

III Combinatorics

Ishan Nath, Michaelmas 2024

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0 Introduction

We have the following list of things.

- 1: Set systems.
- 2: Isoperimetric inequalities.
- 3: Intersection families.

Books include ‘Combinatorics’ by Bollobás, and ‘Combinatorics of Finite Sets’, by Anderson.

1 Set Systems

Let X be a set. A *set system* on X , also called a family of subsets of X , is a family $\mathcal{A} \subseteq \mathcal{P}(X)$. For example,

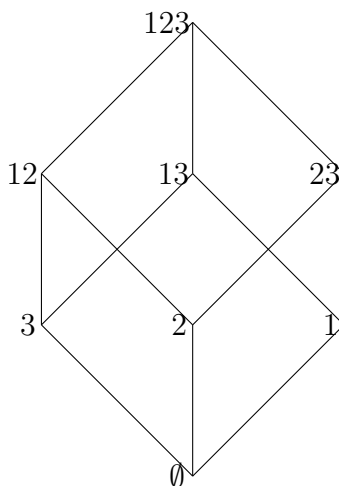
$$X^{(r)} = \{A \subseteq X \mid |A| = r\}.$$

Usually, $X = [n] = \{1, 2, \dots, n\}$, so $|X^{(r)}| = \binom{n}{r}$. Thus,

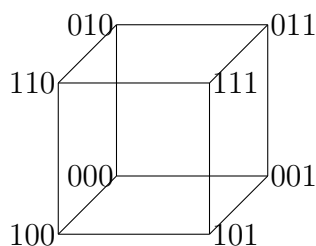
$$[4]^{(2)} = \{12, 13, 14, 23, 24, 34\}.$$

We make $\mathcal{P}(X)$ into a graph by joining A and B if $|A \triangle B| = 1$. This is the *discrete cube* Q_n .

Literally just a cube.



Alternatively, can view Q_n as an n -dimensional unit cube $\{0, 1\}^n$, by identifying e.g. $\{1, 3\}$ with the binary string 101000...



Say $\mathcal{A} \subseteq \mathcal{P}(X)$ is a *chain* if, for all $A, B \in \mathcal{A}$, $A \subseteq B$ or $B \subseteq A$. For example,

$$\mathcal{A} = \{23, 12357, 1235, 123567\}$$

is a chain.

Say \mathcal{A} is an *antichain* if, for all $A, B \in \mathcal{A}$ and $A \neq B$, we have $A \not\subseteq B$. For example, $\mathcal{A} = \{23, 137\}$ is an antichain.

How large can a chain be? We can achieve $|\mathcal{A}| = n + 1$ by taking

$$\mathcal{A} = \{\emptyset, 1, 12, 123, \dots, [n]\}$$

Cannot beat this as each $0 \leq r \leq n$, \mathcal{A} can contain at most one r -set (a member of $X^{(r)}$).

How large can an antichain be? We can achieve $|\mathcal{A}| = n$, e.g. $\mathcal{A} = \{1, 2, \dots, n\}$. More generally, we can take $\mathcal{A} = X^{(r)}$, and the best is when $r = \lfloor n/2 \rfloor$.

Theorem 1.1 (Sperner's Lemma). *Let $\mathcal{A} \subseteq \mathcal{P}(X)$ be an antichain. Then,*

$$|\mathcal{A}| \leq \binom{n}{\lfloor n/2 \rfloor}.$$

The idea is follows: we know that a chain meets a layer in at most one point, since a layer is an antichain. If we decompose the cube into chains, we have at most one element of an antichain in each chain.

Proof: We will decompose $\mathcal{P}(X)$ into $\binom{n}{\lfloor n/2 \rfloor}$ chains, then we are done. To achieve this, it is sufficient to find:

- (i) For each $r < n/2$, a matching from $X^{(r)}$ to $X^{(r+1)}$.
- (ii) For each $r \geq n/2$, a matching from $X^{(r)}$ to $X^{(r-1)}$.

Then we put these together to form our chains; each passing through $X^{(\lfloor n/2 \rfloor)}$.

By taking complements, it is enough to prove (i).

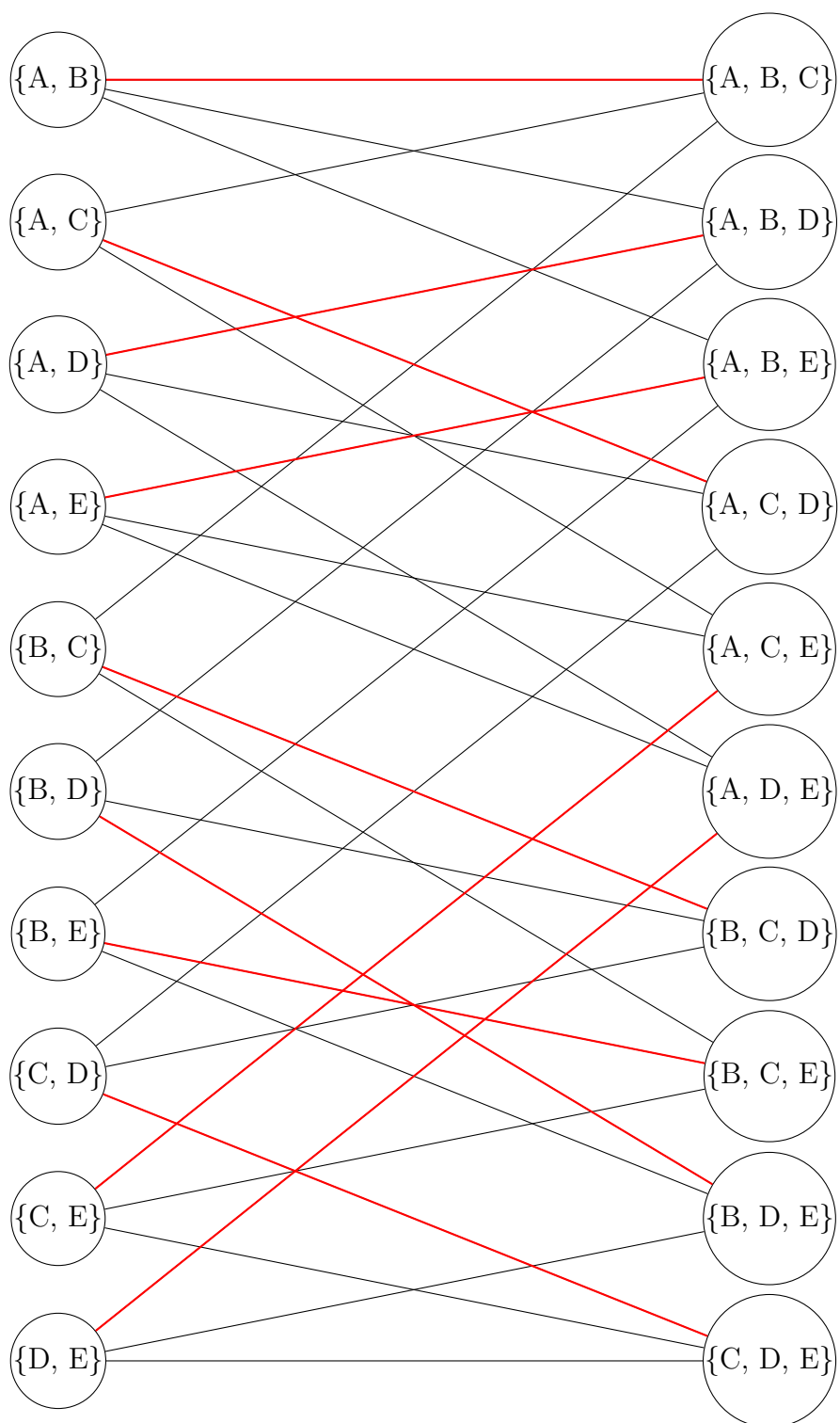
Let G be the bipartite subgraph of Q_n spanned by $X^{(r)} \cup X^{(r+1)}$: we seek a matching from $X^{(r)}$ to $X^{(r+1)}$.

For any $S \subseteq X^{(r)}$, the number of edges from S to $\Gamma(S)$ is $|S|(n - r)$, since each edge in S has $n - r$ edges.

Moreover there are at most $|\Gamma(S)|(r + 1)$ edges, counting from $\Gamma(S)$. Therefore,

$$|\Gamma(S)| \geq \frac{|S|(n - r)}{r + 1} \geq |S|.$$

So we are done, by Hall's matching theorem.



When do we have equality in Sperner's? The above proof tells us nothing.

Our aim is to prove the following: if \mathcal{A} is an antichain, then

$$\sum_{r=0}^n \frac{|\mathcal{A} \cap X^{(r)}|}{\binom{n}{r}} \leq 1.$$

In other words, the percentages of each layer occupied add up to at most 1. This trivially implies Sperner's.

1.1 Shadows

For $\mathcal{A} \subseteq X^{(r)}$, the *shadow* of \mathcal{A} is $\partial\mathcal{A} = \partial^- \mathcal{A} \subseteq X^{(r-1)}$ defined by

$$\partial\mathcal{A} = \{B \in X^{(r-1)} \mid B \subseteq A \text{ for some } A \in \mathcal{A}\}.$$

For example, if $\mathcal{A} = \{123, 124, 134, 137\}$, then

$$\partial\mathcal{A} = \{12, 13, 23, 14, 24, 34, 17, 37\}.$$

Proposition 1.1 (Local LYM). *Let $\mathcal{A} \subseteq X^{(r)}$. Then,*

$$\frac{|\partial\mathcal{A}|}{\binom{n}{r-1}} \geq \frac{|\mathcal{A}|}{\binom{n}{r}}.$$

So, the fraction of the local occupancy by $\partial\mathcal{A}$, is at least the occupancy by \mathcal{A} .

Remark. LYM = Lubell, Meshalkin, Yamamoto.

Proof: We look at the number of \mathcal{A} to $\partial\mathcal{A}$ edges in the bipartite graph Q_n ; counting from above, there are exactly $|\mathcal{A}|r$.

However counting from below, it is at most $|\partial\mathcal{A}|(n - r + 1)$. So,

$$\frac{|\partial\mathcal{A}|}{|\mathcal{A}|} \geq \frac{r}{n - r + 1} = \frac{\binom{n}{r-1}}{\binom{n}{r}}.$$

So we are done.

Remark. When do we have equality? We lose equality if an element in $\partial\mathcal{A}$ is connected to an element not in \mathcal{A} , so for this not to occur, we need that for all $A \in \mathcal{A}$, and $i \in A$, $j \notin \mathcal{A}$, that $A - \{i\} \cup \{j\} \in \mathcal{A}$.

But this is very strong, and in fact either $\mathcal{A} = \emptyset$ or $X^{(r)}$.

Theorem 1.2 (LYM Inequality). *Let $A \subseteq \mathcal{P}(X)$ be an antichain. Then,*

$$\sum_{r=0}^n \frac{|\mathcal{A} \cap X^{(r)}|}{\binom{n}{r}} \leq 1.$$

As a bit of notation, we write \mathcal{A}_r for $\mathcal{A} \cap X^{(r)}$.

We will look at two proofs. The first idea is to bubble down with local LYM.

Proof: Obviously

$$\frac{|\mathcal{A}_n|}{\binom{n}{n}} \leq 1.$$

Now, $\partial\mathcal{A}_n$ and \mathcal{A}_{n-1} are disjoint, as \mathcal{A} is an antichain. So,

$$\frac{|\partial\mathcal{A}_n|}{\binom{n}{n-1}} + \frac{|\mathcal{A}_{n-1}|}{\binom{n}{n-1}} = \frac{|\partial\mathcal{A}_n \cup \mathcal{A}_{n-1}|}{\binom{n}{n-1}} \leq 1,$$

whence we get

$$\frac{|\mathcal{A}_n|}{\binom{n}{n}} + \frac{|\mathcal{A}_{n-1}|}{\binom{n}{n-1}} \leq 1,$$

by local LYM. We now continue again. Notice $\partial(\partial\mathcal{A}_n \cup \mathcal{A}_{n-1})$ is disjoint from \mathcal{A}_{n-2} , we find

$$\frac{|\partial(\partial\mathcal{A}_n \cup \mathcal{A}_{n-1})|}{\binom{n}{n-2}} + \frac{|\mathcal{A}_{n-2}|}{\binom{n}{n-2}} \leq 1,$$

whence

$$\frac{|\partial\mathcal{A}_n \cup \mathcal{A}_{n-1}|}{\binom{n}{n-1}} + \frac{|\mathcal{A}_{n-2}|}{\binom{n}{n-2}} \leq 1.$$

We can now continue inductively.

When do we have equality? We must have had equality in each use of local LYM. Hence equality in LYM needs that the maximum r with $\mathcal{A}_r \neq \emptyset$, then $\mathcal{A}_r = X^{(r)}$.

Hence equality in Sperner needs either $\mathcal{A} = X^{(n/2)}$, if n is even, or $\mathcal{A} = X^{(\lfloor n/2 \rfloor)}$ or $X^{(\lceil n/2 \rceil)}$, for n odd.

Now time for another proof.

Proof: Choose uniformly at random a maximal chain \mathcal{C} . For any r -set A , note that

$$\mathbb{P}(A \in \mathcal{C}) = \frac{1}{\binom{n}{r}}.$$

So for our antichain \mathcal{A} ,

$$\mathbb{P}(\mathcal{C} \text{ meets } \mathcal{A}_r) = \frac{|\mathcal{A}_r|}{\binom{n}{r}},$$

as these events are disjoint. Hence, since \mathcal{C} can meet \mathcal{A} at one point at most,

$$\mathbb{P}(\mathcal{C} \text{ meets } \mathcal{A}) = \sum_{r=0}^n \frac{|\mathcal{A}_r|}{\binom{n}{r}},$$

from which we get

$$\sum_{r=0}^n \frac{|\mathcal{A}_r|}{\binom{n}{r}} \leq 1.$$

Equivalently, the number of maximal chains is $n!$, and the number through any fixed r -set is $r!(n-r)!$, so

$$\sum_r |\mathcal{A}_r| r!(n-r)! \leq n!.$$

We now return to shadows. For $\mathcal{A} \subseteq X^{(r)}$, we have

$$|\partial\mathcal{A}| \geq |\mathcal{A}| \frac{r}{n-r+1}.$$

We know that equality is rare: it only happens for $\mathcal{A} = \emptyset$, or $X^{(r)}$. What happens in between?

In other words, given $|\mathcal{A}|$, how should we choose $\mathcal{A} \subseteq X^{(r)}$ to minimise $|\partial\mathcal{A}|$?

It is believable that if $|\mathcal{A}| = \binom{k}{r}$, then we should take $\mathcal{A} = [k]^{(r)}$. In between adjacent binomials, it is believable that we should take $[k]^{(r)}$, plus some r -sets in $[k+1]^{(r)}$.

Example 1.1.

For $\mathcal{A} \subseteq X^{(3)}$ with

$$|\mathcal{A}| = \binom{8}{3} + \binom{4}{2},$$

we could take

$$\mathcal{A} = [8]^3 \cup \{9 \cup B \mid B \in [4]^{(2)}\}.$$

In some ways our set \mathcal{A} should be of minimal ‘order’, under some ordering on $X^{(r)}$.

1.2 Total Orders

Let A, B be distinct r -sets, and say $A = a_1 \dots a_r$, $B = b_1 \dots b_r$, where $a_1 < \dots < a_r$, $b_1 < \dots < b_r$.

We say that $A < B$ in the *lexographic* or *lex* ordering if for some j we have $a_i = b_i$ for all $i < j$, and $a_j < b_j$. So lex cares about small elements.

Example 1.2.

Lex on $[4]^{(2)}$ orders the elements as 12, 13, 14, 23, 24, 34.

Lex on $[6]^{(3)}$ orders the elements as

$$\begin{aligned} &123, 124, 125, 126, 134, 135, 136, 145, 146, 156, \\ &234, 235, 236, 245, 246, 256, 345, 346, 356, 456. \end{aligned}$$

We say that $A < B$ in the *colexographic* or *colex* ordering if for some j , we have $a_i = b_i$ for all $i > j$, and $a_j < b_j$. So colex cares about big elements.

Example 1.3.

Colex on $[4]^{(2)}$ orders the elements as 12, 13, 23, 14, 24, 34.

Colex on $[6]^{(3)}$ orders the elements as

$$\begin{aligned} &123, 124, 134, 234, 125, 135, 235, 145, 245, 345, \\ &126, 136, 236, 146, 246, 346, 156, 256, 356, 456. \end{aligned}$$

Note that in colex, $[n-1]^{(r)}$ is an initial segment of $[n]^{(r)}$. This is not true in lex. This allows us to view colex as an enumeration of $\mathbb{N}^{(r)}$.

Remark. $A < B$ in colex $\iff A^c < B^c$ in lex, with ground set ordering reversed.

Colex in particular may be the ordering we want to solve the above problem, minimizing $|\partial\mathcal{A}|$. Our aim will then be to show that initial segments of colex are the best for ∂ , i.e. if $\mathcal{A} \subseteq X^{(r)}$ and $\mathcal{C} \subseteq X^{(r)}$ is the initial segment of colex with $|\mathcal{C}| = |\mathcal{A}|$, then

$$|\partial\mathcal{C}| \leq |\partial\mathcal{A}|.$$

In particular,

$$|\mathcal{A}| = \binom{k}{r} \implies |\partial\mathcal{A}| = \binom{k}{r-1}.$$

1.3 Compression

The idea is to try to transform $\mathcal{A} \subseteq X^{(r)}$ into some $\mathcal{A}' \subseteq X^{(r)}$ such that:

- (i) $|\mathcal{A}'| = |\mathcal{A}|$.
- (ii) $|\partial\mathcal{A}'| \leq |\partial\mathcal{A}|$.
- (iii) \mathcal{A}' looks more like \mathcal{C} than \mathcal{A} did.

Ideally, we would like a family of such ‘compressions’

$$\mathcal{A} \rightarrow \mathcal{A}' \rightarrow \dots \rightarrow \mathcal{B},$$

such that either $\mathcal{B} = \mathcal{C}$, or \mathcal{B} is so similar to \mathcal{C} that we can directly check that

$$|\partial\mathcal{B}| \geq |\partial\mathcal{C}|.$$

The fact that colex prefers 1 to 2 inspires the following: fix $1 \leq i < j \leq n$. The *ij-compression* C_{ij} is defined as follows:

For $A \in X^{(r)}$, set

$$C_{ij}(A) = \begin{cases} A \cup i - j & \text{if } j \in A, i \notin A, \\ A & \text{else.} \end{cases}$$

For $\mathcal{A} \subseteq X^{(r)}$, set

$$C_{ij}(\mathcal{A}) = \{C_{ij}(A) \mid A \in \mathcal{A}\} \cup \{A \in \mathcal{A} \mid C_{ij}(A) \in \mathcal{A}\}.$$

So $C_{ij}(\mathcal{A}) \subseteq X^{(r)}$, and $|C_{ij}(\mathcal{A})| = |\mathcal{A}|$. Say \mathcal{A} is *ij-compressed* if $C_{ij}(\mathcal{A}) = \mathcal{A}$.

Lemma 1.1. *Let $\mathcal{A} \subseteq X^{(r)}$, and $1 \leq i < j \leq n$. Then*

$$|\partial C_{ij}(\mathcal{A})| \leq |\partial\mathcal{A}|.$$

Proof: Write \mathcal{A}' for $C_{ij}(\mathcal{A})$, and let $B \in \partial\mathcal{A}' - \partial\mathcal{A}$. We will show that $i \in B, j \notin B$, and $B \cup j - i \in \partial\mathcal{A} - \partial\mathcal{A}'$, which will show that we are done.

We have that $B \cup x \in \mathcal{A}'$, for some x , with $B \cup x \notin \mathcal{A}$. So, $i \in B \cup x$, $j \notin B \cup x$, and $(B \cup x) \cup j - i \in \mathcal{A}$.

We cannot have $x = i$, otherwise $(B \cup x) \cup j - i = B \cup j$, giving $B \in \partial\mathcal{A}$. So $i \in B$, and $j \notin B$.

Also, notice $B \cup j - i \in \partial\mathcal{A}$, since $(B \cup x) \cup j - i \in \mathcal{A}$.

Suppose $B \cup j - i \in \partial\mathcal{A}'$, so $(B \cup j - i) \cup y \in \mathcal{A}'$ for some y . We cannot have $y = i$, else $B \cup j \in \mathcal{A}'$, so $B \cup j \in \mathcal{A}$, contradicting $B \notin \partial\mathcal{A}$. So $j \in (B \cup j - i) \cup y$, and $i \notin (B \cup j - i) \cup y$.

Whence both $(B \cup j - i) \cup y$ and $B \cup y$ belong to \mathcal{A} , by definition of \mathcal{A}' , contradicting $B \notin \partial\mathcal{A}$.

Remark. We have actually shown that $\partial C_{ij}(\mathcal{A}) \subseteq C_{ij}\partial\mathcal{A}$.

Say $\mathcal{A} \subseteq X^{(r)}$ is *left-compressed* if $C_{ij}(\mathcal{A}) = \mathcal{A}$ for all $i \leq j$.

Corollary 1.1. *Let $\mathcal{A} \subseteq X^{(r)}$. Then there exists a left-compressed $\mathcal{B} \subseteq X^{(r)}$ with $|\mathcal{B}| = |\mathcal{A}|$, and $|\partial\mathcal{B}| \leq |\partial\mathcal{A}|$.*

Proof: Define a sequence $\mathcal{A}_0, \mathcal{A}_1, \dots$ as follows. Let $\mathcal{A}_0 = \mathcal{A}$.

Having defined $\mathcal{A}_0, \dots, \mathcal{A}_k$, if \mathcal{A}_k is left-compressed then we can stop the sequence with \mathcal{A}_k .

If not, choose $i < j$ such that \mathcal{A}_k is not ij -compressed, and set $\mathcal{A}_{k+1} = C_{ij}(\mathcal{A}_k)$.

This must terminate, as for example

$$\sum_{A \in \mathcal{A}_k} \sum_{i \in A} i$$

is strictly decreasing in k .

Then the final term $\mathcal{B} = \mathcal{A}_k$ satisfies that $|\mathcal{B}| = |\mathcal{A}|$, and $|\partial\mathcal{B}| \leq |\partial\mathcal{A}|$, by the previous lemma.

Remark.

1. Similarly we may choose all $\mathcal{B} \subseteq X^{(r)}$ with $|\mathcal{B}| = |\mathcal{A}|$, and $|\partial\mathcal{B}| \leq |\partial\mathcal{A}|$, and

then choose one with smallest sum of elements.

2. We can choose the order of the C_{ij} so that no C_{ij} is applied twice.
3. Any initial segment of colex is left-compressed. The converse is false, for example lex: $\{123, 124, 125, 126\}$.

This is not exactly what we want; we want to show that this is colex.

The fact that colex prefers 23 to 14 inspires the following. Let $U, V \subseteq X$ with $|U| = |V|$, $U \cap V = \emptyset$, and $\max V > \max U$.

Define the UV -compression as follows: for $A \subseteq X$,

$$C_{UV}(A) = \begin{cases} A \cup U - V & \text{if } V \subseteq A, U \cap A = \emptyset, \\ A & \text{otherwise.} \end{cases}$$

For $\mathcal{A} \subseteq X^{(r)}$, set

$$C_{UV}(\mathcal{A}) = \{C_{UV}(A) \mid A \in \mathcal{A}\} \cup \{A \in \mathcal{A} \mid C_{UV}(A) \in \mathcal{A}\}.$$

For example if $\mathcal{A} = \{123, 124, 147, 237, 238, 149\}$, then

$$C_{23,14}(\mathcal{A}) = \{123, 124, 147, 237, 238, 239\}.$$

So $C_{UV}(\mathcal{A}) \subseteq X^{(r)}$, and $|C_{UV}(\mathcal{A})| = |\mathcal{A}|$. Say \mathcal{A} is UV -compressed if $C_{UV}(\mathcal{A}) = \mathcal{A}$.

Sadly, we can have $|\partial C_{UV}(\mathcal{A})| > |\partial \mathcal{A}|$. For example if $\mathcal{A} = \{147, 137\}$, then $|\partial \mathcal{A}| = 5$, but $C_{23,14}(\mathcal{A}) = \{237, 147\}$ has $|\partial C_{23,14}(\mathcal{A})| = 6$.

We can prove the following at least:

Lemma 1.2. *Let $\mathcal{A} \subseteq X^{(r)}$ be UV -compressed for all U, V with $|U| = |V|$, $U \cap V = \emptyset$ and $\max V > \max U$. Then \mathcal{A} is an initial segment of colex.*

Proof: Suppose not. Then there exists $A, B \in X^{(r)}$ with $B < A$ in colex, but $A \in \mathcal{A}$, $B \notin \mathcal{A}$.

Set $V = A \setminus B$, $U = B \setminus A$. Then clearly $|V| = |U|$, and U, V are disjoint, with $\max V > \max U$ since $B < A$. So, $C_{UV}(A) = B$, contradicting \mathcal{A} UV -compressed.

But we can show the following:

Lemma 1.3. *Let $U, V \subseteq X$ with $|U| = |V|$, $U \cap V = \emptyset$, and $\max U < \max V$. For $\mathcal{A} \subseteq X^{(r)}$, suppose that for all u , there exists v such that \mathcal{A} is $(U - u, V - v)$ -compressed. Then,*

$$|\partial C_{UV}(\mathcal{A})| \leq |\partial \mathcal{A}|.$$

Proof: Let $\mathcal{A}' = C_{UV}(\mathcal{A})$. For $B \in \partial\mathcal{A}' - \partial\mathcal{A}$, we will show that $U \subseteq B$, $V \cap B = \emptyset$, and $B \cup V - U \in \partial\mathcal{A} - \partial\mathcal{A}'$.

We have that $B \cup x \in \mathcal{A}'$, and $B \cup x \notin \mathcal{A}$. So $U \subseteq (B \cup x)$, $V \cap (B \cup x) = \emptyset$, and $(B \cup x) \cup V - U \in \mathcal{A}$, by the definition of C_{UV} .

If $x \in U$, then there exists $y \in U$ such that \mathcal{A} is $(U - x, V - y)$ -compressed, by assumption. So from $(B \cup x) \cup V - U \in \mathcal{A}$, we have $B \cup y \in \mathcal{A}$, contradicting $B \notin \partial\mathcal{A}$.

Thus $x \notin U$, and so $U \subseteq B$, $V \cap B = \emptyset$.

We certainly have $B \cup V - U \in \partial\mathcal{A}$, as $(B \cup x) \cup V - U \in \mathcal{A}$, so we just need to show that $B \cup V - U \notin \partial\mathcal{A}'$.

Suppose that $B \cup V - U \in \partial\mathcal{A}'$, so that $(B \cup V - U) \cup w \in \mathcal{A}'$, for some w .

If $w \in U$, then we know that \mathcal{A} is $(U - w, V - z)$ -compressed for some $z \in V$, so $B \cup z \in \mathcal{A}$, contradicting $B \notin \partial\mathcal{A}$.

If $w \notin U$, we have that $V \subseteq (B \cup V - U) \cup w$, and $U \cap ((B \cup V - U) \cup w) = \emptyset$, so by definition of C_{UV} , we must have that both $(B \cup V - U) \cup w$ and $B \cup w \in \mathcal{A}$, contradicting $B \notin \partial\mathcal{A}$.

Theorem 1.3 (Kruskal-Katona). *Let $\mathcal{A} \subseteq X^{(r)}$, where $1 \leq r \leq n$, and let \mathcal{C} be the initial sequence of colex on $X^{(r)}$, with $|\mathcal{C}| = |\mathcal{A}|$. Then,*

$$|\partial\mathcal{C}| \leq |\partial\mathcal{A}|.$$

In particular, if $|\mathcal{A}| = \binom{k}{r}$, then

$$|\partial\mathcal{A}| \geq \binom{k}{r-1}.$$

Proof: Let

$$P = \{(U, V) \mid |U| = |V| > 0, U \cap V = \emptyset, \max U < \max V\} \cup \{(\emptyset, \emptyset)\}.$$

Define sets $\mathcal{A}_0, \mathcal{A}_1, \dots$ of sets systems in X as follows: set $\mathcal{A}_0 = \mathcal{A}$.

Having defined $\mathcal{A}_0, \dots, \mathcal{A}_k$, if \mathcal{A}_k is (U, V) -compressed for all $(U, V) \in P$, then we are done.

Otherwise, we have $(U, V) \in P$ with $|U| = |V| > 0$ and disjoint, such that \mathcal{A}_k is not (U, V) -compressed. Choose (U, V) minimal.

Note that for all $u \in U$, there is $v \in V$ such that $(U - u, V - v) \in P$, namely take $v = \min V$. So by the previous lemma, we get

$$|\partial C_{UV}(\mathcal{A}_k)| = |\partial \mathcal{A}_k|.$$

Set $\mathcal{A}_{k+1} = C_{UV}(\mathcal{A}_k)$, and continue. This must terminate, as

$$\sum_{A \in \mathcal{A}_k} \sum_{i \in A} 2^i$$

is strictly decreasing in k . Hence the final term \mathcal{B} satisfies $|\mathcal{B}| = |\mathcal{A}|$, $|\partial \mathcal{B}| \leq |\partial \mathcal{A}|$ and is (U, V) -compressed for all $(U, V) \in P$.

So, $\mathcal{B} = \mathcal{C}$ by lemma 1.2.

Remark.

1. Equivalently, if we write

$$|\mathcal{A}| = \binom{k_r}{r} + \binom{k_{r-1}}{r-1} + \cdots + \binom{k_s}{s},$$

where $k_r > k_{r-1} > \cdots > k_s$, and $s \geq 1$, then

$$|\partial \mathcal{A}| \geq \binom{k_r}{r-1} + \binom{k_{r-1}}{r-2} + \cdots + \binom{k_s}{s-1}.$$

2. When do we have equality in Kruskal-Katona? We can check that if $|\mathcal{A}| = \binom{k}{r}$ and $|\partial \mathcal{A}| = \binom{k}{r-1}$, then $\mathcal{A} = Y^{(r)}$ for some $Y \subseteq X$ with $|Y| = k$.
3. However, it is not true in general that if $|\partial \mathcal{A}| = |\partial \mathcal{C}|$ then \mathcal{A} is isomorphic to \mathcal{C} (isomorphism means the sets are equal up to a permutation of the ground set X).

For $\mathcal{A} \subseteq X^{(r)}$, $0 \leq r \leq n$, the *upper shadow* of \mathcal{A} is

$$\partial^+ \mathcal{A} = \{A \cup x \mid A \in \mathcal{A}, x \notin A\} \subseteq X^{(r+1)}.$$

Corollary 1.2. *Let $\mathcal{A} \subseteq X^{(r)}$, where $0 \leq r \leq n$, and let \mathcal{C} be the initial segment of lex on $X^{(r)}$ with $|\mathcal{C}| = |\mathcal{A}|$. Then,*

$$|\partial^+ \mathcal{A}| \geq |\partial^+ \mathcal{C}|.$$

Proof: From Kruskal-Katona, note $A < B$ in colex $\iff A^c < B^c$ in lex, with the ground set order reversed.

From the fact that the shadow of an initial segment is an initial segment, we get the following:

Corollary 1.3. *Let $\mathcal{A} \subseteq X^{(r)}$, and \mathcal{C} the initial segment of colex on $X^{(r)}$ with $|\mathcal{C}| = |\mathcal{A}|$. Then,*

$$|\partial^t \mathcal{C}| < |\partial^t \mathcal{A}|,$$

for all $1 \leq t \leq r$.

Proof: If $|\partial^t \mathcal{C}| \leq |\partial^t \mathcal{A}|$, then $|\partial^{t+1} \mathcal{C}| \leq |\partial^{t+1} \mathcal{A}|$ by Kruskal-Katona, since $\partial^t \mathcal{C}$ is an initial segment of colex.

So, if $|\mathcal{A}| = \binom{k}{r}$, then

$$|\partial^t \mathcal{A}| \geq \binom{k}{r-t}.$$

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