar{v}_2 \hat{\cdot} \text{apply}(e, \bar{v}_1 H) \\
 \\
 \text{EFFAPP-FILL} \\
 \hline
 \frac{_ \sqcup H' \Vdash \bar{v}_2 : \mathbf{A} \quad \text{names}(H \sqcup \{\mathbf{h} : \bar{n} \mathbf{A}\}) \cap \text{names}(H') = \emptyset}{\text{apply}(\mathbf{h} := \bar{v}_2 \cdot e, \bar{v}_1 H \sqcup \{\mathbf{h} : \bar{n} \mathbf{A}\}) = \text{apply}(e, \bar{v}_1 [\mathbf{h} := \bar{v}_2] H \sqcup \bar{n} \cdot H')}
 \end{array}$$

$t \Downarrow v \mid e$

(Big-step evaluation into commands)

$$\begin{array}{c}
 \text{BIGSTEP-VAL} \\
 \hline
 v \Downarrow v \mid \varepsilon \\
 \\
 \text{BIGSTEP-APP} \\
 \hline
 \frac{t_1 \Downarrow v_1 \mid e_1 \quad t_2 \Downarrow \lambda x. u \mid e_2 \quad u[x := v_1] \Downarrow v_3 \mid e_3}{t_1 \succ t_2 \Downarrow v_3 \mid e_1 \cdot e_2 \cdot e_3} \\
 \\
 \text{BIGSTEP-PATUNIT} \\
 \hline
 \frac{t_1 \Downarrow () \mid e_1 \quad t_2 \Downarrow v_2 \mid e_2}{t_1 \succ \text{case}() \mapsto t_2 \Downarrow v_2 \mid e_1 \cdot e_2} \\
 \\
 \text{BIGSTEP-PATINL} \\
 \hline
 \frac{t \Downarrow \text{Inl } v_1 \mid e_1 \quad u_1[x_1 := v_1] \Downarrow v_2 \mid e_2}{t \succ \text{case} \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \} \Downarrow v_2 \mid e_1 \cdot e_2} \\
 \\
 \text{BIGSTEP-PATINR} \\
 \hline
 \frac{t \Downarrow \text{Inr } v_1 \mid e_1 \quad u_2[x_2 := v_1] \Downarrow v_2 \mid e_2}{t \succ \text{case} \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \} \Downarrow v_2 \mid e_1 \cdot e_2} \\
 \\
 \text{BIGSTEP-MAPAMPAR} \\
 \hline
 \frac{t \Downarrow \langle v_1, \bar{v}_2 \rangle_H \mid e_1 \quad u[x := v_1] \Downarrow v_3 \mid e_2 \quad e_3, \bar{v}_4 H' = \text{apply}(e_2, \bar{v}_2 H)}{t \succ \text{mapL } x \mapsto u \Downarrow \langle v_3, \bar{v}_4 \rangle_{H'} \mid e_1 \cdot e_3} \\
 \\
 \text{BIGSTEP-PATPROD} \\
 \hline
 \frac{t \Downarrow (v_1, v_2) \mid e_1 \quad u[x_1 := v_1, x_2 := v_2] \Downarrow v_2 \mid e_2}{t \succ \text{case}(x_1, x_2) \mapsto u \Downarrow v_2 \mid e_1 \cdot e_2} \\
 \\
 \text{BIGSTEP-ALLOC} \\
 \hline
 \frac{\text{fresh } \mathbf{h}}{\text{alloc}_{\mathbf{A}} \Downarrow \langle @\mathbf{h}, \mathbf{h} \rangle_{\{\mathbf{h} : \bar{\nu} \mathbf{A}\}} \mid \varepsilon} \\
 \\
 \text{BIGSTEP-TOAMPAR} \\
 \hline
 \frac{t \Downarrow v \mid e}{\text{to}_{\mathbf{x}} t \Downarrow \langle (), v \rangle_{\emptyset} \mid e} \\
 \\
 \text{BIGSTEP-FROMAMPAR} \\
 \hline
 \frac{t \Downarrow \langle (), v \rangle_{\emptyset} \mid e}{\text{from}_{\mathbf{x}} t \Downarrow v \mid e} \\
 \\
 \text{BIGSTEP-FILLUNIT} \\
 \hline
 \frac{t \Downarrow @\mathbf{h} \mid e}{t \triangleleft () \Downarrow () \mid e \cdot \mathbf{h} := ()} \\
 \\
 \text{BIGSTEP-FILLINL} \\
 \hline
 \frac{t \Downarrow @\mathbf{h} \mid e \quad \text{fresh } \mathbf{h}'}{t \triangleleft \text{Inl} \Downarrow @\mathbf{h}' \mid e \cdot \mathbf{h} := \text{Inl } \mathbf{h}'} \\
 \\
 \text{BIGSTEP-FILLINR} \\
 \hline
 \frac{t \Downarrow @\mathbf{h} \mid e}{t \triangleleft \text{Inr} \Downarrow @\mathbf{h}' \mid e \cdot \mathbf{h} := \text{Inr } \mathbf{h}'} \\
 \\
 \text{BIGSTEP-FILLPROD} \\
 \hline
 \frac{t \Downarrow @\mathbf{h} \mid e \quad \text{fresh } \mathbf{h}_1 \quad \text{fresh } \mathbf{h}_2}{t \triangleleft (,) \Downarrow (@\mathbf{h}_1, @\mathbf{h}_2) \mid e \cdot \mathbf{h} := (\mathbf{h}_1, \mathbf{h}_2)} \\
 \\
 \text{BIGSTEP-FILLCOMP} \\
 \hline
 \frac{t \Downarrow @\mathbf{h} \mid e_1 \quad u \Downarrow \langle v_1, \bar{v}_2 \rangle_H \mid e_2}{t \triangleleft \bullet u \Downarrow v_1 \mid e_1 \cdot e_2 \cdot \mathbf{h} := \bar{v}_2}
 \end{array}$$