

metavariable, x, y

term, t, u	$::=$	
		v
		x
		$t\ u$
		$t \ ;\ u$
		$\text{case } t \text{ of } \{\star \mapsto u\}$
		$\text{case } t \text{ of } \{\text{Ur } x \mapsto u\}$
		$\text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\}$
		$\text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}$
		$\text{case } t \text{ of } \{\text{@R } x \mapsto u\}$
		$\text{alloc } x . t$
		$t \triangleleft^p \star$
		$t \triangleleft^p \lambda x : A . u$
		$t \triangleleft^p u$
		$t \triangleleft^p \text{Ur } y . u$
		$t \triangleleft^p \text{Inl } y . u$
		$t \triangleleft^p \text{Inr } y . u$
		$t \triangleleft^p \langle y_1, y_2 \rangle . u$
		$t \triangleleft^p \text{@R } y . u$
		(t)
		$t[\text{subs}]$

term

value
variable
application
effect sequencing
pattern-matching on unit
pattern-matching on exponentiated value
pattern-matching on sum
pattern-matching on product
unroll for recursive types
allocate data
fill destination with unit
fill destination with function

fill destination with exponential
fill sum-type destination with variant 1
fill sum-type destination with variant 2
fill product-type destination
fill destination with recursive type

S
M

hole, h

$::=$

val, v

$::=$

| \bullet
| d

unreducible value
no-effect effect
data structure

data, d

$::=$

| $[h]$
| \star
| $\lambda x : A . t$
| $\text{Ur } d$
| $\text{Inl } d$
| $\text{Inr } d$
| $\langle d_1, d_2 \rangle$
| $\text{@R } d$
| (d)

S

multiplicity, p

$::=$

| 1
| ω

multiplicity

sub

$::=$

| $x := t$
| $h := \underline{d}$

substitution

subs

$::=$

| sub

substitutions

		sub, subs	
data_with_hole, \underline{d}	::=	<ul style="list-style-type: none"> d h $\text{Ur } \underline{d}$ $\text{Inl } \underline{d}$ $\text{Inr } \underline{d}$ $\langle \underline{d}_1, \underline{d}_2 \rangle$ $\textcircled{R} \underline{d}$ (\underline{d}) 	S
store_affect, sa	::=	$x : D = \underline{d}$	store cell
store_affects	::=	<ul style="list-style-type: none"> sa sa, store_affects 	store cells
store, \mathbb{S}	::=	<ul style="list-style-type: none"> \emptyset $\{\text{store_affects}\}$ $\mathbb{S}_1 \sqcup \mathbb{S}_2$ $\mathbb{S}[\text{subs}]$ 	M store contents
type, A	::=	<ul style="list-style-type: none"> \perp D 	bottom type data type
data_type, D	::=	<ul style="list-style-type: none"> 1 R $D_1 \otimes D_2$ $D_1 \oplus D_2$ $A_1 \multimap A_2$ $[D]^p$ $!D$ (D) $\underline{D}[X := D]$ 	S M unit type recursive type bound to a name product type sum type linear function type destination type exponential
type_with_var, \underline{A}	::=	<ul style="list-style-type: none"> \perp \underline{D} 	
data_type_with_var, \underline{D}	::=	<ul style="list-style-type: none"> X 1 R $\underline{D}_1 \otimes \underline{D}_2$ 	type variable in recursive definition unit type recursive type bound to a name product type

	$ \begin{array}{ l} \hline \underline{D}_1 \oplus \underline{D}_2 \\ \underline{A}_1 \multimap \underline{A}_2 \\ [\underline{D}]^p \\ !\underline{D} \\ (\underline{D}) \\ \hline \end{array} $	<p>sum type</p> <p>linear function type</p> <p>destination type</p> <p>exponential</p> <p>S</p>
<code>rec_type_bound</code> , R	$::=$	recursive type bound to a name
<code>rec_type_def</code>	$ \begin{array}{ l} \hline \mu X. \underline{D} \\ \hline \end{array} $	
<code>type_affect</code> , ta	$ \begin{array}{ l} \hline x : A \\ h^p : D \\ \hline \end{array} $	<p>type affectation</p> <p>variable</p> <p>hole</p>
<code>type_affects</code>	$ \begin{array}{ l} \hline ta \\ ta, type_affects \\ \hline \end{array} $	type affectations
<code>typing_context</code> , Γ , \mathcal{U} , H , H_P , H_N	$ \begin{array}{ l} \hline \emptyset \\ \{type_affects\} \\ \Gamma_1 \sqcup \Gamma_2 \\ \hline \end{array} $	typing context
<code>command</code>	$ \begin{array}{ l} \hline S \mid t \\ \hline \end{array} $	
<code>terminals</code>	$ \begin{array}{ l} \hline () \\ \mapsto \\ \oplus \\ \multimap \\ := \\ \vdash \\ \dashv \\ \sqcup \\ \emptyset \\ \neq \\ \in \\ \notin \\ \backslash n \\ \langle \\ \rangle \\ Inl \\ Inr \\ Ur \\ Dest \\ \triangleleft \\ \hline \end{array} $	

	$ \begin{array}{ l} \textcolor{blue}{\underline{c}} \\ \Downarrow \\ \underline{\text{fix}} \\ \textcolor{blue}{\perp} \\ \bullet \\ \subset \\ \mathcal{N} \\ \Longrightarrow \\ \textcolor{blue}{@} \\ \wedge \end{array} $	
formula	$ \begin{array}{ l} ::= \\ \text{judgement} \end{array} $	
Ctx	$ \begin{array}{ l} ::= \\ x \in \mathcal{N}(\Gamma) \\ x \notin \mathcal{N}(\Gamma) \\ \text{type_affect} \in \Gamma \\ \mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset \\ \textcolor{blue}{p_1} = \textcolor{blue}{p_2} \Longrightarrow \Gamma_1 = \Gamma_2 \\ \textcolor{blue}{p_1} = \textcolor{blue}{p_2} \Longrightarrow (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4) \end{array} $	Γ_1 and Γ_2 are disjoint typing contexts with no
Store	$ \begin{array}{ l} ::= \\ \text{fresh } \textcolor{brown}{h} \\ \text{store_affect} \in \mathbb{S} \\ x \notin \mathcal{N}(\mathbb{S}) \end{array} $	
Eq	$ \begin{array}{ l} ::= \\ A_1 = A_2 \\ A_1 \neq A_2 \\ t = u \\ \Gamma = D \end{array} $	
Ty	$ \begin{array}{ l} ::= \\ R \stackrel{\text{fix}}{=} \text{rec_type_def} \\ \mathcal{U} ; \Gamma \vdash \text{command} : A \\ H_N \vdash \underline{d}^{\textcolor{blue}{p}} : D \vdash H_P \\ S \vdash H \\ \mathcal{U} ; H \sqcup \Gamma \vdash t : A \end{array} $	H_P stands for "provides", H_N for "needs"
Sem	$ \begin{array}{ l} ::= \\ \text{command} \Downarrow \text{command}' \end{array} $	
judgement	$ \begin{array}{ l} ::= \\ \text{Ctx} \\ \text{Store} \\ \text{Eq} \\ \text{Ty} \\ \text{Sem} \end{array} $	

user_syntax ::=

- metavariable
- term
- hole
- val
- data
- multiplicity
- sub
- subs
- data_with_hole
- store_affect
- store_affects
- store
- type
- data_type
- type_with_var
- data_type_with_var
- rec_type_bound
- rec_type_def
- type_affect
- type_affects
- typing_context
- command
- terminals

$x \in \mathcal{N}(\Gamma)$

$x \notin \mathcal{N}(\Gamma)$

$\text{type_affect} \in \Gamma$

$\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ Γ_1 and Γ_2 are disjoint typing contexts with no clashing variable names or labels

$p_1 = p_2 \implies \Gamma_1 = \Gamma_2$

$p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$

fresh h

$\text{store_affect} \in \mathbb{S}$

$x \notin \mathcal{N}(\mathbb{S})$

$A_1 = A_2$

$A_1 \neq A_2$

$t = u$

$\Gamma = D$

$R \stackrel{\text{fix}}{=} \text{rec_type_def}$

$\mathcal{U} ; \Gamma \vdash \text{command} : A$

$$\frac{\mathbb{S} \vdash H \quad \mathcal{U} ; H \cup \Gamma \vdash t : A}{\mathcal{U} ; \Gamma \vdash \mathbb{S} | t : A} \text{TYCOMM_DEF}$$

$H_N \vdash \underline{d}^P : D \vdash H_P$

H_P stands for "provides", H_N for "needs"

$$\begin{array}{c}
\frac{}{\emptyset \vdash \textcolor{brown}{h} \textcolor{blue}{p}: \textcolor{blue}{D} \vdash \{\textcolor{brown}{h} \textcolor{blue}{p}: \textcolor{blue}{D}\}} \text{TyDh_H} \\
\frac{}{\{\textcolor{brown}{h} \textcolor{blue}{p}: \textcolor{blue}{D}\} \vdash [\textcolor{brown}{h}]^I: [\textcolor{blue}{D}]^p \vdash \emptyset} \text{TyDh_Dest} \\
\frac{}{\emptyset \vdash \star \textcolor{blue}{p}: \textcolor{blue}{1} \vdash \emptyset} \text{TyDh_U} \\
\frac{\emptyset ; H_N \sqcup \{\textcolor{blue}{x}: \textcolor{blue}{A}_1\} \vdash \textcolor{blue}{t}: \textcolor{blue}{A}_2 \quad p = \omega \implies H_N = \emptyset}{H_N \vdash \lambda \textcolor{blue}{x}: \textcolor{blue}{D}. \textcolor{blue}{t} \textcolor{blue}{p}: \textcolor{blue}{A}_1 \multimap \textcolor{blue}{A}_2 \vdash \emptyset} \text{TyDh_Fn} \\
\frac{H_N \vdash \underline{d} \textcolor{blue}{\omega}: \textcolor{blue}{D} \vdash H_P}{H_N \vdash \text{Ur } \underline{d} \textcolor{blue}{p}: \textcolor{blue}{!D} \vdash H_P} \text{TyDh_E} \\
\frac{H_N \vdash \underline{d} \textcolor{blue}{p}: \textcolor{blue}{D}_1 \vdash H_P}{H_N \vdash \text{Inl } \underline{d} \textcolor{blue}{p}: \textcolor{blue}{D}_1 \oplus \textcolor{blue}{D}_2 \vdash H_P} \text{TyDh_Inl} \\
\frac{H_N \vdash \underline{d} \textcolor{blue}{p}: \textcolor{blue}{D}_2 \vdash H_P}{H_N \vdash \text{Inr } \underline{d} \textcolor{blue}{p}: \textcolor{blue}{D}_1 \oplus \textcolor{blue}{D}_2 \vdash H_P} \text{TyDh_Inr} \\
\frac{H_{N_1} \vdash \underline{d}_1 \textcolor{blue}{p}: \textcolor{blue}{D}_1 \vdash H_{P_1} \quad H_{N_2} \vdash \underline{d}_2 \textcolor{blue}{p}: \textcolor{blue}{D}_2 \vdash H_{P_2}}{H_{N_1} \sqcup H_{N_2} \vdash \langle \underline{d}_1, \underline{d}_2 \rangle \textcolor{blue}{p}: \textcolor{blue}{D}_1 \otimes \textcolor{blue}{D}_2 \vdash H_{P_1} \sqcup H_{P_2}} \text{TyDh_P} \\
\frac{R \stackrel{\text{fix}}{=} \mu X. \underline{D} \quad H_N \vdash \underline{d} \textcolor{blue}{p}: \underline{D}[X := R] \vdash H_P}{H_N \vdash \textcolor{blue}{@R} \underline{d} \textcolor{blue}{p}: R \vdash H_P} \text{TyDh_R}
\end{array}$$

$$\boxed{\mathbb{S} \vdash H}$$

$$\begin{array}{c}
\frac{}{\emptyset \vdash \emptyset} \text{TyStore_Empty} \\
\frac{\mathbb{S} \vdash H \sqcup H_N \quad H_N \vdash \underline{d} \textcolor{blue}{I}: \textcolor{blue}{D} \vdash H_P}{\mathbb{S} \sqcup \{\textcolor{blue}{x}: \textcolor{blue}{D} = \underline{d}\} \vdash H \sqcup H_P} \text{TyStore_Root}
\end{array}$$

$$\boxed{\mathbb{U} ; H \sqcup \Gamma \vdash \textcolor{blue}{t}: \textcolor{blue}{A}}$$

$$\begin{array}{c}
\frac{}{\mathbb{U} ; \emptyset \sqcup \emptyset \vdash \bullet: \textcolor{blue}{\perp}} \text{TyTerm_NoEff} \\
\frac{}{\mathbb{U} ; \{\textcolor{brown}{h} \textcolor{blue}{p}: \textcolor{blue}{D}\} \sqcup \emptyset \vdash [\textcolor{brown}{h}]: [\textcolor{blue}{D}]^p} \text{TyTerm_Dest} \\
\frac{}{\mathbb{U} ; \emptyset \sqcup \emptyset \vdash \star: \textcolor{blue}{1}} \text{TyTerm_U} \\
\frac{\emptyset ; H \sqcup \{\textcolor{blue}{x}: \textcolor{blue}{A}_1\} \vdash \textcolor{blue}{t}: \textcolor{blue}{A}_2}{\mathbb{U} ; H \sqcup \emptyset \vdash \lambda \textcolor{blue}{x}: \textcolor{blue}{A}_1. \textcolor{blue}{t}: \textcolor{blue}{A}_1 \multimap \textcolor{blue}{A}_2} \text{TyTerm_Fn} \\
\frac{\emptyset ; \emptyset \sqcup \emptyset \vdash \textcolor{blue}{d}: \textcolor{blue}{D}}{\mathbb{U} ; \emptyset \sqcup \emptyset \vdash \text{Ur } \textcolor{blue}{d}: \textcolor{blue}{!D}} \text{TyTerm_E} \\
\frac{\emptyset ; H \sqcup \emptyset \vdash \textcolor{blue}{d}: \textcolor{blue}{D}_1}{\mathbb{U} ; H \sqcup \emptyset \vdash \text{Inl } \textcolor{blue}{d}: \textcolor{blue}{D}_1 \oplus \textcolor{blue}{D}_2} \text{TyTerm_Inl} \\
\frac{\emptyset ; H \sqcup \emptyset \vdash \textcolor{blue}{d}: \textcolor{blue}{D}_2}{\mathbb{U} ; H \sqcup \emptyset \vdash \text{Inr } \textcolor{blue}{d}: \textcolor{blue}{D}_1 \oplus \textcolor{blue}{D}_2} \text{TyTerm_Inr}
\end{array}$$

$$\begin{array}{c}
\frac{\emptyset ; H_1 \sqcup \emptyset \vdash d_1 : D_1 \quad \emptyset ; H_2 \sqcup \emptyset \vdash d_2 : D_2}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \emptyset \vdash \langle d_1, d_2 \rangle : D_1 \otimes D_2} \text{TYTERM_P} \\
\\
\frac{\begin{array}{c} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ \emptyset ; H \sqcup \emptyset \vdash d : \underline{D}[X := R] \end{array}}{\mathcal{U} ; H \sqcup \emptyset \vdash @R d : R} \text{TYTERM_R} \\
\\
\frac{}{\mathcal{U} ; \emptyset \sqcup \{x : A\} \vdash x : A} \text{TYTERM_ID} \\
\\
\frac{}{\mathcal{U} \sqcup \{x : A\} ; \emptyset \sqcup \emptyset \vdash x : A} \text{TYTERM_ID'} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : A_1 \multimap A_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A_1}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t u : A_2} \text{TYTERM_APP} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \perp \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A_2}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \ ; \ u : A_2} \text{TYTERM_EFFSEQ} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : 1 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\star \mapsto u\} : A} \text{TYTERM_PATU} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : !D \quad \mathcal{U} \sqcup \{x : D\} ; H_2 \sqcup \Gamma_2 \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} : A} \text{TYTERM_PATE} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : D_1 \oplus D_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : D_1\} \vdash u_1 : A \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_2 : D_2\} \vdash u_2 : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : A} \text{TYTERM_PATs} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : D_1 \otimes D_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : D_1, x_2 : D_2\} \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : A} \text{TYTERM_PATP} \\
\\
\frac{\begin{array}{c} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : R \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \underline{D}[X := R]\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{@R x \mapsto u\} : A} \text{TYTERM_PATR} \\
\\
\frac{\mathcal{U} ; H \sqcup \Gamma \sqcup \{x : [\underline{D}]^I\} \vdash t : \perp}{\mathcal{U} ; H \sqcup \Gamma \vdash \text{alloc } x. t : \underline{D}} \text{TYTERM_ALLOC} \\
\\
\frac{\mathcal{U} ; H \sqcup \Gamma \vdash t : [\underline{1}]^p}{\mathcal{U} ; H \sqcup \Gamma \vdash t \triangleleft^p \star : \perp} \text{TYTERM_FILLU} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : [A_1 \multimap A_2]^p \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : A_1\} \vdash u : A_2 \quad p = \omega \implies (H_2 = \emptyset \wedge \Gamma_2 = \emptyset)}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \lambda x : A_1. u : \perp} \text{TYTERM_FILLFN} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : [\underline{D}]^p \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : \underline{D} \quad p = \omega \implies (H_2 = \emptyset \wedge \Gamma_2 = \emptyset)}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p u : \perp} \text{TYTERM_FILLL}
\end{array}$$

$$\begin{array}{c}
\frac{\begin{array}{c} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket !D \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D \rrbracket^\omega\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Ur } x.u : A} \text{TYTERM_FILLE} \\
\\
\frac{\begin{array}{c} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \oplus D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D_1 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inl } x.u : A} \text{TYTERM_FILLINL} \\
\\
\frac{\begin{array}{c} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \oplus D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D_2 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inr } x.u : A} \text{TYTERM_FILLINR} \\
\\
\frac{\begin{array}{c} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \otimes D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : \llbracket D_1 \rrbracket^p, x_2 : \llbracket D_2 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \langle x_1, x_2 \rangle . u : A} \text{TYTERM_FILLP} \\
\\
\frac{\begin{array}{c} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket R \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket \underline{D}[X := R] \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p @R x.u : A} \text{TYTERM_FILLR}
\end{array}$$

command \Downarrow command'

$$\begin{array}{c}
\frac{}{\mathbb{S} \mid v \Downarrow \mathbb{S} \mid v} \text{SEMOP_VAL} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lambda x : A. t' \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \\ \mathbb{S}_2 \mid t'[x := v_2] \Downarrow \mathbb{S}_3 \mid v_3 \end{array}}{\mathbb{S}_0 \mid t u \Downarrow \mathbb{S}_3 \mid v_3} \text{SEMOP_APP} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \bullet \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid t ; u \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_EFFSEQ} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \star \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\star \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATU} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Ur } d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Ur } y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATE} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inl } d \\ \mathbb{S}_1 \mid u_1[y_1 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINL} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inr } d \\ \mathbb{S}_1 \mid u_2[y_2 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINR} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \langle d_1, d_2 \rangle \\ \mathbb{S}_1 \mid u[y_1 := d_1, y_2 := d_2] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\langle y_1, y_2 \rangle \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATP} \\
\\
\frac{\begin{array}{c} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid @R d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{@R y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATR}
\end{array}$$

$$\begin{array}{c}
\text{fresh } h \\
\frac{\mathbb{S}_0 \sqcup \{x : D = h\} \mid t[x := \lfloor h \rfloor] \Downarrow \mathbb{S}_1 \sqcup \{x : D = d\} \mid \bullet}{\mathbb{S}_0 \mid \text{alloc } x . t \Downarrow \mathbb{S}_1 \mid d} \text{ SEMOP_ALLOC} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \star \Downarrow \mathbb{S}_1[h := \star] \mid \bullet} \text{ SEMOP_FILLU} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \lambda x : A . u \Downarrow \mathbb{S}_1[h := \lambda x : A . u] \mid \bullet} \text{ SEMOP_FILLFN} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid d_2}{\mathbb{S}_0 \mid t \triangleleft^p u \Downarrow \mathbb{S}_2[h := d_2] \mid \bullet} \text{ SEMOP_FILLL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Ur } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Ur } y . u \Downarrow \mathbb{S}_2 \mid v_2} \text{ SEMOP_FILLE} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inl } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inl } y . u \Downarrow \mathbb{S}_2 \mid v_2} \text{ SEMOP_FILLINL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inr } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inr } y . u \Downarrow \mathbb{S}_2 \mid v_2} \text{ SEMOP_FILLINR} \\
\\
\text{fresh } h_1 \\
\text{fresh } h_2 \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \langle h_1, h_2 \rangle] \mid u[y_1 := \lfloor h_1 \rfloor, y_2 := \lfloor h_2 \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \langle y_1, y_2 \rangle . u \Downarrow \mathbb{S}_2 \mid v_2} \text{ SEMOP_FILLP} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{OR } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{OR } x . u \Downarrow \mathbb{S}_2 \mid v_2} \text{ SEMOP_FILLR}
\end{array}$$

Definition rules: 57 good 0 bad
 Definition rule clauses: 152 good 0 bad