

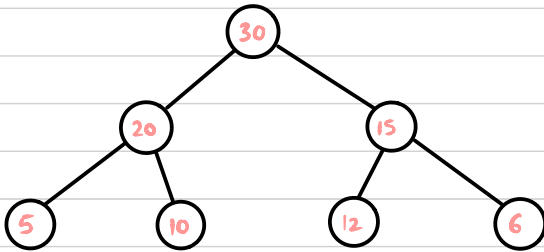


HEAP

↳ It is a complete binary tree

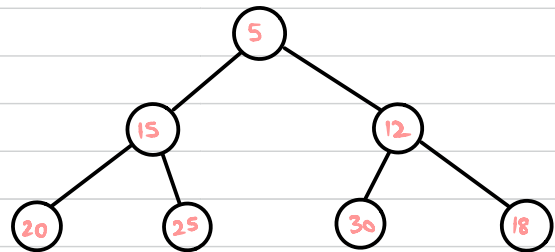
- what is a heap?
- Insert in a heap
- Deleting from heap
- Heap Sort
- Heapify
- Priority Queues

Max Heap



30	20	15	5	10	12	6			
1	2	3	4	5	6	7	8	9	10

Min Heap



- For complete binary tree, there should not be any free space between the elements of array

Node at index i

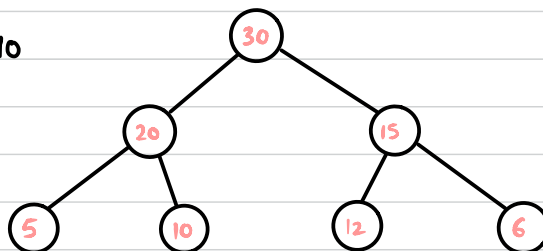
Left child at index $2*i$

Right child at index $2*i+1$

- Every node should have an element greater than or equal to all its descendants (duplicates can be there)

INSERTING IN HEAP

Element to insert = 40



30	20	15	5	10	12	6			
1	2	3	4	5	6	7	8	9	10

Step 1: Insert element at next free index in array

30	20	15	5	10	12	6	40		
1	2	3	4	5	6	7	8	9	10

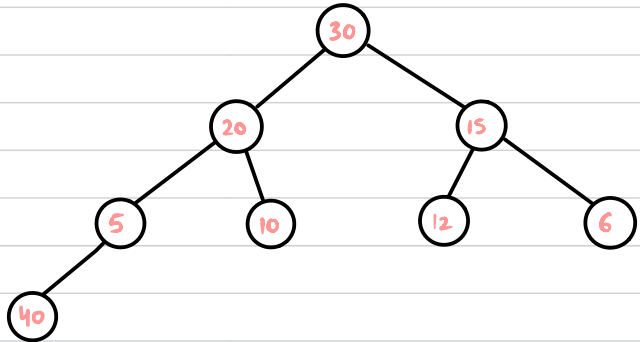
↑

Step 2: Insert the element in the tree

$$\text{Parent of } 40 = \frac{8}{2} = 4$$

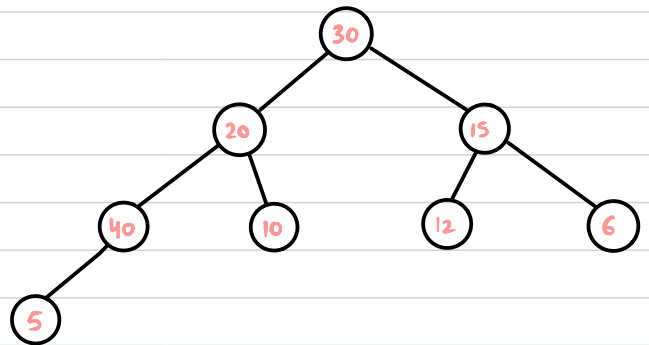
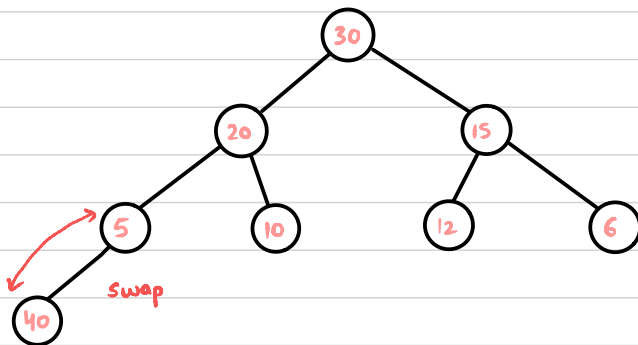
30	20	15	5	10	12	6	40		
1	2	3	4	5	6	7	8	9	10

↑

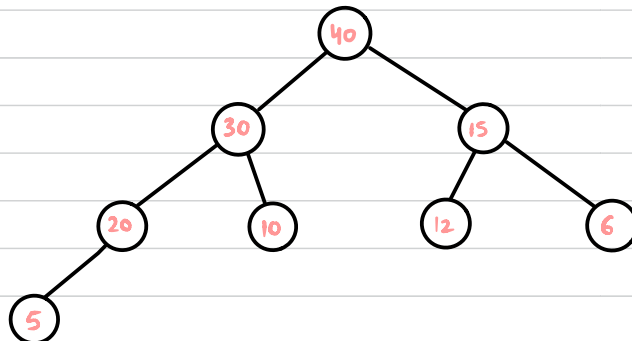


Step 3: Now tree does not satisfy max heap condition

Compare 40 with its parent/ancestor, if node is found greater than ancestor, swap.



Repeat until condition is satisfied



Check this condition in array also. Compare node with parent $(\frac{i}{2})$, if found greater than parent, replace

40	30	15	20	10	12	6	5		
1	2	3	4	5	6	7	8	9	10

PROGRAM

```
void Insert (int A[], int n)     $O(\log n)$ 
{
    int temp, i = n;
    temp = A[n];

    while (i > 1 && temp > A[i/2])
    {
        A[i] = A[i/2];
        i = i/2;
    }
    A[i] = temp;
}
```

```
void createheap()     $O(n \log n)$     1 element -  $\log n$ 
{
                                n elements -  $n \log n$ 
    int A[] = {0, 10, 20, 30, 25, 5, 40, 35};
    int i;

     $\rightarrow$  because first element is the heap and next element onwards,
    for (i = 2; i <= 7; i++)    insertion in heap is done.
        Insert (A, i);
}
```

Heapify: Faster method to implement
creation of binary heap

← Separate Topics



Element Priority

Elements \rightarrow 6, 8, 3, 10, 15, 2, 9, 17, 5, 8

- Element is itself a priority
- Smaller number higher priority

1. Insert in same order $O(1)$ $O(\log n)$
Delete max priority by searching it $O(n)$ $O(\log n)$

} After
implementing
using heap

Heap is used to implement priority queues

DELETION FROM HEAP Only the root node can be deleted from heap i.e the first node

```
void Delete (int A[], int n)
{
```

```
    int x, i, j;
```

```
    x = A[n];
```

```
    A[1] = A[n] // Smallest or last node is copied at first place
```

```
    i = 1; j = 2 * i; // i is at index 1 and j is at its left child
```

```
    while (j < n - 1)
```

```
    {
```

```
        if (A[j+1] > A[j]) // Right child is greater than left child
```

```
            j = j + 1;
```

```
        if (A[i] < A[j]) // Parent is smaller than child
```

```
        {
```

```
            swap(A[i], A[j]);
```

```
            i = j;
```

```
            j = 2 * j;
```

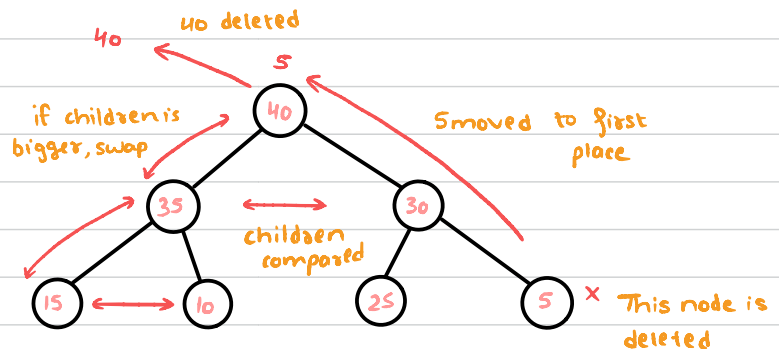
```
        }
```

```
    else
```

```
        break;
```

```
    A[n] = x;
```

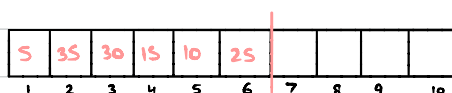
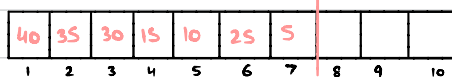
```
}
```



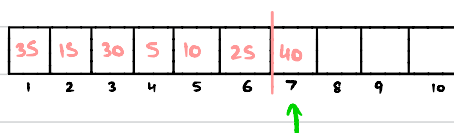
HEAP SORT

While deleting, store the deleted element at vacant position

Through this, you will get sorted array in ascending order after deleting all the elements



$x = 40$



STEPS

$n \log n$ • Create heap of 'n' elements

$n \log n$ • Delete 'n' elements 1 by 1
 $O(n \log n)$