

## Homework 2

1. For constants  $c$  and  $n_0$ , we get the following when  $n \geq n_0$ .

$$\begin{aligned} 2^{2n} &\leq c \cdot 2^n \\ \log(2^{2n}) &\leq \log(c \cdot 2^n) \\ 2n \log(2) &\leq \log(c) + n \log(2) \\ n \log(2) &\leq \log(c) \\ n &\leq \frac{\log(c)}{\log(2)} = c' \end{aligned}$$

Since no constant  $c'$  can be bigger than every possible  $n$ , this results in a contradiction. Thus  $O(2^{2n}) \neq O(2^n)$ .

2. 

```
def PhysIndexReverse(A, i):  
    return (A.front + A.size - i - 1) % A.cap
```

3. 

```
def Sum(L):  
    if L is None:  
        return 0  
    else:  
        return L.data + Sum(L.next)
```

4. When we pop 6 off of the stack we know that 5 has already been pushed onto the stack. It is thus impossible to pop 4 because we know that 5 is above it.

5. 3 of the 10 pop operations fail, so only 7 of them remove items. The peek operations have no effect on the size. 25 pushes - 7 pops = 18 items.

6. 

```
def Transfer(S, T):  
    while not S.isEmpty():  
        T.push(S.pop())
```

7. Transfer  $S$  to temporary stack  $A$  which now contains the items in reverse order. Transfer  $A$  to  $B$  to have the items in the original order. Transfer  $B$  back to  $S$ , which now has the items in the reverse order.

8. After enqueueing 8 items, the front is at index 0 and the back is at index 7 (first empty is index 8). Dequeueing 5 items moves the front to index 5, but leaves the back at index 7. Adding 3 items causes the back to loop around to index 0, and the remaining 3 moves it to 3. The front is thus at index 5 and the back is at index 3.

9. 

```
def Clear(Q):  
    if not Q.isEmpty():  
        Q.dequeue()  
        Clear()
```