

## A 부록: 파이썬 IR의 요약된 문법 및 의미구조

### A.1 요약된 문법

```
Program ::= stmt

stmt ::= Pass
        | Expr(expr)
        | Seq(stmt , stmt)
        | Assign(left-val , expr)
        | If(expr , stmt , stmt)
        | ForIn(left-val , expr , stmt)
        | Return(expr) | Continue | Break
        | Let(left-val, (expr | Undef) , stmt)
        | FuncDef(id , expr* , stmt)
        | Try(stmt, (exceptionclass | Undef , stmt)* , stmt , stmt)
        | Raise(exceptionclass , (expr | Undef))

expr ::= n ∈ ℤ | r ∈ Float | s ∈ String | True | False | None
        | object
        | BinOp(Binary-op , expr , expr)
        | UnaryOp(Unary-op , expr)
        | left-val
        | Call(expr.id , expr*)
        | LibCall(id , expr*)
        | exceptionclass

left-val ::= Name(id)
           | Attribute(expr , id)
           | Subscript(expr , expr)

exceptionclass ::= Except(Exception-id)

id ::= string

object ::= id | List(expr*) | Dictionary((expr , expr)*) | Tuple(expr*)

Numeric-op ::= + | - | / | // | % | **

Bool-op ::= and | or | is | is not | in | not in
```

$$\begin{aligned}
\textit{Compare-op} &::= == \mid != \mid < \mid > \mid <= \mid >= \\
\textit{Binary-op} &::= \textit{Numeric-op} \mid \textit{Bool-op} \mid \textit{Compare-op} \mid \textit{Bit-op} \\
\textit{Bit-op} &::= \mid \mid \& \mid ^ \mid << \mid >> \\
\textit{Unary-op} &::= \texttt{not} \mid - \\
\textit{Exception-id} &::= \textit{id} \mid \texttt{Exception} \mid \texttt{IndexError} \mid \texttt{ZeroDivisionError} \\
&\quad \texttt{KeyError} \mid \texttt{AttributeError} \mid \texttt{NameError}
\end{aligned}$$

## A.2 도메인

$$\begin{aligned}
\sigma &\in \textit{Env} = \textit{Id} \xrightarrow{fin} \textit{Addr} \\
\mathcal{H} &\in \textit{Heap} = \textit{Addr} \xrightarrow{fin} \textit{Value} \\
v &\in \textit{Value} = \textit{ValueExpr} + \textit{Addr} + \textit{Object} + \textit{Func} + \textit{Exception} \\
&\quad + \textit{Cont} + \{ \texttt{None}, \texttt{NotImpl}, \texttt{Undef} \} \\
s &\in \textit{ValueExpr} = \textit{String} + \textit{Boolean} + \textit{Float} + \textit{Int} \\
e &\in \textit{Exception} = \textit{Id} + \{ \texttt{Exception}, \texttt{NameError}, \texttt{AttributeError}, \\
&\quad \texttt{ZeroDivisionError}, \texttt{KeyError}, \texttt{IndexError} \} \\
o &\in \textit{Object} = (\textit{Id} + \textit{Int} + \textit{String}) \xrightarrow{fin} \textit{Value} \\
f &\in \textit{Func} = \textit{Id} \times \textit{Id}^* \times \textit{Stmt} \times \textit{Env} \\
\kappa &\in \textit{Cont} = \{ \texttt{run}, \texttt{cnt}, \texttt{brk} \} \\
l &\in \textit{Addr} = \textit{Int}(\textit{address space}) \\
i, x &\in \textit{Id} = \textit{String}(\textit{identifier}) \\
stmt &\in \textit{Stmt} = \texttt{Pass} + \dots
\end{aligned}$$

$$\text{"length", ""}(\text{empty string}), \dots \in \textit{String}$$

$$\_\_len\_\_, \_\_add\_\_, \_\_getitem\_\_, \dots \in \textit{Id}$$

요약 환경 ( $Env$ ) 는 식별자( $Id$ )와 요약 메모리 주소 ( $Addr$ )의 매핑(mapping)이다.

또한, 요약 메모리 ( $Heap$ ) 는 요약 메모리 주소 ( $Addr$ )와 요약 값 ( $Value$ )의 매핑

(mapping)이다. 요약 값은 String, Float, Int, Boolean 등의 값 (*ValueExpr*) 또는 요약 메모리 주소, 객체 (*Object*), 함수 (*Func*), 예외 클래스 (*Exception*), 프로그램 진행 값 (*Cont*), NULL 값 **None**, 구현 안된 값 **NotImpl**, 정의 안된 값 **Undef** 이다. 객체는 Id 또는 Int, String으로 부터 요약 값으로의 매핑이다. 함수는 함수 이름을 나타내는 식별자, 함수 인자를 나타내는 식별자, 함수 몸체 ((*Cont*)), 요약 환경으로 구성된다. 프로그램 진행 값은 프로그램 진행(**pass**, **run**), **continue(cnt)**, 프로그램 진행 중단(**break**, **brk**) 중 하나이다.

A.3의 의미구조에서 도메인 *Object*는 글씨체의 차이에 의해 구분된다. 도메인 *String* 는 딕셔너리 또는 `__getitem__` 함수의 도메인이다 (e.g., `obj["key"]`). *Id* 는 객체의 속성 및 `__getattr__` 함수의 도메인이다 (e.g., `obj.key`). *Int* 는 리스트나 `__getitem__` 함수의 도메인이다 (e.g., `obj[0]`).

클래스, 리스트, 딕셔너리 등 모든 파이썬의 객체들은 *Object*으로 나타낼 수 있다. 최적화를 위해서, *Object* 와 *Func*의 실제 구현에서는 함수의 디폴트 인자등의 추가적인 정보를 가질 수 있다.

```
class A:      # A.__mro__ = (A, object)
    pass
class B(A):   # B.__mro__ = (B, A)
    pass
b = B()      # b.__mro__ = (B, A)
```

그림 1: 클래스와 객체에서의 `__mro__` 호출 결과의 예

각 *Object*의 값마다 두 개의 빌트인 속성이 있다. *length*속성은 리스트와 같은 객체들에서 원소의 수와 같은 객체의 길이를 나타낸다. `__mro__`는 파이썬에서 Method Resolution Order(MRO)를 의미한다. 본 연구에서는 각 객체의 `__mro__` 속성마다 길이가 2인 튜플을 할당한다. 튜플의 첫번째 원소는 객체의 클래스이다. 객체가 클래스 일 때는 객체 자신이 첫 번째 원소가 된다. 두번째 원소는 첫번째 원소의 부모 클래스 이다. 실제 파이썬의 문법과는 다르게 요약된 문법에서는 다중 상속을 지원하지 않는다. 객체의 `__mro__` 속성들이 클래스

상속 트리에서 간선이 된다는 의미이다.

### A.3 의미구조

의미구조를 간략하게 나타내기 위해서 도우미 함수  $call$ ,  $attr$ ,  $item$ 를 아래와 같이 정의한다.

- $call(\mathcal{H}, f, v)$ : 힙 공간  $\mathcal{H}$ 아래에서 인자  $v$  로 함수  $f$ 를 계산한다 .

$$\frac{f = (x_f, x_1, S, \sigma_f) \quad l_f, l_1 \notin Dom(\mathcal{H}) \quad \sigma_f[x_f \mapsto l_f, x_1 \mapsto l_1], \mathcal{H}[l_f \mapsto f, l_1 \mapsto v_1] \vdash S \Rightarrow v', \mathcal{H}'}{call(\mathcal{H}, f, v) = v', \mathcal{H}'}$$

- $attr(\mathcal{H}, o, i)$ : 힙 공간  $\mathcal{H}$ 아래에서 객체  $o$ 의 속성  $i$  을 검색한다 . 속성  $i$ 가 객체  $o$ 의 속성이 아니면, 함수 `__getattr__`에 인자  $i$ 를 주어 호출한다. 만약 함수 `__getattr__`가  $i$ 에 없으면, 함수 `__mro__`를 활용해 기본 클래스 (base class)를 찾는다.
- $item(\mathcal{H}, o, i)$ : 위와 같으나  $i$  와 함수 `__getitem__`를 사용한다.

$$\begin{array}{c} \frac{o(i) = v}{attr(\mathcal{H}, o, i) = v, \mathcal{H}} \quad (\text{ATTR-DIRECT}) \end{array} \qquad \begin{array}{c} i \notin Dom(o) \\ o(\_\_getattr\_\_) = f \\ \frac{call(\mathcal{H}, f, i) = v, \mathcal{H}'}{attr(\mathcal{H}, o, i) = v, \mathcal{H}'} \\ (\text{ATTR-METHOD}) \end{array}$$

$$\begin{array}{c} i, \_\_getattr\_\_ \notin Dom(o) \\ o(\_\_mro\_\_)(0) \neq o \\ \frac{attr(\mathcal{H}, o(\_\_mro\_\_)(0), i) = v, \mathcal{H}'}{attr(\mathcal{H}, o, i) = v, \mathcal{H}'} \\ (\text{ATTR-INSTANCE}) \end{array} \qquad \begin{array}{c} i, \_\_getattr\_\_ \notin Dom(o) \\ o(\_\_mro\_\_)(0) = o \\ \frac{attr(\mathcal{H}, o(\_\_mro\_\_)(1), i) = v, \mathcal{H}'}{attr(\mathcal{H}, o, i) = v, \mathcal{H}'} \\ (\text{ATTR-CLASS}) \end{array}$$

### Expression의 의미구조

$$\boxed{\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}'}$$

$$\begin{array}{c}
\frac{E \in \mathbb{Z} + \text{Float} + \text{String} + \text{Bool} + \{ \text{None} \}}{\sigma, \mathcal{H} \vdash E \Rightarrow c, \mathcal{H}} \quad (\text{CONSTRAINTS}) \\
\\
\frac{\sigma, \mathcal{H} \vdash E \Rightarrow l, \mathcal{H}' \quad \mathcal{H}'(l) = v, v \notin \{ \text{Object}, \text{Undef} \}}{\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}'} \quad (\text{ADDRESS}) \\
\\
\frac{i \in \text{Dom}(\sigma)}{\sigma, \mathcal{H} \vdash \text{Name}(i) \Rightarrow \sigma(i), \mathcal{H}} \quad (\text{ID}) \\
\\
\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash E_1 \Rightarrow v_1, \mathcal{H}' \\ \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v_2, \mathcal{H}'' \\ o = \{0 \mapsto v_1, 1 \mapsto v_2, \text{"\$length"} \mapsto 2\} \\ l \notin \text{Dom}(\mathcal{H}'') \end{array}}{\sigma, \mathcal{H} \vdash \text{List}(E_1, E_2) \Rightarrow l, \mathcal{H}''[l \mapsto o]} \quad (\text{LIST}) \\
\\
\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash E_1 \Rightarrow v_1, \mathcal{H}' \\ \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v_2, \mathcal{H}'' \\ o = \{0 \mapsto v_1, 1 \mapsto v_2, \text{"\$length"} \mapsto 2\} \\ l \notin \text{Dom}(\mathcal{H}'') \end{array}}{\sigma, \mathcal{H} \vdash \text{Tuple}(E_1, E_2) \Rightarrow l, \mathcal{H}''[l \mapsto o]} \quad (\text{TUPLE}) \\
\\
\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash E_1 \Rightarrow v_1, \mathcal{H}' \\ \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v_2, \mathcal{H}'' \\ \sigma, \mathcal{H} \vdash E_3 \Rightarrow v_3, \mathcal{H}''' \\ \sigma, \mathcal{H}' \vdash E_4 \Rightarrow v_4, \mathcal{H}'''' \\ o = \{0 \mapsto v_1, 1 \mapsto v_2, 2 \mapsto v_3, 3 \mapsto v_4, \text{"\$length"} \mapsto 2\} \\ l \notin \text{Dom}(\mathcal{H}''') \end{array}}{\sigma, \mathcal{H} \vdash \text{Dictionary}((E_1, E_2), (E_3, E_4)) \Rightarrow l, \mathcal{H}''''[l \mapsto o]} \quad (\text{DICTIONARY})
\end{array}$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow v_1, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow v_2, \mathcal{H}'' \\
\hline
v_1, v_2 \in \text{Int} + \text{Float} + \text{Bool} \quad (\text{BINARY OPERATION}) \\
\sigma, \mathcal{H} \vdash \text{BinOp}(+, E_1, E_2) \Rightarrow v_1 + v_2, \mathcal{H}''
\end{array}$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow v_1, \mathcal{H}'' \\
\mathcal{H}''(l) = o, \text{attr}(\mathcal{H}'', o, \_ \_ \text{add} \_ \_) = f, \mathcal{H}''' \\
\text{call}(\mathcal{H}''', f, v_1) = v, \mathcal{H}'''' \quad v \neq \text{NotImpl} \\
\hline
(\text{BINARY OPERATION-LM}) \\
\sigma, \mathcal{H} \vdash \text{BinOp}(+, E_1, E_2) \Rightarrow v, \mathcal{H}''''
\end{array}$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l_1, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow l_2, \mathcal{H}'' \\
\mathcal{H}''(l_1) = o_1, \mathcal{H}''(l_2) = o_2, \_ \_ \text{add} \_ \_ \notin \text{Dom}(o_1) \\
\text{attr}(\mathcal{H}'', o_2, \_ \_ \text{radd} \_ \_) = f, \mathcal{H}''' \\
\text{call}(\mathcal{H}''', f, l_1) = v, \mathcal{H}'''' \\
\hline
(\text{BINARY OPERATION-RM}) \\
\sigma, \mathcal{H} \vdash \text{BinOp}(+, E_1, E_2) \Rightarrow v, \mathcal{H}''''
\end{array}$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l_1, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow l_2, \mathcal{H}'' \\
\mathcal{H}''(l_1) = o_1, \mathcal{H}''(l_2) = o_2, \text{attr}(\mathcal{H}'', o_1, \_ \_ \text{add} \_ \_) = f_1, \mathcal{H}''' \\
\text{call}(\mathcal{H}''', f_1, l_2) = \text{NotImpl}, \mathcal{H}'''' \\
\mathcal{H}''''(l_2) = o'_2, \text{attr}(\mathcal{H}'''', o'_2, \_ \_ \text{radd} \_ \_) = f_2, \mathcal{H}''''' \\
\text{call}(\mathcal{H}''''', f_2, l_1) = v, \mathcal{H}'''''' \\
\hline
(\text{BINARY OPERATION-RMN}) \\
\sigma, \mathcal{H} \vdash \text{BinOp}(+, E_1, E_2) \Rightarrow v, \mathcal{H}''''''
\end{array}$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow v_1, \mathcal{H}' \\
\hline
v_1 \in \text{Int} + \text{Float} + \text{Bool} \\
\sigma, \mathcal{H} \vdash \text{UnryOp}(-, E_1) \Rightarrow -v_1, \mathcal{H}'
\end{array}
\quad (\text{UNARY OPERATION})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}' \\
\mathcal{H}'(l) = o, \text{attr}(\mathcal{H}', o, \_\_ \text{neg} \_\_) = f, \mathcal{H}'' \\
\text{call}(\mathcal{H}'', f, l) = v, \mathcal{H}''' \quad v \neq \text{NotImpl} \\
\hline
\sigma, \mathcal{H} \vdash \text{UnryOp}(-, E_1) \Rightarrow -v, \mathcal{H}'''
\end{array}
\quad (\text{UNARY OPERATION-O})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow f, \mathcal{H}' \quad \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\text{call}(\mathcal{H}'', f, v) = v, \mathcal{H}''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Call}(E_1, E_2) \Rightarrow v, \mathcal{H}'''
\end{array}
\quad (\text{CALL})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow o, \mathcal{H}' \quad \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\text{attr}(\mathcal{H}'', o, \_\_ \text{call} \_\_) = f, \mathcal{H}''' \\
\text{call}(\mathcal{H}''', f, v) = v, \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Call}(E_1, E_2) \Rightarrow v, \mathcal{H}''''
\end{array}
\quad (\text{CALL-M})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow o, \mathcal{H}' \quad \sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\text{attr}(\mathcal{H}'', o, \_\_ \text{call} \_\_) = f, \mathcal{H}''' \\
\text{call}(\mathcal{H}''', f, v) = v, \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{LibCall}(E_1, E_2) \Rightarrow v, \mathcal{H}''''
\end{array}
\quad (\text{LIBCALL})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow l, \mathcal{H}' \\
\mathcal{H}'(l) = o, o(i) = v \\
\hline
\sigma, \mathcal{H} \vdash \text{Attribute}(E, i) \Rightarrow v, \mathcal{H}'
\end{array}
\quad (\text{ATTRIBUTE})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow l, \mathcal{H}' \\
\mathcal{H}'(l) = o, \text{attr}(\mathcal{H}', o, i) = v, \mathcal{H}'' \\
\hline
\sigma, \mathcal{H} \vdash \text{Attribute}(E, i) \Rightarrow v, \mathcal{H}''
\end{array}
\quad (\text{ATTRIBUTE-M})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow x, \mathcal{H}'' \\
\mathcal{H}''(l) = o, \text{ \_\_getitem\_\_ } \notin \text{Dom}(o) \\
\hline
x \in \text{Int} + \text{String}, o(x) = v \\
\hline
\sigma, \mathcal{H} \vdash \text{Subscript}(E_1, E_2) \Rightarrow v, \mathcal{H}''
\end{array}
\quad (\text{SUBSCRIPT})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\mathcal{H}''(l) = o, \text{attr}(\mathcal{H}'', o, \text{ \_\_getitem\_\_ }) = f, \mathcal{H}''' \\
\hline
\text{call}(\mathcal{H}''', f, v) = v', \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Subscript}(E_1, E_2) \Rightarrow v', \mathcal{H}''''
\end{array}
\quad (\text{SUBSCRIPT-M})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash e \Rightarrow v, \mathcal{H}' \\
v \in \{\text{Exception}, \text{NameError}, \text{AttributeError}, \\
\text{ZeroDivisionError}, \text{IndexError}, \text{KeyError}\} \\
\hline
\sigma, \mathcal{H} \vdash \text{Exception}(e) \Rightarrow v
\end{array}
\quad (\text{EXCEPTIONCLASS})$$

## Statement의 의미구조

$$\boxed{\sigma, \mathcal{H} \vdash S \Rightarrow v, \mathcal{H}'}$$

$$\begin{array}{c}
S \in \{\text{Pass}, \text{Continue}, \text{Break}\} \\
\hline
\sigma, \mathcal{H} \vdash S \Rightarrow \kappa, \mathcal{H}
\end{array}
\quad (\text{FLOW CONTROL})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \\
\hline
\sigma, \mathcal{H} \vdash \text{Expr}(E) \Rightarrow \text{run}, \mathcal{H}'
\end{array}
\quad (\text{EXCECUTE})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \\
\hline
\sigma, \mathcal{H} \vdash \text{Return}(E) \Rightarrow v, \mathcal{H}'
\end{array}
\quad (\text{RETURN})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash S_1 \Rightarrow \text{run}, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash S_2 \Rightarrow v, \mathcal{H}'' \\
\hline
\sigma, \mathcal{H} \vdash \text{Seq}(S_1, S_2) \Rightarrow v, \mathcal{H}''
\end{array}
\quad (\text{SEQUENCE-R})$$



$$\frac{\sigma, \mathcal{H} \vdash S_1 \Rightarrow \mathbf{brk}, \mathcal{H}'}{\sigma, \mathcal{H} \vdash \mathbf{Seq}(S_1, S_2) \Rightarrow \mathbf{brk}, \mathcal{H}'} \quad (\text{SEQUENCE-FB})$$

$$\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash S_1 \Rightarrow v, \mathcal{H}' \\ v \neq \mathbf{brk} \quad v \neq \mathbf{cnt} \\ \sigma, \mathcal{H}' \vdash S_2 \Rightarrow \mathbf{brk}, \mathcal{H}'' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{Seq}(S_1, S_2) \Rightarrow \mathbf{brk}, \mathcal{H}''} \quad (\text{SEQUENCE-SB})$$

$$\frac{\sigma, \mathcal{H} \vdash S_1 \Rightarrow \mathbf{cnt}, \mathcal{H}'}{\sigma, \mathcal{H} \vdash \mathbf{Seq}(S_1, S_2) \Rightarrow \mathbf{run}, \mathcal{H}'} \quad (\text{SEQUENCE-FC})$$

$$\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash S_1 \Rightarrow v, \mathcal{H}' \\ v \neq \mathbf{brk} \quad v \neq \mathbf{cnt} \\ \sigma, \mathcal{H}' \vdash S_2 \Rightarrow \mathbf{cnt}, \mathcal{H}'' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{Seq}(S_1, S_2) \Rightarrow \mathbf{cnt}, \mathcal{H}''} \quad (\text{SEQUENCE-SC})$$

$$\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash S_1 \Rightarrow v, \mathcal{H}' \\ v \neq \mathbf{run} \quad v \neq \mathbf{brk} \quad v \neq \mathbf{cnt} \\ \sigma, \mathcal{H}' \vdash S_2 \Rightarrow v', \mathcal{H}'' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{Seq}(S_1, S_2) \Rightarrow v', \mathcal{H}''} \quad (\text{SEQUENCE})$$

$$\frac{\begin{array}{c} f = (x_f, x_i, S_1, \sigma) \quad l \notin \text{Dom}(\mathcal{H}) \\ \sigma[x_f \mapsto l], \mathcal{H}[l \mapsto f] \vdash S_2 \Rightarrow v, \mathcal{H}' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{FunDef}(x_f, x_i, S_1, S_2) \Rightarrow v, \mathcal{H}'} \quad (\text{DEF FUNCTION})$$

$$\frac{\begin{array}{c} \sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \quad l \notin \text{Dom}(\mathcal{H}') \\ \sigma[i \mapsto l], \mathcal{H}'[l \mapsto v] \vdash S \Rightarrow v', \mathcal{H}'' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{Let}(i, E, S) \Rightarrow v', \mathcal{H}''} \quad (\text{LET})$$

$$\frac{\begin{array}{c} l \notin \text{Dom}(\mathcal{H}) \\ \sigma[i \mapsto l], \mathcal{H}[l \mapsto \mathbf{Undef}] \vdash S \Rightarrow v, \mathcal{H}' \end{array}}{\sigma, \mathcal{H} \vdash \mathbf{Let}(i, \mathbf{Undef}, S) \Rightarrow v, \mathcal{H}'} \quad (\text{LET-U})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \\
v \in \text{Int} + \text{Float} + \text{Boolean} + \text{String} + \text{Object} + \text{Func} \\
v \notin \{ \text{False}, \text{None}, 0, 0.0, "" \} \\
\sigma, \mathcal{H}' \vdash S_t \Rightarrow v_t, \mathcal{H}'' \\
\hline
\sigma, \mathcal{H} \vdash \text{If}(E, S_t, S_f) \Rightarrow v_t, \mathcal{H}''
\end{array}
\quad (\text{IF-TC})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \\
v \in \{ \text{False}, \text{None}, 0, 0.0, "" \} \\
\sigma, \mathcal{H}' \vdash S_f \Rightarrow v_f, \mathcal{H}'' \\
\hline
\sigma, \mathcal{H} \vdash \text{If}(E, S_t, S_f) \Rightarrow v_f, \mathcal{H}''
\end{array}
\quad (\text{IF-FC})$$

$$\begin{array}{c}
\sigma(i) = l \\
\sigma, \mathcal{H} \vdash E \Rightarrow v, \mathcal{H}' \\
\hline
\sigma, \mathcal{H} \vdash \text{Assign}(\text{Name}(i), E) \Rightarrow \text{run}, \mathcal{H}'[l \mapsto v]
\end{array}
\quad (\text{ASSIGN-N})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}' \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\mathcal{H}''(l) = o, o' = o[i \mapsto v] \\
\hline
\sigma, \mathcal{H} \vdash \text{Assign}(\text{Attribute}(E_1, i), E_2) \Rightarrow \text{run}, \mathcal{H}''[l \mapsto o']
\end{array}
\quad (\text{ASSIGN-A})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow l, \mathcal{H}_1 \\
\sigma, \mathcal{H}_1 \vdash E_2 \Rightarrow v_i, \mathcal{H}_2 \\
\sigma, \mathcal{H}_2 \vdash E_3 \Rightarrow v, \mathcal{H}_3 \\
\mathcal{H}_3(l) = o, \text{attr}(\mathcal{H}_3, o, \_\_ \text{setitem} \_\_) = \mathcal{H}_4, f \\
\text{call}(\mathcal{H}_4, f, (v_i, v)) = v', \mathcal{H}_5 \\
\hline
\sigma, \mathcal{H} \vdash \text{Assign}(\text{Subscript}(E_1, E_2), E) \Rightarrow \text{run}, \mathcal{H}_5
\end{array}
\quad (\text{ASSIGN-S})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow l, \mathcal{H}' \\
\mathcal{H}'(l) = o \quad \text{call}(\mathcal{H}', \_ \_ \text{len} \_ \_, o) = n, \mathcal{H}_0 \\
l' \notin \text{Dom}(\mathcal{H}_0) \quad \text{item}(\mathcal{H}_0, o, 0) = v_0, \mathcal{H}'_0 \\
\sigma[i \mapsto l'], \mathcal{H}'_0[l' \mapsto v_0] \vdash S \Rightarrow \kappa_1, \mathcal{H}_1 \\
\vdots \\
\text{item}(\mathcal{H}_{n-1}, o, n-1) = v_{n-1}, \mathcal{H}'_{n-1} \\
\sigma[i \mapsto l'], \mathcal{H}'_{n-1}[l' \mapsto v_{n-1}] \vdash S \Rightarrow \kappa_n, \mathcal{H}_n \\
\forall i, \kappa_i \in \{\mathbf{run}, \mathbf{cnt}\} \\
(\text{o가 S안에서 바뀌지 않는다고 가정}) \\
\hline
\sigma, \mathcal{H} \vdash \mathbf{ForIn}(i, E, S) \Rightarrow \mathbf{run}, \mathcal{H}_n
\end{array}
\quad (\mathbf{FORIN})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow l, \mathcal{H}' \\
\mathcal{H}'(l) = o \quad \text{call}(\mathcal{H}', \_ \_ \text{len} \_ \_, o) = n, \mathcal{H}_0 \\
l' \notin \text{Dom}(\mathcal{H}_0) \quad \text{item}(\mathcal{H}_0, o, 0) = v_0, \mathcal{H}'_0 \\
\sigma[i \mapsto l'], \mathcal{H}'_0[l' \mapsto v_0] \vdash S \Rightarrow \kappa_1, \mathcal{H}_1 \\
\vdots \\
\text{item}(\mathcal{H}_{k-1}, o, k-1) = v_{k-1}, \mathcal{H}'_{k-1} \\
\sigma[i \mapsto l'], \mathcal{H}'_{k-1}[l' \mapsto v_{k-1}] \vdash S \Rightarrow \mathbf{brk}, \mathcal{H}_k \\
\forall i, \kappa_i \in \{\mathbf{run}, \mathbf{cnt}\} \\
(\text{o가 S안에서 바뀌지 않는다고 가정, } k \leq n) \\
\hline
\sigma, \mathcal{H} \vdash \mathbf{ForIn}(i, E, S) \Rightarrow \mathbf{run}, \mathcal{H}_k
\end{array}
\quad (\mathbf{FORIN-B})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow e, \mathcal{H}' \\
e \in \{\mathbf{Exception}, \mathbf{NameError}, \mathbf{AttributeError}, \\
\mathbf{ZeroDivisionError}, \mathbf{IndexError}, \mathbf{KeyError}\} \\
\hline
\sigma, \mathcal{H} \vdash \mathbf{Raise}(E, \mathbf{Undef}) \Rightarrow \mathbf{brk}, \mathcal{H}'
\end{array}
\quad (\mathbf{RAISE-U})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E_1 \Rightarrow e, \mathcal{H}' \\
e \in \{\text{Exception}, \text{NameError}, \text{AttributeError}, \\
\text{ZeroDivisionError}, \text{IndexError}, \text{KeyError}\} \\
\sigma, \mathcal{H}' \vdash E_2 \Rightarrow v, \mathcal{H}'' \\
\hline
\sigma, \mathcal{H} \vdash \text{Raise}(E_1, E_2) \Rightarrow \text{brk}, \mathcal{H}''
\end{array}
\quad (\text{RAISE})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow e, \mathcal{H}' \quad e \in \{\text{Undef}, \text{Exception}\} \\
\sigma, \mathcal{H}' \vdash S_1 \Rightarrow \text{brk}, \mathcal{H}'' \\
\sigma, \mathcal{H}'' \vdash S_2 \Rightarrow v', \mathcal{H}''' \\
\sigma, \mathcal{H}''' \vdash S_4 \Rightarrow v'', \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Try}(S_1, E, S_2, S_3, S_4) \Rightarrow v'', \mathcal{H}''''
\end{array}
\quad (\text{TRY-UX})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow e, \mathcal{H}' \quad e \in \{\text{Undef}, \text{Exception}\} \\
\sigma, \mathcal{H}' \vdash S_1 \Rightarrow v, \mathcal{H}'' \\
\sigma, \mathcal{H}'' \vdash S_3 \Rightarrow v', \mathcal{H}''' \\
\sigma, \mathcal{H}''' \vdash S_4 \Rightarrow v'', \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Try}(S_1, E, S_2, S_3, S_4) \Rightarrow v'', \mathcal{H}''''
\end{array}
\quad (\text{TRY-UNX})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow e, \mathcal{H}' \quad e \notin \{\text{Undef}, \text{Exception}\} \\
\sigma, \mathcal{H}' \vdash S_1 \Rightarrow \text{brk}, \mathcal{H}'' \\
\sigma, \mathcal{H}'' \vdash S_2 \Rightarrow v', \mathcal{H}''' \\
\sigma, \mathcal{H}''' \vdash S_4 \Rightarrow v'', \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Try}(S_1, E, S_2, S_3, S_4) \Rightarrow v'', \mathcal{H}''''
\end{array}
\quad (\text{TRY-X})$$

$$\begin{array}{c}
\sigma, \mathcal{H} \vdash E \Rightarrow e, \mathcal{H}' \quad e \notin \{\text{Undef}, \text{Exception}\} \\
\sigma, \mathcal{H}' \vdash S_1 \Rightarrow v, \mathcal{H}'' \\
\sigma, \mathcal{H}'' \vdash S_3 \Rightarrow v', \mathcal{H}''' \\
\sigma, \mathcal{H}''' \vdash S_4 \Rightarrow v'', \mathcal{H}'''' \\
\hline
\sigma, \mathcal{H} \vdash \text{Try}(S_1, E, S_2, S_3, S_4) \Rightarrow v'', \mathcal{H}''''
\end{array}
\quad (\text{TRY-NX})$$