

# Destination $\lambda$ -calculus

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## 1 Term and value syntax

$\text{var}, x, y$  Term-level variable name  
 $k$  Index for ranges

$\text{hdn}, h$	$::=$ $  \quad h + h'$ $  \quad \max(H)$	Hole or destination name ( $\mathbb{N}$ ) M Sum M Maximum of a set of holes
$\text{hdns}, H$	$::=$ $  \quad \{h_1, \dots, h_k\}$ $  \quad H_1 \cup H_2$ $  \quad H \pm h$ $  \quad \text{hnames}(\Gamma)$ $  \quad \text{hnames}(C)$	Set of hole names M Union of sets M Increase all names from $H$ by $h$ . M Hole names of a context (requires $\text{ctx\_NoVar}(\Gamma)$ ) M Hole names of an evaluation context
term, $t, u$	$::=$ $  \quad v$ $  \quad x$ $  \quad t \succ u$ $  \quad t ; u$ $  \quad t \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$ $  \quad t \succ \text{case}_m (x_1, x_2) \mapsto u$ $  \quad t \succ \text{case}_m E^m x \mapsto u$ $  \quad t \succ \text{map } x \mapsto u$ $  \quad \text{to}_x t$ $  \quad \text{from}_x t$ $  \quad \text{alloc}$ $  \quad t \triangleleft ()$ $  \quad t \triangleleft \text{Inl}$ $  \quad t \triangleleft \text{Inr}$ $  \quad t \triangleleft E^m$ $  \quad t \triangleleft (,)$ $  \quad t \triangleleft (\lambda x_m \mapsto u)$ $  \quad t \triangleleft \bullet u$ $  \quad t[x := v]$	Term Value Variable Application Pattern-match on unit Pattern-match on sum Pattern-match on product Pattern-match on exponential Map over the right side of ampar $t$ Wrap $t$ into a trivial ampar Extract value from trivial ampar Return a fresh "identity" ampar object Fill destination with unit Fill destination with left variant Fill destination with right variant Fill destination with exponential constructor Fill destination with product constructor Fill destination with function Fill destination with root of ampar $u$ M
val, $v$	$::=$ $  \quad -h$ $  \quad +h$ $  \quad ()$ $  \quad \lambda^v x_m \mapsto t$ $  \quad \text{Inl } v$ $  \quad \text{Inr } v$ $  \quad E^m v$ $  \quad (v_1, v_2)$ $  \quad H(v_1, v_2)$ $  \quad v \pm h$	Term value Hole Destination Unit Lambda abstraction Left variant for sum Right variant for sum Exponential Product Ampar M Rename hole names inside $v$ by shifting them by $h$

ctx, C	::=		Evaluation context
		$\square$	Identity
		$C \succ u$	Application
		$v \succ C$	Application
		$C ; u$	Pattern-match on unit
		$C \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$	Pattern-match on sum
		$C \succ \text{case}_m (x_1, x_2) \mapsto u$	Pattern-match on product
		$C \succ \text{case}_m E^n x \mapsto u$	Pattern-match on exponential
		$C \succ \text{map } x \mapsto u$	Map over the right side of ampar
		$\text{to}_\times C$	Wrap into a trivial ampar
		$\text{from}_\times C$	Extract value from trivial ampar
		$C \triangleleft ()$	Fill destination with unit
		$C \triangleleft \text{Inl}$	Fill destination with left variant
		$C \triangleleft \text{Inr}$	Fill destination with right variant
		$C \triangleleft E^m$	Fill destination with exponential constructor
		$C \triangleleft ()$	Fill destination with product constructor
		$C \triangleleft (\lambda x_m \mapsto u)$	Fill destination with function
		$C \triangleleft \bullet u$	Fill destination with root of ampar
		$v \triangleleft \bullet C$	Fill destination with root of ampar
		$\overset{\text{op}}{\mathbf{H}} \langle v_1, C \rangle$	Open ampar. <b>Only new addition to term shapes</b>
		$C \circ C'$	M Compose evaluation contexts
		$C[\mathbf{h} :=_{\mathbf{H}} v]$	M Fill $\mathbf{h}$ with value $v$ (that may contain holes)

## 2 Type system

type, T, U	::=		Type
		$1$	Unit
		$T_1 \oplus T_2$	Sum
		$T_1 \otimes T_2$	Product
		$!^m T$	Exponential
		$T_1 \times T_2$	Ampar type (consuming $T_2$ yields $T_1$ )
		$T_1 \xrightarrow{m_1} T_2$	Function
		$[T]^m$	Destination
mode, m, n	::=		Mode (Semiring)
		$pa$	Pair of a multiplicity and age
		$\omega$	Error case (incompatible types, multiplicities, or ages)
		$m_1 \cdot \dots \cdot m_k$	M Semiring product
mul, p	::=		Multiplicity (first component of modality)
		$1$	Linear. Neutral element of the product
		$\omega$	Non-linear. Absorbing for the product
		$p_1 \cdot \dots \cdot p_k$	M Semiring product
age, a	::=		Age (second component of modality)
		$\nu$	Born now. Neutral element of the product
		$\uparrow$	One scope older
		$\infty$	Infinitely old / static. Absorbing for the product
		$a_1 \cdot \dots \cdot a_k$	M Semiring product
bndr, b	::=		Type assignment to either variable, destination or hole
		$x :_m T$	Variable
		$+h :_m [T]^n$	Destination ( $m$ is its own modality; $n$ is the modality for values it accepts)
		$-h : T^n$	Hole ( $n$ is the modality for values it accepts, it doesn't have a modality on its own)
ctx, $\Gamma, \Delta, \Pi$	::=		Typing context
		$\{b_1, \dots, b_k\}$	List of bindings
		$m \cdot \Gamma$	M Multiply each binding by $m$
		$\Gamma_1 \uplus \Gamma_2$	M Sum contexts $\Gamma_1$ and $\Gamma_2$ . Duplicates/incompatible elements will give bindings with
		$-\Gamma$	M Transforms every dest binding into a hole binding (requires <code>ctx_DestOnly</code> $\Gamma$ )

$\boxed{\Gamma \Vdash v : \mathbf{T}}$ 
*(Typing of values (raw))*

$\frac{\text{TYR-VAL-H}}{\{-\mathbf{h} : \mathbf{T}^{\text{lv}}\} \Vdash -\mathbf{h} : \mathbf{T}}$	$\frac{\text{TYR-VAL-D} \quad \text{ctx\_Compatible } \Gamma \quad +\mathbf{h} :_{\text{lv}} [\mathbf{T}]^n}{\Gamma \Vdash +\mathbf{h} : [\mathbf{T}]^n}$	$\frac{\text{TYR-VAL-U}}{\{\} \Vdash () : \mathbf{1}}$	$\frac{\text{TYR-VAL-F} \quad \text{ctx\_DestOnly } \Delta \quad \Delta \uplus \{\mathbf{x} :_m \mathbf{T}_1\} \vdash t : \mathbf{T}_2}{\Delta \Vdash \lambda^v \mathbf{x}_m \mapsto t : \mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2}$
$\frac{\text{TYR-VAL-L} \quad \Gamma \Vdash v : \mathbf{T}_1}{\Gamma \Vdash \text{Inl } v : \mathbf{T}_1 \oplus \mathbf{T}_2}$	$\frac{\text{TYR-VAL-R} \quad \Gamma \Vdash v : \mathbf{T}_2}{\Gamma \Vdash \text{Inr } v : \mathbf{T}_1 \oplus \mathbf{T}_2}$	$\frac{\text{TYR-VAL-P} \quad \Gamma_1 \Vdash v_1 : \mathbf{T}_1 \quad \Gamma_2 \Vdash v_2 : \mathbf{T}_2}{\Gamma_1 \uplus \Gamma_2 \Vdash (v_1, v_2) : \mathbf{T}_1 \otimes \mathbf{T}_2}$	$\frac{\text{TYR-VAL-E} \quad \Gamma \Vdash v : \mathbf{T}}{n \cdot \Gamma \Vdash \mathbf{E}^n v : !^n \mathbf{T}}$
$\frac{\text{TYR-VAL-A} \quad \begin{array}{l} \text{ctx\_DestOnly } \Delta_1 \\ \text{ctx\_DestOnly } \Delta_2 \\ \text{ctx\_DestOnly } \Delta_3 \\ \text{ctx\_Disjoint } \Delta_1 \Delta_2 \\ \text{ctx\_Disjoint } \Delta_1 \Delta_3 \\ \text{ctx\_Disjoint } \Delta_2 \Delta_3 \\ \Delta_1 \uplus (-\Delta_3) \Vdash v_1 : \mathbf{T}_1 \\ \Delta_2 \uplus \Delta_3 \Vdash v_2 : \mathbf{T}_2 \end{array}}{\Delta_1 \uplus \Delta_2 \Vdash \text{hnames}(-\Delta_3)(v_1, v_2) : \mathbf{T}_1 \ltimes \mathbf{T}_2}$			

 $\boxed{\Pi \vdash t : \mathbf{T}}$ 
*(Typing of terms)*

$\frac{\text{TY-TERM-VAL} \quad \text{ctx\_DestOnly } \Delta \quad \Delta \vdash v : \mathbf{T}}{\Delta \vdash v : \mathbf{T}}$	$\frac{\text{TY-TERM-VAR} \quad \text{ctx\_Compatible } \Pi \quad \mathbf{x} :_{\text{lv}} \mathbf{T}}{\Pi \vdash \mathbf{x} : \mathbf{T}}$	$\frac{\text{TY-TERM-APP} \quad \Pi_1 \vdash t : \mathbf{T}_1 \quad \Pi_2 \vdash u : \mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2}{m \cdot \Pi_1 \uplus \Pi_2 \vdash t \succ u : \mathbf{T}_2}$		
$\frac{\text{TY-TERM-PATU} \quad \Pi_1 \vdash t : \mathbf{1} \quad \Pi_2 \vdash u : \mathbf{U}}{\Pi_1 \uplus \Pi_2 \vdash t ; u : \mathbf{U}}$	$\frac{\text{TY-TERM-PATS} \quad \begin{array}{l} \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x}_1 :_m \mathbf{T}_1\} \\ \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x}_2 :_m \mathbf{T}_2\} \\ \Pi_1 \vdash t : \mathbf{T}_1 \oplus \mathbf{T}_2 \\ \Pi_2 \uplus \{\mathbf{x}_1 :_m \mathbf{T}_1\} \vdash u_1 : \mathbf{U} \\ \Pi_2 \uplus \{\mathbf{x}_2 :_m \mathbf{T}_2\} \vdash u_2 : \mathbf{U} \end{array}}{m \cdot \Pi_1 \uplus \Pi_2 \vdash t \succ \text{case}_m \{ \text{Inl } \mathbf{x}_1 \mapsto u_1, \text{Inr } \mathbf{x}_2 \mapsto u_2 \} : \mathbf{U}}$			
$\frac{\text{TY-TERM-PATP} \quad \begin{array}{l} \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x}_1 :_m \mathbf{T}_1\} \\ \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x}_2 :_m \mathbf{T}_2\} \\ \text{ctx\_Disjoint } \{\mathbf{x}_1 :_m \mathbf{T}_1\} \quad \{\mathbf{x}_2 :_m \mathbf{T}_2\} \\ \Pi_1 \vdash t : \mathbf{T}_1 \otimes \mathbf{T}_2 \\ \Pi_2 \uplus \{\mathbf{x}_1 :_m \mathbf{T}_1, \mathbf{x}_2 :_m \mathbf{T}_2\} \vdash u : \mathbf{U} \end{array}}{m \cdot \Pi_1 \uplus \Pi_2 \vdash t \succ \text{case}_m(\mathbf{x}_1, \mathbf{x}_2) \mapsto u : \mathbf{U}}$	$\frac{\text{TY-TERM-PATE} \quad \begin{array}{l} \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x} :_{m \cdot n} \mathbf{T}\} \\ \Pi_1 \vdash t : !^n \mathbf{T} \\ \Pi_2 \uplus \{\mathbf{x} :_{m \cdot n} \mathbf{T}\} \vdash u : \mathbf{U} \end{array}}{m \cdot \Pi_1 \uplus \Pi_2 \vdash t \succ \text{case}_m \mathbf{E}^n \mathbf{x} \mapsto u : \mathbf{U}}$	$\frac{\text{TY-TERM-MAP} \quad \begin{array}{l} \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x} :_{\text{lv}} \mathbf{T}_2\} \\ \Pi_1 \vdash t : \mathbf{T}_1 \ltimes \mathbf{T}_2 \\ \mathcal{I} \uparrow \cdot \Pi_2 \uplus \{\mathbf{x} :_{\text{lv}} \mathbf{T}_2\} \vdash u : \mathbf{U} \end{array}}{\Pi_1 \uplus \Pi_2 \vdash t \succ \text{map } \mathbf{x} \mapsto u : \mathbf{T}_1 \ltimes \mathbf{U}}$		
$\frac{\text{TY-TERM-TOA} \quad \Pi \vdash t : \mathbf{T}}{\Pi \vdash \text{to}_{\ltimes} t : \mathbf{T} \ltimes \mathbf{1}}$	$\frac{\text{TY-TERM-FROMA} \quad \Pi \vdash t : \mathbf{T} \ltimes \mathbf{1}}{\Pi \vdash \text{from}_{\ltimes} t : \mathbf{T}}$	$\frac{\text{TY-TERM-ALLOC}}{\{\} \vdash \text{alloc} : \mathbf{T} \ltimes [\mathbf{T}]^{\text{lv}}}$	$\frac{\text{TY-TERM-FILLU} \quad \Pi \vdash t : [\mathbf{1}]^n}{\Pi \vdash t \triangleleft () : \mathbf{1}}$	$\frac{\text{TY-TERM-FILLL} \quad \Pi \vdash t : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n}{\Pi \vdash t \triangleleft \text{Inl} : [\mathbf{T}_1]^n}$
$\frac{\text{TY-TERM-FILLR} \quad \Pi \vdash t : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n}{\Pi \vdash t \triangleleft \text{Inr} : [\mathbf{T}_2]^n}$	$\frac{\text{TY-TERM-FILLP} \quad \Pi \vdash t : [\mathbf{T}_1 \otimes \mathbf{T}_2]^n}{\Pi \vdash t \triangleleft (,) : [\mathbf{T}_1]^n \otimes [\mathbf{T}_2]^n}$	$\frac{\text{TY-TERM-FILLE} \quad \Pi \vdash t : [!^{n'} \mathbf{T}]^n}{\Pi \vdash t \triangleleft \mathbf{E}^{n'} : [\mathbf{T}]^{n' \cdot n}}$	$\frac{\text{TY-TERM-FILLF} \quad \begin{array}{l} \text{ctx\_Disjoint } \Pi_2 \quad \{\mathbf{x} :_m \mathbf{T}_1\} \\ \Pi_1 \vdash t : [\mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2]^n \\ \Pi_2 \uplus \{\mathbf{x} :_m \mathbf{T}_1\} \vdash u : \mathbf{T}_2 \end{array}}{\Pi_1 \uplus (\mathcal{I} \uparrow \cdot n) \cdot \Pi_2 \vdash t \triangleleft (\lambda \mathbf{x}_m \mapsto u) : \mathbf{1}}$	
$\frac{\text{TY-TERM-FILLC} \quad \Pi_1 \vdash t : [\mathbf{T}_1]^n \quad \Pi_2 \vdash u : \mathbf{T}_1 \ltimes \mathbf{T}_2}{\Pi_1 \uplus (\mathcal{I} \uparrow \cdot n) \cdot \Pi_2 \vdash t \triangleleft \bullet u : \mathbf{T}_2}$				

$$\Delta \vdash C : \mathbf{T}_1 \multimap \mathbf{T}_2$$

(Typing of evaluation contexts)

$$\begin{array}{c} \text{TYR-ECTX-ID} \\ \frac{\text{TYR-ECTX-APPFOC1} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad m \cdot \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{T}_2 \multimap \mathbf{U}_0 \quad \Delta_2 \vdash u : \mathbf{T}_1 \multimap \mathbf{T}_2}{\{\} \vdash \square : \mathbf{U}_0 \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-APPFOC2} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad m \cdot \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{T}_2 \multimap \mathbf{U}_0 \quad \Delta_1 \vdash v : \mathbf{T}_1}{\Delta_2 \vdash C \circ (v \succ \square) : (\mathbf{T}_1 \multimap \mathbf{T}_2) \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-PATUFOC} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{U} \multimap \mathbf{U}_0 \quad \Delta_2 \vdash u : \mathbf{U}}{\Delta_1 \vdash C \circ (\square ; u) : \mathbf{1} \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-PATSFoc} \\ \frac{\text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \wedge \text{ctx\_DestOnly } \Delta_2 \quad m \cdot \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{U} \multimap \mathbf{U}_0 \quad \Delta_2 \uplus \{x_1 : m \mathbf{T}_1\} \vdash u_1 : \mathbf{U} \quad \Delta_2 \uplus \{x_2 : m \mathbf{T}_2\} \vdash u_2 : \mathbf{U}}{\Delta_1 \vdash C \circ (\square \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}) : (\mathbf{T}_1 \oplus \mathbf{T}_2) \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-PATPFoc} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \wedge \text{ctx\_DestOnly } \Delta_2 \quad \text{ctx\_Disjoint } \{x_1 : m \mathbf{T}_1\} \{x_2 : m \mathbf{T}_2\} \quad m \cdot \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{U} \multimap \mathbf{U}_0 \quad \Delta_2 \uplus \{x_1 : m \mathbf{T}_1, x_2 : m \mathbf{T}_2\} \vdash u : \mathbf{U}}{\Delta_1 \vdash C \circ (\square \succ \text{case}_m (x_1, x_2) \mapsto u) : (\mathbf{T}_1 \otimes \mathbf{T}_2) \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-PATEFOC} \\ \frac{\text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad m \cdot \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{U} \multimap \mathbf{U}_0 \quad \Delta_2 \uplus \{x : m \mathbf{T}\} \vdash u : \mathbf{U}}{\Delta_1 \vdash C \circ (\square \succ \text{case}_m E^{m'} x \mapsto u) : !^{m'} \mathbf{T} \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-MAPFOC} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \Delta_1 \uplus \Delta_2 \vdash C : \mathbf{U} \multimap \mathbf{U}_0 \quad I \uparrow \cdot \Delta_2 \uplus \{x : !v \mathbf{T}_2\} \vdash u : \mathbf{U}}{\Delta_1 \vdash C \circ (\square \succ \text{map } x \mapsto u) : (\mathbf{T}_1 \ltimes \mathbf{T}_2) \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-TOAFOC} \quad \Delta \vdash C : (\mathbf{T} \ltimes \mathbf{1}) \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\text{to}_\ltimes \square) : \mathbf{T} \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-FROMAFOC} \\ \frac{\Delta \vdash C : \mathbf{T} \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\text{from}_\ltimes \square) : (\mathbf{T} \ltimes \mathbf{1}) \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLUFOC} \quad \Delta \vdash C : \mathbf{1} \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\square \triangleleft ()) : [\mathbf{1}]^n \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLFOC} \quad \Delta \vdash C : [\mathbf{T}_1]^n \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\square \triangleleft \text{Inl}) : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-FILLRFOC} \\ \frac{\Delta \vdash C : [\mathbf{T}_2]^n \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\square \triangleleft \text{Inr}) : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLPFoc} \quad \Delta \vdash C : ([\mathbf{T}_1]^n \otimes [\mathbf{T}_2]^n) \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\square \triangleleft (,)) : [\mathbf{T}_1 \otimes \mathbf{T}_2]^n \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLEFOC} \quad \Delta \vdash C : [\mathbf{T}]^{m \cdot n} \multimap \mathbf{U}_0}{\Delta \vdash C \circ (\square \triangleleft E^m) : [!^m \mathbf{T}]^n \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-FILLFFoc} \\ \frac{\text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \Delta_1 \uplus (I \uparrow \cdot n) \cdot \Delta_2 \vdash C : \mathbf{1} \multimap \mathbf{U}_0 \quad \Delta_2 \uplus \{x : m \mathbf{T}_1\} \vdash u : \mathbf{T}_2}{\Delta_1 \vdash C \circ (\square \triangleleft (\lambda x_m \mapsto u)) : [\mathbf{T}_1 \multimap \mathbf{T}_2]^n \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLCFoc1} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \Delta_1 \uplus (I \uparrow \cdot n) \cdot \Delta_2 \vdash C : \mathbf{T}_2 \multimap \mathbf{U}_0 \quad \Delta_2 \vdash u : \mathbf{T}_1 \ltimes \mathbf{T}_2}{\Delta_1 \vdash C \circ (\square \triangleleft \bullet u) : [\mathbf{T}_1]^n \multimap \mathbf{U}_0} \quad \frac{\text{TYR-ECTX-FILLCFoc2} \quad \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \Delta_1 \uplus (I \uparrow \cdot n) \cdot \Delta_2 \vdash C : \mathbf{T}_2 \multimap \mathbf{U}_0 \quad \Delta_1 \vdash v : [\mathbf{T}_1]^n}{\Delta_2 \vdash C \circ (v \triangleleft \bullet \square) : \mathbf{T}_1 \ltimes \mathbf{T}_2 \multimap \mathbf{U}_0} \end{array}$$

$$\begin{array}{c} \text{TYR-ECTX-AOPENFOC} \\ \text{ctx\_Disjoint } \Delta_1 \Delta_2 \quad \text{ctx\_Disjoint } \Delta_1 \Delta_3 \quad \text{hdns\_Disjoint } \text{hnames}(C) \text{ hnames}(-\Delta_3) \quad \text{ctx\_DestOnly } \Delta_1 \quad \text{ctx\_DestOnly } \Delta_2 \quad \text{ctx\_DestOnly } \Delta_3 \quad \Delta_1 \uplus \Delta_2 \vdash C : (\mathbf{T}_1 \ltimes \mathbf{U}) \multimap \mathbf{U}_0 \quad \Delta_1 \uplus -\Delta_3 \Vdash v_1 : \mathbf{T}_1 \\ \hline I \uparrow \cdot \Delta_2 \uplus \Delta_3 \vdash C \circ (\overset{\text{op}}{\text{hnames}}(-\Delta_3) \langle v_1, \square \rangle) : \mathbf{U} \multimap \mathbf{U}_0 \end{array}$$

$$\vdash C[t] : \mathbf{T}$$

(Typing of extended terms (pair of evaluation context and term))

$$\begin{array}{c} \text{TY-ETERM-CLOSEDETERM} \\ \frac{\Gamma \vdash C : \mathbf{T} \multimap \mathbf{U}_0 \quad \Gamma \vdash t : \mathbf{T}}{\vdash C[t] : \mathbf{U}_0} \end{array}$$

### 3 Small-step semantics

$$\boxed{C[t] \longrightarrow C'[t']}$$

(Small-step evaluation of terms using evaluation contexts)

$$\frac{\text{SEM-ETERM-APPFOC1} \quad \text{term\_NotVal } t}{C[t \succ u] \longrightarrow (C \circ (\Box \succ u))[t]}$$

$$\frac{\text{SEM-ETERM-APPUNFOC1}}{(C \circ (\Box \succ u))[v] \longrightarrow C[v \succ u]}$$

$$\frac{\text{SEM-ETERM-APPFOC2} \quad \text{term\_NotVal } u}{C[v \succ u] \longrightarrow (C \circ (v \succ \Box))[u]}$$

$$\frac{\text{SEM-ETERM-APPUNFOC2}}{(C \circ (v \succ \Box))[v'] \longrightarrow C[v \succ v']}$$

$$\frac{\text{SEM-ETERM-APPRED}}{C[v \succ (\lambda^v x_m \mapsto t)] \longrightarrow C[t[x := v]]}$$

$$\frac{\text{SEM-ETERM-PATUFOC} \quad \text{term\_NotVal } t}{C[t ; u] \longrightarrow (C \circ (\Box ; u))[t]}$$

$$\frac{\text{SEM-ETERM-PATUNFOC}}{(C \circ (\Box ; u))[v] \longrightarrow C[v ; u]}$$

$$\frac{\text{SEM-ETERM-PATURED}}{C[() ; u] \longrightarrow C[u]}$$

$$\frac{\text{SEM-ETERM-PATLFOC} \quad \text{term\_NotVal } t}{C[t \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \longrightarrow (C \circ (\Box \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}))[t]}$$

$$\frac{\text{SEM-ETERM-PATLUNFOC}}{(C \circ (\Box \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}))[v] \longrightarrow C[v \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}]}$$

$$\frac{\text{SEM-ETERM-PATLRED}}{C[(\text{Inl } v) \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \longrightarrow C[u_1[x := v]]}$$

$$\frac{\text{SEM-ETERM-PATRFoc} \quad \text{term\_NotVal } t}{C[t \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \longrightarrow (C \circ (\Box \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}))[t]}$$

$$\frac{\text{SEM-ETERM-PATRUNFOC}}{(C \circ (\Box \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}))[v] \longrightarrow C[v \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}]}$$

$$\frac{\text{SEM-ETERM-PATRRRED}}{C[(\text{Inr } v) \succ \text{case}_m \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \longrightarrow C[u_2[x := v]]}$$

$$\frac{\text{SEM-ETERM-PATPFOC} \quad \text{term\_NotVal } t}{C[t \succ \text{case}_m (x_1, x_2) \mapsto u] \longrightarrow (C \circ (\Box \succ \text{case}_m (x_1, x_2) \mapsto u))[t]}$$

$$\frac{\text{SEM-ETERM-PATPUNFOC}}{(C \circ (\Box \succ \text{case}_m (x_1, x_2) \mapsto u))[v] \longrightarrow C[v \succ \text{case}_m (x_1, x_2) \mapsto u]}$$

$$\frac{\text{SEM-ETERM-PATPREd}}{C[(v_1, v_2) \succ \text{case}_m (x_1, x_2) \mapsto u] \longrightarrow C[u[x_1 := v_1][x_2 := v_2]]}$$

$$\frac{\text{SEM-ETERM-PATEFOC} \quad \text{term\_NotVal } t}{C[t \succ \text{case}_m E^n x \mapsto u] \longrightarrow (C \circ (\Box \succ \text{case}_m E^n x \mapsto u))[t]}$$

$$\frac{\text{SEM-ETERM-PATEUNFOC}}{(C \circ (\Box \succ \text{case}_m E^n x \mapsto u))[v] \longrightarrow C[v \succ \text{case}_m E^n x \mapsto u]}$$

$$\frac{\text{SEM-ETERM-PATERED}}{C[E^n v \succ \text{case}_m E^n x \mapsto u] \longrightarrow C[u[x := v]]}$$

$$\frac{\text{SEM-ETERM-MAPFOC} \quad \text{term\_NotVal } t}{C[t \succ \text{map } x \mapsto u] \longrightarrow (C \circ (\Box \succ \text{map } x \mapsto u))[t]}$$

$$\frac{\text{SEM-ETERM-MAPUNFOC}}{(C \circ (\Box \succ \text{map } x \mapsto u))[v] \longrightarrow C[v \succ \text{map } x \mapsto u]}$$

$$\frac{\text{SEM-ETERM-MAPREDAOPENFOC} \quad h' = \max(hnames(C))}{C[h \langle v_1, v_2 \rangle \succ \text{map } x \mapsto u] \longrightarrow (C \circ (\overset{OP}{H \pm h'} \langle v_1 \pm h', \Box \rangle))[u[x := v_2 \pm h']]}$$

$$\frac{\text{SEM-ETERM-AOPENUNFOC}}{(C \circ \overset{OP}{H} \langle v_1, \Box \rangle)[v_2] \longrightarrow C[h \langle v_1, v_2 \rangle]}$$

$$\frac{\text{SEM-ETERM-ALLOCRED}}{C[\text{alloc}] \longrightarrow C[\{1\} \langle +1, -1 \rangle]}$$

$$\frac{\text{SEM-ETERM-TOAFoc} \quad \text{term\_NotVal } t}{C[\text{to}_\times t] \longrightarrow (C \circ (\text{to}_\times \Box))[t]}$$

$$\frac{\text{SEM-ETERM-TOAUNFOC}}{(C \circ (\text{to}_\times \Box))[v] \longrightarrow C[\text{to}_\times v]}$$

SEM-ETERM-TOARED	SEM-ETERM-FROMAFOC term_NotVal t	SEM-ETERM-FROMAUNFOC
$\overline{C[\text{to}_{\times} v] \rightarrow C[\{ \} \langle v, () \rangle]}$	$\overline{C[\text{from}_{\times} t] \rightarrow (C \circ (\text{from}_{\times} \square))[t]}$	$\overline{(C \circ (\text{from}_{\times} \square))[v] \rightarrow C[\text{from}_{\times} v]}$
SEM-ETERM-FROMARED	SEM-ETERM-FILLUFOC term_NotVal t	SEM-ETERM-FILLUUNFOC
$\overline{C[\text{from}_{\times} \{ \} \langle v, () \rangle] \rightarrow C[v]}$	$\overline{C[t \triangleleft ()] \rightarrow (C \circ (\square \triangleleft ()))[t]}$	$\overline{(C \circ (\square \triangleleft ()))[v] \rightarrow C[v \triangleleft ()]}$
SEM-ETERM-FILLURED	SEM-ETERM-FILLLFOC term_NotVal t	SEM-ETERM-FILLLUNFOC
$\overline{C[+h \triangleleft ()] \rightarrow C[h := \{ \} ()][()]}$	$\overline{C[t \triangleleft \text{Inl}] \rightarrow (C \circ (\square \triangleleft \text{Inl}))[t]}$	$\overline{(C \circ (\square \triangleleft \text{Inl}))[v] \rightarrow C[v \triangleleft \text{Inl}]}$
SEM-ETERM-FILLRED $h' = \max(\text{hnames}(C) \cup \{h\})$	SEM-ETERM-FILLRFOC term_NotVal t	SEM-ETERM-FILLRUNFOC
$\overline{C[+h \triangleleft \text{Inl}] \rightarrow C[h := \{h'+1\} \text{Inl} - (h'+1)][+(h'+1)]}$	$\overline{C[t \triangleleft \text{Inr}] \rightarrow (C \circ (\square \triangleleft \text{Inr}))[t]}$	$\overline{(C \circ (\square \triangleleft \text{Inr}))[v] \rightarrow C[v \triangleleft \text{Inr}]}$
SEM-ETERM-FILLRRED $h' = \max(\text{hnames}(C) \cup \{h\})$	SEM-ETERM-FILLEFOC term_NotVal t	
$\overline{C[+h \triangleleft \text{Inr}] \rightarrow C[h := \{h'+1\} \text{Inr} - (h'+1)][+(h'+1)]}$	$\overline{C[t \triangleleft E^m] \rightarrow (C \circ (\square \triangleleft E^m))[t]}$	
SEM-ETERM-FILLEUNFOC	SEM-ETERM-FILLERED $h' = \max(\text{hnames}(C) \cup \{h\})$	SEM-ETERM-FILLPFOC term_NotVal t
$\overline{(C \circ (\square \triangleleft E^m))[v] \rightarrow C[v \triangleleft E^m]}$	$\overline{C[+h \triangleleft E^m] \rightarrow C[h := \{h'+1\} E^m - (h'+1)][+(h'+1)]}$	$\overline{C[t \triangleleft (,)] \rightarrow (C \circ (\square \triangleleft (,)))[t]}$
SEM-ETERM-FILLPUNFOC	SEM-ETERM-FILLPRED $h' = \max(\text{hnames}(C) \cup \{h\})$	
$\overline{(C \circ (\square \triangleleft (,)))[v] \rightarrow C[v \triangleleft (,)]}$	$\overline{C[+h \triangleleft (,)] \rightarrow C[h := \{h'+1, h'+2\} (- (h'+1), - (h'+2))][+(h'+1), +(h'+2)]}$	
SEM-ETERM-FILLFFOC term_NotVal t	SEM-ETERM-FILLFUNFOC	
$\overline{C[t \triangleleft (\lambda \times_m \mapsto u)] \rightarrow (C \circ (\square \triangleleft (\lambda \times_m \mapsto u)))[t]}$	$\overline{(C \circ (\square \triangleleft (\lambda \times_m \mapsto u)))[v] \rightarrow C[v \triangleleft (\lambda \times_m \mapsto u)]}$	
SEM-ETERM-FILLFRED	SEM-ETERM-FILLCFOC1 term_NotVal t	SEM-ETERM-FILLCUNFOC1
$\overline{C[+h \triangleleft (\lambda \times_m \mapsto u)] \rightarrow C[h := \{ \} \lambda^v \times_m \mapsto u][()]}$	$\overline{C[t \triangleleft \bullet u] \rightarrow (C \circ (\square \triangleleft \bullet u))[t]}$	$\overline{(C \circ (\square \triangleleft \bullet u))[v] \rightarrow C[v \triangleleft \bullet u]}$
SEM-ETERM-FILLCFOC2	SEM-ETERM-FILLCUNFOC2	SEM-ETERM-FILLCRED $h' = \max(\text{hnames}(C) \cup \{h\})$
$\overline{C[v \triangleleft \bullet u] \rightarrow (C \circ (v \triangleleft \bullet \square))[u]}$	$\overline{(C \circ (v \triangleleft \bullet \square))[v'] \rightarrow C[v \triangleleft \bullet v']}$	$\overline{C[+h \triangleleft \bullet_h \langle v_1, v_2 \rangle] \rightarrow C[h := (\text{H} \pm h') v_1 \pm h'][v_2 \pm h']}$