

metavariable, x, y

term, t, u ::=

- | v
- | x
- | $t\ u$
- | $t \ ; \ u$
- | $\text{case } t \text{ of } \{ \star \mapsto u \}$
- | $\text{case } t \text{ of } \{ \text{Ur } x \mapsto u \}$
- | $\text{case } t \text{ of } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$
- | $\text{case } t \text{ of } \{ \langle x_1, x_2 \rangle \mapsto u \}$
- | $\text{case } t \text{ of } \{ @R\ x \mapsto u \}$
- | $\text{alloc } x . t$
- | $t \triangleleft^p \star$
- | $t \triangleleft^p \lambda x : A . u$
- | $t \triangleleft^p u$
- | $t \triangleleft^p \text{Ur } y . u$
- | $t \triangleleft^p \text{Inl } y . u$
- | $t \triangleleft^p \text{Inr } y . u$
- | $t \triangleleft^p \langle y_1, y_2 \rangle . u$
- | $t \triangleleft^p @R\ y . u$
- | (t)
- | $t[\text{subs}]$

term

- value
- variable
- application
- effect sequencing
- pattern-matching on unit
- pattern-matching on exponentiated value
- pattern-matching on sum
- pattern-matching on product
- pattern-matching on recursive data
- get data from a dest-filling statement
- fill destination with unit
- fill destination with function
- fill destination with value
- fill destination with exponential
- fill destination with sum variant 1
- fill destination with sum variant 2
- fill destination with product
- fill destination with recursive data

S
M

hole, h

::=

hole in the lexical store

val, v

::=

- | \bullet
- | d

value (unreducible term)

- empty effect
- data structure

data, d

::=

- | $[h]$
- | \star
- | $\lambda x : A . t$
- | $\text{Ur } d$
- | $\text{Inl } d$
- | $\text{Inr } d$
- | $\langle d_1, d_2 \rangle$
- | $@R\ d$
- | (d)

- hole reified as a destination
- unit
- lambda abstraction
- exponential
- sum variant 1
- sum variant 2
- product
- recursive data

S

multiplicity, p

::=

- | 1
- | ω

multiplicity

- for holes/destinations not under a Ur
- for holes/destinations under a Ur

sub

::=

- | $x := v$
- | $h := \underline{d}$

substitution

note: $[h]$ won't be replaced in $\mathbb{S}[h := dh]$

subs

::=

- | sub

substitutions

		sub, subs	
data_with_hole, \underline{d}	::=	d h $\text{Ur } \underline{d}$ $\text{Inl } \underline{d}$ $\text{Inr } \underline{d}$ $\langle \underline{d}_1, \underline{d}_2 \rangle$ $\textcircled{R} \underline{d}$ (\underline{d})	S
store_affect, sa	::=		store cell
		$x : D = \underline{d}$	
store_affects	::=		store cells
		sa	
		sa, store_affects	
store, \mathbb{S}	::=		
		\emptyset	
		$\{\text{store_affects}\}$	
		$\mathbb{S}_1 \sqcup \mathbb{S}_2$	
		$\mathbb{S}[\text{subs}]$	M
type, A	::=		
		\perp	bottom (effect) type
		D	data type
data_type, D	::=		
		1	unit type
		R	recursive type bound to a name
		$D_1 \otimes D_2$	product type
		$D_1 \oplus D_2$	sum type
		$A_1 \multimap A_2$	linear function type
		$[D]^p$	destination type
		$!D$	exponential
		(D)	S
		$\underline{D}[X := D]$	M unroll a recursive data type
type_with_var, \underline{A}	::=		
		\perp	
		\underline{D}	
data_type_with_var, \underline{D}	::=		
		X	
		1	
		R	
		$\underline{D}_1 \otimes \underline{D}_2$	

	$ \begin{array}{ l} \hline \underline{D}_1 \oplus \underline{D}_2 \\ \underline{A}_1 \multimap \underline{A}_2 \\ [\underline{D}]^p \\ !\underline{D} \\ (\underline{D}) \\ \hline \end{array} $	S
rec_type_bound, R	::=	name for recursive type
rec_type_def	$ \begin{array}{ l} \hline \mu X. \underline{D} \\ \hline \end{array} $	recursive type definition
type_affect, ta	$ \begin{array}{ l} \hline x : A \\ h^p : D \\ \hline \end{array} $	type affectation variable hole
type_affects	$ \begin{array}{ l} \hline ta \\ ta, type_affects \\ \hline \end{array} $	type affectations
typing_context, Γ , \mathcal{U} , H, H_P , H_N	$ \begin{array}{ l} \hline \emptyset \\ \{type_affects\} \\ \Gamma_1 \sqcup \Gamma_2 \\ \hline \end{array} $	typing context
command	$ \begin{array}{ l} \hline S \mid t \\ \hline \end{array} $	
terminals	$ \begin{array}{ l} \hline () \\ \mapsto \\ \multimap \\ := \\ \vdash \\ \sqcup \\ \emptyset \\ \neq \\ \in \\ \notin \\ \backslash n \\ \langle \\ \rangle \\ Inl \\ Inr \\ Ur \\ Dest \\ \triangleleft \\ \Downarrow \\ \text{fix} \\ \hline \end{array} $	

	\perp \bullet \subset \mathcal{N} \implies $@$ \wedge $;$	
formula	$::=$ \mid judgement	
Ctx	$::=$ $\mid x \in \mathcal{N}(\Gamma)$ $\mid x \notin \mathcal{N}(\Gamma)$ $\mid \text{type_affect} \in \Gamma$ $\mid \mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ $\mid p_1 = p_2 \implies \Gamma_1 = \Gamma_2$ $\mid p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$	Γ_1 and Γ_2 are disjoint typing contexts with n
Store	$::=$ $\mid \text{fresh } h$ $\mid \text{store_affect} \in \mathbb{S}$ $\mid x \notin \mathcal{N}(\mathbb{S})$	
Eq	$::=$ $\mid A_1 = A_2$ $\mid A_1 \neq A_2$ $\mid t = u$ $\mid \Gamma = D$	
Ty	$::=$ $\mid R \stackrel{\text{fix}}{=} \text{rec_type_def}$ $\mid \mathcal{U} ; \Gamma \vdash \text{command} : A$ $\mid H_N \vdash \underline{d}^p : D \vdash H_P$ $\mid \mathbb{S} \vdash H$ $\mid \mathcal{U} ; H \sqcup \Gamma \vdash t : A$	H_N stands for "needs", H_P for "provides"
Sem	$::=$ $\mid \text{command} \Downarrow \text{command}'$	
judgement	$::=$ $\mid \text{Ctx}$ $\mid \text{Store}$ $\mid \text{Eq}$ $\mid \text{Ty}$ $\mid \text{Sem}$	
user_syntax	$::=$	

- | metavariable
- | term
- | *hole*
- | val
- | data
- | *multiplicity*
- | sub
- | subs
- | data_with_hole
- | store_affect
- | store_affects
- | store
- | type
- | data_type
- | type_with_var
- | data_type_with_var
- | rec_type_bound
- | rec_type_def
- | type_affect
- | type_affects
- | typing_context
- | command
- | terminals

$x \in \mathcal{N}(\Gamma)$

$x \notin \mathcal{N}(\Gamma)$

$\text{type_affect} \in \Gamma$

$\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ Γ_1 and Γ_2 are disjoint typing contexts with no clashing variable names or labels

$p_1 = p_2 \implies \Gamma_1 = \Gamma_2$

$p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$

fresh *h*

$\text{store_affect} \in \mathbb{S}$

$x \notin \mathcal{N}(\mathbb{S})$

$A_1 = A_2$

$A_1 \neq A_2$

$t = u$

$\Gamma = D$

$R \stackrel{\text{fix}}{=} \text{rec_type_def}$

$\mathcal{U} ; \Gamma \vdash \text{command} : A$

$$\frac{\begin{array}{l} \mathbb{S} \vdash H \\ \mathcal{U} ; H \sqcup \Gamma \vdash t : A \end{array}}{\mathcal{U} ; \Gamma \vdash \mathbb{S} \mid t : A} \quad \text{TYCOMM_DEF}$$

$H_N \vdash \underline{d} \text{ } ^p \text{ } D \vdash H_P$ H_N stands for "needs", H_P for "provides"

$$\overline{\emptyset \vdash h \text{ } ^p \text{ } D \vdash \{h \text{ } ^p \text{ } D\}} \quad \text{TYDH_H}$$

$$\begin{array}{c}
\overline{\{h \textcolor{brown}{p}: D\} \vdash [h] \textcolor{brown}{t}: [D]^p \vdash \emptyset} \quad \text{TYDH_DEST} \\
\\
\overline{\emptyset \vdash \star \textcolor{blue}{p}: 1 \vdash \emptyset} \quad \text{TYDH_U} \\
\\
\frac{\emptyset ; H_N \sqcup \{x: A_1\} \vdash t: A_2 \quad p = \omega \implies H_N = \emptyset}{H_N \vdash \lambda x: D. t \textcolor{blue}{p}: A_1 \multimap A_2 \vdash \emptyset} \quad \text{TYDH_FN} \\
\\
\frac{H_N \vdash d \textcolor{blue}{\omega}: D \vdash H_P}{H_N \vdash \text{Ur } d \textcolor{blue}{p}: !D \vdash H_P} \quad \text{TYDH_E} \\
\\
\frac{H_N \vdash d \textcolor{blue}{p}: D_1 \vdash H_P}{H_N \vdash \text{Inl } d \textcolor{blue}{p}: D_1 \oplus D_2 \vdash H_P} \quad \text{TYDH_INL} \\
\\
\frac{H_N \vdash d \textcolor{blue}{p}: D_2 \vdash H_P}{H_N \vdash \text{Inr } d \textcolor{blue}{p}: D_1 \oplus D_2 \vdash H_P} \quad \text{TYDH_INR} \\
\\
\frac{H_{N_1} \vdash d_1 \textcolor{blue}{p}: D_1 \vdash H_{P_1} \quad H_{N_2} \vdash d_2 \textcolor{blue}{p}: D_2 \vdash H_{P_2}}{H_{N_1} \sqcup H_{N_2} \vdash \langle d_1, d_2 \rangle \textcolor{blue}{p}: D_1 \otimes D_2 \vdash H_{P_1} \sqcup H_{P_2}} \quad \text{TYDH_P} \\
\\
\frac{\begin{array}{c} R \stackrel{\text{fix}}{=} \mu X. D \\ H_N \vdash d \textcolor{blue}{p}: D[X := R] \vdash H_P \end{array}}{H_N \vdash @R d \textcolor{blue}{p}: R \vdash H_P} \quad \text{TYDH_R}
\end{array}$$

$$\boxed{S \vdash H}$$

$$\begin{array}{c}
\overline{\emptyset \vdash \emptyset} \quad \text{TYSTORE_EMPTY} \\
\\
\frac{S \vdash H \sqcup H_N \quad H_N \vdash d \textcolor{brown}{t}: D \vdash H_P}{S \sqcup \{x: D = d\} \vdash H \sqcup H_P} \quad \text{TYSTORE_ROOT}
\end{array}$$

$$\boxed{\bar{U} ; H \sqcup \Gamma \vdash t: A}$$

$$\begin{array}{c}
\overline{\bar{U} ; \emptyset \sqcup \emptyset \vdash \bullet: \perp} \quad \text{TYTERM_NoEFF} \\
\\
\overline{\bar{U} ; \{h \textcolor{brown}{p}: D\} \sqcup \emptyset \vdash [h]: [D]^p} \quad \text{TYTERM_DEST} \\
\\
\overline{\bar{U} ; \emptyset \sqcup \emptyset \vdash \star: 1} \quad \text{TYTERM_U} \\
\\
\frac{\emptyset ; H \sqcup \{x: A_1\} \vdash t: A_2}{\bar{U} ; H \sqcup \emptyset \vdash \lambda x: A_1. t \textcolor{blue}{p}: A_1 \multimap A_2} \quad \text{TYTERM_FN} \\
\\
\frac{\emptyset ; \emptyset \sqcup \emptyset \vdash d: D}{\bar{U} ; \emptyset \sqcup \emptyset \vdash \text{Ur } d: !D} \quad \text{TYTERM_E} \\
\\
\frac{\emptyset ; H \sqcup \emptyset \vdash d: D_1}{\bar{U} ; H \sqcup \emptyset \vdash \text{Inl } d: D_1 \oplus D_2} \quad \text{TYTERM_INL} \\
\\
\frac{\emptyset ; H \sqcup \emptyset \vdash d: D_2}{\bar{U} ; H \sqcup \emptyset \vdash \text{Inr } d: D_1 \oplus D_2} \quad \text{TYTERM_INR} \\
\\
\frac{\emptyset ; H_1 \sqcup \emptyset \vdash d_1: D_1 \quad \emptyset ; H_2 \sqcup \emptyset \vdash d_2: D_2}{\bar{U} ; H_1 \sqcup H_2 \sqcup \emptyset \vdash \langle d_1, d_2 \rangle: D_1 \otimes D_2} \quad \text{TYTERM_P}
\end{array}$$

$$\begin{array}{c}
\frac{R \stackrel{\text{fix}}{=} \mu X. \underline{D} \quad \emptyset ; H \sqcup \emptyset \vdash d : \underline{D}[X := R]}{\emptyset ; H \sqcup \emptyset \vdash @R d : R} \quad \text{TYTERM_R} \\
\\
\frac{x \notin \mathcal{N}(\mathcal{U})}{\mathcal{U} ; \emptyset \sqcup \{x : A\} \vdash x : A} \quad \text{TYTERM_ID} \\
\\
\frac{}{\mathcal{U} \sqcup \{x : A\} ; \emptyset \sqcup \emptyset \vdash x : A} \quad \text{TYTERM_ID'} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : A_1 \multimap A_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A_1}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash tu : A_2} \quad \text{TYTERM_APP} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \perp \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A_2}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t ; u : A_2} \quad \text{TYTERM_EFFSEQ} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : 1 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\star \mapsto u\} : A} \quad \text{TYTERM_PATU} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : !D \quad \mathcal{U} \sqcup \{x : D\} ; H_2 \sqcup \Gamma_2 \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} : A} \quad \text{TYTERM_PATE} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : D_1 \oplus D_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : D_1\} \vdash u_1 : A \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_2 : D_2\} \vdash u_2 : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : A} \quad \text{TYTERM_PATs} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : D_1 \otimes D_2 \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : D_1, x_2 : D_2\} \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : A} \quad \text{TYTERM_PATP} \\
\\
\frac{R \stackrel{\text{fix}}{=} \mu X. \underline{D} \quad \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : R \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \underline{D}[X := R]\} \vdash u : A}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{@R x \mapsto u\} : A} \quad \text{TYTERM_PATR} \\
\\
\frac{\mathcal{U} ; H \sqcup \Gamma \sqcup \{x : [\underline{D}]^I\} \vdash t : \perp}{\mathcal{U} ; H \sqcup \Gamma \vdash \text{alloc } x. t : \underline{D}} \quad \text{TYTERM_ALLOC} \\
\\
\frac{\mathcal{U} ; H \sqcup \Gamma \vdash t : [\underline{1}]^p}{\mathcal{U} ; H \sqcup \Gamma \vdash t \triangleleft^p \star : \perp} \quad \text{TYTERM_FILLU} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : [A_1 \multimap A_2]^p \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : A_1\} \vdash u : A_2 \quad p = \omega \implies (H_2 = \emptyset \wedge \Gamma_2 = \emptyset)}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \lambda x : A_1. u : \perp} \quad \text{TYTERM_FILLFN} \\
\\
\frac{\mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : [\underline{D}]^p \quad \mathcal{U} ; H_2 \sqcup \Gamma_2 \vdash u : \underline{D} \quad p = \omega \implies (H_2 = \emptyset \wedge \Gamma_2 = \emptyset)}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p u : \perp} \quad \text{TYTERM_FILLL}
\end{array}$$

$$\begin{array}{c}
\frac{\begin{array}{l} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D \rrbracket^\omega\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Ur } x.u : A} \text{TYTERM_FILLE} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \oplus D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D_1 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inl } x.u : A} \text{TYTERM_FILLINL} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \oplus D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D_2 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \text{Inr } x.u : A} \text{TYTERM_FILLINR} \\
\\
\frac{\begin{array}{l} \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket D_1 \otimes D_2 \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x_1 : \llbracket D_1 \rrbracket^p, x_2 : \llbracket D_2 \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p \langle x_1, x_2 \rangle . u : A} \text{TYTERM_FILLP} \\
\\
\frac{\begin{array}{l} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ \mathcal{U} ; H_1 \sqcup \Gamma_1 \vdash t : \llbracket R \rrbracket^p \\ \mathcal{U} ; H_2 \sqcup \Gamma_2 \sqcup \{x : \llbracket D[X := R] \rrbracket^p\} \vdash u : A \end{array}}{\mathcal{U} ; H_1 \sqcup H_2 \sqcup \Gamma_1 \sqcup \Gamma_2 \vdash t \triangleleft^p @R x.u : A} \text{TYTERM_FILLR}
\end{array}$$

command \Downarrow command'

$$\begin{array}{c}
\frac{}{\mathbb{S} \mid v \Downarrow \mathbb{S} \mid v} \text{SEMOP_VAL} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lambda x:A. t' \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \\ \mathbb{S}_2 \mid t'[x := v_2] \Downarrow \mathbb{S}_3 \mid v_3 \end{array}}{\mathbb{S}_0 \mid t u \Downarrow \mathbb{S}_3 \mid v_3} \text{SEMOP_APP} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \bullet \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid t ; u \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_EFFSEQ} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \star \\ \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\star \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATU} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Ur } d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Ur } y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATE} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inl } d \\ \mathbb{S}_1 \mid u_1[y_1 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINL} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \text{Inr } d \\ \mathbb{S}_1 \mid u_2[y_2 := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATINR} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \langle d_1, d_2 \rangle \\ \mathbb{S}_1 \mid u[y_1 := d_1, y_2 := d_2] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{\langle y_1, y_2 \rangle \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATP} \\
\\
\frac{\begin{array}{l} \mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid @R d \\ \mathbb{S}_1 \mid u[y := d] \Downarrow \mathbb{S}_2 \mid v_2 \end{array}}{\mathbb{S}_0 \mid \text{case } t \text{ of } \{@R y \mapsto u\} \Downarrow \mathbb{S}_2 \mid v_2} \text{SEMOP_PATR}
\end{array}$$

$$\begin{array}{c}
\text{fresh } h \\
\frac{\mathbb{S}_0 \sqcup \{x : D = h\} \mid t[x := \lfloor h \rfloor] \Downarrow \mathbb{S}_1 \sqcup \{x : D = d\} \mid \bullet}{\mathbb{S}_0 \mid \text{alloc } x . t \Downarrow \mathbb{S}_1 \mid d} \quad \text{SEMOp_ALLOC} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \star \Downarrow \mathbb{S}_1[h := \star] \mid \bullet} \quad \text{SEMOp_FILLU} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor}{\mathbb{S}_0 \mid t \triangleleft^p \lambda x : A . u \Downarrow \mathbb{S}_1[h := \lambda x : A . u] \mid \bullet} \quad \text{SEMOp_FILLFN} \\
\\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1 \mid u \Downarrow \mathbb{S}_2 \mid d_2}{\mathbb{S}_0 \mid t \triangleleft^p u \Downarrow \mathbb{S}_2[h := d_2] \mid \bullet} \quad \text{SEMOp_FILLL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Ur } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Ur } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOp_FILLE} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inl } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inl } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOp_FILLINL} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{Inr } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{Inr } y . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOp_FILLINR} \\
\\
\text{fresh } h_1 \\
\text{fresh } h_2 \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \langle h_1, h_2 \rangle] \mid u[y_1 := \lfloor h_1 \rfloor, y_2 := \lfloor h_2 \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \langle y_1, y_2 \rangle . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOp_FILLP} \\
\\
\text{fresh } h' \\
\frac{\mathbb{S}_0 \mid t \Downarrow \mathbb{S}_1 \mid \lfloor h \rfloor \quad \mathbb{S}_1[h := \text{OR } h'] \mid u[y := \lfloor h' \rfloor] \Downarrow \mathbb{S}_2 \mid v_2}{\mathbb{S}_0 \mid t \triangleleft^p \text{OR } x . u \Downarrow \mathbb{S}_2 \mid v_2} \quad \text{SEMOp_FILLR}
\end{array}$$

Definition rules: 57 good 0 bad
 Definition rule clauses: 153 good 0 bad