

PROJECT Specification

Task Booklet 04

COS341 2020

INTRODUCTION

- ▶ You must once again perform static semantic analysis on your SPL program's AST. For this part you must ultimately determine whether variables have been assigned a **value** before they are used.
- ▶ You will be crawling your AST again to guarantee that every use of a variable has a value at its usage point. Otherwise you must display an error and continue checking to find all errors.

VALUE CHECKING RULES

- ▶ The following slides will define the rules for your value checker.
- ▶ Each rule will be defined in terms of a grammar production where certain terminal and non-terminals will be given value checking rules.

PROGRAM

- ▶ **PROG \rightarrow CODE ; PROC_DEFS**
 - ▶ if some variable v_i has a value before some call to p_i in CODE then v_i also has a value in the definition of p_i inside PROC_DEFS
- ▶ **CODE \rightarrow INSTR ; CODE**
 - ▶ if some v_i in INSTR has a value then it also has a value in CODE

CALL

- ▶ **CALL** \rightarrow *userDefinedIdentifier*
 - ▶ Value checking must be applied to the procedure referred to by *userDefinedIdentifier* as if the procedure were inlined.

IO

- ▶ IO \rightarrow input (VAR)
 - ▶ VAR **gets** a value at this point
- ▶ IO \rightarrow output (VAR)
 - ▶ VAR must have a value

ASSIGNMENT

- ▶ $ASSIGN \rightarrow VAR = \textit{stringLiteral}$
 - ▶ **VAR** gets a value at this point
- ▶ $ASSIGN \rightarrow VAR = VAR^*$
 - ▶ VAR^* must have a value, consequently VAR gets a value
- ▶ $ASSIGN \rightarrow VAR = NUMEXPR$
 - ▶ NUMEXPR must have a value, consequently VAR gets a value
- ▶ $ASSIGN \rightarrow VAR = BOOL$
 - ▶ BOOL must have a value, consequently VAR gets a value

NUMEXPR

- ▶ **NUMEXPR** \rightarrow **VAR**
 - ▶ if VAR has a value then NUMEXPR gets a value
- ▶ **NUMEXPR** \rightarrow *integerLiteral*
 - ▶ NUMEXPR gets a value
- ▶ **NUMEXPR** \rightarrow **CALC**
 - ▶ if CALC has a value then NUMEXPR gets a value

CALC

- ▶ $CALC \rightarrow add (NUMEXPR , NUMEXPR)$
 - ▶ if both NUMEXPR have a value then CALC gets a value
- ▶ $CALC \rightarrow sub (NUMEXPR , NUMEXPR)$
 - ▶ if both NUMEXPR have a value then CALC gets a value
- ▶ $CALC \rightarrow mult (NUMEXPR , NUMEXPR)$
 - ▶ if both NUMEXPR have a value then CALC gets a value

COND_BRANCH

- ▶ COND_BRANCH \rightarrow if (BOOL) then { CODE }
- ▶ BOOL must have a value
- ▶ COND_BRANCH \rightarrow if (BOOL) then { CODE } else { CODE }
- ▶ BOOL must have a value

BOOL

- ▶ $\text{BOOL} \rightarrow \text{eq} (\text{VAR} , \text{VAR})$
 - ▶ if both VAR have a value then BOOL gets a value
- ▶ $\text{BOOL} \rightarrow (\text{VAR} < \text{VAR})$
 - ▶ if both VAR have a value then BOOL gets a value
- ▶ $\text{BOOL} \rightarrow (\text{VAR} > \text{VAR})$
 - ▶ if both VAR have a value then BOOL gets a value

BOOL

- ▶ $\text{BOOL} \rightarrow \text{not } \text{BOOL}^*$
 - ▶ if BOOL^* has a value then BOOL gets a value
- ▶ $\text{BOOL} \rightarrow \text{and} (\text{BOOL}^* , \text{BOOL}^*)$
 - ▶ if both BOOL^* have a value then BOOL gets a value
- ▶ $\text{BOOL} \rightarrow \text{or} (\text{BOOL}^* , \text{BOOL}^*)$
 - ▶ if both BOOL^* have a value then BOOL gets a value

BOOL

- ▶ **BOOL** → T
 - ▶ **BOOL gets** a value at this point
- ▶ **BOOL** → F
 - ▶ **BOOL gets** a value at this point
- ▶ **BOOL** → VAR
 - ▶ if VAR has a value then **BOOL** gets a value

COND LOOP

- ▶ **COND_LOOP** → `while (BOOL) { CODE }`
 - ▶ **BOOL** must have a value
- ▶ **COND_LOOP** →
`for(VAR=0;VAR*<VAR*;VAR*=add(VAR*;1)){CODE}`
 - ▶ all **VAR*** must have a value

EXAMPLE (1)

- ▶ The following program has no errors and no variables without a value:

```
num a; bool b;  
input(b);  
if (b) then {  
  a = 5  
} else {  
  a = 10  
};
```

```
num c; num d; num i;  
for (i = 0; i < a; i = add(i, 1)) {  
  c = 11  
};
```

0 < (5 or 10) so this will be executed at least once

```
while ((a < c)) { d = 0 };
```

(5 or 10) < 11 so this will be executed at least once

```
num b;  
p;  
output(b)  
proc p {  
  b = 10  
}
```

b was assigned a value in the call to **p**

EXAMPLE (2)

- ▶ The following program has one **error** and has one variable without a known value the rest of the program is valid:

```
num a;  
input(a);  
  
num c;  
num i;  
for (i = 0; i < a; i = add(i, 1)) {  
  c = 11  
};  
  
output(c);  
g;  
  
proc g {  
  output(a)  
}
```

← a's value is determined at runtime

← we cannot guarantee $i < a$ hence we cannot guarantee $c = 11$ is ever executed.

← we cannot determine whether c has a value because of the reasons stated above

← a was assigned a value above so this is okay

YOUR TASK

- ▶ Assume all productions and terminals have no value at the start. Apply the rules given to you above. **Only** the rules that are described with “... **must have a value** ...” should produce an error, such as:
 - ▶ VALUE ERROR [line: 4, col: 3]: cannot assign undefined value to variable
 - ▶ VALUE ERROR [line: 6, col: 7]: if statement condition cannot be undefined
- ▶ You must output messages for all of these errors that occur in the program.
- ▶ You must also mark all nodes that don't have a value in your symbol table with “No-Value” and all nodes with a value must be marked with “Has-Value”.

See slides 6,7,10 and 14 for rules with that produce error messages.

ADDITIONAL NOTES

- ▶ **If** you changed your grammar for the parser, you must find a way to adapt the given rules to you in this spec to match your modified grammar.
- ▶ **Plagiarism is not allowed!** You or your group may not use any code written by someone not within your own group.

And now: **HAPPY CODING** 😊 😊 😊