CSC 220

11/3/2020

Regular Expressions

. A language over a finite alphabet is a set of strings of letters from A.

. A language over A is a subset of Ax.

the language that can be constructed from letters of an alphabet by operations such as union, concatenation and closure is a regular language.

of Sal for all act



If Land M are regular than LUM, LM, LX are also regular.

A =  $\{a,b\}$   $\Rightarrow \phi, \{a\}, \{b\}, \{\lambda\}\}$  are  $\log$ .  $\{\lambda,b\}$  is regular  $\{\lambda,b\}$  So regular  $\{a,ab\} = \{a\}, \{\lambda,b\}$  So regular

 $\{\lambda, b, bb, bbb, ---\} = \{b^n | near \} = \{b\}^k$ is regular

A regular language is often described by means of an algebraic expression called regular expression.

\$, h, a are regular for all acA.

. If R, S are regular expressions then

(R), R+S, R.S, RK are

also regular.

$$\int \int Se^{\frac{1}{2}}$$

R.E

(3)

RE over 
$$A=\{a,b\}$$

$$\lambda, \quad \alpha, \quad \alpha.\alpha, \quad \alpha^*, \quad \alpha^*+b, \quad \alpha^*.b, \dots$$

. For each R.E., we will associate a regular language L(E), where A is an alphabet and R,S are regular

ex pressions

$$L(\alpha) = \begin{cases} L(\lambda) = \lambda \\ L(\alpha) = \begin{cases} q \\ q \end{cases} \end{cases}$$

$$L(R+S) = L(R)UL(S)$$

$$L(R,S) = L(R).L(S)$$

Find the language of the regular expression a+bcx.

$$L(a+bc^{*}) = L(a)UL(bc^{*})$$

$$= L(a)U(L(b).L(c^{*}))$$

$$= \{a\}U\{b\}. \{\lambda, C, CC, CCC, --\}$$

$$= \{a\}U\{b, bC, bC, --\} = \{a, b, bC, bC, --\}$$

Single letter a, or strings of form b followed by zero or more C's.

$$is$$
 regular
$$\lambda + b + ab$$

$$L(a+b) = L(a)UL(b)$$

$$= \{a\}U\{b\} = \{a\}b\} = \{b\}a\}$$

$$= L(b+a)$$

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 $L(ab) \neq L(ba)$ 

ab +ba



 $L = \left\{ \begin{array}{l} a, aa, aaa, -- \\ -- \\ \end{array} \right\} = \left\{ \begin{array}{l} a \\ \end{array} \right\} \left\{ \begin{array}{l} a \\ \end{array}$