

# Advanced Microeconomics

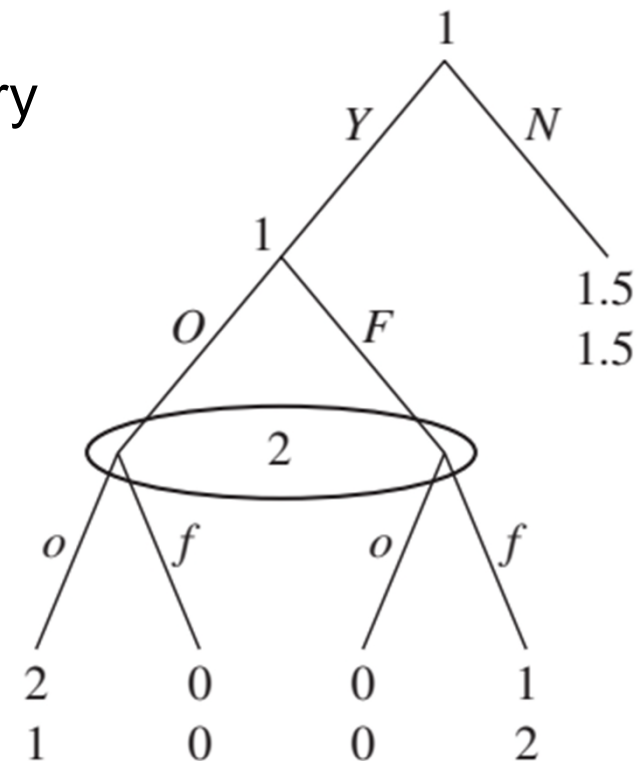
Carsten Helm

## VL 7 - Dynamic games with complete but imperfect information

- Gibbons, Chapters 2.2 and 2.4
  - Tadelis, Chapter 8
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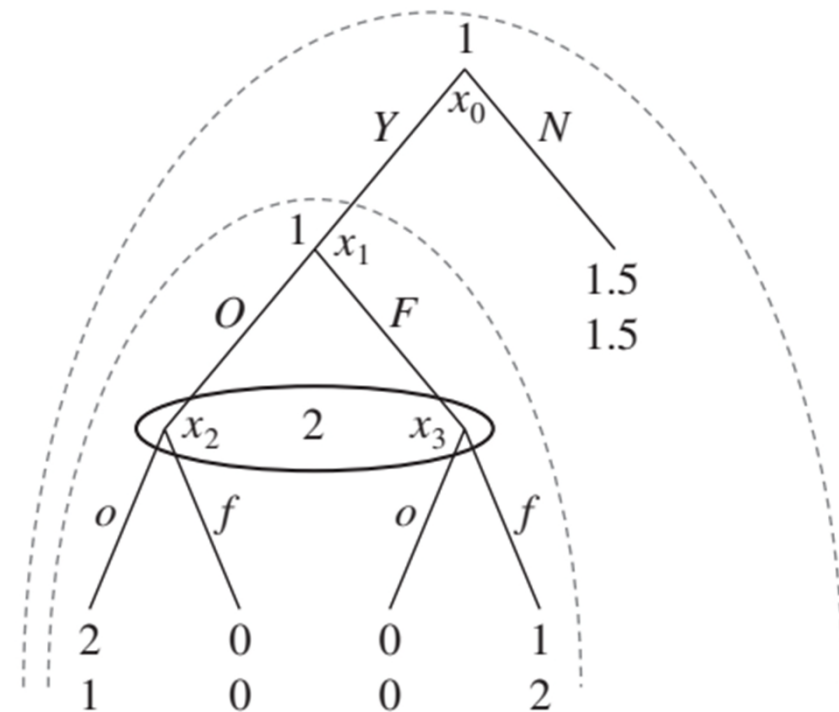
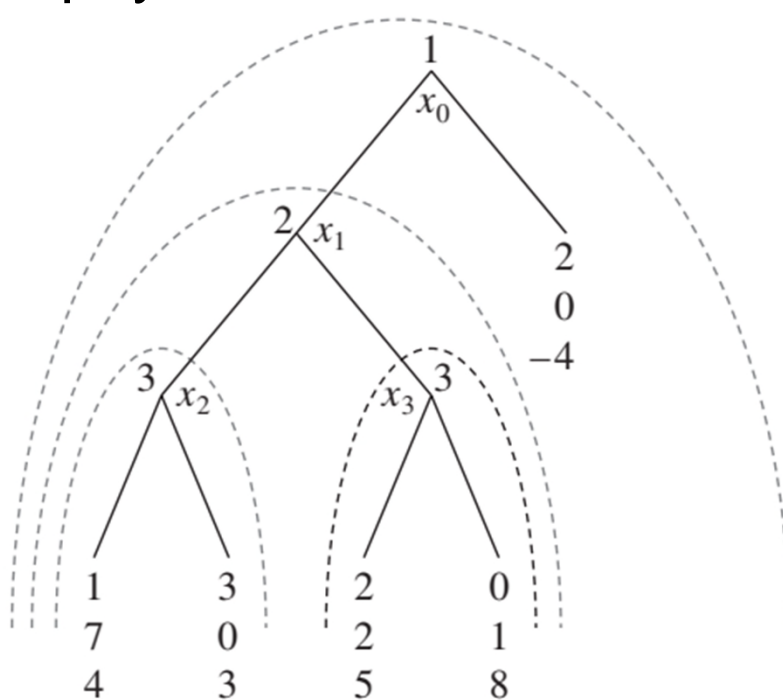
# Games with imperfect information

- Last lecture:
  - „Any finite game of perfect information has a **backward induction solution that is sequentially rational**, i.e., where players are playing rationally at every stage in the sequence of the game, whether it is on or off the equilibrium path of play.”
- Try to apply backwards induction to voluntary Battle of Sexes game
  - New: Player 1 decides Yes ( $Y$ ) or No ( $N$ ) whether to play the game
- Problem: Player 2's best response is not well defined without assigning a belief to this player about what player 1 actually choses to do



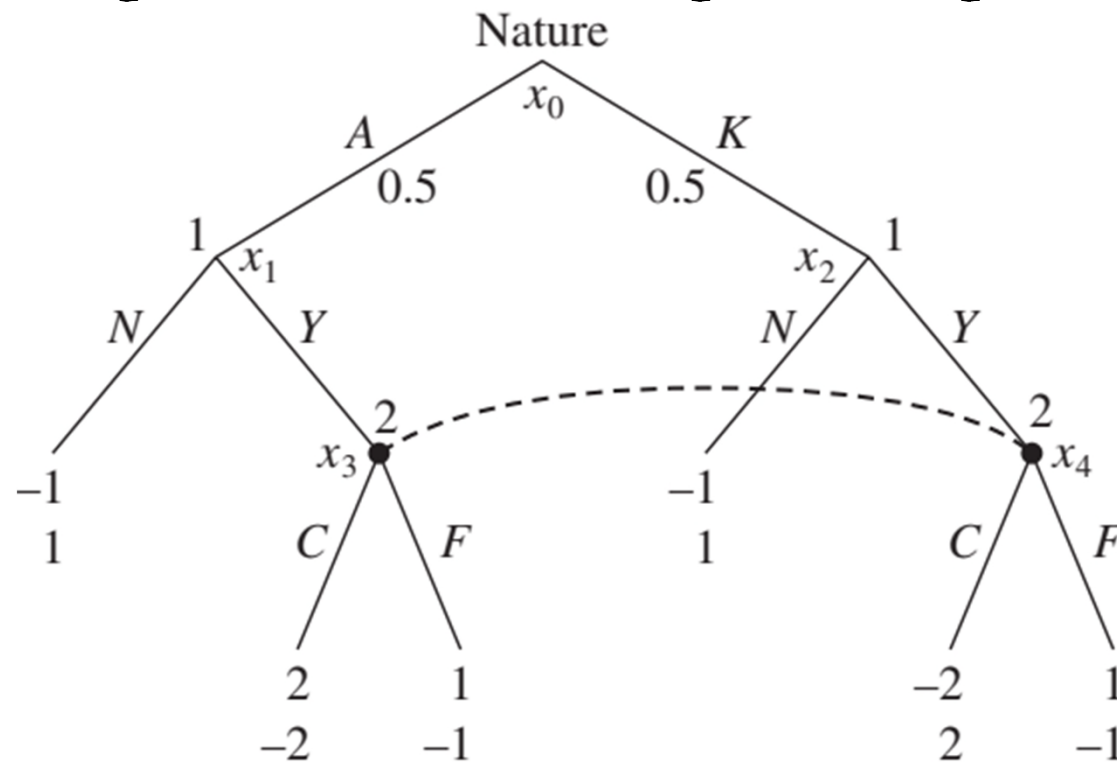
- Another problem arises if a game is not finite
  - Such games cannot be solved from behind by „**backward induction**“, simply because they have no finite set of terminal nodes
- This lecture:
  - General solution concept for **dynamic** games (finite and infinite) with **perfect and imperfect information**

**Definition.** A proper **subgame**  $G$  of an extensive-form game  $\Gamma$  consists of only a single node and all its successors in  $\Gamma$  with the property that if  $x \in G$  and  $x' \in h(x)$ , then  $x' \in G$ . The subgame  $G$  is itself a game tree with its information sets and payoffs inherited from  $\Gamma$ .



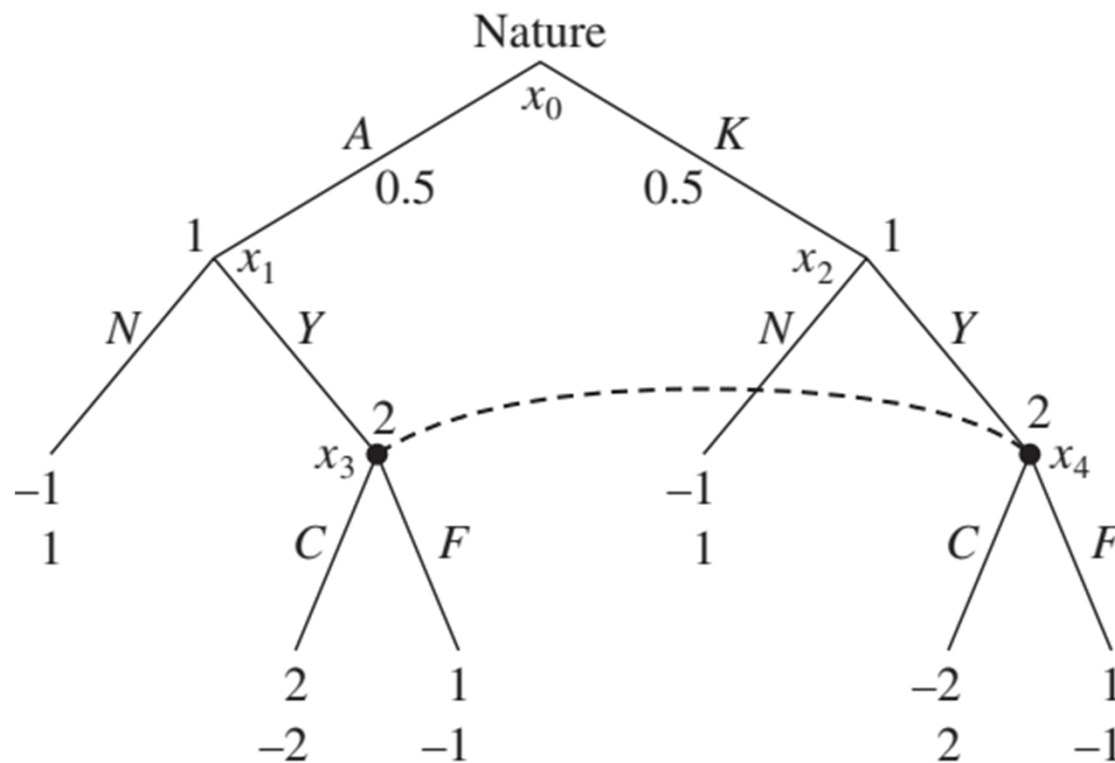
# A game of cards

- Nature pulls a card from deck of kings ( $K$ ) and aces ( $A$ )
- P1 observes card and decides whether to play game ( $Y$ ) or not ( $N$ )
- If P1 proceeds, player 2 can fold ( $F$ ) or call ( $C$ ), without knowing card of P1
  - P2 loses against an ace but wins against a king



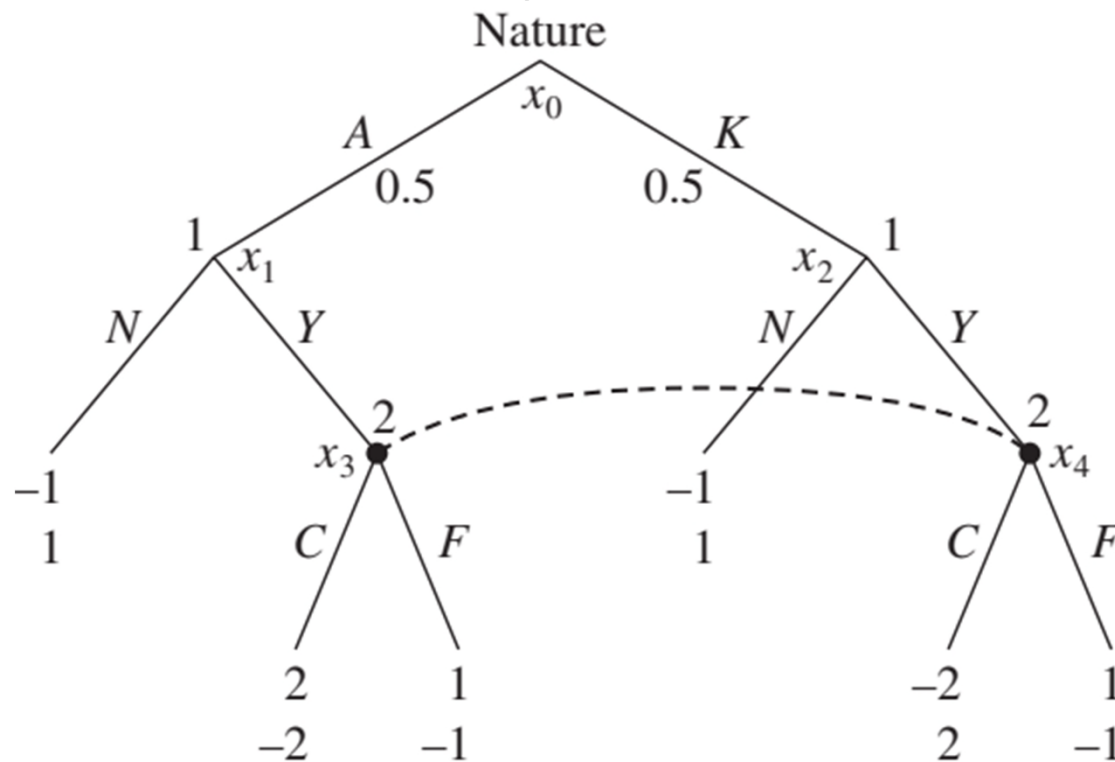
# Subgames of „game of cards“?

A proper **subgame**  $G$  of an extensive-form game  $\Gamma$  consists of only a single node and all its successors in  $\Gamma$  with the property that if  $x \in G$  and  $x' \in h(x)$ , then  $x' \in G$ . The subgame  $G$  is itself a game tree with its information sets and payoffs inherited from  $\Gamma$ .



# Subgames of „game of cards“?

- $x_1$  (and  $x_2$ ) is a single node that could start a subgame  $G$ , but its successors violate property that “if  $x_3 \in G$  and  $x_4 \in h(x_3)$ , then  $x_4 \in G$ ”. Because  $x_4$  is not a successor of  $x_1$
- Hence the only proper subgame is the complete game.

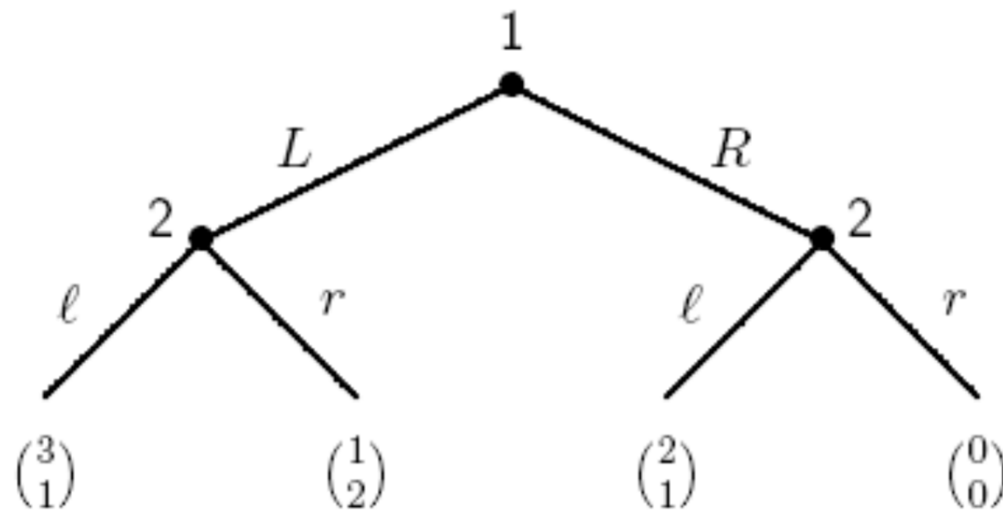


**Definition. 8.3** Let  $\Gamma$  be an  $n$ -player extensive-form game. A behavioral strategy profile  $\sigma^* = (\sigma_1^*, \sigma_2^*, \dots, \sigma_n^*)$  is a **subgame-perfect (Nash) equilibrium** if for every proper subgame  $G$  of  $\Gamma$  the restriction of  $\sigma^*$  to  $G$  is a Nash equilibrium in  $G$ .

- i.e., players' strategies constitute a Nash equilibrium in every subgame
- Hence profile of strategies must consist of mutual best responses **on and off** the equilibrium path
- introduced by Reinhard Selten (1975), who won the Nobel prize

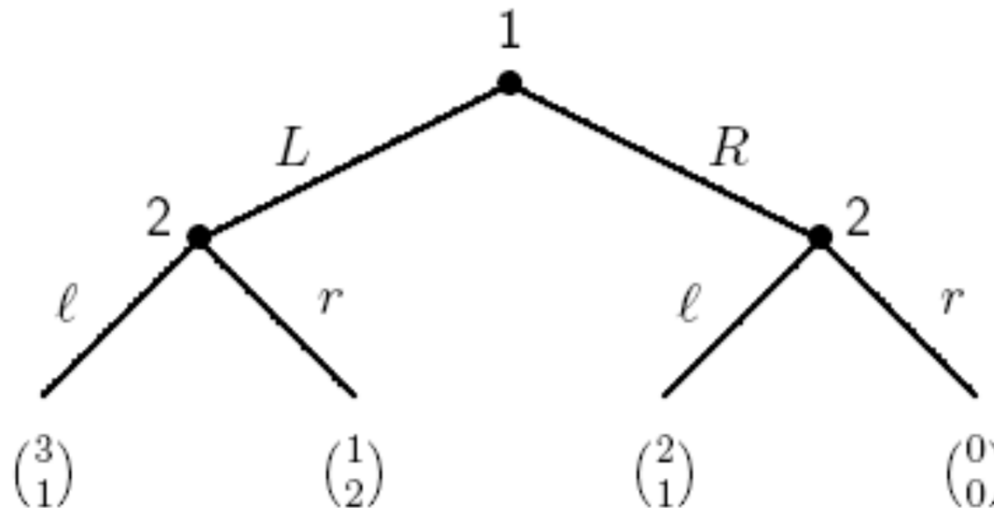


- Subgames?
- Equilibrium path?
- Subgame perfect equilibrium?



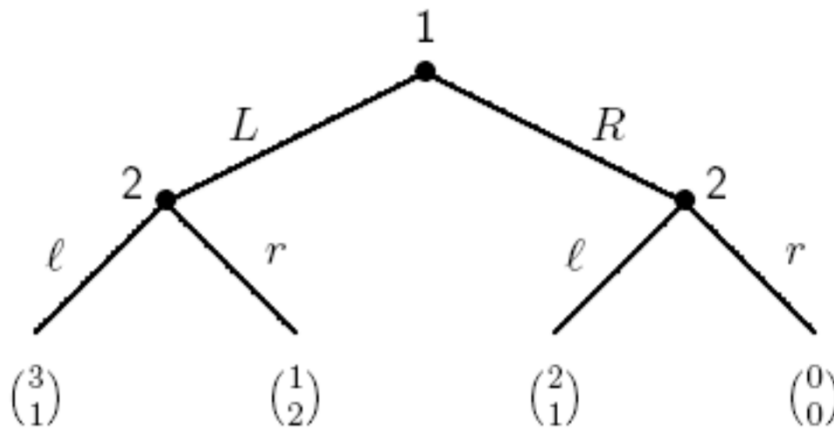
# A simple example

- Subgames: 3
- Equilibrium Path:  $(R, l)$
- subgame perfect equilibrium:  $(R, rl)$ 
  - The equilibrium must also include, what would happen off the equilibrium path
    - so that the equilibrium strategy is a complete plan of actions



# A simple example

- There is a second Nash equilibrium  $(L, rr)$ 
  - BUT this Nash equilibrium is not subgame perfect; it contains the noncredible threat that player 2 will play  $r$  should player 1 play  $R$ 
    - strategy  $ll$  of player 2 means: play  $l$  if player 1 has chosen  $L$  and also play  $l$  if player 1 has chosen  $R$



		2			
		$ll$	$lr$	$rl$	$rr$
1	$L$	3, 1	3, 1	1, 2	1, 2
	$R$	2, 1	0, 0	2, 1	0, 0

# General remarks on subgame perfection

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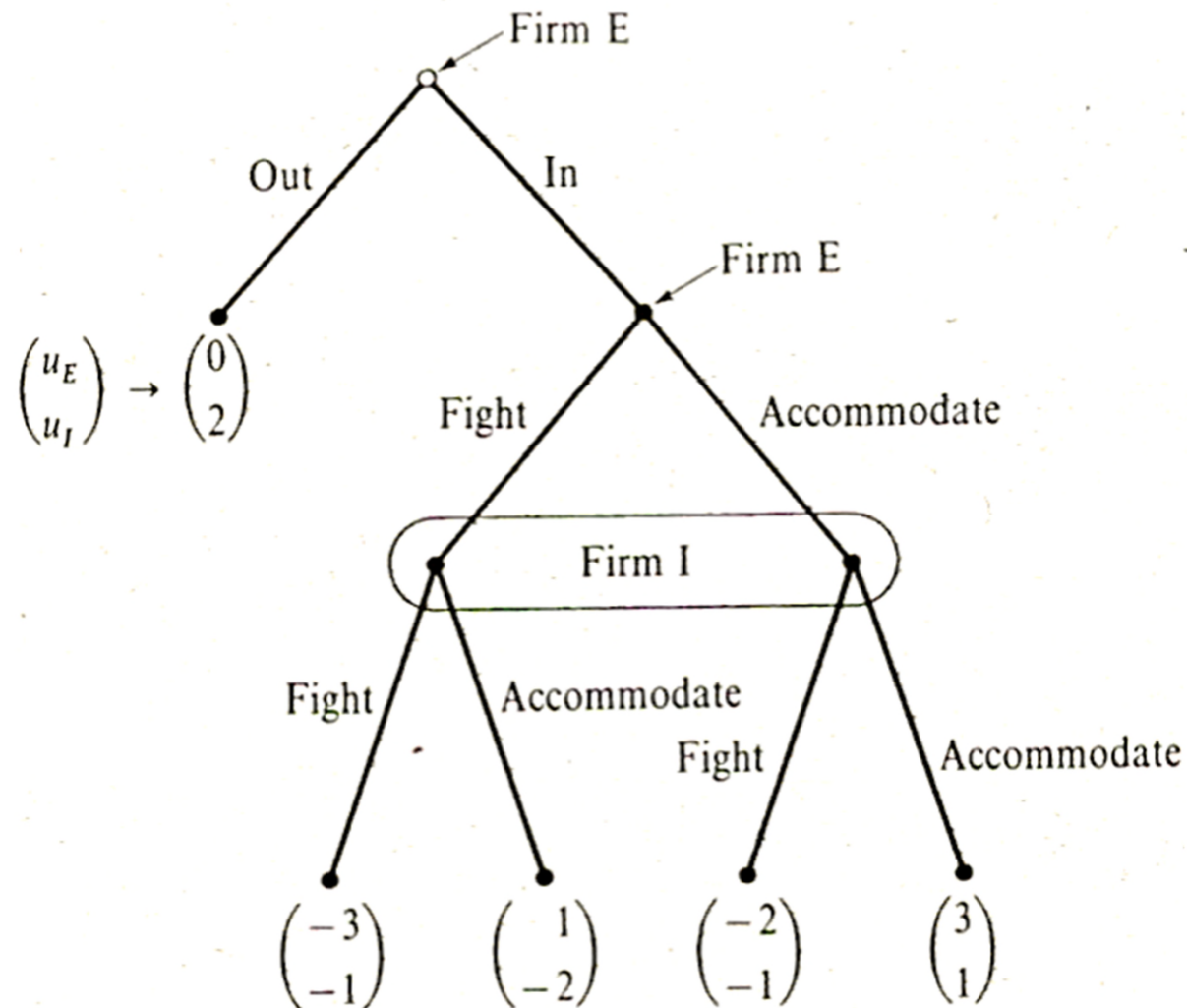
- subgame perfection is well-defined not only for games with perfect information, but also for games with imperfect information
- Every subgame perfect equilibrium is a Nash equilibrium, but not every Nash equilibrium is subgame perfect.
- Every finite game in extensive form has at least one subgame perfect Nash equilibrium
- For any finite game of perfect information, the set of subgame-perfect Nash equilibria coincides with the set of Nash equilibria that survive backward induction.

## Solution procedure: Backwards induction for games with imperfect information

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- Generalization of backward induction
  1. Start at the end of the game tree and identify the Nash equilibrium for each smallest subgame,
    - since these last subgames are finite games, there must be at least one Nash equilibrium in each of them (possibly in mixed strategies)
  2. Choose a Nash equilibrium for each of these smallest subgames and construct the reduced extensive game in which the smallest subgames are replaced by the vector of equilibrium payoffs
  3. Repeat steps 1 and 2 for the reduced game; continue this procedure until every move in the game is determined; thus a subgame perfect Nash equilibrium is found.

# Example: Market entry game



# Extensive market entry game

- Application of the solution procedure described above
  - How many subgames are there?
  - 1. Start at the end of the game tree and identify the Nash equilibrium for each smallest subgame

		Firm I	
		<i>Acc</i>	<i>Fig</i>
Firm E	<i>Acc</i>	3, 1	-2, -1
	<i>Fig</i>	1, -2	-3, -1

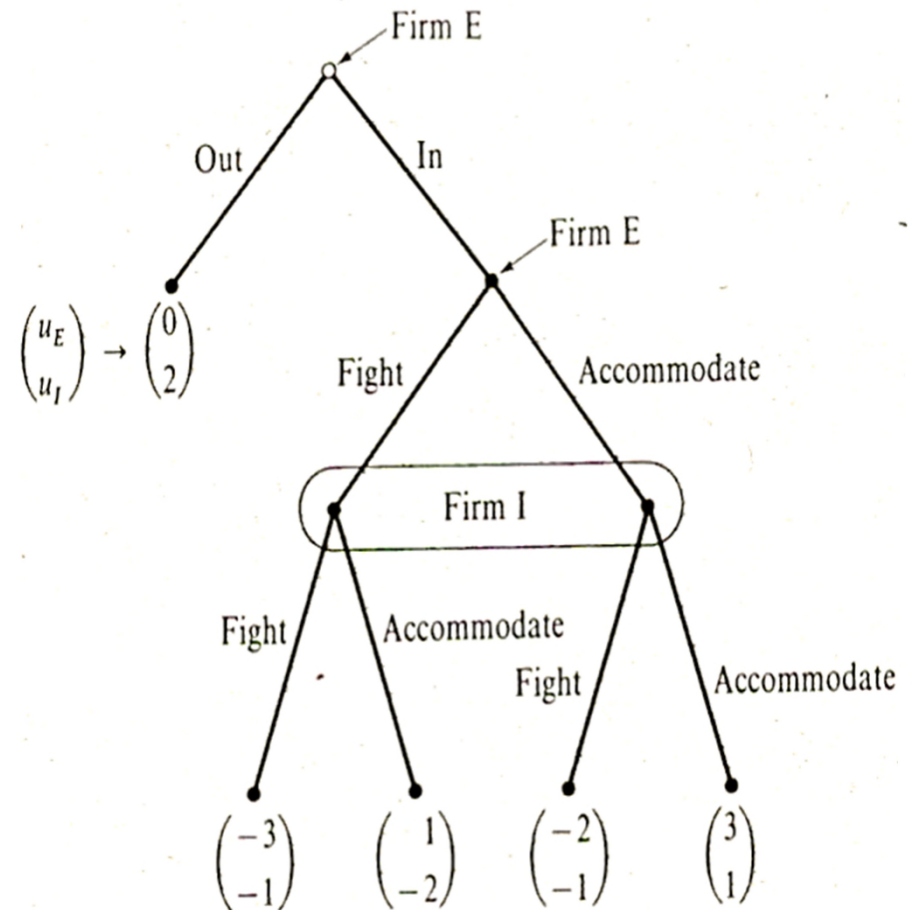
2. Choose a Nash equilibrium for each of these smallest subgames and construct the reduced extensive game in which the smallest subgames are replaced by the vector of equilibrium payoffs
  - The Nash equilibrium of the lower subgame is (*Acc*, *Acc*)

# Extensive market entry game

- Solution upper subgame
  - The best answer is *In*

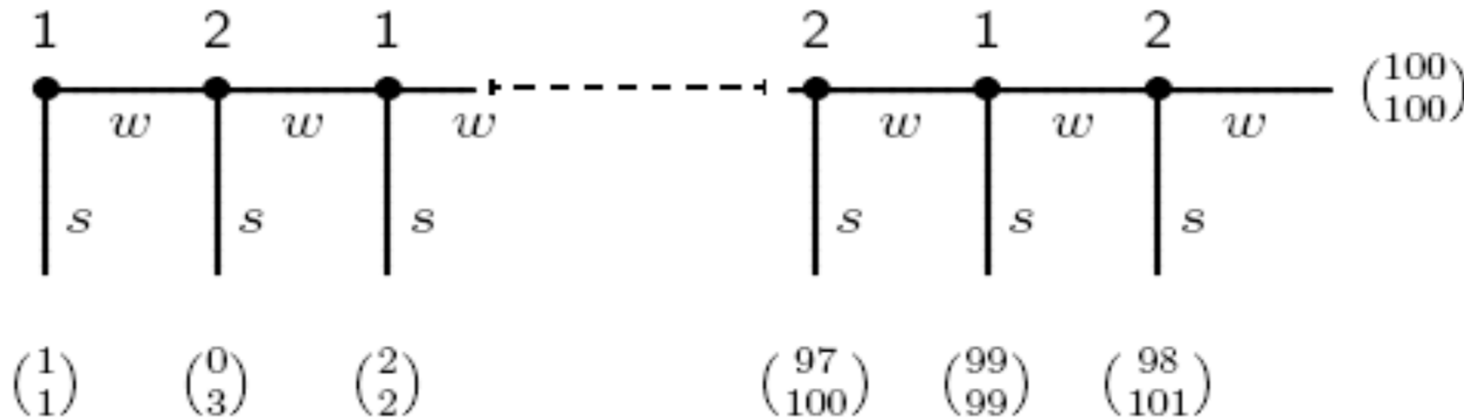
Firm E	<i>In</i>	3, 1
	<i>Out</i>	0, 2

- subgame perfect equilibrium (SPE) of the total game
  - $s^* = (In, Acc; Acc)$





# Example: Rosenthal's centipede



- Players 1 and 2 take turns and can decide each time whether the game should end ( $s$ ) or continue ( $w$ ).
- unique SPE: each player finishes the game when he has to move
- What happens if player 1 does not end the game in period 1?
  - Should player 2 still stick to the equilibrium strategy?
- Whether this SPE is convincing depends crucially on how players interpret deviations from the equilibrium path