

# INTRO TO DYNAMIC PROGRAMMING: EXAMPLES

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QUANTITATIVE ECONOMICS 2024

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December 17, 2024

# INTRODUCTION

- Shortest Path
- Tree-cutting
- Resource Extraction
- Neoclassical Growth Model

SHORTEST PATH

## SHORTEST PATH

- We have a set of nodes,  $S$ :  $s_1, s_2, \dots, s_N$ .
- Let  $F_s$  be a set of nodes that can be reached from the node  $s$  in one step.
- $c(s, s')$  is the cost of moving from node  $s$  to node  $s' \in F_s$ .
- The goal is to find the shortest path from some starting node to some destination  $d$ .

## SHORTEST PATH

- Let  $J(s)$  be the length of the shortest path from the node  $s$  to the destination. Suppose we know  $J(s)$  for all  $s \in N$ .
- Suppose the starting node is  $s$ . We can find the shortest path by choosing a node  $s'$  such that

$$s' = \operatorname{argmin}_{s' \in F_s} c(s, s') + J(s').$$

- It must also be the case that

$$J(s) = \min_{s' \in F_s} c(s, s') + J(s')$$

- The above is the **Bellman equation** for the shortest path problem.

## SHORTEST PATH

- If we know  $J(s)$  for all  $s \in S$ , we can find the shortest path easily. But we don't know it!
- How to find  $J$ ? We only know that  $J(d) = 0$ .
- We will start with a guess  $J_0(s)$  for all  $s \in N$ :
  - $J_0(s) = M$  for all  $s \neq d$ ;
  - $J_0(s) = 0$ , where  $M$  is a very large number.

## SHORTEST PATH

- We will use the following algorithm:
  1. Set  $k = 0$ .
  2. Set  $J_{k+1}(s) = \min_{s' \in F_s} c(s, s') + J_k(s')$  for all  $s \in N$ .
  3. If  $J_{k+1} = J_k$ , stop. Otherwise, set  $k = k + 1$  and go to the previous step.

TREE-CUTTING



## TREE-CUTTING

- There is a tree of size  $s$ .
- You can either **cut down** the tree now and sell the wood or **wait** until next period:
  - If you **cut down** the tree now, you get  $f(s)$  dollars. You do not have the tree anymore.
  - If you **wait** until the next period, the tree grows to size  $s + h$ , unless  $s = \bar{S}$ , which is the maximum size of the tree. If  $s = \bar{S}$  the tree stays the same size.
- Your objective function is to maximize the present discounted value of money you get from cutting down the tree.
- The discount factor is  $\frac{1}{1+r}$ , where  $r > 0$  is the interest rate.

## TREE-CUTTING

- Let  $v(s)$  be the value function. State: the size of the tree,  $s$ .
- Bellman equation:

$$v(s) = \max \left\{ \underbrace{f(s)}_{\text{cut down}}, \underbrace{\frac{1}{1+r} v(\min\{\bar{S}, s+h\})}_{\text{wait}} \right\}.$$

- We can use value function iteration to solve this problem.

## TREE-CUTTING

1. Suppose that cutting down the tree costs  $c > 0$  dollars.
2. Suppose that if you wait until the next period, the tree might not grow at all - it happens with probability  $p$ .
3. Suppose that if you wait until the next period, the tree might die which means that you have to cut it down and get only  $1/2$  of the value of the tree. It happens with probability  $p$ .
4. Suppose that if you wait until the next period, the tree might become sick which means that it will stop growing. It happens with probability  $p$ . With probability  $q$  a sick tree can recover.
  - How to modify the Bellman equation?

## TREE-CUTTING

- Suppose that cutting down the tree costs  $c > 0$  dollars.
- Bellman equation:

$$v(s) = \max \left\{ \underbrace{f(s) - c}_{\text{cut down}}, \underbrace{\frac{1}{1+r} v(\min\{\bar{S}, s+h\})}_{\text{wait}} \right\}.$$

## TREE-CUTTING

- Suppose that if you wait until the next period, the tree might not grow at all - it happens with probability  $p$ .
- Bellman equation:

$$v(s) = \max \left\{ \underbrace{f(s)}_{\text{cut down}}, \underbrace{\frac{1}{1+r} [(1-p) \cdot v(\min\{\bar{S}, s+h\}) + p \cdot v(s)]}_{\text{wait}} \right\}.$$

## TREE-CUTTING

- Suppose that if you wait until the next period, the tree might die which means that you have to cut it down and get only  $1/2$  of the value of the tree. It happens with probability  $p$ .
- Bellman equation:

$$v(s) = \max \left\{ \underbrace{f(s)}_{\text{cut down}}, \underbrace{\frac{1}{1+r} \left[ (1-p) \cdot v(\min\{\bar{S}, s+h\}) + p \cdot \frac{1}{2} f(s) \right]}_{\text{wait}} \right\}.$$

## TREE-CUTTING

- Suppose that if you wait until the next period, the tree might become sick which means that it will stop growing. It happens with probability  $p$ . With probability  $q$  a sick tree can recover.
- Additional state variable - health of the tree.
- Bellman equation ( $H$  is for healthy,  $S$  is for sick):

$$v^H(s) = \max \left\{ \underbrace{f(s)}_{\text{cut down}}, \underbrace{\frac{1}{1+r} \left[ (1-p) \cdot v^H(\min\{\bar{S}, s+h\}) + p \cdot v^S(s) \right]}_{\text{wait}} \right\}$$
$$v^S(s) = \max \left\{ \underbrace{f(s)}_{\text{cut down}}, \underbrace{\frac{1}{1+r} \left[ (1-q) \cdot v^S(s) + q \cdot v^H(\min\{\bar{S}, s+h\}) \right]}_{\text{wait}} \right\}$$

## RESOURCE EXTRACTION



## RESOURCE EXTRACTION

- You have a resource of size  $s$  that you can extract.  $s$  is an integer.
- The current price of the resource is  $p \stackrel{iid}{\sim} \phi$  with  $N$  possible values:  $p_1, p_2, \dots, p_N$ .
- The cost of extracting  $x$  units of the resource is  $c(x)$ . Extracting  $x$  units of the resource reduces the size of the resource to  $s - x$  units next period. It is never possible to extract more than  $s$  units of the resource.
- The objective is to maximize the expected present discounted value of the profit from extracting the resource:

$$\mathbb{E}_0 \sum_{t=0}^{\infty} \beta^t (p_t x_t - c(x_t))$$

## RESOURCE EXTRACTION

- Bellman equation:

$$v(s, p) = \max_{x \in \{0, 1, \dots, s\}} \left\{ px - c(x) + \beta \sum_{p' \in P} v(s - x, p') \phi(p') \right\}.$$

## NEOCLASSICAL GROWTH MODEL

## NEOCLASSICAL GROWTH MODEL

- A representative agent's problem

$$\sum_{t=0}^{\infty} \beta^t u(c_t), \quad \text{s.t. } k_{t+1} = f(k_t) + (1 - \delta) k_t - c_t, \quad k_0 \text{ given.}$$

- $\beta \in (0, 1)$  is the discount factor,  $c_t \geq 0$  is the consumption,  $k_t \geq 0$  is the stock of capital.
- $u(\cdot)$  is the utility function,  $f(\cdot)$  is the production function,  $0 < \delta \leq 1$  is the depreciation rate.
- We will assume  $u(c) = \frac{c^{1-\gamma}}{1-\gamma}$ ,  $f(k) = k^\alpha$ , with  $\gamma > 0$ ,  $\alpha \in (0, 1)$ .
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## NEOCLASSICAL GROWTH MODEL

- Bellman equation:

$$v(k) = \max_{c \in [0, f(k) + (1-\delta)k]} \{u(c) + \beta v(f(k) + (1-\delta)k - c)\}.$$

- Alternatively, we can write it as:

$$v(k) = \max_{k' \in [0, f(k) + (1-\delta)k]} \{u(f(k) + (1-\delta)k - k') + \beta v(k')\}.$$

## NEOCLASSICAL GROWTH MODEL

- Note: in the problem itself there is nothing that says there is only a finite number of  $k$ .
- We will discretize the state space and create a finite grid of  $k$ :  $k_1, k_2, \dots, k_N$ .
- We need to decide on the grid size and the grid points. These are properties of the solution method, not the problem itself!
- What we will find is only an approximation to the true value and policy function.

## NEOCLASSICAL GROWTH MODEL

- Let  $k^*$  be a level of capital that satisfies

$$1 = \beta \left[ \alpha k^{*\alpha-1} + 1 - \delta \right].$$

- In this model we call  $k^*$  the **steady state** level of capital. If  $k_t = k^*$ , then  $k_{t+1} = k^*$ .
- We also know that if we start from  $k_0 > 0$ , then  $k_t \rightarrow k^*$  as  $t \rightarrow \infty$ .
- It makes sense to have a grid that includes  $k^*$ .