## E.Range Sums

题目给出得元素和可以看做成 i[r] - i[l-1] (i[x]代表前x元素之和) 要能够得到完整序列,即i[N] - i[0] 将 0~N 看作顶点,有Q条无向边链接这些顶点只要 0-N 可达即可.可以使用bfs或并查集来解决.

```
const int max_n = 3e5+10;
int N,Q;
struct my_union{
        int father[max_n] = \{0\};
        my_union(int N){
            for(int i = 0; i < N; i++){
                father[i] = i;
            }
        int find(int x){
            if(x != father[x])
                father[x] = find(father[x]);
            return father[x];
        void unionset(int a,int b){
            a = find(a);
            b = find(b);
            if(a == b)
                return;
            father[a] = b;
        bool judge(int a,int b){
            return find(a) == find(b);
        }
    };
int main(){
    scanf(" %d %d",&N,&Q);
    my_union tree(N+1);
    int 1,r;
    for(int i = 1; i <= Q; i++){
        scanf(" %d %d",&1,&r);
        tree.unionset(l-1,r);
    printf("%s\n",tree.judge(0,N) ? "Yes" : "No");
    return 0;
}
```