CSC343 - Assignment 3 Dependcies

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Question 1

(a)

Question 2

Relation T contains attributes CDEFGHIJ and FDs $S_T = \{C \to EH, DEI \to F, F \to D, EH \to CJ, J \to FGI\}$

(a)

- $C^+ = CDEFGHIJ$ so $C \to EH$ does not violate BCNF
- $DEI^+ = DEIF$ so $DEI \rightarrow F$ does violate BCNF
- $F^+ = FD$ so $F \to D$ does violate BCNF
- $EH^+ = CDEFGHIJ$ so $EH \to CJ$ does not violate BCNF
- $J^+ = JFGID$ so $J \to FGI$ does violate BCNF

(b)

- Decompose T using FD $DEI \rightarrow F$. $DEI^+ = DEIF$, so this yields two relations $R_1 = DEFI$ and $R_2 = CDEGHIJ$
- Project the DFs onto $R_1 = DEFI$.

D	Е	F	I	closure	FDs
\checkmark				$D^+ = D$	nothing
	√			$E^+ = E$	nothing
		√		$F^+ = FD$	$F \to D$; violates BCNF; abort

We must decompose R_1 further.

- Decompose R_1 using FD $F \to D$. so this yields two relations $R_3 = FD$ and $R_4 = FEI$
- Project the FDs onto $R_3 = FD$.

F	D	closure	FDs
√		$F^+ = FD$	$F \to D$
	√	$D^+ = D$	nothing

• Project the FDs onto $R_4 = FEI$.

F	Е	I	closure	FDs
\checkmark			$F^+ = FD$	nothing
	✓		$E^+ = E$	nothing
		√	$I^+ = I$	nothing
√	✓		$FE^+ = FED$	nothing
\checkmark		√	$FI^+ = FI$	nothing
	✓	✓	$EI^+ = EI$	nothing

• Return to $R_2 = CDEGHIJ$ and prodject the FDs onto it.

$\begin{array}{c ccccccccccccccccccccccccccccccccccc$							J .		
$\begin{array}{c ccccc} \checkmark & & & & & & & & & & & \\ \hline \checkmark & & & & & & & & & \\ \hline & \checkmark & & & & & & & \\ \hline & & \checkmark & & & & & & \\ \hline & & & \checkmark & & & & \\ \hline & & & & & & & \\ \hline & & & & \checkmark & & & \\ \hline & & & & & & \\ \hline & & & & & & \\ \hline & & & &$	С	D	E	G	Н	I	J	closure	FDs
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	✓							CDEGHIJ	nothing
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		√						$D^+ = D$	nothing
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			√						nothing
\checkmark $I^+=I$ nothing				√					nothing
					√			$H^+ = H$	nothing
$\checkmark \qquad J^{+} = JFGID \qquad J \rightarrow IDG$						√			nothing
							√	$J^+ = JFGID$	$J \rightarrow IDG$

- Decompose R_2 using FD $J \to IDG$. so this yields two relations $R_5 = JIDG$ and $R_6 = CEHJ$
- Project the FDs onto $R_5 = JIDG$.

J	Ι	D	G	closure	FDs
<u> </u>				$J^+ = JFGID$	$J \rightarrow IDG$; dont need to consider
,				0 01 012	super-sets of J
	✓			$I^+ = I$	nothing
		✓		$D^+ = D$	nothing
			✓	$G^+ = G$	nothing
	√	√		$ID^+ = ID$	nothing
	√		✓	$IG^+ = IG$	nothing
		✓	✓	$DG^+ = DG$	nothing

• Project the FDs onto $R_6 = CEHJ$.

С	Е	Н	J	closure	FDs
				$C^+ =$	$C \rightarrow EHJ$; dont need to consider
V				CEHJFGI	super-sets of C
	✓			$E^+ = E$	nothing
		√		$H^+ = H$	nothing
			√	$J^+ = J$	nothing
	(($EH^+ =$	$EH \rightarrow CEHJFGID$
	V	V		CEHJFGID	$EII \rightarrow CEIIJI GID$
	./		./	$EJ^+ =$	nothing
	V		V	EJFGID	nothing
		./		$HJ^+ =$	nothing
		V	V	HJFGID	nouning

So relation T containing attributes CDEFGHIJ decomposes in to relations

- $R_3 = FD$ $FDs = \{F \to D\}$
- $R_4 = EFI$ FDs = {}
- $R_5 = DGIJ$ FDs = $\{J \to DGI\}$
- $R_6 = CEHJ$ FDs = $\{EH \to CJ, C \to EH\}$