### Part I

# Evaluation Page Part B

Name: Signature:
------------------

## 1 Evaluation

This part of the exam has 152 questions, with a total of 416 points and 4 bonus points.

Part	Max. Points	Scored Points
A	60	
В	180	
	Total	

#### Evaluation Part B 2

Page	Points	Bonus Points	Score
5	7	1	
6	12	0	
7	6	0	
8	11	0	
9	9	0	
10	7	0	
11	7	0	
12	10	0	
13	9	0	
14	12	0	
15	10	0	
16	11	0	
17	2	1	
18	10	0	
19	6	0	
20	10	0	
21	4	0	
22	8	0	
23	5	0	
24	16	0	
25	4	0	
26	5	0	
Total:	181	2	

		Г	
Page	Points	Bonus Points	Score
27	5	2	
28	4	0	
29	6	0	
30	4	0	
31	7	0	
32	15	0	
33	13	0	
34	10	0	
35	7	0	
36	9	0	
37	9	0	
38	11	0	
39	8	0	
40	12	0	
41	8	0	
42	9	0	
43	9	0	
44	6	0	
45	3	0	
46	12	0	
47	3	0	
48	9	0	
49	13	0	
50	3	0	
51	37	0	
52	3	0	
Total:	235	2	

### Part II

## Rules

Answer the questions within the space provided. If you do not have enough space, you can use the backside of the sheet. In that case clearly indicate that your answer continues on the backside.

#### 3 Supporting Materials

This is an examination in writing, without the usage of any electronic devices, except a scientific pocket calculator. No restriction on the model of calculator that may be used, but no device with communication capability shall be accepted as a calculator. All other electronic devices are prohibited. Writing paper is available, writing instruments (pencil, pens, etc) have to be organized by the student.

- Part A: Without any supporting material, with calculator.
- Part B: With a self written summary (format A4, 8 sheets or up to 16 pages), with calculator

#### 4 Procedure

1. Duration:

Part A: 1 hour = 60 minutes = 60 points.

(short break)

Part B: 3 hours = 180 minutes = 180 points.

- 2. Sign the first page in the provided space. With this you certify that you are only using permitted support material and you are complying to the rules.
- 3. Write your name on any detached or additional paper sheets. Sheets without a name will not be evaluated.
- 4. Use the provided paper for your solutions. Use the provided space in the forms and tables. If needed use scratch paper. Document your way to your solution as appropriate.
- 5. Each question has a defined number of maximum points associated.
- 6. If a question is unclear, make reasonable assumptions. Document your assumptions and provide a rationale.
- 7. Write clearly and legibly. Unclear or multiple solutions will not be evaluated.
- 8. There is a short break between part A and B. You have to sign into a list for a needed break during the examination parts. Only one person can leave the room for a short time.
- 9. If something is unclear, ask your supervisor in the room.

#### 5 Time Management

Read first all questions. Make sure you distribute your available time to all the questions. To reduce disturbance, ask questions in the first 15 minutes of the exam period.

### 6 Multiple-Choice Questions

- 1. Try to answer all questions if possible. If you are not sure, choose the answer which seems the best one.
- 2. For the questions of type  $\bigcirc$ : Choose **exactly one** option with  $\otimes$  (or  $\sqrt{\ }$ ), which you think is the best match. With a correct answer you get the given number of points for that question.
- 3. For the questions of type ±: After a question or possibly incomplete sentence there are four answers or extensions. Evaluate each of them if they are true or false and mark them accordingly with '+' (true) or '-' (false). Independent if the question is formulated grammatically in singular or plural, it is possible that 0, 1, 2, 3, 4 of the choices are true. For three correct answers out of four you receive half of the points.
- 4. Wrong answers will have no penalty. Each question which has no answer is treated like a wrong answers and will be evaluated with zero points.
- 5. If you are changing your mind: cross out your old answer and clearly mark which answer is the new one.

# May Dilbert be with you! ©

	ould be your own bonus question ☺.
	1
In versi you sho	2
	drivers we implemented Deinit() functions, for example for the shell:  HELL_Deinit(void);
	4
  Question (a) Tl	
uestion  (a) Tl	4
	4
uestion  (a) Tl	4
uestion  (a) Tl	4
  Question (a) Tl	4
	4

[2]

	at all:
<b>+:</b> -	Doints, [10
	n 5Points: [10 n the following interfaces:
bool <b>void</b>	KEY_Get(void); /* return FALSE if key is pressed, TRUE otherwise * WAIT_Waitms(uint16_t ms); /* realtime waiting for the given nilliseconds */
	def enum {
	_KEY, /* no key pressed */ ORT KEY, /* short key press */
LO	NG_KEY /* long key press */
} KE	Y_State;
KEY	_State GetKey(void);
requi	ement the function GetKey() without the usage of interrupts, with followin rements:
2.	Return NO_KEY if the key is not pressed.
	• •
4	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.
4.	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc longer than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not blocklonger than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc longer than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc longer than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not blocklonger than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not blocklonger than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc longer than needed.
	If the key is pressed, debounce it for 50 ms.  If the key is pressed for less or equal than 500 ms, return SHORT_KEY.  For a short key press or long key press detection, the function shall not bloc longer than needed.

ou a	re using GIT as VCS. Explain what the following actions mean. Use 'remote
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
ou a eposi	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
Tou a eeposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
Tou a eeposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
Tou a eeposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
Tou a eeposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
Tou a eeposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.
ou a eposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.  Elone
You a eposi (a) c	are using GIT as VCS. Explain what the following actions mean. Use 'remote itory', 'local repository', 'index' and 'working directory' in your answers.

(d) c	commit
List t	hree different clock sources which can be used for the microcontroller we used
in IN	TRO:
in IN'	TRO:
in IN'	PRO:
in IN'	TRO:
	RO:  1 8
estion	1 8
  uestion	1 8
uestion We us byte	1 8
 estion We us	sed in our shell parser the following xatoi() interface:  UTIL1_xatoi(const unsigned char **str, long *res);
 estion We us	sed in our shell parser the following xatoi() interface:  UTIL1_xatoi(const unsigned char **str, long *res);
uestion We us byte	sed in our shell parser the following xatoi() interface:  UTIL1_xatoi(const unsigned char **str, long *res);
iestion We us byte	sed in our shell parser the following xatoi() interface:  UTIL1_xatoi(const unsigned char **str, long *res);
uestion We us byte	Points: [3] sed in our shell parser the following xatoi() interface:  UTIL1_xatoi(const unsigned char **str, long *res);  in why it is using **str and not *str:

-	n extract of a header fil			Points: [3]
/* drv.h #ifdef DR #define D /* header #endif	RV_H_	s here */		
Using the a	above header file, you fin m.	d out that thing	s are not working pr	operly. Identify
	ntages and disadvantage			
The micro table conta	controller we used in IN ains the address of the reset if the whole vector	TRO are using vector function.	a table for the vect Explain what will	ors: The vector happen during

 $\overline{\text{Page 9 of 52}}$ 

Reached: \_\_\_\_

fundamental differences between using EnterCritical() Exitated xSemaphoreTake() xSemaphoreGive() with respect to interrupts	Expla
M board we used an interrupt to detect button presses on the joystick	)n th hield
	On th shield
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On th shield
M board we used an interrupt to detect button presses on the joystick	On th hield
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On th hield
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	)n th hield
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On the chield (a) V
	On the chield (a) V
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On the chield (a) V
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On the hield (a) V
M board we used an interrupt to detect button presses on the joystick we not know which button was pressed from the interrupt source?	On the chield (a) V

Page 10 of 52

Reached: \_\_\_\_

from 7 points



Figure 1 shows three different circuits how to connect a switch to a microcontroller. Rate them either 'good', 'fair' and 'poor' with a short explanation.

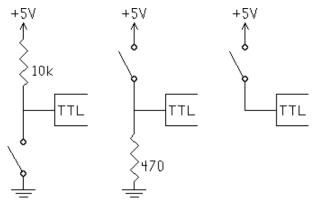


Figure 1: Three Switches

	•		•	•	٠		•	•	•	•	•	•	•	•	•	٠				•	 •	٠	٠		•	•	٠	•	٠	•	•	•	•		•	•		•		•	٠	•	٠	•			•	•	•	٠			•	•		 	•		•	٠	•	•	٠	٠	•	•	٠		
	•										•												•			•	•		•			•									•			•						•						 			•								•		
																										•																											•			 													
																										•										•		•																		 												•	
																										•										•		•																		 												•	
																										•										•		•																		 								•				•	
											•			•								•	•			•	•		•			•									•		•	•						•						 	•		•				•				•		
Qı	s T																																											•	•	•	•	•	•	•	•	•	•	•	· •	•	•	P	'c _	)i	i <b>r</b> .	11	ts —	5:	:	[	3	]	
	π.	1		۰.		_		т	. T/	71	ο.	т	ı,	л	1	ΛT	т	1/	_	١				_		1																																											

Explain the problem and provide a solution how to fix it:

• • • •	
Expla	in why FreeRTOS is using 'number of addressable units' as parameters to create
a task	x, and not 'number of bytes':
n R	18Points: [4] ΓΟS is using following interface for a 'wait' API call:
An R	18Points: [4]
An R' void (a) I	18Points: [4] ΓΟS is using following interface for a 'wait' API call:
An R' void  (a) I	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rou-
An R' 'oid  (a) I	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rou-
An R'	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rou-
An R' void (a) I	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rou-
An R'	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rou-
An R' void  (a) I t	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rouines and not 'number of milliseconds':
An R' void  (a) I t	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rouines and not 'number of milliseconds':  Cou need to wait for 200 milliseconds with vTaskDelay(), but you do not want
(a) I t	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rouines and not 'number of milliseconds':  Cou need to wait for 200 milliseconds with vTaskDelay(), but you do not want
(a) I t	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rouines and not 'number of milliseconds':  You need to wait for 200 milliseconds with vTaskDelay(), but you do not want
An R'  /oid  (a) I  t	TOS is using following interface for a 'wait' API call:  vTaskDelay(portTickType xTicksToDelay);  Explain why FreeRTOS is using 'number ticks' as parameters to the delay rouines and not 'number of milliseconds':  Cou need to wait for 200 milliseconds with vTaskDelay(), but you do not want

Question 24
Question 25Points: [2  Explain reasons when you should use a VCS:
Question 26
Question 27
Question 28
Question 29

Que	estion 30Points: [2]
-	Explain reasons why we used a Shell in INTRO:
Que	estion 31
Que	estion 32
Que	estion 33
Que	estion 34Points: [2] What are stdin, stdout and stderr?


indicate when a task is running. You can ignore the overhead in the task loop and

the overhead in the RTOS/scheduler itself.

```
static portTASK FUNCTION(T3, pvParameters) { /* priority 3 task */
  portTickType xLastTime = xTaskGetTickCount();
  \mathbf{for}(;;) { /* task time is 5 ms including overhead */
    DoWorkFor5ms(); /* this needs 5 ms */
    vTaskDelayUntil(&xLastTime, 25/portTICK RATE MS);
  } /* loop forever */
static portTASK FUNCTION(T2, pvParameters) { /* priority 2 task */
  \mathbf{for}\;(\;;;)\;\;\{\;\;/*\;\;task\;\;time\;\;is\;\;10\;\;ms\;\;including\;\;overhead\;\;*/
    {\rm DoWorkFor 10ms}\,(\,)\,\,;\,\,\,\,/*\,\,\,t\,h\,i\,s\,\,\,t\,a\,k\,e\,s\,\,\,10\,\,\,ms\,\,*/
    vTaskDelay(30/portTICK RATE MS);
  } /* loop forever */
static portTASK_FUNCTION(T1, pvParameters) { /* priority 1 task */
  \mathbf{for}(;;) \ \{ \ /* \ task \ time \ is \ 2 \ ms \ including \ overhead \ */
    {
m DoWorkFor2ms}(); /* this needs 2 ms */
    vTaskDelay (20/portTICK_RATE_MS);
  } /* loop forever */
```

Evaluate following statements for the ANSI strcmp() and strncmp() library functions as used in the INTRO program:

- ± strcmp() has one additional parameter compared to strncmp().
- ± strcmp() and strncmp() return both zero, if the two argument strings are equal.
- ± Using strncmp() allow to compare two strings up to a given offset.
- ± The parameter passing for strncmp() is usually more efficient than the parameter passing for strcmp().

Page 16 of 52 Reached: \_\_\_\_\_

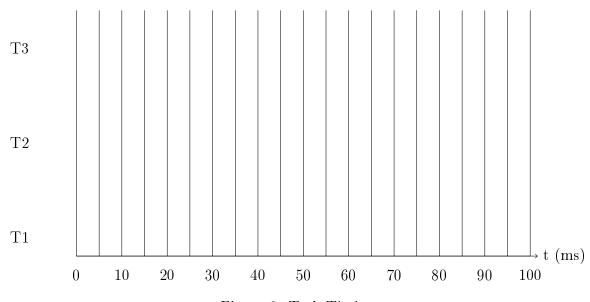


Figure 2: Task Timing

Question 37 1 Point (Bonus)
List a compelling reason why normally only debounced switches should be used at
an interrupt generating input pin only:
Question 38
Determine the value of variable j after calling the function foo():
#define SUM(a,b) (a)+b #define MUL(a,b) a*b
#define CALL(a) bar(a)
<pre>int bar(int i) {   return MUL(i,SUM(3,MUL(5,10))); } int j;</pre>
<pre>void foo(void) {     j = CALL(SUM(10,2)); }</pre>

38. \_\_\_\_\_

Consider following semaphore definition:

A semaphore S denotes an integer variable, which is accessed through the following two atomic operations:

Using the two operations, it is possible for a program to implement a critical section or a mutex, as in the following example:

```
Semaphore S = 1; /* initialize semaphore to 1 */
```

Each process or task is able to protect its critical section:

```
while(1) {
   DOWN(S);
   /* in critical section */
   UP(S);
}
```

(a)	Is the implementation with UP() and DOWN() using a 'busy waiting'? Justify your answer.	[2]
(b)	The tasks are using a priority based scheduling. Is it possible that the above solution leads to a priority inversion problem? Provide an rationale for your answer and provide an example why.	[2]

Page 18 of 52 Reached: \_\_\_\_\_

both D and E can start.

(c) Implement in pseudo code the tasks A, B, C, D and E. Use UP() and DOWN()

on the needed semaphores to implement following scenario: After both A and B have finished their work, task C can start. After task C has finish its work,

[6]

Total: 1
10001. 1
Question 40
Explain the <i>Priority Inversion</i> problem and illustrate it with an example:
••••••••••••••
••••••
Question 41

Page 19 of 52

Reached: \_\_\_\_\_

from 6 points

A device is connected to a microcontro (PTB: Data) and two control signals (P	Deller as in Figure 3. It is using a data bus PTA0: status; PTA1: control). The protocol in Figure 3. Implement the protocol using a
New	Data available Prepare next output
	PTB0-7 Data
HCS08 PTA0 ← Status Output	PTA1 /
PTA1 Control	Control
PTB0-7 Data	PTA0 /
	Status
Figure 3: Ou	utput Device
Gadfly synchronization method.	
Note: use pseudo code as SetPAOasIn naming to make clear what it does	nput(), but make sure with comments and
(a) Implement the function Init() whi variables.	ch initializes the ports and any other required

·		
•		
,		
		Total:
estion	$143.\ldots$ Points	s: [4]
(a) T	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x47}_g$ the corresponding Binary code <sub>b</sub> :	nding
	Solic Dillary codes.	
	(a)	
8 (b) D	·	
8 (b) D	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspong 8bit Binary code $_b$ :	oond-
8 (b) D	(a)	oond-
8 (b) E in	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code $_b$ :	oond- —— Tota
8 (b) E in	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspong 8bit Binary code $_b$ :	oond- —— Tota
(b) E in Consider typed	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	pond- —— Tota
(b) E in Consider typed LED	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	pond- —— Tota
(b) E in Consider typed LED LED	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	pond- —— Tota
(b) Doing to the consideration of the consideration	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code $_b$ :  (b)  A 44. Points ider following ANSI-C implementation: $\mathbf{def} \ \mathbf{enum} \ \mathbf{LED\_Set} \ \{ 0\_0 = 1 \ , \\ 0\_1 = 2 \ , \\ 0\_2 = 4 \ , \\ 0\_3 = 8 \ $	pond- —— Tota
estion Consider LED LED LED LED	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond Sbit Binary code.  (b)	oond- ——— Tota
(b) E in  consider  typed LED LED LED LED	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code $_b$ :  (b)  A 44. Points ider following ANSI-C implementation: $\mathbf{def} \ \mathbf{enum} \ \mathbf{LED\_Set} \ \{ 0\_0 = 1 \ , \\ 0\_1 = 2 \ , \\ 0\_2 = 4 \ , \\ 0\_3 = 8 \ $	oond- —— Tota
(b) E in Consider typed LED LED LED LED LED LED LED	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	oond- —— Tota
(b) E in Consider LED LED LED LED Altern	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	pond- —— Tota
estion Considered LED LED LED LED LED Altern #defin	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	pond- —— Tota
estion Considered LED LED LED LED LED LED LED LED Hdefin #defin #defin	Determine for the 8bit (binary reflected) Gray code $\mathbf{0x74}_g$ Code the correspond 8bit Binary code <sub>b</sub> :  (b)	oond- —— Tota

(b) List pros and cons for using #define's for above example:	
c) Explain the reason why the values 1 2 4 8 are used (and not 0 1 2 3):	
(c) Explain the reason why the values 1, 2, 4, 8 are used (and not 0, 1, 2, 3):	
(c) Explain the reason why the values 1, 2, 4, 8 are used (and not 0, 1, 2, 3):	
(c) Explain the reason why the values 1, 2, 4, 8 are used (and not 0, 1, 2, 3):	
(c) Explain the reason why the values 1, 2, 4, 8 are used (and not 0, 1, 2, 3):	
(c) Explain the reason why the values 1, 2, 4, 8 are used (and not 0, 1, 2, 3):	

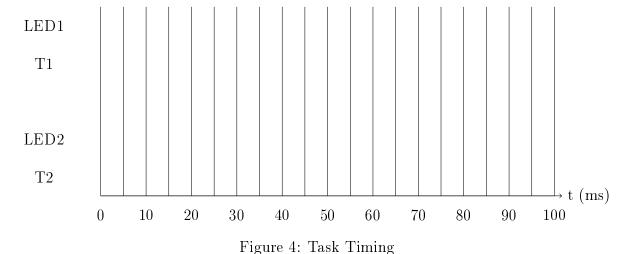
Page 22 of 52

Reached: \_\_\_\_

In this example FreeRTOS is used in priority based pre-emptive mode. Tick timer is set up to 10 ms, and the processor is running at maximum speed. The two LED's are switched off at the beginning. Task T1 has priority 3, and task T2 has priority 2. T1 has to work for 15 ms, and T2 for 2 ms (the overhead within the for loop can be ignored). Both tasks have been created with xTaskCreate() right after each other, and the scheduler has been started with vTaskStartScheduler() at the time t=0 ms.

```
static portTASK FUNCTION(T1, pvParameters) { /* priority 3 task */
  portTickType xLastTime = xTaskGetTickCount();
  LED1 Off();
  \mathbf{for}\;(\;;;)\;\;\{\;\;/*\;\;task\;\;time\;\;is\;\;15\;\;ms\;\;including\;\;overhead\;\;*/
    DoWorkFor15ms();
    LED1 Neg();
    vTaskDelayUntil(&xLastTime, 25/portTICK RATE MS);
  } /* loop forever */
static portTASK_FUNCTION(T2, pvParameters) { /* priority 2 task */
  portTickType xLastTime = xTaskGetTickCount();
 LED2 Off();
  \mathbf{for}(;;) { /* task time is 2 ms including overhead */
    LED2 Neg();
    DoWorkFor2ms();
    vTaskDelayUntil(&xLastTime, 30/portTICK RATE MS);
  } /* loop forever */
```

Draw in Figure 4 a timing diagram for the execution of the two tasks for the first 100 ms (after t0=0ms). Indicate the state of both LED's (LED1 and LED2). Use a bar to indicate when a task is running, and a bar to indicate the time when a LED is on. You can ignore the overhead in the task loop and the overhead in the RTOS/scheduler itself.



Page 23 of 52 Reached: \_\_\_\_\_

Give ity) gett	on 48
•	$J_1 = \{10 \text{ms}, LS_1, 10 \text{ms}, US_1, 10 \text{ms}\}: total 30 ms task time.$
	$J_2 = \{10 \text{ms}, LS_2, 10 \text{ms}, LS_1, 10 \text{ms}, US_1, 10 \text{ms}, US_2, 10 \text{ms}\}: total 50 ms task time.$
•	$J_3 = \{10 \text{ms}, LS_1, 30 \text{ms}, US_1, 10 \text{ms}\}: total 50 ms task time.$
	Calculate the $Priority\ Ceiling\ for\ S_1$ :
( )	
	(a)
(b)	Calculate the Priority Ceiling for $S_2$ :
	(b)
(c)	The delta between $t_n$ and $t_{n+1}$ in Figure 5 is 10 ms. Following times are defined:
	• $J_3$ starts at time $t_0$ .
	• $J_2$ starts at time $t_2$ .
	• $J_1$ starts at time $t_4$ .
	Complete and extend the timing diagram in Figure 5 indicating the task execution times. Mark with an $LS_x$ the request for a semaphore x, and with $US_x$ the release of a semaphore.
	Pon 49
• • • •	

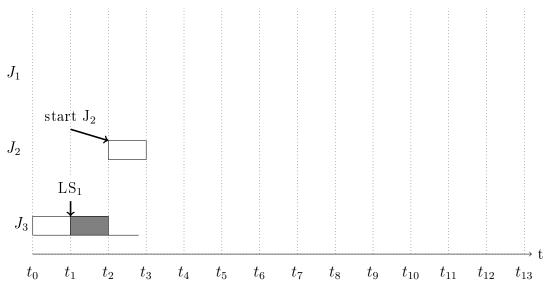


Figure 5: PCP

Question 50......Points: [1] The following code

s = k0\*e + k1\*(e-eprev)/Ta;

$$s = k0*e + k1*(e-eprev)/Ta;$$
  
 $eprev = e;$ 

Implements in software a closed loop control using the following method:

- $\bigcirc$  PD
- $\bigcirc$  ID
- $\bigcirc$  PI
- $\cap$  P
- $\bigcirc$  PID

Consider the H-Bridge we used for our INTRO lab. You have to implement a software which is able to drive the motor, from -100\% (full speed backward) to +100% (full speed forward). Explain how you realize this with the signals you have available.

.....

Question 52......Points: [2]

(a) Determine for the 8bit Binary Code  $0x31_b$  the corresponding 8bit (binary reflected) Gray  $code_q$ :

(a) \_\_\_\_

Question 52 continuous on the next page...

Reached: \_

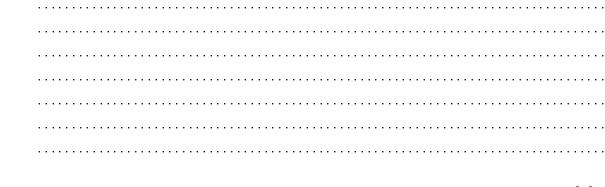
[1]

[1]

	(b)
	53Points: [1] the pseudo code for a closed loop control system:
$\mathbf{s} = \mathbf{a}$	+= e; *e + b*Ta*esum;
This s	oftware controller implements a
	O PD controller
	○ ID controller
	O PI controller
	O P Controller
	O PID Controller
For the	54
	± The quadrature encoder produces with C1 and C2 a 4 bit Gray encoded signal.
	$\pm$ In case C1 is not working or not available (only C2 is producing a signal) then we still can determine the motor direction.
	$\pm$ With minimal four consecutive codes it is possible to determine the motor direction.
	$\pm$ Using the two signals C1 and C2 it is possible to determine the absolute motor position.
Given	following program to debounce a switch on a board. RawKeyPressed() returns the switch is pressed. The function DebounceSwitch() gets called periodically 3 ms.
Sta Sta if (	DebounceSwitch(void) { tic word16 State = 0; te = (State <<1)   !RawKeyPressed()   0xF800; State==0xFC00) return TRUE; urn FALSE;
	mine the switch release debouncing time in milliseconds which is used to dee the switch.

Page 26 of 52 Reached: \_\_\_\_\_

(a)	3 FreeRTOS examples for <i>Methods</i> :
b)	3 FreeRTOS examples for <i>Properties</i> :
(c)	3 FreeRTOS examples for <i>Events</i> :
	connect a push button which is <i>not</i> debounced to interrupt input pin, using an rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?
	rnal pull-up resistor. You configure the pin to generate an interrupt whenever
	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?
	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?
   	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?  on 58
her  stic	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?  on 58
tic	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?  on 58
tio	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?  on 58
stic	rnal pull-up resistor. You configure the pin to generate an interrupt whenever e is a falling edge. What are the things you have to consider?  on 58



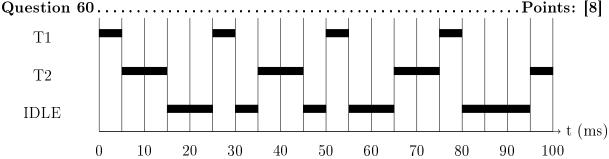


Figure 6: FreeRTOS Task Schedule

Given the FreeRTOS task scheduling pattern as shown in Figure 6 for the first 100 milliseconds. The RTOS is in preemptive scheduling mode. The RTOS IDLE task is running with priority tskIDLE\_PRIORITY. The tasks T1 and T2 are implemented as following:

```
static portTASK_FUNCTION(T1, pvParameters) {
   for (;;) {
      DoWorkT1();
      vTaskDelay(delayT1);
   } /* loop forever */
}

static portTASK_FUNCTION(T2, pvParameters) {
   for (;;) {
      DoWorkT2();
      vTaskDelay(delayT2);
   } /* loop forever */
}
```

- (a) Determine the value of delayT1 and delayT2 in milliseconds:
  - (a) \_\_\_\_\_
- (b) Each task does some work which is *not* using any blocking RTOS calls. Determine the time of work for DoWorkT2() and DoWorkT1() in *milliseconds*:
  - (b) \_\_\_\_\_

[2]

[2]

(c)	What is the tick timer period of the RTOS? Explain the reasoning behind your answer.
(d)	Could it be possible that T1, T2 and IDLE share the same priority level? Justify your answer.
.esti	on 61Points: [2]
Giv	en the following enumeration type:
typ	edef enum { RED, BLUE, GREEN, YELLOW, BLACK } ColorsEnum;
For	edef enum { RED, BLUE, GREEN, YELLOW, BLACK } ColorsEnum; the compiler we used in INTRO, the following applies for ANSI-C (with default spiler settings):
For	the compiler we used in INTRO, the following applies for ANSI-C (with default
For	the compiler we used in INTRO, the following applies for ANSI-C (with default apiler settings):
For	the compiler we used in INTRO, the following applies for ANSI-C (with default apiler settings):
For	the compiler we used in INTRO, the following applies for ANSI-C (with default apiler settings):  \( \text{ sizeof(ColorsEnum)} == 1 \) \( \text{ sizeof(ColorsEnum)} == 2 \)
For	the compiler we used in INTRO, the following applies for ANSI-C (with default spiler settings):  our sizeof(ColorsEnum)==1 our sizeof(ColorsEnum)==2 our sizeof(ColorsEnum)==4
For com	the compiler we used in INTRO, the following applies for ANSI-C (with default spiler settings):  output settings:  outpu
For com	the compiler we used in INTRO, the following applies for ANSI-C (with default apiler settings):  output settings:  outpu
For com  destice Give typ typ u T	the compiler we used in INTRO, the following applies for ANSI-C (with default apiler settings):  outpiler settings
For com  Give  typ  typ  T  T	the compiler we used in INTRO, the following applies for ANSI-C (with default spiler settings):    sizeof(ColorsEnum)==1

The triggers are initialized with following function:

[2]

[2]

```
void Init (void) {
  static volatile uint8_t i;

for(i=sizeof(TriggerList)/sizeof(TriggerDesc);i>0;i--) {
    TriggerList[i].triggerTicks = 0;
    TriggerList[i].callback = NULL;
  }
}
```

- (a) Evaluate the following statements:
  - $\pm$  The function Init() will never finish.
  - $\pm$  Every element of Triggerlist is initialized with with the same values.
  - ± Using static volatile for the variable i ensures that the function Init() is reentrant.
  - ± Using volatile for the variable i ensures that the function Init() is reentrant.

strange	way, it even cr	ashes after a while.	Explain the reason:

(b) Using above implementation, you notice that your application is working in a

Given following implementation of an event module:

```
#define GET_EVENT(event) \
    (bool) (EVNT_Events[(event)/8]&(0x80>>(uint8_t)((event)%8)))
#define CLR_EVENT(event) \
    EVNT_Events[(event)/8] &= ~(0x80>>((uint8_t)((event)%8)))

static uint8_t EVNT_Events[((EVNT_NOF_EVENTS-1)/8)+1];

bool EVNT_GetEvent(EVNT_Handle event) {
    bool isSet;
    EnterCritical();
    isSet = GET_EVENT(event);
    ExitCritical();
    return isSet;
}
```

```
{\bf void} \ \ {\bf EVNT\_ClearEvent} \ ({\bf EVNT\_Handle} \ \ {\bf event} \ ) \ \ \{
  EnterCritical();
  CLR_EVENT(event);
  Exit Critical();
uint8 t EVNT\_CheckEvents(void)  {
  uint8_t i;
  EnterCritical();
```

<pre>for (i = 0; i &lt; size of (EVNT_Events) / size of (EVNT_Events[0]); i++)     if (EVNT_GetEvent(i)) {         EVNT_ClearEvent(i);         break;     } } ExitCritical(); return i; /* return the event which was set */ }</pre>	) {
(a) Using that implementation and calling EVNT_CheckEvents(), you your periodic interrupt timer does not work any more. Explain the how to fix it:	
(b) After you have fixed the previous problem, the application still d correctly. Somehow EVNT_CheckEvents() does not properly return have been set. Explain the problem and how to fix it:	<u>-</u>
uestion 64  Evaluate following statements in the context the RTOS startup, as used i lab (FreeRTOS):  ± All interrupts are disabled during the (ANSI) startup code.	

Page 31 of 52

Question 64 continuous on the next page...

Reached: \_\_\_\_

±	Global interrupts are enabled at the entry point of main().
±	Interrupts for RTOS timers are disabled at the entry point of main().
±	Interrupts for RTOS timers get enabled during startup of the scheduler.
Explain	5
If you ca	6
coding fo	an choose between the delta-time and delta-pos method for quadrature de- or slow (few steps) state changes, which one would you prefer and why?
Explain	8
Name a	9

,	
0,,,	stion 70Points: [3]
-	What is the difference between a Mutex and a normal Semaphore in FreeRTOS?
•	
,	
I	stion 71
,	
,	
i S	Stion 72
,	
,	
I	stion 73
,	
•	
	stion 74

Ωı	uestion 75Points: [3]
<b>پ</b>	Using a network stack like RNet, which things do you have to initialize?
_	
Q۱	RNet uses four different layers in the stack. List the layers with their acronyms.
Qı	What is a CRC and for what it is used for?
Qı	What is a so called Ad-Hoc Network?
Qı	List pros and cons of an Ad-Hoc Network.
Qı	Using Segger JScope, which variables can be monitored by that tool without special extensions of the program itself?

Question 81	[1]
Question 82	[1]
Question 83	
Question 84	
	• •
Question 85	
Question 86	
Question 87	

Question 88
Question 89
Question 90
Question 91
Question 92
Question 93
What is the purpose of the Anti-Windup?

Question 94	
Question 95  Provide an implementation way where you can store calibration dat	
Question 96	
Question 97  Why did we use a digital interface to the reflectance sensor and not	
Question 98 Even if all sensors of the reflectance sensor are exposed to the same measurement will differ from sensor to sensor. Why?	
Question 99  The reflectance sensor array has two red LED's. What's the purpose	

Question 100
Question 101
••••••
Question 102
Question 103
Question 104
Ougstion 105
Question 105Points: [1 List the five memory allocation schemes in FreeRTOS.
Question 106

O114	estion 107Points: [1]
	What are the two different categories of timeliness?
_	
	estion 108
Que	estion $109$ Points: [1]
	In a microcontroller you can have different kinds of clocks, list at least two:
-	estion 110
-	estion 111Points: [1] List pros and cons of using an internal reference clock (compared to an external clock generator):
	estion 112Points: [3] List reasons why to use an RTOS:

	What is the difference between preemptive and non-preemptive scheduling:
-	stion 114Points: [2] List the three standard I/O channels used by the Shell:
-	stion 115
-	stion 116Points: [1] What does the abbreviation $HAL$ mean?
7	stion 117
]	stion 118
r	stion 119

	stion 120Points: [2] Which Processor Expert generated macros can you use to protect a critical section?
	stion 121
1	stion 122
•	stion 123
-	stion 124
-	stion 125

Question 126
#defineCALC1 (2+5) #defineCALC2 (5*3)
Why is it important to use parenthesis?
Question 127Points: [2] List some advantages and disadvantages of using macros:
Question 128
Question 129
What is the fundamental difference between the two VCS: Git and SubVersion?
Question 131

<u> </u>		plain what this file does with such a
* !*.c !*.h !*. gitignore ! dev /		
For multiple Eclipse c:/mySources/commo' Virtual Group', 'Lin in your project. Now that common folder.	projects, you want to use n) which is outside of you ked Folder' or 'Linked File you have added a new so For each of the three about	e source files in a common folder (e.g. ar Eclipse project. You consider to use le' methods to use the common file(s) source file (e.g. accelerometer.c) to we methods, explain what you have to to use the new accelerometer.c file.
Given a Quadrature e signals get sampled maximum frequency	encoder as in Figure 7, gen with a periodic timer in	nerating two signals C0 and C1. Both terrupt. The disk is turning with a minimal sampling frequency in Hz to gnals C0 and C1.

Page 43 of 52 Reached: \_\_\_\_\_

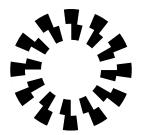


Figure 7: Quad Disk

Qı	You are using a wireless transceiver as used in the lab to transmit sensor data values to another system. Describe strategies to reduce the energy consumption for the communication. Limit your answer to the communication and transceiver only.
Qι	Calculate for the 16bit <b>Bindary Code <math>0x1722_b</math></b> the corresponding 16bit (binary reflected) Gray $code_g$ :

Page 44 of 52

Reached: \_\_\_\_\_

the gaps:

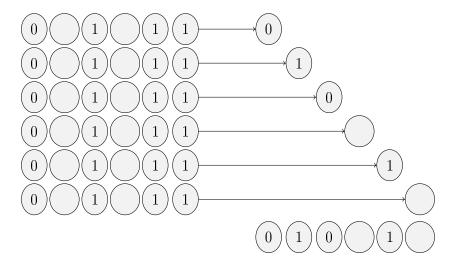


Figure 8: 6bit Gray Code

Communication Interface). The SCI is using interrupts for RX and TX. As other tasks are using the Read() and Write() functions as well, it is necessary to use critical sections.

```
void Read(char *buf, size_t bufSize) {
  /* start critical section */
 SCI Read(buf, bufSize);
 /* end critical section */
void Write(char *buf) {
  /* start critical section */
 SCI Write(&buf[0]);
 /* end critical section */
static portTASK FUNCTION(Task1, pvParameters) {
 unsigned char buf [16];
  (void) pvParameters; /* parameter not used */
 for (;;) {
    /* start critical section */
    Read(buf, sizeof(buf));
    Write(buf, sizeof(buf));
    /* end critical section */
 }
```

You identified 3 different ways for implementing the /\* start critical section

[2]

\*/ and /\* end critical section \*/ in above source. Evaluate following statements for the consequences of each implementation for the above source: [2] (a) With using DisableAllInterrupts() and EnableAllInterrupts() the following applies:  $\pm$  It increases the interrupt latency time.  $\pm$  Is the best solution with respect to RAM/ROM footprint.  $\pm$  Makes the usage of SCI\_Read() and SCI\_Write() reentrant.  $\pm$  Allows the scheduler to run inside every critical section. (b) With using FreeRTOS binary semaphore/mutex the following applies: [2]± Calling the semaphore/mutex API might trigger a context switch.  $\pm$  No context switch will happen within the critical section.  $\pm$  Only one task at a time will be inside the critical section.  $\pm$  Interrupts will occur inside the critical section. [2]

(c) With using FreeRTOS recursive semaphore/mutex the following applies:

± Calling the semaphore/mutex API might trigger a context switch.

± Multiple tasks might be within the same critical section.

± No context switch will happen within the critical section.

± No interrupts will occur inside the critical section.
(d) From the 3 solutions a), b) and c), select the one you consider as the best one for above source, and explain briefly why:

- $\bullet \ J_1 = \{LS_4, \, US_4\}$
- $\bullet \ J_2 = \{ LS_2, \, LS_1, \, US_1, \, US_2, \, LS_3, \, US_3 \}$
- $\bullet \ J_3 = \{LS_1,\, US_1,\, LS_3,\, US_3\}$
- $\bullet \ J_4 = \{LS_1, \, US_1\}$

Page 46 of 52 Reached: \_\_\_\_\_

Determine the <i>Priority Ceiling</i> for each semi- Priority Ceiling for $S_1$ :	aphore:	
	139	
Priority Ceiling for $S_2$ :	1301	
	139	
Priority Ceiling for $S_3$ :		
	139	
Priority Ceiling for $S_4$ :		
	139	
estion 140		Points: [3]
Below is a way to prevent spurious keyboard	l interrupts on a mic	erocontroller:
$egin{array}{lll}  ext{PTADD} &= 0 ext{x}00; & /* & set & port & as & input & */ \  ext{PTAD} &= 0 ext{xFF}; & /* & write & 1 's & to & data & port \  ext{PTAPE} &= 0 ext{xFF}; & /* & enable & pull & ups & */ \ \end{array}$		
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	input values */ synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	synchronization:
PTADD = 0x00; /* set port as input */ PTAD = 0xFF; /* write 1's to data port PTAPE = 0xFF; /* enable pull ups */ EnableInterrupts;  Transform the above code into an implement	to have defined	synchronization:

Page 47 of 52 Reached: \_\_\_\_\_

## 

A preemptive system has 3 tasks  $T_1$ ,  $T_2$  and  $T_3$  where the task index denotes the task priority, with 1 the lowest priority. The system is using *Priority Inheritance* for the semaphores. Create a timing diagram with the tasks in Figure 9 indicating the execution time for each task. The system has following timing:

- $t_0$ :  $T_1$  and  $T_2$  are ready,  $T_3$  is running,
- t<sub>1</sub>: The running task gets suspended
- t<sub>3</sub>: The running task gets suspended
- t<sub>4</sub>: The running task requests the semaphore
- $t_5$ :  $T_2$  and  $T_3$  get ready
- t<sub>7</sub>: The running task requests the semaphore
- t<sub>8</sub>: The task having the semaphore releases the semaphore
- t<sub>9</sub>: The task having the semaphore releases the semaphore
- $t_{10}$ : The running task terminates
- t<sub>11</sub>: The running task terminates
- t<sub>12</sub>: The running task terminates

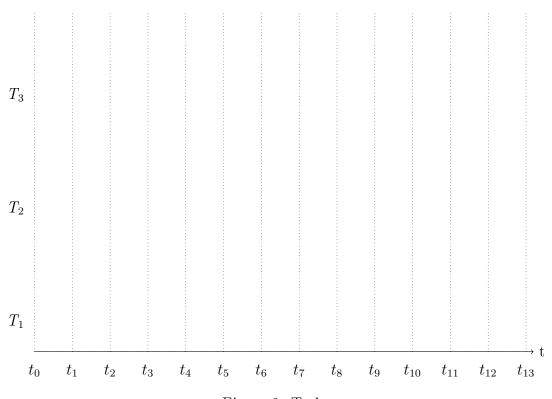


Figure 9: Tasks

Page 48 of 52 Reached: \_\_\_\_\_

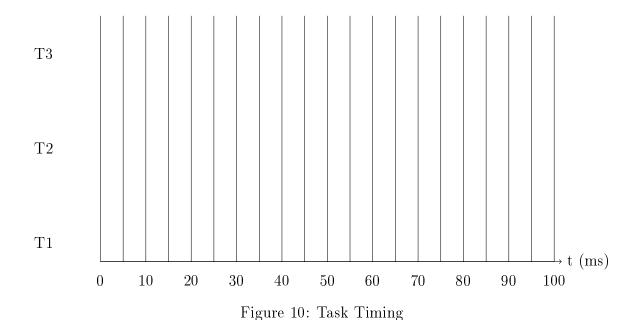
Qu	estion 143Points: [3] Identify the three serious problems in the following source:
	#define QUEUE LENGTH 5
	#define QUEUE_ITEM_SIZE size of (char_t*)
	<pre>void QUEUE_SendMessage(const char *msg) {    char *ptr = FRTOS1_pvPortMalloc(UTIL1_strlen(msg));    UTIL1_strcpy(ptr, msg);</pre>
	if (xQueueSendToBack(queueHandle, ptr, portMAX_DELAY)!=pdPASS)
	<pre>for (;;) {} /* ups? */</pre>
	}
Эu	estion 144 Points: [10]
	FreeRTOS is used in priority based preemptive mode. Tick timer is set up for 5 ms, and the processor is running at maximum speed. Task T3 has priority 3 (highest priority), task T2 has priority 2 and task T1 has priority 1. The tasks have been created with xTaskCreate() before time t=0 ms, and the scheduler is started with vTaskStartScheduler() at the time t=0 ms. Draw in Figure 10 a timing diagram for the execution of the two tasks for the first 100 ms (after t0 = 0 ms). Use a bar to indicate when a task is running. You can ignore the overhead in the task loop and the overhead in the RTOS/scheduler itself.
	<pre>static portTASK_FUNCTION(T3, pvParameters) { /* priority 3 task */ portTickType xLastTime = xTaskGetTickCount(); for(;;) { /* task time is 7 ms including overhead */     DoWorkFor7ms(); /* this needs 7 ms */     vTaskDelayUntil(&amp;xLastTime, 30/portTICK_RATE_MS);</pre>

Page 49 of 52 Reached: \_\_\_\_\_

```
} /* loop forever */

static portTASK_FUNCTION(T2, pvParameters) { /* priority 2 task */
    for (;;) { /* task time is 7 ms including overhead */
        DoWorkFor7ms(); /* this takes 7 ms */
        vTaskDelay(30/portTICK_RATE_MS);
    } /* loop forever */
}

static portTASK_FUNCTION(T1, pvParameters) { /* priority 1 task */
    for (;;) { /* task time is 4 ms including overhead */
        DoWorkFor4ms(); /* this needs 4 ms */
        vTaskDelay(28/portTICK_RATE_MS);
    } /* loop forever */
}
```



Question 145
You are working with an Eclipse based project (you can consider the one we used
for INTRO). Which kind of files are you going to put into a VCS? Which ones not
Explain briefly why.

Page 50 of 52 Reached: \_\_\_\_

## 

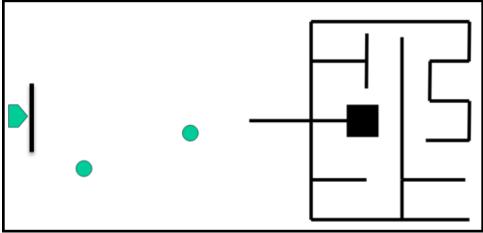


Figure 11: Maze robot

Question 147Points:	[1]
The slide material for this semester week 2 are covering the following topics:	
$\pm$ Systems and Realtime	
$\pm$ GitHub and Version Control	
$\pm$ Processor Expert and Drivers	
$\pm$ Preprocessor and Macros	
Question 148	[1]
Question 149	[1]
onverting, reacting, interacting	
o measuring, controlling, calculating	

transforming, reactive, interactive
interactive, responsive, collective
o soft realtime, hard realtime, true realtime
Question 150
$egin{array}{ll} {f int} & { m a} &= 0{ m x}1234; \ { m a} &\&= \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \$
Question 151
#define CAT(a,b) a##b void foo(void) {     CAT(t,a)(); }
$\bigcirc$ taab();
$\bigcirc$ TaAb();
$\bigcirc$ ta();
$\bigcirc$ aTbA();
$\bigcirc$ ab();
Question 152