

Advanced Microeconomics

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Lecture 11: Dynamic games with incomplete information

Essential reading:

- Gibbons, Chapter 4
 - Tadelis, Chapter 15
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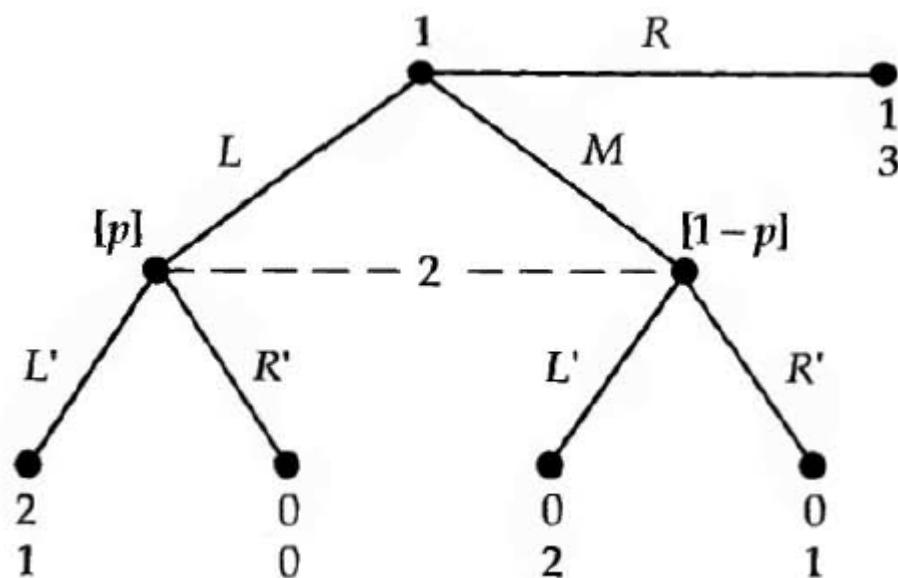
- as in dynamic games with complete information, we need to eliminate equilibria that are based on implausible threats
- We will generalize the idea of subgame perfection by considering continuation games
 - these can also start in an information set that contains more than 1 decision nodes
 - in contrast to subgames
- the player with the move at such a nonsingleton information set must have a "Belief" – i.e. a probability distribution – about which node in the information set has been reached by the play of the game
 - Remember: an information set is a collection of all the possible decision nodes that a player cannot distinguish among when called on to act.

Comparison of equilibrium concepts

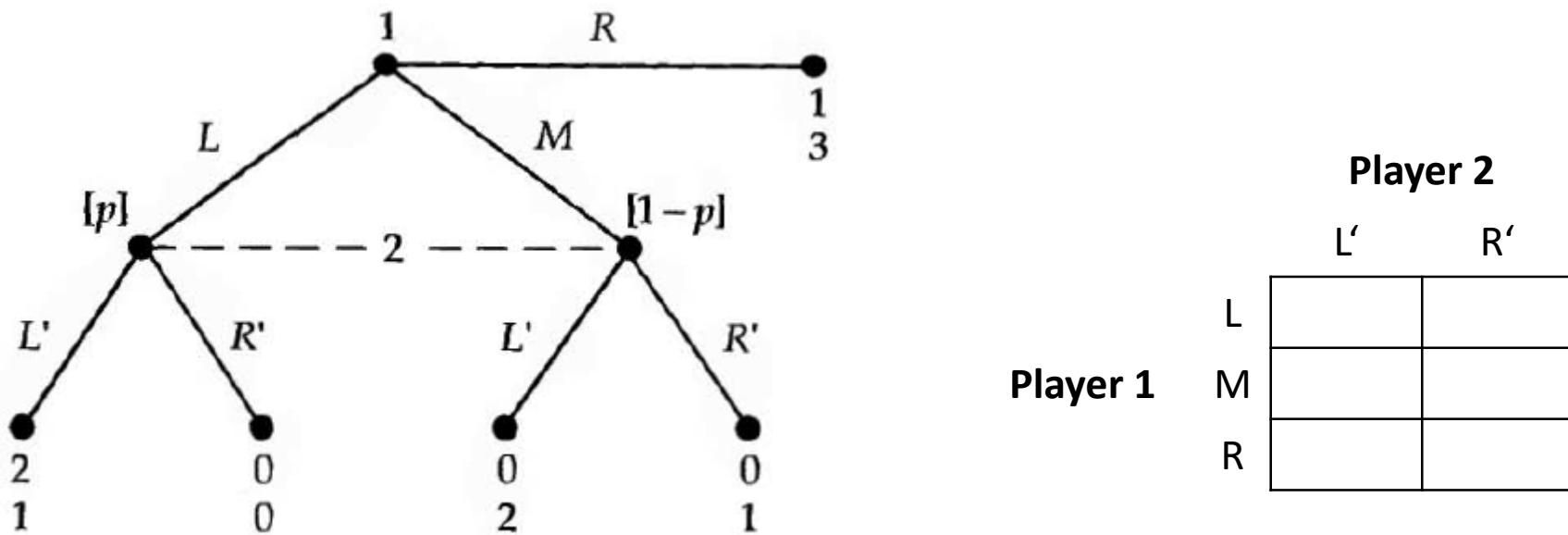
- This will lead us to the **perfect Bayesian equilibrium (PBE)**
- The PBE is consistent with
 - the SPE if it is a dynamic game with complete information,
 - the Bayesian Nash equilibrium if it is a static game with incomplete information, and
 - the Nash equilibrium if it is a static game with complete information.

Simple example: Beliefs

- Definition: A system of **beliefs** p in an extensive form game is an assignment of probabilities $p(x) \in [0,1]$ for each decision node x in the game such that $\sum_{x \in H(x)} p(x) = 1$ for all information sets H .

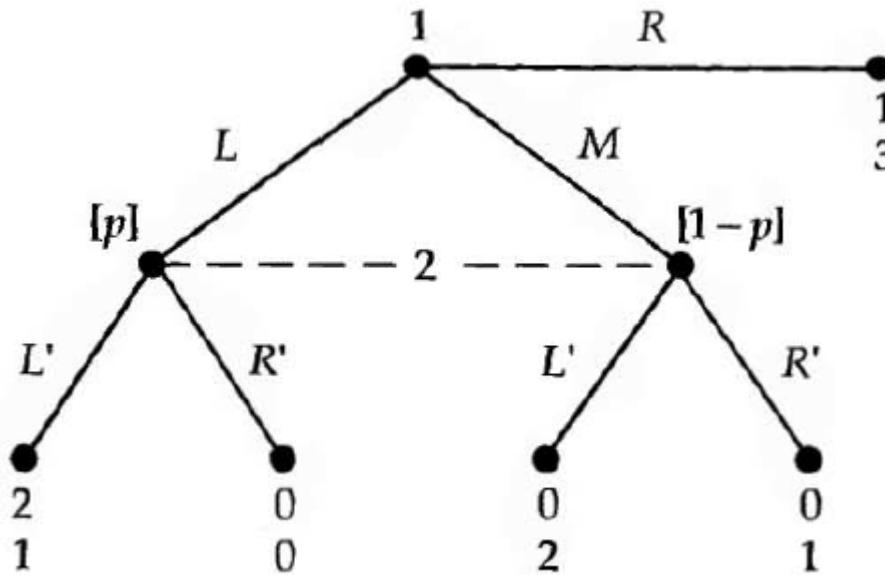


Simple example: Beliefs



- Nash-equilibrium in pure strategies:

Simple example: Beliefs



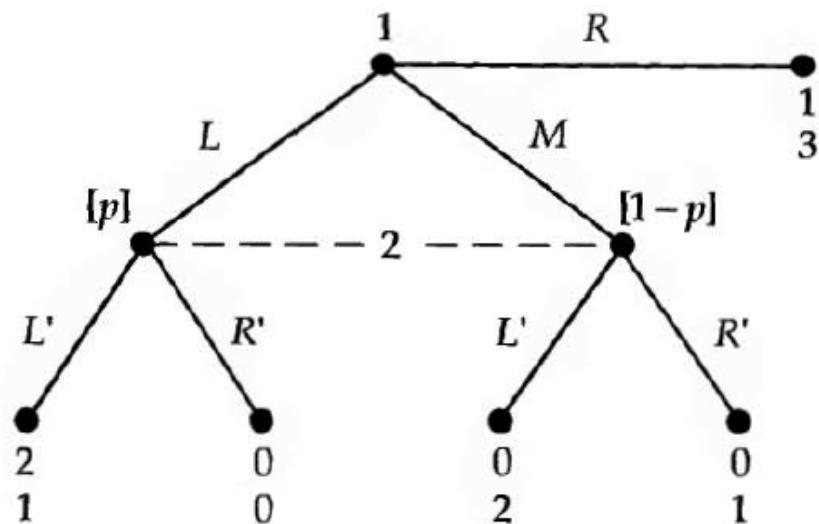
		Player 2	
		L'	R'
		L	<u>0, 0</u>
Player 1		M	0, 1
		R	<u>1, 3</u>

- Nash-equilibrium in pure strategies: (L, L') , (R, R')
- In this game, there is no proper subgame
 - Since (R, R') is a Nash equilibrium of the total game, it is also subgame perfect. But it is not sequentially rational.
 - Player 2 threatens to play R' , but for any Beliefs p , L' leads to a higher payoff.

- Additional equilibrium requirements (to exclude the SPE (R, R'))
 - **Requirement 1:** At each information set, the player with the move must have a belief about which node in the information set has been reached by the play of the game.
 - **Requirement 2:** Players' strategies must be sequentially rational (given their beliefs).
 - That is, given a player's beliefs, his strategy in any continuation game must be a best response against his opponents' strategies.

Application of requirements to example

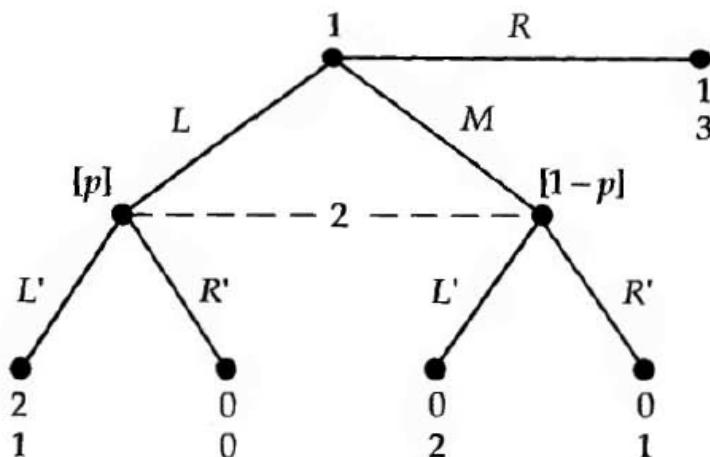
- **Requirement 1:** At each information set, the player with the move must have a belief about which node in the information set has been reached by the play of the game.
- **Requirement 2:** Players' strategies must be sequentially rational (given their beliefs).



- No matter what belief Player 2 has in his information set, it is always better to play L' than R'
 - i.e., L' is a dominant strategy
- these claims eliminate the Nash equilibrium (R, R')

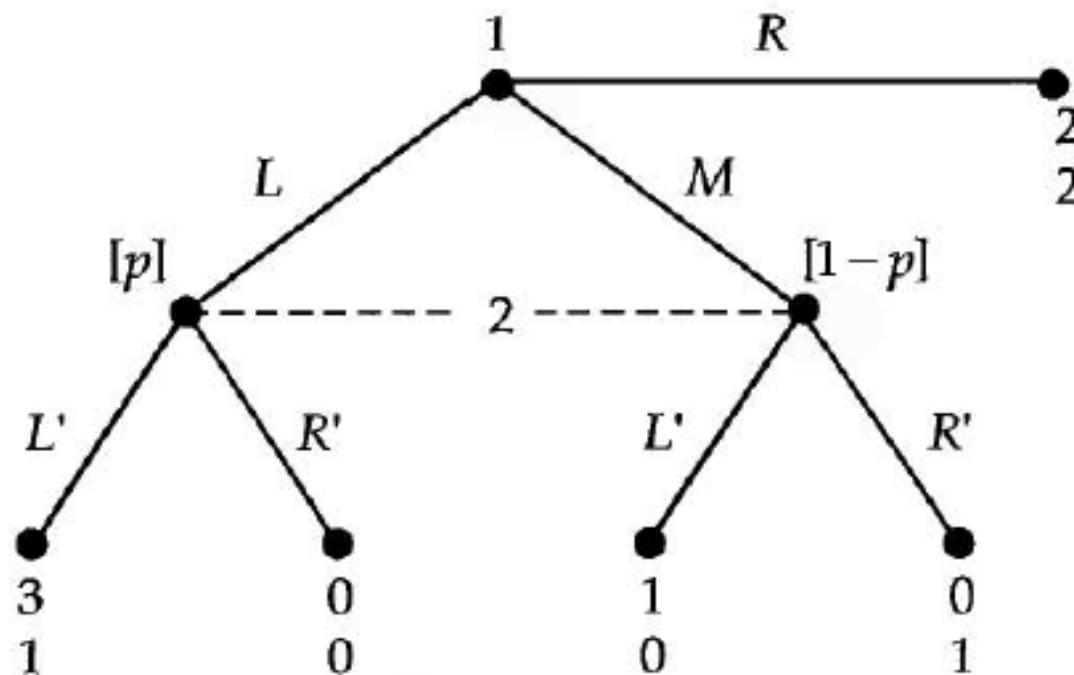
- Extension: players must not only have Beliefs (requirement 1), but beliefs must also be "reasonable".
 - Optimal action of a player depends on his beliefs
 - the beliefs must be consistent with the strategies of the players and the history of the game
- *Definition:* For a given equilibrium in a given extensive-form game, an information set is
 - **on the equilibrium path** if it will be reached with positive probability if the game is played according to the equilibrium strategies, and is
 - **off the equilibrium path** if it is certain not to be reached if the game is played according to the equilibrium strategies

- **Requirement 3:** In information sets on the equilibrium path, beliefs are determined by players' equilibrium strategies and by Bayes' rule.
- Consider the SPE (L, L') in the example. Player 2's belief must be $p = 1$
 - given the equilibrium strategy L of player 1, player 2 can anticipate which node is reached



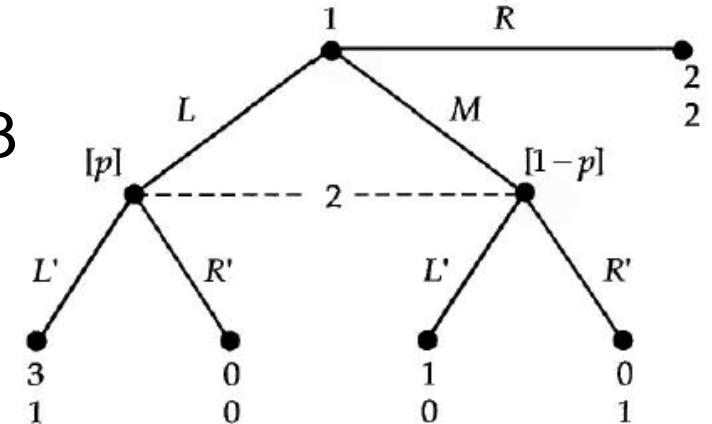
Example with new payoffs

- Suppose $p > \frac{1}{2}$. Player 2's sequentially rational strategy is L'
 - Player 1's optimal strategy in this case is L
 - Requirement 3 stipulates that beliefs are determined by players' equilibrium strategies. Thus: $p = 1$



Example with new payoffs

- Suppose $p > \frac{1}{2}$. Player 2's sequentially rational strategy is L'
 - Player 1's optimal strategy in this case is L
 - Requirement 3 stipulates that beliefs are determined by players' equilibrium strategies. Thus: $p = 1$
- Suppose $p < \frac{1}{2}$. Player 2's sequentially rational strategy is R'
 - Player 1's optimal strategy in this case is R
- The information set of the game is not reached at all, **therefore requirement 3 has no effect** (it only refers to information sets on the GG path)
- Two equilibria meet requirements 1 to 3
 - $s^* = ((L, L'), p = 1), ((R, R'), p < \frac{1}{2})$



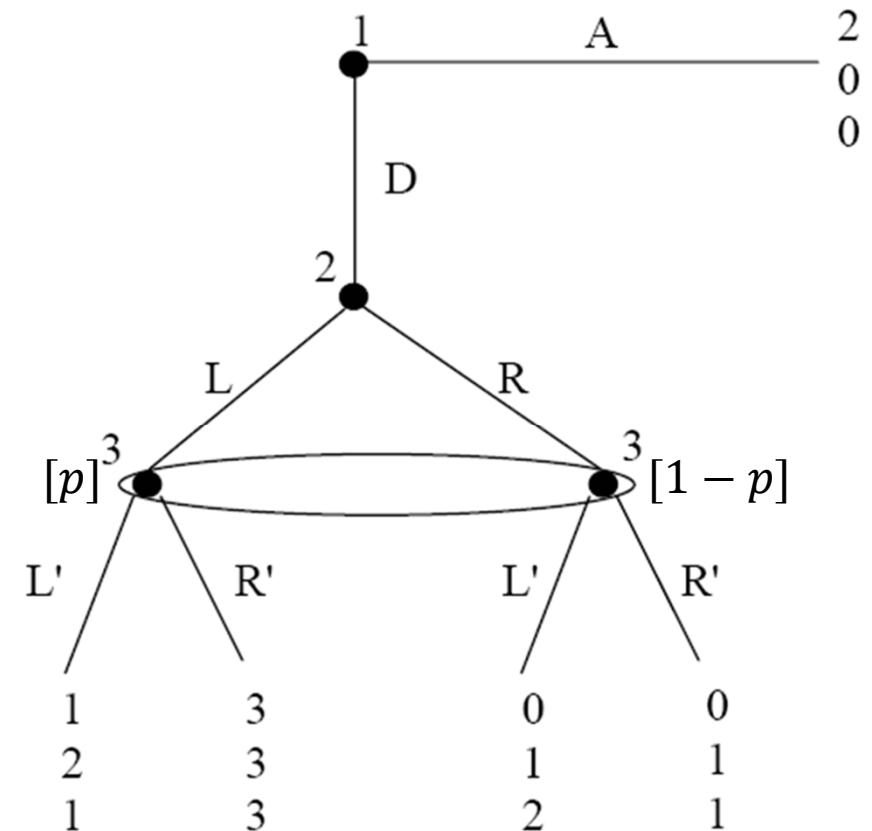
- *Definition:* A **weakly perfect Bayesian equilibrium** (weak PBE) consists of a strategy vector and a system of beliefs satisfying requirements 1 to 3.
- When strategies are not completely mixed, certain information sets are reached with probability 0.
 - In such information sets, the concept of weak PBE allows for **arbitrary beliefs**, so that often very many results can be supported as weak PBE.
 - such as $\left((R, R'), p < \frac{1}{2} \right)$ in the example above

Perfect Bayesian Equilibrium

- refinement of the equilibrium concept if some information sets are reached with probability 0
 - **Requirement 4:** At information sets *off the equilibrium path*, beliefs are determined by Bayes' rule and the players' equilibrium strategies where possible.
- *Definition:* A **perfect Bayesian equilibrium** (PBE) consists of a strategy vector and a system of beliefs satisfying requirements 1 to 4.

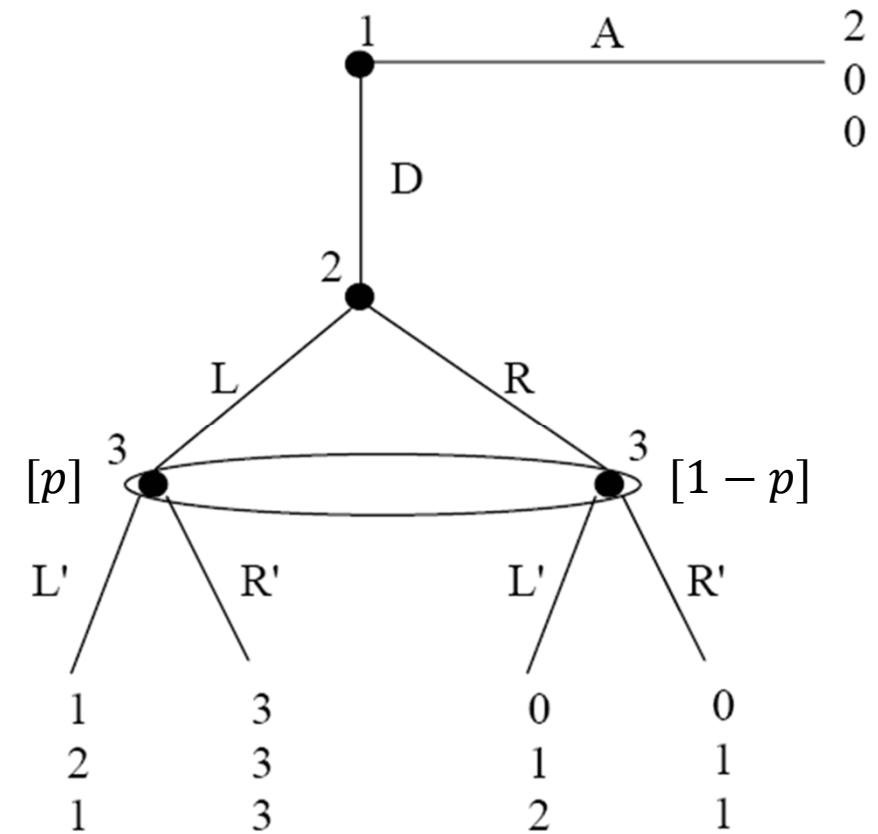
Beliefs off the GG path

- one proper subgame with a unique Nash equilibrium (L, R') :
 - since $(3,3)$ highest payoff for P2 & P3
- unique SPE: (D, L, R')
- Unique PBE: $((D, L, R'), p = 1)$: Requirements 1 to 3 are fulfilled
- Requirement 4 is trivially satisfied, since there is no information set off the equilibrium path



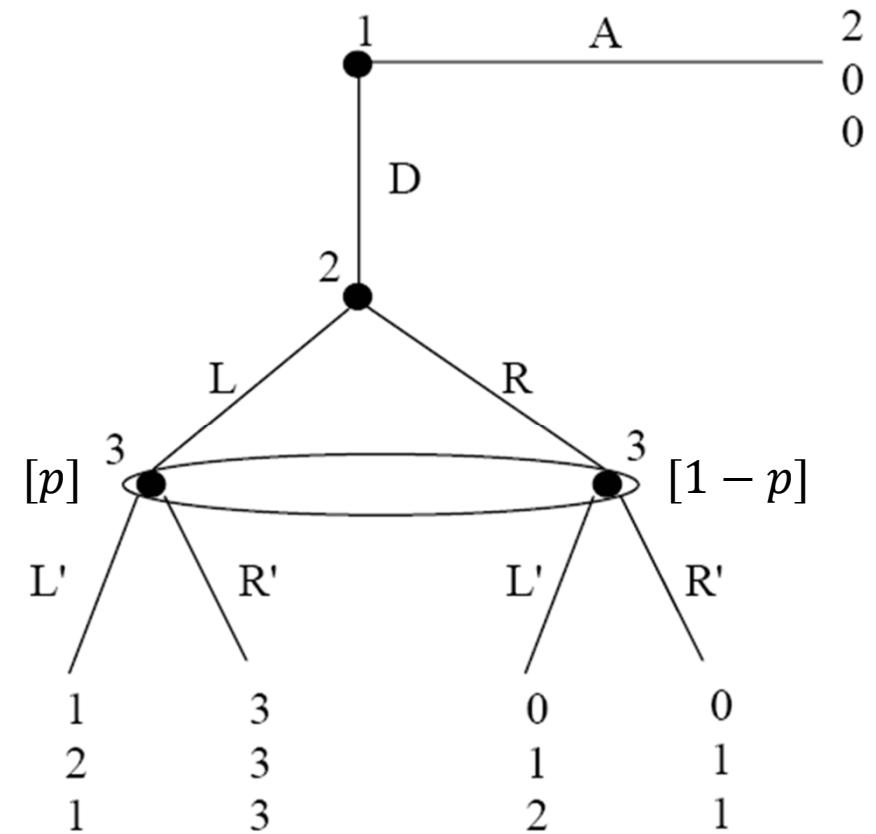
Beliefs off the GG path

- Further Nash GG: (A, L, L')
 - no one gets higher payoff from deviating unilaterally
- $((A, L, L'), p = 0)$ is a *weak PBE* as it meets requirements 1 to 3
 - Problem: belief $p = 0$ of player 3 is inconsistent with the strategy of player 2
 - but the Belief $p = 0$ concerns an information set *off the equilibrium-path*



Beliefs off the GG path

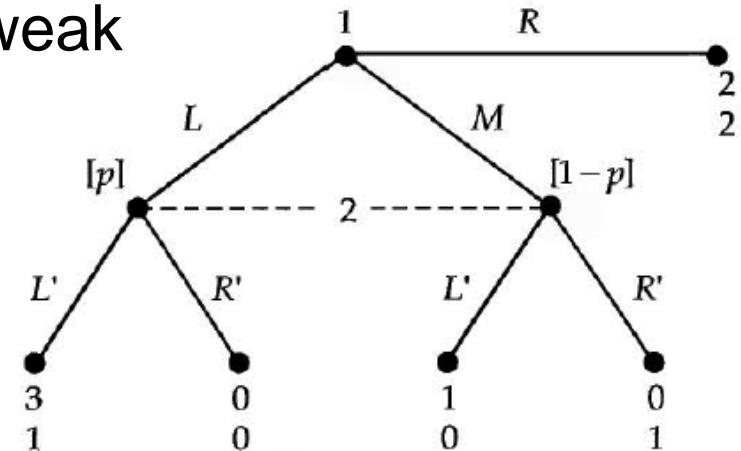
- $((A, L, L'), p = 0)$ is a weak PBE
- Requirement 4: also off the equilibrium path, Player 3 must derive his belief from player 2's strategy
- this implies $p = 1$
 - For $p = 1$, player 3 must play R' (requirement 2)
 - $((A, L, L'), p = 0)$ violates requirement 4



What does requirement 4 imply for preceding example

- In preceding example, there were 2 weak PBE: $((L, L'); p = 1)$, $((R, R'), p < \frac{1}{2})$

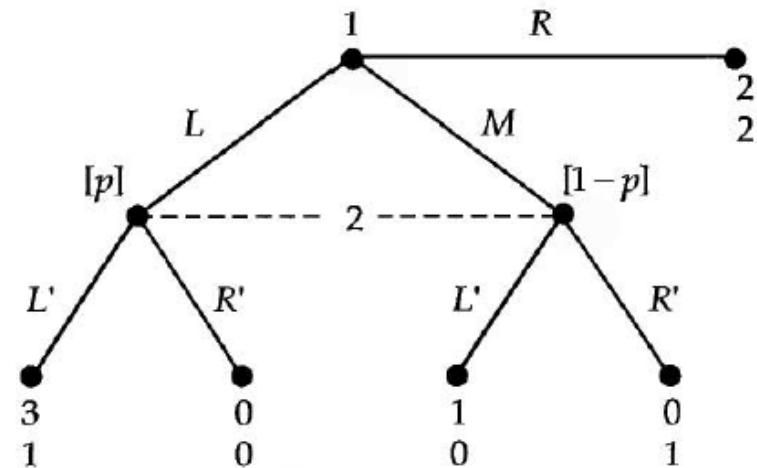
- requirement 4
 - also off the equilibrium path, beliefs must be consistent with players' strategies



- the weak PBE $((R, R'), p < \frac{1}{2})$ has an information set off the equilibrium path that concerns the beliefs of player 2
- But: in this weak PBE $((R, R'), p < \frac{1}{2})$ the strategy of player 1 has no implications for the beliefs of player 2
- Therefore, also the PBE $((R, R'), p < \frac{1}{2})$ satisfies **all** 4 requirements

Equilibrium refinements

- The literature has considered „equilibrium refinements“ to exclude equilibrium candidates that appear less plausible
 - Requirement 5:** If possible, each player's beliefs off the equilibrium path should place zero probability on nodes that are reached only if another player plays a strategy that is strictly dominated beginning at some information set.
- Consider again our PBE $\left((R, R'), p < \frac{1}{2} \right)$



Equilibrium refinements

- **Requirement 5:** If possible, each player's beliefs off the equilibrium path should place zero probability on nodes that are reached only if another player plays a strategy that is strictly dominated beginning at some information set.
- Consider again our PBE $((R, R'), p < \frac{1}{2})$
- For player 1, M is strictly dominated by R
- Thus, requirement 5 implies that $1 - p = 0$
 - As that node is only reached if player 1 plays his strictly dominated strategy M
- The weak PBE $((R, R'), p < \frac{1}{2})$ therefore violates requirement 5
- only one PBE that satisfies all 5 requirements: $s^* = ((L, L'); p = 1)$

