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# Convex hull

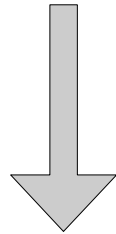
*(divide and conquer algorithm)*

*Prerequisites:*

- *recursion.*

# Divide and conquer

24	10	6	18	9	3	25	32	16
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3	6	9	10	16	18	24	25	32
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Many techniques which work for sorting, also work for finding convex hull.

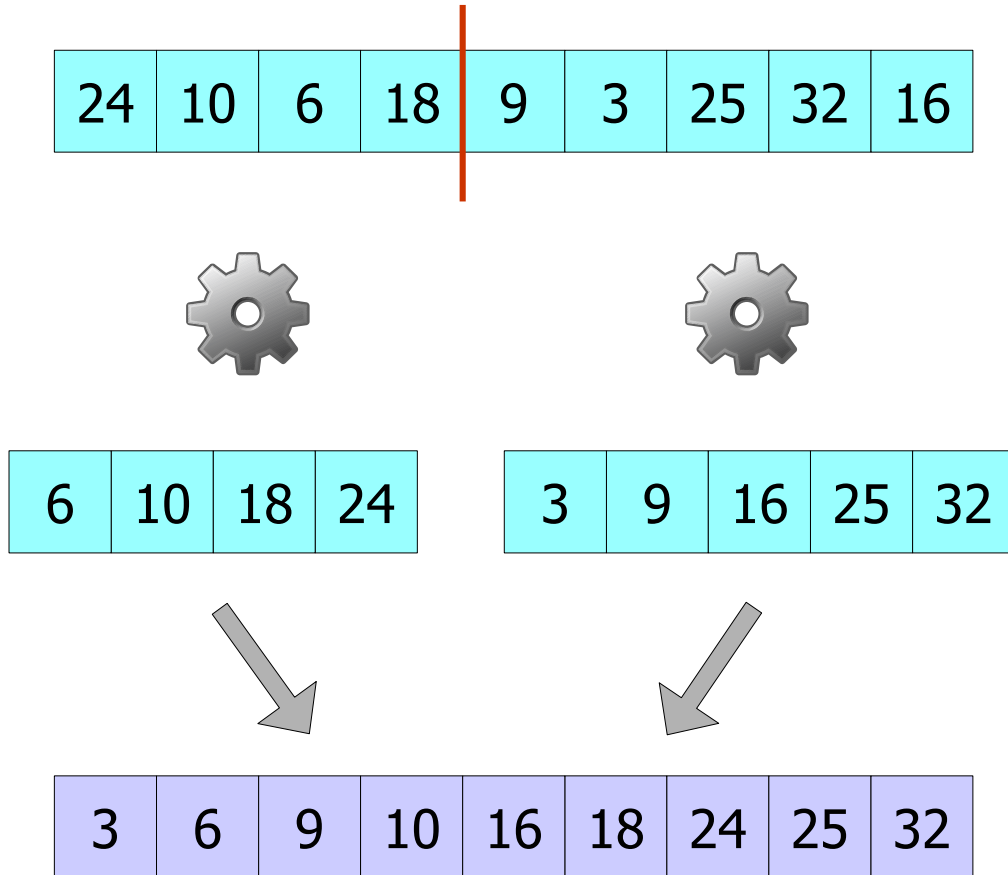
Merge sort uses divide and conquer technique.

... let's apply it for finding the convex hull too.

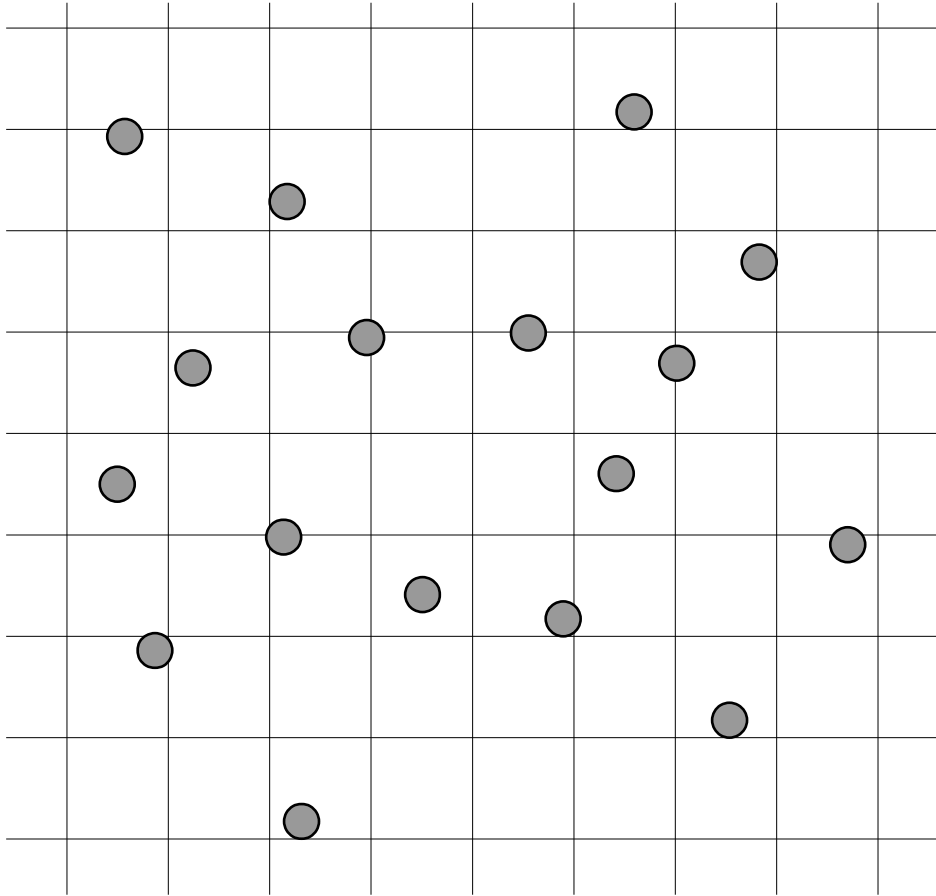
# Divide and conquer

Let's recall how Merge sort works:

- Divide the array in **2** equal halves,
- Sort them independantly (recursively),
- Merge them.

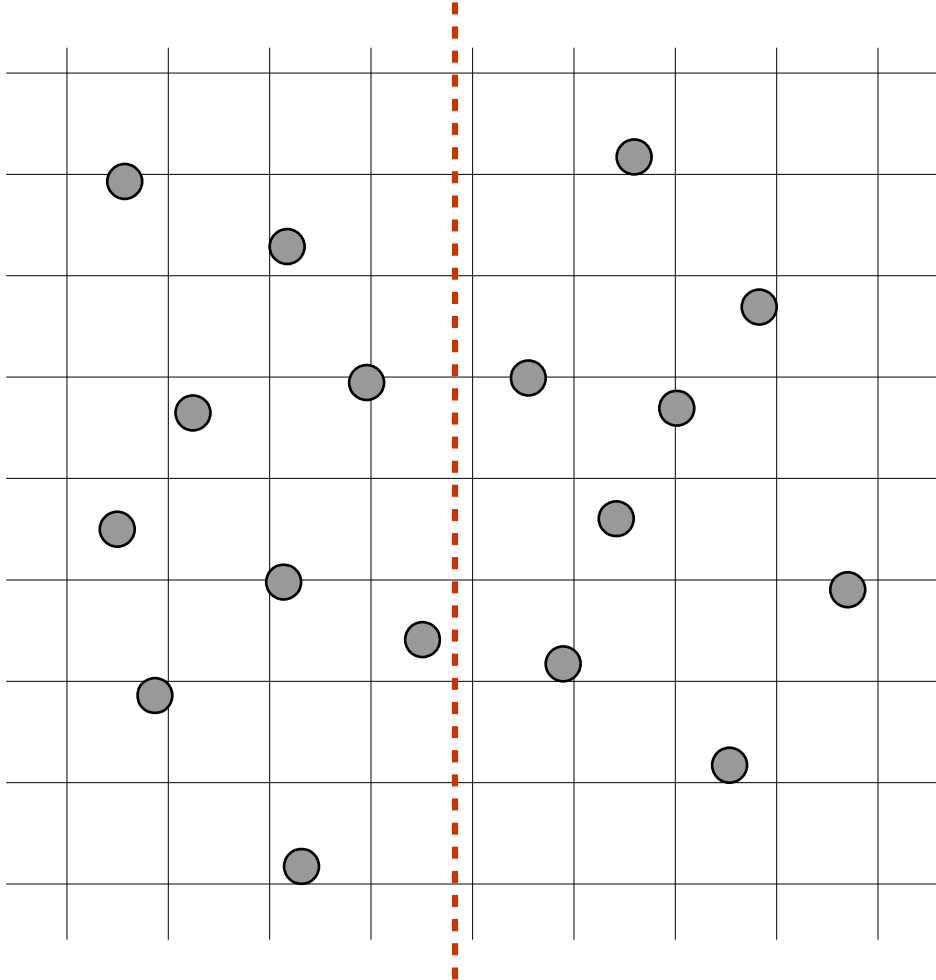


# Divide and conquer



**Question:** How can we apply the same technique to convex hull finding?

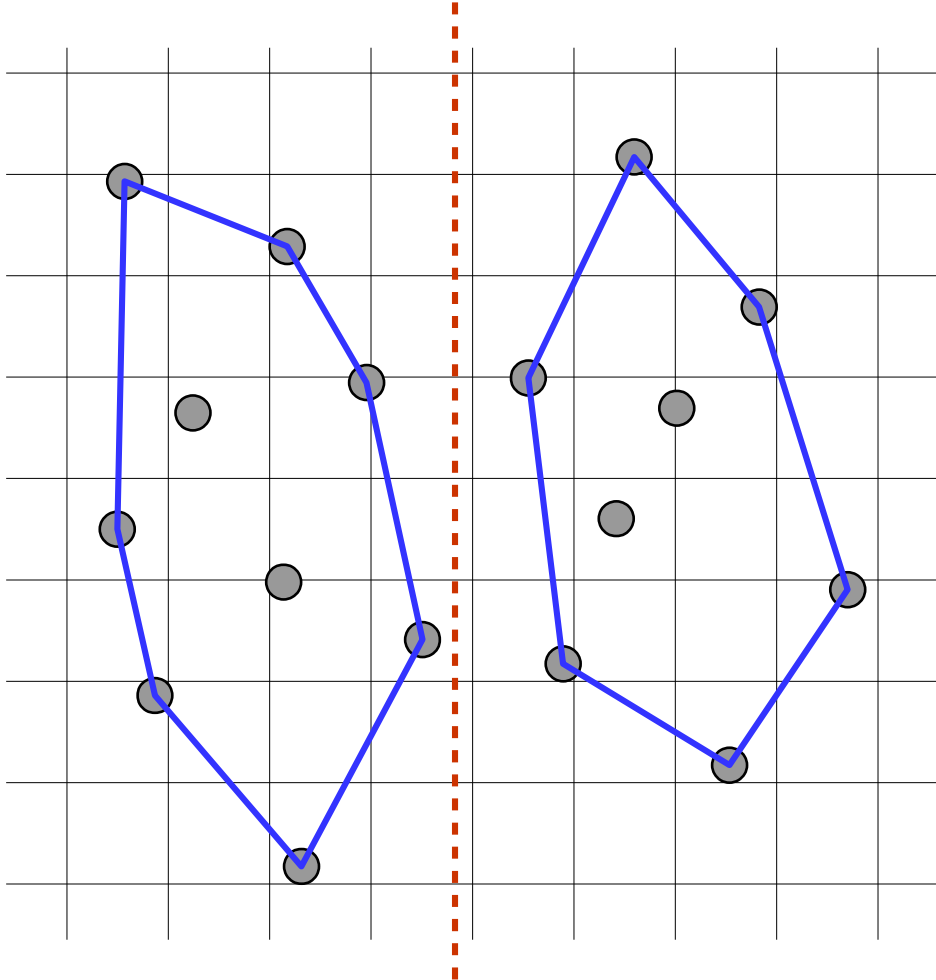
# Divide and conquer



***Answer.***

- Divide the points into 2 equal parts,

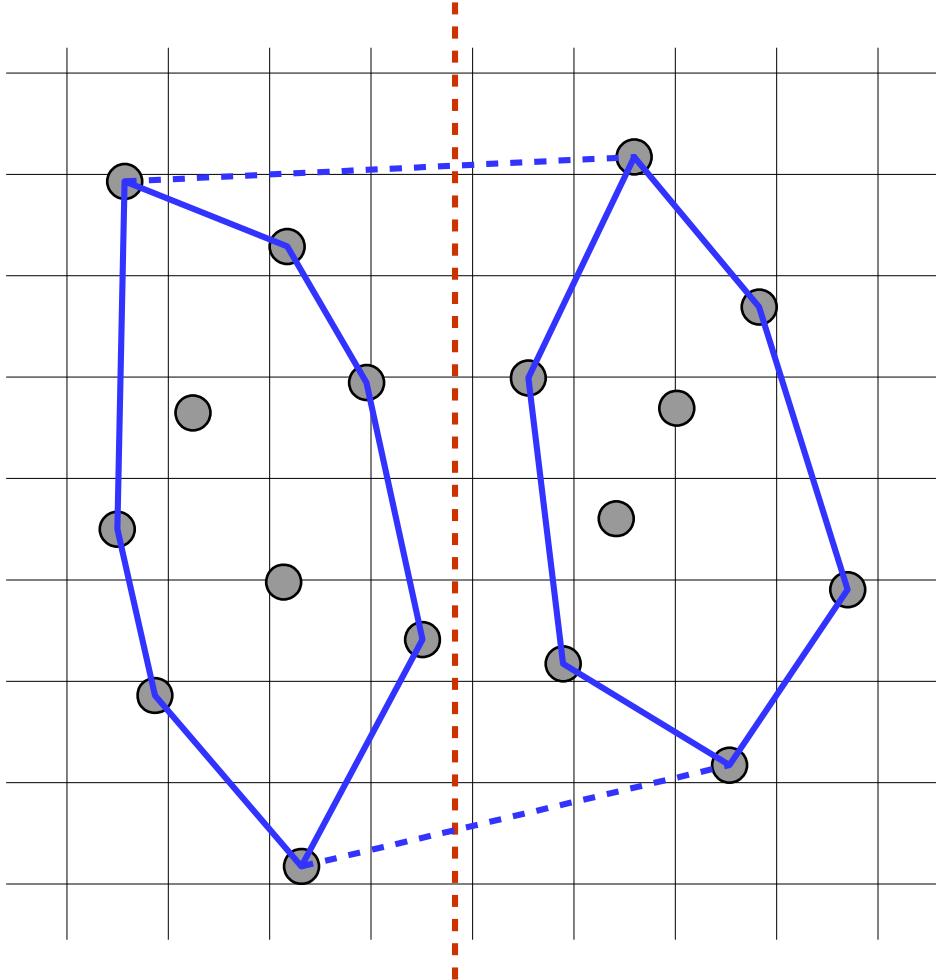
# Divide and conquer



## ***Answer.***

- Divide the points into 2 equal parts,
- Construct convex hulls for every part (recursively),

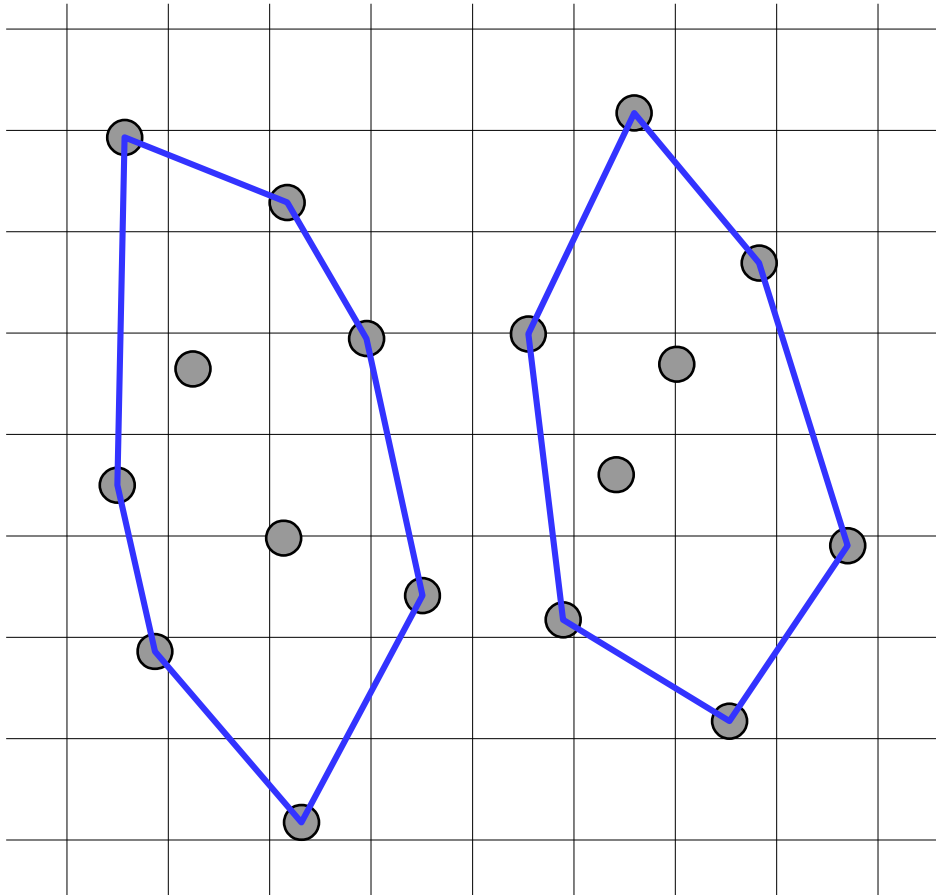
# Divide and conquer



## ***Answer.***

- Divide the points into 2 equal parts,
- Construct convex hulls for every part (recursively),
- Merge them.

# Divide and conquer

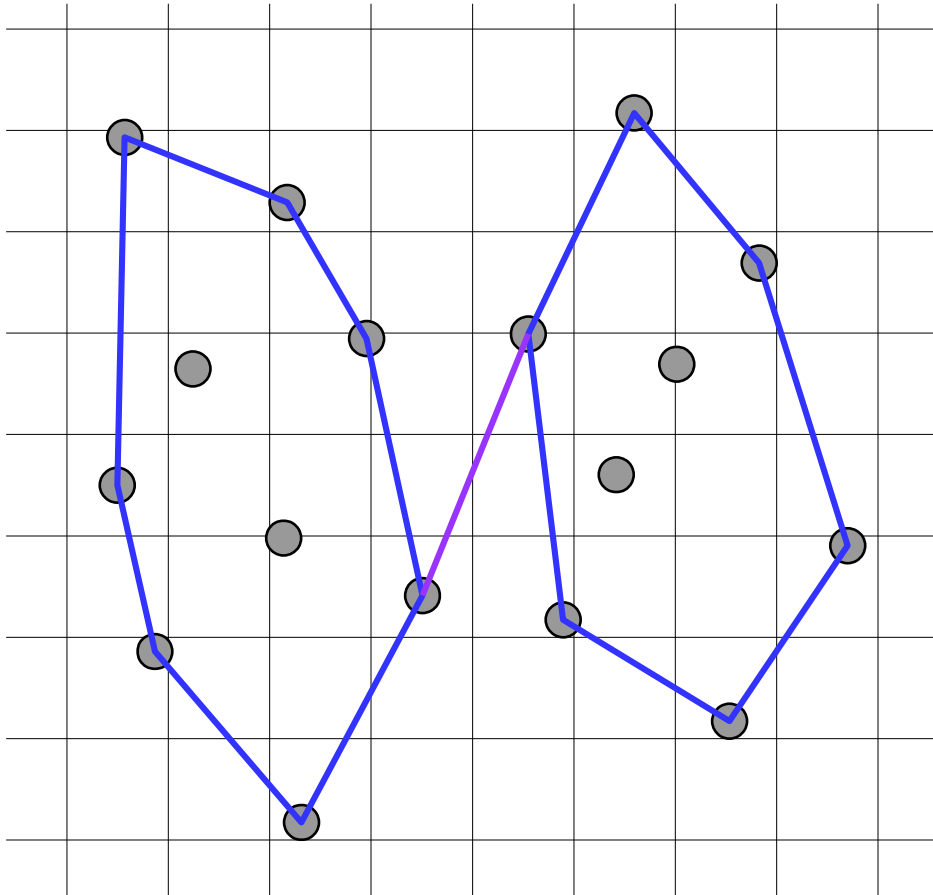


How to merge the **2** hulls?

... note, they always don't intersect.



# Divide and conquer



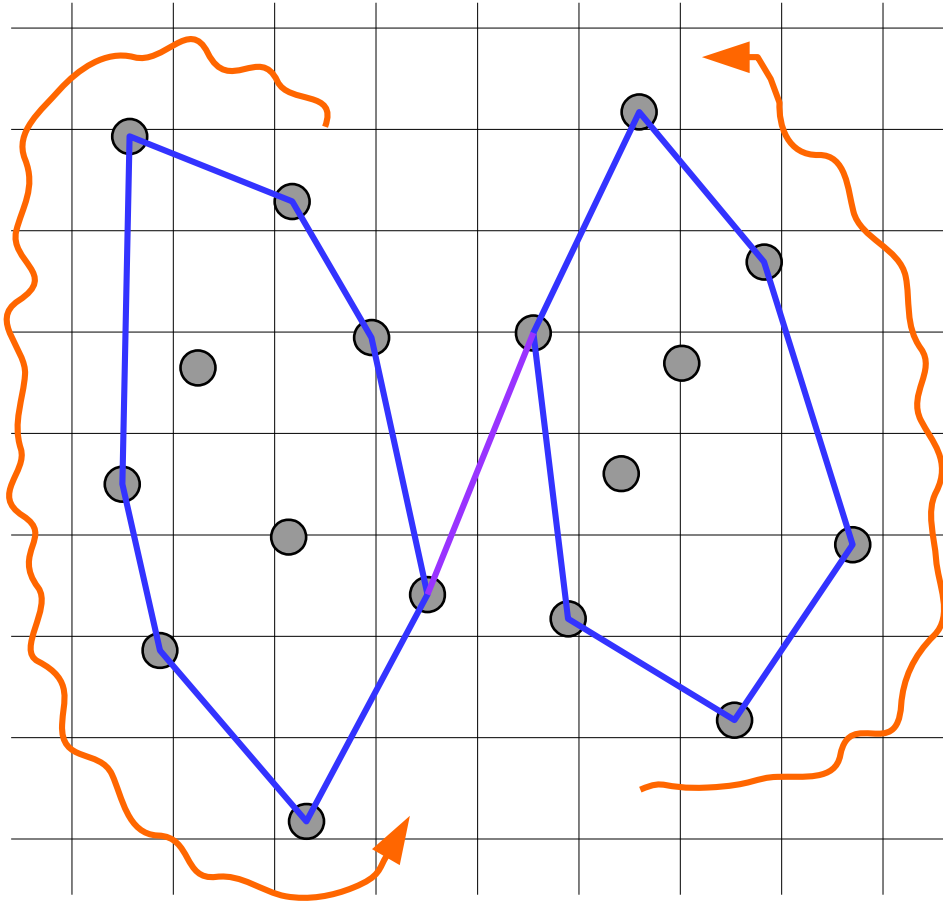
How to merge the **2** hulls?

... note, they always don't intersect.

Let's connect them:

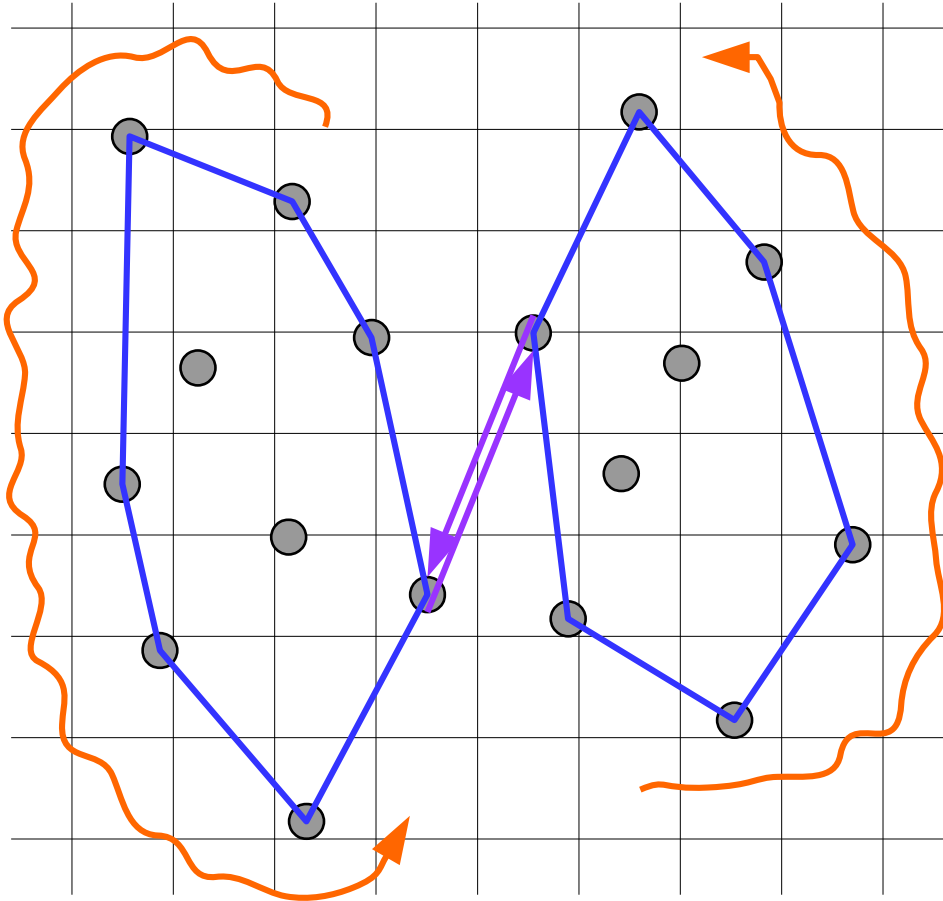
- > rightmost point of left hull,
- > leftmost point of right hull.

# Divide and conquer



Remember, that every hull can be interpreted also as cycling around all the input points.

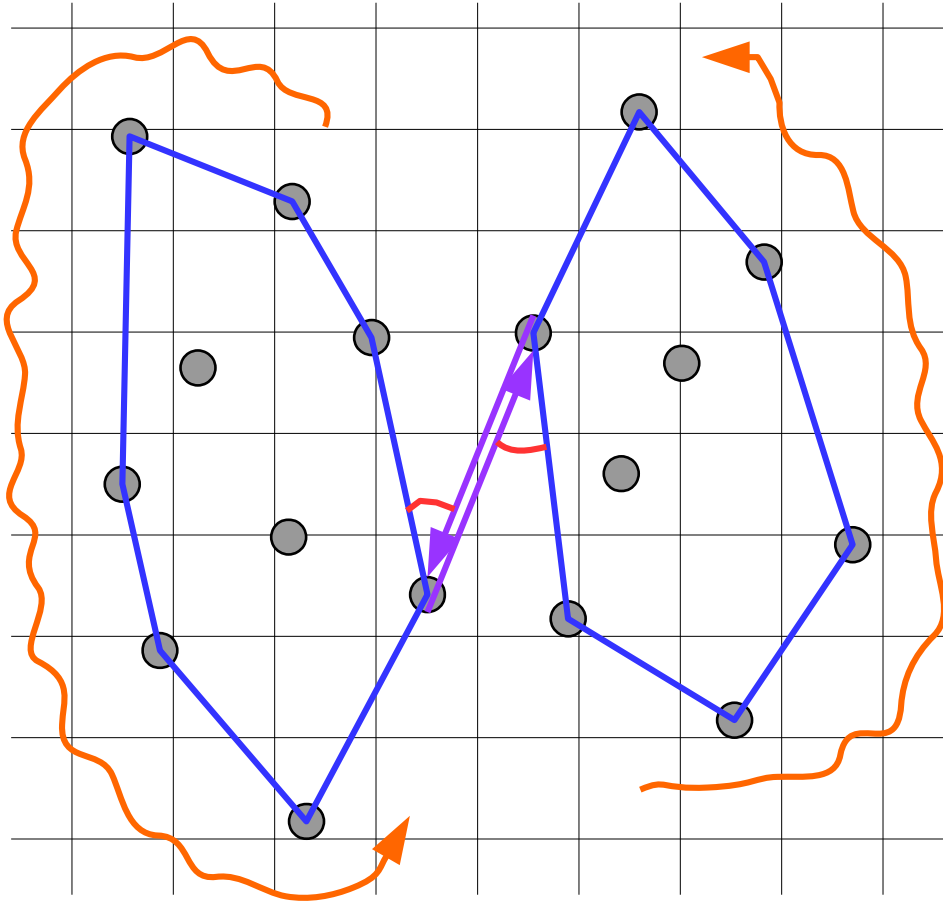
# Divide and conquer



Remember, that every hull can be interpreted also as cycling around all the input points.

Now, adding the "bridge" can be interpreted this way:

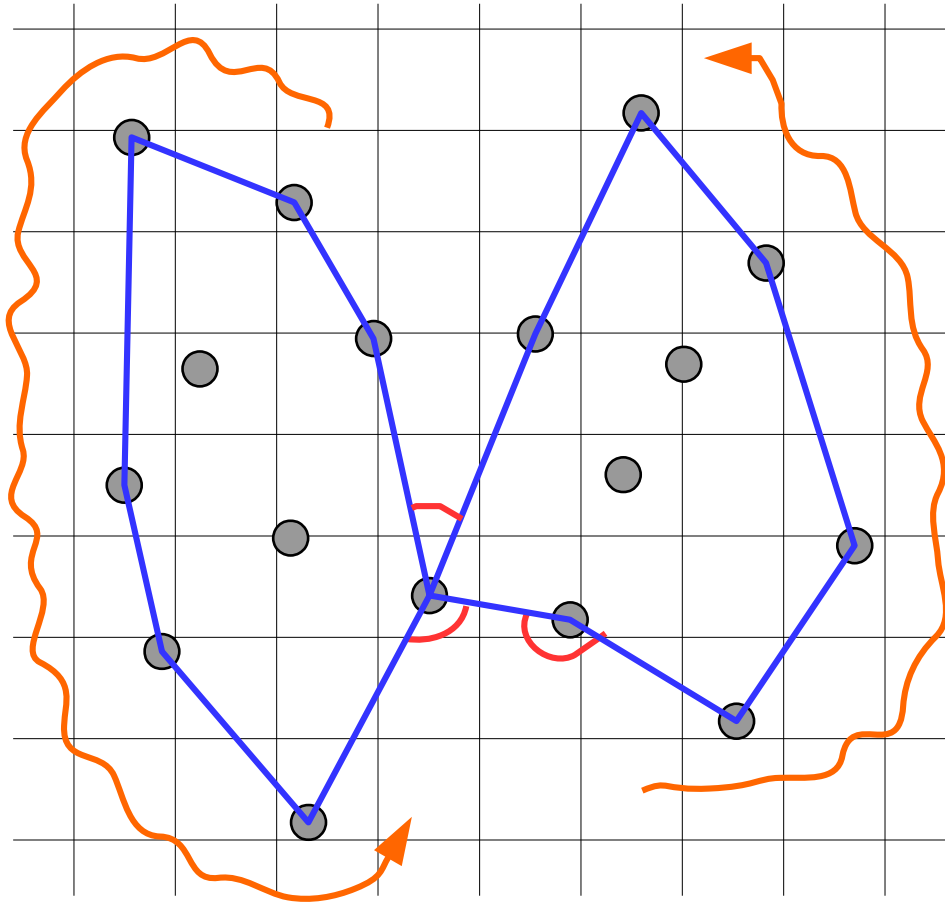
# Divide and conquer



Having this, it will remain only to exclude "right turns",

... the way we were doing in the previous algorithm.

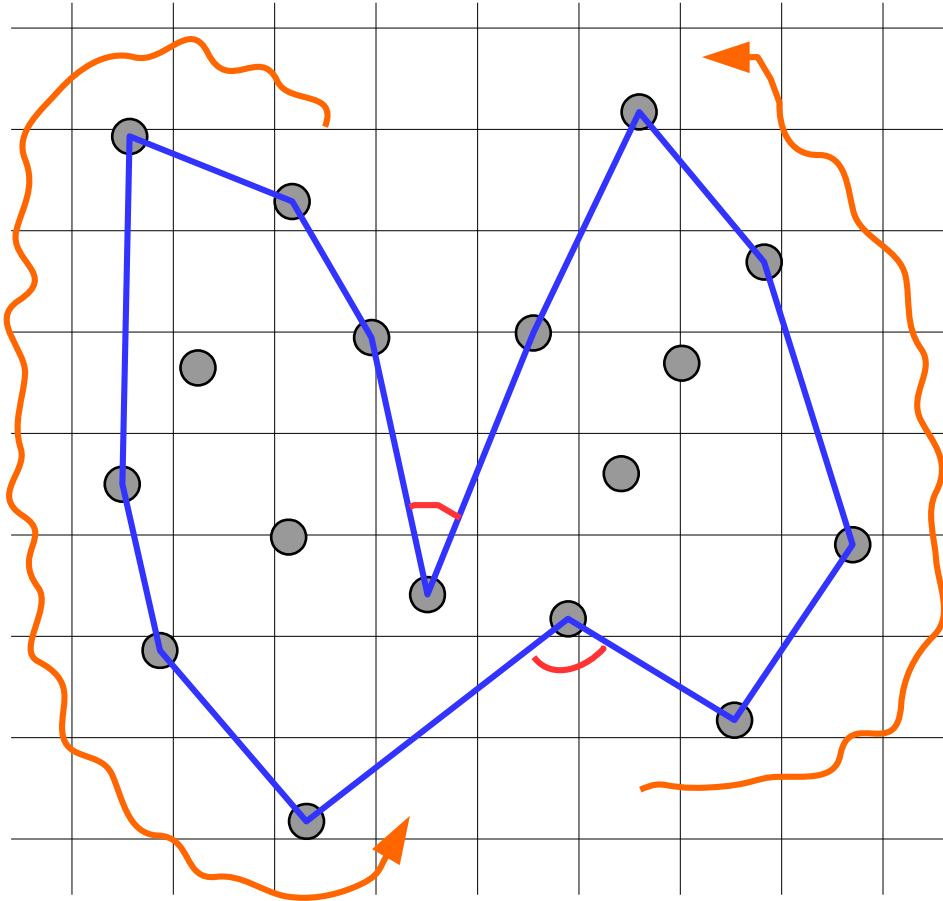
# Divide and conquer



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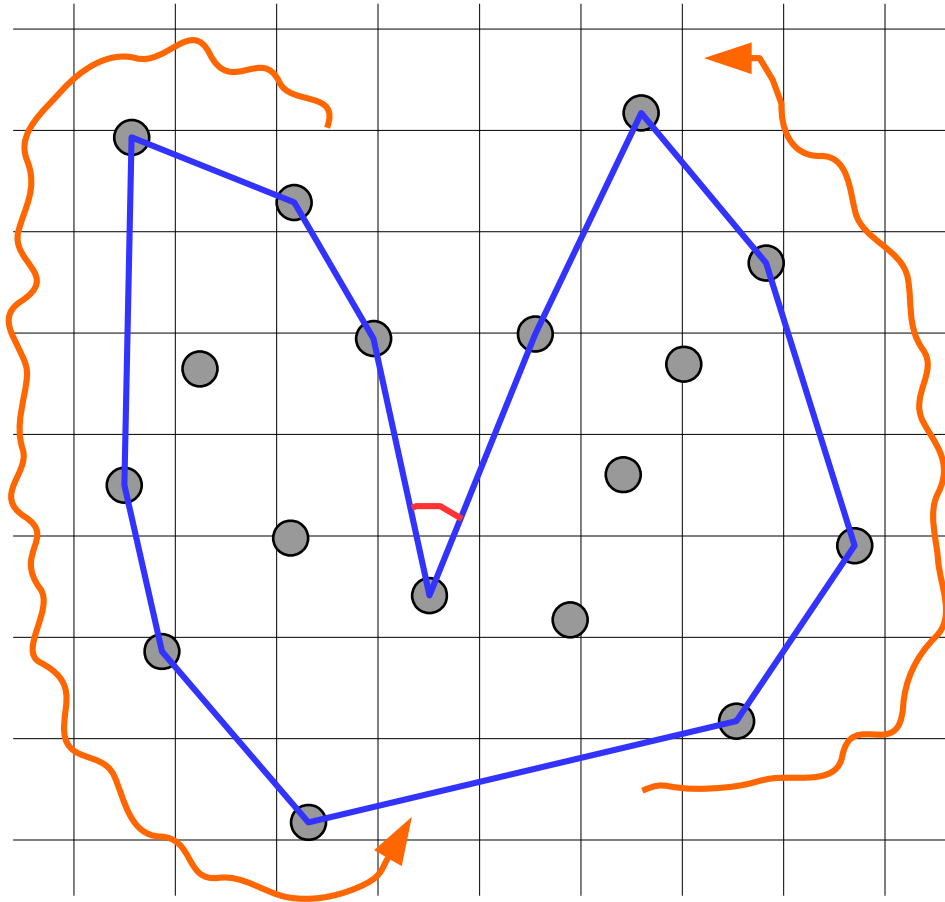
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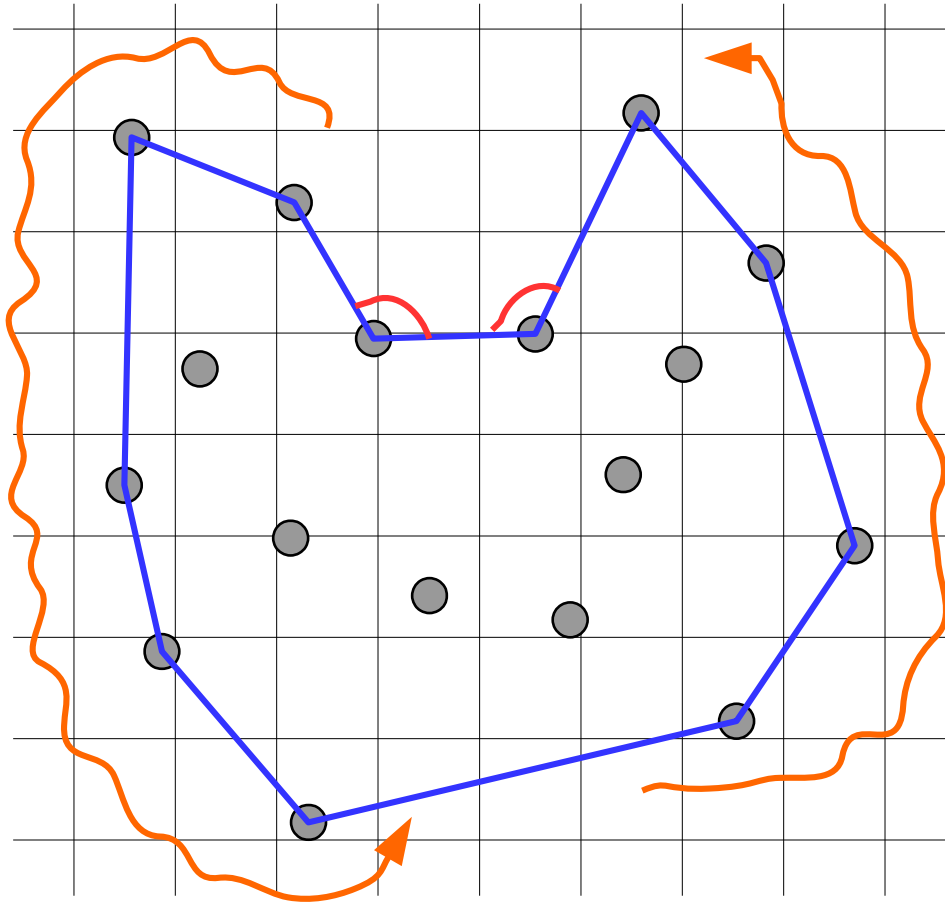
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# Divide and conquer

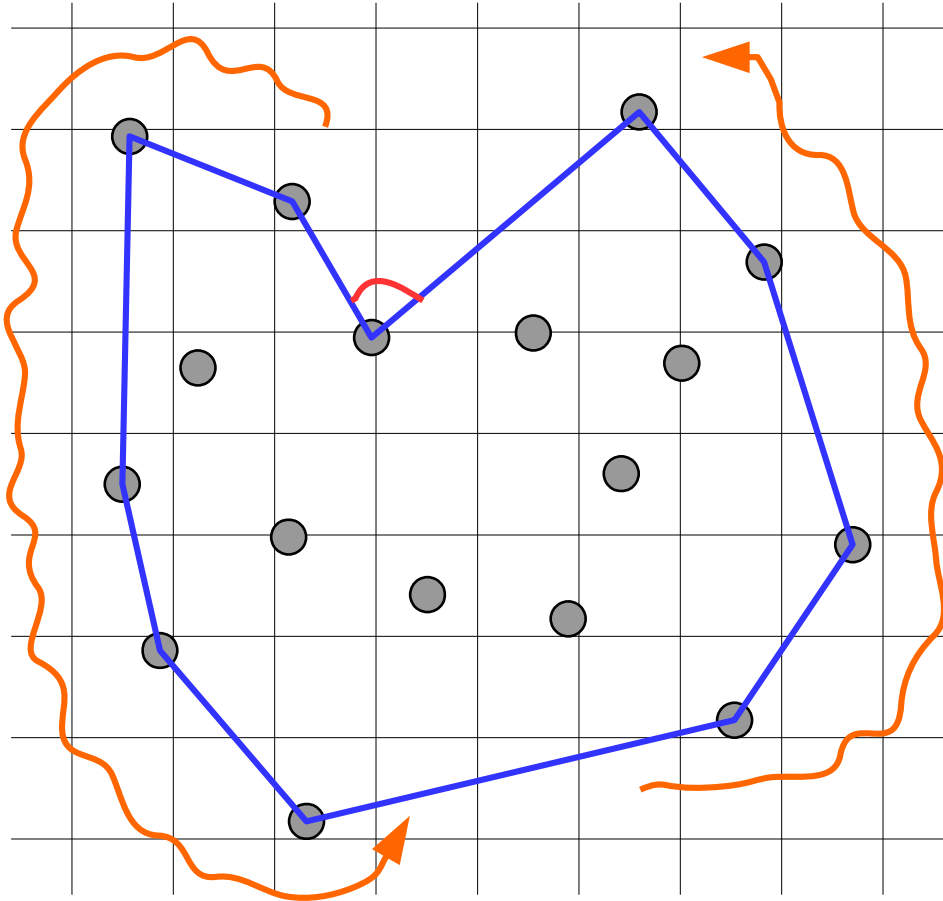


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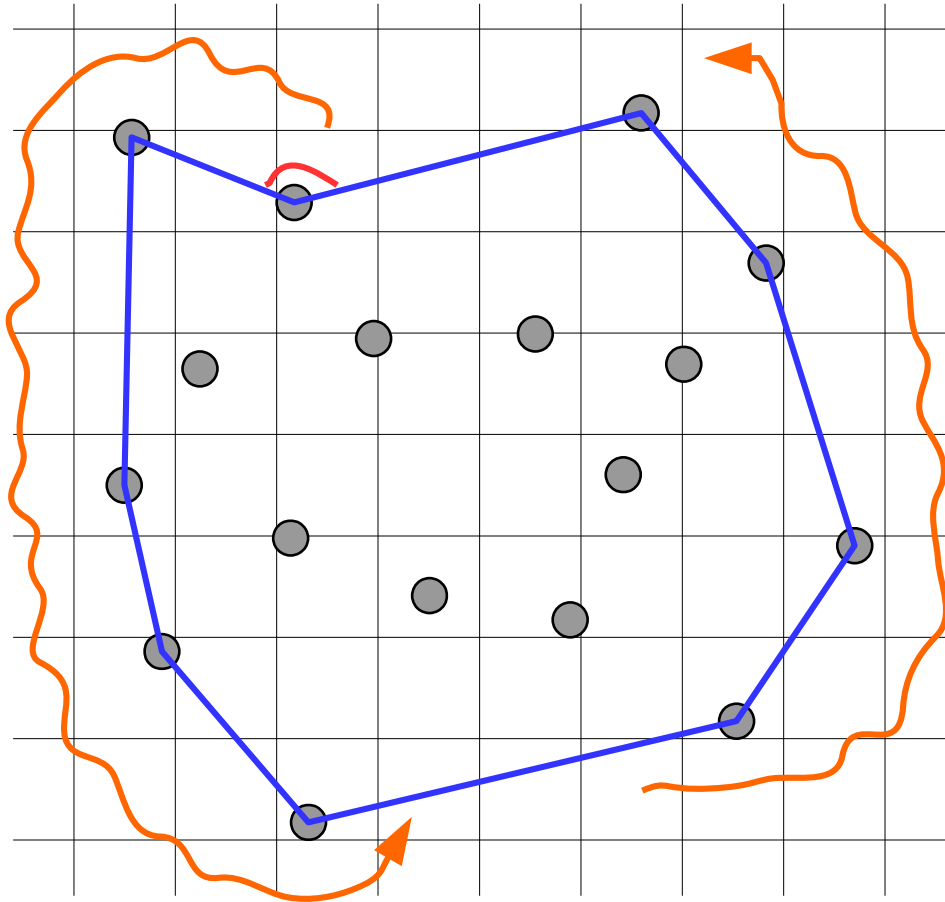
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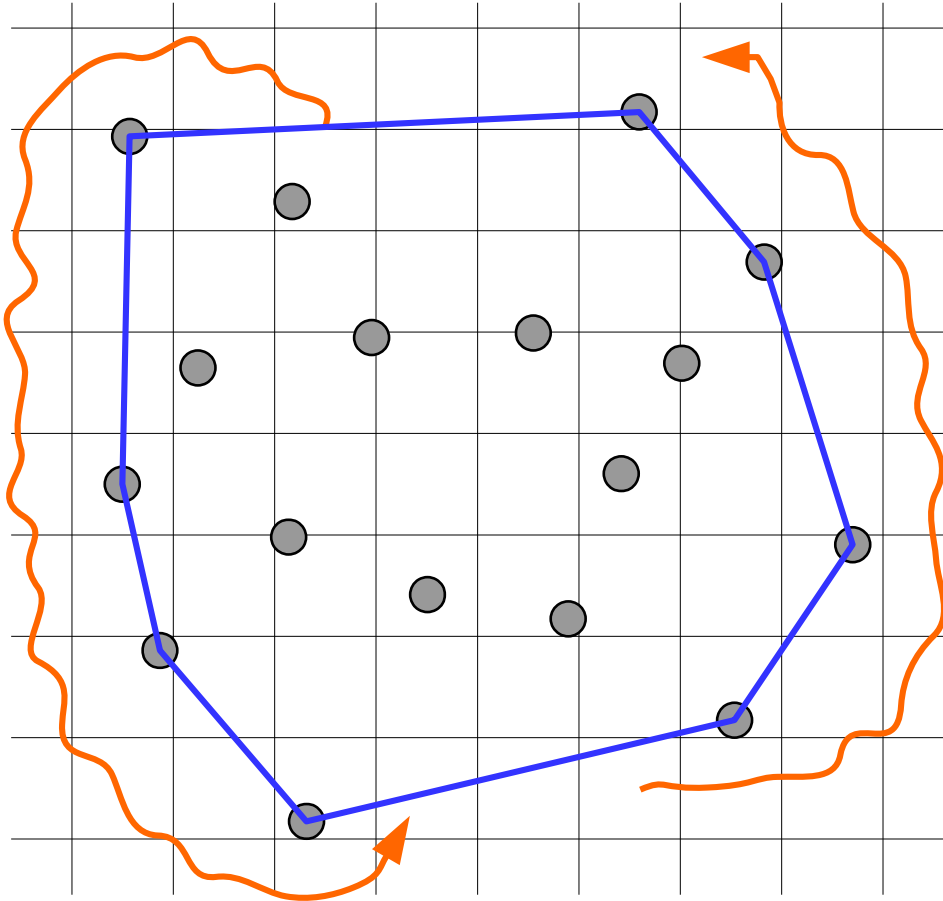
# Divide and conquer



Having this, it will remain only to  
exclude "right turns",

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# Divide and conquer



Having this, it will remain only to  
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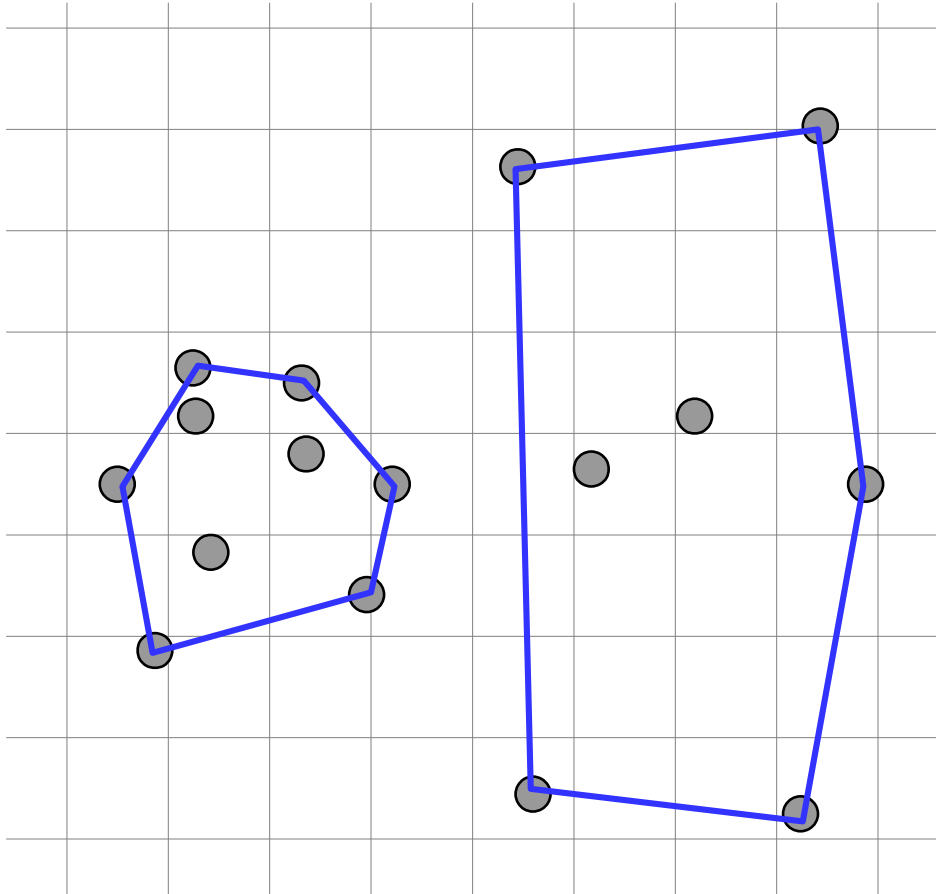
# Time complexity

Calculating time complexity:

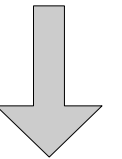
$$\begin{aligned} f(n) &= f(n/2) + f(n/2) + n = 2*f(n/2) + n = \\ &= 2*( 2*f(n/4)+n/2 ) + n = 4*f(n/4) + 2n = \\ &= 4*( 2*f(n/8)+n/4 ) + 2n = 8*f(n/8) + 3n = \\ &= \dots = \\ &= n * \log n \end{aligned}$$

So here the logarithm originates not from sorting, but from partitioning.

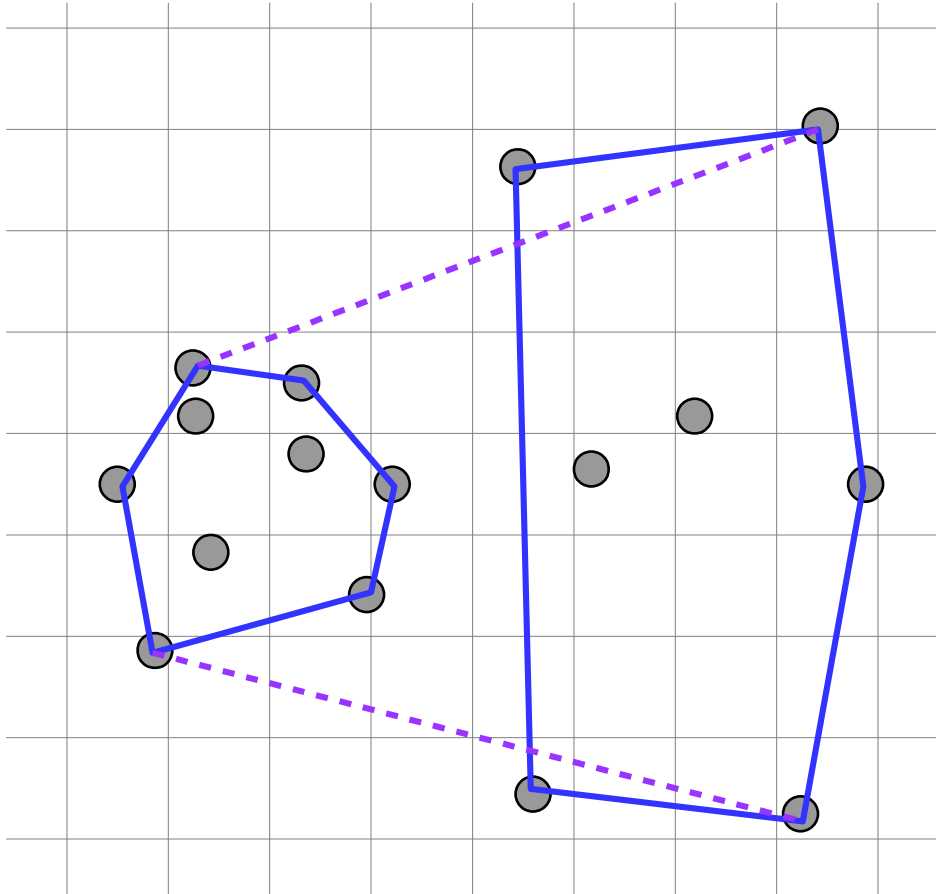
# Divide and conquer



Note, in order to merge the **2** convex hulls, we can't just take the bottom-most pair and the top-most pair.

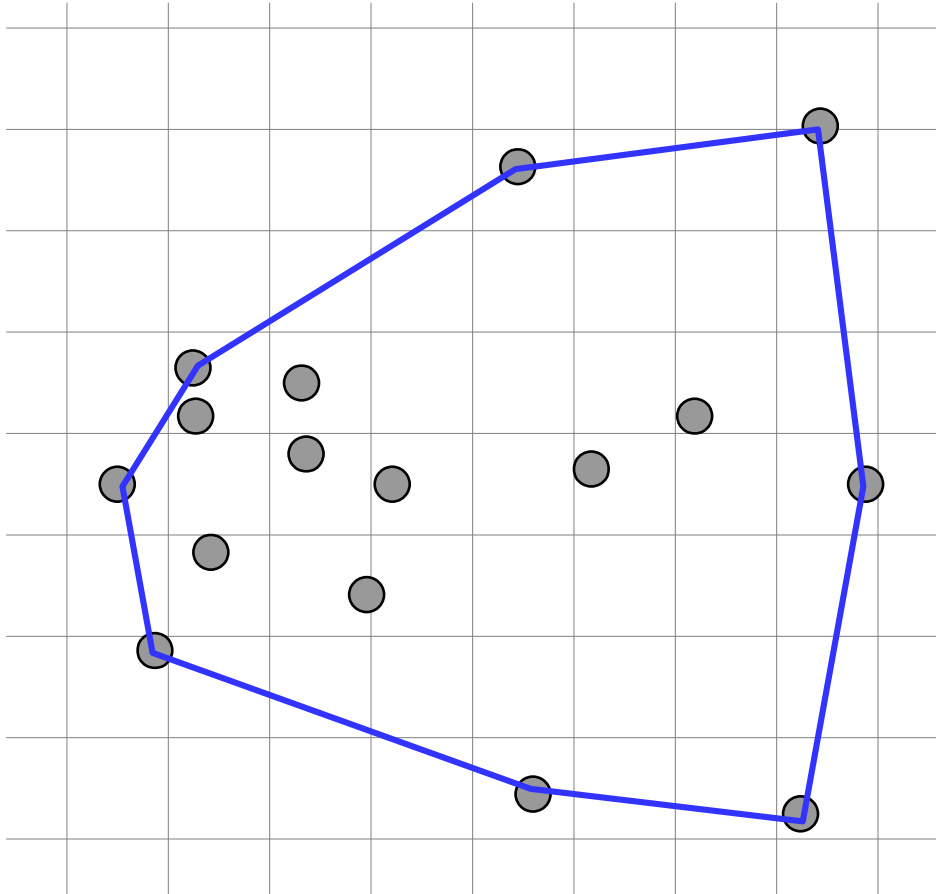


# Divide and conquer



Note, in order to merge the **2** convex hulls, we can't just take the bottom-most pair and the top-most pair.

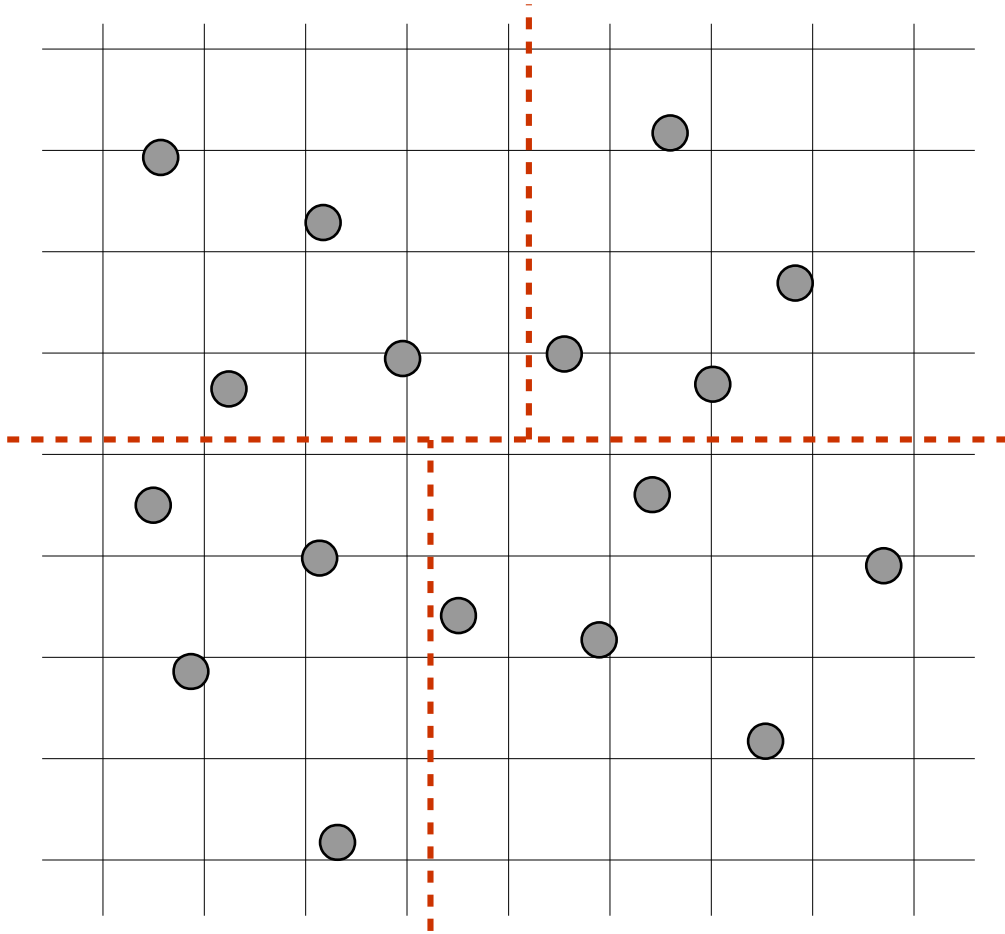
# Divide and conquer



Note, in order to merge the **2** convex hulls, we can't just take the bottom-most pair and the top-most pair.

... actual result of merging is this.

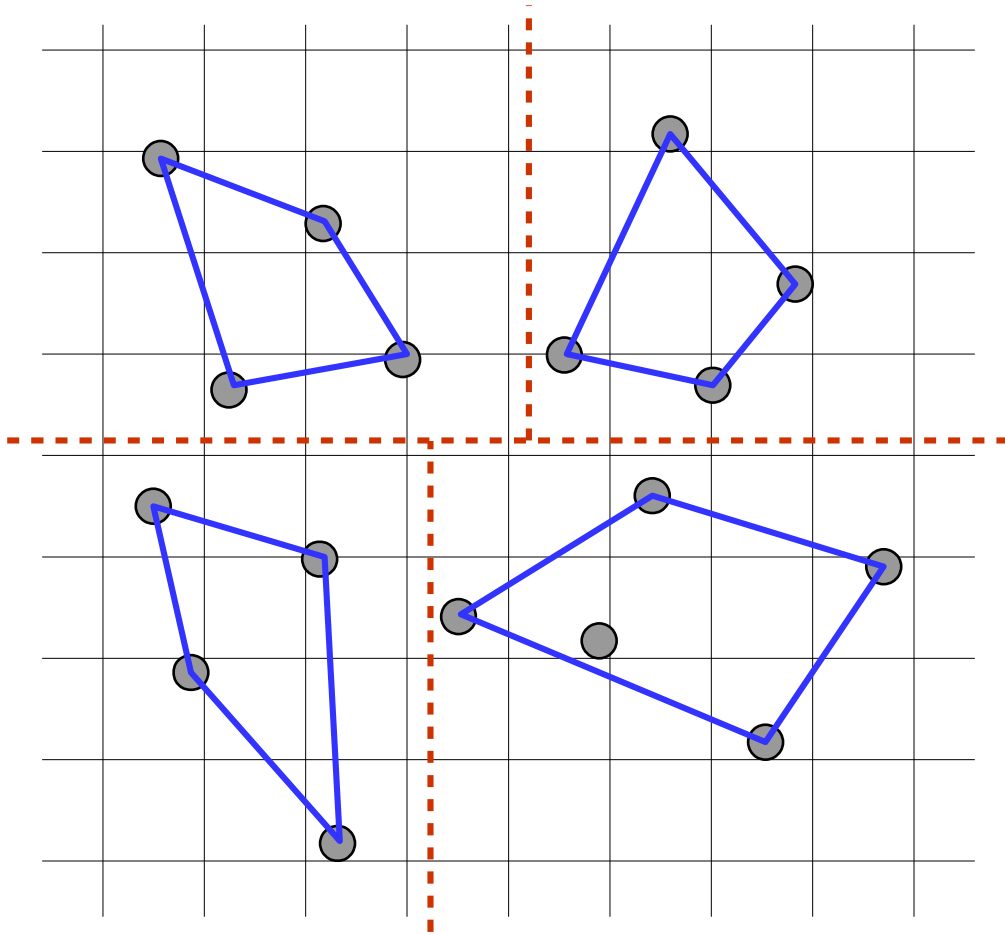
# Divide and conquer



***Question:*** Can we interleave directions of the split: once horizontally, then vertically, and so on...



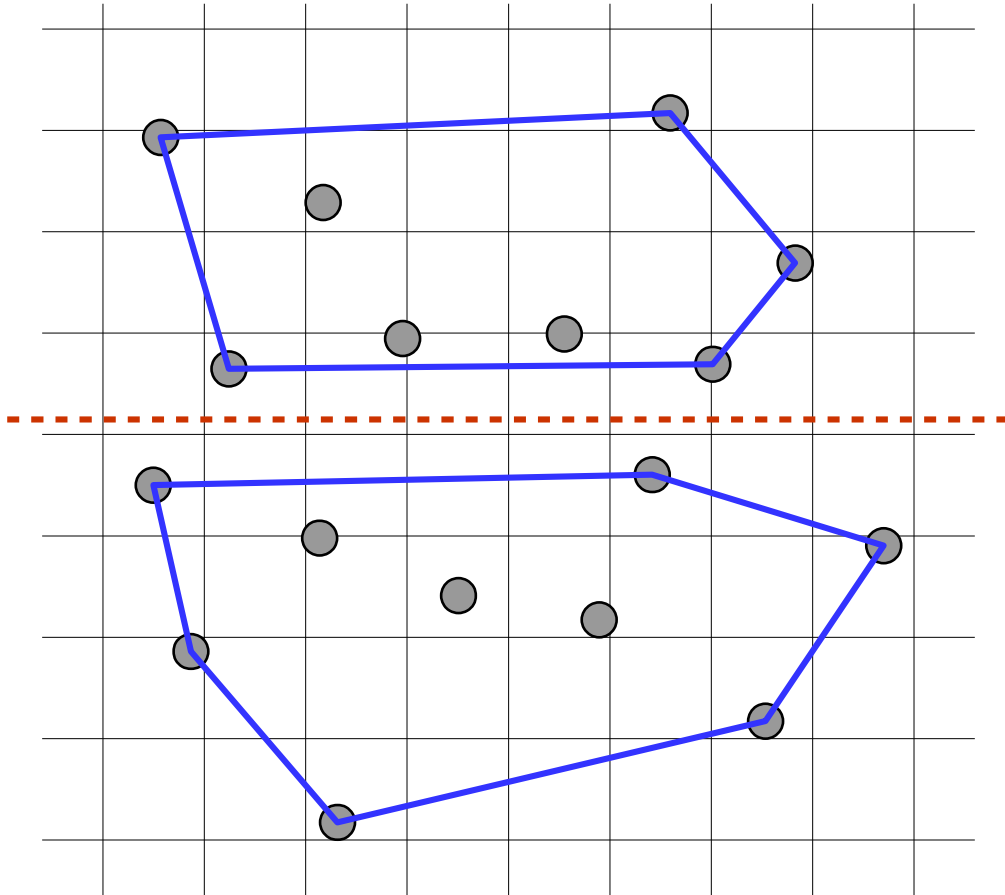
# Divide and conquer



**Question:** Can we interleave directions of the split: once horizontally, then vertically, and so on...

**Answer.** Yes.

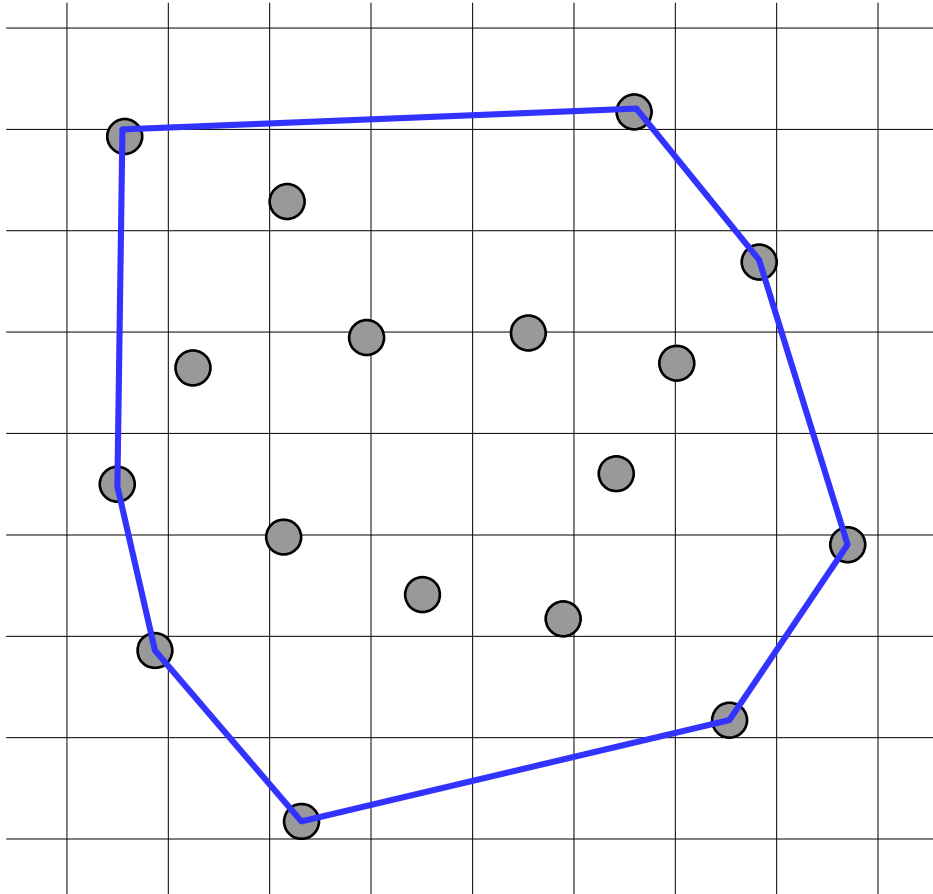
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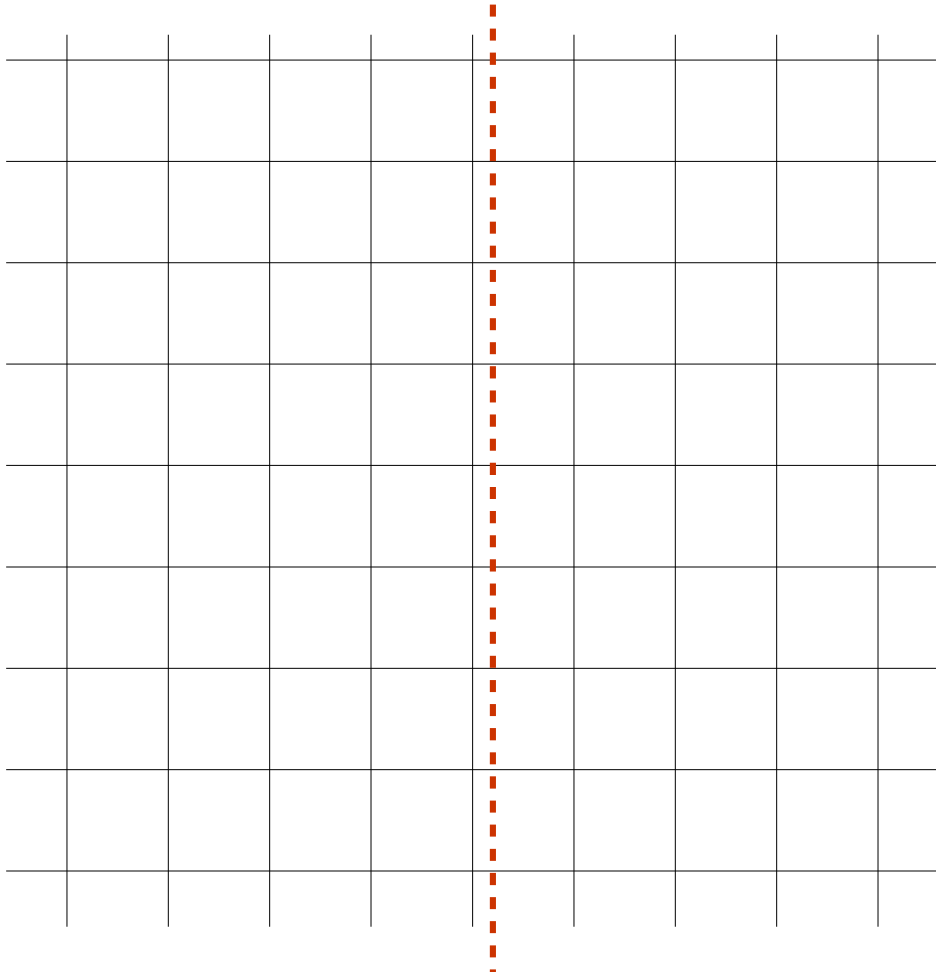
# Divide and conquer



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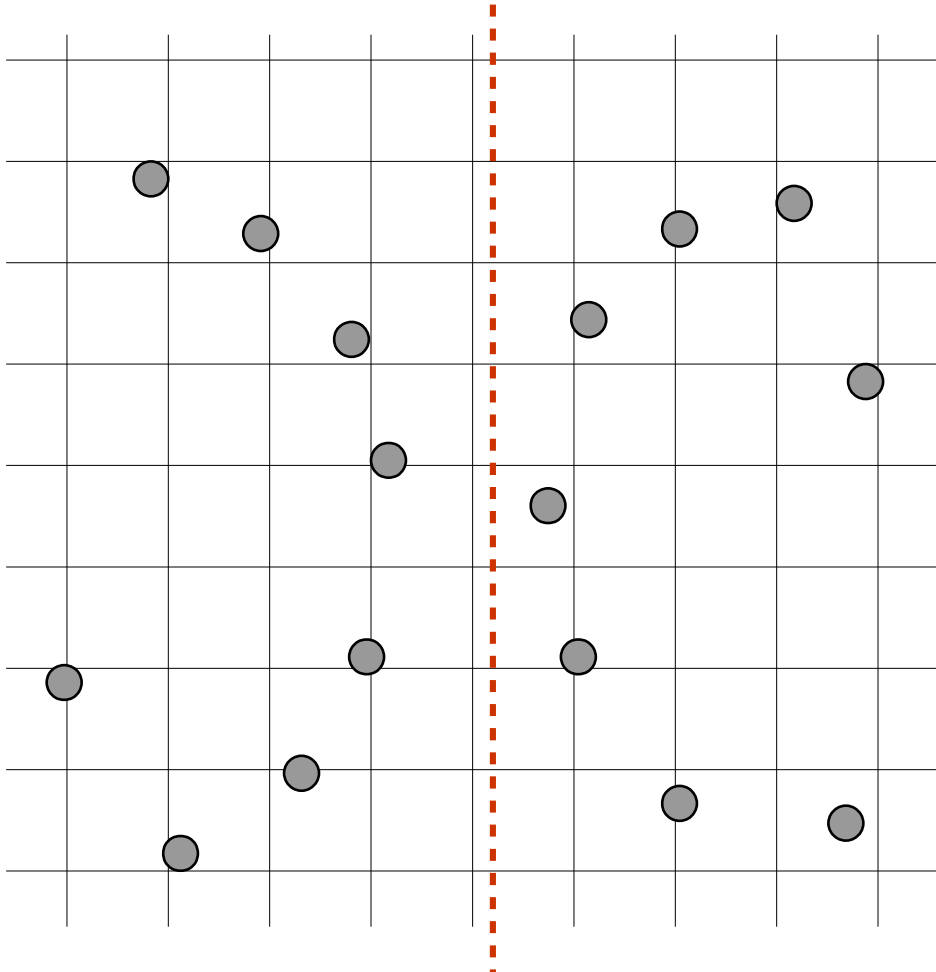
# Divide and conquer



**Question:** Describe a case, when:

- final '**h**' is small,
- much time is spent on the final merge.

# Divide and conquer

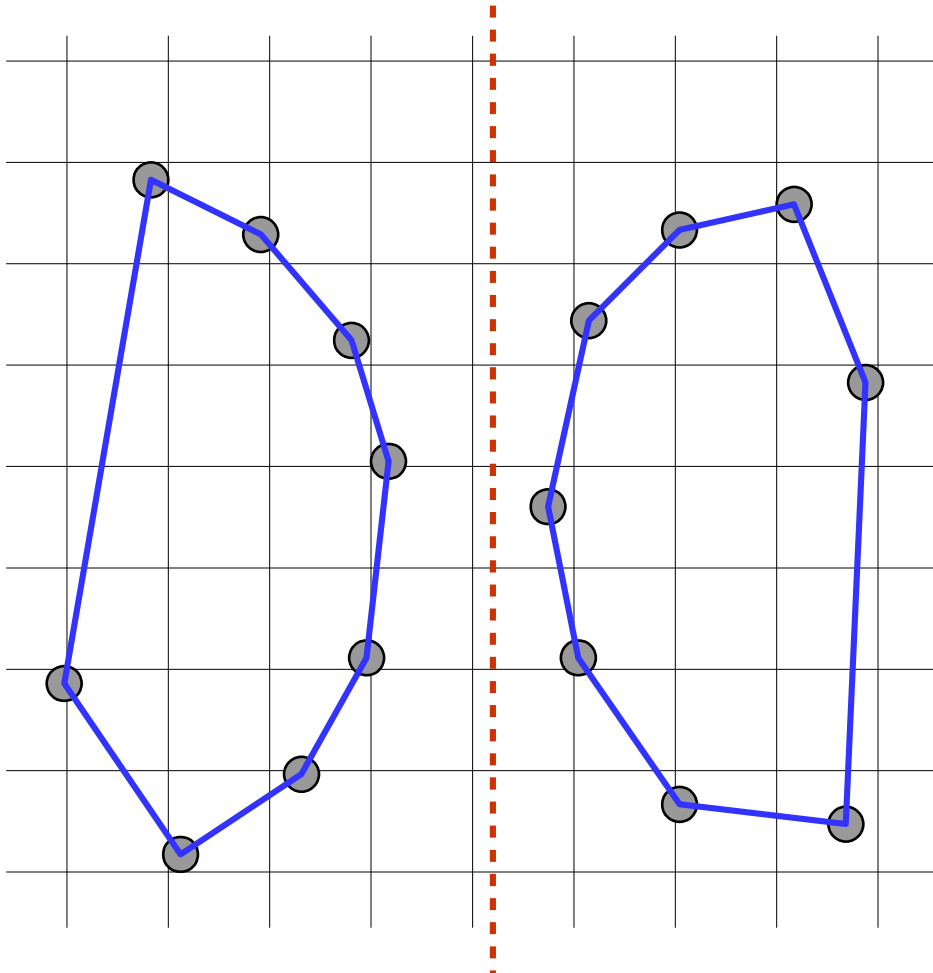


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**Answer.** <---

# Divide and conquer



**Question:** Describe a case, when:

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- much time is spent on the final merge.

**Answer.** <---

# Exercise

Draw on paper  **$n=16$**  points, and perform first **2** levels of the algorithm.

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# Thank you!

Convex hull

*(divide and conquer algorithm)*