

# The Tyr Programming Language

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## Abstract

This document defines Tyr, a research language for type-oriented programming. Type-oriented programming is a paradigm that extends on object-oriented programming. In type-oriented languages, types are first order values like integers and objects. An existing but primitive form of type orientation is the Java reflection API.

## Acknowledgements

For Pony!

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## Part I

# Core Language

## 1 Introduction

Type-oriented programming (TOP) is a paradigm that states that types are objects. In consequence, it is possible to perform calculations on types like any other calculation. As it is true for objects in object-oriented programming (OOP), types can be copied and may have mutable state. The mutable state of a type can be bounded by static knowledge in the same way as pointer can be restricted to point to objects of a certain type. As such, TOP implies OOP.

Tyr is a programming language created to explore this idea in practice. Tyr as a language is a descendant of C++ and Scala. In order to examine consequences of TOP for resource management, Tyr features manual memory management.

Type-level functions are descendants of C++ templates and Ada generics.

### 1.1 On OOP and TOP

TOP has no point without OOP for the following reason:

```
type A;  
type B;
```

```
val x = new (if(phi) A else B)();
```

What could be the type of  $x$  if not  $A \sqcap B$ ?

## 2 Syntax

The syntax of Tyr is inspired by Scala and SKILL.

### 2.1 Literals

### 2.2 Grammar

#### 2.2.1 Top Level Structure

#### 2.2.2 Members

#### 2.2.3 Expressions

## 3 Semantics

The semantics of Tyr is loosely based on C++ and Scala.

### 3.1 Types of literal Values

The type of `<Integer>` is `int`, if no `'i'` or no number behind the `'i'` is supplied. If a number is supplied, the number will be used as argument for `Integer(n)`. If the

To do (1)

$\langle int \rangle$	::= ('0'-'9') <sup>+</sup>
$\langle hex \rangle$	::= ('0'-'0' 'A'-'F' 'a'-'f') <sup>+</sup>
$\langle Integer \rangle$	::= '-'? $\langle int \rangle$ ('i' $\langle int \rangle$ )?
$\langle HexInteger \rangle$	::= 0x $\langle hex \rangle$ ('i' $\langle int \rangle$ )?
$\langle long \rangle$	::= '-'? $\langle int \rangle$ 'L'
$\langle Float \rangle$	::= '-'? $\langle int \rangle$ ? '.' $\langle int \rangle$ (('e' 'E') '-'? $\langle int \rangle$ )? ('f' $\langle int \rangle$ )?
$\langle string \rangle$	::= '"' ~['\']* '"'
$\langle Identifier \rangle$	::= '~' ~['~'] <sup>+</sup> '~'

Figure 1: Literals

$\langle file \rangle$	::= ???
------------------------	---------

Figure 2: Literals

resulting Type has a named subtype in `tyr.lang`, the subtype will be chosen. Hence, `0i8` is a byte of value 0. Also, `0`, `0i` and `0i32` are indistinguishable.

The type of `<HexInteger>` is `UnsignedInteger(n)`, where `n` is the number supplied via 'i' defaulting to 32.

The type of `<long>` is `long`.

The type of `<Float>` is analogous to `<Integer>` except that it is based on `FloatingPoint(n)` and defaults to `double`. The type is `float` if a single `f` is supplied. This rule is designed to be compatible with common programming languages.

The type of `<string>` is `LiteralString`.

`<Identifier>` literals yield identifiers. An identifier is neither a type nor a value.

### 3.2 Unescaping of String and Identifier Literals

Tyr uses the same unescaping mechanism as Java. Unescaping happens for `<string>` and `<identifier>` before further processing.

### 3.3 Access Paths introduced by local with

Any expression that yields a scope can be imported with a `with`. This import results in an access path to the names imported into the local scope. If that access has a side effect, the effect may be executed at every access to the scope. Function with an implicit `this` parameter start with an implicit `with this;` expression.

`var/val`: fields type `var` -> type field (in vtable)

`defs`: `def` -> virtual static `def` -> static type (ada non-poly pointer) type `def` -> type method

Typen: Any (top) void (<: Any) bool Integer int byte long UnsignedInteger FloatingPoint float double pointer

class Object <: pointer String <: Object IterableOnce <: String Iterable <: IterableOnce Option <: Iterable Seq <: Iterable Array <: Iterable

## Part II

# Compilation

modules, source paths, modules scopes, default scopes,

module naming convention: <organization>.<project> tyr.lang tyr.system tyr.collection skill.common

## Part III

# Libraries

## 4 IO

- Path (VFS) - File (cfile) - MappedFile (mmap) - Console

## 5 Collections

IterableOnce(T : Type) - static def for (p, b) - def foreach (f : LocalLambda[-> T])  
Iterator <: IterableOnce - empty() - move() : bool - get() - for (p, b) = if(!empty) do  
EquivalenceRelation(T : Type) - equals(T, T) : bool - hash(T) : int  
MinimalEQR <: EquivalenceRelation - equals := == - hash := .to(Int)  
Iterable <: IterableOnce  
Seq <: Iterable  
Array <: Seq  
ArrayBuffer <: Seq  
StringBuffer <: Seq(String)  
List <: Seq  
LinkedList <: List  
Set <: Seq  
HashSet(T : Type <: CT, Eq : Type(EquivalenceRelation) <: CT = MinimalEQR) <:  
Set  
Map <: Iterable  
HashMap(K : Type <: CT, V : Type <: CT, Eq : Type(EquivalenceRelation(K)) <:  
CT) <: Set

## 6 Threads

- Thread - ThreadPool - Semaphore - Mutex - Barrier

## 7 Native

-C method placement -C++ method placement?

Part IV

# Appendix





### **To do...**

- ☐ 1 (p. 3): sind die ints nicht in wahrheit LiteralInt?