DATA MINING

SEQUENCE MINING AND GRAPH MINING

OVERVIEW

- Apriori refresher
- Frequent-sequence mining
- Graphs in data mining
- Introduction to graph mining

APRIORI

APRIORI IN A NUTSHELL

- Problem: we cannot test the frequency of all possible itemsets
- Needed: algorithm that reduces the number of candidates!
- But: we cannot miss interesting candidates in the process
- Apriori property: if an itemset is not frequent, its supersets won't be either

APRIORI IN A NUTSHELL

Generate k-candidates using (only) frequent (k-1)-itemsets

- Join if only last element is different
- **Prune** if any k-subset is not in the list of frequent k-itemsets

Test frequency of candidates

TID	Items
10	A, C, D
20	B, C
30	A, B, C, E
40	B, E

Let us apply Apriori to this data set, containing a list of items for each transaction ID.

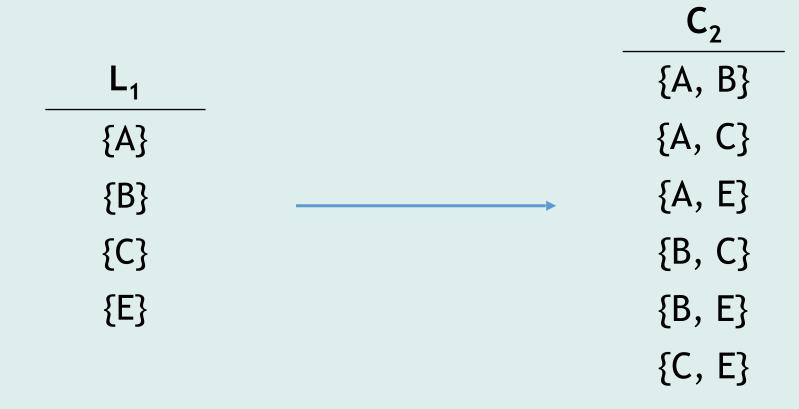
1. Scan for frequent 1-itemsets

TID	Items	C ₁	Support
10	A, C, D	{A}	2
20	B, C	 {B}	3
30	A, B, C,	{C}	3
40	E B, E	{D}	1
70	, L	{E}	2

2. Discard non frequent candidates: we have L₁

C ₁	Support	L ₁
{A}	2	{A}
{B}	3	 {B}
{C}	3	{C}
{D}	1	{E}
{E}	2	

3. Join step to produce k = 2 candidates



4. Pruning (trivial for k = 2)

C_2	_
{A, B}	
{A, C}	
{A, E}	
{B, C}	
{B, E}	
{C, E}	

C ₂	Subset	In L₁?
{A, B}	{A}	yes
	{B}	yes
{A, C}	{A}	yes
	{C}	yes
{A, E}	{A}	yes
	{E}	yes
{B, C}	{B}	yes
	{C}	yes
{B, E}	{B}	yes
	{E}	yes
{C, E}	{C}	yes
	{E}	yes

5. Scan and discard non-frequent candidates: we have L₂

C ₂	Support	L_2	Support
{A, B}	1	 A, C}	2
{A, C}	2	B, C}	2
{A, E}	1	B, E}	2
{B, C}	2	,	_
{B, E}	2		
{C, E}	1		

6. Join step to produce k = 3 candidates

L ₂	Support	C ₃
{A, C}	2	{B, C, E}
{B, C}	2	
{B, E}	2	

7. Prune

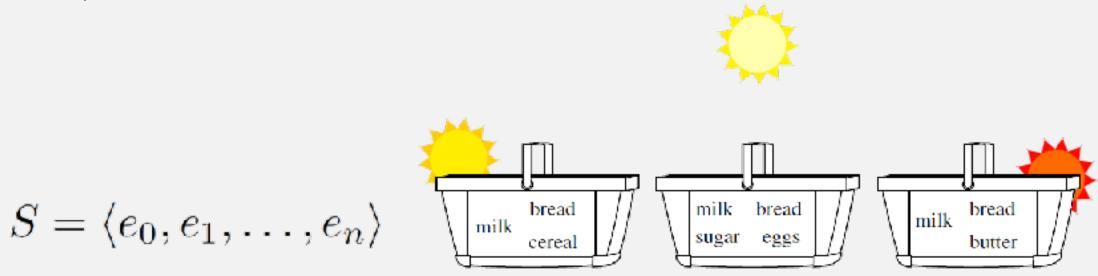
C_3	Subset	In L ₂ ?
{B, C, E}	{B, C}	yes
	{B, E}	yes
	{C, E}	no

8. All C₃ candidates have been rejected. End and return L₁, L₂

TID	Items	L ₂	L ₁
	A, C, D	{A, C}	{A}
20	В, С	{B, C}	{B}
30	A, B, C, E	{B, E}	{C}
40	B, E		{E}

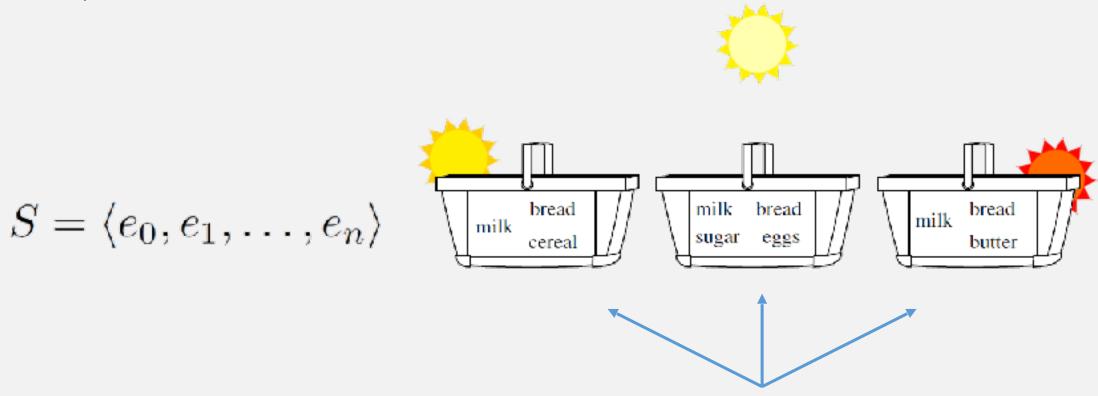
FREQUENT-SEQUENCE MINING

SEQUENCE



- Ordered list of events
- Events are also called elements
- E.g., in customer purchase each event is an itemset

SEQUENCE



The items in an itemset are considered **simultaneous**

SEQUENCE

$$S = \langle e_0, e_1, \dots, e_n \rangle$$

- Itemsets in a sequence are written with parentheses: $(x_1, x_2, ..., x_a)$
- An item can only happen once in an itemset, but may appear in different events!
- Sequence length: number of items (I-sequence)

SEQUENCE—EXAMPLE

Sequence

```
(bread, milk, cereal) (milk, bread, sugar, eggs) (bread, milk, butter)
```

Elements

```
(bread, milk, cereal)(milk, bread, sugar, eggs)(bread, milk, butter)
```

Items

bread, milk, cereal, sugar, eggs, butter

Length: 10

SUB/SUPERSEQUENCE

 S_1 is a subsequence of S_2 if:

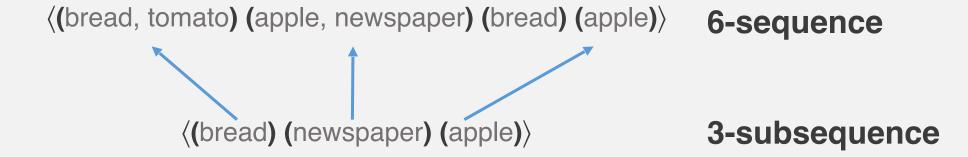
- Every itemset in S₁ has a superset in S₂
- The order of the itemsets in S₁ is the same
 as the order of their supersets in S₂

((bread, tomato) (apple, newspaper) (bread) (apple) (bread) (hewspaper) (apple)

SUB/SUPERSEQUENCE

 S_1 is a subsequence of S_2 if:

- Every itemset in S₁ has a superset in S₂
- The order of the itemsets in S₁ is the same
 as the order of their supersets in S₂



- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- **3.** ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

- 1. ((tomato) (newspaper)) Yes
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

- 1. ((tomato) (newspaper)) Yes
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- **3.** ((newspaper, tomato))
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

No (missing super-itemset)

- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

No (missing super-itemset)

- 1. \(\text{(tomato) (newspaper)}\) 2. ((newspaper) (tomato))
- **3.** ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

No (wrong order)

No (missing super-itemset)

Yes

- (tomato) (newspaper)
 (newspaper) (tomato)
 (newspaper, tomato)

 No (wrong order)
 No (missing super-itemset)
- 4. ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

((tomato) (newspaper))
 ((newspaper) (tomato))
 ((newspaper, tomato))
 ((tomato) (apple))
 ((apple) (apple))
 ((apple, apple))

7. ((newspaper) (apple))

8. ((bread) (tomato))

- 1. ((tomato) (newspaper))
- 2. ((newspaper) (tomato))
- **3.** ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- **6. ((apple, apple))**
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

No (missing super-itemset)

Yes

Yes

- 1. ((tomato) (newspaper))
- 2. ((newspaper) (tomato))
- **3.** ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- **6. ((apple, apple))**
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

No (missing super-itemset)

Yes

/es

No (repetition within itemset)

No (repetition within itemset)

(tomato) (newspaper)
 (newspaper) (tomato)
 (newspaper, tomato)
 (tomato) (apple)

Yes

Yes

8. ((bread) (tomato))

7. ((newspaper) (apple))

5. ((apple) (apple))

6. ((apple, apple))

1. \(\text{(tomato) (newspaper)}\) No (wrong order) 2. ((newspaper) (tomato)) No (missing super-itemset) **3.** ((newspaper, tomato)) 4. ((tomato) (apple)) **5.** ((apple) (apple)) **No** (repetition within itemset) 6. ((apple, apple)) Yes 7. ((newspaper) (apple))

8. ((bread) (tomato))

1. \(\text{(tomato) (newspaper)}\) No (wrong order) 2. ((newspaper) (tomato)) No (missing super-itemset) **3.** ((newspaper, tomato)) 4. ((tomato) (apple)) **5.** ((apple) (apple)) **No** (repetition within itemset) 6. ((apple, apple)) 7. ((newspaper) (apple))

8. ((bread) (tomato))

Exercise: Which is a valid subsequence of ...?

(bread, tomato) (apple, newspaper) (bread) (apple)

- 1. \(\text{(tomato) (newspaper)}\)
- 2. ((newspaper) (tomato))
- 3. ((newspaper, tomato))
- **4.** ((tomato) (apple))
- **5.** ((apple) (apple))
- 6. ((apple, apple))
- 7. ((newspaper) (apple))
- 8. ((bread) (tomato))

Yes

No (wrong order)

No (missing super-itemset)

Yes

Yes

No (repetition within itemset)

Yes

No (wrong order/missing super-itemset)

DATABASE

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

- Each element: data sequence
- Data sequence: ordered list of itemsets
 May or may not have timestamps

DATABASE

Two examples

ds 1	1 2 3 4	(Basement) (Atrium) (Elevator) (Class)	ds 1	02/02/2015 02/02/2015 07/02/2015	15:02 15:17 09:20	(bread, tomato) (apple, newspaper) (apple)
ds 2	1 2 3	(Basement) (Scrollbar) (Basement)	ds 2	03/02/2015 07/02/2015 10/02/2015	16:04 18:11 00:01	(bread) (newspaper) (apple)
	1	(Atrium)	ds 3	01/02/2015	19:45	(apple, bread)
ds 3	3	(Class) (Atrium)	ds 4	01/02/2015	19:50	(apple)

- Is supported by a minimum number of data sequences: minimum support
- If a sequence is a subsequence of a data sequence,

it is **supported** by it. Ignore time, not order

A sequence can only be supported
 only once by each data sequence

ds 1	02/02/2015 02/02/2015 07/02/2015	15:02 15:17 09:20	(bread, tomato) (apple, newspaper) (apple)
ds 2	03/02/2015 07/02/2015 10/02/2015	16:04 18:11 00:01	(bread) (newspaper) (apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	
(apple, bread)	
(bread)	
(apple)	

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	2
(apple, bread)	
(bread)	
(apple)	

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	2
(apple, bread)	1
(bread)	
(apple)	

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	2
(apple, bread)	1
(bread)	3
(apple)	

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	2
(apple, bread)	1
(bread)	3
(apple)	4

Only counts once!

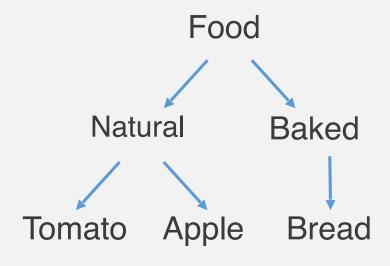
ds 1	02/02/2015 02/02/2015 07/02/2015	15:02 15:17 09:20	(bread, tomato) (apple, newspaper) (apple)
ds 2	03/02/2015 07/02/2015 10/02/2015	16:04 18:11 00:01	(bread) (newspaper) (apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(bread) (apple)	2
(apple, bread)	1
(bread)	3
(apple)	4

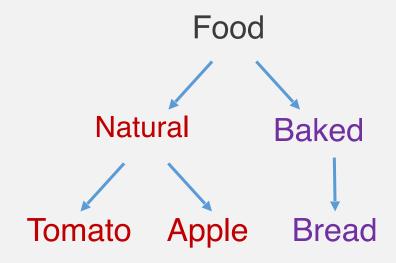
CONSTRAINTS—TAXONOMIES

When counting support an item is equivalent to all its ancestors

ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)



CONSTRAINTS—TAXONOMIES



ds 1	02/02/2015	15:02	(bread, tomato)
	02/02/2015	15:17	(apple, newspaper)
	07/02/2015	09:20	(apple)
ds 2	03/02/2015	16:04	(bread)
	07/02/2015	18:11	(newspaper)
	10/02/2015	00:01	(apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(apple)

Sequence	Support
(food) (newspaper)	2
(baked) (newspaper)	2
(natural, newspaper)	1
(food)	4
(natural)	4

CONSTRAINTS—SLIDING WINDOW

- Close itemsets may be considered simultaneous
- A super-itemset may be distributed across itemsets!

ds 1	02/02/2015 02/02/2015 07/02/2015 07/02/2015	15:02 15:17 09:20 09:35	(bread, tomato) (apple, newspaper) (apple) (tomato)
ds 2	03/02/2015 07/02/2015 10/02/2015	16:04 18:11 00:01	(bread) (newspaper) (apple)
ds 3	01/02/2015	19:45	(apple, bread)
ds 4	01/02/2015	19:50	(bread)

Window size = 30 min

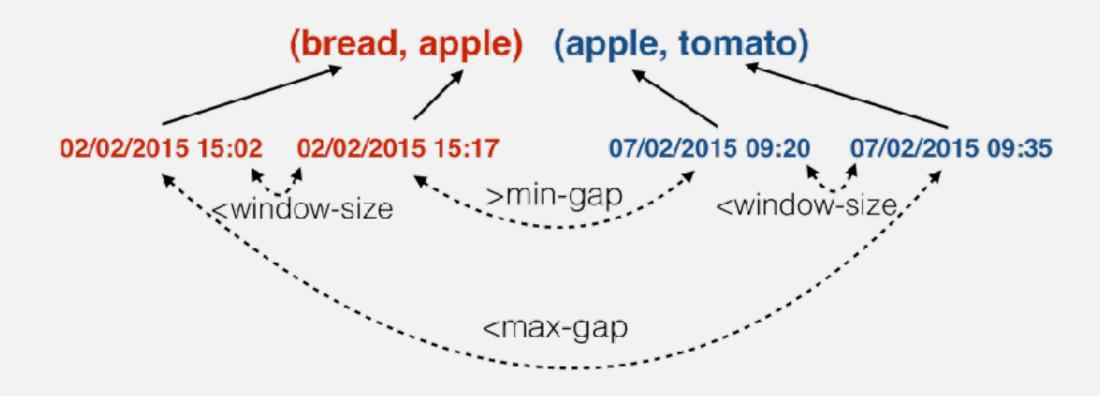
Sequence	Support
(bread) (apple)	2
(apple, bread)	2
(apple, tomato)	1
(bread)	3
(apple)	4

CONSTRAINTS—GAPS

- Window size: itemsets closer than this may be considered simultaneous
- Min-gap: minimum time between last item in one itemset and the first in the next
- Max-gap: maximum time between the first item in one itemset and the last in the next

ds 1 02/02/2015 15:02 (bread, tomato)
02/02/2015 15:17 (apple, newspaper)
07/02/2015 09:20 (apple)
07/02/2015 09:35 (tomato)

```
ds 1 02/02/2015 15:02 (bread, tomato)
02/02/2015 15:17 (apple, newspaper)
07/02/2015 09:20 (apple)
07/02/2015 09:35 (tomato)
```



CONSTRAINTS

 Anti-monotonic: if a sequence does not satisfy it, neither do its supersequences

E.g., Apriori property, "an itemset is only frequent if its subsets are"

 Monotonic: if a sequence satisfies it, so do its supersequences

E.g., duration more than 10 min (supersequences are longer!)

• Succint: we can enumerate all and only those sequences that are guaranteed to satisfy the constraint

E.g., price lower than 5€

GSP ALGORITHM

GSP: GENERALISED SEQUENTIAL PATTERNS

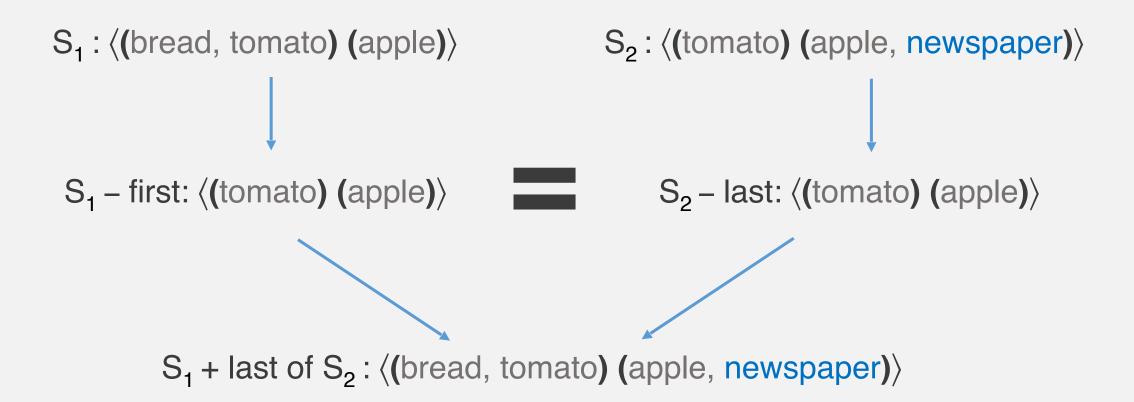
Extension of the Apriori algorithm to sequences!

- 1. Generate (k+1)-candidates using frequent k-sequences
 - Generate (join)
 - Prune
- 2. Count support of candidates
- 3. Drop non-frequent candidates

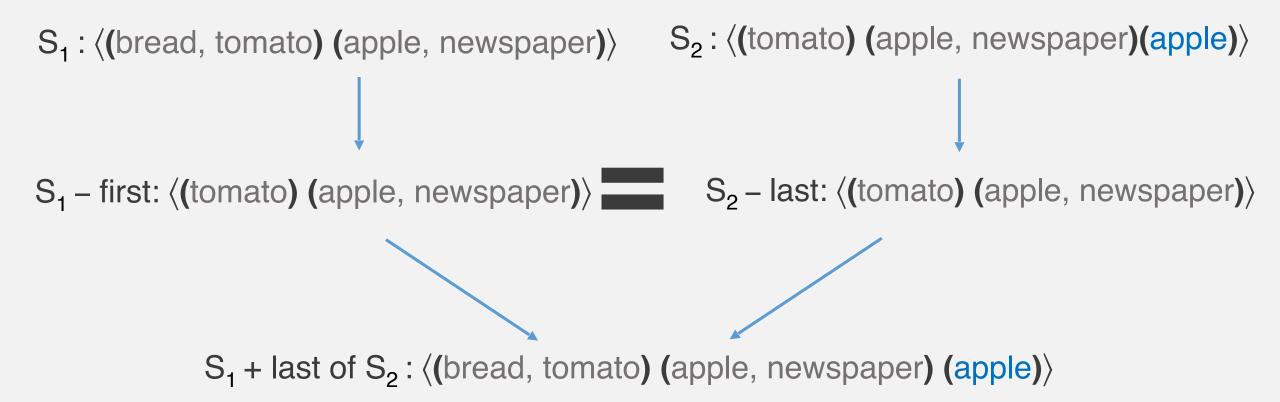
GSP—CANDIDATE GENERATION

- Join sequences S₁ and S₂ if removing the first item of S₁ and the last of S₂ produces the same sequence
 - This is a contiguous subsequence
 - **Items** need to be **sorted**! (E.g., lexicographical.)
- Add last item of S₂ to S₁
 - If last item is a separate itemset, append as new itemset.
 Otherwise add to the last itemset of
 - For k = 2 generate 1 and 2 itemsets (E.g., $\langle (A) (B) \rangle$ and $\langle (AB) \rangle$)

GSP—CANDIDATE GENERATION



GSP—CANDIDATE GENERATION



GSP—CANDIDATE PRUNING

- A candidate may only be frequent if all its contiguous subsequences are
- Generation of contiguous k-subsequences from (k+1)sequence
 - Drop item from any itemset with more than one item
 - Dropping from first and last itemsets was used in join!

GSP—CANDIDATE PRUNING

5-candidate: ((bread, tomato) (apple, newspaper) (apple)

4-contiguous subsequences:

S₁: ((bread, tomato) (apple, newspaper))

 S_2 : $\langle \text{(tomato) (apple, newspaper)(apple)} \rangle$

BEYOND GSP

- GSP requires many very costly scans for support count
- Improvements possible. E.g., hash-tree technique
- Alternative algorithms

SPADE: vertical format

PrefixSpan: Pattern-growth method, avoids candidate generation!

WHY SEQUENTIAL PATTERNS?

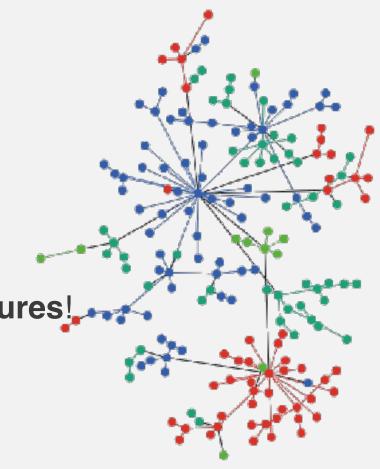
- Useful to understand data
- It is possible to extract association rules as with itemsets!

GRAPH MINING

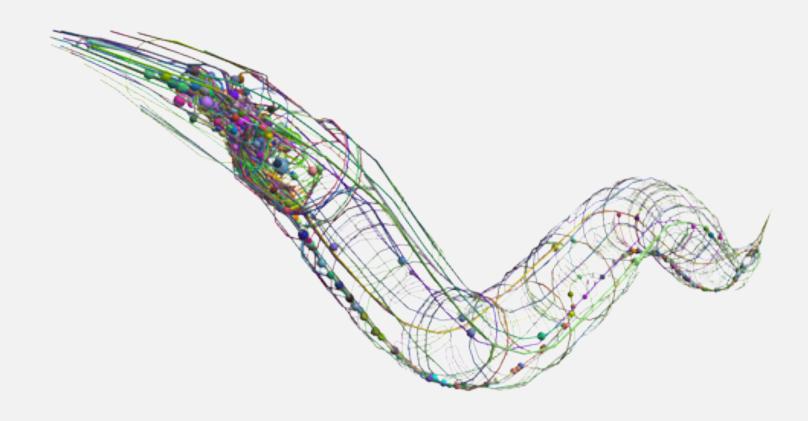
WHY GRAPHS?

Graphs are great for more complex data structures!

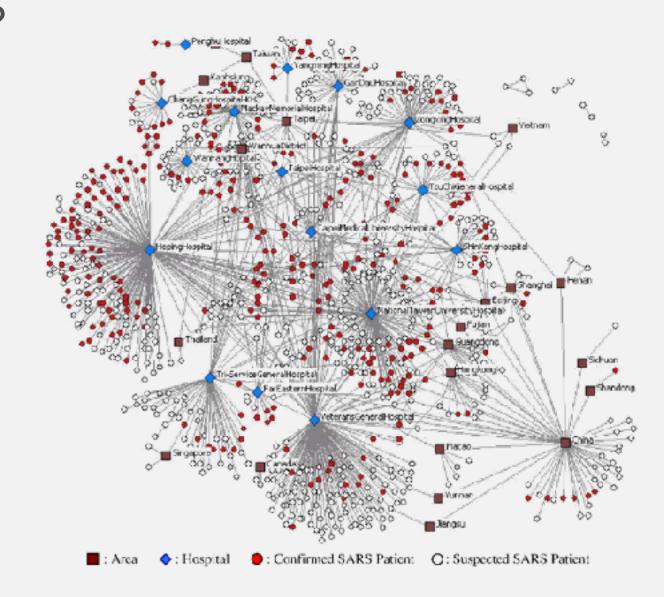
- Social networks
- Computational biology
- Long "etc."



WHY GRAPHS?



WHY GRAPHS?

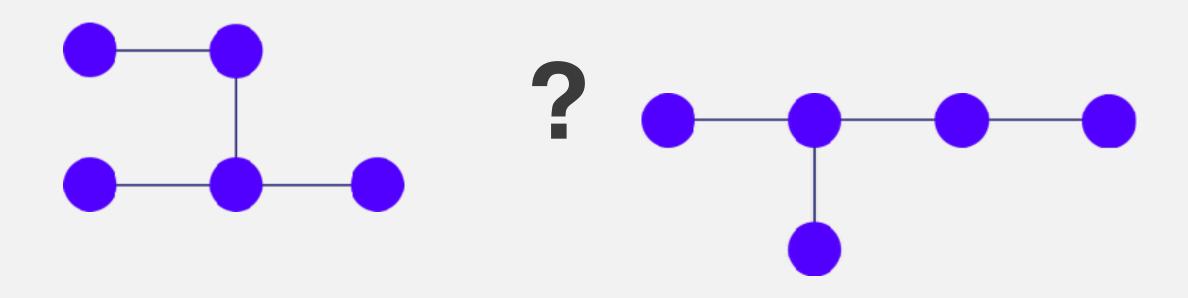


DOI: 10.1007/978-3-540-72608-1

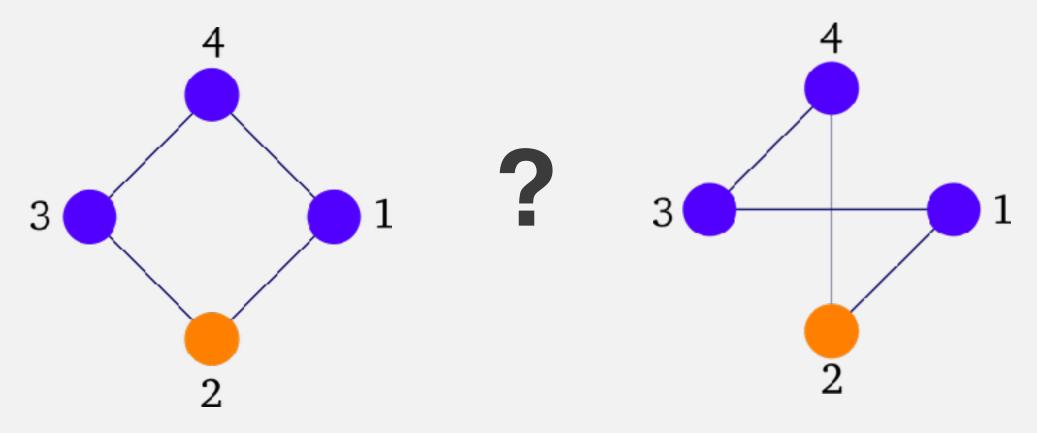
FREQUENT GRAPH MINING

- Adaptation of Apriori ideas
- Key question: how do you compare two graphs?

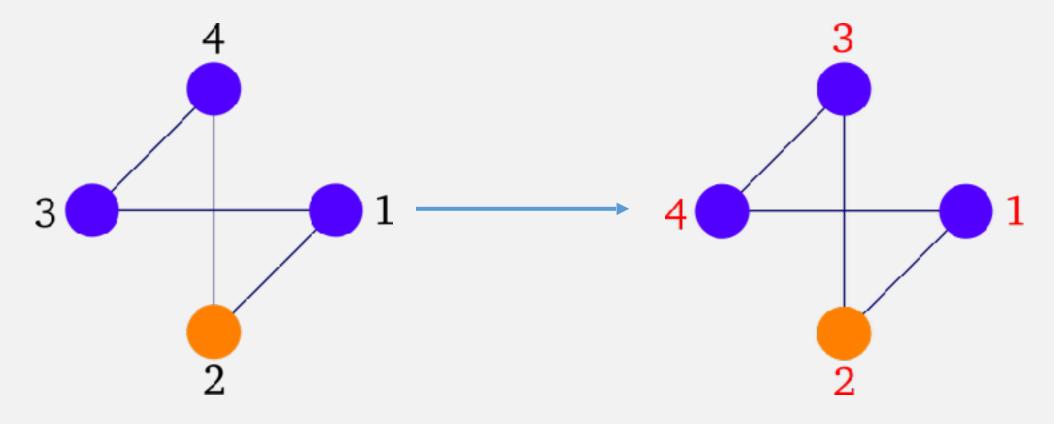
GRAPH COMPARISON



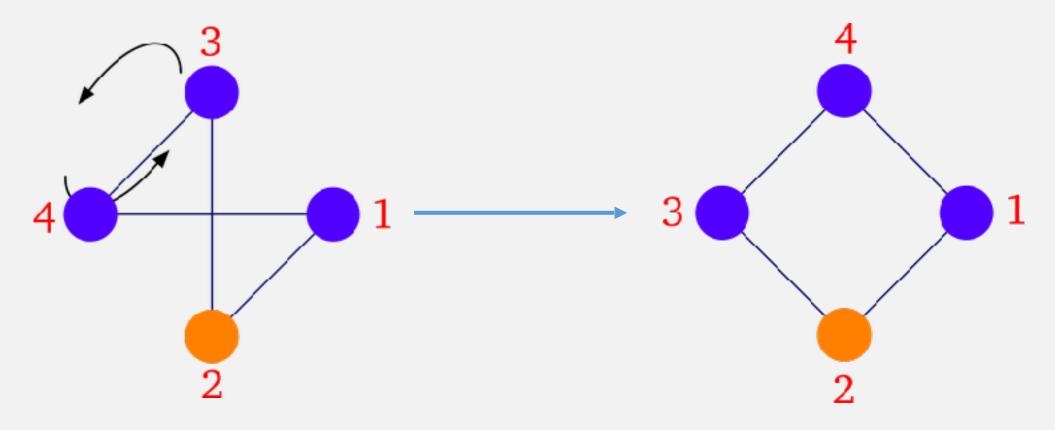
Problem: ISOMORPHISM ("equal shape")



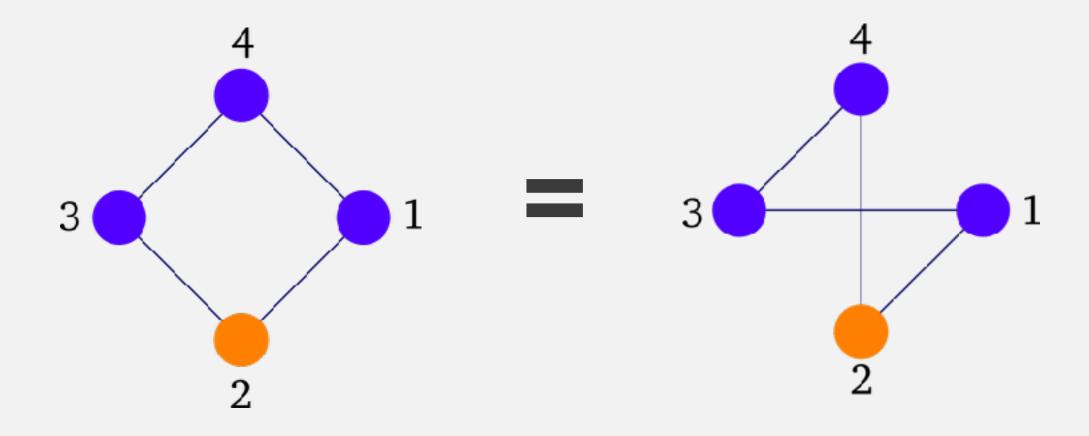
Node identifiers are arbitrary; are these graphs the same?



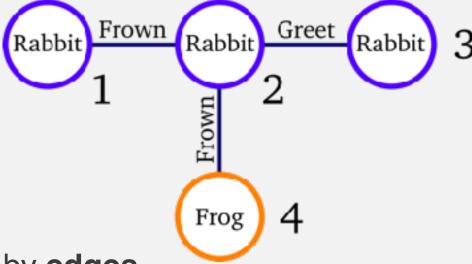
Node identifiers are arbitrary!



Position of nodes is not fixed! (Not in this case at least.)

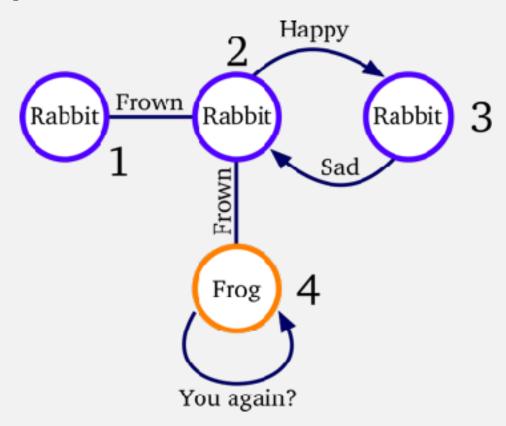


GRAPH REPRESENTATION



- Graph: set of vertices (nodes) connected by edges
- Vertices and edges have non-unique labels
- Vertices have unique but arbitrary identifiers
 Edges are defined by their label and the nodes they connect

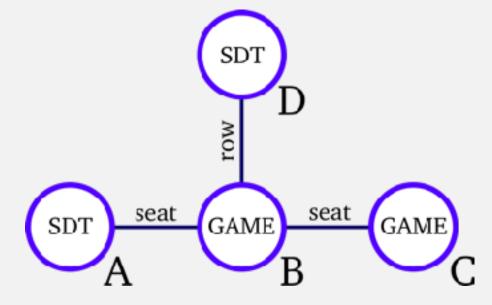
GRAPH REPRESENTATION



Note: We will work talk about simple, undirected graphs

GRAPH REPRESENTATION

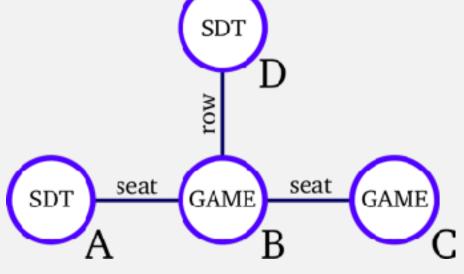
- The first step towards comparison is unique representation of isomorphisms
- Let's begin by representing graphs as a matrix

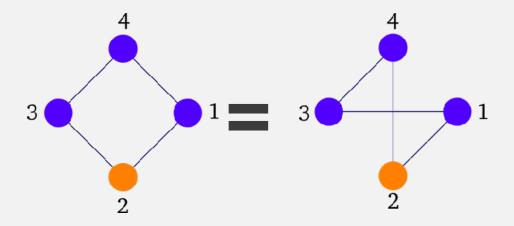


ADJACENCY MATRIX

	A (SDT)	B (GAME)	C (GAME)	D (SDT)
Α	0	seat	0	0
В	seat	0	seat	row
C	0	seat	0	0
D	0	row	0	0

- Columns and rows for vertices
- Cells represent the edge between the row and column nodes
- Zero if no edge

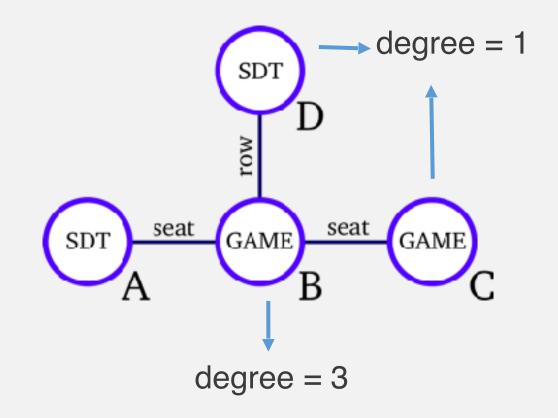




- Unique code that represents a graph and all its isomorphic graphs
- Flattening (processing) adjacency matrix
 - 1. Sort vertices by **degree** (number of edges)
 - 2. Sort vertices (of same degree) by label
 - 3. Compare edge permutations

1. Sort vertices by **degree** (number of edges)

	A	В	С	D
	(SDT)	(GAME)	(GAME)	(SDT)
Α	0	seat	0	0
В	seat	0	seat	row
С	0	seat	0	0
D	0	row	0	0



1. Sort vertices by degree (number of edges)

	Α	B C D		D
	(SDT)	(GAME)	(GAME)	(SDT)
Α	0	seat	0	0
В	seat	0	seat	row
С	0	seat	0	0
D	0	row	0	0

	Α	С	D	В
	(SDT)	(GAME)	(SDT)	(GAME)
Α	0	0	0	seat
C	0	0	0	seat
D	0	0	0	row
В	seat	seat	row	0

2. Sort vertices (of same degree) by label

	Α	C	D	В
	(SDT)	(GAME)	(SDT)	(GAME)
Α	0	0	0	seat
C	0	0	0	seat
D	0	0	0	row
В	seat	seat	row	0

	С	Α	D	В
	(GAME)	(SDT)	(SDT)	(GAME)
С	0	0	0	seat
Α	0	0	0	seat
D	0	0	0	row
В	seat	seat	row	0

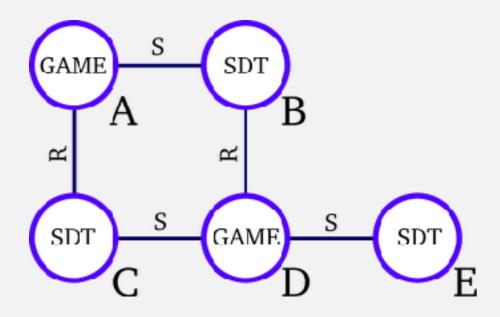
3. Compare edge **permutations**

Permutation 1

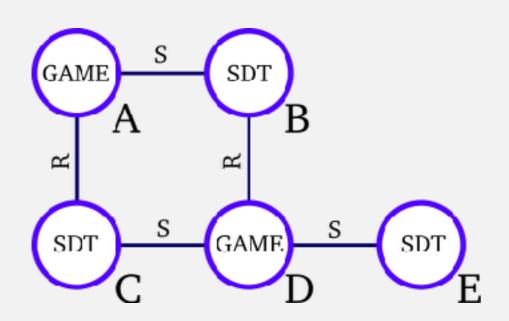
Α D В (GAME) (SDT) (SDT) (GAME) C 0 seat 0 seat D 0 0 row В 0 seat seat row

Permutation 2

	С	D	Α	В
	(GAME)	(SDT)	(SDT)	(GAME)
С	0	0	0	seat
D	0	0	0	row
Α	0	0	0	seat
В	seat	row	seat	0

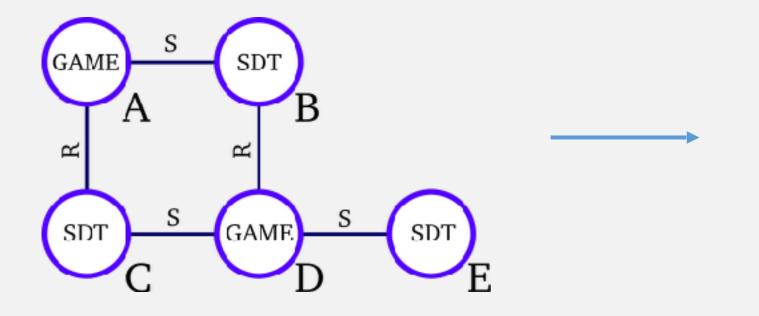


We start with the adjacency matrix



	A	В	C	D	E
A	0	S	R	0	0
В	S	0	0	R	0
С	R	0	0	S	0
D	0	R	S	0	S
E	0	0	0	S	0

1. Sort vertices by **degree** (number of edges)



Node	Degree
Α	2
В	2
С	2
D	3
E	1

STEP BY STEP EXAMPLE

1. Sort vertices by **degree** (number of edges)

Node	Degree
A	2
В	2
C	2
D	3
E	1

	A	В	C	D	E
A	0	S	R	0	0
В	S	0	0	R	0
С	R	0	0	S	0
D	0	R	S	0	S
E	0	0	0	S	0

	Е	Α	В	C	D
Е	0	0	0	0	S
A	0	0	S	R	0
В	0	S	0	0	R
С	0	R	0	0	S
D	S	0	R	S	0

2. Sort vertices (of same degree) by label

$$A = GAME$$
, $B = SDT$, $C = SDT$

 \Rightarrow A before B, C \checkmark



	Е	Α	В	С	D
Е	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
С	0	R	0	0	S
D	S	0	R	S	0

3. Compare permutations

For each permutation we generate a **code** reading **below the diagonal**

	Е	Α	В	С	D
Е	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
С	0	R	0	0	S
D	S	0	R	S	0

3. Compare permutations

For each permutation we generate a **code**

reading below the diagonal

	Ш	Α	В	C	D
Ш	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
C	0	R	0	0	S
D	S	0	R	S	0

Label: 0

3. Compare permutations

For each permutation we generate a **code**

reading below the diagonal

	Ш	Α	В	С	D
Ш	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
C	0	R	0	0	S
D	S	0	R	S	0

Label: 0; 0S

3. Compare permutations

For each permutation we generate a **code**

reading below the diagonal

	Е	Α	В	С	D
Е	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
C	0	R	0	0	S
D	S	0	R	S	0

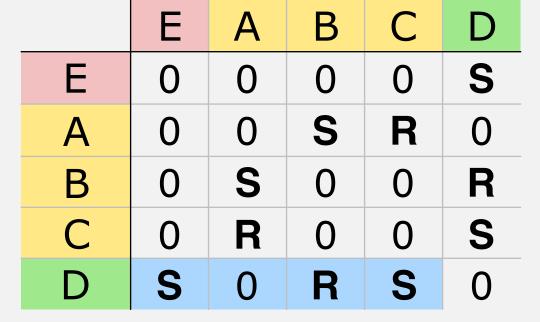
Label: 0; 0S; 0R0

3. Compare permutations

For each permutation we generate a **code**

reading below the diagonal

Label:	0;	0S ;	0R0;	SORS





- 3. Compare permutations
- Rules 1 and 2 have fixed the position of E, A and

$$A = GAME > B = C = SDT$$

• Compare (E A B C D) and (E A C B D)

	Ш	Α	В	С	D
Е	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
C	0	R	0	0	S
D	S	0	R	S	0

Permutation (E A B C D)

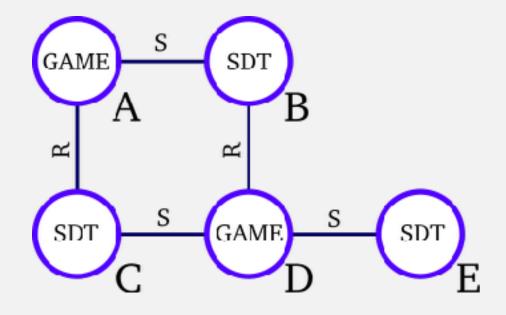
	Е	Α	В	С	D
Е	0	0	0	0	S
Α	0	0	S	R	0
В	0	S	0	0	R
С	0	R	0	0	S
D	S	0	R	S	0

Permutation (E A C B D)

	E	Α	С	В	D
Е	0	0	0	0	S
Α	0	0	R	S	0
С	0	R	0	0	R
В	0	S	0	0	S
D	S	0	S	R	0

Label: 0; 0**S**; 0R0; S0RS

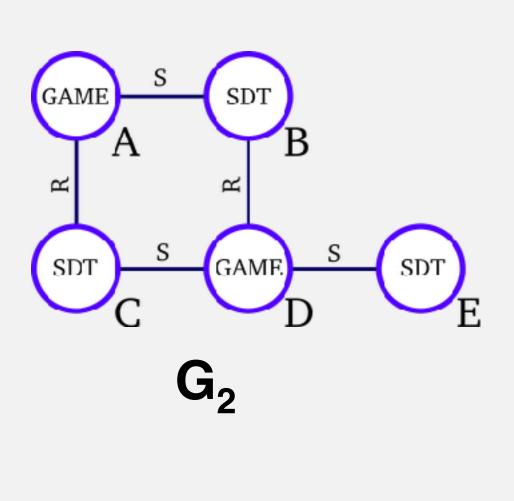
Label: 0; 0R; 0S0; S0SR



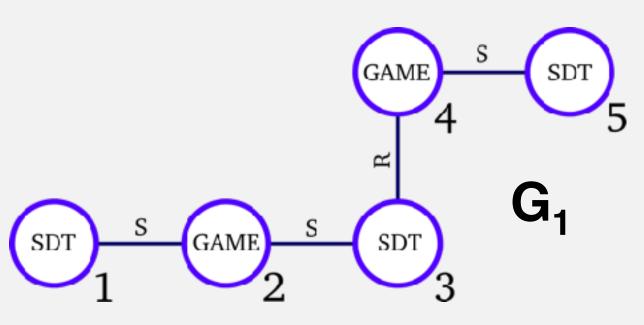
Canonical label: 0; 0R; 0S0; S0SR

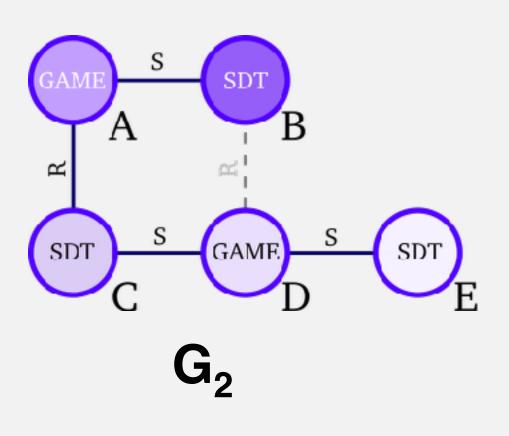
- G₁ is a subgraph of G₂ if we can find an isomorphic graph of G₁ within
 G₂
- If G₁ and G₂ are the same the operation is an **automorphism** (and returns all the isomorphisms of the graph!)
- A subgraph with k edges is called a k-subgraph (but some methods count nodes!)

Testing if a graph is a subgraph of another is a **costly operation**!



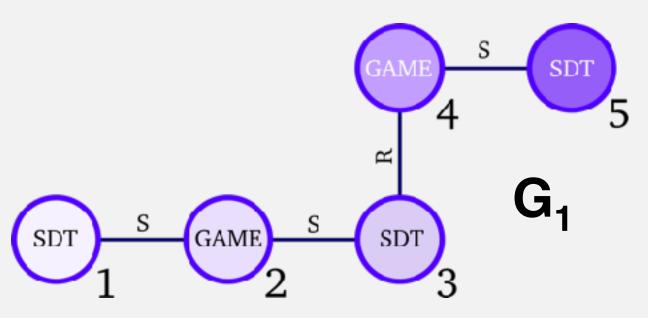
Is G₁ a subgraph of G₂?





Is G₁ a subgraph of G₂?

yes!



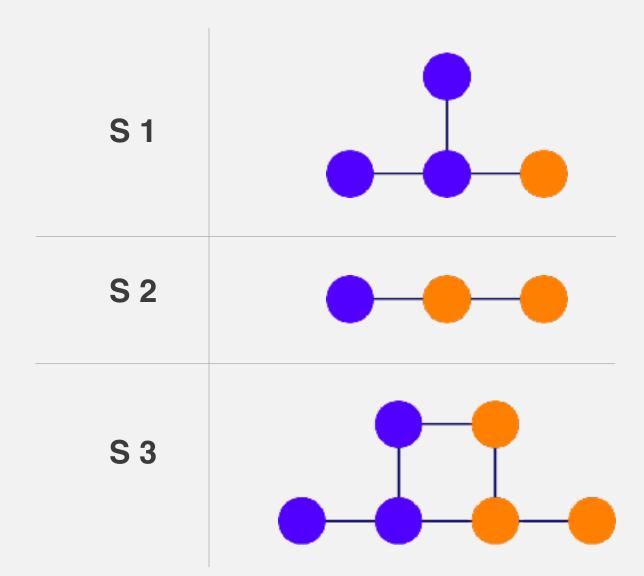
APRIORI FOR GRAPHS

APRIORI FOR GRAPHS

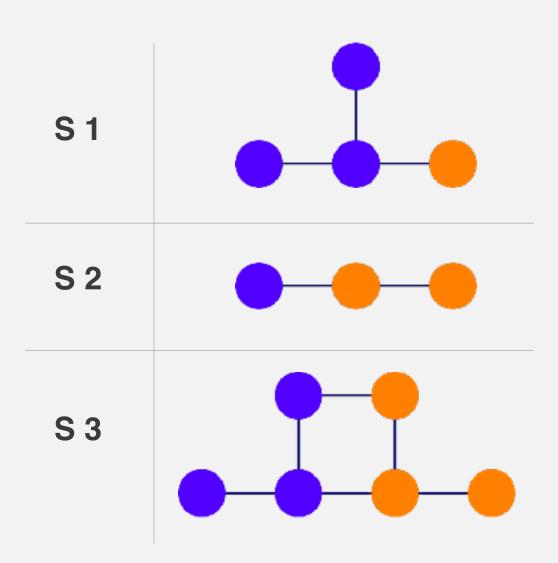
- Adapting Apriori for graphs is not trivial
- (k+1)-candidates: increase one edge or one vertex?
- Adding two graphs does not create unique results

DATASET

- Each sample is a graph
- Independent samples (social network of different users) or from a larger graph (distance-2 social networks starting at different nodes)

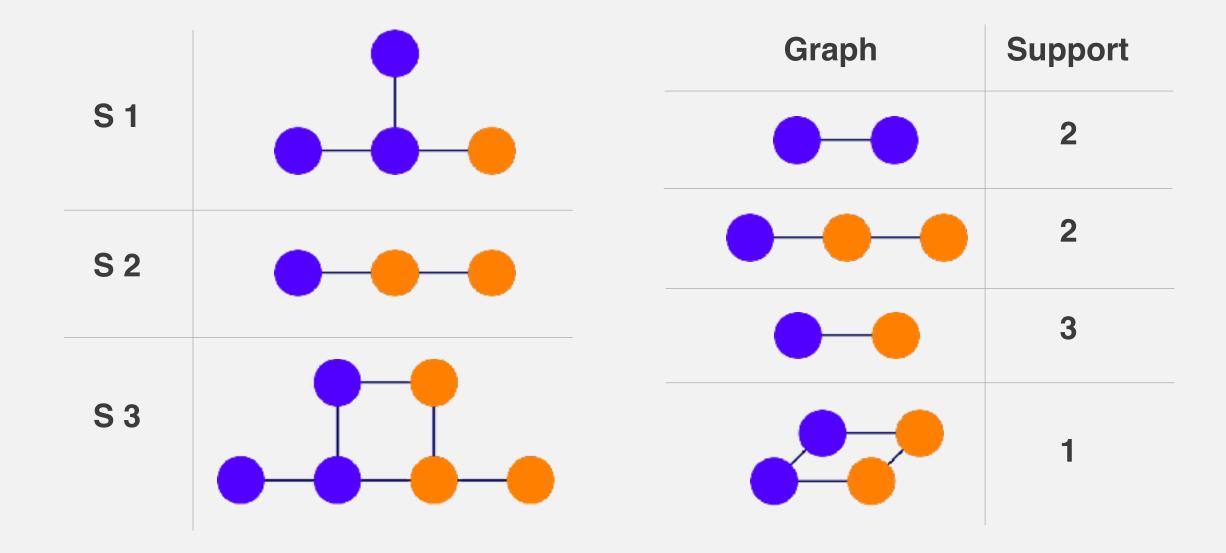


SUPPORT COUNT



- A graph is supported by a sample if it contains one of its subgraphs
- A sample supports at most once each graph

SUPPORT COUNT



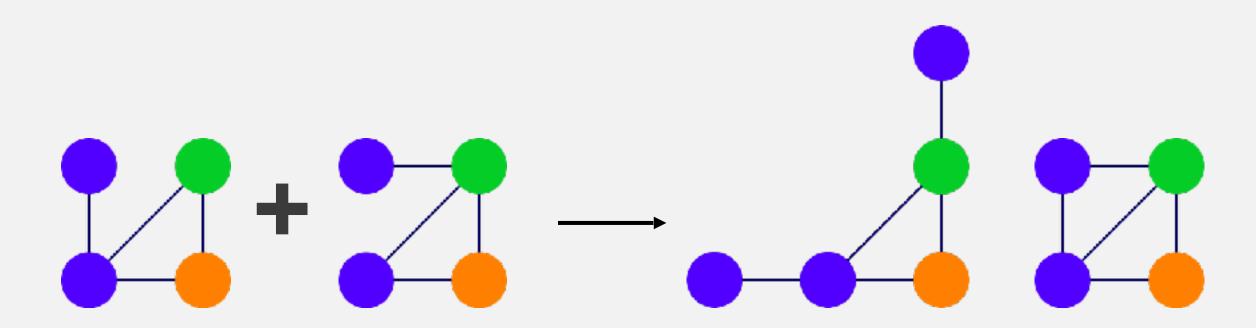
FSG ALGORITHM

- Generate (k+1)-candidates adding a node
 AGM adds one vertex instead!
 - Join
 - Prune
- 2. Count support of candidates
- 3. Drop non-frequent candidate graphs

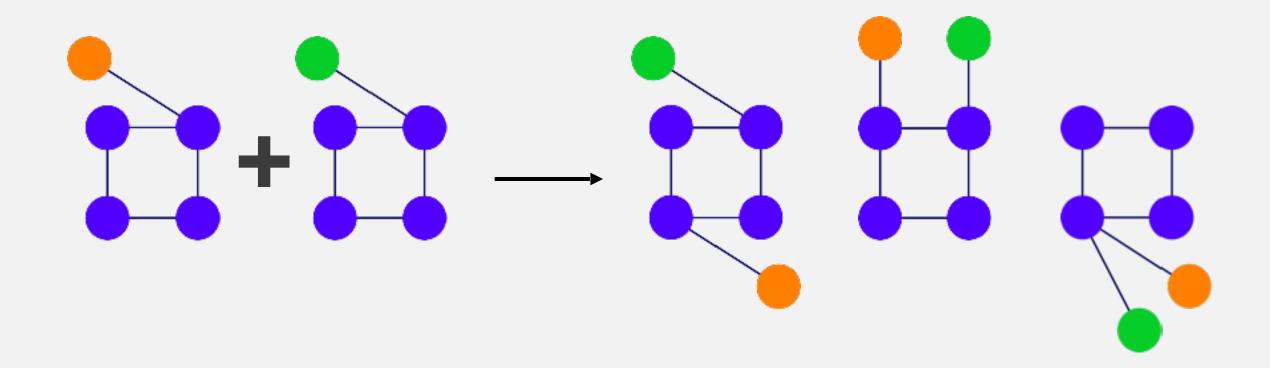
- We add G₁ and G₂ if removing an edge from G₁
 results in a subgraph in G₂
- This subgraph is called a core. G₁ and G₂ may share more than one!
- We add the removed edge from G₁ to G₂

Both candidates have one more vertex than the core

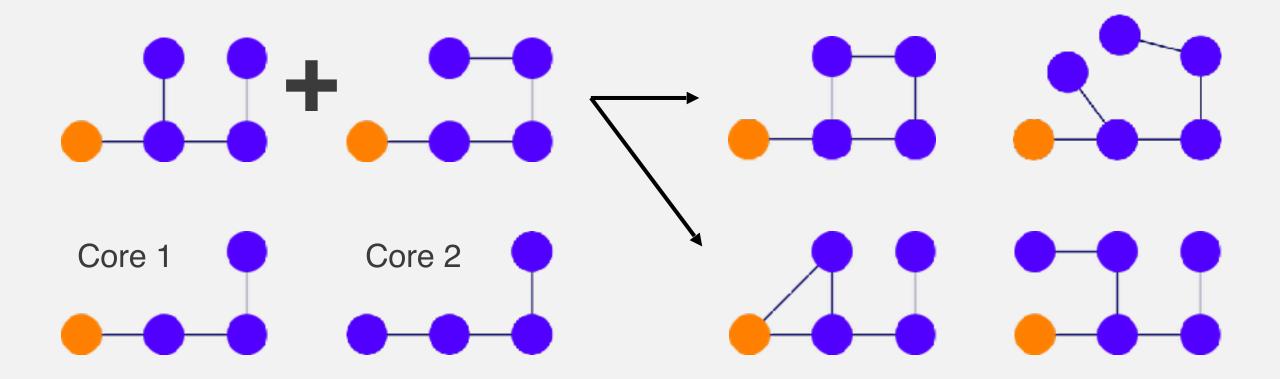
If these vertices have the same label, we add two candidates



We have to consider all automorphisms (versions) of the core!



We have to consider all cores!



FSG-PRUNE

- A priori property: a graph may only be frequent if all its subgraphs are!
- After pruning we need to scan to count the support of the candidates
- As always, there are possible alternatives, such as vertical data format

APRIORI—SUMMARY

Most data mining algorithms may be adapted to different problems!

- GSP: Apriory + sequences
 - Support count: time and order constraints
 - Candidate generation: contiguous sequences
- FSM: Apriori + graphs
 - Support count: isomorphisms
 - Candidate generation: core of two graphs

THANKS FOR LISTENING!