

metavariable,  $x$ ,  $y$

term, $t$ , $u$	$::=$ <ul style="list-style-type: none"> <li><math>v</math></li> <li><math>t\ u</math></li> <li><math>\text{case } t \text{ of } \{\star \mapsto u\}</math></li> <li><math>\text{case } t \text{ of } \{\text{Ur } x \mapsto u\}</math></li> <li><math>\text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\}</math></li> <li><math>\text{case } t \text{ of } \{\text{@R } x \mapsto u\}</math></li> <li><math>\text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\}</math></li> <li><math>\text{extract } t</math></li> <li><math>\text{flip } t</math></li> <li><math>\text{reassoc } t</math></li> <li><math>\text{redL } t</math></li> <li><math>\text{mapL } t \text{ with } u</math></li> <li><math>\text{alloc } D</math></li> <li><math>t \triangleleft^p u</math></li> <li><math>t \triangleleft^p \star</math></li> <li><math>t \triangleleft^p \lambda x:A. u</math></li> <li><math>t \triangleleft^p \text{Ur}</math></li> <li><math>t \triangleleft^p \text{Inl}</math></li> <li><math>t \triangleleft^p \text{Inr}</math></li> <li><math>t \triangleleft^p \text{@R}</math></li> <li><math>t \triangleleft^p \langle, \rangle</math></li> <li><math>t \triangleleft^p \langle \odot \rangle</math></li> <li><math>(t)</math></li> <li><math>t[e]</math></li> <li><math>E[t]</math></li> </ul>	term <ul style="list-style-type: none"> <li>value</li> <li>application</li> <li>pattern-matching on unit</li> <li>pattern-matching on unrestricted complete</li> <li>pattern-matching on sum</li> <li>pattern-matching on recursive data</li> <li>pattern-matching on product</li> <li>remove C wrapper</li> <li>flip mpar sides</li> <li>reassociated nested mpar</li> <li>extract right when left is bottom</li> <li>map function on left side</li> <li>generate new hole and associated dest</li> <li>move into destination</li> <li>fill destination with unit</li> <li>fill destination with function</li> <li>fill destination with exponential</li> <li>fill destination with sum variant 1</li> <li>fill destination with sum variant 2</li> <li>fill destination with recursive type label</li> <li>fill destination with product</li> <li>fill destination with data mpar</li> </ul>
value, $v$	$::=$ <ul style="list-style-type: none"> <li><math>\bullet</math></li> <li><math>d</math></li> </ul>	value (unreducible term) <ul style="list-style-type: none"> <li>data structure</li> </ul>
<i>effect</i> , $e$	$::=$ <ul style="list-style-type: none"> <li><math>\varepsilon</math></li> <li>subs</li> </ul>	empty effect
data_value, $d$	$::=$ <ul style="list-style-type: none"> <li><math>[x]</math></li> <li><math>x</math></li> <li><math>\star</math></li> <li><math>\lambda x:A. t</math></li> <li><math>\text{Ur } d</math></li> <li><math>\text{C } d</math></li> <li><math>\text{Inl } d</math></li> <li><math>\text{Inr } d</math></li> <li><math>\text{@R } d</math></li> <li><math>\langle d_1, d_2 \rangle</math></li> <li><math>\langle t_1 \odot t_2 \rangle</math></li> <li><math>(d)</math></li> </ul>	destination var or hole unit lambda abstraction exponential  sum variant 1 sum variant 2 recursive data product mpar
<i>multiplicity</i> , $p$	$::=$	multiplicity

	$\begin{array}{ l} 1 \\ \omega \end{array}$	for holes/destinations not under a $\text{Ur}$ for holes/destinations under a $\text{Ur}$
sub	$\begin{array}{ l} ::= \\   \text{var\_sub} \end{array}$	substitution
subs	$\begin{array}{ l} ::= \\   \text{var\_subs} \end{array}$	substitutions
var_sub	$\begin{array}{ l} ::= \\   x := v \end{array}$	variable substitution
var_subs	$\begin{array}{ l} ::= \\   \text{var\_sub} \\   \text{var\_sub}, \text{var\_subs} \end{array}$	variable substitutions
type, $A$	$\begin{array}{ l} ::= \\   \perp \\   D \end{array}$	bottom (effect) type
data_type, $D$	$\begin{array}{ l} ::= \\   1 \\   N \\   R \\   D_1 \otimes D_2 \\   A_1 \wp A_2 \\   D_1 \oplus D_2 \\   A_1 \multimap A_2 \\   [D]^p \\   \omega N \\   \epsilon D \\   (D) \\   \underline{D}[X := D] \end{array}$	unit type  recursive type bound to a name product type  sum type linear function type destination type exponential  $S$ $M$ unroll a recursive data type
nodest_data_type, $N$	$\begin{array}{ l} ::= \end{array}$	Data type with no dest in its tree
type_with_var, $\underline{A}$	$\begin{array}{ l} ::= \\   \perp \\   \underline{D} \end{array}$	
data_type_with_var, $\underline{D}$	$\begin{array}{ l} ::= \\   X \\   1 \\   \underline{N} \\   R \\   \underline{D}_1 \otimes \underline{D}_2 \\   \underline{D}_1 \oplus \underline{D}_2 \\   \underline{A}_1 \multimap \underline{A}_2 \\   \underline{A}_1 \wp \underline{A}_2 \end{array}$	

	$ \begin{array}{ l} [D]^p \\ \omega N \\ \epsilon D \\ (D) \end{array} $	S
nodest_data_type_with_var, $\underline{N}$	::=	
rec_type_bound, R	::=	name for recursive type
rec_type_def	::=	recursive type definition
	$ \begin{array}{ l} \mu X. \underline{D} \end{array} $	
sign, s	::=	sign
	$ \begin{array}{ l} + \\ - \end{array} $	
type_affect, ta	::=	type affectation
	$ \begin{array}{ l} x : A \\ -x :^p D \end{array} $	variable hole
type_affects	::=	type affectations
	$ \begin{array}{ l} ta \\ ta, type\_affects \end{array} $	
typing_context, $\mathcal{U}, \Gamma, \Gamma^-$	::=	typing context
	$ \begin{array}{ l} \emptyset \\ \{type\_affects\} \\ \Gamma_1 \sqcup \Gamma_2 \\ \Gamma_1 \boxplus \Gamma_2 \\ \Gamma_1 \boxdot \Gamma_2 \end{array} $	
evaluation_context, $E$	::=	evaluation context
	$ \begin{array}{ l} [] \\ \langle E \odot t_2 \rangle \\ \langle v_1 \odot E \rangle \\ \text{case } E \text{ of } \{ \star \mapsto u \} \\ \text{case } E \text{ of } \{ \text{Ur } x \mapsto u \} \\ \text{case } E \text{ of } \{ \text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2 \} \\ \text{case } E \text{ of } \{ @R x \mapsto u \} \\ \text{case } E \text{ of } \{ \langle x_1, x_2 \rangle \mapsto u \} \\ \text{extract } E \\ \text{flip } E \\ \text{reassoc } E \\ \text{redL } E \\ \text{mapL } E \text{ with } t_2 \\ \text{mapL } v_1 \text{ with } E \\ E \triangleleft^p \star \\ E \triangleleft^p \lambda x:A. u \\ E \triangleleft^p u \end{array} $	

		$E \triangleleft^p \mathbf{u}$ $E \triangleleft^p \mathbf{Ur}$ $E \triangleleft^p \mathbf{Inl}$ $E \triangleleft^p \mathbf{Inr}$ $E \triangleleft^p @R$ $E \triangleleft^p \langle , \rangle$ $E \triangleleft^p \langle \odot \rangle$
terminals	::=	$\mapsto$ $\text{---}\circ$ $:=$ $\vdash$ $\sqcup$ $\boxplus$ $\sqcap$ $\emptyset$ $\neq$ $\in$ $\notin$ $\backslash n$ $\langle$ $\rangle$ $\mathbf{Inl}$ $\mathbf{Inr}$ $\mathbf{Ur}$ $\mathbf{C}$ $\text{!}$ $\omega \text{!}$ $\gamma$ $\mathbf{Dest}$ $\triangleleft$ $ $ $\Downarrow$ $\text{fix}$ $\equiv$ $\perp$ $\bullet$ $\subset$ $\mathcal{N}$ $\Rightarrow$ $@$ $\wedge$ $;$ $\star$ $\odot$ $\longrightarrow$ $\rightsquigarrow$
formula	::=	judgement

Ctx	$::=$ <ul style="list-style-type: none"> <li>  <math>x \in \mathcal{N}(\Gamma)</math></li> <li>  <math>x \notin \mathcal{N}(\Gamma)</math></li> <li>  <math>\text{type\_affect} \in \Gamma</math></li> <li>  <math>\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset</math></li> <li>  <math>p_1 = p_2 \implies \Gamma_1 = \Gamma_2</math></li> <li>  <math>p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)</math></li> <li>  <math>\text{fresh } x</math></li> </ul>	$\Gamma_1$ and $\Gamma_2$ are disjoint typing contexts with n
Eq	$::=$ <ul style="list-style-type: none"> <li>  <math>A_1 = A_2</math></li> <li>  <math>A_1 \neq A_2</math></li> <li>  <math>t = u</math></li> <li>  <math>\Gamma = D</math></li> </ul>	
Ty	$::=$ <ul style="list-style-type: none"> <li>  <math>R \stackrel{\text{fix}}{=} \text{rec\_type\_def}</math></li> <li>  <math>\mathcal{U} ; \Gamma \vdash t : A</math></li> </ul>	
Sem	$::=$ <ul style="list-style-type: none"> <li>  <math>t \longrightarrow e \mid t'</math></li> <li>  <math>t \rightsquigarrow t'</math></li> </ul>	
judgement	$::=$ <ul style="list-style-type: none"> <li>  Ctx</li> <li>  Eq</li> <li>  Ty</li> <li>  Sem</li> </ul>	
user_syntax	$::=$ <ul style="list-style-type: none"> <li>  metavariable</li> <li>  term</li> <li>  value</li> <li>  <i>effect</i></li> <li>  data_value</li> <li>  <i>multiplicity</i></li> <li>  sub</li> <li>  subs</li> <li>  var_sub</li> <li>  var_subs</li> <li>  type</li> <li>  data_type</li> <li>  nodest_data_type</li> <li>  type_with_var</li> <li>  data_type_with_var</li> <li>  nodest_data_type_with_var</li> <li>  rec_type_bound</li> <li>  rec_type_def</li> <li>  sign</li> <li>  type_affect</li> </ul>	

	type_affects
	typing_context
	<i>evaluation_context</i>
	terminals

$$\boxed{x \in \mathcal{N}(\Gamma)}$$

$$\boxed{x \notin \mathcal{N}(\Gamma)}$$

$$\boxed{\text{type\_affect} \in \Gamma}$$

$$\boxed{\mathcal{N}(\Gamma_1) \cap \mathcal{N}(\Gamma_2) = \emptyset} \quad \Gamma_1 \text{ and } \Gamma_2 \text{ are disjoint typing contexts with no clashing variable names or labels}$$

$$\boxed{p_1 = p_2 \implies \Gamma_1 = \Gamma_2}$$

$$\boxed{p_1 = p_2 \implies (\Gamma_1 = \Gamma_2 \wedge \Gamma_3 = \Gamma_4)}$$

$$\boxed{\text{fresh } x}$$

$$\boxed{A_1 = A_2}$$

$$\boxed{A_1 \neq A_2}$$

$$\boxed{t = u}$$

$$\boxed{\Gamma = D}$$

$$\boxed{R \stackrel{\text{fix}}{=} \text{rec\_type\_def}}$$

$$\boxed{\mathcal{U}; \Gamma \vdash t : A}$$

$$\frac{}{\mathcal{U}; \emptyset \vdash \bullet : \perp} \text{TyTERM\_NOEFF}$$

$$\frac{}{\mathcal{U}; \emptyset \vdash \star : \mathbf{1}} \text{TyTERM\_U}$$

$$\frac{\emptyset; \Gamma^- \sqcup \{x : A_1\} \vdash t : A_2}{\mathcal{U}; \Gamma^- \vdash \lambda x : A_1. t : A_1 \multimap A_2} \text{TyTERM\_FN}$$

$$\frac{\emptyset; \emptyset \vdash d : \mathbf{N}}{\mathcal{U}; \emptyset \vdash \text{Ur } d : \mathcal{N}} \text{TyTERM\_E}$$

$$\frac{\emptyset; \Gamma^- \vdash d : D}{\mathcal{U}; \Gamma^- \vdash \text{C } d : \mathcal{D}} \text{TyTERM\_C}$$

$$\frac{\emptyset; \Gamma \vdash d : D_1}{\mathcal{U}; \Gamma \vdash \text{Inl } d : D_1 \oplus D_2} \text{TyTERM\_INL}$$

$$\frac{\emptyset; \Gamma \vdash d : D_2}{\mathcal{U}; \Gamma \vdash \text{Inr } d : D_1 \oplus D_2} \text{TyTERM\_INR}$$

$$\frac{R \stackrel{\text{fix}}{=} \mu X. D \quad \emptyset; \Gamma \vdash d : \underline{D}[X := R]}{\mathcal{U}; \Gamma \vdash \text{OR } d : R} \text{TyTERM\_R}$$

$$\frac{\emptyset; \Gamma_1 \vdash d_1 : D_1 \quad \emptyset; \Gamma_2 \vdash d_2 : D_2}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \langle d_1, d_2 \rangle : D_1 \otimes D_2} \text{TyTERM\_P}$$

$$\frac{\mathcal{U}; \Gamma_1 \vdash t_1 : A_1 \quad \mathcal{U}; \Gamma_2 \vdash t_2 : A_2}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \langle t_1 \odot t_2 \rangle : A_1 \wp A_2} \text{TyTERM\_M}$$

$$\frac{}{\mathcal{U}; \{-x : {}^p D\} \vdash [x] : [D]^p} \text{TyTERM\_D}$$

$$\begin{array}{c}
\frac{}{\overline{\mathcal{U}; \{x : A\} \vdash x : A}} \text{TYTERM\_VAR} \\
\\
\frac{}{\overline{\mathcal{U} \sqcup \{x : A\}; \emptyset \vdash x : A}} \text{TYTERM\_VAR'} \\
\\
\frac{\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(A_1 \multimap A_2) \quad \mathcal{U}; \Gamma_2 \vdash u : A_1}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash tu : A_2} \text{TYTERM\_APP} \\
\\
\frac{\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}1 \quad \mathcal{U}; \Gamma_2 \vdash u : A}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\star \mapsto u\} : A} \text{TYTERM\_PATU} \\
\\
\frac{\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(\mathcal{N}) \quad \mathcal{U} \sqcup \{x : \mathcal{d}\mathcal{N}\}; \Gamma_2 \vdash u : A}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Ur } x \mapsto u\} : A} \text{TYTERM\_PATE} \\
\\
\frac{\begin{array}{l} x_1 \notin \mathcal{N}(\Gamma_1) \\ x_2 \notin \mathcal{N}(\Gamma_2) \\ \mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(D_1 \oplus D_2) \\ \mathcal{U}; \Gamma_2 \sqcup \{x_1 : \mathcal{d}D_1\} \vdash u_1 : A \\ \mathcal{U}; \Gamma_2 \sqcup \{x_2 : \mathcal{d}D_2\} \vdash u_2 : A \end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{Inl } x_1 \mapsto u_1, \text{Inr } x_2 \mapsto u_2\} : A} \text{TYTERM\_PATs} \\
\\
\frac{\begin{array}{l} R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\ x \notin \mathcal{N}(\Gamma_1) \\ \mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}R \\ \mathcal{U}; \Gamma_2 \sqcup \{x : \mathcal{d}\underline{D}[X := R]\} \vdash u : A \end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\text{OR } x \mapsto u\} : A} \text{TYTERM\_PATR} \\
\\
\frac{\begin{array}{l} x_1 \notin \mathcal{N}(\Gamma_1) \\ x_2 \notin \mathcal{N}(\Gamma_2) \\ \mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(D_1 \otimes D_2) \\ \mathcal{U}; \Gamma_2 \sqcup \{x_1 : \mathcal{d}D_1, x_2 : \mathcal{d}D_2\} \vdash u : A \end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{case } t \text{ of } \{\langle x_1, x_2 \rangle \mapsto u\} : A} \text{TYTERM\_PATP} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}D}{\mathcal{U}; \Gamma \vdash \text{extract } t : D} \text{TYTERM\_EX} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(A_1 \wp A_2)}{\mathcal{U}; \Gamma \vdash \text{flip } t : \mathcal{d}(A_2 \wp A_1)} \text{TYTERM\_FLIPM} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(A_1 \wp \mathcal{d}(A_2 \wp A_3))}{\mathcal{U}; \Gamma \vdash \text{reassoc } t : \mathcal{d}(\mathcal{d}(A_1 \wp A_2) \wp A_3)} \text{TYTERM\_REASSOCM} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}(\perp \wp D)}{\mathcal{U}; \Gamma \vdash \text{redL } t : \mathcal{d}D} \text{TYTERM\_REDLM} \\
\\
\frac{\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}(A_1 \wp A_2) \quad \mathcal{U}; \Gamma_2 \vdash u : \mathcal{d}(A_1 \multimap A_3)}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash \text{mapL } t \text{ with } u : \mathcal{d}(A_3 \wp A_2)} \text{TYTERM\_MAPLM} \\
\\
\frac{}{\overline{\mathcal{U}; \emptyset \vdash \text{alloc } D : \mathcal{d}(\mathcal{d}[D]^I \wp D)}} \text{TYTERM\_ALLOC} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[1]^p}{\mathcal{U}; \Gamma \vdash t \triangleleft^p \star : \perp} \text{TYTERM\_FILLU}
\end{array}$$



$$\begin{array}{c}
\frac{\begin{array}{l}
x \notin \mathcal{N}(\Gamma_1) \\
\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}[A_1 \multimap A_2]^p \\
\mathcal{U}; \Gamma_2 \sqcup \{x : A_1\} \vdash u : A_2 \\
p = \omega \implies \Gamma_2 = \emptyset
\end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash t \overset{p}{\triangleleft} \lambda x : A_1. u : \perp} \text{TYTERM\_FILLFN} \\
\\
\frac{\begin{array}{l}
\mathcal{U}; \Gamma_1 \vdash t : \mathcal{d}[D]^p \\
\mathcal{U}; \Gamma_2 \vdash u : D \\
p = \omega \implies \Gamma_2 = \emptyset
\end{array}}{\mathcal{U}; \Gamma_1 \sqcup \Gamma_2 \vdash t \overset{p}{\triangleleft} u : \perp} \text{TYTERM\_FILLL} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[\omega N]^p}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \text{Ur} : \mathcal{d}[N]^\omega} \text{TYTERM\_FILLE} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \oplus D_2]^p}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \text{Inl} : \mathcal{d}[D_1]^p} \text{TYTERM\_FILLINL} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \oplus D_2]^p}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \text{Inr} : \mathcal{d}[D_2]^p} \text{TYTERM\_FILLINR} \\
\\
\frac{\begin{array}{l}
R \stackrel{\text{fix}}{=} \mu X. \underline{D} \\
\mathcal{U}; \Gamma \vdash t : \mathcal{d}[R]^p
\end{array}}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \text{@R} : \mathcal{d}[\underline{D}[X := R]]^p} \text{TYTERM\_FILLR} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[D_1 \otimes D_2]^p}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \langle, \rangle : \mathcal{d}(\mathcal{d}[D_1]^p \times \mathcal{d}[D_2]^p)} \text{TYTERM\_FILLP} \\
\\
\frac{\mathcal{U}; \Gamma \vdash t : \mathcal{d}[\mathcal{d}D_1 \times \mathcal{d}D_2]^p}{\mathcal{U}; \Gamma \vdash t \overset{p}{\triangleleft} \langle \odot \rangle : \mathcal{d}(\mathcal{d}[D_1]^p \times \mathcal{d}[D_2]^p)} \text{TYTERM\_FILLM}
\end{array}$$

$$\boxed{t \longrightarrow e \mid t'}$$

$$\begin{array}{c}
\frac{}{\overline{(\mathcal{C}(\lambda x : A. t)) \text{ d} \longrightarrow \varepsilon \mid t[x := \text{d}]}} \text{RLOCAL\_APP} \\
\\
\frac{}{\overline{\text{case } \star \text{ of } \{ \star \mapsto t \} \longrightarrow \varepsilon \mid t}} \text{RLOCAL\_PATU} \\
\\
\frac{}{\overline{\text{case } \mathcal{C}(\text{Ur } d) \text{ of } \{ \text{Ur } x \mapsto t \} \longrightarrow \varepsilon \mid t[x := \mathcal{C} d]}} \text{RLOCAL\_PATE} \\
\\
\frac{}{\overline{\text{case } \mathcal{C}(\text{Inl } d) \text{ of } \{ \text{Inl } x_1 \mapsto t_1, \text{Inr } x_2 \mapsto t_2 \} \longrightarrow \varepsilon \mid t_1[x_1 := \mathcal{C} d]}} \text{RLOCAL\_PATINL} \\
\\
\frac{}{\overline{\text{case } \mathcal{C}(\text{Inr } d) \text{ of } \{ \text{Inl } x_1 \mapsto t_1, \text{Inr } x_2 \mapsto t_2 \} \longrightarrow \varepsilon \mid t_2[x_2 := \mathcal{C} d]}} \text{RLOCAL\_PATINR} \\
\\
\frac{}{\overline{\text{case } \mathcal{C}(\text{@R } d) \text{ of } \{ \text{@R } x \mapsto t \} \longrightarrow \varepsilon \mid t[x := \mathcal{C} d]}} \text{RLOCAL\_PATR} \\
\\
\frac{}{\overline{\text{case } \mathcal{C} \langle d_1, d_2 \rangle \text{ of } \{ \langle x_1, x_2 \rangle \mapsto t \} \longrightarrow \varepsilon \mid t[x_1 := \mathcal{C} d_1, x_2 := \mathcal{C} d_2]}} \text{RLOCAL\_PATP} \\
\\
\frac{}{\overline{\text{extract } (\mathcal{C} d) \longrightarrow \varepsilon \mid d}} \text{RLOCAL\_EX} \\
\\
\frac{}{\overline{\text{flip } (\mathcal{C} \langle v_1 \odot v_2 \rangle) \longrightarrow \varepsilon \mid \mathcal{C} \langle v_2 \odot v_1 \rangle}} \text{RLOCAL\_FLIPM} \\
\\
\frac{}{\overline{\text{reassoc } (\mathcal{C} \langle v_1 \odot \mathcal{C} \langle v_2 \odot v_3 \rangle \rangle) \longrightarrow \varepsilon \mid \mathcal{C} \langle \mathcal{C} \langle v_1 \odot v_2 \rangle \odot v_3 \rangle}} \text{RLOCAL\_REASSOCM}
\end{array}$$

$$\begin{array}{c}
\frac{}{\text{redL } (C \langle \bullet \odot d \rangle) \longrightarrow \varepsilon \mid C d} \quad \text{RLOCAL\_REDLM} \\
\\
\frac{}{\text{mapL } (C \langle v_1 \odot v_2 \rangle) \text{ with } (C (\lambda x : A. t)) \longrightarrow \varepsilon \mid C \langle t[x := v_1] \odot v_2 \rangle} \quad \text{RLOCAL\_MAPLM} \\
\\
\frac{\text{fresh } x}{\text{alloc } D \longrightarrow \varepsilon \mid C \langle C[x] \odot x \rangle} \quad \text{RLOCAL\_ALLOC} \\
\\
\frac{}{(C[x]) \stackrel{p}{\triangleleft} \star \longrightarrow x := \star \mid \bullet} \quad \text{RLOCAL\_FILLU} \\
\\
\frac{}{(C[x]) \stackrel{p}{\triangleleft} \lambda y : A. t \longrightarrow x := \lambda y : A. t \mid \bullet} \quad \text{RLOCAL\_FILLFN} \\
\\
\frac{}{(C[x]) \stackrel{p}{\triangleleft} v \longrightarrow x := v \mid \bullet} \quad \text{RLOCAL\_FILLL} \\
\\
\frac{\text{fresh } x'}{(C[x]) \stackrel{p}{\triangleleft} \text{Ur} \longrightarrow x := \text{Ur } x' \mid C[x']} \quad \text{RLOCAL\_FILLE} \\
\\
\frac{\text{fresh } x'}{(C[x]) \stackrel{p}{\triangleleft} \text{Inl} \longrightarrow x := \text{Inl } x' \mid C[x']} \quad \text{RLOCAL\_FILLINL} \\
\\
\frac{\text{fresh } x'}{(C[x]) \stackrel{p}{\triangleleft} \text{Inr} \longrightarrow x := \text{Inr } x' \mid C[x']} \quad \text{RLOCAL\_FILLINR} \\
\\
\frac{\text{fresh } x'}{(C[x]) \stackrel{p}{\triangleleft} @R \longrightarrow x := @R x' \mid C[x']} \quad \text{RLOCAL\_FILLR} \\
\\
\frac{\text{fresh } x_1 \quad \text{fresh } x_2}{(C[x]) \stackrel{p}{\triangleleft} \langle, \rangle \longrightarrow x := \langle x_1, x_2 \rangle \mid C \langle C[x_1] \odot C[x_2] \rangle} \quad \text{RLOCAL\_FILLP} \\
\\
\frac{\text{fresh } x_1 \quad \text{fresh } x_2}{(C[x]) \stackrel{p}{\triangleleft} \langle \odot \rangle \longrightarrow x := \langle C x_1 \odot C x_2 \rangle \mid C \langle C[x_1] \odot C[x_2] \rangle} \quad \text{RLOCAL\_FILLM}
\end{array}$$

$$\boxed{t \rightsquigarrow t'}$$

$$\frac{t \longrightarrow e \mid t'}{E[t] \rightsquigarrow (E[t'])[e]} \quad \text{RGLOBAL\_GLOBAL}$$

Definition rules: 57 good 0 bad  
 Definition rule clauses: 118 good 0 bad