

# Implementation of Interpolation Algorithms for Software Verification

by

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B.Tech., Universidad de las Américas Puebla, 2015

THESIS

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# Dedication

*To my family.*

# Acknowledgments

TODO.

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## Abstract

This thesis discusses theoretical aspects and an implementation for the interpolation problem of the following theories: (quantifier-free) equality with uninterpreted functions (EUF), unit two variable per inequality (UTVPI), and the combination of the two previous theories. The interpolation algorithms implemented in this thesis are originally proposed in [25].

Refutational proof-based solutions are the usual approach of many interpolation algorithms [20, 33, 32]. The approach taken in [25] relies on quantifier-elimination heuristics to construct an interpolant using one of the two formulas involved in the interpolation problem. The latter makes possible to study the complexity of the algorithms obtained. On the other hand, the combination method implemented in this thesis uses a Nelson-Oppen framework, thus we still require for this particular situation a refutational proof in order to guide the construction of the interpolant for the combined theory.

The implementation uses Z3 [14] for parsing purposes and satisfiability checking in the combination component of the thesis. Minor modifications were applied to

the Z3's enode data structure in order to label and distinguish formulas efficiently (i.e. distinguish A-part, B-part). Thus, the project can easily be integrated to the Z3 solver to extend its functionality for verification purposes using the Z3 plug-in module.

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# Chapter 1

## Introduction

Modern society is witness of the impact computer software has done in recent years. The benefits of this massive automation is endless. On the other hand, when software fails, it becomes a catastrophe ranging from economic loss to threats for human life.

Due to strict and ambitious agendas, many software products are shipped with unseen and unintentional bugs which might potentially put at risk people's life like critical systems. Several approaches have been used to improve software quality. However, many of these approaches offer partial coverage or might take an abyssal amount of human effort to provide such solutions. Thus, these approaches cannot be considered practical. On the other hand, for certain applications these contributions are relevant and their proper use fit such workflows uniquely.

Formal methods aim to bring a unique combination of automation, rigor, and efficiency (whenever efficient algorithms exist for the verification task). This thesis discusses a particular problem in software verification known as *the interpolation problem* for the theories of the quantifier-free fragment of equality with uninterpreted functions (EUF) and unit two variable per inequality (UTVPI) and their combination. These two theories have been studied extensively and researchers have found

several applications for them.

## 1.1 Background

An interpolant of a pair of logical formulas  $(\alpha, \beta)$  is a logical formula  $\gamma$  such that  $\alpha$  implies  $\gamma$ ,  $\beta \wedge \gamma$  is logically inconsistent, and  $\gamma$  only has common symbols of  $\alpha$  and  $\beta$ . Informally this means that the interpolant ‘belongs’ to the consequence of  $\alpha$ , and ‘avoids’ being part of the consequence of  $\beta$ . Not surprisingly, this intuition is behind many software verification routines where the first formula models a desirable state/property (termination, correctness) of a computer program, and the second formula models the set of undesirable states (non-termination, errors, crashes, etc) of such software. In Chapter 2, an extensive review of the formal concept is provided.

Eventhough interpolants are not a direct concern in verification problems, these problems are found in the core algorithms of the following two applications:

- Refinement of abstract models: In order to improve coverage and decrease the complexity in verification problems, abstract interpretation has become a proper technique to accomplish the latter together with model checking. Eventhough the methodology provides sound results, it is certainly not complete. Additionally, several abstractions do not capture the semantical meaning of programs due to the *over-approximation* approach. Hence, interpolants are used to strength predicate abstractions by using interpolants constructed from valid traces in the abstract model but not valid in the actual model (*spurious counterexamples*) [10, 31, 23].
- Invariant generation: following the same idea as in the previous case, *if a fix point is obtained in the refinement process*, we can obtain a logical invariant

of computer programs <sup>1</sup> Situations where this happens might be due to the finiteness of possible states in the program. It is worth mentioning that the invariant problem as stated in [24] is undecidable [42, 3].

## 1.2 Related work

There are several interpolation algorithms for the theories involved in the thesis work. The approaches can be classified into the following categories:

- **Proof-based approach:** This category relies on the availability of a refutational proof. The interpolant is constructed using a recursive function over the structure of the proof tree. In [32, 33] the author defines an interpolation calculus. This particular approach uses a proof tree produced by the SMT solver Z3 and does not need to modify any of Z3's internal mechanisms. Among the advantages of this approach is that theory combination is given for free since the SMT solver takes care of this problem. On the other hand, Z3's satellite theories are not sufficiently integrated with its proof-producing mechanism. Hence, one can find  *$\mathcal{T}$ -lemmas* as black-boxes which introduces incompleteness in the interpolation calculus. For completeness, these lemmas are solved separately by another interpolation algorithm for the respective theory. In [46] the authors provide a Nelson-Oppen framework to compute interpolants. For the convex case, the approach only exchange equalities as required by the Nelson-Oppen framework. For the non-convex case, the authors require a resolution-based refutational proof to compute interpolants using Pudlak's algorithm. The introduction of the class of *equality interpolating theories* is considered among

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<sup>1</sup>The interpolation generation approach discussed can be understood as a *lazy framework* similarly to SAT/SMT algorithms. The former is about the production of interpolants, the latter is for assignments/models respectively. Both *block/learn* the formulas in order to find their results.

the most relevant contributions of the paper by the verification community. This is property about theories which states that if a theory is capable of proving a mixed equality  $a = b$  (an equality which contains symbols from the two formulas in the interpolating problem) then it exists a common term in the language of the theory  $t$  (known as the interpolating-term) such that the theory can prove  $a = t$  and  $t = b$ . The property facilitates formula-splitting for interpolation purposes. In [20] the authors modify a resolution-based refutational proof by introducing common-terms in the proof in order to produce interpolants in what is called colorable-proofs, which are proof trees which do not contained AB-mixed literals. This is pointed as an improvement to the approach followed in [46] which executes a similar idea but done progressively as the proof-tree is built and does not require the theories solver to be *equality propagating*<sup>2</sup>. However, this results is not generalizable to non-convex theories due to internal constraints.

- Reduction-based approach: This method transforms the interpolation problem into a query for some solver related to the theory. An example of this approach can be found in [43] where the authors use a linear-inequality solver to provide an interpolant for the theory of linear inequalities over the rational/real numbers ( $LIA(\mathbb{Q})/LIA(\mathbb{R})$ ). Additionally, they integrate the procedure with a *hierarchical reasoning* approach in order to incorporate the signature of theory for (quantifier-free) equalities with uninterpreted functions.

### 1.3 Outline of the thesis

- Chapter 2 provides an extensive background of fundamentals ideas, definitions and decision procedures used in the thesis work.

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<sup>2</sup>The authors in [46] require that the theory solvers keep track of the interpolating-term and propagate this term whenever possible

- Chapters 3, 4, and 5 explain implementation details of the interpolating algorithms for the theories EUF, UTVPI, and their combination respectively. These chapters share the same structure. They start with the algorithms used to solve the interpolation problem, discuss about implementation details including diagrams of the architecture of the implemented system, and show a performance comparison with the iZ3 interpolation tool available in the SMT solver Z3 until version 4.7.0.

## 1.4 Contributions

The contributions of the thesis can be summarized as follow:

1. Implementation of the interpolation algorithm for the theory EUF proposed in [25].
2. Formulation and implementation of a new procedure for checking unsatisfiability of grounded equations in Horn clauses using a congruence closure algorithm with explanations used in the implementation of item 1.
3. Implementation of the interpolation algorithm for the theory UTVPI proposed in [25].
4. Implementation of the combination procedure for interpolating algorithm proposed in [46] in order to combine the implementations of items 1. and item 3.



# Chapter 2

## Preliminaries

This chapter discusses basic concepts from first-order logic that are used in the rest of this thesis. We will pay particular attention to their language and semantics since these are fundamental concepts necessary to understand the algorithmic constructions to compute interpolants. For a comprehensive treatment on the topic, the reader is suggested the following references [34, 19].

### 2.1 First-Order Predicate Logic

#### 2.1.1 Language

A language is a collection of symbols of different sorts equipped with a rule of composition that effectively tells us how to recognize elements that belong to the language [44]. In particular, a first-order language is a language that expresses boolean combinations of predicates using terms (constant symbols and function applications). In mathematical terms,

**Definition 2.1.1.** *A first-order language (also denoted signature) is a triple  $\langle \mathfrak{C}, \mathfrak{P}, \mathfrak{F} \rangle$*

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of non-logical symbols where:

- $\mathfrak{C}$  is a collection of constant symbols
- $\mathfrak{P}$  is a collection of  $n$ -place predicate symbols
- $\mathfrak{F}$  is a collection of  $n$ -place function symbols

including logical symbols like quantifiers (universal ( $\forall$ ), existential ( $\exists$ )) logical symbols like parenthesis, propositional connectives (implication ( $\rightarrow$ ), conjunction ( $\wedge$ ), disjunction ( $\vee$ ), negation ( $\neg$ )), and a countably number of variables  $\text{Vars}$  (i.e.  $\text{Vars} = \{v_1, v_2, \dots\}$ ).

The rules of composition distinguish two objects, terms and formulas, which are defined recursively as follows:

- Any variable symbol or constant symbol is a term.
- If  $t_1, \dots, t_n$  are terms and  $f$  is an  $n$ -ary function symbol, then  $f(t_1, \dots, t_n)$  is also a term.
- If  $t_1, \dots, t_n$  are terms and  $P$  is an  $n$ -ary predicate symbol, then  $P(t_1, \dots, t_n)$  is formula.
- If  $x$  is a variable and  $\psi, \varphi$  are formulas, then  $\neg\psi, \forall x.\psi, \exists x.\psi$  and  $\psi \square \varphi$  are formulas where  $\square \in \{\rightarrow, \wedge, \vee\}$

No other expression in the language can be considered terms nor formulas if such expressions are not obtained by the previous rules.

### 2.1.2 Semantics

In order to define a notion of truth in a first-order language is necessary to associate for each non-logical symbol (since logical symbols have established semantics from propositional logic) a denotation or mathematical object and an *assignment* to the collection of variables. The two previous components are part of an *structure* [19] (or interpretation [44]) for a first-order language.

**Definition 2.1.2.** *Given a first-order language  $\mathfrak{L}$ , an interpretation  $\mathfrak{I}$  is a pair  $(\mathfrak{A}, \mathfrak{J})$ , where  $\mathfrak{A}$  is a non-empty domain (set of elements) and  $\mathfrak{J}$  is a map that associates to the non-logical symbols from  $\mathfrak{L}$  the following elements:*

- $c^{\mathfrak{J}} \in \mathfrak{A}$  for each  $c \in \mathfrak{C}$
- $f^{\mathfrak{J}} \in \{\mathfrak{A}^n \rightarrow \mathfrak{A}\}$  for each  $n$ -ary function symbol  $f \in \mathfrak{F}$
- $P^{\mathfrak{J}} \in \{\mathfrak{A}^n\}$  for each  $n$ -ary predicate symbol  $P \in \mathfrak{P}$

An assignment  $s : \text{Vars} \rightarrow \mathfrak{A}$  is a map between Vars to elements from the domain of the interpretation.

With the definition of interpretation  $\mathfrak{I}$  and assignment  $s$ , we can recursively define a notion of *satisfiability* (denoted by the symbol  $\models_{\mathfrak{I}, s}$ ) as a free extension from atomic predicates (function application of predicates) to general formulas as described in [19]. For the latter, we need to extend the assignment function to all terms in the language.

**Definition 2.1.3.** *Let  $\mathfrak{I} = (\mathfrak{A}, \mathfrak{J})$  be an interpretation and  $s$  an assignment for a given language, Let  $\bar{s} : \text{Terms} \rightarrow \mathfrak{A}$  be defined recursively as follows:*

- $\bar{s}(c) = c^{\mathfrak{J}}$

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- $\bar{s}(f(t_1, \dots, t_n)) = f^{\mathfrak{J}}(\bar{s}(t_1), \dots, \bar{s}(t_n))$

Notice that the extension of  $s$  depends on the interpretation used.

**Definition 2.1.4.** *Given an interpretation  $\mathfrak{J} = (\mathfrak{A}, \mathfrak{J})$ , an assignment  $s$ , and  $\psi$  a formula, we define  $\mathfrak{J} \models_s \psi$  (read  $\psi$  is satisfiable under interpretation  $\mathfrak{J}$  and assignment  $s$ ) recursively as follows:*

- $\models_{\mathfrak{J},s} P(t_1, \dots, t_n)$  if and only if  $P^{\mathfrak{J}}(\bar{s}(t_1), \dots, \bar{s}(t_n))$
- $\models_{\mathfrak{J},s} \neg\psi$  if and only if it is not the case that  $\models_{\mathfrak{J},s} \psi$
- $\models_{\mathfrak{J},s} \psi \wedge \varphi$  if and only if  $\models_{\mathfrak{J},s} \psi$  and  $\models_{\mathfrak{J},s} \varphi$
- $\models_{\mathfrak{J},s} \psi \vee \varphi$  if and only if  $\models_{\mathfrak{J},s} \psi$  or  $\models_{\mathfrak{J},s} \varphi$
- $\models_{\mathfrak{J},s} \psi \rightarrow \varphi$  if and only if  $\models_{\mathfrak{J},s} \neg\psi$  or  $\models_{\mathfrak{J},s} \varphi$
- $\models_{\mathfrak{J},s} \forall x.\psi$  if and only if for every  $d \in \mathfrak{A}$ ,  $\models_{\mathfrak{J},s_{x \mapsto d}} \psi$ , where  $s_{x \mapsto d} : \text{Vars} \rightarrow \mathfrak{A}$  is reduct of  $s$  under  $\text{Vars} \setminus \{x\}$  and  $s_{x \mapsto d}(x) = d$
- $\models_{\mathfrak{J},s} \exists x.\psi$  if and only if exists  $d \in \mathfrak{A}$ ,  $\models_{\mathfrak{J},s_{x \mapsto d}} \psi$ , where  $s_{x \mapsto d}$  is defined as in the previous item.

If an interpretation and assignment satisfies a formula, then we say that the interpretation and the assignment are a model for the respective formula. A collection of formulas are satisfied by an interpretation and assignment if these model each formula in the collection.

A formula  $\psi$  is said to be a valid formula of the interpretation  $\mathfrak{J}$  when  $\models_{\mathfrak{J},s} \psi$  for all possible assignments  $s$ .

Additionally, if all the models  $(\mathfrak{J}, s)$  in a language of a collection of formulas  $\Gamma$  satisfy a formula  $\psi$ , then we say that  $\Gamma$  logically implies  $\psi$  (written  $\Gamma \models \psi$ ). For the latter,  $\psi$  is said to be a valid formula of the model  $(\mathfrak{J}, s)$ .

## 2.2 Mathematical Theories

A theory  $\mathcal{T}$  is a collection of formulas that are closed under logical implication, i.e. if  $\mathcal{T} \models \psi$  then  $\psi \in \mathcal{T}$ . This concept is quite relevant for our thesis work since we will focus on two theories, the quantifier-free fragment of the theory of equality with uninterpreted functions (EUF), and the theory of unit two variable per inequality (UTVPI).

For some theories it is enough to provide a collection of formulas (known as the axioms of the theory). For the case of the theories of interest for the thesis, the axiomatization is the following:

### 2.2.1 Equality with uninterpreted functions

**Definition 2.2.1.** Let  $\mathfrak{L}_{EUF} = \{\{\}, \{=\}, \{f_1, \dots, f_n\}\}$  be the language of EUF. The axioms of the theory are:

- (Reflexivity)  $\forall x. x = x$
- (Symmetry)  $\forall x. \forall y. x = y \rightarrow y = x$
- (Transitivity)  $\forall x. \forall y. \forall z. (x = y \wedge y = z) \rightarrow x = z$
- (Congruence)  $\forall x_1 \dots \forall x_n. \forall y_1 \dots \forall y_n. (x_1 = y_1 \wedge \dots \wedge x_n = y_n) \rightarrow f(x_1, \dots, x_n) = f(y_1, \dots, y_n)$

We notice that the congruence axiom is not a first-order logic axiom, but rather an axiom-scheme since it is necessary to *instantiate* such axiom for every arity of the function symbols in a given language.

## 2.2.2 Ordered commutative rings

In order to describe the UTVPI theory we will first introduce the language and theory of an ordered commutative ring.

**Definition 2.2.2.** *Let  $\mathfrak{L}_{Ord-R} = \{, \{0, 1\}, \{=, \leq\}, \{+, -, *\}, \}$  be the language of an ordered commutative ring  $R$ . The axioms of the theory are:*

- $\forall x. \forall y. \forall z. x + (y + z) = (x + y) + z$
- $\forall x. \forall y. x + y = y + x$
- $\forall x. x + 0 = x$
- $\forall x. x + (-x) = 0$
- $\forall x. \forall y. \forall z. x * (y * z) = (x * y) * z$
- $\forall x. x * 1 = x$
- $\forall x. \forall y. x * y = y * x$
- $\forall x. \forall y. \forall z. x * (y + z) = x * y + x * z$
- $\forall x. \forall y. \forall z. (y + z) * x = y * x + z * x$
- $\forall x. \forall y. \forall z. x \leq y \rightarrow x + z \leq y + z$
- $\forall x. \forall y. (0 \leq x \wedge 0 \leq y) \rightarrow 0 \leq x * y.$
- $0 \neq 1 \wedge 0 \leq 1$

Section 2.4 discusses computability aspects for the theories of interest that are relevant for verification.

## 2.3 Interpolants

Following the notation in [46], we denote  $\mathcal{V}(\psi)$  to be the set of non-logical symbols, variables and constants of formula  $\psi$ . Given an instance for the interpolation problem  $(A, B)$ <sup>1</sup>, we distinguish the following categories:

- $\psi$  is *A-local* if  $\mathcal{V}(\psi) \in \mathcal{V}(A) \setminus \mathcal{V}(B)$
- $\psi$  is *B-local* if  $\mathcal{V}(\psi) \in \mathcal{V}(B) \setminus \mathcal{V}(A)$
- $\psi$  is *AB-common* if  $\mathcal{V}(\psi) \in \mathcal{V}(A) \cap \mathcal{V}(B)$
- $\psi$  is *AB-pure* when either  $\mathcal{V}(\psi) \subseteq \mathcal{V}(A)$  or  $\mathcal{V}(\psi) \subseteq \mathcal{V}(B)$ , otherwise  $\psi$  is *AB-mixed*

**Example 2.3.0.1.** Consider the following interpolation pair:  $(f(a+2)+1 = c+1 \wedge f(a+2) = 0, f(c) \leq b \wedge b < f(0))$ . With respect to the previous interpolation pair, we can tell that:

- The formula  $f(a+2) = c$  is *AB-pure* but not *A-local* nor *B-local* nor *AB-common*
- The formula  $\neg(a \leq f(f(b)+1))$  is an *AB-mixed literal*
- The formula  $a+1 = 1$  is *A-local*.
- The formula  $c+1 = 1$  is *AB-pure* but not *AB-common*.
- The formula  $c = 0$  is *AB-common*.
- In general, *AB-common* formulas are not *AB-pure* formulas.

---

<sup>1</sup>For the rest of the thesis, we will denote the first formula of an interpolation problem as the A-part and the second component as the B-part

### 2.3.1 Craig interpolation theorem

Let  $\alpha, \beta, \gamma$  be logical formulas in a given theory. If  $\models_{\mathcal{T}} \alpha \rightarrow \beta$ , we say that  $\gamma$  is an interpolant for the interpolation pair  $(\alpha, \beta)$  if the following conditions are met:

- $\models_{\mathcal{T}} \alpha \rightarrow \gamma$
- $\models_{\mathcal{T}} \gamma \rightarrow \beta$
- Every non-logical symbol in  $\gamma$  occurs both in  $\alpha$  and  $\beta$ .

The *interpolation problem* can be stated naturally as follows: given two logical formulas  $\alpha, \beta$  such that  $\models_{\mathcal{T}} \alpha \rightarrow \beta$ , find the interpolant for the pair  $(\alpha, \beta)$ .

In his celebrated result [12], Craig proved that for every pair  $(\alpha, \beta)$  of formulas in first-order logic such that  $\models \alpha \rightarrow \beta$ , an interpolation formula exists. Nonetheless, there are many logics and theories that this result does not hold [27].

Usually, we see the interpolation problem defined differently in the literature, where it is considered  $\beta'$  to be  $\neg\beta$  and the problem requires that the pair  $(\alpha, \beta')$  is mutually contradictory (unsatisfiable). This definition was popularized by McMillan [32]. This shift of attention explains partially the further development in interpolation generation algorithms since many of these relied on SMT solvers that provided refutation proofs in order to (re)construct interpolants for different theories (and their combination) [28, 8, 33].

Relaxed definitions are considered to the interpolation problem when dealing with specific theories [46] in a way that interpreted function can be also part of the interpolant. The latter is justified since otherwise, many interpolation formulas might not exist in different theories or the interpolants obtained might not be relevant (for example, lisp programs). This is formalized as follows:



**Definition 2.3.1.** [46] Let  $\mathcal{T}$  be a first-order theory of a signature  $\Sigma$  and let  $\mathcal{L}$  be the class of quantifier-free  $\Sigma$  formulas. Let  $\Sigma_{\mathcal{T}} \subseteq \Sigma$  denote a designated set of interpreted symbols in  $\mathcal{T}$ . Let  $A, B$  be formulas in  $\mathcal{L}$  such that  $A \wedge B \models_{\mathcal{T}} \perp$ . A theory-specific interpolant for  $(A, B)$  in  $\mathcal{T}$  is a formula  $I$  in  $\mathcal{L}$  such that  $A \models_{\mathcal{T}} I$ ,  $B \wedge I \models_{\mathcal{T}} \perp$ , and  $I$  refers only to  $AB$ -common symbols and symbols in  $\Sigma_{\mathcal{T}}$ .

**Example 2.3.1.1.** In example 2.3.0.1 we can tell  $c + 1 = 1$  is not an interpolant simplify because the symbol 1 only appears on the  $A$ -part. However, if  $\Sigma_{\mathcal{LIA}(\mathbb{Z})}$  contains the interpreted symbols of  $LIA(\mathbb{Z})$  (i.e.  $+, *, 0, 1, 2, \dots$ ), then  $c + 1 = 1$  becomes a theory-specific interpolant.

Notice that  $c = 0$  is an interpolant even if the set of interpreted symbols used for interpolation is empty.

## 2.4 Decision Procedures

Given a theory  $\mathcal{T}$  in a formula  $\psi$  in the language of the theory, is it possible to know  $\models_{\mathcal{T}} \psi$ ? The last question is known as the verification of the decision problem for the respective theory  $\mathcal{T}$ . This question has been studied extensively for many theories of interest [5].

Regarding the decidability of the theories mentioned in Section 2.2, it is known that EUF is undecidable [5], and the theory of ordered commutative rings is undecidable when the structure uses integers as the domain and the semantics of the arithmetical operations [19] (on the other hand, this theory can be decidable if we keep the same structure but use the reals as the domain [19]). Nonetheless, the quantifier-free fragment of EUF and the restriction imposed in the decision problem for the UTVPI theory allow efficient algorithms to decide validity and satisfiability in their respective theories [38, 18, 30].

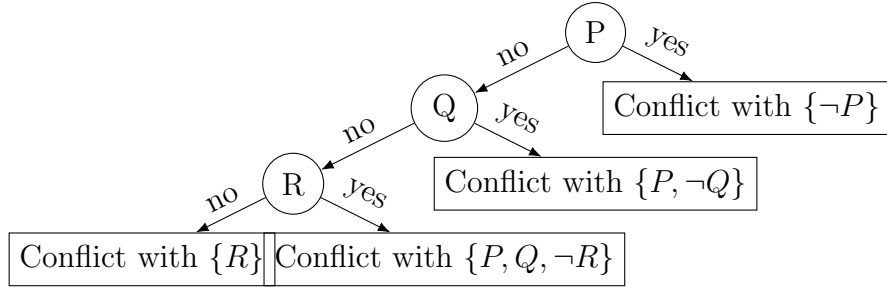


Figure 2.1: Example of DPLL execution on  $\{\{\neg P\}, \{P, Q, \neg R\}, \{R\}, \{P, \neg Q\}\}$

In the rest of this section we review some decision problems and provide references to their respective decision procedures used in the implementation work of the thesis.

### 2.4.1 Satisfiability and Satisfiability Modulo Theories

The satisfiability problem consists on finding a propositional assignment for a propositional formula. This problem is at the core level of complexity theory, defining an important class of problems known as NP, which is a class whose algorithms seem to be intractable. Developments in algorithms and heuristics [26, 36] have made possible to use satisfiability algorithms to solve real-world problems in verification <sup>2</sup>.

The DPLL algorithm [13] (an other extensions) is the frequent algorithm found in many SAT solvers. Fundamentally, it is a search-based algorithm which implements of operators (decide, unit-propagation, backtrack) to find a satisfiable assignment. If the algorithm is not able to find a satisfying assignment for a formula, then it is possible to extract a *resolution proof* based on the traces of the search operations.

**Example 2.4.0.1.** *We can see that the following resolution proof resembles the structure of the DPLL execution on figure 2.1, i.e. if we rotate the proof-tree and mark*

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<sup>2</sup>These advances do not provide an answer to the well-known P vs. NP problem. There are results indicating a class of problem instances for many of the SAT algorithms which cannot be solve in less that  $\mathcal{O}(2^n)$  steps [26].



## 2.4.2 Congruence Closure

The congruence closure problem consists of given a conjunction of equalities and disequalities  $\psi$  determine if an equality  $u = v$  follows from the consequence generated by  $\models_{EUF} \psi$ .

As noted in [11], it is just sufficient to compute the minimal relation containing the initial relation defined by the equalities in  $\psi$  closed under reflexivity, symmetry, transitivity and congruence considering all the subterms in the formulas  $\psi$  and  $u = v$ .

The authors in [18, 38] independently formulated an optimized version of the algorithm afore mentioned. Their key observation was to introduce a list of pointers keeping track of the antecedents nodes in the abstract syntax tree induced by the formulas. The latter allows a fast signature checking in order to determine if two nodes are equivalent under the equivalence relation of the formulas. The algorithm in [18] has better runtime complexity  $\mathcal{O}(n \log n)$  since it also implements a ‘modify the smaller half’ (using the union-find data structure). The congruence closure algorithm in [18] has a runtime complexity of  $\mathcal{O}(n^2)$ . Nonetheless, the authors reported no significant advantage of the first approach because, for verification purposes, the list of antecedents is usually small. Both approaches provides the FIND, MERGE operations.

In [41], the authors introduced a Union-Find data structure that supports the additional Explanation operation. This operation receives as input an equation between constants. If the input equation is a consequence of the current equivalence relation defined in the Union Find data structure, the Explanation operation returns the minimal sequence of equations used to build such equivalence relation, otherwise it returns ‘Not provable’. A proper implementation of this algorithm extends the traditional Union-Find data structure with a *proof-forest*, which consists of an additional representation of the underlying equivalence relation that does not compress

paths whenever a call to the Find operation is made. For efficient reasons, the Find operation uses the path compression and weighted union.

The main observation in [41] is that, in order to recover an explanation between two terms, by traversing the path between the two nodes in the proof tree, the last edge in the path guarantees to be part of the explanation. Intuitively, this follows because only the last Union operation was responsible of merging the two classes into one. Hence, we can recursively recover the rest of the explanation by recursively traversing the subpaths found.

Additionally, the authors in [41] extended the Congruence Closure algorithm [40] using the above data structure to provide Explanations for the theory of EUF. The congruence closure algorithm is a simplification of the congruence closure algorithm in [18]. The latter combines the traditional *pending* and *combine* list into one single list, hence removing the initial *combination* loop in the algorithm in [18].

The implementation work utilizes the latter congruence closure with explanations for the interpolation algorithm of the theory EUF. The idea was to use the explanation operator to construct uncommon-free Horn clauses.

### 2.4.3 Satisfiability of Horn clauses with ground equations

In [22] it was proposed an algorithm for testing the unsatisfiability of ground Horn clauses with equality. The main idea was to interleave two algorithms: *implicational propagation* (propositional satisfiability of Horn clauses) that updates the truth value of equations in the antecedent of the input Horn clauses [17]; and *equational propagation* (congruence closure for grounded equations) to update the state of a Union-Find data structure [21] that keeps the minimal equivalence relation defined by grounded equations in the input Horn clauses.

The author in [22] defined two variations of his algorithms by adapting the Congruence Closure algorithms in [18, 38]. Additionally, modifications in the data structures used by the original algorithms were needed to make the interleaving mechanism more efficient.

Our implementation uses the equality propagation mechanism in the algorithm proposed by Gallier when we have to deal with Horn clauses with ground equations. In addition, we also needed to design some modifications of the original formulation so it can integrate with the congruence closure with explanation algorithm mentioned in the previous section.

#### 2.4.4 Nelson-Oppen framework for theory and interpolation combination

The theory combination problem consists on taking a formula from the union of two (or more) disjoint languages and tell if such formula is satisfiable or not in the combined theory, i.e. a theory resulting after putting together two (or more) axiomatizations.

In [37] the authors defined a procedure to achieve the above problem. The key idea is to *purify* ( or separate) the subformulas by including additional constant symbols equating subterms such that the resulting formula can be splitted into components of the appropriate language for each theory solvers to work with. The separation naturally will hide relevant information to the solvers and they might not be able to decide satisfiability correctly. The authors noticed that to solve the above problem it is enough to share disjunction of equalities between the combined theories of shared terms. In addition, they proved that some theories have the following property:

**Definition 2.4.1.** *Let  $\mathcal{T}$  be a theory. We say that  $\mathcal{T}$  is a convex theory if a finite conjunction of formulas in  $\mathcal{T}$   $\psi = \bigwedge_{i=1}^m \psi_i$  satisfies  $\psi \models_{\mathcal{T}} \bigvee_{j=1}^n x_j = y_j$ , then exists*

$k \in \{1, \dots, n\}$  such that  $\psi \models_{\mathcal{T}} x_k = y_k$ .

Hence, it is important to detect whether the theories involved are convex or not since this can improve performance since convex theories do not need to share disjunctions of equalities as mentioned before (since all these disjunctions imply a single equality).

**Example 2.4.1.1.**    • *The conjunctive fragment of equality logic is convex since it can always decide the membership of an equation in the equivalence relation.*

- *The theory of UTVPI over the integers is not convex. To see the latter consider  $1 \leq x \wedge x \leq 2 \models_{UTVPI(\mathbb{Z})} 1 = x \vee 2 = x$ . However, it is not the case that  $1 \leq x \wedge x \leq 2 \models_{UTVPI(\mathbb{Z})} 1 = x$  nor  $1 \leq x \wedge x \leq 2 \models_{UTVPI(\mathbb{Z})} 2 = x$ .*

An interpolation combination framework as proposed in [46] follow the same idea towards theory combination. Inductively, they define *partial interpolants* for each shared equality/disjunction of equalities until some theory reaches the unsatisfiable state, which is expected since an interpolant is a pair a mutually contradicting formulas.

This framework was implemented in the thesis work. This framework was chosen in particular since it allows to work with non-convex theories (in our case for the theory of UTVPI over  $\mathbb{Z}$ ).

## 2.5 General system description

The algorithms implemented in this thesis used the C++ programming language. The overall architecture of the system is the following:

All the decision procedures mentioned in this chapter were implemented with the exception of the SAT/SMT algorithms. For the latter, PicoSAT/TraceCheck[2] and

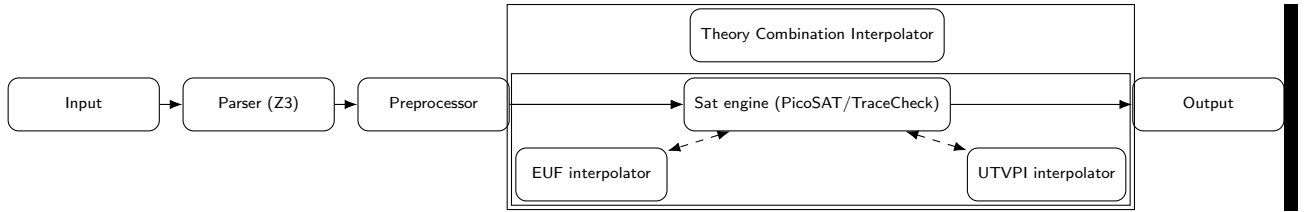


Figure 2.3: General System Diagram

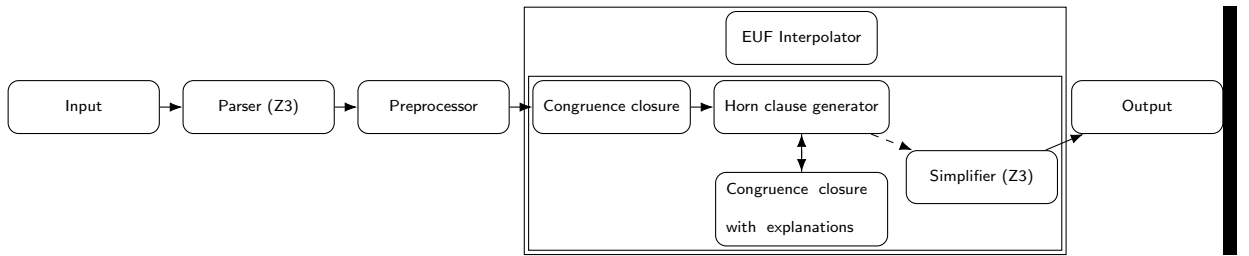


Figure 2.4: EUF Interpolator Diagram

Z3 [14] were chosen as the libraries to work with these algorithms. The rest of this section discusses some minor modifications implemented in the above mentioned Z3 and the proof format output by PicoSAT/TraceCheck.

### 2.5.1 Minor modifications to Z3

Z3 standard input is SMTLib2 [1]. This grammar does not provide a standard specification regarding a suitable format to work with interpolants. Interpolation

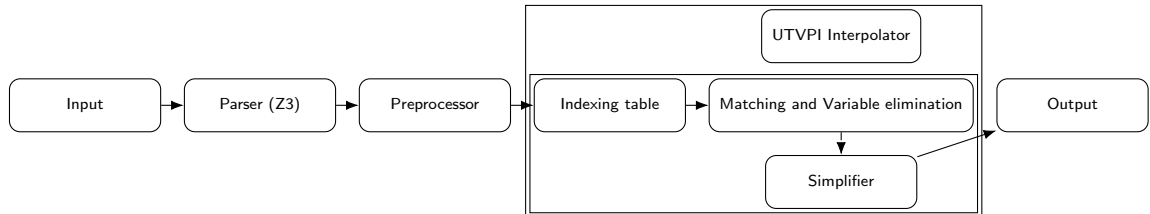


Figure 2.5: UTVPI Interpolator Diagram



software read interpolant formulas based on the order of appearance in a conjunction [33]. In our case we require two conjuncts of conjunctions of literals in the EUF theory, UTVPI theory or combined theory.

As we can notice in figures 2.3, 2.4, and 2.5, there is preprocessor component which prefixes the names of uninterpreted symbols with the strings `a_`, `b_`, `c_` to indicate that the symbol name is either an A-local, B-local, or common symbol respectively. We extended Z3's API with functions that test if a formula is A-local, B-local, AB-pure, AB-common based on the definitions in [46] because it is necessary to constantly check these conditions for splitting purposes. Another reason for the latter is justified because the implemented congruence closure algorithm takes as an additional criteria to maintain as representative term an AB-common term. A similar change was implemented in the congruence closure implementation of Z3. Nonetheless, it was irrelevant since Z3's internal structure separates the abstract syntax tree, which is part of its API with the enode data structures, which does not allow the user to modify it. This is the reason why it was not possible to work directly with Z3 congruence closure implementation and a separate implementation was necessary.

### 2.5.2 PicoSAT/TraceCheck Proof Format

We used PicoSAT/TraceCheck to reconstruct a resolution-based proof necessary for Pudlak's algorithm in the interpolation combination component. Given that Z3 provides a user friendly proof-producing API [16], why did the implementation work require another SAT solver to obtain the resolution-proof?

There are several reasons for the latter. First, the author of the thesis was not able to find an appropriate configuration of parameters for the SMT solver to provide such proofs. In order to grasp an idea of the latter, it was implemented a Z3 proof

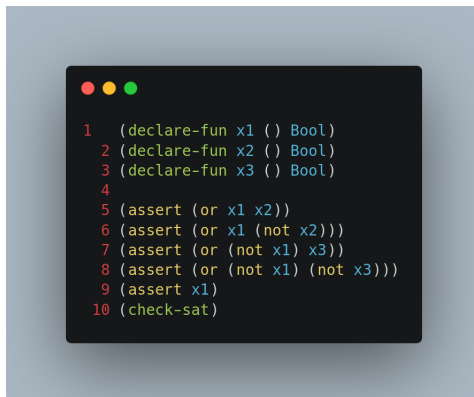


Figure 2.6: Problematic SMT query for resolution proofs

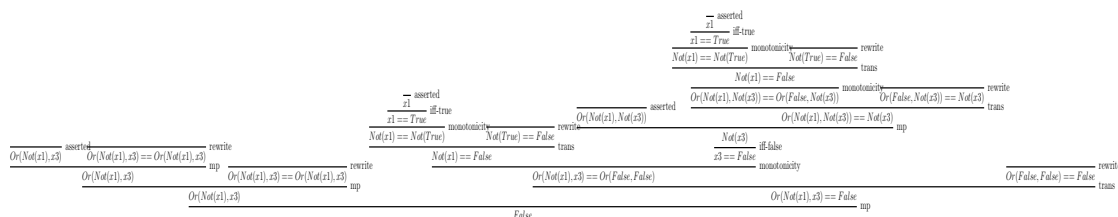


Figure 2.7: Z3 proof of figure 2.6

parser that render a pdf rendered by L<sup>A</sup>T<sub>E</sub>X. Many examples indicated that Z3 selects more convenient theories to work with some problems. For instance, the formula in pure propositional logic shown in figure 2.6 used proof rules from EUF <sup>3</sup>.

Thus, we opted to use the PicoSAT SAT solver which implements the DPLL algorithm. PicoSAT, as many SAT solvers, take as input a DIMACS cnf file which consists of a straightforward language to denote clauses. (Optional) Lines starting with the character ‘c’ denote comments, and a single line starting the string ‘p cnf’ followed by two positive integers denote the number of variables and number of clauses respectively. The next lines encode clauses. Each line is a list of integers numbers terminating with 0. Positive number encode non-negated literals whose identifier is the number, and negative numbers encode the negated literal whose

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<sup>3</sup>Z3 uses the term monotonicity instead of congruence

identifier is the absolute value of the number.

The proof format produced by PicoSAT consists of an ordered list of clauses which are either from the input problem (facts) or learned during the search. Each line is encoded following the DIMACS format, but extended to include an identifying number at the beginning of each line (which identifies the clause), and a second list of integers after the terminating zero of the clause which denotes the non-ordered list of clause identifiers involved in a hyper resolution step; the resulting clause obtained from this step is precisely the clauses encoded by the line. An additional zero is used to terminate the second list of integers.

**Example 2.5.0.1.** *The following extended proof trace is obtained after running the PicoSAT solver on the cnf file encoding the formula in example 2.1:*

1	-1	0	0		
2	-3	2	1	0	0
3	3	0	0		
4	1	-2	0	0	
6	-2	0	1	4	0
7	0	1	6	2	3

Figure 2.8: Example of extended proof trace by PicoSAT for problem 2.1

*As we can notice, the first four lines encode the original input clauses; 1 denotes variable  $P$ , 2 denotes variable  $Q$ , and 3 denotes variable  $R$ . Also, these lines do not include any clause identifier in the second list of number since they are facts. The last two lines denote clauses with identifier 6 and 7. Clause 6 is  $\neg Q$ , which is obtained after the hyper resolution step using clauses 1 and 4. The last line does not contain elements in the first list since this encodes the contradiction clause.*

From the previous example we can observe the list of clause identifiers for the hyper resolution step are not ordered in the sense that we cannot produce multiple resolution steps from the representation. In order to obtain the latter one can

implement a unit propagation algorithm and discover such ordering.

TraceCheck is a tool incorporated in the previous version of PicoSAT [39] to verify proof traces. It is also able to obtain a pure resolution proof. The presentation of the proof trace is restricted to contain only two clause identifiers in the second list for each clause.

**Example 2.5.0.2.** *The following extended proof trace is obtained after running the TraceCheck on the previous proof trace:*

3	3	0	0				
2	1	2	-3	0	0		
4	1	-2	0	0			
1	-1	0	0				
6	-2	0	1	4	0		
7	2	-3	0	1	2	0	
8	-3	0	7	6	0		
9	0	8	3	0			

Figure 2.9: Example of extended proof trace by PicoSAT for problem 2.1

## Chapter 3

# Interpolation algorithm for the theory of EUF

Interpolation algorithms for the theory of equality with uninterpreted functions are relevant as the core component of verification algorithms. Many useful techniques in software engineering like bounded/unbounded model checking and invariant generation benefit directly from this technique. In [7], the authors introduced a methodology to debug/verify the control logic of pipelined microprocessors by encoding its specification and a logical formula denoting the implementation of the circuit into a EUF solver.

Previous work addressing the interpolation problem for EUF has involved techniques ranging from interpolant-extraction from refutation proof trees [32, 33, 45], and colored congruence closure graphs [20]. Kapur's algorithm uses a different approach by using approximated quantifier-elimination, a procedure that given a formula, it produces a logically equivalent formula without a variable in particular [19].

### 3.1 Algorithm

Kapur's interpolation algorithm for the EUF theory uses quantifier-elimination techniques to remove symbols in the first formula of an inconsistent pair of formulas that are not common with the second formula of the latter. Hence, the input for this algorithm is a conjunction of equalities in the EUF theory and a set of symbols to eliminate, also known as uncommon symbols. In preparation to discuss Kapur's algorithm we need to provide the following definitions.

**Definition 3.1.1.** *Let  $f$  be a  $n$  – ary function symbol and  $a_1, \dots, a_n, b$  terms from the EUF language. We say*

$$f(a_1, \dots, a_n) = b$$

*is an  $f$  – equation if the terms  $a_1, \dots, a_n, b$  are constants in the EUF language. We refer to the outermost symbol of the  $f$ -equation as the function symbol appearing in such  $f$  – equation.*

$f$  – equations are used in Kapur's algorithm to simplify the structure of terms and in Phase II to expose hidden arguments and eliminate uncommon function symbols.

As part of the input of Kapur's algorithm, there is a set of uncommon symbols given directly or computed from the inconsistent pair of formulas by inspection. Using these symbols we can define the following recursive definition of *uncommon terms*:

**Definition 3.1.2.** *A term  $t$  in the EUF language is uncommon if:*

- *$t$  is an uncommon constant*
- *$t$  is a function application of the form  $f(t_1, \dots, t_n)$  where either  $f$  is an uncommon symbol or any  $t_i$  where  $1 \leq i \leq n$  is an uncommon term*

### Chapter 3. Interpolation algorithm for the theory of EUF

Similarly, we can extend homomorphically the notion of uncommon terms to uncommon predicates following a similar construction.

- Let  $t_1, t_2$  be terms in the EUF language.  $t_1 = t_2$ ,  $t_1 \neq t_2$  are uncommon predicates if either  $t_1$  or  $t_2$  are uncommon terms.
- Let  $\psi, \varphi$  be predicates in the EUF language.  $\psi \star \varphi$  are uncommon predicates if either  $\psi$  or  $\varphi$  are uncommon predicates where  $\star \in \{\wedge, \vee, \rightarrow\}$ .  $\neg\psi$  is an uncommon predicate if  $\psi$  is an uncommon predicate.

Regarding extensions of the language with new constants, the *uncommon* property is preserved under equalities. Formally, we mean the following:

**Definition 3.1.3.** Let  $\mathfrak{L}$  be a EUF language and  $\mathfrak{a}$  a constant symbol not belonging to  $\mathfrak{L}$ . We say  $\mathfrak{a}$  is a common constant under the theory  $\mathcal{T}$  in the extended language  $\mathfrak{L} \cup \{\mathfrak{a}\}$  if there exists a common term  $t$  in the language  $\mathfrak{L}$  such that  $\models_{\mathcal{T}} t = \mathfrak{a}$ , otherwise  $\mathfrak{a}$  is an uncommon constant.

Since congruence closure algorithms are relevant to Kapur's algorithm, we introduce the following definition for notation purposes:

**Definition 3.1.4.** Let  $\mathcal{E}$  be an equivalence relation between grounded terms of some language  $\mathfrak{A}$ . The function  $\text{repr}_{\mathcal{E}} : \mathfrak{A} \rightarrow \mathfrak{A}$ ,  $\text{repr}_{\mathcal{E}} : a \mapsto b$  where  $b$  is the representative element in  $\mathcal{E}$  for  $a$ .

The interpolating formula produced by Kapur's algorithm is a conjunction of equations and Horn clauses. A Horn clause is a disjunction of literals which contain at most one non-negated literal. Due to its relevance in the procedure in the EUF theory, for a Horn clause  $h$  we denote  $\text{antecedent}(h)$  to be the conjunction of disequations in the disjunction of  $h$  and  $\text{head}(h)$  to be either the equation in the disjunction if such is present in  $h$  or the particle  $\perp$  otherwise.

The main steps in Kapur's algorithm for interpolant generation for the EUF theory are the following:

- **Flattening:** For each subterm  $t$  in the input formula assign a fresh unique constant  $\mathbf{a}_t$ . Additionally, for each subterm  $t$  generate new equations of the form:

- $c = \mathbf{a}_c$ , if  $t$  is a constant  $c$
- $f(\mathbf{a}_{t_1}, \dots, \mathbf{a}_{t_n}) = \mathbf{a}_{f(t_1, \dots, t_n)}$ , if  $t$  is a function application of the form  $f(t_1, \dots, t_n)$

Clearly, we can see that this step generates  $f$  – equations.

- **Elimination of uncommon terms using congruence closure:** This step builds an equivalence relation  $\mathcal{E}$  of the  $f$  – equations introduced in the Flattening step using a congruence closure algorithm such that the representatives are common terms whenever possible. Uncommon terms appearing in the current conjunction of equations are replaced by their representatives.
- **Horn clause generation by exposure:** This step produces for all pairs of  $f$  – equations ( $f(\mathbf{a}_1, \dots, \mathbf{a}_n) = \mathbf{c}, f(\mathbf{b}_1, \dots, \mathbf{b}_n) = \mathbf{d}$ ) Horn clauses of the form  $\bigwedge_{i=1}^n (repr_{\mathcal{E}}(\mathbf{a}_i) = repr_{\mathcal{E}}(\mathbf{b}_i)) \rightarrow repr_{\mathcal{E}}(\mathbf{c}) = repr_{\mathcal{E}}(\mathbf{d})$  when any of the two following situations happen <sup>1</sup>:

- The outermost symbol of the  $f$  – equations is an uncommon symbol.
- There is at least one constant argument in any of the  $f$  – equations that is an uncommon constant.

---

<sup>1</sup>Trivial equations in the antecedent of a Horn clause are removed; if the head equation of Horn clause produced in this step is trivial then such Horn clause is discarded



- **Conditional elimination:** We identify the Horn clauses  $h := \bigwedge_i (c_i = d_i) \rightarrow a = b$  that have *common antecedents* and uncommon head equations. We perform the following procedure:
  - if  $a$  and  $b$  are both uncommon terms: replace the equation  $a = b$  appearing in the antecedents of all the current Horn clauses by  $\text{antecedent}(h)$ .
  - if either  $a$  is common and  $b$  uncommon: replace  $b$  by  $a$  in all the current Horn clauses  $h'$  and append  $\text{antecedent}(h)$  to  $\text{antecedent}(h')$ .
  - if either  $a$  is uncommon and  $b$  common: Proceed similarly as in the previous case.

We repeat this step until we cannot produce any new Horn clauses.

- **Conditional replacement:** For each Horn clause of the form  $\bigwedge_i (a_i = b_i) \rightarrow u = c$  where the antecedent is common, the term  $u$  in its head equation is an uncommon term, and the term  $c$  is a common term, replace every instance of  $u$  appearing in each  $f$  – *equation* by  $c$  to generate Horn clauses with antecedent  $\bigwedge_i a_i = b_i$ .

Return the conjunction of formulas obtained as the interpolant.

If the user is not interested in an explicit interpolant, we can present a **lazy/pseudo interpolant** which is an ordered sequence of the original equations together with the Horn clauses produced in Phase II using an appropriate order between the uncommon terms. In order to provide a description of a procedure for the latter we need to introduce the following definitions:

**Definition 3.1.5.** Let  $\succ$  be a partial order between terms such that  $a \succ b$  whenever  $a$  is common and  $b$  is uncommon. A dependency pair for a horn clause  $h := \bigwedge_i (a_i = b_i) \rightarrow c = d$  is a pair  $(\min(c, d, \succ), \{\max(c, d, \succ)\} \cup \{u \mid u \text{ is an uncommon term}\})$ .

### Chapter 3. Interpolation algorithm for the theory of EUF

appearing in  $\text{antecedent}(h)$  }). The first element of a dependency pair is denoted as the target of  $h$  and the second element the source of  $h$ .

A valid dependency pair is a dependency pair for some Horn clause  $h$  which its target is not included in its source.

We can notice from definition 3.1.5 than even for equations, its source is never empty.

**Definition 3.1.6.** Let  $\succ$  be a partial order between terms such that  $a \succ b$  whenever  $a$  is common and  $b$  is uncommon. Given a set of Horn clauses  $H$ , a dependency graph for  $H$  is the directed graph  $G_H = (V, E)$  where

- $V := \{(target_i, source_i) \mid (target_i, source_i) \text{ is a valid dependency pair from } H\}$
- $E := \{(target_i, source_i) \rightarrow (target_j, source_j) \mid \{target_i\} \cup source_i \subseteq source_j\}$

Using *valid dependency pairs* of the Horn clauses produced in Phase II we can construct an acyclic directed graph as shown by the following theorem:

**Theorem 3.1.1.** For any set of Horn clauses  $H$ , its dependency graph never contains a cycle between its nodes.

*Proof.* Suppose there exists a sequence of  $n$  nodes such that  $(target_1, source_1) \rightarrow \dots \rightarrow (target_n, source_n) \rightarrow (target_1, source_1)$ . Since  $\subseteq$  is transitive we can conclude that  $target_1 \in source_1$ , which leads to a contradiction since all the nodes in a dependency graph are valid dependency pairs.  $\square$

Thus, given a set of Horn clauses  $H$ , we can compute its dependency graph and use a topological sort algorithm to produce the ordered sequence required in the lazy interpolant representation. Lazy interpolants avoid the possible exponential

size of the formal interpolant. This representation is useful because it provides a more compact representation of the interpolant that the user might be able quicker to obtain. Additionally, the user might be just interested in a particular subformula of the interpolant, so the latter representation offers such feature. This algorithm allows a flexible implementation which can lead several optimizations based on the nature and applications of the interpolant.

## 3.2 Implementation

The description of the interpolation algorithm presented in the previous section suggests a straight forward implementation of the first two stages using well-known algorithms [38, 18] and data structures from the SMT solver Z3 [14] to represent elements from the EUF language. One particular change was required in the congruence closure algorithm since Kapur’s algorithm keeps common terms as representatives whenever common terms belong to a partition of terms induced by the equivalence relation.

The algorithm in [18] uses a union-find data structure to encode the equivalence classes of the nodes in the abstract syntax tree of the input formula with a ‘modify the smaller subtree’ strategy. This means that when two nodes  $u, v$  in the abstract syntax tree are meant to be merged, the representative of the new combined equivalence class is the node which has a bigger numbers of predecessors nodes in the abstract syntax tree pointing to the equivalence class of the node. The idea was to update the least amount of nodes that possibly can change representatives due to the most recent merge operation of the equivalence classes and congruence.

**Notation 3.2.1.** *Given a theory  $\mathcal{E}$  and a term  $u$  in the language of  $\mathcal{E}$ , we can indicate by  $[u]$  the equivalence class induced by  $\mathcal{E}$ , i.e.  $[u]_{\mathcal{E}} = \{v \in TERMS \mid \models_{\mathcal{E}} u = v\}$ .<sup>2</sup>*

---

<sup>2</sup>If the theory is clear from context, the notation  $[u]$  denotes the equivalence class of  $u$ .

Our algorithm uses a difference partial order to maintain common terms as representatives of the equivalence classes. The non-reflexive relation  $\succ_{common}$ <sup>3</sup> is defined for all nodes  $u, v$  in the abstract syntax tree of terms as:

$$u \succ_{common} v = \begin{cases} |list(u)| > |list(v)| & \text{if } (u \text{ is a common term} \Leftrightarrow \\ & v \text{ is a common term}) \\ u \text{ is a common term} & \text{otherwise} \end{cases}$$

where  $list(u) = \{f(u_1, \dots, u_n) \in \text{TERMS} \mid \exists i \in \{1, \dots, n\} \text{ such that } u_i \in [u]\}$

In the next section we will discuss the changes proposed to Phase III in Kapur's algorithm.

### 3.2.1 New conditional elimination step in Kapur's algorithm

The modification of Phase III implemented in this thesis work combines and extends the algorithms and data structures introduced in [22, 41]. The algorithm in [22] is a direct extension of [17] which adapts a congruence closure algorithm from [38, 18] in order to update the union-find data structure maintaining the equivalence relation between all the subterms in the input formula. The implementation of the congruence closure algorithm in [41] extends the usual *Find*, *Merge* operations on the union-find data structure with the *Explain* operator, which accomplishes the following:

*Explain*( $e, e'$ ): if a sequence  $U$  of unions of pairs  $(e_1, e'_1), \dots, (e_p, e'_p)$  has taken place, it returns a minimal subset  $E$  of  $U$  such that  $(e, e')$  belongs to the equivalence relation generated by  $E$  and it returns  $\perp$  if no such  $E$  exists.

---

<sup>3</sup>The reflexive relation  $u \succeq_{common} v$  is defined as  $u = v \vee u \succ_{common} v$ , where the equality between nodes is defined as  $|list(u)| = |list(v)| \wedge u \text{ is a common term} \iff v \text{ is a common term}$ .

The motivation behind the combination of Gallier’s data structure and the congruence closure algorithm with explanations is twofold:

1. First, we want to introduce common equations from the antecedents of the Horn clauses obtained during Phase II to a conditional equivalence relation as well as updating the conditional equivalence relation structure by using the equation propagation mechanism from the congruence closure algorithm and the implicational propagation component from Gallier’s structure.
2. Additionally, we want to find all the common Horn clauses provable from the original input of equations. The Explain operator in [41] is recursively used to construct the antecedent of such Horn clauses during the conditional replacement step. The Explain operator traverses a proof-tree data structure containing the nodes that were used to combine equivalence classes in the underlying union-find data structure. Thus, the MERGE operation in [22] is required to update the proof tree <sup>4</sup> from [41] as well.

The thesis work accomplishes the previous points by implementing the following:

1. In addition to the parsing procedure and initialization of the Gallier data structure, we assert into the union-find data structure every common equation in the antecedent of a Horn clause.
2. The implemented C++ class for the congruence closure with explanation <sup>5</sup> includes a pointer as data member to the class implementation for the Gallier data structure in order to propagate the equational information achieved during merges and updates due to congruence. The modified algorithms appears below in pseudo-code notation:

---

<sup>4</sup>This structure is used in order to retrieve explanations.

<sup>5</sup>A fragment of the actual code is shown at Section 6.3.

---

**Algorithm 1** Modified Unsatisfiability Testing for Ground Horn Clauses

---

```

1: procedure SATISFIABLE(var H : Hornclause; var queue, combine: queueType;
   var GT(H) : Graph; var consistent : boolean)
2:   while queue not empty do
3:     node := pop(queue);
4:     for clause1 in H[node].clauselist do
5:       if numargs[clause1] = 0 then
6:         nextnode := poslitlist[clause1];
7:         if ¬ H[nextnode].val then
8:           if nextnode ≠ ⊥ then
9:             queue := push(nextnode, queue);
10:            H[nextnode].val := true;
11:            u := left(H[nextnode].atom);
12:            v := right(H[nextnode].atom);
13:            if FIND(R, u) ≠ FIND(R, v) then
14:              combine := push((u, v), combine);
15:            end if
16:          else
17:            consistent := false;
18:          end if
19:        end if
20:      end for
21:    end while
22:    if queue is empty then
23:      closure(combine, queue, R);
24:    end if
25:  end while
26: end procedure

27: procedure CLOSURE(var combine, queue : queueType; var R : partition)
28:   while combine is not empty do
29:     (u, v) = pop(combine)
30:     MERGE(R, u, v, queue)
31:   end while
32: end procedure

```

---

3. We implement the ExtendedExplain procedure which returns a list of common equations when an equation given as input is provable in the conditional equivalence relation. The pseudo-code for the ExtendedExplain procedure is shown below:

---

**Algorithm 2** Modified Congruence Closure with Explanation Algorithms - Merge

---

```

procedure MERGE( $R$  : partition,  $u, v$  : node; queue, combine : queue type)
2:   if  $u$  and  $v$  are constants  $a$  and  $b$  then
      add  $a = b$  to Pending;
4:   Propagate();
      else  $\triangleright u = v$  is of the form  $apply(a_1, a_2) = a$ 
6:     if  $Lookup(Representative(a_1), Representative(a_2))$  is some
           $apply(b_1, b_2) = b$  then
          add  $(apply(a_1, a_2) = a, apply(b_1, b_2) = b)$  to Pending;
8:     Propagate();
          else
10:      set  $Lookup(Representative(a_1), Representative(a_2))$  to
           $apply(a_1, a_2) = a$ ;
          add  $apply(a_1, a_2) = a$  to  $UseList(Representative(a_1))$  and to
           $UseList(Representative(a_2))$ ;
12:    end if
      end if
14: end procedure

```

---

### 3.2.2 Invariants of the proposed conditional elimination step

We can prove that if a pair of terms  $t_1, t_2$  belong to the same equivalence class, then `ExtendedExplain` returns a list of common equations using the axiomatization introduced in 2.2.1 as inference rules with the following lemmas:

**Lemma 3.2.1.** *Let  $H$  be a set of Horn clauses produced by Phase II of Kapur's algorithm in some EUF language  $\mathcal{L}$ ,  $\mathcal{E}$  the conditional equivalence relation over the terms of  $\mathcal{L}$  produced by the above algorithm on  $H$ , and  $t_1, t_2$  terms in  $\mathcal{L}$  such that  $\models_{\mathcal{E}} t_1 = t_2$ . If  $\exists$  uncommon  $a = b \in Explain(t_1, t_2)$ , then  $\exists h \in H$  such that  $h$  is of the form  $\bigwedge_i (c_i = d_i) \rightarrow a = b$ .*

*Proof.* By the definition of the `Explain` operator in 3.2.1, we see that the `Explain` operator will return a list of all the asserted equations in the conditional equivalence relation. By inspection of the proposed algorithm, assertions into the conditional equivalence relation only happen in two places: when an equation in the antecedent of a horn clause of  $H$  is common, and while updating the underlying union-find data

---

**Algorithm 3** Modified Congruence Closure with Explanation Algorithms - Propagate

---

```

procedure PROPAGATE( )
2:   while Pending is non-empty do
      Remove E of the form  $a=b$  or  $(\text{apply}(a_1, a_2) = a, \text{apply}(b_1, b_2) = b)$  from
      Pending
4:   if Representative( $a$ )  $\neq$  Representative( $b$ ) and w.l.o.g.
       $|\text{ClassList}(\text{Representative}(a))| \leq |\text{ClassList}(\text{Representative}(b))|$  then
       $\text{old}_{repr_a} := \text{Representative}(a)$ ;
6:   Insert edge  $a \rightarrow b$  labelled with E into the proof forest;
      for each  $c$  in  $\text{ClassList}(\text{old}_{repr_a})$  do
8:     set Representative( $c$ ) to Representative( $b$ )
      move  $c$  from  $\text{ClassList}(\text{old}_{repr_a})$  to  $\text{ClassList}(\text{Representative}(b))$ 
10:    for each pointer L in  $\text{ClassList}(u)$  do
      if  $H[L].\text{val} = \text{false}$  then
12:      set the field  $H[L].\text{lclass}$  or  $H[L].\text{rclass}$  pointed to by p to
      Representative( $b$ )
      if  $H[L].\text{lclass} = H[L].\text{rclass}$  then
14:       queue := push(L, queue);
        $H[L].\text{val} := \text{true}$ 
16:     end if
      end if
18:    end for
      end for
20:    for each  $\text{apply}(c_1, c_2) = c$  in  $\text{UseList}(\text{old}_{repr_a})$  do
      if Lookup(Representative( $c_1$ ), Representative( $c_2$ )) is some
       $\text{apply}(d_1, d_2) = d$  then
22:      add  $(\text{apply}(c_1, c_2) = c, \text{apply}(d_1, d_2) = d)$  to Pending;
      remove  $\text{apply}(c_1, c_2) = c$  from  $\text{UseList}(\text{old}_{repr_a})$ ;
24:    else
      set Lookup(Representative( $c_1$ ), Representative( $c_2$ )) to  $\text{apply}(c_1,$ 
       $c_2) = c$ ;
26:    move  $\text{apply}(c_1, c_2) = c$  from  $\text{UseList}(\text{old}_{repr_a})$  to
       $\text{UseList}(\text{Representative}(b))$ ;
      end if
28:    end for
      end if
30:  end while
end procedure

```

---

structure when the *numargs* field of a Horn clause becomes zero. The latter assert the consequent of such Horn clause. Since  $u$  is not common, then such Horn clause should exists in  $H$ . □



---

**Algorithm 4** Auxiliary function - ExtendedExplain

---

```

procedure EXTENDED_EXPLAIN(  $t_1 : \text{TERMS}, t_2 : \text{TERMS}, H : \text{Horn clauses},$ 
 $\mathcal{E} : \text{Conditional Equivalence Relation}$ )
2:   if  $\text{Find}(t_1, \mathcal{E}) \neq \text{Find}(t_2, \mathcal{E})$  then
      throw Error:  $t_1, t_2$  do not belong to the same equivalence class
4:   end if
      if  $t_1.\text{id}() = t_2.\text{id}()$  then
6:     return  $\{\}$ 
      end if
8:    $\{c, u\} = \text{Explain}(t_1, t_2, \mathcal{E})$  where  $c$  is a list of common equation and  $u$  is a
      list of uncommon equations
      return  $c \cup \bigcup \{ \text{ExtendedExplain}(a, b, H, \mathcal{E}) \mid \text{consequent} \in u, a = b \in$ 
 $\text{antecedent}, (\bigwedge \text{antecedent}) \rightarrow \text{consequent} \in H \}$ 
10: end procedure

```

---

**Lemma 3.2.2.** *Let  $H$  be a set of Horn clauses produced by Phase II of Kapur's algorithm in some EUF language  $\mathcal{L}$ ,  $\mathcal{E}$  the conditional equivalence relation over the terms of  $\mathcal{L}$  produced by the above algorithm on  $H$ , and  $t_1, t_2$  terms in  $\mathcal{L}$  such that  $\models_{\mathcal{E}} t_1 = t_2$ .*

*The list of equations returned by  $\text{ExtendedExplain}(t_1, t_2, H, \mathcal{E})$  contains only common equations.*

*Proof.* If  $t_1, t_2$  belong to the same equivalence relation, then there exists a closed derivation (proof tree) with  $t_1 = t_2$  as root node and the asserted equations introduced by the modified algorithm as leaves.

The proof proceeds by induction on the complexity of the derivation:

- Case: the last inference rule in the proof tree was Reflexivity. The ExtendedExplain algorithm returns an empty list. Thus, the statement is vacuously true.
- Case: the last inference rule in the proof tree was Symmetry. Line 8 of ExtendedExplain returns a list of common equations  $c$  and a list of uncommon equations  $u$ . By Lemma 3.2.1 we see that there exists a Horn clause in  $H$  for

every uncommon equation in  $u$ . Line 9 of `ExtendedExplain` applies recursively to each of these equations. Since these equations belong to the conditional equivalence relation, there exists a derivation for the latter. By the inductive invariant, `ExtendedExplain` should return a list of common equation for these subproofs. Thus, the final list contains only common equations.

- Case: the last inference rule in the proof tree was Transitivity. The proof follows similarly to the previous case applied to the two subproofs for the Transitivity rule.
- Case: the last inference rule in the proof tree was Congruence. The proof follows similarly to the previous case applied  $n$  times for each subproof in the Congruence rule.

□

**Corollary 3.2.2.1.** *Let  $H$  be a set of Horn clauses produced by Phase II of Kapur's algorithm,  $\mathcal{E}$  the equivalence relation obtained in Phase I of Kapur's algorithm, and  $\mathcal{E}'$  the conditional equivalence relation obtained after the modified conditional elimination algorithm. Then  $H \models_{\mathcal{E}'} a = b$  if and only if  $\exists$  common  $x \subseteq \bigcup_{h \in H} \text{antecedent}(h)$  such that  $H \models_{\mathcal{E}} \bigwedge x \rightarrow a = b$ .*

### 3.2.3 New conditional replacement step in Kapur's algorithm

Once the conditional equivalence relation is built after the execution of the previous step, we can compute conditional replacements as follows. For the latter, we will require the following auxiliary functions:

By inspection is easy to notice that *Candidates* and *allCandidates* return a set of common terms and a set of sets of common terms respectively.

---

**Algorithm 5** Auxiliary function - Candidates

---

```

procedure CANDIDATES(z3::expr const & t)
2:   if t is common then
       return {t}
4:   else
       return {t'|t' ∈ Class(t), t' is common}
6:   end if
end procedure

```

---



---

**Algorithm 6** Auxiliary function - allCandidates

---

```

procedure ALLCANDIDATES(z3::expr const & t)
2:   if t is a constant then
       undefined
4:   end if
       if t has f-symbol uncommon then
6:       return {{}};
       end if
8:   if t has f-symbol common and is of the form  $f(t_1, \dots, t_n)$  then
       return {Candidates( $t_1$ ), ..., Candidates( $t_n$ )};
10:  end if
end procedure

```

---

The proposed conditional replacement step produces common Horn clauses from previous uncommon equations and uncommon Horn clauses obtained in Phase II of Kapur's algorithm. The pseudo-code of the algorithms to process the original equations are shown below:

**Lemma 3.2.3.** *Let  $Eqs := \{f(\mathbf{a}_{i,1}, \dots, \mathbf{a}_{i,n}) = \mathbf{a}_i\}$  be the set of equations produced in Phase I of Kapur's algorithms,  $H$  the set of Horn clauses obtained in Phase II, and  $\mathcal{E}'$  the conditional equivalence relation obtained by the proposed conditional replacement step. Executing the proposed conditional replacement step on  $Eqs, H, \mathcal{E}$  produces common Horn clauses.*

*Proof.* By inspecting the conditional replacement algorithm in 3.2.3, we noticed the antecedents of the Horn clauses constructed are obtained using the ExtendedExplain procedure, which from lemma 3.2.2 it produces a list of only common equations.

Additionally, the equations in the consequent part of the Horn clauses produced

---

**Algorithm 7** Conditional Replacement - Part 1

---

```

procedure CONDITIONAL REPLACEMENT( $z3::\text{expr const} \& x, z3::\text{expr const} \&$ 
 $y, H : \text{Horn clauses}, \mathcal{E} : \text{conditional equivalence relation}$  )
2:   if  $x$  is constant and  $y$  is constant then
      for  $\sigma_x$  in  $\text{Candidates}(x)$  do
4:       for  $\sigma_y$  in  $\text{Candidates}(y)$  do
            horn_clause.add( $\text{ExtendedExplain}(x, \sigma_x, H, \mathcal{E}) + \text{ExtendedEx-}$ 
plain( $y, \sigma_y, H, \mathcal{E}$ ),  $\sigma_x = \sigma_y$ )
6:       end for
      end for
8:   end if
      if  $x$  is constant and  $y$  is of the form  $f_y(t'_1, \dots, t'_{k_2})$  then
10:      for  $\sigma_x$  in  $\text{Candidates}(x)$  do
            for  $\sigma_{f_y}$  in  $\text{Candidates}(f_y(t'_1, \dots, t'_{k_2}))$  do
12:                horn_clause.add( $\text{ExtendedExplain}(x, \sigma_x, H, \mathcal{E}) +$ 
ExtendedExplain( $f_y(t'_1, \dots, t'_{k_2}, H, \mathcal{E}), \sigma_y$ ),  $\sigma_x = \sigma_{f_y}$ )
            end for
14:      for  $\text{arguments}_{f_y}$  in  $\text{CartesianProd}(\text{AllCandidates}(f_y(t'_1, \dots, t'_{k_2})))$  do
            horn_clause.add( $\text{ExtendedExplain}(x, \sigma_x, H, \mathcal{E}) + \sum_{i=1}^{k_2}$ 
ExtendedExplain( $t'_i, \text{arguments}_{f_y}[i], H, \mathcal{E}$ ),  $\sigma_x = f_y(\text{arguments}_{f_y})$ )
            end for
16:      end for
      end if
18:   if  $x$  is of the form  $f_x(t_1, \dots, t_{k_1})$  and  $y$  is a constant then
20:       return  $\text{CONDITIONAL ELIMINATION}(y, x)$ ;
      end if
22: end procedure

```

---

in the conditional replacement procedure are of the form:

- $\sigma_x = \sigma_y$ , where  $\exists t_1, t_2 \in \text{TERMS}$  such that  $\sigma_x \in \text{Candidates}(t_1)$  and  $\sigma_y \in \text{Candidates}(t_2)$
- $\sigma_x = f(\text{arguments}_f)$ , where  $\exists t, f(t_1, \dots, t_n) \in \text{TERMS}$  such that  $\sigma_x \in \text{Candidates}(t)$  and  $f(\text{arguments}_f) \in \{f(\hat{t}_1, \dots, \hat{t}_n) \mid f \text{ is a common functional symbol, } \hat{t}_i \in \text{Candidates}(t_i) \text{ for } 1 \leq i \leq n\}$
- $f_x(\text{arguments}_x) = f_y(\text{arguments}_y)$ , using the same definition for  $f(\text{arguments}_f)$  from previous item

---

**Algorithm 8** Conditional Replacement - Part 2

---

```

procedure CONDITIONAL REPLACEMENT( $z3::\text{expr const} \& x, z3::\text{expr const} \&$ 
 $y, H : \text{Horn clauses}, \mathcal{E} : \text{conditional equivalence relation}$ )
2:   if  $x$  is of the form  $f_x(t_1, \dots, t_{k_1})$  and  $y$  is of the form  $f_y(t'_1, \dots, t'_{k_2})$  then
      for  $\sigma_{f_x}$  in  $\text{Candidates}(f_x(t_1, \dots, t_{k_1}))$  do
4:       for  $\sigma_{f_y}$  in  $\text{Candidates}(f_y(t'_1, \dots, t'_{k_2}))$  do
             $\text{horn\_clause.add}(\text{ExtendedExplain}(f_x(t_1, \dots, t_{k_1}), \sigma_{f_x}, H, \mathcal{E}) +$ 
 $\text{ExtendedExplain}(f_y(t'_1, \dots, t'_{k_2}), \sigma_y, H, \mathcal{E}), \sigma_{f_x} = \sigma_{f_y})$ 
6:       end for
            for  $\text{arguments}_{f_y}$  in  $\text{CartesianProd}(\text{AllCandidates}(f_y(t'_1, \dots, t'_{k_2})))$  do
8:              $\text{horn\_clause.add}(\text{ExtendedExplain}(f_x(t_1, \dots, t_{k_1}), \sigma_{f_x}, H, \mathcal{E}) +$ 
 $\sum_{i=1}^{k_2} \text{ExtendedExplain}(t'_i, \text{arguments}_{f_y}[i], H, \mathcal{E}), \sigma_{f_x} = f_y(\text{arguments}_{f_y}))$ 
            end for
10:      end for
            for  $\text{arguments}_{f_x}$  in  $\text{CartesianProd}(\text{AllCandidates}(f_x(t_1, \dots, t_{k_1})))$  do
12:             for  $\sigma_{f_y}$  in  $\text{Candidates}(f_y(t'_1, \dots, t'_{k_2}))$  do
                   $\text{horn\_clause.add}(\sum_{i=1}^{k_1} \text{ExtendedExplain}(t_i, \text{arguments}_{f_x}[i], H, \mathcal{E})$ 
 $+ \text{ExtendedExplain}(f_y(t'_1, \dots, t'_{k_2}), \sigma_y, H, \mathcal{E}), f_x(\text{arguments}_{f_x}) = \sigma_{f_y})$ 
14:             end for
                  for  $\text{arguments}_{f_y}$  in  $\text{CartesianProd}(\text{AllCandidates}(f_y(t'_1, \dots, t'_{k_2})))$  do
16:                    $\text{horn\_clause.add}(\sum_{i=1}^{k_1} \text{ExtendedExplain}(t_i, \text{arguments}_{f_x}[i], H, \mathcal{E})$ 
 $+ \sum_{i=1}^{k_2} \text{ExtendedExplain}(t'_i, \text{arguments}_{f_y}[i], H, \mathcal{E}), f_x(\text{arguments}_{f_x}) =$ 
 $f_y(\text{arguments}_{f_y}))$ 
                  end for
18:             end for
            end if
20: end procedure

```

---

It is easy to see that the above equations are common if any of the involved sets are not empty. Therefore, the algorithm only produces common Horn clauses.  $\square$

In order to process Horn clauses obtained from Phase II of Kapur's algorithm, we apply the conditional replacement step for equations over the equations in the consequent of all the Horn clauses which *numargs* entries are equal to 0.

**Lemma 3.2.4.** *Let  $H$  be the set of Horn clauses obtained at the end of Phase II of Kapur's algorithm with *numargs* entry equal to 0 and  $\mathcal{E}'$  the associated conditional equivalence relation obtained at the end of the conditional elimination pro-*

*cedure.* If  $h \in H$  is of the form  $\bigwedge_i (a_i = b_i) \rightarrow c = d$ , then the Horn clauses  $\{\bigwedge_i \text{ExtendedExplain}(a_i, b_i, H, \mathcal{E}') \rightarrow h' \mid h' \in \text{ConditionalReplacement}(c, d, H, \mathcal{E}')\}$  are common Horn clauses.

*Proof.* We observe that if the *numargs* entry of a Horn clause  $h \in H$  is equal to 0, then all the equations in the antecedent of  $h$  belong to the conditional equivalence relation. Thus, there exists a proof in  $\mathcal{E}'$  for each of these equations. From the latter, the consequent of  $h$  also belongs to the conditional equivalence relation. Using Lemmas 3.2.2 and 3.2.3 we can conclude that the Horn clause in the statement is a common Horn clause.  $\square$

### 3.2.4 An additional application: checking membership of ground Horn clauses with Explanations

By adding common equations from the antecedent of the Horn clauses during the conditional elimination step we are able to extend an equivalence relation to a conditional equivalence relation that includes the set of consequences from the aforementioned common equations. In a more general setting, we can specify an arbitrary set of equations, or property, that are meant to enlarge the theory to a conditional equivalence relation for specific purposes. In order to reuse the *ExtendedExplain* procedure and associated lemmas, we replace the *common* predicate for any predicate of interest.

In particular, given an equivalence relation  $\mathcal{E}$ , a set of Horn clauses  $H$ , and Horn clause  $h$  of the form  $\bigwedge (a_i = b_i) \rightarrow c = d$ , we can find an explanation to the query  $H \models_{\mathcal{E}} h$  by conditionally extending  $H$  with the equations from the antecedent of  $h$  and then retrieving an explanation of the membership of the consequent of  $h$  in the conditional equivalence relation.

**Lemma 3.2.5.** *Let  $H$  be a set of Horn clauses,  $h$  a Horn clause of the form  $(\bigwedge_i a_i = b_i) \rightarrow c = d$ ,  $\mathcal{E}$  an equivalence relation between terms of a language in EUF, and  $\mathcal{E}'$  the asociated conditional equivalence relation obtained from extending  $\mathcal{E}$  with the equations in the antecedent of  $h$ .*

*If there exists an explanation  $\mathbf{p} := \{\mathbf{a}_1 = \mathbf{b}_1 \dots \mathbf{a}_n = \mathbf{b}_n\}$  for the query  $H \models_{\mathcal{E}'} c = d$ , then the explanation for the query  $H \models_{\mathcal{E}} (\bigwedge_i a_i = b_i) \rightarrow c = d$  is the subset of  $\mathbf{p}$  that does not contain equations from  $\bigcup_{i=1}^n \{a_i = b_i\}$ .*

*Proof.* If there exists an explanation  $\mathbf{p}$  for  $H \models_{\mathcal{E}'} c = d$  then there exists a proof tree with  $c = d$  as root node and the equations from  $\mathbf{p}$  as leaves. From lemma 3.2.1 and the definition 3.2.1 we notice that  $\mathbf{p}$  only contains equations derived from  $\mathcal{E}$  and the antecedent of  $h$ .

Using the *discharge rule* from basic propositional logic, i.e. 
$$\frac{\begin{array}{c} [A] \\ \vdots \\ B \end{array}}{A \rightarrow B},$$
 then we can obtain a proof tree with leaves only containing terms that do not belong to the antecedent of  $h$ .

Hence, the new proof tree will have  $(\bigwedge_{i=1}^m a'_i = b'_i) \rightarrow c = d$  as root node and derived equations from  $\mathcal{E}$  as leaves, where  $\{a'_i = b'_i\} \subseteq \text{antecedent}(h)$  for each  $1 \leq i \leq m$ . Thus  $\models \text{antecedent}(h) \rightarrow a'_i = b'_i$  for each  $1 \leq i \leq m$ . Let  $\mathbf{p}'$  be  $\mathbf{p} \setminus \bigcup_{i=1}^m \{a_i = b_i\}$ . Therefore, by transitivity, we have that  $\mathbf{p}'$  is an explanation in  $\mathcal{E}$  for  $H \models_{\mathcal{E}} h$ .

□

## 3.3 Evaluation

### 3.3.1 Detailed evaluation of examples

In this section we discuss in full detail the execution trace of the implementation of some examples.

#### Symbol elimination example

Let us consider the following example from [25]  $\alpha_1 = \{f(z_1, v) = s_1, f(z_2, v) = s_2, f(f(y_1, v), f(y_2, v)) = t\}$  with the set of symbols to eliminate  $U_1 = \{u\}$ . The implementation produces the following trace (slightly modified for presentation purposes) in order to compute the interpolant of  $\alpha_1; U_1$ :

```
Before conditionalEliminationEqs
Horn clauses produced
0. 0x5618a6c9c740 (Leader) (= c_y1 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y1 a_v))
1. 0x5618a6c9c410 (Leader) (= c_z2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_s2 c_t)
2. 0x5618a6c7c6a0 (Leader) (= c_z2 c_y1) -> (= c_s2 (c_f c_y1 a_v))
3. 0x5618a6c4db90 (Leader) (= c_y2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y2 a_v))
4. 0x5618a6c9cd30 (Leader) (= a_v (c_f c_y2 a_v)) and (= c_z1 (c_f c_y1 a_v)) -> (= c_s1 c_t)
5. 0x5618a6ca0c10 (Leader) (= c_z1 c_y2) -> (= c_s1 (c_f c_y2 a_v))
6. 0x5618a6ca0b50 (Leader) (= c_z1 c_y1) -> (= c_s1 (c_f c_y1 a_v))
7. 0x5618a6c9c6b0 (Leader) (= c_y1 c_y2) -> (= (c_f c_y1 a_v) (c_f c_y2 a_v))
8. 0x5618a6c7c710 (Leader) (= c_z2 c_y2) -> (= c_s2 (c_f c_y2 a_v))
9. 0x5618a6ca0960 (Leader) (= c_z1 c_z2) -> (= c_s1 c_s2)
Number of horn clauses: 10
Executing conditionalElimination
After conditionalEliminationEqs/Before conditionalEliminationHcs
Horn clauses produced
0. 0x5618a6ca4340 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s1 c_s2))
1. 0x5618a6ca0960 (Leader) (= c_z1 c_z2) -> (= c_s1 c_s2)
2. 0x5618a6ca0270 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s2 c_s2))
3. 0x5618a6c7c710 (Leader) (= c_z2 c_y2) -> (= c_s2 (c_f c_y2 a_v))
4. 0x5618a6c9c6b0 (Leader) (= c_y1 c_y2) -> (= (c_f c_y1 a_v) (c_f c_y2 a_v))
5. 0x5618a6ca5b80 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s2 c_s1))
6. 0x5618a6ca0b50 (Leader) (= c_z1 c_y1) -> (= c_s1 (c_f c_y1 a_v))
7. 0x5618a6ca0c10 (Leader) (= c_z1 c_y2) -> (= c_s1 (c_f c_y2 a_v))
8. 0x5618a6c4db90 (Leader) (= c_y2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y2 a_v))
9. 0x5618a6c7c6a0 (Leader) (= c_z2 c_y1) -> (= c_s2 (c_f c_y1 a_v))
10. 0x5618a6ca55a0 (Leader) (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
11. 0x5618a6c9c410 (Leader) (= c_z2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_s2 c_t)
12. 0x5618a6c9cd30 (Leader) (= a_v (c_f c_y2 a_v)) and (= c_z1 (c_f c_y1 a_v)) -> (= c_s1 c_t)
13. 0x5618a6ca62c0 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s1 c_s1))
```



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```

14. 0x5618a6c9c740 (Leader) (= c_y1 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y1 a_v))
Number of horn clauses: 15
Executing conditionalEliminationfor Horn clauses
After conditionalEliminationHcs
Horn clauses produced
0. 0x5618a6cabb10 (Leader) (= c_y1 c_y2) and (= c_z2 c_y1) and (= c_z1 c_y2) -> (= c_s1 c_s2)
1. 0x5618a6c9fac0 (Leader) (= c_y1 c_y2) and (= c_z1 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
2. 0x5618a6ca4340 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s1 c_s2))
3. 0x5618a6cab830 (Leader) (= c_z2 c_y1) and (= c_z1 c_y1) -> (= c_s1 c_s2)
4. 0x5618a6ca0960 (Leader) (= c_z1 c_z2) -> (= c_s1 c_s2)
5. 0x5618a6ca0270 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s2 c_s2))
6. 0x5618a6c7c710 (Leader) (= c_z2 c_y2) -> (= c_s2 (c_f c_y2 a_v))
7. 0x5618a6c9c6b0 (Leader) (= c_y1 c_y2) -> (= (c_f c_y1 a_v) (c_f c_y2 a_v))
8. 0x5618a6ca5b80 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s2 c_s1))
9. 0x5618a6ca0b50 (Leader) (= c_z1 c_y1) -> (= c_s1 (c_f c_y1 a_v))
10. 0x5618a6ca0c10 (Leader) (= c_z1 c_y2) -> (= c_s1 (c_f c_y2 a_v))
11. 0x5618a6c4db90 (Leader) (= c_y2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y2 a_v))
12. 0x5618a6ca65b0 (Leader) (= c_z1 c_y2) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
13. 0x5618a6c7c6a0 (Leader) (= c_z2 c_y1) -> (= c_s2 (c_f c_y1 a_v))
14. 0x5618a6ca55a0 (Leader) (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
15. 0x5618a6ca6300 (Leader) (= c_z2 c_y2) and (= c_z1 c_y2) -> (= c_s1 c_s2)
16. 0x5618a6c9c410 (Leader) (= c_z2 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_s2 c_t)
17. 0x5618a6c9cd30 (Leader) (= a_v (c_f c_y2 a_v)) and (= c_z1 (c_f c_y1 a_v)) -> (= c_s1 c_t)
18. 0x5618a6ca62c0 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s1 c_s1))
19. 0x5618a6c9c740 (Leader) (= c_y1 (c_f c_y1 a_v)) and (= a_v (c_f c_y2 a_v)) -> (= c_t (c_f c_y1 a_v))
Number of horn clauses: 20
Horn clauses produced
0. 0x5618a6cabb10 (Not leader) (= c_y1 c_y2) and (= c_z2 c_y1) and (= c_z1 c_y2) -> (= c_s1 c_s2)
1. 0x5618a6c9fac0 (Not leader) (= c_y1 c_y2) and (= c_z1 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
2. 0x5618a6ca4340 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s1 c_s2))
3. 0x5618a6cab830 (Not leader) (= c_z2 c_y1) and (= c_z1 c_y1) -> (= c_s1 c_s2)
4. 0x5618a6ca0960 (Leader) (= c_z1 c_z2) -> (= c_s1 c_s2)
5. 0x5618a6ca0270 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_t (c_f c_s2 c_s2))
6. 0x5618a6ca5b80 (Leader) (= c_z2 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s2 c_s1))
7. 0x5618a6ca65b0 (Not leader) (= c_z1 c_y2) and (= c_z1 c_y2) and (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
8. 0x5618a6ca55a0 (Not leader) (= c_z1 c_y1) and (= c_z2 c_y1) -> (= c_s1 c_s2)
9. 0x5618a6ca6300 (Not leader) (= c_z2 c_y2) and (= c_z1 c_y2) -> (= c_s1 c_s2)
10. 0x5618a6ca62c0 (Leader) (= c_z1 c_y1) and (= c_z1 c_y2) -> (= c_t (c_f c_s1 c_s1))
Number of horn clauses: 11
Interpolant:
(ast-vector
  (=> (and (= z1 y1) (= z1 y2) (= z1 y1) (= z2 y1)) (= t (f s1 s2))))
  (=> (= z1 z2) (= s1 s2))
  (=> (and (= z2 y1) (= z1 y2) (= z1 y1) (= z2 y1)) (= t (f s2 s2))))
  (=> (and (= z2 y1) (= z1 y2)) (= t (f s2 s1))))
  (=> (and (= z1 y1) (= z1 y2)) (= t (f s1 s1))))

```

The interface offered by SMT solvers with interpolation features usually provide the conventional A-part, B-part format. In order to compare the results with the implementation two instances were tested so we can obtain interpolants from both systems:

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- Problem instance: A-part :  $\{f(z_1, v) = s_1, f(z_2, v) = s_2, f(f(y_1, v), f(y_2, v)) = t\}$ ; B-part :  $\{z_1 = z_2, s_1 \neq s_2\}$ 
  - Z3:  $(\text{or } (= s2 s1) (\text{not } (= z2 z1)))$
  - Mathsat:  $(\text{not } (\text{and } (= z1 z2) (\text{not } (= s1 s2))))$
  - Our implemenation:  $(\rightarrow (= z1 z2) (= s1 s2))$
- Problem instance: A-part :  $\{f(z_1, v) = s_1, f(z_2, v) = s_2, f(f(y_1, v), f(y_2, v)) = t\}$ ; B-part :  $\{z_1 = y_1, z_1 = y_2, f(s_1, s_1) \neq t\}$ 
  - Z3:  $(\text{or } (= (f s1 s1) t) (\text{not } (= z1 y1)) (\text{not } (= y2 y1)))$
  - Mathsat:  $(\text{not } (\text{and } (\text{not } (= t (f s1 s1))) (\text{and } (= z1 y1) (= z1 y2))))$
  - Our implemenation:  $(\rightarrow (\text{and } (= y1 y2) (= z1 y2) (= z1 y2)) (= t (f s1 s1))))$

Clearly, the interpolant obtained by just eliminating the symbol  $v$  implies the outputs produced by Z3 and Mathsat for the previous problem instances.

### Simple example with disequality

Let us consider another example from [25]  $\alpha_2 = \{f(x_1) \neq f(x_2)\}$  with the set of symbols to eliminate  $U_2 = \{f\}$ . The implementation produces the following trace for  $\alpha_2; U_2$ :

```
Before conditionalEliminationEqs
Horn clauses produced
0. 0x5648882bc710 (Leader) (= c_x2 c_x1) -> (= (a_f c_x2) (a_f c_x1))
1. 0x5648882d7dd0 (Leader) (= (a_f c_x2) (a_f c_x1)) -> false
Number of horn clauses: 2
Executing conditionalElimination
After conditionalEliminationEqs/Before conditionalEliminationHcs
Horn clauses produced
0. 0x5648882bc710 (Leader) (= c_x2 c_x1) -> (= (a_f c_x2) (a_f c_x1))
1. 0x5648882d7dd0 (Leader) (= (a_f c_x2) (a_f c_x1)) -> false
Number of horn clauses: 2
```

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```

Executing conditionalEliminationfor Horn clauses
After conditionalEliminationHcs
Horn clauses produced
0. 0x5648882bc710 (Leader) (= c_x2 c_x1) -> (= (a_f c_x2) (a_f c_x1))
1. 0x5648882d7dd0 (Leader) (= (a_f c_x2) (a_f c_x1)) -> false
Number of horn clauses: 2
Horn clauses produced
0. 0x5648882deea0 (Leader) (= c_x2 c_x1) -> false
Number of horn clauses: 1
(ast-vector
  (=> (= x2 x1) false))

```

To compare our result with Z3 and Mathsat we included the B-part formula to be  $\{x_1 = x_2\}$ . The interpolants obtained by these systems were the same, which was  $(\text{not } (= x1 x2))$ .

#### 3.3.2 Performance comparison with iZ3 and MathSat

This section will discuss a parametrized problem which will allow to test our implementation and contrast the execution with other interpolant generation algorithms from Z3 and MathSat.

**Lemma 3.3.1.** *Let  $x$  be a constant and  $f$  an unary function in the EUF language. For every  $n \in \mathbb{N}$  the following conjunction is unsatisfiable:*

$$f^n(x) = f^{n+1}(x), f^2(x) = x, f(x) \neq x$$

where  $f^n(x)$  denotes the application of the function  $f$   $n$ -times to the constant  $x$ .

*Proof.* Applying the congruence and transitive rules to the equation  $f^2(x) = x$ , we can prove that for any number  $n$  the two following equations hold:  $x = f^{2n}(x)$  and  $f(x) = f^{2n+1}(x)$ .

At this point we distinguish two cases for  $n$ :

- Case  $n$  is even: Then  $\exists m \in \mathbb{N}. n = 2m$ . So by choosing  $n = m$  to the above sequence and by the axiom  $f^{2m}(x) = f^{2m+1}(x)$  we have that  $x = f(x)$ .

- Case  $n$  is odd: Similarly to the previous reasoning, we can also infer the equation  $f^{2m+1}(x) = f^{2m+2}(x)$  by congruence, so  $x = f(x)$ .

In both cases the equation  $x = f(x)$  reaches contradiction with the disequality  $x \neq f(x)$ . □

The interpolation pair for our performance test is  $(f^n(x) = f^{n+1}(x) \wedge f^2(x) = x \wedge f(a) \neq a, x = a)$  for fixed  $n \in \mathbb{N}$ . It is easy to see that the pair is inconsistent due to the lemma 3.3.1.

By lemma 3.3.1 we can also see that the formula  $x = a \rightarrow \perp$  is an interpolating formula for all fixed  $n$  in this parametrized problem. We designed this problem because it is easy to verify the correctness of output of the algorithms and to measure the time used by the algorithms for large values of  $n$ . We executed instances of this problem for values of  $n$  in the range  $\{1, \dots, 10000\}$  using a computer desktop equipped with an Intel i7-9700 @ 4.70 GHz and 16Gb of memory. The output produced by the interpolant generation algorithms were the expected formula. The following graph reports the time measured by the UNIX utility *times* of our implementation, iZ3, and the interpolation generation algorithm from Mathsat.

We noticed a quadratic behaviour in the graph for the time required by our implementation and a linear behaviour by iZ3 and Mathsat. However, according to the following polynomial fitting computation, we noticed that for the case of our implementation, a linear fitting obtained a smaller quadratic error compared to the quadratic fitting; for the case of the Mathsat results, a quadratic fitting obtained a smaller quadratic error.

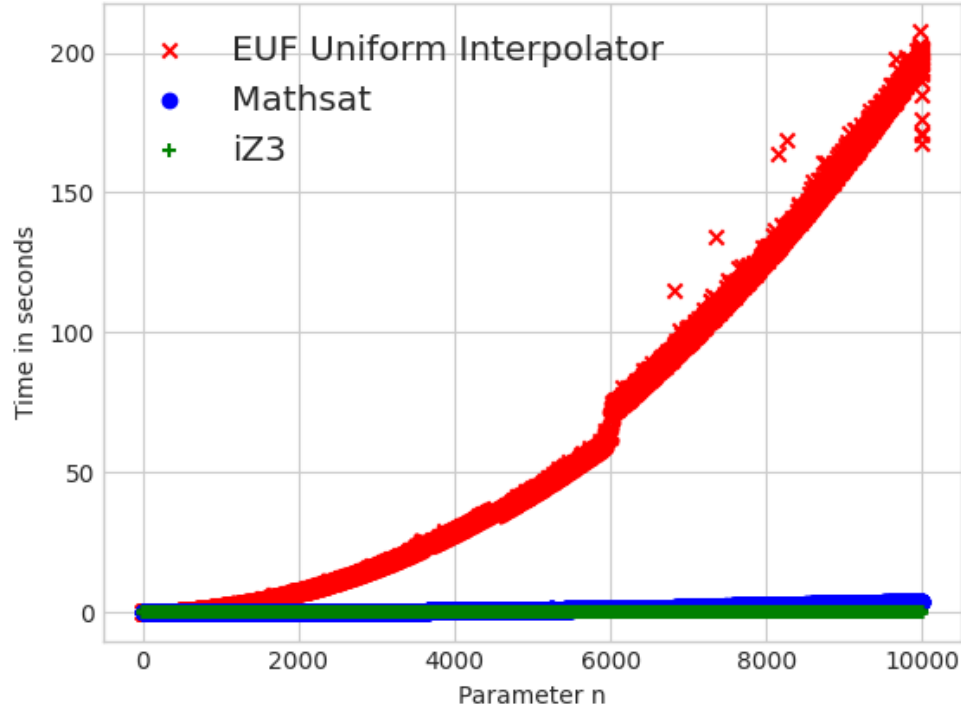


Figure 3.1: Performance comparison graph of EUF interpolant generation algorithms for parametrized problem from section 3.3.2

### 3.4 Conclusions

This chapter discussed the new approach for interpolant generation introduced by Prof. Kapur for the EUF theory. We proposed a modification of his algorithm and proved its correctness. The implementation and testing work confirms that the approach produces stronger interpolants compared to other interpolant generation algorithms (iZ3, Mathsat).

The performance comparison section indicates that of our implementation appears to be slower than iZ3 and Mathsat. This can be due the quadratic step by the

	Error of linear fitting	Error of quadratic fitting
iZ3	1.8723744548656934e-27	2.3456535579201815e-26
Mathsat	8.690566310635247e-23	7.576220133742597e-23
Our implementation	2.5622746654442585e-20	6.648236698726508e-20

Table 3.1: Table of polynomial fitting errors of degrees 1, 2 for the EUF experiments of section 3.3.2.

Phase II in Kapur’s algorithm and both iZ3 and Mathsat might be able to find the linear refutational proof discussed in 3.3.1.

On the other hand, it is important to highlight that the algorithm have different output specification since Prof. Kapur algorithm is able to find a uniform interpolant.

There are several places where the current implementation can be improved. Currently, the data structures for flattening introduce additional constants that might not be needed since the implementation used a recursive scan of the arguments of each expression without checking for term sharing. Moreover, the explanation mechanism might be improved/changed by the proof producing mechanism already available in Z3.

## Chapter 4

# Interpolation algorithm for UTVPI Formulas

This theory appears heavily in formal methods dealing with abstract domains introduced in [35]. The decision problem consists of checking the satisfiability of a particular fragment of  $LIA(\mathbb{Z})$ . The fragment consists on conjunctions of inequalities with at most two variables which integers coefficients are restricted to  $\{-1, 0, 1\}$ . Efficient algorithms are found in the literature for both the satisfiability problem [30] as well as for interpolation [9] of the theory. Eventhough this fragment looks severely constrained, it has been used for a range of applications where problems can be modelled using this particular kind of literals.

The algorithm in [25] follows a similar approach to the interpolation algorithm for EUF in the sense that the attention is given to one of the formulas in the interpolation pair <sup>1</sup> Other approaches towards interpolation follow graph-based algorithms. The idea combines the reduction of UTVPI formulas to difference logic [35] and a cycle detection of maximal size [9].

---

<sup>1</sup>This implementation uses the first formula of the pair.

## 4.1 Algorithm

The algorithm proposed [25] uses inference rules to close the relation and eliminate the uncommon variables from the A-part of the input formula. The rules are the following:

$$\frac{s_1x_1 + s_2x_2 \leq c \quad s_1 = s_2 \in \{-1, 0, 1\} \text{ and } x_1 = x_2 \in Vars}{s_1x_1 \leq \lfloor \frac{c}{2} \rfloor} \text{ Normalize}$$

$$\frac{s_1x_1 + s_2x_2 \leq c_1 \quad -s_2x_2 + s_3x_3 \leq c_2}{s_1x_1 + s_3x_3 \leq c_1 + c_2} \text{ Elim}$$

The algorithm normalizes the inequalities at the beginning as a preprocessing step and applies the Elim rule whenever it is possible until no more uncommon variables remain in the input formula. Hence, having an efficient representation of the inequalities and detecting matches (like pivots for resolution steps) is important for an efficient implementation. To achieve this goal we implemented to an encoding of the inequalities using natural numbers, an array of numbers indexed by the numeral representation of the inequalities which keeps track of the minimum bound of the encoded inequality, and a data structure to keep track of the signs of variables in the inequalities for efficient matching.

## 4.2 Implementation

In order to obtain a bijection between UTVPI inequalities and natural numbers, first we define an ordering on the inequalities and notice some invariants of the latter.

We encode the term  $\pm x_m \pm x_n$  using the point  $(\pm m, \pm n) \in \mathbb{Z}^2$ . Let *TermToPoint* be the map that  $\pm x_m \pm x_n \mapsto (\pm m, \pm n)$ . The variable  $x_0$  is a *dummy variable* that acts as a place holder for 0. Additionally, we will restrict the terms/pairs such that



## Chapter 4. Interpolation algorithm for UTVPI Formulas

the absolute value of the first index variable is strictly greater than the absolute value of the second index variable <sup>2</sup> since addition is commutative, except for the point  $(0, 0)$  which encodes the inequality with no variables.

We define the following orderings relevant for the terms of the form  $\pm x_m \pm x_n$ .

**Definition 4.2.1.** *Let  $\succ_m$  be an ordering on the integers such that  $a \succ_m b$  if and only  $|a| > |b|$  or  $(|a| = |b| \text{ and } a > b)$  where  $>$  is the standard ordering on integers.*

*Let  $\succ_p$  be an ordering on pair of integers such that  $(m_1, n_1) \succ_p (m_2, n_2)$  if and only if  $m_1 \succ_m m_2$  or  $(m_1 = m_2 \text{ and } n_1 \succ_m n_2)$ .*

*Let  $\succ_t$  be an ordering on terms of the form  $\pm x_m \pm x_n$  such that  $t_1 \succ_t t_2$  if and only if  $\text{TermToPoint}(t_1) \succ_p \text{TermToPoint}(t_2)$*

**Example 4.2.1.1.** *The first 32 elements (in ascending order w.r.t.  $\succ_t$ ) of UTVPI inequalities <sup>3</sup> are the following:*

$$\begin{aligned}
 &x_0 + x_0 \leq b_0 \\
 &-x_1 + x_0 \leq b_1, x_1 + x_0 \leq b_2 \\
 &-x_2 + x_0 \leq b_3, -x_2 - x_1 \leq b_4, -x_2 + x_1 \leq b_5, x_2 + x_0 \leq b_6, x_2 - x_1 \leq b_7, x_2 + x_1 \leq b_8 \\
 &-x_3 + x_0 \leq b_9, -x_3 - x_1 \leq b_{10}, -x_3 + x_1 \leq b_{11}, -x_3 - x_2 \leq b_{12}, -x_3 + x_2 \leq b_{13}, \\
 &x_3 + x_0 \leq b_{14}, x_3 - x_1 \leq b_{15}, x_3 + x_1 \leq b_{16}, x_3 - x_2 \leq b_{17}, x_3 + x_2 \leq b_{18}
 \end{aligned}$$

From the example, we notice that we can group/order the inequalities by groups using the first index. The first element of the  $i^{th}$  group corresponds to the  $2(i-1)^2 + 1$  element in the  $\succ_t$  order. The observation follows from an inductive argument since there are  $2(1 + 2(i-1))$  elements in the  $i^{th}$  group. It is also straight forward to

---

<sup>2</sup>This condition also avoids the problem of keeping track of variable duplication in the inequality.

<sup>3</sup>For readability purposes we include the bound for the UTVPI term

find the position of the first element in the second half of any group. With the above information it is possible to find the map between UTVPI terms and naturals numbers.

This bijection allows us to implement a data structure based on a vector of integers extended with  $\pm\infty$  which encodes the upper bounds for the  $i^{th}$  inequality present in the input formula. For initialization purpose all the entries in this vector are set to  $\infty$  and these values are updated accordingly to keep track to the minimum possible value for the inequality after the application of the inference rules mentioned at the introduction of the section.

## 4.3 Evaluation

### 4.3.1 Detailed evaluation of examples

Let us consider the following example:  $\alpha_1 = \{x_1 - x_2 \geq -4, -x_2 - x_3 \geq 5, x_2 + x_6 \geq 4, x_2 + x_5 \geq -3\}; \beta_1 = \{-x_1 + x_3 \geq -2, -x_4 - x_6 \geq 0, -x_5 + x_4 \geq 0\}$ . Our implementation produces the following trace:

```
Processing
(+ (- c_x1) a_x2)
Updating structure with
x_2 - x_1 <= 4
Processing
(+ a_x2 c_x3)
Updating structure with
x_3 + x_2 <= -5
Processing
(- (- a_x2) c_x6)
Updating structure with
- x_4 - x_2 <= -4
Processing
(- (- a_x2) c_x5)
Updating structure with
- x_5 - x_2 <= 3
Removing this var: x_0
Removing this var: x_1
Removing this var: x_2
```

---

**Algorithm 9** UTVPI constructor

---

```

procedure UTVPI CONSTRUCTOR(position : integer)
2:   coefficient1 = 0
   coefficient2 = 0
4:   varindex1 = 0
   varindex2 = 0
6:   if position = 0 then
       return
8:   end if
   varindex1 =  $\sqrt{\frac{position-1}{2}} + 1$ 
10:  initial_group_position =  $2 * (varindex1 - 1)^2 + 1$ 
   half_size_group =  $2 * varindex1 - 1$ 
12:  if position  $\leq$  initial_group_position + half_size_group then
       coefficient1 = -1
14:     if position = initial_group_position then
           coefficient2 = 0
16:       varindex2 = 0
           return
18:     end if
       separation = position - initial_group_position
20:       varindex2 =  $\frac{separation-1}{2} + 1$ 
       if mod separation 2 = 0 then
22:         coefficient2 = 1
           return
24:       end if
       coefficient2 = -1
26:       return
   end if
28:   coefficient1 = 1
   if position = initial_group_position + half_size_group + 1 then
30:     coefficient2 = 0
       varindex2 = 0
32:     return
   end if
34:   separation = position - initial_group_position - half_size_group - 1
   varindex2 =  $\frac{separation-1}{2} + 1$ 
36:   if mod separation 2 = 0 then
       coefficient2 = 1
38:     return
   end if
40:   coefficient2 = -1
       return
42: end procedure

```

---

```

Reducing x_2 - x_1 and - x_4 - x_2
Result: - x_4 - x_1
Reducing x_2 - x_1 and - x_5 - x_2
Result: - x_5 - x_1
Reducing x_3 + x_2 and - x_4 - x_2
Result: - x_4 + x_3

```

---

**Algorithm 10** UTVPI position

---

```

procedure UTVPI POSITION(  $s_1x_{m_1} + s_2x_{m_2}$  : UTVPI term)
2:   initial_group_position =  $2 * (m_1 - 1)^2 + 1$ 
   if  $s_1 = -1$  then
4:     sign_a_offset = 0
   else
6:     if  $s_1 = 0$  then
       return 0
8:     else
       if  $s_1 = 1$  then
10:      sign_a_offset =  $2 * (m_1 - 1) + 1$ 
       end if
12:    end if
   end if
14:   if  $s_2 = -1$  then
       sign_b_offset =  $1 + 2 * (m_2 - 1)$ 
16:   else
       if  $s_2 = 0$  then
18:         sign_b_offset = 0
       else
20:         if  $s_2 = 1$  then
           sign_b_offset =  $2 * m_2$ 
22:         end if
       end if
24:   end if
   return initial_group_position + sign_a_offset + sign_b_offset
26: end procedure

```

---

```

Reducing x_3 + x_2 and - x_5 - x_2
Result: - x_5 + x_3
Removing this var: x_3
Removing this var: x_4
Removing this var: x_5
Interpolant:
(ast-vector
  (<= (- (- x6) x1) 0)
  (<= (+ (- x6) x3) (- 9))
  (<= (- (- x5) x1) 7)
  (<= (+ (- x5) x3) (- 2)))

```

The Z3 SMT solver provides mechanisms to modified the arithmetic engine and several plugins to specialize algorithm for specific theories. It contains a proper specialization to work with Utpvi queries for satisfiability checking. However, the interpolation APIs do not include mechanisms to specialize the interpolation algorithm for Utpvi formulas. Thus, the interpolant obtained by Z3 for the above problem is :  $(\text{and } (\leq 9 (+ (* (- 1) x3) (* 2 x6) x1)) (\leq (- 5) (+ (* (- 1) x3) (* 2 x5) x1)))$ . We can

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notice by the coefficients in the result that the interpolant is not a Utpvi formula, thus Z3 must have reduced the problem to linear integer arithmetic.

The result obtained by Mathsat is  $(\text{and } (\leq (-5) (+ x1 (+ (* (-1) x3) (* 2 x5)))) (\leq 9 (+ x1 (+ (* (-1) x3) (* 2 x6))))))$  which is the same as Z3 modulo the commutativity of the additions in the expression. Despite this difference, the following query to Z3 verifies that the interpolation produced by our implementation implies the interpolation produced by the SMT solver above mentioned; at the same time, the interpolant produced by the SMT solver does not imply our interpolant.

```
(declare-fun x1 () Int)
(declare-fun x3 () Int)
(declare-fun x5 () Int)
(declare-fun x6 () Int)

(define-fun implementation_interpolant () Bool
  (and
    (<= (- (- x6) x1) 0)
    (<= (+ (- x6) x3) (- 9))
    (<= (- (- x5) x1) 7)
    (<= (+ (- x5) x3) (- 2))
  )
)

(define-fun z3_interpolant () Bool
  (and
    (<= 9 (+ (* (- 1) x3) (* 2 x6) x1))
    (<= (- 5) (+ (* (- 1) x3) (* 2 x5) x1))
  )
)

(push)
(assert (not (implies z3_interpolant implementation_interpolant)))
(check-sat) ;; This check returns sat, i.e. z3_interpolant does not imply implementation_interpolant
(pop)
(push)
(assert (not (implies implementation_interpolant z3_interpolant)))
(check-sat) ;; This check returns unsat, i.e. implementation_interpolant implies z3_interpolant
(pop)
```

For the next example, let us consider  $\alpha_2 = \{-x_2 - x_1 + 3 \geq 0, x_1 + x_3 + 1 \geq 0, -x_3 - x_4 - 6 \geq 0, x_5 + x_4 + 1 \geq 0\}$ ;  $\beta_2 = \{x_2 + x_3 + 3 \geq 0, x_6 - x_5 - 1 \geq 0, x_4 - x_6 + 4 \geq 0\}$ . Our implementation produces the following trace:

Processing

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```

(+ c_x2 a_x1)
Updating structure with
x_2 + x_1 <= 3
Processing
(- (- c_x3) a_x1)
Updating structure with
- x_3 - x_2 <= 1
Processing
(+ c_x4 c_x3)
Updating structure with
x_4 + x_3 <= -6
Processing
(- (- c_x5) c_x4)
Updating structure with
- x_5 - x_4 <= 1
Removing this var: x_0
Removing this var: x_1
Removing this var: x_2
Reducing x_2 + x_1 and - x_3 - x_2
Result: - x_3 + x_1
Removing this var: x_3
Removing this var: x_4
Removing this var: x_5
(ast-vector
  (<= (+ (- x3) x2) 4)
  (<= (+ x4 x3) (- 6))
  (<= (- (- x5) x4) 1))

```

The interpolant obtained by Z3 is  $(\text{and } (\leq (- 4) (+ x3 (* (- 1) x2))) (\leq (+ x3 x4) (- 6)) (\geq (+ x4 x5) (- 1)))$ ; and the interpolant produced by Mathsat is  $(\leq (+ x2 (+ x3 (+ x4 (* (- 1) x5)))) (- 7))$ . Using Z3, we were able to verify that:

- the interpolation obtained by our implementation is equivalent to the interpolant obtained by Z3.
- the interpolant obtained by our implementation implies the interpolant obtained by Mathsat, but the converse does not hold.

```

(declare-fun x1 () Int)
(declare-fun x2 () Int)
(declare-fun x3 () Int)
(declare-fun x4 () Int)
(declare-fun x5 () Int)
(declare-fun x6 () Int)

(define-fun implementation_interpolant () Bool
  (and
    (<= (+ (- x3) x2) 4)

```

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```

(<= (+ x4 x3) (- 6))
(<= (- (- x5) x4) 1))
)

(define-fun z3_interpolant () Bool
  (and
    (<= (- 4) (+ x3 (* (- 1) x2)))
    (<= (+ x3 x4) (- 6))
    (>= (+ x4 x5) (- 1)))
  )

(define-fun mathsat_interpolant () Bool
  (<= (+ x2 (+ x3 (+ x4 (* (- 1) x5)))) (- 7))
  )

(push)
(assert (not (iff z3_interpolant implementation_interpolant)))
(check-sat) ;; This check returns unsat, i.e. z3_interpolant is equivalent to implementation_interpolant
(pop)

(push)
(assert (not (implies mathsat_interpolant implementation_interpolant)))
(check-sat) ;; This check returns sat, i.e. mathsat_interpolant does not imply implementation_interpolant
(pop)

(push)
(assert (not (implies implementation_interpolant mathsat_interpolant)))
(check-sat) ;; This check returns unsat, i.e. implementation_interpolant implies mathsat_interpolant
(pop)

```

### 4.3.2 Performance comparison with iZ3 and MathSat

This section discusses a parametrized problem in the UTVPI theory and compares the times required by our implementation, iZ3, and Mathsats.

**Lemma 4.3.1.** *Let  $x_i$  for  $1 \leq i \leq n$  where  $n \in \mathbb{N}$  is fixed. The following conjunction of inequalities is unsatisfiable:*

$$x_1 + x_2 \leq 1 \wedge \bigwedge_{i=2}^{n-1} (x_{i+1} - x_i \leq 1) \wedge x_1 - x_n \leq 1 \wedge x_1 > \lceil n/2 \rceil$$

*Proof.* Using the inequalities  $x_1 + x_2 \leq 1 \wedge \bigwedge_{i=2}^{n-1} (x_{i+1} - x_i \leq 1)$  we can conclude that  $2x_1 \leq n$ . The latter contradicts the last inequality  $x_1 > \lceil n/2 \rceil$ .  $\square$

The interpolation pair for our performance test is  $(x_1 + x_2 \leq 1 \wedge \bigwedge_{i=2}^{n-1} (x_{i+1} - x_i \leq$

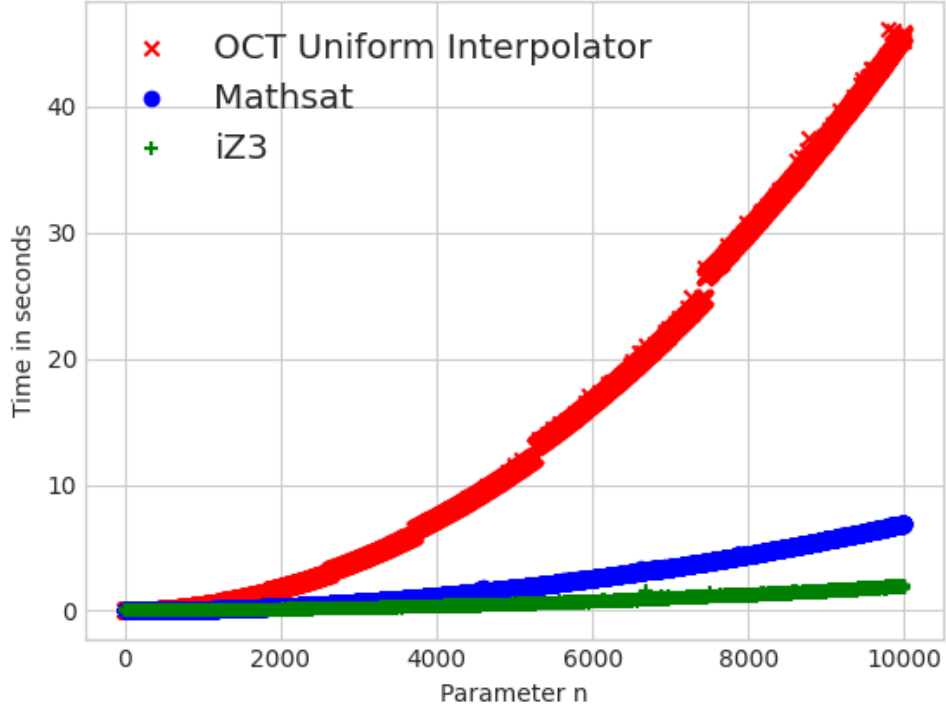


Figure 4.1: Performance comparison graph of Oct interpolant generation algorithms for parametrized problem from section 4.3.2

1)  $\wedge x_1 - x_n \leq 1, x_1 > \lceil n/2 \rceil$ ) for fixed  $n \in \mathbb{N}$ . It is easy to see that the pair is inconsistent due to the lemma 4.3.1.

By lemma 4.3.1 we can also see that the formula  $x_1 \leq \lceil n/2 \rceil$  is an interpolating formula for all fixed  $n$  in this parametrized problem. Similarly to the parametric problem in 3.3.2, we designed this problem because it is easy to verify the correctness of output of the algorithms and to measure the time for various values of  $n$ . We executed instances of this problem for values of  $n$  in the range  $\{1, \dots, 10000\}$  using the same setup used in the performance comparison for the parametrized problem discussed in 3.3.1. The follow graph shows similar results as in 3.3.1.



We noticed a quadratic behaviour in the graph for the time required by all the algorithms. Indeed, the latter was validated since a quadratic fitting has a smaller quadratic error compared to a linear fitting in all the implementations as show in the following table:

	Error of linear fitting	Error of quadratic fitting
iZ3	3.2311742677852644e-25	1.269128969536517e-23
Mathsat	2.3207586060940883e-22	1.1983240884566356e-22
Our implementation	2.9778502051908996e-21	1.455131710005814e-20

Table 4.1: Table of polynomial fitting errors of degrees 1, 2 for the Octagon experiments of section 4.3.2.

## 4.4 Conclusions

This chapter presented the approach by Prof. Kapur for computing the uniform interpolant of an inconsistent pair of formulas in the UTVPI theory. Several examples and testing of the implementation indicated that the output produced by Kapur's algorithm produces a stronger interpolant than the othe ones produced by iZ3 and Mathsat.

The performance comparison showed a closer difference in times compared to the performance comparison obtained for the EUF theory. It remains unclear why the performance comparison data have a better quadratic fit when a linear proof of the instance problem exists for all the algorithms. It will be interesting to explore the reasons of the latter in the current implementation for further improvements.

In order to improve the current implementation it will be interesting to explore the use of heuristics to determine the order in which the uncommon variable should be eliminated. The current implementation uses a lexicographic approach based on the

#### *Chapter 4. Interpolation algorithm for UTVPI Formulas*

occurrence of the symbols in the input formula. Different heuristic can consider the number of occurrences of the uncommon symbols in different inequalities, a topological order based on the dependency of the linear combinations of the inequalities, etc.

## Chapter 5

# Interpolation algorithm for the theory combination of EUF and UTVPI

Theory combination techniques involve reusing the algorithms for verification problems of some of the theories involved by either purifying terms over the theories, or reducing the original problem into a base theory, or a combination of these two. Usually the approaches following the first approach aforementioned rely on a Nelson-Oppen framework [46, 4, 6]

Despite possibly more efficient approaches using model-based theory combination techniques<sup>1</sup> [4], the latter require operations which many SMT solvers does not provide proper API to implement these. Hence, the approach used in the thesis work follows [46] since it does not require extensive modification to the decision procedures used. Many of the necessary modifications were implemented on top of

---

<sup>1</sup>An interesting property of model-based theory combination is that it is not needed to propagate disjunctions, even if the theories are not convex. The reason for this is because the *disjunctiveness* of the problem is handled *lazily* by the SAT solver.

Z3 and PicoSAT/TraceCheck.

The propagation of equalities and disjunction of equalities requires the additional step to split these formulas into the correct part of the interpolant pair. Among the major contributions of [46] was the introduction of the class of equality interpolating theories, which facilitates the splitting of a propagated formula since it is clear that A-local/B-local terms should be included in the A-part/B-part of the interpolant pair respectively.

**Definition 5.0.1.** *A theory  $\mathcal{T}$  is equality interpolating if for every  $a$  (A-local term) and  $b$  (B-local term) and  $\psi$  a formula in  $\mathcal{T}$  such that  $\psi \models_{\mathcal{T}} a = b$ , then there exists a term  $t$  in  $\mathcal{T}$  (called interpolating term) such that  $\psi \models_{\mathcal{T}} a = t$  and  $\psi \models_{\mathcal{T}} b = t$*

The relevance of the existence of the interpolating term for a deduced AB-mixed equality becomes relevant in the context of splitting a formula into a suitable A-part and B-part respectively. If  $t$

Deciding where to include AB-common terms to either the A-part or the B-part of the interpolation pair affects the final result since the interpolant will be *closer* to the A-part or to the B-part respectively. The original paper includes AB-common terms to the B-part. However, the implementation work includes AB-common terms to the A-part since the interpolation algorithms implemented focus the attention on this part.

## 5.1 Algorithm

The algorithm implements a Nelson-Oppen framework. In order to integrate interpolants to the latter, the algorithm includes a *partial interpolant* every time a disjunction of equalities (conflict clause) is propagated. These *partial interpolants*

are computed from an unsatisfiability proof obtained by including the negation of the disjunction to the formula using the following definition:

**Definition 5.1.1.** [46] Let  $\langle A, B \rangle$  be a pair of clause sets such that  $A \wedge B \vdash \perp$ . Let  $\mathcal{T}$  be a proof of unsatisfiability of  $A \wedge B$ . The propositional formula  $p(c)$  for a clause  $c$  in  $\mathcal{T}$  is defined by induction on the proof structure:

- if  $c$  is one of the input clauses then
  - if  $c \in A$ , then  $p(c) := \perp$
  - if  $c \in B$ , then  $p(c) := \top$
- otherwise,  $c$  is a result of resolution, i.e.  $\frac{c_1[= x \vee c'_1] \quad c_2[= \neg x \vee c'_2]}{c[= c'_1 \vee c'_2]}$ 
  - if  $x \in A$  and  $x \notin B$  ( $x$  is  $A$ -local), then  $p(c) := p(c_1) \vee p(c_2)$
  - if  $x \notin A$  and  $x \in B$  ( $x$  is  $B$ -local), then  $p(c) := p(c_1) \wedge p(c_2)$
  - otherwise ( $x$  is  $AB$ -common), then  $p(c) := (x \vee p(c_1)) \wedge (\neg x \vee p(c_2))$

## 5.2 Implementation

The implementation maintains a map data structure that keeps track of the *partial interpolants*. This ensures that the base case for the above formula  $p(c)$  is replaced by previous clauses as required in [46].

Since introducing negations is necessary to compute partial interpolants, we noticed the following interaction with the theories involved in the thesis work:

- EUF case: negations of literals in this theory are just disequalities, which the interpolant algorithm implemented handles as Horn clauses with a false head term.

- UTVPI case: negations of literals in this theory are either disequalities or strict inequalities. The disequalities are purified and appended to the EUF component, and the strict inequalities are appended as a disjunction of two non-strict inequalities following the axioms of number theory with addition and a total order relation [19].

The main loop of the procedure is shown below:

---

**Algorithm 11** Nelson-Oppen Propagation

---

```

procedure NELSON-OPPEN PROPAGATION ( z3::expr_vector const & part_A,
z3::expr_vector const & part_B )
2:    $T_1, T_2 = Purify(part\_A, part\_B)$ 
   DisjunctionEqualitiesIterator  $\psi()$ 
4:    $\psi.init()$ 
   while true do
6:     if  $T_1 \models_{EUF} \perp$  then
       return  $T_1$ 
8:     end if
     if  $T_2 \models_{UTVPI} \perp$  then
10:      return  $T_2$ 
     end if
12:    if  $T_1 \models_{EUF} \psi.current()$  then
      if  $T_2 \models_{UTVPI} \psi.current()$  then
14:        continue
      else
16:        append  $\psi.current()$  to  $T_2$ 
         $\psi.init()$ 
18:      end if
     else
20:      if  $T_2 \models_{UTVPI} \psi.current()$  then
        continue
22:      else
        append  $\psi.current()$  to  $T_1$ 
24:         $\psi.init()$ 
      end if
26:    end if
    UpdatePartialInterpolant( $\psi.current()$ )
28:     $\psi.next()$ 
   end while
30: end procedure

```

---

### 5.3 Octagonalization of input formula

This section discusses a minor issue that occurs after purifying an input formula from the combined language of any theory with UTVPI. It is trivial to see that, in the language of UTVPI, any linear combination of more than 2 variables does not belong to the UTVPI language by definition. On the other hand, purification will introduce equations with an additional fresh variable, so some valid input formulas in the UTVPI language might not be recognized as such.

**Example 5.3.0.1.** *Let us consider the following A-part input formula for the combined theory of EUF and UTVPI:  $f(x) + g(y, 29) \leq 120$ .*

*The formula introduced by purification does not conform the UTVPI language:*  
 $\mathbf{a}_1 = g(y, 29), \mathbf{a}_2 = f(x), \mathbf{a}_3 = \mathbf{a}_1 + \mathbf{a}_2, \mathbf{a}_3 \leq 120$

In some cases, the purified formula can be modified by propagating/eliminating unnecessary constants introduced by purification. An equisatisfiable formula to the above shown in example 5.3.0.1:  $\mathbf{a}_1 = g(y, 29), \mathbf{a}_2 = f(x), \mathbf{a}_1 + \mathbf{a}_2 \leq 120$ .

**Definition 5.3.1.** *An input formula in the combined language of theories including the UTVPI theory is octagonizable if the UTVPI component of the purified formula conforms the restrictions of the UTVPI language.*

**Lemma 5.3.1.** *An input formula  $\varphi$  in the combined language of EUF and UTVPI is octanizable if and only if for every subterm of  $\varphi$ , the purified equation introduced by purification containing three constants can be replaced by an equivalent equation with only two constants.*

*Proof.* The left to right implication is straightforward to prove since there are no equalities with three constants and the equalities with at most two constants can be

replaced by two inequalities with at most two constants, thus the formula conforms the language constraint in the UTVPI theory.

Let us prove the right to left implication. Since the EUF language does not contain addition nor subtraction, all the equalities produced by purification will contain at most three constants. By hypothesis, these equations can be replaced by equations with at most two constants, which can be replaced by two inequalities respectively by the UTVPI theory. Hence the UTVPI part of the purified output in additions with the latter inequalities conform the UTVPI language constraint.  $\square$

Our implementation propagates constants in the integer domain and eliminates unnecessary constants using the Z3 tactic called *solver-eqs*. The tactic *solver-eqs* eliminates variables using Gaussian elimination [15]. This transformation happens before the formula is applied to the UTVPI interpolant algorithm.

## 5.4 Evaluation

### 5.4.1 Detailed evaluation of a complete example

#### A simple example

Let us consider the following example:  $\alpha = \{f(x_1) = 0, x_1 = a, y_1 \leq a\}$ ;  $\beta = \{x_1 \leq b, y_1 = b, f(y_1) \neq 0\}$ . The implementation produces the following interpolant <sup>2</sup>:

```
(let ((a!1 (and (<= (+ (* (- 1) x1) y1) 0) (<= (+ x1 (* (- 1) y1)) (- 1)))))
      (let ((a!2 (or (= x1 y1) a!1 (<= (+ (* (- 1) x1) y1) (- 1)))))
            (let ((a!3 (or (not (= (f x1) 0)) (= 0 (f x1)) (and a!2 (not (= x1 y1)))))
                  (a!5 (or (= x1 y1) (and a!2 (not (= x1 y1)))))
                (let ((a!4 (or (not (= x1 y1)) (= (f y1) 0) (and (= (f x1) 0) a!3)))
                      (and a!4 a!5))))))
```

---

<sup>2</sup>A trace of the execution can be found at the Appendix section since it is considerably large to include it in this section



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The following SMT query verifies that the previous result obtained is an interpolant of the input formula:

```
(declare-fun x1 () Int)
(declare-fun x2 () Int)
(declare-fun x3 () Int)
(declare-fun y1 () Int)
(declare-fun y2 () Int)
(declare-fun y3 () Int)
(declare-fun a () Int)
(declare-fun b () Int)
(declare-fun f (Int) Int)
(declare-fun g (Int) Int)

(define-fun part_a () Bool
  (and
    (= (f x1) 0)
    (= x1 a)
    (<= y1 a)
  ))

(define-fun part_b () Bool
  (and
    (<= x1 b)
    (= y1 b)
    (distinct (f y1) 0)
  ))

(define-fun my_interpolant () Bool

  (let ((a!1 (and (<= (+ (* (- 1) x1) y1) 0) (<= (+ x1 (* (- 1) y1)) (- 1)))))
    (let ((a!2 (or (= x1 y1) a!1 (<= (+ (* (- 1) x1) y1) (- 1)))))
      (let ((a!3 (or (not (= (f x1) 0)) (= 0 (f x1)) (and a!2 (not (= x1 y1)))))
        (a!5 (or (= x1 y1) (and a!2 (not (= x1 y1)))))
      )
      (let ((a!4 (or (not (= x1 y1)) (= (f y1) 0) (and (= (f x1) 0) a!3))))
        (and a!4 a!5))))
  )

(push)
;; This returns unsat, which verifies that the input
;; is an inconsistent pair of formulas
(assert (and part_a part_b))
(check-sat)
(pop)
(push)
;; This returns unsat, which verifies that the output
;; is implied by the A-part
(assert (not (implies part_a my_interpolant)))
(check-sat)
(pop)
(push)
```

## Chapter 5. Interpolation algorithm for the theory combination of EUF and UTVPI

```
;; This returns unsat, which verifies that the output
;; is inconsistent with the B-part
(assert (and my_interpolant part_b))
(check-sat)
(pop)
```

For this example these are the interpolants reported by Z3 and Mathsats respectively:

- Z3 interpolant:  $(\text{and } (\geq (+ x1 (* (- 1) y1)) 0) (= (f x1) 0))$
- Mathsats interpolant:  $(\text{or } (= (f y1) 0) (\leq 1 (+ x1 (* (- 1) y1))))$

The following SMT query verifies the strength relation between the interpolants produced:

```
(declare-fun x1 () Int)
(declare-fun x2 () Int)
(declare-fun x3 () Int)
(declare-fun y1 () Int)
(declare-fun y2 () Int)
(declare-fun y3 () Int)
(declare-fun a () Int)
(declare-fun b () Int)
(declare-fun f (Int) Int)
(declare-fun g (Int) Int)

(define-fun my_interpolant () Bool

  (let ((a!1 (and (<= (+ (* (- 1) x1) y1) 0) (<= (+ x1 (* (- 1) y1)) (- 1)))))
    (let ((a!2 (or (= x1 y1) a!1 (<= (+ (* (- 1) x1) y1) (- 1)))))
      (let ((a!3 (or (not (= (f x1) 0)) (= 0 (f x1)) (and a!2 (not (= x1 y1)))))
        (a!5 (or (= x1 y1) (and a!2 (not (= x1 y1)))))
        (let ((a!4 (or (not (= x1 y1)) (= (f y1) 0) (and (= (f x1) 0) a!3)))
          (and a!4 a!5))))))
  )

(define-fun z3_interpolant () Bool
  (and (>= (+ x1 (* (- 1) y1)) 0) (= (f x1) 0))
  )

(define-fun mathsats_interpolant () Bool
  (or (= (f y1) 0) (<= 1 (+ x1 (* (- 1) y1))))
  )

(push)
;; The following returns sat, which means
```

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```
;; that my_interpolant does not imply z3_interpolant
(assert (not (implies my_interpolant z3_interpolant)))
(check-sat)
(pop)
(push)
;; The following returns unsat, which means
;; that z3_interpolant implies my_interpolant
(assert (not (implies z3_interpolant my_interpolant)))
(check-sat)
(pop)
(push)
;; The following returns unsat, which means
;; that my_interpolant implies mathsat_interpolant
(assert (not (implies my_interpolant mathsat_interpolant)))
(check-sat)
(pop)
(push)
;; The following returns sat, which means
;; that mathsat_interpolant does not imply my_interpolant
(assert (not (implies mathsat_interpolant my_interpolant )))
(check-sat)
(pop)
(push)
;; The following returns unsat, which means
;; that z3_interpolant implies mathsat_interpolant
(assert (not (implies z3_interpolant mathsat_interpolant)))
(check-sat)
(pop)
(push)
;; The following returns sat, which means
;; that mathsat_interpolant does not imply z3_interpolant
(assert (not (implies mathsat_interpolant z3_interpolant )))
(check-sat)
(pop)
```

### 5.4.2 A parametrized example

Let us consider the following parametrized example:  $\alpha = \{1 \leq x, x \leq n\}; \beta = \{f(x) = a, f(1) \neq a, f(2) \neq a, \dots, f(n-1) \neq a, f(n) \neq a\}$  where  $n > 1, a$  are fixed integers. We can see that an interpolant for this parametrized input problem is  $x = 1 \vee x = 2 \vee \dots \vee x = n - 1 \vee x = n$ .

For the case  $n = 3, 4, 5$ , both the implementation, Z3, and Mathsats were able to compute the expected result discussed above. This particular test was useful to check disjunction of equalities propagation mechanism in the implementation work.

### 5.4.3 Performance evaluation

This section discusses a parametrized problem in the EUF + UTVPI theory and compares the times required by our implementation, iZ3, and Mathsat.

**Lemma 5.4.1.** *Let  $x, a$  be constants and  $f$  a unary function in the language of the combined theory of EUF and UTVPI. The following conjunction is unsatisfiable for any fixed  $n \in \mathbb{N}$ :*

$$1 \leq x \wedge x \leq n \wedge f(x) = a \wedge \bigwedge_{i=1}^n f(i) \neq a$$

*Proof.* Using the inequalities  $1 \leq x \wedge x \leq n$  we can conclude that  $x = 1 \vee x = 2 \vee \dots \vee x = n$ . since we are in the theory of UTVPI for integers. We can prove that  $x \neq i$  where  $i \in \{1, \dots, n\}$ . Suppose, by contradiction, that  $x = i$ , so by congruence we have that  $f(x) = f(i)$ , thus  $f(i) = a$ . This will contradict the conjunction  $\bigwedge_{i=1}^n f(i) \neq a$ . Repeating the above step we can prove that  $\bigwedge_{i=1}^n x \neq i$ . Using  $n$  resolution steps, the first disjunction will entail  $\perp$  with  $\bigwedge_{i=1}^n x \neq i$ .  $\square$

The interpolation pair for our performance test is  $(1 \leq x \wedge x \leq n, f(x) = a \wedge \bigwedge_{i=1}^n f(i) \neq a)$  for fixed  $n \in \mathbb{N}$ . It is easy to see that the pair is inconsistent due to the lemma 5.4.1.

The interpolant in this problem is pretty simple since the A-part is already a common conjunction in the language of the combined theory. For the latter, we tested two versions of our implementation, one that executes the algorithm proposed in this section and the second includes an additional check in the algorithm described in this chapter which returns as result the formula in the A-part if this formula is a conjunction of common symbols.

The first experiment tested values for  $n$  from  $\{2, \dots, 8\}$ . Since the second approach returns a result immediately after checking if the input is common or not, we

tested the parametrized problem for values in  $\{2, \dots, 2500\}$ . We show a table and a graph for the first and second experiments respectively.

Parameter n	iZ3	Mathsat	Our implementation
2	0.006	0.004	0.026
3	0.006	0.008	0.030
4	0.003	0.005	0.082
5	0.01	0.004	0.706
6	0.007	0.004	11.347
7	0.013	0.004	213.431
8	0.009	0.008	5355.216

Table 5.1: Time in seconds obtained for the first experiment.

## 5.5 Conclusions

This chapter showed the implementation of the interpolation combination algorithm in [46] and the necessary changes for the EUF and UTVPI algorithms presented in previous chapters.

The performance comparison shows a slower performance of the implementation with respect to iZ3 and Mathsat when the common check is not enable. The latter was done on purpose to see the behaviour of the implementation on the interpolation combination component. On the other hand, if the common check is enabled then our implementation is faster than the algorithms in iZ3 and Mathsat. The latter suggests that iZ3 and Mathsat might not check for common symbols globally and the computation to distinguish symbols is done locally.

Currently the implementation does not incrementally the consistency of the current formula in the main loop of the implementation 5.2. Changing the latter to

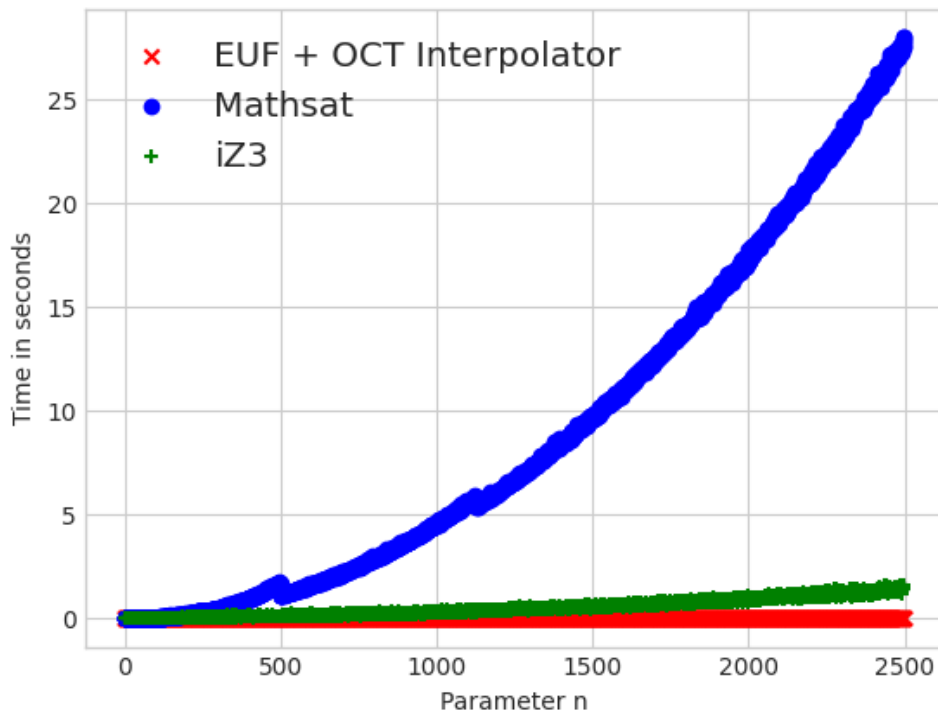


Figure 5.1: Performance comparison graph of EUF + UTVPI interpolant generation algorithms for parametrized problem from section 5.4.3

an incrementally approach will maintain lemmas so the computation of intermediate results will not be performed repetitively.

Additionally, recent approaches in theory combination might benefit the interpolation combination problem. In [4] it was proved that using a model-based combination approach might not require the propagation of disjunction since this happens internally at the core level of the model-based approach. This approach was not investigated in the thesis work since it required to modify the SMT solver at a non superficial level.

# Chapter 6

## Future Work

Regarding implementation work, there are several improvements that were not explored in the thesis work since many of them do not change the overall complexity of the implementation, but definitely these consume more resources. These particular improvements are the following:

- Improve hash function for terms: during testing the congruence closure algorithm, it was evident that dealing with a large number of terms, collisions happened to the point that some terms were merged since the signature function indicated to do so, but this should not happen. Nonetheless, for a typical verification problem the implementation will not find any problems.
- The space needed to encode curry nodes *can be significantly reduced* up to an order of  $\mathcal{O}(n)$ . The reason for this current allocation schema is that it exists a double bonding effect while performing curryfication of all the terms in the arguments of function applications.

Regarding potential extension of the systems, it will be interesting to explore a more specific interpolation combination approach as in [43]. The Nelson-Oppen framework

## *Chapter 6. Future Work*

seems to be adequate if the goal is to combine several theories or be as general as possible. The reason why the hierarchical reasoning approach was not implemented was because the authors did not make any claim about the possibility to use non-convex theories. Perhaps, this is not a real limitation of the approach.



# Appendix A: additional proofs

## 6.1 Theorems about Interpolants

### 6.1.1 Interpolants are closed under conjunction and disjunction

**Theorem 6.1.1.** *If  $I_1, I_2$  are interpolants of  $\langle A, B \rangle$ , then  $I_1 \wedge I_2, I_1 \vee I_2$  are also interpolants of  $\langle A, B \rangle$ .*

*Proof.* We notice that  $\text{vars}(I_1 \wedge I_2) \subseteq \text{vars}(A) \cap \text{vars}(B)$  and  $\text{vars}(I_1 \vee I_2) \subseteq \text{vars}(A) \cap \text{vars}(B)$ , otherwise  $I_1, I_2$  couldn't be interpolants.

Since  $I_1, I_2$  are interpolants, we have that  $A \vdash I_1$ ,  $B \wedge I_1 \vdash \perp$  and  $A \vdash I_2$ ,  $B \wedge I_2 \vdash \perp$

Here are formal proofs for  $A \vdash I_1 \wedge I_2$  and  $B \wedge I_1 \wedge I_2 \vdash \perp$ :

$$\begin{array}{ccc} A & A & B \wedge I_1 \wedge I_2 \\ \vdots & \vdots & \hline I_1 & I_2 & B \wedge I_1 \\ \hline I_1 \wedge I_2 & & \vdots \\ & & \perp \end{array}$$

Here are formal proofs for  $A \vdash I_1 \vee I_2$  and  $B \wedge (I_1 \vee I_2) \vdash \perp$ :

$$\begin{array}{c}
 \begin{array}{c} A \\ \vdots \\ I_1 \\ \hline I_1 \vee I_2 \end{array}
 \quad
 \begin{array}{c} B \wedge (I_1 \vee I_2) \\ \vdots * \\ (B \wedge I_1) \vee (B \wedge I_2) \end{array}
 \quad
 \begin{array}{c} \overline{B \wedge I_1} \\ \vdots \\ \perp \end{array}
 \quad
 \begin{array}{c} \overline{B \wedge I_2} \\ \vdots \\ \perp \end{array} \\
 \hline
 \perp
 \end{array}$$

Where  $*$  is any proof applying the distributivity property of the conjunction symbol over the disjunction symbol.

□

### 6.1.2 Interpolants distribute conjunctions over disjunctions in the A-part

**Theorem 6.1.2.** *Let  $\mathcal{F}_i$  be a set of formulas and  $I_i$  an interpolant for each  $\langle A \wedge \mathcal{F}_i, B \rangle$  respectively. Then  $\bigvee_i I_i$  is an interpolant for  $\langle A \wedge (\bigvee_i \mathcal{F}_i), B \rangle$*

*Proof.* Let  $A_i = A \wedge \mathcal{F}_i$  and  $\hat{A} = A \wedge (\bigvee_i \mathcal{F}_i)$ .

We see that  $\text{vars}(I_i) \subseteq \text{vars}(A_i) \subseteq \text{vars}(\hat{A})$ , hence  $\text{vars}(\bigvee_i I_i) = \bigcup_i \text{vars}(I_i) \subseteq \bigcup_i \text{vars}(A_i) \subseteq \text{vars}(\hat{A})$ . Similarly we can prove that  $\text{vars}(\bigvee_i I_i) \subseteq \text{vars}(B)$ . Thus,  $\text{vars}(\bigvee_i I_i) \subseteq \text{vars}(A \wedge (\bigvee_i \mathcal{F}_i)) \cap \text{vars}(B)$ .

To prove  $\hat{A} \vdash \bigvee_i I_i$ : We notice that  $\hat{A} = \bigvee_i A_i$ . From the latter and using the generalized version of the disjunction elimination rule in logic, i.e.

$$\begin{array}{c}
 \begin{array}{c} \overline{\alpha_1} \\ \vdots \\ I_1 \end{array}
 \quad
 \begin{array}{c} \overline{\alpha_n} \\ \vdots \\ I_n \end{array} \\
 \hline
 \alpha_1 \vee \dots \vee \alpha_n \quad \frac{I_1 \vee \dots \vee I_n}{I_1 \vee \dots \vee I_n} \quad \dots \quad \frac{I_1 \vee \dots \vee I_n}{I_1 \vee \dots \vee I_n} \\
 \hline
 I_1 \vee \dots \vee I_n
 \end{array}$$

and distributing disjunctions over conjunctions in  $\bigvee_i A_i$  the statement holds.

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To prove  $B \wedge \bigvee_i I_i \vdash \perp$ : Since each  $B \wedge I_i \vdash \perp$ , by using the generalized version of the disjunction elimination rule from above, the result holds.

Therefore,  $\bigvee_i I_i$  is an interpolant for  $\langle A \wedge (\bigvee_i \mathcal{F}_i), B \rangle$ .

□

# Appendix B: additional traces

## 6.2 Trace for example in theory combination chapter

```
Part a
EUF-component
(= (c_f c_x1) c_oct_1)
(= c_x1 a_a)
Octagon-component
(= 0 c_oct_1)
(<= c_y1 a_a)
Part b
EUF-component
(= c_y1 b_b)
(distinct (c_f c_y1) c_oct_2)
Octagon-component
(<= c_x1 b_b)
(= 0 c_oct_2)
Shared variables
(ast-vector
 c_x1
 a_a
 c_y1
 b_b
 c_oct_1
 c_oct_2)
(= 0 c_oct_1)
(<= c_y1 a_a)
(<= c_x1 b_b)
(= 0 c_oct_2)
(= (c_f c_x1) c_oct_1)
(= c_x1 a_a)
(= c_y1 b_b)
(not (= (c_f c_y1) c_oct_2))

Disjunction implied in EUF: (= c_x1 a_a)
Partial interpolant already computed
```

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Disjunction implied in EUF: (= c\_y1 b\_b)  
 Partial interpolant already computed

Disjunction implied in OCT: (= c\_x1 c\_y1)  
 Clause Id: 2 (Fact) Predicate: (<= c\_y1 a\_a) Interpolant(old): false  
 Clause Id: 3 (Fact) Predicate: (<= c\_x1 b\_b) Interpolant(old): true  
 Clause Id: 5 (Fact) Predicate: (= c\_x1 a\_a) Interpolant(old): false  
 Clause Id: 6 (Fact) Predicate: (= c\_y1 b\_b) Interpolant(old): true  
 Clause Id: 7 (Fact) Predicate: (not (= c\_x1 c\_y1)) Interpolant(new): false  
 Clause Id: 8 (Conflict Clause) Predicate: (or (not (<= c\_y1 a\_a))  
     (not (<= c\_x1 b\_b))  
     (not (= c\_x1 a\_a))  
     (not (= c\_y1 b\_b))  
     (= c\_x1 c\_y1))

Inside partialInterpolantConflict

Case OCT

(ast-vector

    (<= c\_y1 a\_a)  
     (= c\_x1 a\_a)  
     (not (= c\_x1 c\_y1)))

(ast-vector

    (<= c\_x1 b\_b)  
     (= c\_y1 b\_b))

DNF: (let ((a!1 (and (<= (- c\_y1 a\_a) 0)  
     (<= (- c\_x1 a\_a) 0)  
     (<= (+ (- c\_x1) a\_a) 0)  
     (<= (- c\_x1 c\_y1) (- 1)))))  
   (a!2 (and (<= (- c\_y1 a\_a) 0)  
     (<= (- c\_x1 a\_a) 0)  
     (<= (+ (- c\_x1) a\_a) 0)  
     (<= (+ (- c\_x1) c\_y1) (- 1)))))  
   (or a!1 a!2))

Input for OctagonInterpolant (ast-vector

    (<= (- c\_y1 a\_a) 0)  
     (<= (- c\_x1 a\_a) 0)  
     (<= (+ (- c\_x1) a\_a) 0)  
     (<= (- c\_x1 c\_y1) (- 1)))

Input for OctagonInterpolant (ast-vector

    (<= (- c\_y1 a\_a) 0)  
     (<= (- c\_x1 a\_a) 0)  
     (<= (+ (- c\_x1) a\_a) 0)  
     (<= (+ (- c\_x1) c\_y1) (- 1)))

Theory-specific interpolant: (let ((a!1 (and (<= (+ (- c\_x1) c\_y1) 0) (<= (- c\_x1 c\_y1) (- 1)))))  
   (a!2 (and (<= (+ (- c\_x1) c\_y1) (- 1)))))  
   (or a!1 a!2))

Interpolant for OCT: (let ((a!1 (and (<= (+ (- c\_x1) c\_y1) 0) (<= (- c\_x1 c\_y1) (- 1)))))  
   (a!2 (and (<= (+ (- c\_x1) c\_y1) (- 1)))))  
   (and (or a!1 a!2 false false false) true true))

Interpolant((from conflict)new): (let ((a!1 (and (<= (+ (\* (- 1) c\_x1) c\_y1) 0)  
     (<= (+ c\_x1 (\* (- 1) c\_y1)) (- 1)))))  
   (or a!1 (<= (+ (\* (- 1) c\_x1) c\_y1) (- 1)))))

Clause Id: 9 (Derived(2,8)) Predicate: (or (not (<= c\_x1 b\_b)) (not (= c\_x1 a\_a)) (not (= c\_y1 b\_b)) (= c\_x1 c\_y1)) Pivot: (<= c\_y1 a\_a) ■  
 Pivot is A-local

Partial interpolant (let ((a!1 (and (<= (+ (\* (- 1) c\_x1) c\_y1) 0)

    (<= (+ c\_x1 (\* (- 1) c\_y1)) (- 1)))))

    (or false a!1 (<= (+ (\* (- 1) c\_x1) c\_y1) (- 1)))))

Interpolant((from derived)new): (let ((a!1 (and (<= (+ (\* (- 1) c\_x1) c\_y1) 0)

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```

      (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))
    (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
Clause Id: 10 (Derived(9,3)) Predicate: (or (= c_x1 c_y1) (not (= c_x1 a_a)) (not (= c_y1 b_b))) Pivot: (<= c_x1 b_b)
Pivot is B-local
Partial interpolant (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 true)))
Interpolant((from derived)new): (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
Clause Id: 11 (Derived(10,5)) Predicate: (or (= c_x1 c_y1) (not (= c_y1 b_b))) Pivot: (= c_x1 a_a)
Pivot is A-local
Partial interpolant (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)) false))
Interpolant((from derived)new): (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
Clause Id: 12 (Derived(11,6)) Predicate: (= c_x1 c_y1) Pivot: (= c_y1 b_b)
Pivot is B-local
Partial interpolant (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 true)))
Interpolant((from derived)new): (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (or a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
Clause Id: 13 (Derived(12,7)) Predicate: false Pivot: (= c_x1 c_y1)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 (or (not (= c_x1 c_y1)) false))))
Interpolant((from derived)new): (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 (not (= c_x1 c_y1)))))
Final interpolant for conflict clause: (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 (not (= c_x1 c_y1)))))

Disjunction implied in OCT: (= c_oct_1 c_oct_2)
Clause Id: 1 (Fact) Predicate: (= 0 c_oct_1) Interpolant(old): false
Clause Id: 4 (Fact) Predicate: (= 0 c_oct_2) Interpolant(old): true
Clause Id: 7 (Fact) Predicate: (not (= c_oct_1 c_oct_2)) Interpolant(new): false
Clause Id: 8 (Conflict Clause) Predicate: (or (not (= 0 c_oct_1)) (not (= 0 c_oct_2)) (= c_oct_1 c_oct_2))
Inside partialInterpolantConflict
Case OCT
-----It was sat!
(ast-vector
  (= 0 c_oct_1)
  (not (= c_oct_1 c_oct_2)))
(ast-vector
  (= 0 c_oct_2))
DNF: (let ((a!1 (and (<= c_oct_1 0)
      (<= (- c_oct_1) 0)

```

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```

      (<= (+ (- c_oct_1) c_oct_2) (- 1))))
(or (and (<= c_oct_1 0) (<= (- c_oct_1) 0) (<= (- c_oct_1 c_oct_2) (- 1)))
    a!1))
Input for OctagonInterpolant (ast-vector
  (<= c_oct_1 0)
  (<= (- c_oct_1) 0)
  (<= (- c_oct_1 c_oct_2) (- 1)))
Input for OctagonInterpolant (ast-vector
  (<= c_oct_1 0)
  (<= (- c_oct_1) 0)
  (<= (+ (- c_oct_1) c_oct_2) (- 1)))
Theory-specific interpolant: (let ((a!1 (and (<= (- c_oct_1) 0)
      (<= c_oct_1 0)
      (<= (+ (- c_oct_2) c_oct_1) (- 1)))))
  (or a!1
    (and (<= (- c_oct_1) 0) (<= c_oct_1 0) (<= (- c_oct_2 c_oct_1) (- 1)))))
Interpolant for OCT: (let ((a!1 (and (<= (- c_oct_1) 0)
  (<= c_oct_1 0)
  (<= (+ (- c_oct_2) c_oct_1) (- 1)))))
  (let ((a!2 (or a!1
    (and (<= (- c_oct_1) 0)
      (<= c_oct_1 0)
      (<= (- c_oct_2 c_oct_1) (- 1)))
    false
    false))))
    (and a!2 true)))
Interpolant((from conflict)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
  (or a!1 a!2))
Clause Id: 9 (Derived(1,8)) Predicate: (or (not (= 0 c_oct_2)) (= c_oct_1 c_oct_2)) Pivot: (= 0 c_oct_1)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
  (and (or (= 0 c_oct_1) false) (or (not (= 0 c_oct_1)) a!1 a!2)))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
  (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
Clause Id: 10 (Derived(9,4)) Predicate: (= c_oct_1 c_oct_2) Pivot: (= 0 c_oct_2)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
  (a!4 (or (not (not (= 0 c_oct_2))) true)))

```

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```

(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (and (or (not (= 0 c_oct_2)) a!3) a!4)))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (or (not (= 0 c_oct_2)) a!3)))
Clause Id: 11 (Derived(10,7)) Predicate: false Pivot: (= c_oct_1 c_oct_2)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
  (or (not (= c_oct_1 c_oct_2)) false))))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
  (not (= c_oct_1 c_oct_2)))))
Final interpolant for conflict clause: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
  (not (= c_oct_1 c_oct_2)))))
EUF solver found a contradiction
(ast-vector
  (= (c_f c_x1) c_oct_1)
  (= c_x1 a_a)
  (= c_y1 b_b)
  (not (= (c_f c_y1) c_oct_2))
  (= c_x1 c_y1)
  (= c_oct_1 c_oct_2))
Clause Id: 1 (Fact) Predicate: (= (c_f c_x1) c_oct_1) Interpolant(old): false
Clause Id: 4 (Fact) Predicate: (not (= (c_f c_y1) c_oct_2)) Interpolant(old): true
Clause Id: 5 (Fact) Predicate: (= c_x1 c_y1) Interpolant(old): (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
  (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (and a!2 (not (= c_x1 c_y1)))))
Clause Id: 6 (Fact) Predicate: (= c_oct_1 c_oct_2) Interpolant(old): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))

```



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```

      (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))
  (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
    (not (= c_oct_1 c_oct_2)))))
Clause Id: 7 (Conflict Clause) Predicate: (or (not (= (c_f c_x1) c_oct_1))
  (= (c_f c_y1) c_oct_2)
  (not (= c_x1 c_y1))
  (not (= c_oct_1 c_oct_2)))
Inside partialInterpolantConflict
Case EUF
Part a: (ast-vector
  (= (c_f c_x1) c_oct_1)
  (= c_x1 c_y1)
  (= c_oct_1 c_oct_2))
Part b: (ast-vector
  (not (= (c_f c_y1) c_oct_2)))
Theory-specific Interpolant for EUF: (and (= c_oct_1 (c_f c_x1)))
Interpolant for EUF: (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
  (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))
  (a!3 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!4 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
    (a!5 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!3 a!4))))
  (let ((a!6 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!5)
    (not (= c_oct_1 c_oct_2)))))
  (let ((a!7 (or (and (= c_oct_1 (c_f c_x1))
    false
    (and a!2 (not (= c_x1 c_y1)))
    a!6)))
    (and a!7 true))))))
Interpolant((from conflict)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
    (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
    (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))
  (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
    (not (= c_oct_1 c_oct_2)))))
    (or a!4 (= c_oct_1 (c_f c_x1)) (and a!6 (not (= c_x1 c_y1))))))
Clause Id: 8 (Derived(1,7)) Predicate: (or (= (c_f c_y1) c_oct_2) (not (= c_x1 c_y1)) (not (= c_oct_1 c_oct_2))) Pivot: (= (c_f c_x1) c_oct_1)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!2 (and (>= c_oct_1 0)
    (<= c_oct_1 0)
    (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
    (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2))))

```

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```

(a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
(let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2))) a!3)
      (not (= c_oct_1 c_oct_2)))))
(let ((a!7 (or (not (= (c_f c_x1) c_oct_1))
              a!4
              (= c_oct_1 (c_f c_x1))
              (and a!6 (not (= c_x1 c_y1))))))
      (and (or (= (c_f c_x1) c_oct_1) false) a!7))))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
      (a!2 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
      (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
      (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
            (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
            (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2))) a!3)
                  (not (= c_oct_1 c_oct_2)))))
              (let ((a!7 (or a!4
                (not (= (c_f c_x1) c_oct_1))
                (= c_oct_1 (c_f c_x1))
                (and a!6 (not (= c_x1 c_y1))))))
                (and (= (c_f c_x1) c_oct_1) a!7))))))
      Clause Id: 9 (Derived(8,4)) Predicate: (or (not (= c_oct_1 c_oct_2)) (not (= c_x1 c_y1))) Pivot: (= (c_f c_y1) c_oct_2)
      Pivot is AB-common
      Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
      (a!2 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
      (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
      (a!9 (or (not (= (c_f c_y1) c_oct_2)) true)))
      (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
            (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
            (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2))) a!3)
                  (not (= c_oct_1 c_oct_2)))))
              (let ((a!7 (or a!4
                (not (= (c_f c_x1) c_oct_1))
                (= c_oct_1 (c_f c_x1))
                (and a!6 (not (= c_x1 c_y1))))))
                (let ((a!8 (or (= (c_f c_y1) c_oct_2) (and (= (c_f c_x1) c_oct_1) a!7))))
                  (and a!8 a!9))))))
              Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1)))))
      (a!2 (and (>= c_oct_1 0)
      (<= c_oct_1 0)
      (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1)))))
      (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
      (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
      (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
            (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
            (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2))) a!3)
                  (not (= c_oct_1 c_oct_2)))))
              (not (= c_oct_1 c_oct_2)))))

```

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```

(let ((a!7 (or a!4
  (not (= (c_f c_x1) c_oct_1))
  (= c_oct_1 (c_f c_x1))
  (and a!6 (not (= c_x1 c_y1)))))))
  (or (= (c_f c_y1) c_oct_2) (and (= (c_f c_x1) c_oct_1) a!7))))))
Clause Id: 10 (Derived(9,5)) Predicate: (not (= c_oct_1 c_oct_2)) Pivot: (= c_x1 c_y1)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
  (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
    (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))))
  (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
    (not (= c_oct_1 c_oct_2))))
    (a!9 (or (not (not (= c_x1 c_y1))) (and a!6 (not (= c_x1 c_y1)))))))
  (let ((a!7 (or a!4
    (not (= (c_f c_x1) c_oct_1))
    (= c_oct_1 (c_f c_x1))
    (and a!6 (not (= c_x1 c_y1)))))))
    (let ((a!8 (or (not (= c_x1 c_y1))
      (= (c_f c_y1) c_oct_2)
      (and (= (c_f c_x1) c_oct_1) a!7))))
      (and a!8 a!9))))))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
  (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
    (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1))))))
  (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
    (not (= c_oct_1 c_oct_2))))
    (a!9 (or (= c_x1 c_y1) (and a!6 (not (= c_x1 c_y1)))))))
  (let ((a!7 (or a!4
    (not (= (c_f c_x1) c_oct_1))
    (= c_oct_1 (c_f c_x1))
    (and a!6 (not (= c_x1 c_y1)))))))
    (let ((a!8 (or (= (c_f c_y1) c_oct_2)
      (not (= c_x1 c_y1))
      (and (= (c_f c_x1) c_oct_1) a!7))))
      (and a!8 a!9))))))
Clause Id: 11 (Derived(10,6)) Predicate: false Pivot: (= c_oct_1 c_oct_2)
Pivot is AB-common
Partial interpolant: (let ((a!1 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
  (a!2 (and (>= c_oct_1 0)
  (<= c_oct_1 0)
  (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
  (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
  (<= (+ c_x1 (* (- 1) c_y1)) (- 1))))))

```

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```

(= (+ c_x1 (* (- 1) c_y1) (- 1))))))
(let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
      (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
(let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
                (not (= c_oct_1 c_oct_2)))))
      (a!9 (or (= c_x1 c_y1) (and a!6 (not (= c_x1 c_y1)))))
(let ((a!7 (or a!4
                (not (= (c_f c_x1) c_oct_1))
                (= c_oct_1 (c_f c_x1))
                (and a!6 (not (= c_x1 c_y1)))))
      (a!10 (or (not (not (= c_oct_1 c_oct_2))) a!4)))
(let ((a!8 (= (c_f c_y1) c_oct_2)
      (not (= c_x1 c_y1))
      (and (= (c_f c_x1) c_oct_1) a!7))))
  (and (or (not (= c_oct_1 c_oct_2)) (and a!8 a!9)) a!10))))))
Interpolant((from derived)new): (let ((a!1 (and (>= c_oct_1 0)
          (<= c_oct_1 0)
          (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
      (a!2 (and (>= c_oct_1 0)
          (<= c_oct_1 0)
          (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
      (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
          (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
        (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
  (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
                (not (= c_oct_1 c_oct_2)))))
        (a!9 (or (= c_x1 c_y1) (and a!6 (not (= c_x1 c_y1)))))
  (let ((a!7 (or a!4
                (not (= (c_f c_x1) c_oct_1))
                (= c_oct_1 (c_f c_x1))
                (and a!6 (not (= c_x1 c_y1)))))
        (let ((a!8 (or (= (c_f c_y1) c_oct_2)
                (not (= c_x1 c_y1))
                (and (= (c_f c_x1) c_oct_1) a!7))))
          (and (or (not (= c_oct_1 c_oct_2)) (and a!8 a!9))
                (or (= c_oct_1 c_oct_2) a!4))))))
Final interpolant for conflict clause: (let ((a!1 (and (>= c_oct_1 0)
          (<= c_oct_1 0)
          (<= (+ (* (- 1) c_oct_2) c_oct_1) (- 1))))
      (a!2 (and (>= c_oct_1 0)
          (<= c_oct_1 0)
          (<= (+ c_oct_2 (* (- 1) c_oct_1)) (- 1))))
      (a!5 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
          (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!3 (and (= 0 c_oct_1) (or (not (= 0 c_oct_1)) a!1 a!2)))
        (a!6 (or (= c_x1 c_y1) a!5 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
  (let ((a!4 (and (or (= c_oct_1 c_oct_2) (not (= 0 c_oct_2)) a!3)
                (not (= c_oct_1 c_oct_2)))))
        (a!9 (or (= c_x1 c_y1) (and a!6 (not (= c_x1 c_y1)))))
  (let ((a!7 (or a!4
                (not (= (c_f c_x1) c_oct_1))
                (= c_oct_1 (c_f c_x1))
                (and a!6 (not (= c_x1 c_y1)))))
        (let ((a!8 (or (= (c_f c_y1) c_oct_2)
                (not (= c_x1 c_y1))
                (and (= (c_f c_x1) c_oct_1) a!7))))
          (and (or (not (= c_oct_1 c_oct_2)) (and a!8 a!9))
                (or (= c_oct_1 c_oct_2) a!4))))))

```

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```
(or (= c_oct_1 c_oct_2) a!4))))))
-> Final Interpolant: (let ((a!1 (and (<= (+ (* (- 1) c_x1) c_y1) 0)
                                     (<= (+ c_x1 (* (- 1) c_y1)) (- 1)))))
  (let ((a!2 (or (= c_x1 c_y1) a!1 (<= (+ (* (- 1) c_x1) c_y1) (- 1)))))
    (let ((a!3 (or (not (= (c_f c_x1) 0))
                  (= 0 (c_f c_x1))
                  (and a!2 (not (= c_x1 c_y1)))))
      (a!5 (or (= c_x1 c_y1) (and a!2 (not (= c_x1 c_y1)))))
    (let ((a!4 (or (not (= c_x1 c_y1)) (= (c_f c_y1) 0) (and (= (c_f c_x1) 0) a!3)))
      (and a!4 a!5))))
Interpolant:
(let ((a!1 (and (<= (+ (* (- 1) x1) y1) 0) (<= (+ x1 (* (- 1) y1)) (- 1)))))
  (let ((a!2 (or (= x1 y1) a!1 (<= (+ (* (- 1) x1) y1) (- 1)))))
    (let ((a!3 (or (not (= (f x1) 0)) (= 0 (f x1)) (and a!2 (not (= x1 y1)))))
      (a!5 (or (= x1 y1) (and a!2 (not (= x1 y1)))))
    (let ((a!4 (or (not (= x1 y1)) (= (f y1) 0) (and (= (f x1) 0) a!3)))
      (and a!4 a!5))))
rm -rf tests/*.o tests/basic_test
```

# Appendix C: additional implementation code of relevant data structures

## 6.3 Congruence Closure with Explanation operation header

```
1  class Hornsat;  
2  
3  class CongruenceClosureExplain : public CongruenceClosure {  
4  
5      Hornsat * hsat;  
6  
7      PendingElements pending_elements;  
8      PendingPointers equations_to_merge;  
9      PendingPointers pending_to_propagate;  
10  
11      FactoryCurryNodes const & factory_curry_nodes;  
12  
13      LookupTable lookup_table;
```

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```
14  UseList      use_list;
17
18  void pushPending(PendingPointers &, const PendingElement &);
19  void merge();
20  void merge(EquationCurryNodes const &);
21  void propagate();
22  void propagateAux(CurryNode const &, CurryNode const &,
23                  EqClass, EqClass, PendingElement const &);
24
24  EqClass      highestNode(EqClass, UnionFind &);
25  EqClass      nearestCommonAncestor(EqClass, EqClass,
26                                     UnionFind &);
26  PendingPointers explain(EqClass, EqClass);
27  void          explainAlongPath(EqClass, EqClass, UnionFind
28                                &, ExplainEquations &, PendingPointers &);
28  std::ostream & giveExplanation(std::ostream &, EqClass,
29                                EqClass);
29
30  public:
31  CongruenceClosureExplain(Hornsat *, CongruenceClosureExplain
32                            const &, UnionFindExplain &);
32  CongruenceClosureExplain(Z3Subterms const &,
33                            UnionFindExplain &, FactoryCurryNodes &, IdsToMerge const
34                            &);
33  ~CongruenceClosureExplain();
34
35  bool areSameClass(EqClass, EqClass);
36  bool areSameClass(z3::expr const &, z3::expr const &);
37
38  EqClass constantId(EqClass);
```

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```
39 EqClass find(EqClass);
42 z3::expr z3Repr(z3::expr const &);
43
44 void merge(EqClass, EqClass);
45 void merge(z3::expr const &, z3::expr const &);
46
47 PendingPointers explain(z3::expr const &, z3::expr const &);
48 std::ostream & giveExplanation(std::ostream &, z3::expr
    const &, z3::expr const &);
49
50 z3::expr_vector z3Explain(z3::expr const &, z3::expr const
    &);
51 std::ostream & z3Explanation(std::ostream &, const z3::expr
    &, const z3::expr &);
52
53 friend std::ostream & operator << (std::ostream &, const
    CongruenceClosureExplain &);
54 };
```

## 6.4 Hornsat (Gallier's data structure) header

```
1 class Hornsat {
2
3     friend class EUFInterpolant;
4
5     unsigned num_hcs, num_literals;
6     // This structure is only used in our approach
7     // for conditional-elimination
8     std::unordered_map<unsigned, HornClause *> head_term_indexer
    ;
```



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```
9
12 UnionFindExplain      ufe;
13 CongruenceClosureExplain equiv_classes;
14
15 std::vector<Literal>    list_of_literals;
16 ClassList               class_list;
17 std::vector<unsigned>   num_args;
18 std::vector<LiteralId> pos_lit_list;
19 // 'facts' is a queue of all the (temporary)
20 // literals that have value true
21 std::queue<LiteralId>   facts;
22 std::queue<TermIdPair> to_combine;
23
24 bool consistent;
25
26 void satisfiable();
27 void closure();
28
29 public:
30 Hornsat(CongruenceClosureExplain &, HornClauses const &);
31 ~Hornsat();
32
33 void build(CongruenceClosureExplain &, HornClauses const &);
34 bool isConsistent() const ;
35 void unionupdate(LiteralId, LiteralId);
36 friend std::ostream & operator << (std::ostream &, Hornsat
    const &);
37 };
```

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