

metavariable, x, y

term, t, u	$::=$	
		v
		x
		$t\ u$
		$t \ ;\ u$
		$\text{case } t \text{ of } \{\star \mapsto u\}$
		$\text{case } t \text{ of } \{\text{Ur } x \mapsto u\}$
		$\text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\}$
		$\text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}$
		$\text{case } t \text{ of } \{\text{@R } x \mapsto u\}$
		$\text{alloc } x . t$
		$t \triangleleft^p \star$
		$t \triangleleft^p \lambda x : A . u$
		$t \triangleleft^p u$
		$t \triangleleft^p \text{Ur } y . u$
		$t \triangleleft^p \text{Inl } y . u$
		$t \triangleleft^p \text{Inr } y . u$
		$t \triangleleft^p \langle y_1, y_2 \rangle . u$
		$t \triangleleft^p \text{@R } y . u$
		(t)
		$t[\text{subs}]$

term

value
variable
application
effect sequencing
pattern-matching on unit
pattern-matching on exponentiated value
pattern-matching on sum
pattern-matching on product
unroll for recursive types
allocate data
fill destination with unit
fill destination with function

fill destination with exponential
fill sum-type destination with variant 1
fill sum-type destination with variant 2
fill product-type destination
fill destination with recursive type

S
M

hole, h

$::=$

val, v

$::=$

| \bullet
| d

unreducible value
no-effect effect
data structure

data, d

$::=$

| $[h]$
| \star
| $\lambda x : A . t$
| $\text{Ur } d$
| $\text{Inl } d$
| $\text{Inr } d$
| $\langle d_1, d_2 \rangle$
| $\text{@R } d$
| (d)

S

multiplicity, p

$::=$

| 1
| ω

multiplicity

sub

$::=$

| $x := t$
| $h := d_h$

substitution

subs

$::=$

| sub

substitutions

		sub, subs	
data_with_hole, d_h	::=		
		d	
		h	
		Ur d_h	
		Inl d_h	
		Inr d_h	
		$\langle d_{h1}, d_{h2} \rangle$	
		$\textcircled{R} d_h$	
		(d_h)	S
store_affect, sa	::=		store cell
		$x : D = d_h$	
store_affects	::=		store cells
		sa	
		sa, store_affects	
store, \mathbb{S}	::=		store contents
		\emptyset	
		$\{\text{store_affects}\}$	
		$\mathbb{S}_1 \sqcup \mathbb{S}_2$	
		$\mathbb{S}[\text{subs}]$	M
type, A	::=		
		\perp	bottom type
		D	data type
data_type, D	::=		
		1	unit type
		R	recursive type bound to a name
		$D_1 \otimes D_2$	product type
		$D_1 \oplus D_2$	sum type
		$A_1 \multimap A_2$	linear function type
		$[D_1]^p$	destination type
		$!D$	exponential
		(D)	S
		$D_V[X := D]$	M
type_with_var, A_V	::=		
		\perp	
		D_V	
data_type_with_var, D_V	::=		
		X	type variable in recursive definition
		1	unit type
		R	recursive type bound to a name
		$D_{V1} \otimes D_{V2}$	product type

	$ \begin{array}{l} \quad D_{V_1} \oplus D_{V_2} \\ \quad A_{V_1} \multimap A_{V_2} \\ \quad [D_V]^p \\ \quad !D_V \\ \quad (D_V) \end{array} $	<p>sum type</p> <p>linear function type</p> <p>destination type</p> <p>exponential</p>
	S	
<code>rec_type_bound, R</code>	$::=$	recursive type bound to a name
<code>rec_type_def</code>	$ \begin{array}{l} \quad \mu X. D_V \end{array} $	
<code>type_affect, ta</code>	$ \begin{array}{l} \quad x : A \\ \quad h^p : D \end{array} $	<p>type affectation</p> <p>variable</p> <p>hole</p>
<code>type_affects</code>	$ \begin{array}{l} \quad ta \\ \quad ta, type_affects \end{array} $	type affectations
<code>typing_context, Γ, \mathcal{U}, Γ_h, P, N</code>	$ \begin{array}{l} \quad \emptyset \\ \quad \{type_affects\} \\ \quad \Gamma_1 \sqcup \Gamma_2 \end{array} $	typing context
<code>command</code>	$ \begin{array}{l} \quad S t \end{array} $	
<code>terminals</code>	$ \begin{array}{l} \quad () \\ \quad \mapsto \\ \quad \oplus \\ \quad \multimap \\ \quad := \\ \quad \vdash \\ \quad \dashv \\ \quad \dashv\vdash \\ \quad \sqcup \\ \quad \emptyset \\ \quad \neq \\ \quad \in \\ \quad \notin \\ \quad \backslash n \\ \quad \langle \\ \quad \rangle \\ \quad \text{Inl} \\ \quad \text{Inr} \\ \quad \text{Ur} \\ \quad \text{Dest} \\ \quad \triangleleft \end{array} $	

	$\frac{}{\vdash}$ \Downarrow $\frac{\text{fix}}{=}$ \perp \bullet \subset \mathcal{N} \Rightarrow $\textcircled{\bullet}$ \wedge	
formula	$::=$ \mid judgement	
Ctx	$::=$ $\mid x \in \mathcal{N}(\Gamma)$ $\mid x \notin \mathcal{N}(\Gamma)$ $\mid \text{type_affect} \in \Gamma$ $\mid \mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ $\mid p_1 = p_2 \Rightarrow \Gamma_1 = \Gamma_2$ $\mid p_1 = p_2 \Rightarrow (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$	Γ_1 and Γ_2 are disjoint typing contexts with no
Store	$::=$ $\mid \text{fresh } h$ $\mid \text{store_affect} \in \mathbb{S}$ $\mid x \notin \mathcal{N}(\mathbb{S})$	
Eq	$::=$ $\mid A_1 = A_2$ $\mid A_1 \neq A_2$ $\mid t = u$ $\mid \Gamma = D$	
Ty	$::=$ $\mid R \stackrel{\text{fix}}{=} \text{rec_type_def}$ $\mid \mathcal{U} ; \Gamma \vdash \text{command} : A$ $\mid P ; N \dashv\vdash_{dh} p : D$ $\mid \Gamma_h \dashv \mathbb{S}$ $\mid \mathcal{U} ; \Gamma_h \sqcup \Gamma \vdash t : A$	P stands for "provides", N for "needs"
Sem	$::=$ $\mid \text{command} \Downarrow \text{command'}$	
judgement	$::=$ $\mid \text{Ctx}$ $\mid \text{Store}$ $\mid \text{Eq}$ $\mid \text{Ty}$	

| Sem

user_syntax ::=

- | metavariable
- | term
- | *hole*
- | val
- | data
- | *multiplicity*
- | sub
- | subs
- | data_with_hole
- | store_affect
- | store_affects
- | store
- | type
- | data_type
- | type_with_var
- | data_type_with_var
- | rec_type_bound
- | rec_type_def
- | type_affect
- | type_affects
- | typing_context
- | command
- | terminals

$x \in \mathcal{N}(\Gamma)$

$x \notin \mathcal{N}(\Gamma)$

$\text{type_affect} \in \Gamma$

$\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ Γ_1 and Γ_2 are disjoint typing contexts with no clashing variable names or labels

$p_1 = p_2 \implies \Gamma_1 = \Gamma_2$

$p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$

fresh *h*

$\text{store_affect} \in \mathbb{S}$

$x \notin \mathcal{N}(\mathbb{S})$

$A_1 = A_2$

$A_1 \neq A_2$

$t = u$

$\Gamma = D$

$R \stackrel{\text{fix}}{=} \text{rec_type_def}$

$\mathbb{U} ; \Gamma \vdash \text{command} : A$

$$\frac{\Gamma_h \dashv \mathbb{S} \quad \mathbb{U} ; \Gamma_h \sqcup \Gamma \vdash t : A}{\mathbb{U} ; \Gamma \vdash \mathbb{S} \mid t : A} \text{TYCOMM_DEF}$$

$P ; N \dashv\vdash d_h \text{ }^p\text{: } D$ P stands for "provides", N for "needs"

$$\begin{array}{c}
\frac{}{\overline{\{h^p: D\}; \emptyset \dashv\vdash h^p: D}} \text{TyDH_H} \\
\frac{}{\overline{\{h^p: D\}; \{h^p: D\} \dashv\vdash [h]^I: [D]^p}} \text{TyDH_DEST} \\
\frac{}{\overline{\emptyset; \emptyset \dashv\vdash \star^p: 1}} \text{TyDH_U} \\
\frac{\emptyset; \Gamma_h \sqcup \{x: A_1\} \vdash t: A_2 \quad p = \omega \implies \Gamma_h = \emptyset}{\emptyset; \Gamma_h \dashv\vdash \lambda x: D. t^p: A_1 \multimap A_2} \text{TyDH_FN} \\
\frac{P; N \dashv\vdash d_h^\omega: D}{P; N \dashv\vdash \text{Ur } d_h^p: !D} \text{TyDH_E} \\
\frac{P; N \dashv\vdash d_h^p: D_1}{P; N \dashv\vdash \text{Inl } d_h^p: D_1 \oplus D_2} \text{TyDH_INL} \\
\frac{P; N \dashv\vdash d_h^p: D_2}{P; N \dashv\vdash \text{Inr } d_h^p: D_1 \oplus D_2} \text{TyDH_INR} \\
\frac{P_1; N_1 \dashv\vdash d_{h1}^p: D_1 \quad P_2; N_2 \dashv\vdash d_{h2}^p: D_2}{P_1 \sqcup P_2; N_1 \sqcup N_2 \dashv\vdash \langle d_{h1}, d_{h2} \rangle^p: D_1 \otimes D_2} \text{TyDH_P} \\
\frac{R \stackrel{\text{fix}}{=} \mu X. D_V \quad P; N \dashv\vdash d_h^p: D_V[X := R]}{P; N \dashv\vdash @R_{d_h}^p: R} \text{TyDH_R}
\end{array}$$

$$\boxed{\Gamma_h \dashv\vdash \mathbb{S}}$$

$$\begin{array}{c}
\overline{\emptyset \dashv\vdash \emptyset} \text{TySTORE_EMPTY} \\
\frac{\Gamma_h \sqcup N \dashv\vdash \mathbb{S} \quad P; N \dashv\vdash d_h^I: D}{\Gamma_h \sqcup P \dashv\vdash \mathbb{S} \sqcup \{x: D = d_h\}} \text{TySTORE_ROOT}
\end{array}$$

$$\boxed{\mathcal{U}; \Gamma_h \sqcup \Gamma \vdash t: A}$$

$$\begin{array}{c}
\overline{\mathcal{U}; \emptyset \sqcup \emptyset \vdash \bullet: \perp} \text{TyTERM_NoEff} \\
\frac{}{\overline{\mathcal{U}; \{h^p: D\} \sqcup \emptyset \vdash [h]: [D]^p}} \text{TyTERM_DEST} \\
\frac{}{\overline{\mathcal{U}; \emptyset \sqcup \emptyset \vdash \star: 1}} \text{TyTERM_U} \\
\frac{\emptyset; \Gamma_h \sqcup \{x: A_1\} \vdash t: A_2}{\mathcal{U}; \Gamma_h \sqcup \emptyset \vdash \lambda x: A_1. t: A_1 \multimap A_2} \text{TyTERM_FN} \\
\frac{\emptyset; \emptyset \sqcup \emptyset \vdash d: D}{\mathcal{U}; \emptyset \sqcup \emptyset \vdash \text{Ur } d: !D} \text{TyTERM_E} \\
\frac{\emptyset; \Gamma_h \sqcup \emptyset \vdash d: D_1}{\mathcal{U}; \Gamma_h \sqcup \emptyset \vdash \text{Inl } d: D_1 \oplus D_2} \text{TyTERM_INL} \\
\frac{\emptyset; \Gamma_h \sqcup \emptyset \vdash d: D_2}{\mathcal{U}; \Gamma_h \sqcup \emptyset \vdash \text{Inr } d: D_1 \oplus D_2} \text{TyTERM_INR}
\end{array}$$

$$\begin{array}{c}
\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \emptyset \vdash d_1 : D_1 \quad \emptyset ; \Gamma_{h2} \sqcup \emptyset \vdash d_2 : D_2}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \emptyset \vdash \langle d_1, d_2 \rangle : D_1 \otimes D_2} \text{TYTERM_P}} \\
\\
\frac{\frac{R \stackrel{\text{fix}}{=} \mu X. D_V \quad \emptyset ; \Gamma_h \sqcup \emptyset \vdash d : D_V[X := R]}{\emptyset ; \Gamma_h \sqcup \emptyset \vdash @R d : R} \text{TYTERM_R}} \\
\\
\frac{}{\emptyset ; \emptyset \sqcup \{x : A\} \vdash x : A} \text{TYTERM_ID} \\
\\
\frac{}{\emptyset \sqcup \{x : A\} ; \emptyset \sqcup \emptyset \vdash x : A} \text{TYTERM_ID'} \\
\\
\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : A_1 \multimap A_2 \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \vdash u : A_1}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t u : A_2} \text{TYTERM_APP}} \\
\\
\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : \perp \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \vdash u : A_2}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t ; u : A_2} \text{TYTERM_EFFSEQ}} \\
\\
\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : 1 \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \vdash u : A}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\star \mapsto u\} : A} \text{TYTERM_PATU}} \\
\\
\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : !D \quad \emptyset \sqcup \{x : D\} ; \Gamma_{h2} \sqcup \Gamma_2 \vdash u : A}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} : A} \text{TYTERM_PATE}} \\
\\
\frac{\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : D_1 \oplus D_2 \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \sqcup \{x_1 : D_1\} \vdash u_1 : A \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \sqcup \{x_2 : D_2\} \vdash u_2 : A}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : A} \text{TYTERM_PATs}} \\
\\
\frac{\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : D_1 \otimes D_2 \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \sqcup \{x_1 : D_1, x_2 : D_2\} \vdash u : A}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : A} \text{TYTERM_PATP}} \\
\\
\frac{\frac{\frac{R \stackrel{\text{fix}}{=} \mu X. D_V \quad \emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : R \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \sqcup \{x : D_V[X := R]\} \vdash u : A}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{@R x \mapsto u\} : A} \text{TYTERM_PATR}} \\
\\
\frac{\frac{\emptyset ; \Gamma_h \sqcup \Gamma \sqcup \{x : [D]^I\} \vdash t : \perp}{\emptyset ; \Gamma_h \sqcup \Gamma \vdash \text{alloc } x. t : D} \text{TYTERM_ALLOC}} \\
\\
\frac{\frac{\emptyset ; \Gamma_h \sqcup \Gamma \vdash t : [1]^p}{\emptyset ; \Gamma_h \sqcup \Gamma \vdash t \triangleleft^p \star : \perp} \text{TYTERM_FILLU}} \\
\\
\frac{\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : [A_1 \multimap A_2]^p \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \sqcup \{x : A_1\} \vdash u : A_2 \quad p = \omega \implies (\Gamma_{h2} = \emptyset \wedge \Gamma_2 = \emptyset)}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \lambda x : A_1. u : \perp} \text{TYTERM_FILLFN}} \\
\\
\frac{\frac{\frac{\emptyset ; \Gamma_{h1} \sqcup \Gamma_1 \vdash t : [D]^p \quad \emptyset ; \Gamma_{h2} \sqcup \Gamma_2 \vdash u : D \quad p = \omega \implies (\Gamma_{h2} = \emptyset \wedge \Gamma_2 = \emptyset)}{\emptyset ; \Gamma_{h1} \sqcup \Gamma_{h2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p u : \perp} \text{TYTERM_FILLL}}
\end{array}$$

$$\begin{array}{c}
\frac{\begin{array}{l} \mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_1 \vdash t : [!D]^p \\ \mathcal{U} ; \Gamma_{h_2} \sqcup \Gamma_2 \sqcup \{x : [D]^\omega\} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_{h_2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Ur } x.u : A} \text{TYTERM_FILLE} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_1 \vdash t : [D_1 \oplus D_2]^p \\ \mathcal{U} ; \Gamma_{h_2} \sqcup \Gamma_2 \sqcup \{x : [D_1]^p\} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_{h_2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inl } x.u : A} \text{TYTERM_FILLINL} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_1 \vdash t : [D_1 \oplus D_2]^p \\ \mathcal{U} ; \Gamma_{h_2} \sqcup \Gamma_2 \sqcup \{x : [D_2]^p\} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_{h_2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inr } x.u : A} \text{TYTERM_FILLINR} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_1 \vdash t : [D_1 \otimes D_2]^p \\ \mathcal{U} ; \Gamma_{h_2} \sqcup \Gamma_2 \sqcup \{x_1 : [D_1]^p, x_2 : [D_2]^p\} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_{h_2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \langle x_1, x_2 \rangle . u : A} \text{TYTERM_FILLP} \\
\\
\frac{\begin{array}{l} R \stackrel{\text{fix}}{=} \mu X. D_V \\ \mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_1 \vdash t : [R]^p \\ \mathcal{U} ; \Gamma_{h_2} \sqcup \Gamma_2 \sqcup \{x : [D_V[X := R]]^p\} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_{h_1} \sqcup \Gamma_{h_2} \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p @R x.u : A} \text{TYTERM_FILLR}
\end{array}$$

command	\Downarrow	command'
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$$\begin{array}{c}
\frac{}{\mathbb{S} \mid v \Downarrow \mathbb{S} \mid v} \text{SEMOP_VAL} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lambda x : A. t' \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \\ \mathbb{S}_2 \mid t'[x := v_2] \Downarrow \mathbb{S}_3 \mid v_3 \end{array}}{\mathbb{S}_0 \mid t u \Downarrow \mathbb{S}_3 \mid v_3} \text{SEMOP_APP} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \bullet \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid t ; u \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_EFFSEQ} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \star \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\star \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATU} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Ur } d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Ur } y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATE} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inl } d \\ \mathbb{S}_1 \mid u_1[y_1 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINL} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inr } d \\ \mathbb{S}_1 \mid u_2[y_2 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINR} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \langle d_1, d_2 \rangle \\ \mathbb{S}_1 \mid u[y_1 := d_1, y_2 := d_2] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\langle y_1, y_2 \rangle \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATP} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid @R d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{@R y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATR}
\end{array}$$

$$\begin{array}{c}
\text{fresh } h \\
\frac{\mathbb{S}_0 \sqcup \{x : D = h\} \mid t[x := \lfloor h \rfloor] \Downarrow \mathbb{S}_1 \sqcup \{x : D = d\} \mid \bullet}{\mathbb{S}_0 \mid \text{alloc } x . t \Downarrow \mathbb{S}_1 \mid d} \quad \text{SEMOP_ALLOC} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \star \Downarrow \mathbb{S}_1[h := \star] \mid \bullet} \quad \text{SEMOP_FILLU} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \lambda x : A . u \Downarrow \mathbb{S}_1[h := \lambda x : A . u] \mid \bullet} \quad \text{SEMOP_FILLFN} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid d_2}{\mathbb{S}_0 \mid t \triangleleft^p u \Downarrow \mathbb{S}_2[h := d_2] \mid \bullet} \quad \text{SEMOP_FILLL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Ur } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Ur } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOP_FILLE} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inl } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inl } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOP_FILLINL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inr } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inr } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOP_FILLINR} \\
\\
\text{fresh } h_1 \\
\text{fresh } h_2 \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \langle h_1, h_2 \rangle] \mid u[y_1 := \lfloor h_1 \rfloor, y_2 := \lfloor h_2 \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \langle y_1, y_2 \rangle . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOP_FILLP} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{OR } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{OR } x . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOP_FILLR}
\end{array}$$

Definition rules: 57 good 0 bad
 Definition rule clauses: 152 good 0 bad