# Elliptic Curve Cryptography

Fast implentation of the Diffie-Hellman key exchange

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# Elliptic Curve

Algorithm

▶ An elliptic curve  $E(\mathbb{F}_p)$  consists of the set of the points P(x,y),  $x, y \in \mathbb{F}_p$  satisfying

$$y^2 \equiv x^3 + ax + b \pmod{p}$$

Possible to define an addition rule to add points on E



# Diffie Hellman key exchange

## **Public parameters**

Algorithm

$$y^2 \equiv x^3 + ax + b \pmod{p}$$

 $a, b \in \mathbb{F}_p$ , a prime p and a base point G are known

### **Private computations**

Alice Bob

Compute P = uGCompute Q = vG

## Public exchange of values

Alice  $\xrightarrow{P}$  Bob Bob

## **Further private computations**

Alice Bob

Compute uQCompute *vP* 

The shared secret is uQ = u(vG) = v(uG) = vP

Adaption of Table 2.2 J. Hoffstein et al., An Introduction to Mathematical Cryptography

## Double-and-add-Method

- ▶ Input  $P \in E(F_p)$ ,  $d \in \mathbb{N}$
- ▶ Output:  $d \cdot P \in E(F_p)$

Algorithm

```
N < -P
Q < - \mathcal{O}
for i from 0 to m do
  if d_i = 1 then
     Q \leftarrow point_add(Q, N)
N \leftarrow point_double(N)
return Q
```

where 
$$d = d_0 + d_1 2 + ... + d_m 2^m$$
  $d_i \in \{0, 1\}$ 

https://en.wikipedia.org/wiki/Elliptic\_curve\_point\_multiplication#Double-and-add



# **Implementation**

Bigint

```
typedef uint64_t block;
2
  typedef struct
3
       uint64_t significant_blocks;
       block blocks[BIGINT_BLOCKS_COUNT];
6
7
     __BigInt;
```

Corresponding operations (Addition, Multiplication, Division, ...)

- Elliptic Curve and ECDH
  - ► Elliptic Curve definition and key exchange mechanism
  - 5 predefined curves: 192, 224, 256, 384, 521 bits



# Cost Anaylsis

- Index integer operation matters
- Cost measure
  - $ightharpoonup C = C_{add} + C_{mult} + C_{shift}$
  - Code generated operations counts

# **Optimizations**

#### Stages

Baseline - Implementation without memory optimization

Memory optimization - Implementation with memory optimization

Comparison with OpenSSL

Algorithm

Performance

Algorithm

Performance

Final - Final performance optimization



# **Optimizations**

#### Overview

#### Baseline

- Improved memory allocation
- ► Change the base from 8 bit to 64 bit for the big integers

## Memory optimization

- Change the base from 2 bit to 64 bit for Montogomery
- ▶ Precomputation of  $2^k \cdot G$  where  $k \in \{1, ..., \log_2(p)\}$  and  $G \in E$
- Introduction of Jacobian coordinates
- ▶ Vectorize shifting using AVX2

#### Jacobian coordinates

- function stiching
- Inlining functions
- Unrolling



# Intel ADX

#### C Code

```
low_m1 = _mulx_u64(a->blocks[i], b, &hi_m1);
add.carry_m1 = _addcarryx_u64(add_carry_m1, carry_m1, low_m1, &temp_m1);
add_carry_1 = _addcarryx_u64(add_carry_1, res->blocks[i], temp_m1, tmp->blocks[i]);
carry_m1 = hi_m1;
```

## Created Assembly code

x86 icc 13.0.1 -m64 -march=CORE-AVX2 -O3

http://gcc.godbolt.org/



# Intel ADX

#### C Code

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add_carry_1 = _addcarryx_u64(add_carry_1, res->blocks[i], temp_m1, tmp->blocks[i]);
carry_m1 = hi_m1;
```

## Created Assembly code

x86 gcc 5.3 -m64 -march=haswell -O3

```
mulx
            48(% rsi), %r9, %r10
            %r9. %r11
   adda
            %r10, %r9
   movq
            %bpl
   setc
          $-1. %bl
   addb
6
           48(% rdi), %r11
   adca
7
            %r11, 2288(%rax)
   movq
    setc
            %b1
            -1. \%bpl
   addh
```

http://gcc.godbolt.org/



## ADX vs AVX2

- Bottleneck operation: BigInt block multiplication
- Unavoidable dependecies in carry chain -> vectorization by processing 4 multiplications in parallel

| Approach Lower bound |                     | Bottleneck         |  |
|----------------------|---------------------|--------------------|--|
| ADX                  | 8 cycles/iteration  | ADX throughput     |  |
| AVX2 (base 32)       | 10 cycles/iteration | Emulation of carry |  |
| AVX2 (base 64)       | 24 cycles/iteration | flag               |  |

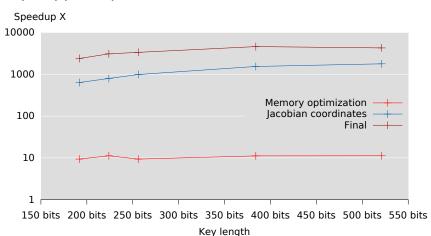
- Further AVX2 downsides
  - higher mul latencies
  - unfriendly data layout
  - multiplications not always independent



# Experiment result

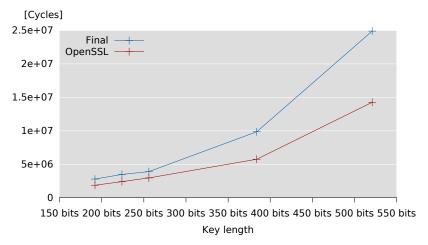
- Environment
  - Platform:Arch Linux 64 bit, GCC 6.1.1 compiler
  - Skylake i7-6600U CPU @ 3GHz
  - ▶ 64 bit multiplication (mul, mulx): 1 op/cycle
  - ▶ 64 bit addition/subtraction (add, sub): 4 op/cycle
  - ▶ 64 bit addition with carry (adc, adcx, adox): 1 op/cycle
  - Carry addition only: peak performance of 2 ops/cycle 6 Gops/s on 1 core
  - Compile flag: -O3 -mavx2 -mbmi2 -madx



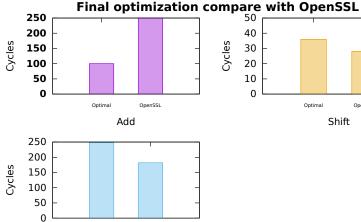


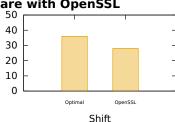


#### **ECDH** execution cycles comparison







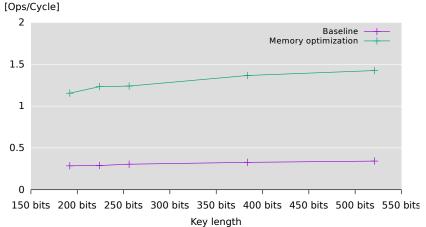


Optimal

Montgomery Mult

OpenSSL

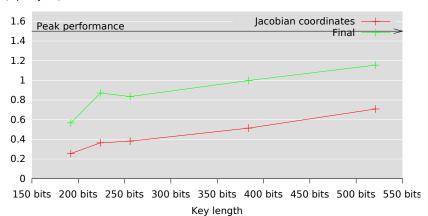






### Performance plot Part 2

[Ops/Cycle]



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## Performance plot operation counts

| Key size | Baseline    | Memory optimization | Precomputation | Jacobian<br>coordinates | Final    |
|----------|-------------|---------------------|----------------|-------------------------|----------|
| 192      | 2117237004  | 912515528           | 16277780       | 2977446                 | 1585686  |
| 224      | 3666675252  | 1384425854          | 25816754       | 5796852                 | 3035341  |
| 256      | 4901227919  | 2122930314          | 35137355       | 6206895                 | 3257705  |
| 384      | 18873323391 | 7047549105          | 109331889      | 19264827                | 9848295  |
| 521      | 48749705798 | 18063182851         | 282776551      | 56794800                | 28765815 |



#### Speedup plot with varied private key form, compared to OpenSSL

Speedup/OpenSSL X

