

ITC Tutorial 3

Symbol	Probability $P(X_i)$	Code	Code length $L(X_i)$	$P(X_i) \times L(X_i)$
A	0.3	0	1	0.30
B	0.25	10	2	0.50
C	0.2	110	3	0.60
D	0.15	1110	4	0.60
E	0.1	11110	5	0.50

$$\begin{aligned}
 \text{Avg Code length} &= \sum P(X_i) \times L(X_i) \\
 &= 0.3 + 0.5 + 0.6 + 0.5 + 0.6 \\
 &= 2.5 \text{ bits}
 \end{aligned}$$

$$\begin{aligned}
 \text{Entropy} &= \sum P(X_i) \log_2 \left(\frac{1}{P(X_i)} \right) \\
 &= 0.3 \log_2 \frac{1}{0.3} + 0.25 \log_2 \frac{1}{0.25} \\
 &\quad + 0.2 \log_2 \frac{1}{0.2} + 0.15 \log_2 \frac{1}{0.15} \\
 &\quad + 0.1 \log_2 \frac{1}{0.1} \\
 &= 2.228
 \end{aligned}$$

$$\begin{aligned}
 \text{Code Efficiency} &= \frac{\text{Entropy}}{\text{Avg Code length}} \\
 &= \frac{2.228}{2.5} \\
 &= \underline{0.8912}
 \end{aligned}$$

$$\begin{aligned}
 \text{Code efficiency \%} &= 0.8912 \times 100 \\
 &= 89.12 \%
 \end{aligned}$$

$$\begin{aligned}\text{Code redundancy} &= 1 - \text{code efficiency} \\ &= 1 - 0.8912 \\ &= 0.1088\end{aligned}$$

$$\begin{aligned}\text{Code redundancy \%} &= 0.1088 \times 100 \\ &= 10.88\%\end{aligned}$$

Q2) Difference between lossy and lossless compression

Lossy Compression	Lossless Compression
1) Lossy compression is the method which eliminates that data which is not noticeable.	While lossless does not eliminate that data is not noticeable.
2) In this a file does not restore or rebuilt in its original form.	In this file can be restored in its original form.
3) Data quality is compromised.	Does not compromise the data's quality.
4) This has more data holding capacity.	This has less data holding capacity comparatively.

When average code length approaches the entropy of DMS it is said to be entropy encoding. It is a lossless data compressing scheme.

Examples :-

- i) Shannon Fano coding
 - (i) List the symbols in order of decreasing probability.
 - (ii) Split the 2 sets that are as close to equiprobable as possible. Upper set - 0, lower set 1

ii) Continue this process each time positioning them into sets until further splitting is not possible

② Huffman coding

- i) List the source symbols in decreasing probability order
- ii) Combine probabilities of 2 symbols with resultant probabilities. This is a reduction. Continue this step until possibilities are remaining
- iii) Start encoding the last reduction, first digit - 0 last digit - 1
- iv) Now assign 0 & 1 to 2nd digit in the previous step and keep doing. All first column is reached.