

metavariable, x , xs , y , uf , f , d

term, t, u	$ \begin{array}{l} ::= \\ \textcolor{violet}{x} \\ v \\ t \ u \\ t \ ; \ u \\ \text{case } t \text{ of } \{ () \mapsto u \} \\ \text{case } t \text{ of } \{ 1.\textcolor{violet}{x}_1 \mapsto u_1, 2.\textcolor{violet}{x}_2 \mapsto u_2 \} \\ \text{case } t \text{ of } \{ \langle \textcolor{violet}{x}_1, \textcolor{violet}{x}_2 \rangle \mapsto u \} \\ \text{case } t \text{ of } \{ \text{Ur } \textcolor{violet}{x} \mapsto u \} \\ \text{case } t \text{ of } \{ \text{roll } \textcolor{violet}{x} \mapsto u \} \\ \text{alloc } \textcolor{blue}{d} . t \\ t \triangleleft () \\ t \triangleleft u \\ t \triangleleft 1.\textcolor{violet}{d} . u \\ t \triangleleft 2.\textcolor{violet}{d} . u \\ t \triangleleft \langle \textcolor{violet}{d}_1, \textcolor{violet}{d}_2 \rangle . u \\ t \triangleleft \text{Ur } \textcolor{violet}{d} . u \\ t \triangleleft \text{roll } \textcolor{violet}{d} . u \\ (t) \\ t[\text{var_subs}] \end{array} $	<p>term</p> <ul style="list-style-type: none"> variable value application effect execution pattern-matching on unit pattern-matching on sum pattern-matching on product pattern-matching on exponentiated value unroll for recursive types allocate data fill destination with unit fill terminal-type destination fill sum-type destination with variant 1 fill sum-type destination with variant 2 fill product-type destination fill destination with exponential fill destination with recursive type <p>S M</p>
var_sub, vs	$ \begin{array}{l} ::= \\ \textcolor{violet}{x} := t \end{array} $	variable substitution
var_subs	$ \begin{array}{l} ::= \\ \text{vs} \\ \text{vs}, \text{var_subs} \end{array} $	variable substitutions
heap_val, h	$ \begin{array}{l} ::= \\ () \\ 1.\textcolor{brown}{l} \\ 2.\textcolor{brown}{l} \\ \langle \textcolor{brown}{l}_1, \textcolor{brown}{l}_2 \rangle \\ \text{Ur } \textcolor{brown}{l} \\ \text{roll } \textcolor{brown}{l} \\ \textcolor{blue}{C} \textcolor{brown}{l} \end{array} $	<p>M generic for all the cases above</p>
val, v	$ \begin{array}{l} ::= \\ \bullet \\ \textcolor{blue}{l} \\ \lambda \textcolor{violet}{x} : \textcolor{blue}{A} . t \\ h \end{array} $	<p>unreducible value</p> <ul style="list-style-type: none"> no-effect effect address of an allocated memory area lambda abstraction heap value
<i>label</i> , <i>l</i>	$::= $	memory address
label_stmt, s	$ \begin{array}{l} ::= \\ \textcolor{brown}{l} \triangleleft v \\ \textcolor{brown}{l} \triangleleft \emptyset \\ \textcolor{brown}{l} \triangleleft \bar{v} \end{array} $	<p>label statement</p> <p>M generic for multiple occurrences</p>

label_stmts	$::=$ $\begin{array}{ l} s \\ s, \text{label_stmts} \end{array}$	label statements
heap_context, \mathbb{H}	$::=$ $\begin{array}{ l} \emptyset \\ \{\text{label_stmts}\} \\ \mathbb{H}_1 \sqcup \mathbb{H}_2 \end{array}$	label statements
type, A, B	$::=$ $\begin{array}{ l} \perp \\ 1 \\ R \\ A \otimes B \\ A \oplus B \\ A \multimap B \\ [A] \\ !A \\ (A) \\ W[r := A] \end{array}$	bottom type unit type recursive type bound to a name product type sum type linear function type destination type exponential S M
type_with_hole, W	$::=$ $\begin{array}{ l} \perp \\ r \\ 1 \\ R \\ W_1 \otimes W_2 \\ W_1 \oplus W_2 \\ W_1 \multimap W_2 \\ [W] \\ !W \\ (W) \end{array}$	bottom type type hole in recursive definition unit type recursive type bound to a name product type sum type linear function type destination type exponential S
rec_type_bound, R	$::=$	recursive type bound to a name
rec_type_def	$::=$ $\begin{array}{ l} \mu r. W \end{array}$	
type_affect, ta	$::=$ $\begin{array}{ l} x : A \\ l : A \\ \bar{l} : \bar{A} \end{array}$	type affectation var label generic for multiple occurrences
type_affects	$::=$ $\begin{array}{ l} \text{ta} \\ \text{ta}, \text{type_affects} \end{array}$	type affectations
typing_context, $\Gamma, \Delta, \mathcal{U}, \Phi$	$::=$ $\begin{array}{ l} \emptyset \end{array}$	typing context

		{type_affects}	
		$\Gamma \sqcup \Delta$	
types, \bar{A}	::=		empty type list
		.	
		A	
		A types	
heap_constructor, C	::=		
		{() }	
		{1. }	
		{2. }	
		{⟨,⟩ }	
		{Ur }	
		{roll R }	
judg	::=		
		$\Phi \vdash \mathbb{H}$	
		$\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A$	
		$\Phi ; \mathcal{U} ; \Gamma \vdash t : A$	
		$C : \bar{A} \hookrightarrow A$	
		$A = B$	
		$t = u$	
		$\Gamma = \Delta$	
		$l \notin \text{names}(\Phi)$	
		$\text{type_affect} \in \Gamma$	
		$\text{label_stmt} \in \mathbb{H}$	
		$\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset$	
terminals	::=		
		()	
		\mapsto	
		\star	
		\otimes	
		\oplus	
		\multimap	
		$:=$	
		\vdash	
		\sqcup	
		;	
		\cap	
		\emptyset	
		\rightarrow	
		\triangleright	
		\neq	
		\in	
		\notin	
		$\backslash n$	
		\langle	
		\rangle	

	$ \begin{array}{ l} 1. \\ 2. \\ \mathbf{U_r} \\ \triangleleft \\ \\ \emptyset \\ \textcolor{blue}{\hookrightarrow} \\ = \\ \Downarrow \\ \dots \\ \text{fix} \\ \text{---} \\ \textcolor{blue}{\perp} \\ \bullet \end{array} $	
formula	$ \begin{array}{ l} ::= \\ \text{judgement} \end{array} $	
Ctx	$ \begin{array}{ l} ::= \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\ \textcolor{red}{l} \notin \text{names}(\Phi) \\ \text{type_affect} \in \Gamma \end{array} $	Γ and Δ are disjoint typing contexts with no clashing
Heap	$ \begin{array}{ l} ::= \\ \text{label_stmt} \in \mathbb{H} \end{array} $	
Eq	$ \begin{array}{ l} ::= \\ \textcolor{blue}{A} = \textcolor{blue}{B} \\ t = u \\ \Gamma = \Delta \end{array} $	
Ty	$ \begin{array}{ l} ::= \\ \textcolor{blue}{R} \stackrel{\text{fix}}{=} \text{rec_type_def} \\ \Phi \vdash \mathbb{H} \\ \Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : \textcolor{blue}{A} \\ \Phi ; \mathcal{U} ; \Gamma \vdash t : \textcolor{blue}{A} \\ C : \bar{\textcolor{blue}{A}} \textcolor{blue}{\hookrightarrow} \textcolor{blue}{A} \end{array} $	\mathbb{H} is a well-typed heap given heap typing context Φ t is a well-typed term of type $\textcolor{blue}{A}$ given heap typing con Heap constructor C builds a value of type $\textcolor{blue}{A}$ given arg
Sem	$ \begin{array}{ l} ::= \\ \mathbb{H} \mid t \Downarrow \mathbb{H}' \mid t' \end{array} $	t reduces to t' , with heap growing from \mathbb{H} to \mathbb{H}'
judgement	$ \begin{array}{ l} ::= \\ \text{Ctx} \\ \text{Heap} \\ \text{Eq} \\ \text{Ty} \\ \text{Sem} \end{array} $	
user_syntax	$ \begin{array}{ l} ::= \\ \text{metavariable} \end{array} $	

- | term
- | var_sub
- | var_subs
- | heap_val
- | val
- | *label*
- | label_stmt
- | label_stmts
- | heap_context
- | type
- | type_with_hole
- | rec_type_bound
- | rec_type_def
- | type_affect
- | type_affects
- | typing_context
- | types
- | heap_constructor
- | judg
- | terminals

$$\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset$$

Γ and Δ are disjoint typing contexts with no clashing variable names or labels

$$l \notin \text{names}(\Phi)$$

$$\text{type_affect} \in \Gamma$$

$$\text{label_stmt} \in \mathbb{H}$$

$$A = B$$

$$t = u$$

$$\Gamma = \Delta$$

$$R \stackrel{\text{fix}}{=} \text{rec_type_def}$$

$$\Phi \vdash \mathbb{H} \quad \mathbb{H} \text{ is a well-typed heap given heap typing context } \Phi$$

$$\overline{\emptyset \vdash \emptyset} \quad \text{TYHEAP_EMPTY}$$

$$\frac{\Phi \vdash \mathbb{H}}{\Phi \sqcup \{l : \perp\} \vdash \mathbb{H} \sqcup \{l \triangleleft \bullet\}} \quad \text{TYHEAP_NOEFF}$$

$$\frac{\Phi \vdash \mathbb{H} \quad l' \notin \text{names}(\Phi)}{\Phi \sqcup \{l : [A]\} \vdash \mathbb{H} \sqcup \{l \triangleleft [A]\}} \quad \text{TYHEAP_LDEST}$$

$$\frac{\Phi \vdash \mathbb{H} \quad \Phi ; \mathcal{U} ; \Gamma \vdash \lambda x:A. t : A \multimap B}{\Phi \sqcup \{l : A \multimap B\} \vdash \mathbb{H} \sqcup \{l \triangleleft \lambda x:A. t\}} \quad \text{TYHEAP_LAM}$$

$$\frac{\Phi \sqcup \{\bar{l} : \bar{A}\} \vdash \mathbb{H} \sqcup \{\bar{l} \triangleleft \bar{v}\} \quad C : \bar{A} \stackrel{\leq}{\hookrightarrow} A}{\Phi \sqcup \{l : A, \bar{l} : \bar{A}\} \vdash \mathbb{H} \sqcup \{l \triangleleft C\bar{l}, \bar{l} \triangleleft \bar{v}\}} \quad \text{TYHEAP_HEAPVAL}$$

$$\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} \mid t : A$$

$$\frac{\Phi \vdash \mathbb{H} \quad \Phi ; \mathcal{U} ; \Gamma \vdash t : A}{\Phi ; \mathcal{U} ; \Gamma \vdash \mathbb{H} | t : A} \text{TYCOMMAND_DEF}$$

$\Phi ; \mathcal{U} ; \Gamma \vdash t : A$ t is a well-typed term of type A given heap typing context Φ , unrestricted typing context

$$\frac{}{\Phi ; \mathcal{U} ; \emptyset \vdash \bullet : \perp} \text{TYTERM_NOEFF}$$

$$\frac{l \notin \text{names}(\Phi)}{\Phi ; \mathcal{U} ; \emptyset \vdash \llbracket A \rrbracket : \llbracket A \rrbracket} \text{TYTERM_LDEST}$$

$$\frac{\Phi ; \mathcal{U} ; \Gamma \sqcup \{x : A\} \vdash t : B}{\Phi ; \mathcal{U} ; \Gamma \vdash \lambda x : A. t : A \multimap B} \text{TYTERM_LAM}$$

$$\frac{C : \bar{A} \sqsubseteq A}{\Phi \sqcup \{\bar{l} : \bar{A}\} ; \mathcal{U} ; \emptyset \vdash C \bar{l} : A} \text{TYTERM_HEAPVAL}$$

$$\frac{}{\Phi ; \mathcal{U} ; \{x : A\} \vdash x : A} \text{TYTERM_ID}$$

$$\frac{}{\Phi ; \mathcal{U} \sqcup \{x : A\} ; \emptyset \vdash x : A} \text{TYTERM_ID'}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : A \multimap B \\ \Phi ; \mathcal{U} ; \Delta \vdash u : A \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t u : B} \text{TYTERM_APP}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : \perp \\ \Phi ; \mathcal{U} ; \Delta \vdash u : B \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t ; u : B} \text{TYTERM_Witheff}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : 1 \\ \Phi ; \mathcal{U} ; \Delta \vdash u : A \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{() \mapsto u\} : A} \text{TYTERM_PATU}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : A_1 \oplus A_2 \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{x_1 : A_1\} \vdash u_1 : B \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{x_2 : A_2\} \vdash u_2 : B \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{1.x_1 \mapsto u_1, 2.x_2 \mapsto u_2\} : B} \text{TYTERM_PATs}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : A_1 \otimes A_2 \\ \Phi ; \mathcal{U} ; \Delta \sqcup \{x_1 : A_1, x_2 : A_2\} \vdash u : B \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : B} \text{TYTERM_PATP}$$

$$\frac{\begin{array}{l} \Phi ; \mathcal{U} ; \Gamma \vdash t : !A \\ \Phi ; \mathcal{U} \sqcup \{x : A\} ; \Delta \vdash u : B \\ \text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \end{array}}{\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{ \text{Ur } x \mapsto u \} : B} \text{TYTERM_PATE}$$

$$\begin{array}{c}
R \stackrel{\text{fix}}{=} \mu r. W \\
\Phi ; \mathcal{U} ; \Gamma \vdash t : R \\
\Phi ; \mathcal{U} ; \Delta \sqcup \{x : W[r := R]\} \vdash u : B \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash \text{case } t \text{ of } \{\text{roll } x \mapsto u\} : B \quad \text{TYTERM_PATR}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \sqcup \{d : [A]\} \vdash t : \perp \\
\hline
\Phi ; \mathcal{U} ; \Gamma \vdash \text{alloc } d . t : A \quad \text{TYTERM_ALLOC}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [1] \\
\hline
\Phi ; \mathcal{U} ; \Gamma \vdash t \triangleleft () : \perp \quad \text{TYTERM_FILLU}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [A] \\
\Phi ; \mathcal{U} ; \Delta \vdash u : A \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft u : \perp \quad \text{TYTERM_FILLL}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [A_1 \oplus A_2] \\
\Phi ; \mathcal{U} ; \Delta \sqcup \{d' : [A_1]\} \vdash u : B \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft 1.d'.u : B \quad \text{TYTERM_FILLV1}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [A_1 \oplus A_2] \\
\Phi ; \mathcal{U} ; \Delta \sqcup \{d' : [A_2]\} \vdash u : B \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \vdash t \triangleleft 2.d'.u : B \quad \text{TYTERM_FILLV2}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [A_1 \otimes A_2] \\
\Phi ; \mathcal{U} ; \Delta \sqcup \{d_1 : [A_1], d_2 : [A_2]\} \vdash u : B \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft \langle d_1, d_2 \rangle . u : B \quad \text{TYTERM_FILLP}
\end{array}$$

$$\begin{array}{c}
\Phi ; \mathcal{U} ; \Gamma \vdash t : [!A] \\
\Phi ; \mathcal{U} ; \emptyset \sqcup \{d : [A]\} \vdash u : B \\
\hline
\Phi ; \mathcal{U} ; \Gamma \vdash t \triangleleft \text{Ur } d . u : B \quad \text{TYTERM_FILLE}
\end{array}$$

$$\begin{array}{c}
R \stackrel{\text{fix}}{=} \mu r. W \\
\Phi ; \mathcal{U} ; \Gamma \vdash t : [R] \\
\Phi ; \mathcal{U} ; \Delta \sqcup \{d : [W[r := R]]\} \vdash u : B \\
\text{names}(\Gamma) \cap \text{names}(\Delta) = \emptyset \\
\hline
\Phi ; \mathcal{U} ; \Gamma \sqcup \Delta \vdash t \triangleleft_{\text{R}} \text{roll } d . u : B \quad \text{TYTERM_FILLR}
\end{array}$$

$C : \bar{A} \xrightarrow{\text{c}} A$

Heap constructor C builds a value of type A given arguments of type \bar{A}

$$\frac{}{\{()\} : \cdot \xrightarrow{\text{c}} 1} \quad \text{TYCTOR_U}$$

$$\frac{}{\{1.\} : A \xrightarrow{\text{c}} A \oplus B} \quad \text{TYCTOR_V1}$$

$$\frac{}{\{2.\} : B \xrightarrow{\text{c}} A \oplus B} \quad \text{TYCTOR_V2}$$

$$\frac{}{\{\langle, \rangle\} : A \ B \xrightarrow{\text{c}} A \otimes B} \quad \text{TYCTOR_P}$$

$$\frac{}{\{ \text{Ur} \} : A \xrightarrow{\text{blue}} !A} \text{TyCTOR_E}$$

$$\frac{R \stackrel{\text{fix}}{=} \mu \text{r} . W}{\{ \text{roll } R \} : W[r := R] \xrightarrow{\text{blue}} R} \text{TyCTOR_R}$$

$$\boxed{\mathbb{H} | t \Downarrow \mathbb{H}' | t'} \quad t \text{ reduces to } t', \text{ with heap growing from } \mathbb{H} \text{ to } \mathbb{H}'$$

$$\frac{}{\mathbb{H} | \bullet \Downarrow \mathbb{H} | \bullet} \text{SEMOP_NOEFF (value)}$$

$$\frac{}{\mathbb{H} | \text{[A]} \Downarrow \mathbb{H} | \text{[A]}} \text{SEMOP_LDEST (value)}$$

$$\frac{}{\mathbb{H} | \lambda \text{x:A} . t \Downarrow \mathbb{H} | \lambda \text{x:A} . t} \text{SEMOP_LAM (value)}$$

$$\frac{}{\mathbb{H} | \text{C}\bar{l} \Downarrow \mathbb{H} | \text{C}\bar{l}} \text{SEMOP_HEAPVAL (value)}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \lambda \text{x:A} . t' \\ \mathbb{H}_1 | u \Downarrow \mathbb{H}_2 | v_2 \\ \mathbb{H}_2 | t'[\text{x} := v_2] \Downarrow \mathbb{H}_3 | v_3 \end{array}}{\mathbb{H}_0 | t u \Downarrow \mathbb{H}_3 | v_3} \text{SEMOP_APP}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | () \\ \mathbb{H}_1 | u \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ () \mapsto u \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATU}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | 1.l \\ \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | u_1[\text{x}_1 := v] \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ 1.\text{x}_1 \mapsto u_1, 2.\text{x}_2 \mapsto u_2 \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATV1}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | 2.l \\ \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | u_2[\text{x}_2 := v] \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ 1.\text{x}_1 \mapsto u_1, 2.\text{x}_2 \mapsto u_2 \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATV2}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 \sqcup \{ l_1 \triangleleft v_{11}, l_2 \triangleleft v_{12} \} | \langle l_1, l_2 \rangle \\ \mathbb{H}_1 \sqcup \{ l_1 \triangleleft v_{11}, l_2 \triangleleft v_{12} \} | u[\text{x}_1 := v_1, \text{x}_2 := v_2] \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ \langle \text{x}_1, \text{x}_2 \rangle \mapsto u \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATP}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | \text{Ur } l \\ \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | u[\text{x} := v_1] \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ \text{Ur } \text{x} \mapsto u \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATE}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | \text{roll } l \\ \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | u[\text{x} := v_1] \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | \text{case } t \text{ of } \{ \text{roll } \text{x} \mapsto u \} \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_PATR}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \bullet \\ \mathbb{H}_1 | u \Downarrow \mathbb{H}_2 | v_2 \end{array}}{\mathbb{H}_0 | t ; u \Downarrow \mathbb{H}_2 | v_2} \text{SEMOP_Witheff}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t[\text{d} := \text{[A]}] \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | \bullet \\ \mathbb{H}_0 | \text{alloc } \text{d} . t \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft v_1 \} | v_1 \end{array}}{\text{SEMOP_ALLOC}}$$

$$\frac{\begin{array}{c} \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \text{[1]} \\ \mathbb{H}_0 | t \triangleleft () \Downarrow \mathbb{H}_1 \sqcup \{ l \triangleleft () \} | \bullet \end{array}}{\text{SEMOP_FILLU}}$$

$$\begin{array}{c}
\frac{\mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!A} \quad \mathbb{H}_1 | u \Downarrow \mathbb{H}_2 | v_2}{\mathbb{H}_0 | t \triangleleft u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft v_2\} | \bullet} \text{SEMOp_FILL} \\
\\
\frac{\mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!A \oplus B} \quad \mathbb{H}_1 | u[d := \underline{!A}] \Downarrow \mathbb{H}_2 \sqcup \{l' \triangleleft v_1\} | v_2}{\mathbb{H}_0 | t \triangleleft 1.d.u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft 1.l', l' \triangleleft v_1\} | v_2} \text{SEMOp_FILLV1} \\
\\
\frac{\mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!A \oplus B} \quad \mathbb{H}_1 | u[d := \underline{!B}] \Downarrow \mathbb{H}_2 \sqcup \{l' \triangleleft v_1\} | v_2}{\mathbb{H}_0 | t \triangleleft 2.d.u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft 2.l', l' \triangleleft v_1\} | v_2} \text{SEMOp_FILLV2} \\
\\
\frac{\mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!A \otimes B} \quad \mathbb{H}_1 | u[d_1 := \underline{!A}, d_2 := \underline{!B}] \Downarrow \mathbb{H}_2 \sqcup \{l_1 \triangleleft v_{11}, l_2 \triangleleft v_{12}\} | v_2}{\mathbb{H}_0 | t \triangleleft \langle d_1, d_2 \rangle . u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft \langle l_1, l_2 \rangle, l_1 \triangleleft v_{11}, l_2 \triangleleft v_{12}\} | v_2} \text{SEMOp_FILLP} \\
\\
\frac{\mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!A} \quad \mathbb{H}_1 | u[d := \underline{!A}] \Downarrow \mathbb{H}_2 \sqcup \{l' \triangleleft v_1\} | v_2}{\mathbb{H}_0 | t \triangleleft \text{Ur } d.u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft \text{Ur } l', l' \triangleleft v_1\} | v_2} \text{SEMOp_FILLE} \\
\\
\frac{\begin{array}{c} R \stackrel{\text{fix}}{=} \mu r. W \\ \mathbb{H}_0 | t \Downarrow \mathbb{H}_1 | \underline{!R} \\ \mathbb{H}_1 | u[d := \underline{!W[r := R]}] \Downarrow \mathbb{H}_2 \sqcup \{l' \triangleleft v_1\} | v_2 \end{array}}{\mathbb{H}_0 | t \triangleleft \text{roll}_R d.u \Downarrow \mathbb{H}_2 \sqcup \{l \triangleleft \text{roll}_R l', l' \triangleleft v_1\} | v_2} \text{SEMOp_FILLR}
\end{array}$$

Definition rules: 53 good 0 bad

Definition rule clauses: 141 good 0 bad