

Mutex Locks

Mutex Locks

- Used to protect critical regions and thus prevent race conditions.
- A process must acquire the lock before entering a critical section; it releases the lock when it exits the critical section.
- The `acquire()` function acquires the lock, and the `release()` function releases the lock.

```
do {  
    acquire lock  
    critical section  
    release lock  
    remainder section  
} while (true);
```

Solution to the critical-section problem using mutex locks.

Mutex Locks

- The definition of `acquire()` is as follows:

```
acquire()
{
    while (!available); /* busy wait */
    available = false;
}
```

- The definition of `release()` is as follows:

```
release()
{
    available = true;
}
```

Mutex Locks

- Calls to either `acquire()` or `release()` must be performed atomically.

Disadvantage

- Busy waiting.
- While a process is in its critical section, any other process that tries to enter its critical section must loop continuously in the call to `acquire()`.
- This type of mutex lock is also called a **spinlock** because the process “spins” while waiting for the lock to become available.
- Busy waiting wastes CPU cycles that some other process might be able to use productively.