

metavariable, x , y

term, t , u	$::=$ <ul style="list-style-type: none"> v $t\ u$ $\text{case } t \text{ of } \{\star \mapsto u\}$ $\text{case } t \text{ of } \{\text{Ur } x \mapsto u\}$ $\text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\}$ $\text{case } t \text{ of } \{\text{@R } x \mapsto u\}$ $\text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}$ $\text{extract } t$ $\text{flip } t$ $\text{reassoc } t$ $\text{redL } t$ $\text{mapL } t \text{ with } u$ $\text{alloc } D$ $t \overset{p}{\prec} \star$ $t \overset{p}{\prec} \lambda x:A. u$ $t \overset{p}{\prec} u$ $t \overset{p}{\prec} \text{Ur}$ $t \overset{p}{\prec} \text{Inl}$ $t \overset{p}{\prec} \text{Inr}$ $t \overset{p}{\prec} \text{@R}$ $t \overset{p}{\prec} \langle, \rangle$ $t \overset{p}{\prec} \langle \odot \rangle$ (t) $t[\text{subs}]$ 	term <ul style="list-style-type: none"> value application pattern-matching on unit pattern-matching on unrestricted complete pattern-matching on sum pattern-matching on recursive data pattern-matching on product get data from a dest-filling statement fill destination with unit fill destination with function fill destination with value fill destination with exponential fill destination with sum variant 1 fill destination with sum variant 2 fill destination with recursive data fill destination with product
val, v	$::=$ <ul style="list-style-type: none"> \bullet d 	value (unreducible term) <ul style="list-style-type: none"> empty effect data structure
data, d	$::=$ <ul style="list-style-type: none"> $[x]$ x \star $\lambda x:A. t$ $\text{Ur } d$ $\text{C } d$ $\text{Inl } d$ $\text{Inr } d$ $\text{@R } d$ $\langle d_1, d_2 \rangle$ $\langle v_1 \odot v_2 \rangle$ (d) 	<ul style="list-style-type: none"> destination var or hole unit lambda abstraction exponential sum variant 1 sum variant 2 recursive data product mpar
<i>multiplicity</i> , p	$::=$ <ul style="list-style-type: none"> 1 ω 	multiplicity <ul style="list-style-type: none"> for holes/destinations not under a Ur for holes/destinations under a Ur
sub	$::=$ <ul style="list-style-type: none"> var_sub 	substitution

subs	$\begin{array}{l} ::= \\ \text{var_subs} \end{array}$	substitutions
var_sub	$\begin{array}{l} ::= \\ x := v \end{array}$	variable substitution
var_subs	$\begin{array}{l} ::= \\ \text{var_sub} \\ \text{var_sub}, \text{var_subs} \end{array}$	variable substitutions
type, A	$\begin{array}{l} ::= \\ \perp \\ D \end{array}$	bottom (effect) type
data_type, D	$\begin{array}{l} ::= \\ 1 \\ N \\ R \\ D_1 \otimes D_2 \\ A_1 \wp A_2 \\ D_1 \oplus D_2 \\ A_1 \multimap A_2 \\ [D]^p \\ \omega N \\ \epsilon D \\ (D) \\ \underline{D}[X := D] \end{array}$	unit type recursive type bound to a name product type sum type linear function type destination type exponential S M unroll a recursive data type
nodest_data_type, N	$::=$	Data type with no dest in its tree
type_with_var, \underline{A}	$\begin{array}{l} ::= \\ \perp \\ \underline{D} \end{array}$	
data_type_with_var, \underline{D}	$\begin{array}{l} ::= \\ X \\ 1 \\ \underline{N} \\ R \\ \underline{D}_1 \otimes \underline{D}_2 \\ \underline{D}_1 \oplus \underline{D}_2 \\ \underline{A}_1 \multimap \underline{A}_2 \\ \underline{A}_1 \wp \underline{A}_2 \\ [\underline{D}]^p \\ \omega \underline{N} \\ \epsilon \underline{D} \\ (\underline{D}) \end{array}$	S
nodest_data_type_with_var, \underline{N}	$::=$	

<code>rec_type_bound, R</code>	$::=$	name for recursive type
<code>rec_type_def</code>	$::=$ $\mu X. \underline{D}$	recursive type definition
<code>sign, s</code>	$::=$ $+$ $-$	sign
<code>type_affect, ta</code>	$::=$ $x : A$ $-x :^P D$	type affectation variable hole
<code>type_affects</code>	$::=$ ta $ta, type_affects$	type affectations
<code>typing_context, \mathcal{U}, Γ, Γ^-</code>	$::=$ \emptyset $\{type_affects\}$ $\Gamma_1 \sqcup \Gamma_2$ $\Gamma_1 \boxplus \Gamma_2$ $\Gamma_1 \boxdot \Gamma_2$	typing context
<code>terminals</code>	$::=$ \mapsto $\text{---} \circ$ $:=$ \vdash \sqcup \boxplus \boxdot \emptyset \neq \in \notin $\backslash n$ \langle \rangle Inl Inr Ur C $\text{c}^!$ $\text{e}^!$ γ Dest \triangleleft 	

	\Downarrow fix \equiv \perp \bullet \subset \mathcal{N} \implies $@$ \wedge $;$ \star \odot
formula	$::=$ \mid judgement
Ctx	$::=$ $\mid x \in \mathcal{N}(\Gamma)$ $\mid x \notin \mathcal{N}(\Gamma)$ $\mid \text{type_affect} \in \Gamma$ $\mid \mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ $\mid p_1 = p_2 \implies \Gamma_1 = \Gamma_2$ $\mid p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$ $\mid \text{fresh } x$
Eq	$::=$ $\mid A_1 = A_2$ $\mid A_1 \neq A_2$ $\mid t = u$ $\mid \Gamma = D$
Ty	$::=$ $\mid R \stackrel{\text{fix}}{=} \text{rec_type_def}$ $\mid \mathcal{U} ; \Gamma \vdash t : A$
Sem	$::=$ $\mid t \Downarrow t'$
judgement	$::=$ $\mid \text{Ctx}$ $\mid \text{Eq}$ $\mid \text{Ty}$ $\mid \text{Sem}$
user_syntax	$::=$ $\mid \text{metavariable}$ $\mid \text{term}$ $\mid \text{val}$ $\mid \text{data}$

Γ_1 and Γ_2 are disjoint typing contexts with n

- | *multiplicity*
- | sub
- | subs
- | var_sub
- | var_subs
- | type
- | data_type
- | nodest_data_type
- | type_with_var
- | data_type_with_var
- | nodest_data_type_with_var
- | rec_type_bound
- | rec_type_def
- | sign
- | type_affect
- | type_affects
- | typing_context
- | terminals

$x \in \mathcal{N}(\Gamma)$

$x \notin \mathcal{N}(\Gamma)$

$\text{type_affect} \in \Gamma$

$\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset$ Γ_1 and Γ_2 are disjoint typing contexts with no clashing variable names or labels

$p_1 = p_2 \implies \Gamma_1 = \Gamma_2$

$p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)$

fresh x

$A_1 = A_2$

$A_1 \neq A_2$

$t = u$

$\Gamma = D$

$R \stackrel{\text{fix}}{=} \text{rec_type_def}$

$\mathcal{U} ; \Gamma \vdash t : A$

$\frac{}{\mathcal{U} ; \emptyset \vdash \bullet : \perp}$ TyTERM_NoEff

$\frac{}{\mathcal{U} ; \emptyset \vdash \star : 1}$ TyTERM_U

$\frac{\emptyset ; \Gamma^- \sqsupset \{x : A_1\} \vdash t : A_2}{\mathcal{U} ; \Gamma^- \vdash \lambda x : A_1 . t : A_1 \multimap A_2}$ TyTERM_FN

$\frac{\emptyset ; \emptyset \vdash d : N}{\mathcal{U} ; \emptyset \vdash \text{Ur } d : \text{d}N}$ TyTERM_E

$\frac{\emptyset ; \Gamma^- \vdash d : D}{\mathcal{U} ; \Gamma^- \vdash \text{C } d : \text{c}D}$ TyTERM_C

$\frac{\emptyset ; \Gamma \vdash d : D_1}{\mathcal{U} ; \Gamma \vdash \text{Inl } d : D_1 \oplus D_2}$ TyTERM_INL

$$\begin{array}{c}
\frac{\emptyset ; \Gamma \vdash d : D_2}{\mathcal{U} ; \Gamma \vdash \text{Inr } d : D_1 \oplus D_2} \quad \text{TYTERM_INR} \\
\\
\frac{R \stackrel{\text{fix}}{=} \mu X. \underline{D} \quad \emptyset ; \Gamma \vdash d : \underline{D}[X := R]}{\mathcal{U} ; \Gamma \vdash @R d : R} \quad \text{TYTERM_R} \\
\\
\frac{\emptyset ; \Gamma_1 \vdash d_1 : D_1 \quad \emptyset ; \Gamma_2 \vdash d_2 : D_2}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \langle d_1, d_2 \rangle : D_1 \otimes D_2} \quad \text{TYTERM_P} \\
\\
\frac{\emptyset ; \Gamma_1 \vdash d_1 : D_1 \quad \emptyset ; \Gamma_2 \vdash d_2 : D_2}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \langle d_1 \odot d_2 \rangle : D_1 \wp D_2} \quad \text{TYTERM_M} \\
\\
\frac{}{\mathcal{U} ; \{ -x : ^p D \} \vdash [x] : [D]^p} \quad \text{TYTERM_D} \\
\\
\frac{}{\mathcal{U} ; \{ x : A \} \vdash x : A} \quad \text{TYTERM_VAR} \\
\\
\frac{}{\mathcal{U} \sqcup \{ x : A \} ; \emptyset \vdash x : A} \quad \text{TYTERM_VAR'} \\
\\
\frac{\mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}(A_1 \multimap A_2) \quad \mathcal{U} ; \Gamma_2 \vdash u : A_1}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash t u : A_2} \quad \text{TYTERM_APP} \\
\\
\frac{\mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}1 \quad \mathcal{U} ; \Gamma_2 \vdash u : A}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{ \star \mapsto u \} : A} \quad \text{TYTERM_PATU} \\
\\
\frac{\mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}(\omega N) \quad \mathcal{U} \sqcup \{ x : \mathcal{d}N \} ; \Gamma_2 \vdash u : A}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{ \text{Ur } x \mapsto u \} : A} \quad \text{TYTERM_PATE} \\
\\
\frac{\begin{array}{l} x_1 \notin \mathcal{N}(\Gamma_1) \\ x_2 \notin \mathcal{N}(\Gamma_2) \\ \mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}(D_1 \oplus D_2) \\ \mathcal{U} ; \Gamma_2 \sqcup \{ x_1 : \mathcal{d}D_1 \} \vdash u_1 : A \\ \mathcal{U} ; \Gamma_2 \sqcup \{ x_2 : \mathcal{d}D_2 \} \vdash u_2 : A \end{array}}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \} : A} \quad \text{TYTERM_PATs} \\
\\
\frac{\begin{array}{l} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ x \notin \mathcal{N}(\Gamma_1) \\ \mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}R \\ \mathcal{U} ; \Gamma_2 \sqcup \{ x : \mathcal{d}\underline{D}[X := R] \} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{ @R x \mapsto u \} : A} \quad \text{TYTERM_PATR} \\
\\
\frac{\begin{array}{l} x_1 \notin \mathcal{N}(\Gamma_1) \\ x_2 \notin \mathcal{N}(\Gamma_2) \\ \mathcal{U} ; \Gamma_1 \vdash t : \mathcal{d}(D_1 \otimes D_2) \\ \mathcal{U} ; \Gamma_2 \sqcup \{ x_1 : \mathcal{d}D_1, x_2 : \mathcal{d}D_2 \} \vdash u : A \end{array}}{\mathcal{U} ; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{ \langle x_1, x_2 \rangle \mapsto u \} : A} \quad \text{TYTERM_PATP} \\
\\
\frac{\mathcal{U} ; \Gamma \vdash t : \mathcal{d}D}{\mathcal{U} ; \Gamma \vdash \text{extract } t : D} \quad \text{TYTERM_EX}
\end{array}$$

$$\begin{array}{c}
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(A_1 \wp A_2)}{\mathcal{U}; \Gamma \vdash \text{flip } t : \mathcal{d}(A_2 \wp A_1)} \quad \text{TYTERM_FLIPM} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(A_1 \wp \mathcal{d}(A_2 \wp A_3))}{\mathcal{U}; \Gamma \vdash \text{reassoc } t : \mathcal{d}(\mathcal{d}(A_1 \wp A_2) \wp A_3)} \quad \text{TYTERM_REASSOCM} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(\perp \wp D)}{\mathcal{U}; \Gamma \vdash \text{redL } t : \mathcal{d}D} \quad \text{TYTERM_REDRM} \\
\\
\frac{\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(A_1 \wp A_2) \quad \mathcal{U}; \Gamma_2 \vdash u : \mathcal{d}(A_1 \multimap A_3)}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{mapL } t \text{ with } u : A_3 \wp A_2} \quad \text{TYTERM_MAPRM} \\
\\
\frac{}{\mathcal{U}; \emptyset \vdash \text{alloc } D : \mathcal{d}(\mathcal{d}[D]^I \wp D)} \quad \text{TYTERM_ALLOC} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[1]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \star : \perp} \quad \text{TYTERM_FILLU} \\
\\
\frac{\begin{array}{l} \mathbf{x} \notin \mathcal{N}(\Gamma_1) \\ \mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}[A_1 \multimap A_2]^p \\ \mathcal{U}; \Gamma_2 \sqcup \{\mathbf{x} : A_1\} \vdash u : A_2 \\ p = \omega \implies \Gamma_2 = \emptyset \end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash t \prec^p \lambda \mathbf{x} : A_1 . u : \perp} \quad \text{TYTERM_FILLFN} \\
\\
\frac{\begin{array}{l} \mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}[D]^p \\ \mathcal{U}; \Gamma_2 \vdash u : D \\ p = \omega \implies \Gamma_2 = \emptyset \end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash t \prec^p u : \perp} \quad \text{TYTERM_FILLL} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[\omega N]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \text{Ur} : \mathcal{d}[N]^\omega} \quad \text{TYTERM_FILLE} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \oplus D_2]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \text{Inl} : \mathcal{d}[D_1]^p} \quad \text{TYTERM_FILLINL} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \oplus D_2]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \text{Inr} : \mathcal{d}[D_2]^p} \quad \text{TYTERM_FILLINR} \\
\\
\frac{\begin{array}{l} R \stackrel{\text{fix}}{=} \mu X . \underline{D} \\ \mathcal{U}; \Gamma \vdash t : \mathcal{d}[R]^p \end{array}}{\mathcal{U}; \Gamma \vdash t \prec^p @R : \mathcal{d}[\underline{D}[X := R]]^p} \quad \text{TYTERM_FILLR} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \otimes D_2]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \langle, \rangle : \mathcal{d}(\mathcal{d}[D_1]^p \wp \mathcal{d}[D_2]^p)} \quad \text{TYTERM_FILLP} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[\mathcal{d}D_1 \wp \mathcal{d}D_2]^p}{\mathcal{U}; \Gamma \vdash t \prec^p \langle \odot \rangle : \mathcal{d}([D_1]^p \otimes [D_2]^p)} \quad \text{TYTERM_FILLM}
\end{array}$$

$$\boxed{t \Downarrow t'}$$

$$\frac{}{v \Downarrow v} \quad \text{SEMOP_VAL}$$

$$\begin{array}{c}
\frac{
\begin{array}{c}
t \Downarrow C(\lambda x:A. t') \\
u \Downarrow v_2 \\
t'[x := v_2] \Downarrow v_3
\end{array}
}{t u \Downarrow v_3} \text{ SEMOP_APP} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C \star \\
u \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{\star \mapsto u\} \Downarrow v_2} \text{ SEMOP_PATU} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C(\text{Ur } d) \\
u[y := C d] \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{\text{Ur } y \mapsto u\} \Downarrow v_2} \text{ SEMOP_PATE} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C(\text{Inl } d) \\
u_1[y_1 := C d] \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow v_2} \text{ SEMOP_PATINL} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C(\text{Inr } d) \\
u_2[y_2 := C d] \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{\text{Inl } y_1 \mapsto u_1, \text{Inr } y_2 \mapsto u_2\} \Downarrow v_2} \text{ SEMOP_PATINR} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C(@R d) \\
u[y := C d] \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{@R y \mapsto u\} \Downarrow v_2} \text{ SEMOP_PATR} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C \langle d_1, d_2 \rangle \\
u[y_1 := C d_1, y_2 := C d_2] \Downarrow v_2
\end{array}
}{\text{case } t \text{ of } \{\langle y_1, y_2 \rangle \mapsto u\} \Downarrow v_2} \text{ SEMOP_PATP} \\
\\
\frac{
t \Downarrow C d
}{\text{extract } t \Downarrow d} \text{ SEMOP_EX} \\
\\
\frac{
t \Downarrow C \langle v_1 \odot v_2 \rangle
}{\text{flip } t \Downarrow C \langle v_2 \odot v_1 \rangle} \text{ SEMOP_FLIPM} \\
\\
\frac{
t \Downarrow C \langle v_1 \odot C \langle v_2 \odot v_3 \rangle \rangle
}{\text{reassoc } t \Downarrow C \langle C \langle v_1 \odot v_2 \rangle \odot v_3 \rangle} \text{ SEMOP_REASSOCM} \\
\\
\frac{
t \Downarrow C \langle \bullet \odot d \rangle
}{\text{redL } t \Downarrow C d} \text{ SEMOP_REDLM} \\
\\
\frac{
\begin{array}{c}
t \Downarrow C \langle v_1 \odot v_2 \rangle \\
u v_1 \Downarrow v_3
\end{array}
}{\text{mapL } t \text{ with } u \Downarrow C \langle v_3 \odot v_2 \rangle} \text{ SEMOP_MAPLM} \\
\\
\frac{
\text{fresh } x
}{\text{alloc } D \Downarrow C \langle C [x] \odot x \rangle} \text{ SEMOP_ALLOC}
\end{array}$$

Definition rules: 48 good 0 bad
 Definition rule clauses: 121 good 0 bad