

Destination λ -calculus

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1 Term and value syntax

tvar, x, y, d Term-level variable
 hvar, h Hole
 n

v ::=

- $\langle v_1, \overline{v_2} \rangle \Delta$
- $@h$
- $()$
- $\text{Inl } v$
- $\text{Inr } v$
- (v_1, v_2)
- $\mathcal{J}^M v$
- $\lambda x. t$

Term value

- Ampar
- Destination
- Unit
- Left variant for sum
- Right variant for sum
- Product
- Exponential
- Linear function

\bar{v} ::=

- v
- h
- $\text{Inl } \bar{v}$
- $\text{Inr } \bar{v}$
- (\bar{v}_1, \bar{v}_2)
- $\mathcal{J}^M \bar{v}$

Pseudo-value that may contain holes

- Term value
- Hole
- Left variant with val or hole
- Right variant with val or hole
- Product with val or hole
- Exponential with val or hole

t, u ::=

- v
- x
- $t \succ u$
- $t \succ \text{case } () \mapsto u$
- $t \succ \text{case } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$
- $t \succ \text{case } (x_1, x_2) \mapsto u$
- $t \succ \text{case } \mathcal{J}^M x \mapsto u$
- $t \succ \text{mapL } x \mapsto u$
- $\text{to}_x t$
- $\text{from}_x t$
- alloc_T
- $t \triangleleft ()$
- $t \triangleleft \text{Inl}$
- $t \triangleleft \text{Inr}$
- $t \triangleleft (,)$
- $t \triangleleft \mathcal{J}^M$
- $t \triangleleft \bullet u$

Term

- Term value
- Variable
- Application
- Pattern-match on unit
- Pattern-match on sum
- Pattern-match on product
- Pattern-match on exponential
- Map over the left side of the ampar
- Wrap t into a trivial ampar
- Extract value from trivial ampar
- Return a fresh "identity" ampar object
- Fill destination with unit
- Fill destination with left variant
- Fill destination with right variant
- Fill destination with product constructor
- Fill destination with exponential constructor
- Fill destination with root of ampar u

2 Type system

\mathbf{T}, \mathbf{U}	$::=$	$\mathbf{1}$ $\mathbf{T}_1 \oplus \mathbf{T}_2$ $\mathbf{T}_1 \otimes \mathbf{T}_2$ $!^{\mathbf{M}} \mathbf{T}$ $\mathbf{T}_1 \ltimes \mathbf{T}_2$ $\mathbf{T}_1 \multimap \mathbf{T}_2$ $^{\mathbf{M}}[\mathbf{T}]$	Type Unit Sum Product Exponential Ampar type (consuming \mathbf{T}_1 yields \mathbf{T}_2) Linear function Destination
\mathbf{M}, \mathbf{N}	$::=$	ν \uparrow ∞ $\mathbf{M}_1 \cdot \dots \cdot \mathbf{M}_n$	Multiplicity (Semiring with product \cdot) Born now. Identity of the product One scope older Infinitely old / static. Absorbing for product Semiring product
PC, Γ	$::=$	$\{\mathbf{PA}_1, \dots, \mathbf{PA}_n\}$ $\mathbf{M} \cdot \Gamma$ $\Gamma_1 \sqcup \Gamma_2$	Positive typing context Increase age of bindings by \mathbf{M}
PA	$::=$	$\mathbf{x} :_{\mathbf{M}} \mathbf{T}$ $\textcircled{\mathbf{h}} :_{\mathbf{M}} {}^{\mathbf{N}}[\mathbf{T}]$	Positive type assignment Variable Destination (\mathbf{M} is its own age; \mathbf{N} is the age of values it accepts)
NC, Δ	$::=$	$\{\mathbf{NA}_1, \dots, \mathbf{NA}_n\}$ $\mathbf{M} \cdot \Delta$ $\textcircled{\mathbf{a}}^{-1} \Gamma$ $\Delta_1 \sqcup \Delta_2$	Negative typing context Increase age of bindings by \mathbf{M} Invert the sign of the context
NA	$::=$	$\mathbf{h} :^{\mathbf{N}} \mathbf{T}$	Negative type assignment Hole (\mathbf{N} is the age of values it accepts, its own age is undefined)

$$\boxed{\Gamma \cup \Delta \Vdash \bar{v} : \mathbf{T}}$$

(Typing of extended values (require both positive and negative contexts))

$$\text{TY-W-H} \quad \frac{}{\{\} \cup \{\mathbf{h} :_{\nu} \mathbf{T}\} \Vdash \mathbf{h} : \mathbf{T}}$$

$$\text{TY-W-D} \quad \frac{}{\{\textcircled{\mathbf{h}} :_{\nu} {}^{\mathbf{N}}[\mathbf{T}]\} \cup \{\} \Vdash \textcircled{\mathbf{h}} : {}^{\mathbf{N}}[\mathbf{T}]}$$

$$\text{TY-W-U} \quad \frac{}{\{\} \cup \{\} \Vdash () : \mathbf{1}}$$

$$\text{TY-W-L} \quad \frac{\Gamma \cup \Delta \Vdash \bar{v} : \mathbf{T}_1 \quad \mathbf{dom}(\Gamma) \cap \mathbf{dom}(\Delta) = \emptyset}{\Gamma \cup \Delta \Vdash \text{Inl } \bar{v} : \mathbf{T}_1 \oplus \mathbf{T}_2}$$

$$\text{TY-W-P} \quad \frac{\begin{array}{l} \Gamma_1 \cup \Delta_1 \Vdash \bar{v}_1 : \mathbf{T}_1 \\ \Gamma_2 \cup \Delta_2 \Vdash \bar{v}_2 : \mathbf{T}_2 \\ \mathbf{dom}(\Gamma_1) \cap \mathbf{dom}(\Gamma_2) = \emptyset \\ \mathbf{dom}(\Gamma_1) \cap \mathbf{dom}(\Delta_1) = \emptyset \\ \mathbf{dom}(\Gamma_1) \cap \mathbf{dom}(\Delta_2) = \emptyset \\ \mathbf{dom}(\Gamma_2) \cap \mathbf{dom}(\Delta_1) = \emptyset \\ \mathbf{dom}(\Gamma_2) \cap \mathbf{dom}(\Delta_2) = \emptyset \\ \mathbf{dom}(\Delta_1) \cap \mathbf{dom}(\Delta_2) = \emptyset \end{array}}{\Gamma_1 \sqcup \Gamma_2 \cup \Delta_1 \sqcup \Delta_2 \Vdash (\bar{v}_1, \bar{v}_2) : \mathbf{T}_1 \otimes \mathbf{T}_2}$$

$$\text{TY-W-R} \quad \frac{\Gamma \cup \Delta \Vdash \bar{v} : \mathbf{T}_2 \quad \mathbf{dom}(\Gamma) \cap \mathbf{dom}(\Delta) = \emptyset}{\Gamma \cup \Delta \Vdash \text{Inr } \bar{v} : \mathbf{T}_1 \oplus \mathbf{T}_2}$$

$$\text{TY-W-E} \quad \frac{\Gamma \cup \Delta \Vdash \bar{v} : \mathbf{T} \quad \mathbf{dom}(\Gamma) \cap \mathbf{dom}(\Delta) = \emptyset}{\mathbf{M} \cdot \Gamma \cup {}^{\mathbf{M}} \cdot \Delta \Vdash \mathbf{!}^{\mathbf{M}} \bar{v} : !^{\mathbf{M}} \mathbf{T}}$$

$$\text{TY-W-A} \quad \frac{\begin{array}{l} \Gamma_1 \cup \{\} \Vdash v_1 : \mathbf{T}_1 \\ \Gamma_2 \cup \textcircled{\mathbf{a}}^{-1} \Gamma_1 \Vdash \bar{v}_2 : \mathbf{T}_2 \\ \mathbf{dom}(\Gamma_1) \cap \mathbf{dom}(\Gamma_2) = \emptyset \end{array}}{\Gamma_2 \cup \{\} \Vdash \langle v_1, \bar{v}_2 \rangle_{\textcircled{\mathbf{a}}^{-1} \Gamma_1} : \mathbf{T}_1 \ltimes \mathbf{T}_2}$$

$$\text{TY-W-F} \quad \frac{\begin{array}{l} \Gamma \sqcup \{\mathbf{x} :_{\mathbf{M}} \mathbf{T}_1\} \vdash t : \mathbf{T}_2 \\ \mathbf{dom}(\Gamma) \cap \mathbf{dom}(\{\mathbf{x} :_{\mathbf{M}} \mathbf{T}_1\}) = \emptyset \end{array}}{\Gamma \cup \{\} \Vdash \lambda \mathbf{x} . t : \mathbf{T}_1 \multimap \mathbf{T}_2}$$

$$\boxed{\Gamma \vdash t : \mathbf{T}}$$

(Typing of terms (only a positive context is needed))

$$\frac{\text{TY-T-V} \quad \Gamma \sqcup \{ \} \Vdash v : \mathbf{T}}{\Gamma \vdash v : \mathbf{T}}$$

$$\frac{\text{TY-T-X0}}{\{x :_{\nu} \mathbf{T}\} \vdash x : \mathbf{T}}$$

$$\frac{\text{TY-T-XINF}}{\{x :_{\infty} \mathbf{T}\} \vdash x : \mathbf{T}}$$

$$\frac{\text{TY-T-APP} \quad \Gamma_1 \vdash t : \mathbf{T}_1 \quad \Gamma_2 \vdash u : \mathbf{T}_1 \xrightarrow{\mathbf{M}} \mathbf{T}_2 \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset}{\mathbf{M} \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ u : \mathbf{T}_2}$$

TY-T-PATS

$$\frac{\text{TY-T-PATU} \quad \Gamma_1 \vdash t : \mathbf{1} \quad \Gamma_2 \vdash u : \mathbf{U} \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset}{\Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case}() \mapsto u : \mathbf{U}}$$

$$\frac{\Gamma_1 \vdash t : \mathbf{T}_1 \oplus \mathbf{T}_2 \quad \Gamma_2 \sqcup \{x_1 :_{\mathbf{M}} \mathbf{T}_1\} \vdash u_1 : \mathbf{U} \quad \Gamma_2 \sqcup \{x_2 :_{\mathbf{M}} \mathbf{T}_2\} \vdash u_2 : \mathbf{U} \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x_1 :_{\mathbf{M}} \mathbf{T}_1\}) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x_2 :_{\mathbf{M}} \mathbf{T}_2\}) = \emptyset}{\mathbf{M} \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case} \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \} : \mathbf{U}}$$

TY-T-PATP

$$\frac{\Gamma_1 \vdash t : \mathbf{T}_1 \otimes \mathbf{T}_2 \quad \Gamma_2 \sqcup \{x_1 :_{\mathbf{M}} \mathbf{T}_1, x_2 :_{\mathbf{M}} \mathbf{T}_2\} \vdash u : \mathbf{U} \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x_1 :_{\mathbf{M}} \mathbf{T}_1\}) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x_2 :_{\mathbf{M}} \mathbf{T}_2\}) = \emptyset \quad \text{dom}(\{x_1 :_{\mathbf{M}} \mathbf{T}_1\}) \cap \text{dom}(\{x_2 :_{\mathbf{M}} \mathbf{T}_2\}) = \emptyset}{\mathbf{M} \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case}(x_1, x_2) \mapsto u : \mathbf{U}}$$

TY-T-PATE

$$\frac{\Gamma_1 \vdash t : !^{\mathbf{N}} \mathbf{T} \quad \Gamma_2 \sqcup \{x :_{\mathbf{M} \cdot \mathbf{N}} \mathbf{T}\} \vdash u : \mathbf{U} \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x :_{\mathbf{M} \cdot \mathbf{N}} \mathbf{T}\}) = \emptyset}{\mathbf{M} \cdot \Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{case}^{\mathbf{N}} x \mapsto u : \mathbf{U}}$$

TY-T-MAP

$$\frac{\Gamma_1 \vdash t : \mathbf{T}_1 \rtimes \mathbf{T}_2 \quad \uparrow \cdot \Gamma_2 \sqcup \{x :_{\nu} \mathbf{T}_1\} \vdash u : \mathbf{U} \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset \quad \text{dom}(\Gamma_2) \cap \text{dom}(\{x :_{\nu} \mathbf{T}_1\}) = \emptyset}{\Gamma_1 \sqcup \Gamma_2 \vdash t \succ \text{mapL } x \mapsto u : \mathbf{U} \rtimes \mathbf{T}_2}$$

TY-T-FILLC

$$\frac{\Gamma_1 \vdash t : {}^{\mathbf{N}}[\mathbf{T}_2] \quad \Gamma_2 \vdash u : \mathbf{T}_1 \rtimes \mathbf{T}_2 \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_2) = \emptyset}{\Gamma_1 \sqcup (\uparrow \cdot \mathbf{N}) \cdot \Gamma_2 \vdash t \triangleleft \bullet u : \mathbf{T}_1}$$

$$\frac{\text{TY-T-FILLU} \quad \Gamma \vdash t : {}^{\mathbf{N}}[\mathbf{1}]}{\Gamma \vdash t \triangleleft () : \mathbf{1}}$$

$$\frac{\text{TY-T-FILLL} \quad \Gamma \vdash t : {}^{\mathbf{N}}[\mathbf{T}_1 \oplus \mathbf{T}_2]}{\Gamma \vdash t \triangleleft \text{Inl} : {}^{\mathbf{N}}[\mathbf{T}_1]}$$

$$\frac{\text{TY-T-FILLR} \quad \Gamma \vdash t : {}^{\mathbf{N}}[\mathbf{T}_1 \oplus \mathbf{T}_2]}{\Gamma \vdash t \triangleleft \text{Inr} : {}^{\mathbf{N}}[\mathbf{T}_2]}$$

TY-T-FILLP

$$\frac{\Gamma \vdash t : {}^{\mathbf{N}}[\mathbf{T}_1 \otimes \mathbf{T}_2]}{\Gamma \vdash t \triangleleft (,) : {}^{\mathbf{N}}[\mathbf{T}_1] \otimes {}^{\mathbf{N}}[\mathbf{T}_2]}$$

TY-T-FILLE

$$\frac{\Gamma \vdash t : {}^{\mathbf{M}}[!^{\mathbf{N}} \mathbf{T}]}{\Gamma \vdash t \triangleleft \mathbf{N} : {}^{\mathbf{M} \cdot \mathbf{N}}[\mathbf{T}]}$$

TY-T-ALLOC

$$\frac{}{\{ \} \vdash \text{alloc}_{\mathbf{T}} : {}^{\nu}[\mathbf{T}] \rtimes \mathbf{T}}$$

TY-T-TOA

$$\frac{\Gamma \vdash t : \mathbf{T}}{\Gamma \vdash \text{to}_{\rtimes} t : \mathbf{1} \rtimes \mathbf{T}}$$

TY-T-FROMA

$$\frac{\Gamma \vdash t : \mathbf{1} \rtimes \mathbf{T}}{\Gamma \vdash \text{from}_{\rtimes} t : \mathbf{T}}$$

$$\boxed{\Gamma \vdash v \diamond e : \mathbf{T}}$$

(Typing of commands (only a positive context is needed))

TY-C

$$\frac{\Gamma_{11} \sqcup \Gamma_{12} \vdash v : \mathbf{T} \quad \Gamma_2 u @^{-1} \Gamma_{12} \Vdash e \quad \text{dom}(\Gamma_{11}) \cap \text{dom}(\Gamma_{12}) = \emptyset \quad \text{dom}(\Gamma_{11}) \cap \text{dom}(\Gamma_2) = \emptyset \quad \text{dom}(\Gamma_{12}) \cap \text{dom}(\Gamma_2) = \emptyset}{\Gamma_{11} \sqcup \Gamma_2 \vdash v \diamond e : \mathbf{T}}$$

3 Effects and big-step semantics

e	$::=$	Effect
	ε	
	$\mathbf{h} := \bar{v}$	
	e_1, \dots, e_n	Chain effects

$$\boxed{\Gamma \sqcup \Delta \Vdash e}$$

(Typing of effects (require both positive and negative contexts))

<p>TY-E-N</p> $\frac{\{\} \sqcup \{\} \Vdash \varepsilon}{\{\} \sqcup \{\} \Vdash \varepsilon}$	<p>TY-E-A</p> $\frac{\Gamma \sqcup \Delta \Vdash \bar{v} : \mathbf{T} \quad \text{dom}(\Gamma) \cap \text{dom}(\{\mathbf{@h} :_{\nu} {}^N[\mathbf{T}]\}) = \emptyset \quad \text{dom}(\Gamma) \cap \text{dom}(\Delta) = \emptyset \quad \text{dom}(\{\mathbf{@h} :_{\nu} {}^N[\mathbf{T}]\}) \cap \text{dom}(\Delta) = \emptyset}{(\uparrow \mathbf{M} \cdot \mathbf{N}) \cdot \Gamma \sqcup \mathbf{M} \cdot \{\mathbf{@h} :_{\nu} {}^N[\mathbf{T}]\} \sqcup (\mathbf{M} \cdot \mathbf{N}) \cdot \Delta \Vdash \mathbf{h} := \bar{v}}$	<p>TY-E-X</p> $\frac{\Gamma_1 \sqcup \Delta_1 \sqcup \mathbf{@}^{-1} \Gamma_{22} \Vdash e_1 \quad \Gamma_{21} \sqcup \Gamma_{22} \sqcup \Delta_2 \Vdash e_2 \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_{21}) = \emptyset \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Gamma_{22}) = \emptyset \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Delta_1) = \emptyset \quad \text{dom}(\Gamma_1) \cap \text{dom}(\Delta_2) = \emptyset \quad \text{dom}(\Gamma_{21}) \cap \text{dom}(\Gamma_{22}) = \emptyset \quad \text{dom}(\Gamma_{21}) \cap \text{dom}(\Delta_1) = \emptyset \quad \text{dom}(\Gamma_{21}) \cap \text{dom}(\Delta_2) = \emptyset \quad \text{dom}(\Gamma_{22}) \cap \text{dom}(\Delta_1) = \emptyset \quad \text{dom}(\Gamma_{22}) \cap \text{dom}(\Delta_2) = \emptyset \quad \text{dom}(\Delta_1) \cap \text{dom}(\Delta_2) = \emptyset}{\Gamma_1 \sqcup \Gamma_{21} \sqcup \Delta_1 \sqcup \Delta_2 \Vdash e_1, e_2}$
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$$\boxed{\bar{v}_1 \Delta_1 \mid e_1 \rightsquigarrow \bar{v}_2 \Delta_2 \mid e_2}$$

(Big-step evaluation of effects on extended values)

<p>SEM-E-S</p> $\frac{\mathbf{h} \notin \text{dom}(\Delta_1) \quad \bar{v}_1 \Delta_1 \mid e_1 \rightsquigarrow \bar{v}_2 \Delta_2 \mid e_2}{\bar{v}_1 \Delta_1 \mid \mathbf{h} := \bar{v}', e_1 \rightsquigarrow \bar{v}_2 \Delta_2 \mid \mathbf{h} := \bar{v}', e_2}$	<p>SEM-E-F</p> $\frac{\Gamma'_1 \sqcup \Delta'_1 \Vdash \bar{v}' : \mathbf{T} \quad \text{dom}(\Gamma'_1) \cap \text{dom}(\Delta'_1) = \emptyset \quad \text{dom}(\Delta_1) \cap \text{dom}(\{\mathbf{h} :_{\nu} {}^N \mathbf{T}\}) = \emptyset \quad \text{dom}(\Delta_1) \cap \text{dom}(\Delta'_1) = \emptyset \quad \bar{v}_1[\mathbf{h} := \bar{v}'](\Delta_1 \sqcup {}^N \Delta'_1) \mid e_1 \rightsquigarrow \bar{v}_2 \Delta_2 \mid e_2}{\bar{v}_1 \Delta_1 \sqcup \{\mathbf{h} :_{\nu} {}^N \mathbf{T}\} \mid \mathbf{h} := \bar{v}', e_1 \rightsquigarrow \bar{v}_2 \Delta_2 \mid e_2}$
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$$\boxed{t \Downarrow v \diamond e}$$

(Big-step evaluation into commands)

<p>SEM-T-V</p> $\frac{}{v \Downarrow v \diamond \varepsilon}$	<p>SEM-T-APP</p> $\frac{t_1 \Downarrow v_1 \diamond e_1 \quad t_2 \Downarrow \lambda \mathbf{x} . u \diamond e_2 \quad u[\mathbf{x} := v_1] \Downarrow v_3 \diamond e_3}{t_1 \succ t_2 \Downarrow v_3 \diamond e_1, e_2, e_3}$	<p>SEM-T-PATU</p> $\frac{t_1 \Downarrow () \diamond e_1 \quad t_2 \Downarrow v_2 \diamond e_2}{t_1 \succ \text{case } () \mapsto t_2 \Downarrow v_2 \diamond e_1, e_2}$	
<p>SEM-T-PATL</p> $\frac{t \Downarrow \text{Inl } v_1 \diamond e_1 \quad u_1[\mathbf{x}_1 := v_1] \Downarrow v_2 \diamond e_2}{t \succ \text{case } \{\text{Inl } \mathbf{x}_1 \mapsto u_1, \text{Inr } \mathbf{x}_2 \mapsto u_2\} \Downarrow v_2 \diamond e_1, e_2}$	<p>SEM-T-PATR</p> $\frac{t \Downarrow \text{Inr } v_1 \diamond e_1 \quad u_2[\mathbf{x}_2 := v_1] \Downarrow v_2 \diamond e_2}{t \succ \text{case } \{\text{Inl } \mathbf{x}_1 \mapsto u_1, \text{Inr } \mathbf{x}_2 \mapsto u_2\} \Downarrow v_2 \diamond e_1, e_2}$		
<p>SEM-T-PATP</p> $\frac{t \Downarrow (v_1, v_2) \diamond e_1 \quad u[\mathbf{x}_1 := v_1][\mathbf{x}_2 := v_2] \Downarrow v_2 \diamond e_2}{t \succ \text{case } (\mathbf{x}_1, \mathbf{x}_2) \mapsto u \Downarrow v_2 \diamond e_1, e_2}$	<p>SEM-T-MAP</p> $\frac{t \Downarrow \langle v_1, \bar{v}_2 \rangle_{\Delta} \diamond e_1 \quad u[\mathbf{x} := v_1] \Downarrow v_3 \diamond e_2 \quad \bar{v}_2 \Delta \mid e_2 \rightsquigarrow \bar{v}_4 \Delta' \mid e_3}{t \succ \text{mapL } \mathbf{x} \mapsto u \Downarrow \langle v_3, \bar{v}_4 \rangle_{\Delta'} \diamond e_1, e_3}$		
<p>SEM-T-TOA</p> $\frac{t \Downarrow v \diamond e}{\text{to}_{\mathbf{x}} t \Downarrow \langle (), v \rangle_{\{\}} \diamond e}$	<p>SEM-T-FROMA</p> $\frac{t \Downarrow \langle (), v \rangle_{\{\}} \diamond e}{\text{from}_{\mathbf{x}} t \Downarrow v \diamond e}$	<p>SEM-T-FILLU</p> $\frac{t \Downarrow \mathbf{@h} \diamond e}{t \triangleleft () \Downarrow () \diamond e, \mathbf{h} := ()}$	<p>SEM-T-FILLL</p> $\frac{t \Downarrow \mathbf{@h} \diamond e \quad \text{fresh } \mathbf{h}'}{t \triangleleft \text{Inl} \Downarrow \mathbf{@h}' \diamond e, \mathbf{h} := \text{Inl } \mathbf{h}'}$
<p>SEM-T-FILLR</p> $\frac{t \Downarrow \mathbf{@h} \diamond e}{t \triangleleft \text{Inr} \Downarrow \mathbf{@h}' \diamond e, \mathbf{h} := \text{Inr } \mathbf{h}'}$	<p>SEM-T-FILLP</p> $\frac{t \Downarrow \mathbf{@h} \diamond e \quad \text{fresh } \mathbf{h}_1 \quad \text{fresh } \mathbf{h}_2}{t \triangleleft (,) \Downarrow (\mathbf{@h}_1, \mathbf{@h}_2) \diamond e, \mathbf{h} := (\mathbf{h}_1, \mathbf{h}_2)}$		<p>SEM-T-FILLC</p> $\frac{t \Downarrow \mathbf{@h} \diamond e_1 \quad u \Downarrow \langle v_1, \bar{v}_2 \rangle_{\Delta} \diamond e_2}{t \triangleleft \bullet u \Downarrow v_1 \diamond e_1, e_2, \mathbf{h} := \bar{v}_2}$

4 Type safety

Theorem 1 (Type safety). *If **destOnly** Γ and $\Gamma \vdash t : \mathbf{T}$ then $t \Downarrow v \diamond e$ and $\Gamma \vdash v \diamond e : \mathbf{T}$.*

Theorem 2 (Type safety for complete programs). *If $\{\} \vdash t : \mathbf{T}$ then $t \Downarrow v \diamond \varepsilon$ and $\{\} \vdash v : \mathbf{T}$*