Elliptic Curve Cryptography

Fast implentation of the Diffie-Hellman key exchange

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Algorithm

▶ An elliptic curve $E(\mathbb{F}_p)$ consists of the set of the points P(x,y), $x, y \in \mathbb{F}_p$ satisfying

$$y^2 \equiv x^3 + ax + b \pmod{p}$$

Possible to define an addition rule to add points on E



Optimizations

Diffie Hellman key exchange

Public parameters

Algorithm

$$y^2 \equiv x^3 + ax + b \pmod{p}$$

 $a, b \in \mathbb{F}_p$, a prime p and a base point G are known

Private computations

Alice Bob

Compute P = uGCompute Q = vG

Public exchange of values

Alice \xrightarrow{P} Bob Bob

Further private computations

Alice Bob

Compute uQCompute *vP*

The shared secret is uQ = u(vG) = v(uG) = vP

Adaption of Table 2.2 J. Hoffstein et al., An Introduction to Mathematical Cryptography



Double-and-add-Method

▶ Input $P \in E(F_p)$, $d \in \mathbb{N}$

Algorithm

▶ Output: $d \cdot P \in E(F_p)$

```
 \begin{array}{|c|c|c|c|c|} \hline 1 & N & < & P \\ \hline 2 & Q & < & \mathcal{O} \\ \hline 3 & for & i & from & 0 & to & m & do \\ \hline 4 & & if & d_i & = & 1 & then \\ \hline 5 & & Q & < & point\_add(Q, N) \\ \hline 6 & N & < & point\_double(N) \\ \hline 7 & return & Q \\ \hline \end{array}
```

where
$$d = d_0 + d_1 2 + ... + d_m 2^m$$
 $d_i \in \{0, 1\}$

https://en.wikipedia.org/wiki/Elliptic_curve_point_multiplication#Double-and-add



Implementation

Bigint

```
typedef uint64_t block;

typedef struct
{
    uint64_t significant_blocks;
    block blocks[BIGINT_BLOCKS_COUNT];
} __BigInt;
```

Corresponding operations (Addition, Multiplication, Division, ...)

- Elliptic Curve and ECDH
 - ▶ Elliptic Curve definition and key exchange mechanism
 - ▶ 5 predefined curves : 192, 224, 256, 384, 521 bits



Cost Anaylsis

- Index integer operation matters
- Cost measure
 - $ightharpoonup C = C_{add} + C_{mult} + C_{shift}$
 - Code generated operations counts

Optimizations

Stages

Baseline - Implementation without memory optimization Performance Memory optimization - Implementation with memory optimization Comparison with OpenSSL Algorithm Jacobian coordinates - Algorithmic changes, jacobian coordinates Performance **Final** - Final performance optimization



Optimizations

Overview

Baseline

- Improved memory allocation
- ► Change the base from 8 bit to 64 bit for the big integers

Memory optimization

- Change the base from 2 bit to 64 bit for Montogomery
- ▶ Precomputation of $2^k \cdot G$ where $k \in \{1, ..., \log_2(p)\}$ and $G \in E$
- Introduction of Jacobian coordinates
- Vectorize shifting using AVX2

Jacobian coordinates

- function stiching
- Inlining functions
- Unrolling



Intel ADX

C. Code

```
low_m1 = _mulx_u64(a \rightarrow blocks[i], b, &hi_m1);
add_carry_m1 = _addcarryx_u64(add_carry_m1, carry_m1, low_m1, &temp_m1);
add_carry_1 = _addcarryx_u64(add_carry_1, res->blocks[i], temp_m1, tmp->blocks[i]);
carry_m1 = hi_m1;
```

Created Assembly code

x86 icc 13.0.1 -m64 -march=CORE-AVX2 -O3

```
mov
          rdx, QWORD PTR [48+rsi]
          rdx. rbx. rax
mulx
adox
          rbp, rdx
adcx
          rbp, QWORD PTR [48+rdi]
          OWORD PTR [2288+r9], rbp
mov
```

http://gcc.godbolt.org/



Intel ADX

C Code

```
low_m1 = _mulx_u64(a->blocks[i], b, &hi_m1);
add.carry_m1 = _addcarryx_u64(add_carry_m1, carry_m1, low_m1, &temp_m1);
add_carry_1 = _addcarryx_u64(add_carry_1, res->blocks[i], temp_m1, tmp->blocks[i]);
carry_m1 = hi_m1;
```

Created Assembly code

x86 gcc 5.3 -m64 -march=haswell -O3

```
mulx
            48(% rsi), %r9, %r10
            %r9. %r11
   adda
            %r10, %r9
   movq
            %bpl
   setc
          $-1. %bl
   addb
6
           48(% rdi), %r11
   adca
7
            %r11, 2288(%rax)
   movq
    setc
            %b1
            -1. \%bpl
   addh
```

http://gcc.godbolt.org/



ADX vs AVX2

- Bottleneck operation: BigInt block multiplication
- Unavoidable dependecies in carry chain -> vectorization by processing 4 multiplications in parallel

Approach	Lower bound	Bottleneck	
ADX	8 cycles/iteration	ADX throughput	
AVX2 (base 32)	10 cycles/iteration	Emulation of carry	
AVX2 (base 64)	24 cycles/iteration	flag	

- Further AVX2 downsides
 - higher mul latencies
 - unfriendly data layout
 - multiplications not always independent

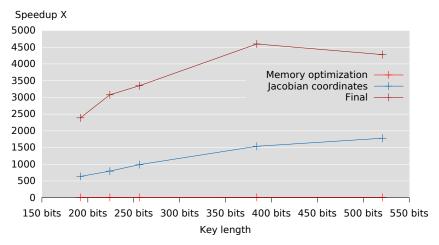


Experiment result

- Environment
 - ▶ Platform:Arch Linux 64 bit, GCC 6.1.1 compiler
 - Skylake i7-6600U CPU @ 3GHz
 - ▶ 64 bit multiplication (mul, mulx): 1 op/cycle
 - ▶ 64 bit addition/subtraction (add, sub): 4 op/cycle
 - ▶ 64 bit addition with carry (adc, adcx, adox): 1 op/cycle
 - Carry addition only: peak performance of 2 ops/cycle 6 Gops/s on 1 core
 - Compile flag: -O3 -mavx2 -mbmi2 -madx

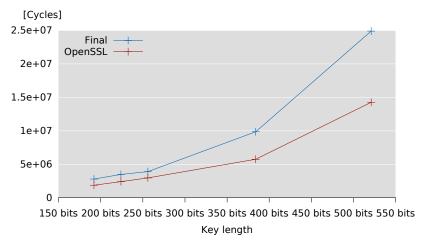


Speedup plot compared to Baseline

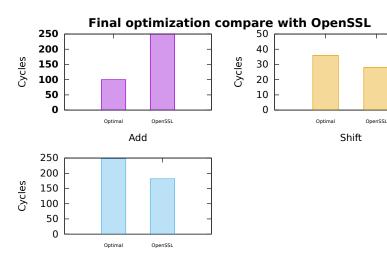




ECDH execution cycles comparison





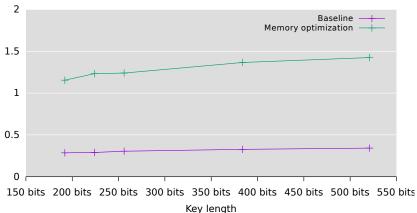




Montgomery Mult

Performance plot Part 1

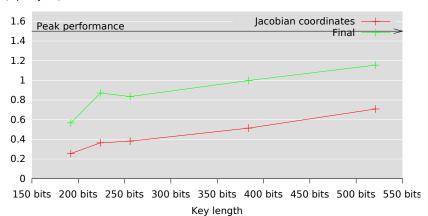






Performance plot Part 2

[Ops/Cycle]





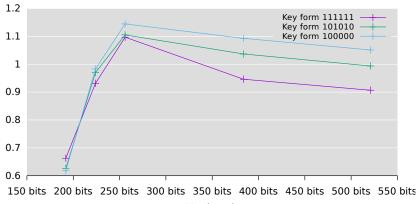
Performance plot operation counts

Key size	Baseline	Memory optimization	Precomputation	Jacobian coordinates	Final
192	2117237004	912515528	16277780	2977446	1585686
224	3666675252	1384425854	25816754	5796852	3035341
256	4901227919	2122930314	35137355	6206895	3257705
384	18873323391	7047549105	109331889	19264827	9848295
521	48749705798	18063182851	282776551	56794800	28765815



Speedup plot with varied private key form, compared to OpenSSL

Speedup/OpenSSL X



Key length

