This document includes solution hints, but **NOT** necessarily the full solution!

Part I

Evaluation Page Part A

N T	G. 1
Name:	Signature:
ivanic.	Digitature.

1 Evaluation

This part of the exam has 270 questions, with a total of 447 points and 4 bonus points.

Part	Max. Points	Scored Points
A	60	
В	180	
	Total	

Points:	240-220	219-201	200-182	181-162	161-144	143-0
Grade:	A	В	С	D	Е	F
Score:						

Evaluation Part A 2

Page	Points	Bonus Points	Score
6	5	0	
7	7	2	
8	13	0	
9	12	0	
10	6	0	
11	8	0	
12	4	0	
13	13	0	
14	10	0	
15	8	0	
16	5	0	
17	3	2	
18	7	0	
19	2	0	
20	3	0	
21	2	0	
23	6	0	
24	5	0	
26	3	0	
27	5	0	
28	5	0	
29	9	0	
30	11	0	
31	13	0	
32	9	0	
33	6	0	
Total:	180	4	

Page	Points	Bonus Points	Score
34	9	0	
35	7	0	
36	6	0	
37	7	0	
38	1	0	
39	8	0	
40	10	0	
41	11	0	
42	10	0	
43	8	0	
44	8	0	
45	8	0	
46	13	0	
47	17	0	
48	15	0	
49	14	0	
50	13	0	
51	10	0	
52	16	0	
53	12	0	
54	8	0	
55	8	0	
56	8	0	
57	7	0	
58	6	0	
59	6	0	
60	7	0	
61	7	0	
62	6	0	
63	1	0	
	1		

Page 3 of of that:

267

0

Part II

Rules

Answer the questions within the space provided. If you do not have enough space, you can use the backside of the sheet. In that case clearly indicate that your answer continues on the backside.

3 Supporting Materials

This is an examination in writing, without the usage of any electronic devices, except a scientific pocket calculator. No restriction on the model of calculator that may be used, but no device with communication capability shall be accepted as a calculator. All other electronic devices are prohibited. Writing paper is available, writing instruments (pencil, pens, etc) have to be organized by the student.

- Part A: Without any supporting material, with calculator.
- Part B: With a self written summary (format A4, 8 sheets or up to 16 pages), with calculator

4 Procedure

1. Duration:

Part A: 1 hour = 60 minutes = 60 points.

(short break)

Part B: 3 hours = 180 minutes = 180 points.

- 2. Sign the first page in the provided space. With this you certify that you are only using permitted support material and you are complying to the rules.
- 3. Write your name on any detached or additional paper sheets. Sheets without a name will not be evaluated.
- 4. Use the provided paper for your solutions. Use the provided space in the forms and tables. If needed use scratch paper. Document your way to your solution as appropriate.
- 5. Each question has a defined number of maximum points associated.
- 6. If a question is unclear, make reasonable assumptions. Document your assumptions and provide a rationale.
- 7. Write clearly and legibly. Unclear or multiple solutions will not be evaluated.
- 8. There is a short break between part A and B. You have to sign into a list for a needed break during the examination parts. Only one person can leave the room for a short time.
- 9. If something is unclear, ask your supervisor in the room.

5 Time Management

Read first all questions. Make sure you distribute your available time to all the questions. To reduce disturbance, ask questions in the first 15 minutes of the exam period.

6 Multiple-Choice Questions

- 1. Try to answer all questions if possible. If you are not sure, choose the answer which seems the best one.
- 2. For the questions of type \bigcirc : Choose **exactly one** option with \otimes (or $\sqrt{\ }$), which you think is the best match. With a correct answer you get the given number of points for that question.
- 3. For the questions of type ±: After a question or possibly incomplete sentence there are four answers or extensions. Evaluate each of them if they are true or false and mark them accordingly with '+' (true) or '-' (false). Independent if the question is formulated grammatically in singular or plural, it is possible that 0, 1, 2, 3, 4 of the choices are true. For three correct answers out of four you receive half of the points.
- 4. Wrong answers will have no penalty. Each question which has no answer is treated like a wrong answers and will be evaluated with zero points.
- 5. If you are changing your mind: cross out your old answer and clearly mark which answer is the new one.

May Dilbert be with you! ©

Question 1
$\sqrt{\text{Learning summary with 5 questions.}}$
Sumo robot PCB capacitor.
Collection of slides.
Line sensor capacitance.
O Tips for students in next semester.
Question 2
O Variable Capacity System.
O Volatile Control Status.
$\sqrt{\text{Version Control System.}}$
○ Variable Computer Software.
O Volatile Client Storage.
Question 3
$\sqrt{ m ARM~Cortex-M0}+$
○ HCS08
\bigcirc MMA8780Q
○ ARM Cortex M4
○ MCP4728
Question 4
<pre>. gitignore readme.txt list.txt src\rotor.c src\rotor.h obj\rotor.o obj\rotor.txt</pre>
The .gitignore file has following content:
/ o b j /* . t x t !/ r*

In above directory listing, strike through the files which are ignored.

Page 6 of 63 Reached: _____

Solution: not ignored are: .gitignore, readme.txt, src/main.c, src/main.h Explain in a single sentence what each of the following basic VCS actions mean in Git: (a) Committing $\frac{1}{2}$ **Solution:** Putting a change into the local repository. (b) Reverting $[1/_{2}]$ **Solution:** Undo a local change. (c) Pushing $[\frac{1}{2}]$ **Solution:** Moving a local change into the remote repository. (d) Cloning $[1/_{2}]$ **Solution:** Copy a repository and create a new local one. What is the fundamental difference between SVN and Git? Solution: SVN: centralized VCS, Git: distributed VCS. Explain the difference between the optimistic and pessimistic approach in a VCS. Explain it with an example. **Solution:** Optimistic: assumes that two developers do not work on the same file, so system potentially allows conflicts. Conflicts have to be resolved later. Pessimistic: assumes that conflict will happen, and whenever a developer wants to edit a file, the file gets locked so no conflict can occur. This could be your bonus question you have submitted... ②. 8. ves or no? [1] (a) Provide an example for a hard real-time system: **Solution:** Air bag, pacemaker or aircraft control system. (b) Provide an example for a *soft* real-time system: [1]

Page 7 of 63 Reached: ____

Solution: Video streamer.
Question 10
(a) Name some benefits implementing a state machine:
Solution: well structured, easy to implement, reusable design pattern.
(b) What should be the first steps when implementing a state machine?
Solution: Define all states with input/output and transitions. Draw a diagram.
Question 11
(a) Why did we implement it as an array of bits?
Solution: To save RAM.
(b) What is the fundamental disadvantage of such an array of bits?
Solution: Costs runtime performance for bit manipulation.
(c) It implements critical section (e.g. to set an event bit) with EnterCritical() and ExitCritical(). Under which conditions such a critical section would no be required?
Solution: If there is not a possiblity that the operation gets interrupted, or if the operation is atomic.
(d) List reasons why an interrupt service routine should use such an Event module
Solution: Reduce interrupt latency time, only setting a bit and let the main application do the heavy lifting.
Question 12
Solution: Pre-emptive: Always runs the highest available task. Tasks of identical priority share CPU time Cooperative: Context switches only occur if a task blocks, or explicitly calls yield.
Question 13
(a) In an RTOS, each task can be in one of 5 fundamental states: List them:
Solution: New, Ready, Running, Waiting, Stopped.
(b) What's the purpose of the scheduler in an RTOS?

[2]

[2]

Solution: To determine which tasks gets executed next to minimize waiting time.

Provide a short definition of the term *Interrupt Latency*, and which factors/aspects are contributing to it:

Solution: Time between the event itself and until the ISR executes. Factors are stopping/finishing the current interrupt, interrupt destination calculation/arbitration, pushing state and diverting to the ISR.

Question 15.......Points: [2]

Provide an example of a typical *Reactive System*, and explain why this is a reactive system:

Solution: Airbag, it reacts on external events.

(a) A PWM signal on a H-Bridge is labeled as *low active*. Explain what this means and how this impacts the speed of a DC motor:

Solution: low active means that the motor is active when the signal is low. It means for the PWM duty cycle: the longer in the low state, the higher the voltage, the faster the motor turns.

(b) Draw a timing diagram for that PWM signal: the PWM period is 5 ms, and the motor shall at 20% speed. Indicate how many milliseconds the signal is high and low.

Solution: [Timing diagram drawn here, with a frequency of 5 ms and 20% low duty cycle (1 ms low, 4 ms high)].

Question 17......Points: [3] Given the following program:

```
#define ADC_CONFIG (*(volatile uint8_t*)0x123)

static void Interrupt(void) {
    uint8_t i;

    while (ADC_CONFIG & ~0x10);
    for (i = 0; i < 10; i++) {
        __asm("nop");
    }
}</pre>
```

This program is using

Page 9 of 63

Question 17 continuous on the next page...

Reached: ____

Your Eclipse project stores the make files, object files and the final (binary) application file in a sub folder inside your project. Are you going to store this folder and files in a version control system? Justify your answer: Solution: No, as the content of this folder is generated. It does not make sense to store derived content in a version control system, as it can be generated from the sources.
 Gadfly synchronization. Realtime synchronization. ✓ Realtime and Gadfly synchronization. ○ No synchronization. Question 18
 ○ Realtime synchronization. ✓ Realtime and Gadfly synchronization. ○ No synchronization. Question 18
✓ Realtime and Gadfly synchronization. ○ No synchronization. Question 18
O No synchronization. Question 18
Your Eclipse project stores the make files, object files and the final (binary) application file in a sub folder inside your project. Are you going to store this folder and files in a version control system? Justify your answer: Solution: No, as the content of this folder is generated. It does not make sense to store derived content in a version control system, as it can be generated from the sources. Question 19
store derived content in a version control system, as it can be generated from the sources. Question 19
Given the source of a PID control loop implementation. Identify in the source lines for the P, I and D part: mark them clearly and label it with P, I and D. Mark/circle this in the following source listing: #define max 0x33ff static int32_t old=0, b=0;
\mathbf{static} int32_t old=0, b=0;
<pre>int32_t f, s, a; v = 0; f = should-actual; a = f-old; old = f; v += a/10; v += f*35; b += f; if (b > max) { b = max; } v += b/4; setAcuator(v); }</pre>

Solution: D, then P, then I with anti-windup.

(a) 2 typical *Methods* for an ADC component:

[1]

What is the difference between the two following usages

```
sizeof(string)
```

```
strlen(string)
```

in respect to the result and the expected code generated?

Solution: sizeof gives the size in memory, which is here the size of a pointer (2 or 4 bytes, depending on the machine), while strlen() the length without the zero byte. sizeof() is calculated at compile time (constant), while strlen() is a library routine call.

Question 22......Points: [3] Given following interface implementation for a DC motor driver:

```
/* motor.h */
#include "LED.h" /* LED interface */
#include "PWM.h" /* PWM interface */

static uint16_t MOT_motorSpeed;

void MOT_Init(void);
/* end of motor.h */
```

Identify three issues with such an interface implementation (issues which could lead to linker/compiler failure, or things which are not considered as good programming style):

Page 11 of 63 Reached: _____

Solution: No #ifndef...#define, not necessary includes, static definition in header file.

```
Given following program:
```

```
#define DEC(i) {int b=0; i--;}

void main(void) {
  int a = 5, b = 5;
  DEC(a);
  DEC(b);
  printf("a is %d, b is %d\n", a, b);
}
```

What is the output of printf()?

```
Solution: a is 4, b is 5
```

Question 24.......Points: [3]

Given the *Mealy Sequential State Machine* in Figure 1 with five states, two input values and two output values:

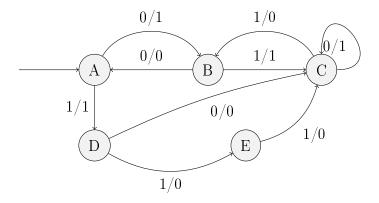


Figure 1: Mealy Sequential State Machine

(a) The state machine in Figure 1 is in the state ${}^{\prime}\mathbf{B}{}^{\prime}$. Determine the input sequence in order to generate the following output:

1, 1, 1, 0

```
Solution: B: input1 output1 -> C
C: input0 output1 -> C
C: input0 output1 -> C
C: input1 output0 -> B
Solution: 1001
```

Page 12 of 63

Question 24 continuous on the next page...

Reached: _____

[2]

(b)	The State Machine in Figure 1 is not complete and has an undefined transition from one state to another: fix this with a solution in Figure 1.
	Solution: Node E needs to have another outgoing arrow with $0/$? (e.g. to node E).
ıesti	on 25Points: [4]
(a)	Provide an example of a typical <i>Transforming System</i> , and explain why this is a Transforming System:
	Solution: Network router, it transform packes and distributes them.
(b)	Explain why $Optimized\ Memory\ Usage$ is a typical attribute for a $Transforming\ System$:
	Solution: Such systems transform an input stream into an output stream, and this usually involves larger amount of memory for buffering and transforming. As memory is expensive, such systems need to be optimized for this.
Exp	on 26
stac	ck, and some only push a subset of the registers.
So	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow.
So reg	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. on 27
So reg	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. On 27
So reg testi (a)	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. on 27
So reg nesti (a)	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. on 27
So regularization (a)	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. on 27
So reg	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. On 27
reguesti (a) (b)	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. On 27
So reguesti (a) (b) (c)	lution: In order not to increase the interrupt latency in case there are many gisters, and as well to reduce the chance for stack overflow. On 27

Page 13 of 63

Reached: _____

estion 29 List the things a processor has to do in ord and to return from it.	L
Solution: Stop actual instruction (or unconvector, store status on stack, branch thing in reverse order to return from the I	o ISR and context switch. Do the same
estion 30 Explain multiple things which affects the i	L
Solution: Latency time depends on the spisters/stack to be changed for the context until the ISR routine can start.	,
estion 31 What does it mean, if somebody says "I have been says to be says to be says to be says."	•
Solution: The interrupts are disabled.	
Doog the ADM Context MO	Points: [2
Does the ARM Cortex M0+ support neste	d interrupts? 32No
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number).	d interrupts? 32. No terrupt priority number (higher interrupt
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	d interrupts? 32. No terrupt priority number (higher interrupt an interrupt source with higher priority priority number (higher priority an interrupt source with higher priority number (higher prio
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number).	d interrupts? 32. No terrupt priority number (higher interrupt an interrupt source with higher priority
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	d interrupts? 32. No terrupt priority number (higher interrupt an interrupt source with higher priority
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	an interrupt source with higher priority
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	an interrupt source with higher priority
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	an interrupt source with higher priority
Solution: Yes, an interrupt with lower in priority number) can be interrupted by (lower interrupt priority number). estion 33	an interrupt source with higher priority

(d) What gives sizeof(MyVar) for the Robot bo	${ m ard/project}$?	
	(d)	3	
(e) Which value has YELLOW?			
	(e)	7	
estion 34		Points	: [3]
uint16_t abcd[16]; uint8_t buf[10]; static uint16_t values[3];			
(a) Determine the value of following expression: sizeof("abcd")			
	(a)	5	
(b) Determine the value of following expression: sizeof(buf)			
	(b)	10	
(c) Determine the value of following expression: sizeof(values)			
	(c)	6	
estion 35			
1. Linked Folder			
2. Linked Files			
3. Virtual Group			
List pros and cons for each approach:			
Solution: Linked folder: Pros: new files in that the project. Cons: you get all or nothing. Linked Files: Pros: You can decide for each remote you need to do this for every file. Virtual Group: Pro: arbitrary group of files. Confor the build tools settings	e file if it is	included or not. C	ons:

Page 15 of 63 Reached: _____

A hard realtime system or a soft realtime system: which do you consider easier to implement and test? List one pro and one cons for each:

Solution: A hard realtime system is probably harder to implement, but it is easier to test, as if you can make it fail a deadline, it is clear that it fails. A soft realtime system is probably easier to implement as it does not have to stick to hard deadlines, but it will be more difficult to test.

Consider following source:

```
#define MACRO(a,b) a = j \
=b
```

Write down the text which would be produced by the preprocessor of the compiler, if you call the MACRO as following:

```
MACRO(i,5);
```

Solution:

```
MACRO(i,5);

i = j = 5;
```

Given following C source:

```
#define MACRO(var, mask1, mask2) \
  (var = (var & (~(uint8_t)(mask1))) | (uint8_t)(mask2))
  static uint8_t var;

void foo(void) {
  var = 0x22;
  MACRO(var, 16, 0x13);
}
```

What is the value of var after execution of foo()?

38	0x33 oder	51
	UNOU OUCI	OI

Page 16 of 63 Reached: _____

Solution:

```
\begin{array}{l} var = 0x22;\\ (var = (var \& (~(uint8\_t)(16))) \mid (uint8\_t)(0x13));\\ var = (var \& (~(uint8\_t)(0x10))) \mid (uint8\_t)(0x13);\\ \Longrightarrow clear \ bits \ in \ mask1, \ set \ bits \ in \ mask2\\ \Longrightarrow clear \ bit \ 0x10, \ set \ bits \ 0x13\\ \Longrightarrow clear \ has \ no \ effect \ , \ so \ it \ is \ 0x22 \mid 0x13 \Longrightarrow 0x33 \\ \end{array}
```

Enable Pull-Up Resistors;

√ Set Port direction register as input;
Enable Pull-Up Resistors;
Acknowledge Pending Interrupt;
Enable Keyboard Interrupts:

Enable Keyboard Interrupts;

Enable Pull-Up Resistors;
 Enable Keyboard Interrupts;
 Acknowledge Pending Interrupt;
 Set Port direction register as input;

- \ominus ± Interrupts have to be enabled globally during the driver initialization.
- \oplus \pm The driver shall reset the device interrupt flag during initialization.
- \oplus ± After a power-on reset, it might be necessary to wait a certain time until the hardware signals have stabilized.
- \ominus ± The interrupt handler shall be as efficient as possible in order to increase the interrupt latency time.

Question 42......Points: [1] Given following program:

```
void main(void) {
unsigned char *src=(unsigned char*)0x100, buffer[0x100], i;
for(i=0;i<100; i++) {
   buffer[i]=*src;
}
}</pre>
```

For the above program, following applies:

- \ominus ± It reads the values from the address 256 and 512 and stores it in a buffer.
- \oplus ± It reads 100 times the value at the address 0x100 and stores the values one after each other in a buffer.
- \ominus ± At termination of the program, the whole buffer is filled with the values from address 0x100.
- \oplus ± With disabled interrupts, the program behaves in a deterministic way.

Solution: 3rd Answer: only part of the buffer ([0]..[99] is filled with *0x100, but the buffer has the size of [0x100].

Question 43......Points: [7] Given the Mealy Sequential State Machine in Figure 2.

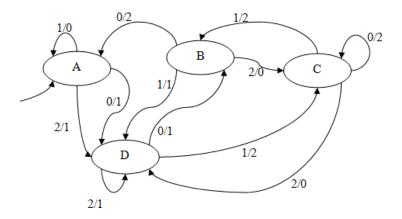


Figure 2: Mealy Machine

(a) The machine in Figure 2 is currently in state 'C'. Determine the output sequence for following input values: 0, 1, 0, 1, 1, 0

Solution: 2, 2, 2, 0, 0, 1

(b) Given following Mealy program:

typedef enum {A=0, B, C, D} States;
void Run(void) {
 char j, i = 0
 for (;;) {

Page 18 of 63 Reached: _____

[1]

[6]

```
j = Input();
Output(tbl[i][j][1]);
i = tbl[i][j][0];
}
```

To implement the machine in Figure 2, complete the initialization of table tbl:

```
const char tbl[4][3][2] = {
```

Solution:

For realtime systems following applies:

- \ominus ± Realtime systems have to have reaction times below 1 ms in order to be realtime compliant.
- \oplus \pm For a realtime system not the average system load matters, but the highest possible system load.
- \ominus ± Hard realtime systems are more difficult to verify, because the realtime conditions are not exactly specified.
- \ominus ± A system can be a realtime system, if it is using true random number generator for its decision instead of a pseudo random number generator.

Given following program:

```
char buf[0x100];
int i,j;

static void test(void) {
  for(i=0; i<sizeof(buf); i++) {
    CFG = 0x80; PORTB = 4;
    buf[i] = PORTA;
    PORTB = 0;
  }
}</pre>
```

For this program following applies:

- $\ominus \pm$ Implements an interrupt synchronization.
- $\ominus \pm$ Implements a gadfly synchronization.
- \ominus ± Implements a realtime synchronization.
- $\oplus \pm$ None of above.

For all reentrant functions in C, following has to apply:

- \ominus ± The function shall not be recursive.
- \ominus ± The function shall not be called from an ISR.
- \oplus \pm The access to shared data has to be protected from mutual access.
- \oplus ± The function shall not modify itself (self modifying code).

Question 47......Points: [1]

The following program gets compiled for the FRDM board with default compiler options:

```
static char ch;  /* Linker places this variable at address 0x10 */
void foo(void) {
    static char i, j=4;
    volatile char v;
    v = i;
    v = j;
    ch = v;
}
```

Following applies:

- \ominus ± The variables i, j and v are allocated on the stack.
- \oplus \pm The compiler cannot optimize the two assignments to v because of volatile.
- \ominus ± At execution time of foo(), the variable v gets initialized with a value of 4.
- \oplus ± After execution of foo(), the memory at address 0x10 will have a value of 4.

Question 48.......Points: [1]

For the interrupt system of the ARM Cortex-M0+/M4 following applies:

- \oplus \pm The interrupt latency is the sum of execution time of the current instruction, pushing of the registers, calculating the ISR PC address and the branching to the ISR itself.
- \ominus ± With 'masking the interrupts' we are enabling the interrupts.
- \oplus ± In order for the ISR program to return to the interrupted program, the return address of the interrupted program is stored on the stack by the hardware.

 \oplus ± In order to reduce the interrupt latency time, the core can decide not to push all registers on the stack.

Question 49.......Points: [2]

The diagram in Figure 3 shows an interrupt system with multiple interrupts (IRQ1 and IRQ2) and the corresponding interrupt service routines (ISR) #1 and #2). The lines on the time axis denote the execution time boundaries of the instructions. Fol-

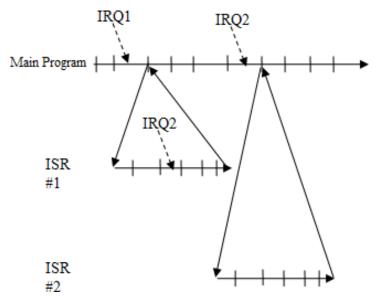


Figure 3: Interrupts

lowing applies:

- At the beginning of ISR #1 all interrupts get disabled, and at the end of ISR #1 the flag for IRQ1 gets acknowledged.
- \bigcirc The main program has at the beginning all interrupts disabled and has the IRQ1 flag acknowledged. After execution of ISR #1 the main program enables all interrupts.
- $\sqrt{\ \text{ISR}\ \#1}\ \text{turns}$ off all interrupts at the beginning. At the end of ISR #1 it acknowledged the IRQ1 and IRQ2 flag and enables all interrupts again.
- \bigcirc At the beginning of ISR #1 the flags for IRQ1 and IRQ2 are acknowledged. All interrupts get disabled at the end of ISR #1.
- \bigcirc ISR #1 has not acknowledged the IRQ1 flag. ISR #2 acknowledged the flags for IRQ1 and IRQ2 at the beginning of ISR #2.

A punched paper ticket is used in a parking system. The punched paper ticket is using following format for each data line in Figure 4:

- 1 guidance bit (small holes)
- 8 data bits (large holes)

The punched paper tape gets pulled into the machine with constant speed of 50 ms for each data line. The data lines are scanned with an optical sensor, and the

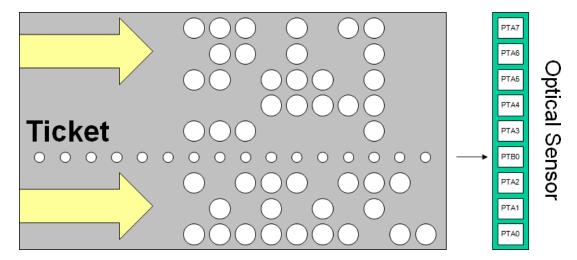


Figure 4: Parking Ticket

sensor digital output is attached to the port of a microcontroller. The state of the sensor/holes is available on the microcontroller PORTA, bit 0 to 7:

- Value of bit is 0: no hole, light does not go through
- Value of bit is 1: hole, light goes through

The state of the guidance hole is available on bit 0 of PORTB. The bit 0 of PORTB is configured to raise an keyboard interrupt on falling edge.

Given following program:

```
extern WaitMs(unsigned int ms); /* wait for the given ms */
unsigned char buffer [16]; /* contains the data read */
void Read(void) {
 uint8 t i;
  for(i=0; i < sizeof(buffer); i++) {
    WaitMs(50);
    buffer[i] = PORTA;
  }
}
void KBI Interrupt(void) {
  /* Guidance Hole Sensor */
  AcknowledgeKBI();
  DisableInterrupts();
 Read();
  {\bf Enable Interrupts}
void main(void) {
  for (;;);
```

parking ticket? Combination of interrupt and realtime synchronization. ✓ Interrupt synchronization. ○ Realtime synchronization. ○ Combination of gadfly synchronization and realtime synchronization. ○ Gadfly synchronization. (b) Which synchronization method is used for the synchronization on the first data hole? ○ Combination of interrupt and realtime synchronization. ○ Interrupt synchronization. ○ Realtime synchronization. ○ Combination of gadfly synchronization und realtime synchronization. ○ Gadfly synchronization. ○ Gadfly synchronization.			
V Interrupt synchronization. ○ Realtime synchronization. ○ Combination of gadfly synchronization and realtime synchronization. ○ Gadfly synchronization. ○ Badfly synchronization method is used for the synchronization on the first data hole? ○ Combination of interrupt and realtime synchronization. ○ Interrupt synchronization. ○ Combination of gadfly synchronization und realtime synchronization. ○ Gadfly synchronization. ○ Gadfly synchronization. ○ Gadfly synchronization. ○ Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. ▼ void ReadGadfly(void) { ■ uint8 t i; ■ for (i = 0; i < size of (buffer); i + +) { ■ while ((PORTB&1) == 0); ■ buffer i - PORTA; ■ while ((PORTB&1) == 1); ■ ** * * * * * * * * * * * * * * * * *	(a)	-	[1
Realtime synchronization. Combination of gadfly synchronization and realtime synchronization. Gadfly synchronization. (b) Which synchronization method is used for the synchronization on the first data hole? Combination of interrupt and realtime synchronization. Interrupt synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. (c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { Solution: void ReadGadfly(void) { Solution: void ReadGadfly(void) { Solution: Void ReadGadfly(void) { while ((PORTB&1)==0); buffer[i] - PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: Exertion 51		•	
Combination of gadfly synchronization and realtime synchronization. Gadfly synchronization. (b) Which synchronization method is used for the synchronization on the first data hole? Combination of interrupt and realtime synchronization. Interrupt synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. (c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { uint8_t i; for (i=0; i <sizeof(buffer); (portb&1)="=1);" *="" 51<="" buffer[i]="PORTA;" end="" estion="" for="" i++)="" readgadfly="" td="" total:="" while="" {="" }="" ₹=""><td></td><td>•</td><td></td></sizeof(buffer);>		•	
Gadfly synchronization. (b) Which synchronization method is used for the synchronization on the first data hole? Combination of interrupt and realtime synchronization. Interrupt synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. (c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { uint8_t i; for(i=0; i <sizeof(buffer); ((portb&1)="=1);" (exp<="0)" *="" 1;<="" 5="" buffer[i]="PORTA;" end="" exp)="" following="" for="" i++)="" if="" int="" louble="" power(double="" program:="" readgadfly="" return="" siven="" td="" total:="" while="" x,="" {="" }=""><td></td><td></td><td></td></sizeof(buffer);>			
(b) Which synchronization method is used for the synchronization on the first data hole? Combination of interrupt and realtime synchronization. Interrupt synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. (c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { uints_t i; for(i=0; i <sizeof(buffer); ((portb&1)="=1);" *="" 5="" 51<="" buffer[i]="PORTA;" end="" for="" i++)="" readgadfly="" sition="" td="" total:="" while="" {="" }=""><td></td><td></td><td></td></sizeof(buffer);>			
<pre></pre>	(b)	Which synchronization method is used for the synchronization on the first data	[1]
<pre></pre>		Combination of interrupt and realtime synchronization.	
Combination of gadfly synchronization und realtime synchronization. Gadfly synchronization. C) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { Solution: void ReadGadfly(void) { uint8_t i; for (i = 0; i < sizeof(buffer); i++) { while ((PORTB&1) == 0); buffer[i] = PORTA; while ((PORTB&1) == 1); } /* end for */ } /* end ReadGadfly */ Total: 5 tion 51		○ Interrupt synchronization.	
Gadfly synchronization. (c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) { Solution: void ReadGadfly(void) { uint8_t i; for (i = 0; i < sizeof(buffer); i++) { while ((PORTB&1) == 0); buffer i = PORTA; while ((PORTB&1) == 1); } /* end for */ } /* end ReadGadfly */ Total: 5 tion 51		$\sqrt{\text{Realtime synchronization}}$.	
(c) Implement a new function ReadGadfly() which does the same as Read(), but uses a gadfly synchronization method. void ReadGadfly(void) {			
uses a gadfly synchronization method. void ReadGadfly(void) { Solution: void ReadGadfly(void) { uint8_t i; for (i = 0; i < sizeof(buffer); i++) { while ((PORTB&1)==0); buffer [i] = PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: 5 tion 51		○ Gadfly synchronization.	
Solution: void ReadGadfly(void) { uint8_t i; for (i = 0; i < sizeof(buffer); i++) { while ((PORTB&1)==0); buffer[i] = PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: 5 tion 51	(c)	· · · · · · · · · · · · · · · · · · ·	[3]
<pre>void ReadGadfly(void) { uint8_t i; for (i = 0; i < sizeof(buffer); i++) { while ((PORTB&1)==0); buffer[i] = PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: 5 ition 51</pre>		void ReadGadfly(void) {	
<pre>uint8_t i; for (i = 0; i < size of (buffer); i++) { while ((PORTB&1) == 0); buffer [i] = PORTA; while ((PORTB&1) == 1); } /* end for */ } /* end ReadGadfly */ Total: 5 stion 51</pre>			
<pre>while ((PORTB&1)==0); buffer[i] = PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: 5 stion 51</pre>			
buffer[i] = PORTA; while ((PORTB&1)==1); } /* end for */ } /* end ReadGadfly */ Total: 5 Stion 51			
\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \		buffer[i] = PORTA;	
\[\text{\cond} \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \			
stion 51		} /* end ReadGadfly */	
stion 51			
Given following program: ouble power(double x, int exp) { if (exp<=0) return 1;		T	otal: 5
if (exp<=0) return 1;		• •	
return(x*power(x, exp-1)); }	doul	ble power(double x, int exp) { (exp<=0) return 1;	
	re }	eturn(x*power(x, exp-1));	

Evaluate following:

 \ominus ± In order to have this program reentrant, it is sufficient that **x** and **exp** are variables on a hardware stack.

Page 23 of 63

- \oplus \pm It depends on the compiler and the generated code, if this program is reentrant or not.
- \ominus ± The program is reentrant if it is called from an interrupt service routine only.
- \ominus ± The recursive implementation of this program ensures that it is reentrant.

Solution: Notice that the compiler routines for float operations (runtime routines) might not be reentrant. And not every hardware architecture is using a hardware stack for local variables, e.g. some PIC controlles or ST5 do not have a hardware stack (variables are on a software stack which prevent recursion or reentrant code).

Program	Main Priority	Sub Priority	Time
HP	0	0	5 ms
UP1	1	1	$2 \mu s$
UP2	1	2	$3 \ \mu s$
UP3	2	1	$5 \ \mu \mathrm{s}$
UP4	2	2	$2 \mu s$

Table 1: Interrupt System

The timing required for a context switch is given in table 2, which is illustrated in Figure 5.

Context Switch	Time
Total time for the interrupt, switch to a new program and starting execution	$1 \ \mu s$
of the waiting program	
Total time for the interrupt, switch to the interrupted program, immediate	$1 \ \mu s$
interruption of this program and switching and starting execution of the	
waiting program	

Table 2: Context Switch Timing

The interrupt system is using following rules (as used in the lecture):

$$if(MP(s) \le MP(fn)) \to ws = ws \cup s$$
 (1)

$$if(MP(s) > MP(fn)) \to INT(fn)$$
 (2)

$$if(MP(s) \le MP(fn)) \to ws = ws \cup s$$
 (3)

$$if(SP(s) > SP(fn)) \to ws = ws \cup s$$
 (4)

$$if(ws \neq \{\}) \rightarrow fn(MAX(SP(ws)))$$
 (5)

$$if(MP(ws) > MP(in)) \rightarrow fn = in \rightarrow ws = in$$
 (6)

The programs run according following information:

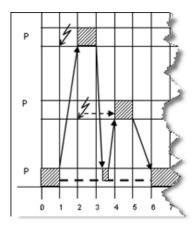


Figure 5: Example Context Switch

- 1. At the time 0 μ s HP starts.
- 2. At the time 2 μ s an interrupt for UP1 is raised.
- 3. At the time 4 μ s an interrupt for UP4 is raised
- 4. At the time 6 μ s an interrupt for UP2 is raised
- 5. At the time 9 μ s an interrupt for UP3 is raised
- 6. At the time 23 μ s an interrupt for UP2 is raised

Show the sequence of programs and interrupts in Figure 6. Use the same notation as in Figure 5 for interrupts (Exception, Pending), program switches, program (aktive, suspended).

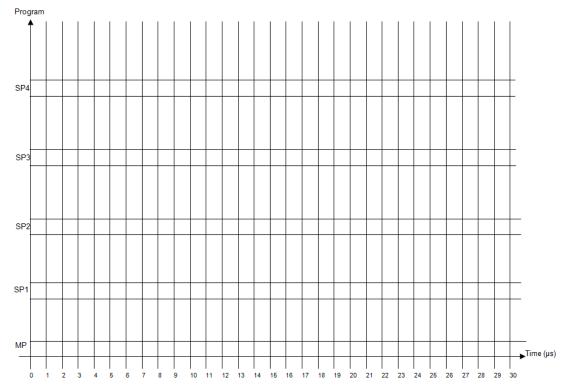


Figure 6: Program Timing

Page 25 of 63 Reached: ____

Solution: See Figure below.

- 1. A signal with a higher main priority interrupts always a running program.
- 2. A signal with same main priority interrupts never.
- 3. A signal with higher sub priority has to wait for an already started program with same main priority, even if it had been interrupted.
- 4. If there are multiple signals waiting with the same main priority, then the sub priority decides which one will be handled first. With sub priorities it is possible to influence the sequence of execution and you are not depending on the signal raise time.
- 5. If there is a pending signal with a higher main priority than the last interrupted program, then the control goes back to that interrupted program, but gets interrupted directly again.

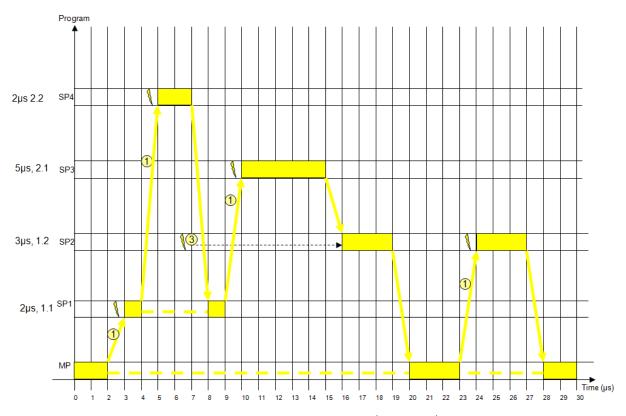


Figure 7: Program Timing (Solution)

Question 53......Points: [3]
Consider following implementation:

```
#define FUNC(a,b) i+a+b
int foo(int i, int j) {
   return FUNC(i,j);
}
```

Page 26 of 63 Reached: _____

	53	16	
Solution:			
return FUNC(i,j);			
return i+i+j;			
return 5+5+6;			
return 16;			
stion 54		Poi	nts: [3]
Given following source code:			
uint16_t abcd[8]; uint32_t buf[10]; static uint16 t values[2];			
static unitio_t values[2];			
(a) Determine the value of following expression:			
sizeof("abcd")			
	()	5	
4	(a)	<u> </u>	
(b) Determine the value of following expression:			
sizeof(buf)			
	(b)	40	
(c) Determine the value of following expression:	()		
sizeof(values)			
2			
	(c)	4	
stion 55		Poi	nts: [2]
Consider following program:			
void delay(void) {			
uint8_t i;			
${f for}\ (\ {f i=0}; {f i<50}; {f i++});$			
)			

- \oplus ± can be optimized by a smart compiler to a function which only contains a return; statement
- \oplus ± will wait for a certain time which is depending on the speed of the microcontroller used
- $\ominus \pm$ will never terminate

```
typedef enum {A, B, C, D, E, } States;
const char tbl[3][2][2] =
{ {{A,0}, {B,1}},
    {{C,3}, {A,4}},
    {{C,0}, {B,5}}
};

void Run(void) {
    char j, i = 0

    for (;;) {
        j = Input();
        Output(tbl[i][j][1]);
        i = tbl[i][j][0];
    }
}
```

- (a) Given following sequence of Input() values: 0, 1, 0, 1, 1, 1. Determine the sequence of Output() values: [1]
 - (a) 0, 1, 3, 5, 4, 1
- (b) Draw the corresponding state diagram:

Solution: See diagram.

Total: 2

[1]

Consider following program:

```
        void main (void) {

        char buf [0x100];
```

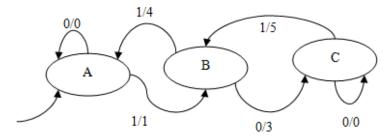


Figure 8: Solution Mealy Diagram Drawing

```
int i, j;

PORTB = 0;
for (i = 0; i < size of (buf); i++) {
    CFG = 0x80; PORTB = 4;
    while (CFG!=0);
    buf [i] = PORTA;
    PORTB = 0;
}</pre>
```

The following applies:

- $\ominus \pm$ It implements an interrupt synchronization.
- $\oplus \pm$ It implements a Gadfly synchronization.
- $\ominus \pm$ It implements a Realtime synchronization.
- $\ominus \pm$ It implements no synchronization.

Evaluate following statements about reentrancy:

- \oplus \pm A function which modifies its own code is not reentrant.
- $\ominus \pm A$ function which calls an interrupt service routine is not reentrant.
- \ominus ± Recursive functions are always reentrant.
- \ominus ± Interrupt service routines are always reentrant if they do not call another routine.

If discussing interactive, reactive and transforming systems, then

- \oplus ± relative short answer times are typical for interactive systems.
- \oplus ± reactive systems are common in systems which do measurement and control.
- \oplus ± transforming systems are typically optimized for high throughput.
- \oplus ± an example for an transforming system could be a network router.

In the context of real time following applies:

 \ominus ± Realtime means to produce a result as fast as possible. \ominus ± A computer is realtime, if is is able to produce at average system load the correct result as fast as possible. \ominus ± For realtime it is sufficient to have an accurate timing system. $\ominus \pm$ An RTOS is required for a realtime system. Question 61......Points: [3] For all reentrant functions implemented in C the following applies: $\ominus \pm$ A reentrant function shall not be interrupted. \ominus ± Interrupt functions does not have to be reentrant, but all functions called from that interrupt routine. $\oplus \pm$ A function which modify itself is reentrant, as long the self modification happens with disable interrupts. \oplus ± On the ARM Cortex-M0+/M4F the usage of local stack variables does not violate reentrancy. Question 62......Points: [2] Explain in a few words the reasons why a switch (like a button) needs a resistor. Illustrate it with a small drawing. **Solution:** Without a resistor, the signal is undetermined or open. It is needed to define the signal to a defined level (low or high) if the switch is not closed. Question 63......Points: [2] List important points to be considered for the implementation of an ISR: Solution: As short and as fast as possible. Need to acknowledge the interrupt at the beginning (or additionally at the end. Need to care about reentrancy, and that shared functions are reentrant and shared variables are protected. Explain two different ways how a microcontroller can implement interrupts: Solution: Using a vector table: the interrupt source gets translated into a vector number which is used as index into a vector table. The other way is that the processor directly jumps to an address and executes the code there (stub based approach). What happens if two developers work at the same project but in different files or at different parts in one file and commit on Git?

Page 30 of 63 Reached: ____

Solution: Git merge automatically the changes together.
Question 66
Solution: SVN: Centralized VCS, Git: Distributed VCS.
Question 67
Solution: There will be a conflict which cannot be resolved by the VCS. The developers have to solve this problem manually (solving the conflict).
Question 68
Solution: Every file is load in the repository except those are mentioned in the .gitignore.
Question 69
Solution: TS: Video encoder, RS: PID Controller, IS: ticket selling machine
Question 70
#define CALC1 (2+5) #define CALC2 (5*3)
Solution: To ensure proper usage in other macros or in calculations. $10*CALC1$ is not the same as $(10*2)+5$
Question 71
Solution: Advantages: Faster code, smaller code. Disadvantages: Interface, Encapsulation, Debugging

	Solution:
	<pre>#ifndefLED_H_ #defineLED_H_ /* content of header file */ #endif</pre>
	estion 73Points: [2] How do you declare global variables? In what kind of file?
	Solution: In a header file with *.h extension. Using extern for variable declaration, e.g. extern int LED_global; I use a good name with prefix because the name is visible in the whole project.
•	estion 74Points: [1] Is it possible to include other files than *.h with #include? Can you give an example?
	Solution: Yes, any text files can be included (as long as understood by the compiler). For example I can have an array of hex values in an array and then include it as bitmap.txt.
Que	estion 75
	Solution: It blocks further execution.
	estion 76
	Solution: Each ISR needs to ensure that the other ISR is not executed. This can be with priorities, or with disabling the interrupt for the other ISR.
	estion 77
	Solution: Each ISR needs to ensure that the other ISR is not executed. This can be with priorities, or with disabling the interrupt for the other ISR.
Que	estion 78Points: [1] What have you to do with unneeded interrupts?

handler, order of events has impact about priority, need mutual exclusion for shared data.

Solution: Advantage: simple. Disadvantages: long if-elsif-else/switch in the event

Page 33 of 63 Reached: _____

Question 82
Solution: ARM Cortex-M0+
Question 83
Solution: Main Stack Pointer (MSP) and Process Stack Pointer (PSP).
Question 84
Solution: Absolute (e.g. 13:50) or relative (e.g. after 50 ms).
Question 85
Solution: 8 MHz crystal oscillator.
Question 86
Solution: Gadfly, interrupt, realtime.
Question 87Points: [1] How can the logic level of a pin be enforced?
Solution: Internal (port configuration, pull resistor), external circuit (pull resistor).
Question 88
Solution: To have defined voltage levels.
Question 89
Solution: functional, hierarchical, else-if state machine.
Question 90

Colution if alifala and tall
Solution: if-elsif-else, switch, table
Question 91
Solution: determined by current state and current output.
Question 92
Solution: Handle multiple 'interrupt like things' with just one timer.
Question 93
Solution: yes.
Question 94
Solution: To avoid multiple interrupts, and to have reached a stable state.
Question 95
Solution: hardware (R-C) or software (timer, delay).
Question 96
Solution: USB is a much more complex protocol, and therefore, has a bigger overhead. If the USB performance is not needed, a simpler UART connection is more efficient. At the Freedom board, there is an extra MCU to convert the USB to a UART protocol. This means that there is an extra MCU to take over the USB overhead.
Question 97

Page 35 of 63

Reached: ____

Solution: Serial means that it sends the data bits after bits on a single data line (as a sample Rx or Tx line). Asynchronous means that there is no clock supported to read the data. The start of the data has to be detected by the protocol (start

bits, stop bits). Question 98.......Points: [1] What's the command parser table? **Solution:** It's a list with function pointers. In this list the parser selects the method which is going to be executed. What's the difference between Memory Scheme 1 and 2 in FreeRTOS? Solution: Memory Scheme 1 only allocates Memory. It's not possible to delete Tasks. With Scheme 2 it's possible to free space and reuse. There would be another Scheme where you can merge freed blocks. Provide a good example how FreeRTOS queues can be used between multiple tasks: Solution: To send messages between tasks. What is the difference between xQueueReceive() and xQueuePeek() in FreeRTOS? Solution: xQueuePeek() only checks if there is an item in the queue. xQueue-Receive() does remove the item. What are the H-bridges of the motors needed for? **Solution:** They are used to set the direction of the engine by simply twisting the two wires of the motor. Why can the PWM channels not run with different frequencies?

Page 36 of 63 Reached: ____

a separate value register which all compare with the same common timer.

Solution: There is only one Timer for all the different PWM-channels and therefore the frequency can only set for this common timer. The different channels only have

Question 104
Solution: The coast modus of the engines is only possible if "Mode 0" is used on the motor drivers. On the SUMO robot, this pin is not routet to the FRDM board and just hardware wise set to HIGH (Mode 1).
Question 105
Solution: More write cycles, durable, defined read/write times.
Question 106
Solution: Add a checksum for the blocks.
Question 107
Solution: The Binary Code several bit changes between the counting steps are possible whereas the Grey Code changes only one bit between each step. It's possible to build the Grey Code recursive and it is permutable.
Question 108
$\textbf{Solution:} \ \ \text{Interrupts, Sampling, input capture, dedicated quadrature peripheral/IC}$
Question 109
Solution: The sampling method causes a constant system load whereas the interrupt method causes a speed dependent system load. It's easier to handle a constant system load to test a system.
Question 110

Page 37 of 63 Reached: _____

Solution: If you use the interrupt method the system freezes (to many interrupts). If you use the sampling method the system still works normally, but can't recognize every encoder step (errors).

Solution: The input signals are the two hardware signals and the output signals are the number of errors and the current position/steps.

Page 38 of 63 Reached: _____

Question 112
Solution: It is used as a sentinel, so the code can iterate through the table until there is a NULL entry.
Question 113
Solution: sizeof() returns the size of the string including zero byte, but we want to compare the string without it.
Question 114
Solution: baudrate, number of data bits, stopbit, parity
Question 115
Solution: volatile
Question 116
Solution: Because they operate with different speeds.
Question 117
Solution: Because an interrupt can happen any time, it must be ensured that there is no data corruption.
Question 118
Solution: No, for example the HCS08 does not have nested interrupts.
Question 119

Solution: Reset, NMI, HardFault
Question 120
Solution: Yes, e.g. by a Reset or NMI as they have numerically lower interrupt priorities.
Question 121
Solution: Measure the delta-time (duration of period) if frequency is low. Use counting the steps or periods if time intervals are short.
Question 122
Solution: The speed can be estimated more frequently than the measurement interval. Additionally it allows an averaging of the speed. With the ring buffer a configurable and dynamic time span for the estimate can be used.
Question 123
Solution: Is the system control value (e.g. PWM for the motor) limited in range, then the integral part can go out of this range. Then the integral sum will increase without having an additional impact on the control value, and delays the impact of the integral on the control value. Additionally the limit avoids a numerical overflow of the integral sum.
Question 124
Solution: Nyquist/Shannon: need to sample it with at least 50 μ s.
Question 125

Solution: There are several options: one way is to create a drive task and call the PID from there. Or to use a periodic interrupt to call the PID.

Why is it not possible to directly measure the output signal of the optical quadrature encoder we have used?

Solution: The signal is an analog sinus-like signal, without a 50%-50% high-low signal distribution. To get a clean quadrature signal a DAC with comparator devices are used.

Discuss the pros and cons of using either sampling or interrupt method for a quadrature signal:

Solution: Usually sampling is usually preferred as it will create a constant system load. However, this creates a high load of the system even if the wheel is not moving or only slowly moving. Using interrupts can cause problems if there are too many interrupts. It would allow a low system load if the wheel is not moving or slowly moving.

What fundamental problem exists for absolute position encoders, and how can it be solved:

Solution: Because of mechanical tolerances, multiple bits can change from one step to another. The solution is to use a Gray code/encoder, as with this only one bit changes from one sector to another. In addition the Gray code is cyclic and is therefore ideal for wheel position measurement

Can you list the main features of the MCP4728?

Solution: 12bit DA-Converter, includes an EEPROM to store DAC values and settings, I²C bus and protocol, 4 DAC output signals.

Question 130.......Points: [2]

What are the special things or attributes of the Gray code?

Solution: Hamming distance of 1 (only one bit changes), the code has permutation (every code only occurs once), it is cyclic (last and first code confirm to the rules) and recursive (codes of lower order are embedded in code of higher order), and it is simple to transform a binary code into a Gray code.

Page 41 of 63 Reached: _____

Question 131
with 100 Hz, what would be an estimated way distance over the white line until you detect the white sumo line in the application?
Solution: If robo is moving with 1 m/s and we measure with 100 Hz, then the robot will move in average 2 cm until the motors get stopped.
Question 132
Solution: To limit the current through the photo diodes.
Question 133
Solution: This generates a digital signal we can measure, and we do not need an A/D converter pin.
Question 134
Solution: Predictability, precise timing, speed.
Question 135
Solution: Running multiple things in parallel, using RTOS services (queues, semaphores, mutex,), scalability of application.
Question 136
Solution: Preemptive: the RTOS distributes the processing time, tasks get suspended by the RTOS. Non-preemptive: the task are cooperative, they pas the control back to the kernel.
Question 137

Page 42 of 63

Reached: _____

Solution: To maximize the CPU utilization among different tasks.
Question 138
Solution: The RTOS needs to adhere to strict timing and needs to produce output with given timing constraints.
Question 139
Solution: IDLE, PRESSED, RELEASE.
Question 140
Solution: Software (state machine, time delay/low pass filter) or hardware (capacitor, Schmitt-Trigger.
Question 141
Solution: Bouncing is a mechanical problem. To process the state of a bouncing push button, it needs first to stabilize.
Question 142
Solution: Add additional events to message an inter-click. Extend the state machine so it either continues or uses different states in the state machine.
Question 143
Solution: The fast decay principle is to revert the current from the previous movement. E.g. if the motor is turning forward, to put current into the H-Bridge to in reverse order (turning it backward). This will bring the energy stored in the inductor down fast. So the 'fast' is about how fast the current reaches zero. With 'slow decay' either the lower half or upper half of the transistors are on, allowing the inductive current to flow back. Because this takes longer than with the 'fast' method, this is called 'slow decay'. With 'fast decay', the motor coasts down the speed, while with 'slow decay' the H-Bridge using an active break. Hint: .

Question 144
Solution: No, this depends if the signal is LOW or HIGH active.
Question 145
Solution: Increase the PWM frequency.
Question 146
Solution: A half-H Bridge has two transistors, while a full H-Bridge has 4 transistors. With a full H-Bridge the direction of a DC motor can be changed.
Question 147
Solution: Files which can be used for multiple projects, like files of a library.
Question 148
Solution: A lot of duplicated files, maintenance problem could occur as I need to change files in different places.
Question 149
Solution: easier to read, re-usability, easier to understand.
Question 150
Solution: Either with '//' or e.g. with '{PARENT-1 PROJECT_LOC}
Question 151

Page 44 of 63

Reached: ____

Solution: Because with sharing the project the paths might be different on another machine.
Question 152
Solution: To debug it a preprocessor listing needs to be generated (-E for gcc).
Question 153
Solution: With double quotes 'user' header files are included, and with <> library header files. The compiler uses two different search settings for user and library header files.
Question 154
Solution: The processor is using muxing: a single pin can have different purposes, like as touch sensing pin, as a input or output port pin or as a timer channel pin.
Question 155
Solution: That it is only possible to use the pin in one mode at a given time. And that muxing a pin might have impacts and side effects on other pins, e.g. certain functions are not possible any more.
Question 156
Solution: printf() is not a save function and can cause a stack overflow, and is subject of security issues. Additionally it needs a lot of code space and $stack/RAM$.
Question 157
Solution: Because the string is terminated with zero byte.
Question 158

static int counter = 0;Now you remove the static. What is the effect? **Solution:** The variable now gets initialized with zero at every interrupt. You consider to handle the event bits set from the main loop, instead of using a check/clear in several places in your application. Discuss the pros and cons of this approach: Solution: Pros: events are all handled in a single place and centralized. Complexity is simple. Cons: can be a long switch/if-else-if, event handling depends on frequency of main loop. Cannot handle events independently/concurrently. Can you give reasons why the KL25Z128 bus clock is limited to 24 MHz, while the K22FX512 can run a bus clock of 60 MHz? Solution: Maximum core clock of KL25Z is 48 MHz, while the K22 can run up to 120 MHz. The K22 core is a Cortex-M4F and it can run at a higher speed than the Cortex-M0+. Additionally, the bus clock can run at maximum half of the speed of the core clock. What are the pros and cons of using an external clock vs. internal clock? Solution: Pros: more accurate, higher overall speed possible. Cons: more costs, more PCB space needed. Briefly explain the purpose of CPU clock, Bus Clock and System Clock: **Solution:** The CPU is clocked by the CPU clock, the bus clock is used for data transfer and to access memory, and the system clock is used to clock the peripherals.

Page 46 of 63 Reached: ____

using different platforms?

Solution: Different clock rates, different instructions and register size, different compiler and different optimizations, different CPU speed and performance, different

interrupts affecting the system; they all can make the delay loop different and make the timing fail. Can you provide reasons why using an interrupt synchronization might not be the best option? **Solution:** There is overhead associated with switching the context (saving registers, context switch), this all can cost too much time. Can you give examples for where synchronization is necessary? Solution: I/O, HMI, whenever multiple processes need access to a shared resource What is the HardFault component doing? Solution: It copies the pushed registers from the stack to local variables so they can be easier inspected. What do you have to consider in the Events module when you clear/set event bits? Solution: It has to be done in an reentrant way, otherwise race conditions and wrong data could occur. How can you quickly check the interrupt level of your interrupts of your application? Solution: Check the vectors.c file where it lists all vectors in the vector table with their interrupt levels. Explain the principle of the Event handling in the Event module: Solution: Events can be set from everywhere in the application. In the 'main' loop an event handler gets called which checks for any pending events and then calls a callback for events found. What are Thumb and Thumb-2 instructions, and what is special about them?

Page 47 of 63

Solution: Thumb are reduced/condensed 16bit instructions on ARM Cortex devices for better code density. In Thumb-2 has both 16bit and 32bit instructions.

They are backwards compatible, so it allows to run M0+ code on an M4 core. List advantages of using an I/O structure: **Solution:** It adds flexibility: the different input and output streams can be remapped, re-assigned, chained/piped or be used to distribute the date to different channels. The separation between stdout and stderr allows to separate normal output from error output. List reasons why you should use the Utility component methods: Solution: It provides commonly used functions (copy, compare, etc), plus all string operations are implemented that no buffer overflow can happen. Question 173......Points: [3] Which possibilities do you have on the FRDM board to communicate to a COM port on the computer? Solution: Using the (debug) K20 with USB CDC. The KL25Z uses UART to communicate with the K20 which then is using USB CDC to communicate with the host. Another way would be to implement the USB CDC stack on the KL25Z. For what are the callbacks used in the Trigger data structure? Solution: The callback is used to call the user functionality when the trigger has expired. What is the purpose of TRG_NOF_TRIGGERS and why does it have to be at the end of the list? Solution: It counts the number of elements in the list so it can be used to allocate the correct size of the array. Question 176.......Points: [1] How many FreeRTOS tasks can be in the RUNNING state?

Solution: As we have only one core: 1
Question 177
Solution: For timing reasons, the embedded OS application typically can directly access the hardware. On a standard OS the hardware can be protected by the OS with an additional layer. Typically the drivers on an embedded OS system are part of the application (not part of the OS as with a standard OS).
Question 178
Solution: In principle, every mechanical switch needs to be debounced. And the debouncing timing shall be as such to cover a broad range of different bouncing time, as even within a series buttons will be different.
Question 179
Solution: Either by hardware (filter, Schmitt Trigger) or with software (state machine, low pass filter). A basic idea is to filter out the bounces (wait for some time).
Question 180
Solution: Cooperative and preemptive. In preemptive mode the scheduler runs the task with the highest prio, and usually time slices between tasks of same priority. In cooperative mode the application has to pass back control to the scheduler, or it has to be blocked.
Question 181
Solution: vTaskDelay() waits for the given time at the calling time. vTaskDelay-Until() delays from the time which is passed to the function, thus allowing a constant task frequency.
Question 182

Page 49 of 63

Solution: If higher prio tasks are available, then this high priority task might be scheduled. Otherwise the scheduler tries to make the timing happen as much as possible, and so it will schedule the calling task again.

Solution: Heap size is in bytes, stack size is in stack units (usually 32bit on a 32bit microcontroller).

Question 184......Points: [2] What is the purpose of 'max syscall interrupt priority'?

Solution: This defines the interrupt priorities masked out (blocked) by the Kernel.

Solution: Alloc only, no merge, malloc/free wrapper, merge blocks, multiple blocks

Question 186......Points: [1] What does IPC stand for?

Solution: InterProcess Communication

Solution: Needs more space for data structure/information, requires more computation time

Solution:

```
static xSemaphoreHandle sem = NULL;
static void vSlaveTask(void* pvParameters) {
  for (;;) {
    if (sem != NULL) {
      if (xSemaphoreTake(sem, portMAX_DELAY) == pdTRUE) {
        LED1_Neg();
    }
```

Page 50 of 63 Reached: ____

```
}
}

static void vMasterTask(void* pvParameters) {
  for (;;) {
    if (sem != NULL) {
       (void)xSemaphoreGive(sem);
       FRTOS1_vTaskDelay(300/portTICK_RATE_MS);
    }
}
```

Solution: Interrupts needs to be masked.

Solution: No, not always needed. In FreeRTOS semaphores can be used like tokens and don't have to be returned. Mutexes have to be returned.

Solution: With volatile memory (VM), the data is lost when power is removed (typically RAM or SRAM). NVM stands for Non-Volatile-Memory and here the data remains even if power is removed. Usually EEPROM or FLASH is used for NVM.

Question 192......Points: [2] Why is it not possible e.g. to write only 1023 bytes in FLASH?

Solution: Typically FLASH memory is block oriented (e.g. 1 kByte). And that whole block needs to be erased first. So in order to write 1023 bytes without loosing the 1024th byte, a backup has to be done.

Page 51 of 63 Reached: ____

	Solution: All bits are 1: 0xffff	
	estion 194	[2]
	Solution: FLASH memory has limited erase/programming cycles, e.g. 10k or 100)k.
	estion 195	[2]
	$\textbf{Solution:} \ \ \text{They have a data sheet and guaranteed access and erase/program tim}$	es.
-	estion 196	[2]
	Solution: In the Gray code only one bit changes from one code to another.	
-	estion 197	[2]
	Solution: $Value_Gray = (Value_Bin >> 1)EXORValue_Bin$	
-	estion 198	[2]
	Solution: We needed to tune the DAC in order to have a 50% - 50% signal, overlapping properly in the middle.	er-
-	estion 199	
	Solution: $f_s ampling > 4 * N_H oles * 2 * Revolution = 4 * 100 * 2 * 6000/60 = 80kH$	Hz
	estion 200	[2]
	Solution: Using a table: using the previous value as first index, the current value as the second index. The value in the array gives the step change.	ue
	estion 201Points: What is the problem with the quantization and the change of bits?	[2]

Solution: There are always quantization errors with converting an analog value into a bit value. The bit values making jumps between the steps, and are not continuous.
Question 202
Solution: Δpos approach is better for higher speed. $\Delta time$ approach is better for slow speeds.
Question 203
Solution: Have flexibility to estimate over different time periods and to average the position changes. Disadvantage is that it requires more memory.
Question 204
Solution: The frequency should be high enough (e.g. 50 kHz) to produce a 'constant' voltage, but it has to be lower than the maximum input frequency of the IC.
Question 205
Solution: Turn off the motors (speed and duty cycle set to zero), release any memory allocated (if any), turn off PWM signals.
Question 206
Solution: The driver caps it to 100%.
Question 207
Solution: If negative, the motor direction is set to 'backward' and it uses the absolute number of the duty for the speed.
Question 208

Solution: Use a capacitor to buffer the voltage (to avoid the voltage drop). Either

on the battery connectors or on the 5V supply lines. In C, how can you make sure a variable is always created when you enter a function? Solution: Using normal local variables in the function. They get created on the stack when the program executes the function. How can you increase your confidence that your implementation is reentrant? Solution: Carefully review your code and make sure it is really reentrant. Other than that, test your system with a high interrupt load as this can increase the chances of race conditions. Question 211......Points: [1] Discuss 'self modifying code' for embedded systems: Solution: Usually code gets executed in FLASH which cannot be modified (except the FLASH gets reprogrammed). So typically code in Embedded Systems is not subject of self modification. But if code (like for the FLASH programming) is running in RAM, it easily can do self-modification. Which problems can happen if for a critical section the interrupts get disabled? **Solution:** It increases the interrupt latency time, or interrupts might be missed. Question 213.......Points: [2] What is the advantage of using sub-priorities in an interrupt system? Solution: It allows greater flexibility, you can group interrupts and ensure that interrupts in a group are not interrupted by the group interrupts. Which points you have to consider and check if using a LED with a microcontroller pin? Solution: Check the schematics on which pin the LED is connected, check if the anode or the cathode is connected to the pin, verify that the pin is able to drive or

Page 54 of 63 Reached: _____

sink the current required for the LED.

Question 215
Solution: 0 Volt.
Question 216
Solution: Init() to initialize the driver and allocate memory structures, Deinit() for deinitializing it, Open() to reserve it for usage and to get a handle, and Close() for releasing the reservation.
Question 217Points: [1] Why should a driver deinitialized with Deinit() if it is not used any more?
Solution: Deinit() would free up any allocated resources (e.g. memory or pins) for it, so it can be used otherwise.
Question 218
Solution: Use braces for macro definitions.
Question 219
Solution: Use #ifndef \rightarrow #define \rightarrow #endif
Question 220
Solution: Using a function without a prototype.
Question 221Points: [1] What does the keyword volatile?
Solution: Volatile disables compiler optimization.
Question 222

Solu	ion: Active waiting needs processing power, blocks further execution.
-	223
Solu	ion: Realtime synchronization, Polling, Interrupt synchronization
	224
Solu	ion: Realtime Operating System
•	225Points: [1 ne advantage and one disadvantage using an RTOS:
	ion: Advantage: quasi-concurrent execution, Disadvantage: uses resources nemory
	226Points: [1 is Preemptive Scheduling?
	ion: Scheduler always runs the highest priority ready task, multiple tasks the same priority are time-sliced.
-	227Points: [1 e one method a task can use to request a context switch:
Solu	ion: taskYield()
	228
Solu	ion: Use an endless loop for the task code, e.g. for(;;) { /* code */ }
	229Points: [1 ne four conditions for a realtime program:
	ion: The correct result at the correct time, independent of the current system in a deterministic and foreseeable way.
	230Points: [1 is the main difference between hard and soft realtime?

Solution: Failing the timing condition for a hard realtime system means failure. For a soft realtime system it means degradation.
Question 231
Solution: Methods, Events, Properties and inheritance.
Question 232
Solution: With absolute things have to happen at a given absolute time. With relative it means things shall happen after a certain time (relative).
Question 233
Solution: The ARM Cortex-M0+ has no hardware integer division instruction.
Question 234
Solution: No, this will only mask maskable interrupts, and reset is not maskable.
Question 235
Solution: Ringbuffer or Queue: enqueue(), dequeue(), push(), pull()
Question 236
Solution: non-reentrant shared subroutines, incorrect priorities of interrupts, buffer or stack overflows, dangling pointers, not initialized variables.
Question 237

Page 57 of 63

Reached: _____

Solution: Keep the ISR as small and fast as possible. Only handle things in ISRs which cannot wait. Question 238.......Points: [1] List three different design patterns to implement a state machine: Solution: IF-ELSEIF-ELSE, SWITCH-CASE, Table-driven Question 239......Points: [1] On the INTRO V2 robot you are getting sporadic interrupts for the push button, but the button has not been pressed. What is the most likely reason for this? **Solution:** Internal pull-ups are not enabled for the push button pin. You are using polling in the INTRO project to check the status of the push button on the robot. This works fine at the beginning of the semester, but later in the project not any more and push button presses are not always detected. What could be the problem? **Solution:** The polling gets delayed due other higher priority task, so the polling is not executed frequently enough. How can you find out on a linux system which port/device you can use for communication with the robot? **Solution:** Use 'dmesg' or 'dmesg | grep tty' and look for a device matching the USB CDC of the robot. Why should you not use printf() e.g. for printing some strings? **Solution:** printf() provides functionality for complex string printing and is therefore huge and needs a lot of stack space. For printing normal text smaller implementations like the ones in the utility or shell modules should be used. Question 243.......Points: [1] Which three types of information are stored in a Trigger TriggerDesc structure? Solution: When something shall happen (time, triggerTick counter), what shall happen (callback function pointer) with what kind of data (data pointer).

Page 58 of 63 Reached: ____

Question 244
rupt pin.
Question 245
Solution: The plugin cannot read memory from the target while it is running. You need to pause/suspend the application in the debugger first.
Question 246
Solution: In post mortem mode the data is not streamed while the target is running. In post mortem mode you have to read the recorded data from the viewer using a menu item.
Question 247
Solution: 4, 3, 2, 1, 0
Question 248
Solution: configMAX_SYSCALL_INTERRUPT_PRIORITY has no effect on ARM Cortex-M0+, so all interrupts are disabled.
Question 249

Page 59 of 63

Reached: ____

Solution: Only the ones with a _FromISR suffix in the method name.
Question 250Points: [1] Using FreeRTOS queues, can you change the size of a queue?
Solution: No, the size of the queue is fixed at the queue creation time.
Question 251Points: [1] What are FreeRTOS hooks? List three example hooks:
Solution: Hooks are optional callbacks which can be used by the application. Hook for stack overflow, for malloc/heap failed, tick hook.
Question 252
Solution: CLS1_ParseWithCommandTable() directly parses the given command string, while CLS1_ReadAndParseWithCommandTable() is reading and appending from standard input to a buffer and only parses the command if a line end is read into the buffer.
Question 253
Solution: strcmp() compares two strings regardless their size, while strncmp() only compares up to a given number of characters.
Question 254
Solution: The clock gets divided by two in the USB peripheral (48 MHz to 24 MHz. The USB signal itself has a frequency of 12 Mhz and it needs to be sampled by the peripheral with the a double frequency of 24 MHz).
Question 255
Solution: 3 V.
Question 256

Solution: Wrong clock to the USB module, USB module not enabled or USB interrupts not turned on, application does not call the USB application hook/task function.
Question 257
Solution: No, unless you create a new queue.
Question 258
Solution: By value
Question 259
Solution: A lower priority task is blocking a higher priority task because of a resource held by a lower priority task.
Question 260
Solution: A semaphore can have a counter, while a mutex is a special type of a binary semaphore to implement mutual exclusive access to a resource. In FreeRTOS a semaphore token has not to be returned/released, while a mutex always has to be released.
Question 261
Solution: Semphore are a tool to implement a critical section and is typically a feature of an operating system. A critical section is a section of code which is protected that not multiple program flows can be active in that section.
Question 262
Solution: No supply voltage for the sensor (no batteries inserted, not turned on, batteries empty).
Question 263

Solution: Increase frequency of the timer counting the time.
Question 264
Solution: The Mutex implements the Priority Inheritance protocol, while the Semaphore does not.
Question 265
Solution: Because on some microcontrollers the FLASH memory is disconnected from the internal bus while the flash is programmed.
Question 266
Solution: The currently programmed flash block could be completely erased, or only part of it would be programmed.
Question 267
Solution: Check the data sheet for minimum and maximum allowed values. Set the value to somewhat above 20 kHz to avoid audible sound. Set it to something possible by your microcontroller PWM hardware. Don't set it to high as a high frequency means higher transistor switching and more power loss.
Question 268
Solution: The struct has a field (currSpeedPercent) which is signed. Negative values mean backward and positive forward.
Question 269

Solution: This setup will only get half of a quadrature signal, so only speed, but no direction information.

Question 270.......Points: [1]

With your robot you are driving with a constant speed. The problem is that you get about the same number of 0, 1, -1 and ERR steps from your quadrature decoder. What could be the problem?

Solution: The sampling frequency is far too low.