

Advanced Microeconomics

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VL 10 - Repeated games

- Gibbons, Chapter 2
 - Osborne, Chapters 14 and 15
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- Games in which *the same stage-game* is being played at every stage
- All players are informed about the result of the last round (stage) before the next round (stage) starts
- Repeated games are actually a special case of dynamic games
- However, they have some interesting properties that do not apply to general dynamic games
 - In particular, threats or promises about behavior in the future can influence present behavior

Finitely often repeated games

- Definition
 - Let $G = (I, \{S_i\}, \{u_i\})$ be the stage-game in normal form. Then $G(T, \delta)$ denotes the **finitely repeated game** in which the stage-game G is played T consecutive times and δ is the **common discount factor**.
- Example: Prisoner's dilemma is the stage game G
 - (O, L) is the unique Nash equilibrium of the stage-game with payoff $(1, 1)$

		Player 2	
		L	R
Player 1	O	1, 1	5, 0
	U	0, 5	4, 4

Finitely often repeated games

- Prisoner's dilemma played 2 times (no discounting)
 - Backward induction: unique Nash equilibrium of the lowest subgame (stage-game) is (O_2, L_2) with payoff $(1, 1)$

		Player 2	
		L_2	R_2
Player 1	O_2	1, 1	5, 0
	U_2	0, 5	4, 4

- then the „reduced“ game at the first stage is (index stands for stage)

		Player 2	
		L_1	R_1
Player 1	O_1	2, 2	6, 1
	U_1	1, 6	5, 5

- Prisoner's dilemma played 2 times
 - The unique Nash equilibrium of the reduced game is (O_1, L_1)
 - the unique SPE of the once repeated prisoner's dilemma is $((O_1, O_2), (L_1, L_2))$
- If the stage-game G has a unique Nash equilibrium, then the finitely often repeated game $G(T, \delta)$ has a unique SPE
 - The Nash equilibrium of the stage-game is played at every stage
 - In this class of games, SPE are independent of the "history" of the game.

- Definition

- Let $G = (I, \{S_i\}, \{u_i\})$ be the stage-game in normal form. Then the payoffs of the **infinitely repeated game** are defined as follows: Let $\pi_1, \pi_2, \pi_3, \dots$ be the payoffs at each stage. Then the payoff for the total game (i.e., the present value of the infinite sequence of payoffs $\pi_1, \pi_2, \pi_3, \dots$) is:

$$\pi_1 + \delta\pi_2 + \delta^2\pi_3 \dots = \sum_{t=1}^{\infty} \delta^{t-1}\pi_t,$$

where $\delta \in (0,1)$ is the discount factor.

- Sometimes such a game is written as $G(\infty, \delta)$, or simply $G(\delta)$

- In the above Prisoner's Dilemma, can cooperation (U, R) occur at every stage of the infinitely repeated game?
 - Even though the only Nash equilibrium of the stage-game is (O, L) ?
- Claim: For a sufficiently large discount factor δ , the reciprocal play of the following **(grim) trigger strategy** is a SPE of the infinitely repeated prisoner's dilemma:
 - Player 2: "Play R at the first stage ($t = 1$). Play R at the stage $t \geq 2$ if the result of all previous $t - 1$ stages was (U, R) , otherwise play L ."
 - Player 1 with equivalent strategy for actions U and O

		Player 2	
		L	R
Player 1	O	1, 1	5, 0
	U	0, 5	4, 4

- Proof - Step 1

- Show that there is a δ (sufficiently large) such that reciprocal play of this strategy is a Nash equilibrium of the infinitely repeated prisoner's dilemma
 - i.e., of the total game
- a) Suppose both players stick to the **(grim) trigger strategy**, then (U, R) is played at each stage.
 - Each player receives the sequence of payoffs 4, 4, 4,
 - The present value of this payoff stream is

Infinitely repeated games

- Proof - Step 1

- Show that there is a δ (sufficiently large) such that reciprocal play of this strategy is a Nash equilibrium of the infinitely repeated prisoner's dilemma

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- a) Suppose both players stick to the **(grim) trigger strategy**, then (U, R) is played at each stage.

- Each player receives the sequence of payoffs 4, 4, 4,

- The present value of this payoff stream is

$$4 + 4\delta + 4\delta^2 + \dots = 4 \sum_{t=1}^{\infty} \delta^{t-1} = \frac{4}{1-\delta}$$

- Proof - Step 1 continued

- b) Assuming player 1 sticks to the (grim) trigger strategy, but player 2 deviates

- In the event of deviation, her payoff sequence would be 5, 1, 1, ...
 - The present value of its payoff sequence is

- Deviation is not worthwhile if

$$\frac{4}{1-\delta} \geq 5 + \frac{\delta}{1-\delta}$$

- Proof - Step 1 continued

- b) Assuming player 1 sticks to the (grim) trigger strategy, but player 2 deviates

- In the event of deviation, her payoff sequence would be 5, 1, 1, ...
 - The present value of its payoff sequence is

$$5 + \delta + \delta^2 + \dots = 5 + \sum_{t=2}^{\infty} \delta^{t-1} = 5 + \frac{\delta}{1 - \delta}$$

- Deviation is not beneficial if

$$\frac{4}{1 - \delta} \geq 5 + \frac{\delta}{1 - \delta} \Rightarrow 4 \geq 5 - 5\delta + \delta \Rightarrow \delta \geq \frac{1}{4}$$

Proof - Step 2: Show that the (grim) trigger strategies induce a Nash equilibrium in each subgame.

- Each subgame of the infinitely repeated game is identical to the whole game.
- these subgames can be split into two groups:
 - 1) Subgames in which the results of all previous periods were (U, R) :
 - If $\delta \geq \frac{1}{4}$, then the (grim) trigger strategy is a Nash equilibrium of this subgame.
 - this is exactly what we have shown above
 - 2) subgames in which the result of at least one previous stage was not (U, R) :
 - in these subgames, the (grim) trigger strategies of the players dictate playing (O, L) at each stage.
 - Deviating unilaterally isn't worth it, so the (grim) trigger strategy is also a Nash equilibrium of this subgame

- Alternative popular strategy: "**Tit-for-tat**"
 - Play 'cooperation' in period 1 and then always replicate the behavior of the other player in the previous period, i.e.
 - if the other player is cooperative, you behave cooperatively in the following round
 - if the other player is non-cooperative, you also behave non-cooperatively in the following round
 - For δ sufficiently large, these punishment strategies also lead to mutual cooperation in the SPE

- What payoffs can you get in an equilibrium of any infinitely repeated game?
- The Folk Theorem or Friedman Theorem (1971) gives information about this
 - Definition: feasible payoffs
 - Definition: Average payoffs
 - Folk Theorem
 - Proof

- Definition *feasible*

- The payoffs (x_1, \dots, x_n) of players $i = 1, \dots, n$ are called feasible if they can be described as a **convex combination** of the payoff combinations of the stage-game
 - Convex combinations are the **weighted average**, where the weights must be non-negative and sum to 1
- Fig. shows **convex combination** of the payoff combinations of the stage-game

		Player 2	
		<i>L</i>	<i>R</i>
Player 1	<i>O</i>	1, 1	5, 0
	<i>U</i>	0, 5	4, 4

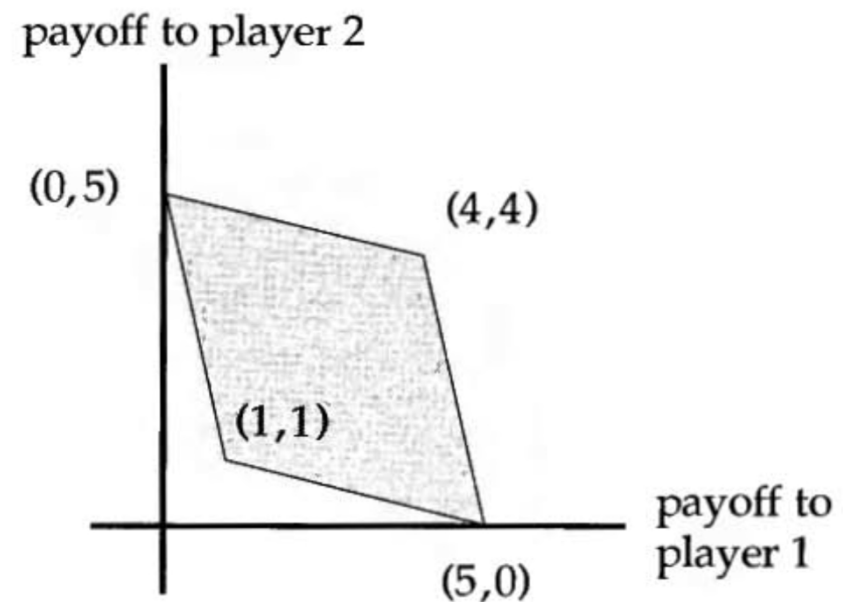
- Definition *feasible*

- The payoffs (x_1, \dots, x_n) of players $i = 1, \dots, n$ are called feasible if they can be described as a **convex combination** of the payoff combinations of the stage-game

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		<i>L</i>	<i>R</i>
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	<i>U</i>	0, 5	4, 4



- Definition of **average payoff**

- The **present value** of the payoff sequence π_1, π_2, \dots is

$$\pi_1 + \delta\pi_2 + \delta^2\pi_3 + \dots = \sum_{t=1}^{\infty} \delta^{t-1}\pi_t$$

- For $\pi_t = \pi = \text{constant}$, we have....

$$\text{so that } (1 - \delta) \sum_{t=1}^{\infty} \delta^{t-1}\pi = \pi$$

- In general, $(1 - \delta) \sum_{t=1}^{\infty} \delta^{t-1}\pi_t$ is called the **average payoff** of the sequence of payoffs π_1, π_2, \dots
- The average payoff is therefore merely a rescaling of the present value $\sum_{t=1}^{\infty} \delta^{t-1}\pi_t$ with the factor $1 - \delta$
 - Therefore, maximizing the average payoff is equivalent to maximizing the present value
- Advantage of using the average payoff: it is directly comparable to the payoffs from the stage-game

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$$\pi_1 + \delta\pi_2 + \delta^2\pi_3 + \dots = \sum_{t=1}^{\infty} \delta^{t-1}\pi_t$$

- For $\pi_t = \pi = \text{constant}$, we have $\sum_{t=1}^{\infty} \delta^{t-1}\pi_t = \pi \sum_{t=1}^{\infty} \delta^{t-1} = \frac{\pi}{1-\delta}$

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- **Folk Theorem (Friedman, 1971)**

- Let G be a finite, static game with complete information. Let (e_1, \dots, e_n) be the payoffs of a Nash equilibrium of G and let (x_1, \dots, x_n) be a feasible payoff of G . If $x_i > e_i$ for all $i = 1, \dots, n$ and δ is sufficiently close to 1, then there exists a SPE of the infinitely repeated game $G(\infty, \delta)$ that achieves an average payoff (x_1, \dots, x_n) .

- **Implication**

- With δ sufficiently large, we can achieve a SPE with better payoffs than in the Nash-equilibrium of the stage-game
 - E.g. mutual cooperation in the prisoner's dilemma

- An infinitely repeated game has an infinite number of SPE with extremely different outcomes
 - Therefore, the Folk theorem has little predictive power
 - "Almost all" payoffs can be achieved as a SPE if players are sufficiently patient
- Question: Is the game with alternating offers from the last lecture a repeated game?
- Nice link to play: <https://ncase.me/trust/>

- The sum of a geometric sequence can also be represented as a fraction
- Consider the following infinite geometric sequence

$$\pi + \pi\delta + \pi\delta^2 + \pi\delta^3 + \dots = \pi(1 + \delta + \delta^2 + \dots) = \pi \sum_{t=0}^{\infty} \delta^t.$$

- Define x as the value of the geometric sequence

$$x := 1 + \delta + \delta^2 + \dots$$

- Multiply both sides by δ

$$\delta x = \delta + \delta^2 + \delta^3 + \dots$$

- Subtraction of the two equations gives:

$$(1 - \delta)x = 1 \rightarrow x = \frac{1}{1 - \delta}$$