

Destination λ -calculus

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1 Term and value syntax

var, x, y Term-level variable name
 k Index for ranges

hdn, h	$::=$ $ $ 1 $ $ 2 $ $ $h+h'$ $ $ $\max(H)$	Hole or destination name Sum Maximum of a set of holes
hdns, H	$::=$ $ $ $\{h_1, \dots, h_k\}$ $ $ $H_1 \cup H_2$ $ $ $H \pm h$ $ $ $\text{hnames}(\Gamma)$	Set of hole or destination names Set of holes Union of sets Increase all names from H by h . Hole names of a context (requires $\text{ctx_NoVar}(\Gamma)$)
val, v	$::=$ $ $ $-h$ $ $ $+h$ $ $ $()$ $ $ $\lambda_v x \mapsto t$ $ $ $\text{Inl } v$ $ $ $\text{Inr } v$ $ $ $E^m v$ $ $ (v_1, v_2) $ $ $H(v_1, v_2)$ $ $ $v \pm h$	Term value Hole Destination Unit Lambda abstraction Left variant for sum Right variant for sum Exponential Product Ampar Rename hole names inside v by shifting them by h
term, t, u	$::=$ $ $ v $ $ x $ $ $t \succ u$ $ $ $t ; u$ $ $ $t \succ \text{case } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}$ $ $ $t \succ \text{case } (x_1, x_2) \mapsto u$ $ $ $t \succ \text{case } E^m x \mapsto u$ $ $ $t \succ \text{map } x \mapsto u$ $ $ $\text{to}_{\times} t$ $ $ $\text{from}_{\times} t$ $ $ alloc_{\top} $ $ $t \triangleleft ()$ $ $ $t \triangleleft \lambda x \mapsto u$ $ $ $t \triangleleft \text{Inl}$ $ $ $t \triangleleft \text{Inr}$ $ $ $t \triangleleft (,)$ $ $ $t \triangleleft E^m$ $ $ $t \triangleleft \bullet u$	Term Value Variable Application Pattern-match on unit Pattern-match on sum Pattern-match on product Pattern-match on exponential Map over the right side of the ampar Wrap t into a trivial ampar Extract value from trivial ampar Return a fresh "identity" ampar object Fill destination with unit Fill destination with function Fill destination with left variant Fill destination with right variant Fill destination with product constructor Fill destination with exponential constructor Fill destination with root of ampar u

		$t[x := v]$	
eterm, j	::=		Pseudo-term
		t	
		$C[j]$	
ectx, C	::=		Evaluation context
		$[]$	Identity
		${}^o_{\mathbf{h}}(v_1, C$	Open ampar
		$C[C']$	Compose evaluation contexts
		$C[\mathbf{h} :=_{\mathbf{h}} v]$	Fill \mathbf{h} with value v (that may contain holes)

2 Type system

type, T, U	$::=$	$\mathbf{1}$ $\mathsf{T}_1 \oplus \mathsf{T}_2$ $\mathsf{T}_1 \otimes \mathsf{T}_2$ $!^m \mathsf{T}$ $\mathsf{T}_1 \ltimes \mathsf{T}_2$ $\mathsf{T}_1 \xrightarrow{m_1} \mathsf{T}_2$ $[\mathsf{T}]^m$ $\mathsf{T}_1 \multimap \mathsf{T}_2$	Type Unit Sum Product Exponential Ampar type (consuming T_2 yields T_1) Function Destination Evaluation contexts
mode, m, n	$::=$	pa ω $m_1 \cdot \dots \cdot m_k$	Mode (Semiring) Pair of a multiplicity and age Error case (incompatible types, multiplicities, or ages) Semiring product
mul, p	$::=$	1 ω $p_1 \cdot \dots \cdot p_k$	Multiplicity (first component of modality) Linear. Neutral element of the product Non-linear. Absorbing for the product Semiring product
age, a	$::=$	ν \uparrow ∞ $a_1 \cdot \dots \cdot a_k$	Age (second component of modality) Born now. Neutral element of the product One scope older Infinitely old / static. Absorbing for the product Semiring product
ctx, Γ, Δ	$::=$	$\{\mathsf{b}_1, \dots, \mathsf{b}_k\}$ $m \cdot \Gamma$ $\Gamma_1 \uplus \Gamma_2$ $-\Gamma$	Typing context List of bindings Multiply each binding by m Sum contexts Γ_1 and Γ_2 . Duplicates/incompatible elements will give bindings with modality ω Transforms every dest binding into a hole binding (requires <code>ctx_DestOnly</code> Γ)
bndr, b	$::=$	$x :_m \mathsf{T}$ $+h :_m [\mathsf{T}]^n$ $-h : \mathsf{T}^n$	Type assignment to either variable, destination or hole Variable Destination (m is its own modality; n is the modality for values it accepts) Hole (n is the modality for values it accepts, it doesn't have a modality on its own)

$$\boxed{\Gamma \Vdash v : \mathbf{T}}$$

(Typing of values (raw))

TYR-VAL-H

$$\frac{}{\{-h : \mathbf{T}^{fa}\} \Vdash -h : \mathbf{T}}$$

TYR-VAL-D

$$\frac{\text{ctx_Compatible } \Gamma \text{ } +h :_{lv} [\mathbf{T}]^n}{\Gamma \Vdash +h : [\mathbf{T}]^n}$$

TYR-VAL-U

$$\frac{}{\{\} \Vdash () : \mathbf{1}}$$

TYR-VAL-F

$$\frac{\text{ctx_DestOnly } \Gamma \quad \Gamma \uplus \{x :_m \mathbf{T}_1\} \vdash t : \mathbf{T}_2}{\Gamma \Vdash \lambda_v x \mapsto t : \mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2}$$

TYR-VAL-L

$$\frac{\Gamma \Vdash v : \mathbf{T}_1}{\Gamma \Vdash \text{Inl } v : \mathbf{T}_1 \oplus \mathbf{T}_2}$$

TYR-VAL-R

$$\frac{\Gamma \Vdash v : \mathbf{T}_2}{\Gamma \Vdash \text{Inr } v : \mathbf{T}_1 \oplus \mathbf{T}_2}$$

TYR-VAL-P

$$\frac{\Gamma_1 \Vdash v_1 : \mathbf{T}_1 \quad \Gamma_2 \Vdash v_2 : \mathbf{T}_2}{\Gamma_1 \uplus \Gamma_2 \Vdash (v_1, v_2) : \mathbf{T}_1 \otimes \mathbf{T}_2}$$

TYR-VAL-E

$$\frac{\Gamma \Vdash v : \mathbf{T}}{m \cdot \Gamma \Vdash E^m v : !^m \mathbf{T}}$$

TYR-VAL-A

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_1 \Gamma_2 \\ \text{ctx_DestOnly } \Gamma_2 \uplus \Gamma_3 \\ \text{ctx_DestOnly } \Gamma_1 \\ \Gamma_1 \uplus (-\Gamma_3) \Vdash v_1 : \mathbf{T}_1 \\ \Gamma_2 \uplus \Gamma_3 \Vdash v_2 : \mathbf{T}_2 \end{array}}{\Gamma_1 \uplus \Gamma_2 \Vdash \text{hnames}(-\Gamma_3) \langle v_1, v_2 \rangle : \mathbf{T}_1 \ltimes \mathbf{T}_2}$$

$$\boxed{\Gamma \vdash j : \mathbf{T}}$$

(Typing of terms)

TY-TERM-VAL

$$\frac{\text{ctx_NoHole } \Gamma \quad \Gamma \Vdash v : \mathbf{T}}{\Gamma \vdash v : \mathbf{T}}$$

TY-TERM-VAR

$$\frac{\text{ctx_Compatible } \Gamma \text{ } x :_{lv} \mathbf{T}}{\Gamma \vdash x : \mathbf{T}}$$

TY-TERM-APP

$$\frac{\Gamma_1 \vdash t : \mathbf{T}_1 \quad \Gamma_2 \vdash u : \mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2}{m \cdot \Gamma_1 \uplus \Gamma_2 \vdash t \succ u : \mathbf{T}_2}$$

TY-TERM-PATS

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_2 \{x_1 :_m \mathbf{T}_1\} \\ \text{ctx_Disjoint } \Gamma_2 \{x_2 :_m \mathbf{T}_2\} \\ \Gamma_1 \vdash t : \mathbf{T}_1 \oplus \mathbf{T}_2 \\ \Gamma_2 \uplus \{x_1 :_m \mathbf{T}_1\} \vdash u_1 : \mathbf{U} \\ \Gamma_2 \uplus \{x_2 :_m \mathbf{T}_2\} \vdash u_2 : \mathbf{U} \end{array}}{m \cdot \Gamma_1 \uplus \Gamma_2 \vdash t \succ \text{case } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : \mathbf{U}}$$

TY-TERM-PATU

$$\frac{\Gamma_1 \vdash t : \mathbf{1} \quad \Gamma_2 \vdash u : \mathbf{U}}{\Gamma_1 \uplus \Gamma_2 \vdash t ; u : \mathbf{U}}$$

TY-TERM-PATP

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_2 \{x_1 :_m \mathbf{T}_1\} \\ \text{ctx_Disjoint } \Gamma_2 \{x_2 :_m \mathbf{T}_2\} \\ \text{ctx_Disjoint } \{x_1 :_m \mathbf{T}_1\} \{x_2 :_m \mathbf{T}_2\} \\ \Gamma_1 \vdash t : \mathbf{T}_1 \otimes \mathbf{T}_2 \\ \Gamma_2 \uplus \{x_1 :_m \mathbf{T}_1, x_2 :_m \mathbf{T}_2\} \vdash u : \mathbf{U} \end{array}}{m \cdot \Gamma_1 \uplus \Gamma_2 \vdash t \succ \text{case } (x_1, x_2) \mapsto u : \mathbf{U}}$$

TY-TERM-PATE

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_2 \{x :_{m \cdot n} \mathbf{T}\} \\ \Gamma_1 \vdash t : !^n \mathbf{T} \\ \Gamma_2 \uplus \{x :_{m \cdot n} \mathbf{T}\} \vdash u : \mathbf{U} \end{array}}{m \cdot \Gamma_1 \uplus \Gamma_2 \vdash t \succ \text{case } E^m x \mapsto u : \mathbf{U}}$$

TY-TERM-MAP

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_2 \{x :_{lv} \mathbf{T}_2\} \\ \Gamma_1 \vdash t : \mathbf{T}_1 \ltimes \mathbf{T}_2 \\ \uparrow \cdot \Gamma_2 \uplus \{x :_{lv} \mathbf{T}_2\} \vdash u : \mathbf{U} \end{array}}{\Gamma_1 \uplus \Gamma_2 \vdash t \succ \text{map } x \mapsto u : \mathbf{T}_1 \ltimes \mathbf{U}}$$

TY-TERM-FILLF

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_2 \{x :_m \mathbf{T}_1\} \\ \Gamma_1 \vdash t : [\mathbf{T}_1 \xrightarrow{m} \mathbf{T}_2]^n \\ \Gamma_2 \uplus \{x :_m \mathbf{T}_1\} \vdash u : \mathbf{T}_2 \end{array}}{\Gamma_1 \uplus (\uparrow \cdot n) \cdot \Gamma_2 \vdash t \triangleleft \lambda x \mapsto u : \mathbf{1}}$$

TY-TERM-FILLL

$$\frac{\Gamma \vdash t : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n}{\Gamma \vdash t \triangleleft \text{Inl} : [\mathbf{T}_1]^n}$$

TY-TERM-FILLR

$$\frac{\Gamma \vdash t : [\mathbf{T}_1 \oplus \mathbf{T}_2]^n}{\Gamma \vdash t \triangleleft \text{Inr} : [\mathbf{T}_2]^n}$$

TY-TERM-FILLU

$$\frac{\Gamma \vdash t : [\mathbf{1}]^n}{\Gamma \vdash t \triangleleft () : \mathbf{1}}$$

TY-TERM-FILLP

$$\frac{\Gamma \vdash t : [\mathbf{T}_1 \otimes \mathbf{T}_2]^n}{\Gamma \vdash t \triangleleft (,) : [\mathbf{T}_1]^n \otimes [\mathbf{T}_2]^n}$$

TY-TERM-FILLE

$$\frac{\Gamma \vdash t : [!^m \mathbf{T}]^n}{\Gamma \vdash t \triangleleft E^m : [\mathbf{T}]^{m \cdot n}}$$

TY-TERM-FILLC

$$\frac{\Gamma_1 \vdash t : [\mathbf{T}_1]^n \quad \Gamma_2 \vdash u : \mathbf{T}_1 \ltimes \mathbf{T}_2}{\Gamma_1 \uplus (\uparrow \cdot n) \cdot \Gamma_2 \vdash t \triangleleft \bullet u : \mathbf{T}_2}$$

TY-TERM-ALLOC

$$\frac{}{\{\} \vdash \text{alloc}_{\mathbf{T}} : \mathbf{T} \ltimes [\mathbf{T}]^{lv}}$$

TY-TERM-TOA

$$\frac{\Gamma \vdash t : \mathbf{T}}{\Gamma \vdash \text{to}_{\ltimes} t : \mathbf{T} \ltimes \mathbf{1}}$$

TY-TERM-FROMA

$$\frac{\Gamma \vdash t : \mathbf{T} \ltimes \mathbf{1}}{\Gamma \vdash \text{from}_{\ltimes} t : \mathbf{T}}$$

$$\boxed{\Gamma \Vdash C : \mathbf{T}}$$

(Typing of evaluation contexts)

TYR-ECTX-T

$$\frac{\begin{array}{c} \text{ctx_Disjoint } \Gamma_1 \Gamma_2 \\ \text{ctx_Disjoint } \Gamma_2 \Gamma_3 \\ \text{ctx_NoVar } \Gamma_2 \uplus \Gamma_3 \\ \text{ctx_NoVar } \Gamma_1 \\ \Gamma_2 \uplus \Gamma_3 \vdash t : \mathbf{T}_1 \quad \Gamma_1 \vdash C[t] : \mathbf{T}_2 \end{array}}{\Gamma_1 \uplus (-\Gamma_2) \Vdash C : \mathbf{T}_1 \xrightarrow{\triangleright} \mathbf{T}_2}$$

3 Effects and big-step semantics

$$\boxed{j \xrightarrow{h}{}_{h'} j'}$$

(Small-step evaluation of terms using evaluation contexts)

SEM-ETERM-APP

$$\frac{}{C[v \succ (\lambda_v x \mapsto t)] \xrightarrow{h}{}_h C[t[x := v]]}$$

SEM-ETERM-PATU

$$\frac{}{C[(\) ; t_2] \xrightarrow{h}{}_h C[t_2]}$$

SEM-ETERM-PATL

$$\frac{}{C[(\text{Inl } v) \succ \text{case } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \xrightarrow{h}{}_h C[u_1[x := v]]}$$

SEM-ETERM-PATR

$$\frac{}{C[(\text{Inr } v) \succ \text{case } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \}] \xrightarrow{h}{}_h C[u_2[x := v]]}$$

SEM-ETERM-PATP

$$\frac{}{C[(v_1, v_2) \succ \text{case } (x_1, x_2) \mapsto u] \xrightarrow{h}{}_h C[u[x_1 := v_1][x_2 := v_2]]}$$

SEM-ETERM-PATE

$$\frac{}{C[E^n v \succ \text{case } E^m x \mapsto u] \xrightarrow{h}{}_h C[u[x := v]]}$$

SEM-ETERM-MAPOPEN

$$\frac{}{C[_H \langle v_1, v_2 \rangle \succ \text{map } x \mapsto u] \xrightarrow{h}{}_{\max(H)+h} C[_{H \pm h}^o \langle v_1 \overset{\pm}{h}, u[x := v_2 \overset{\pm}{h}] \rangle]}$$

SEM-ETERM-MAPCLOSE

$$\frac{}{C[_H^o \langle v_1, v_2 \rangle] \xrightarrow{h}{}_h C[_H \langle v_1, v_2 \rangle]}$$

SEM-ETERM-ALLOC

$$\frac{}{\text{alloc}_T \xrightarrow{h}{}_h \{1\} \langle +1, -1 \rangle}$$

SEM-ETERM-TOA

$$\frac{}{C[\text{to}_\times v] \xrightarrow{h}{}_h C[\{ \} \langle v, () \rangle]}$$

SEM-ETERM-FROMA

$$\frac{}{C[\text{from}_\times \{ \} \langle v, () \rangle] \xrightarrow{h}{}_h v}$$

SEM-ETERM-FILLU

$$\frac{}{C[+h \triangleleft ()] \xrightarrow{h}{}_h C[h := \{ \} ()]()]}$$

SEM-ETERM-FILLL

$$\frac{}{C[+h \triangleleft \text{Inl}] \xrightarrow{h}{}_{h+1} C[h := \{h+1\} \text{Inl} - (h+1)][+(h+1)]}$$

SEM-ETERM-FILLR

$$\frac{}{C[+h \triangleleft \text{Inr}] \xrightarrow{h}{}_{h+1} C[h := \{h+1\} \text{Inr} - (h+1)][+(h+1)]}$$

SEM-ETERM-FILLE

$$\frac{}{C[+h \triangleleft E^m] \xrightarrow{h}{}_{h+1} C[h := \{h+1\} E^m - (h+1)][+(h+1)]}$$

SEM-ETERM-FILLP

$$\frac{}{C[+h \triangleleft (,)] \xrightarrow{h}{}_{h+2} C[h := \{h+1, h+2\} (- (h+1), - (h+2))][+(h+1), +(h+2)]}$$

SEM-ETERM-FILLC

$$\frac{}{C[+h \triangleleft \bullet \langle v_1, v_2 \rangle] \xrightarrow{h}{}_{\max(H)+h} C[h := \binom{H \pm h}{(H \pm h)} v_1 \overset{\pm}{h}][v_2 \overset{\pm}{h}]}$$