빅데이터 분석 및 응용

L03: MapReduce Data Flow

Summer 2020 Kookmin University

Map-Reduce: Environment

Map-Reduce environment takes care of:

- Partitioning the input data
- Scheduling the program's execution across a set of machines
- Performing the group by key step
- Handling machine failures
- Managing required inter-machine communication

Map-Reduce: A diagram

MAP:

Read input and produces a set of key-value pairs

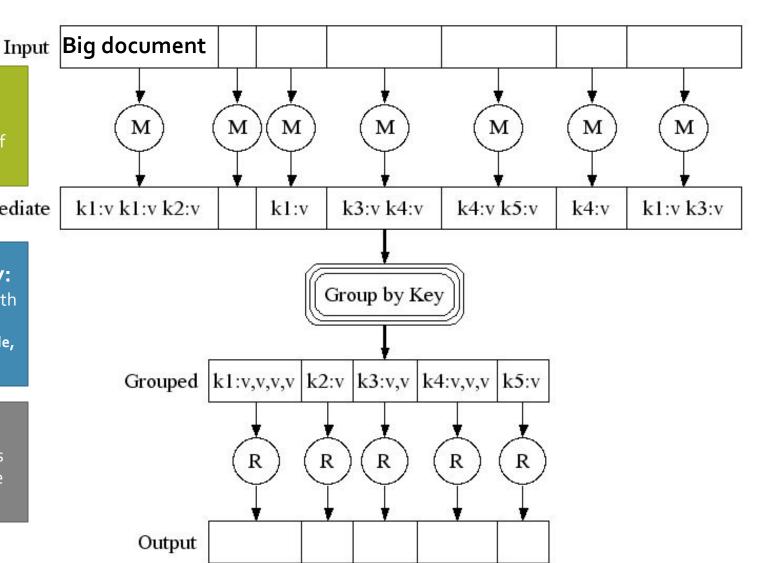
Intermediate

Group by key:

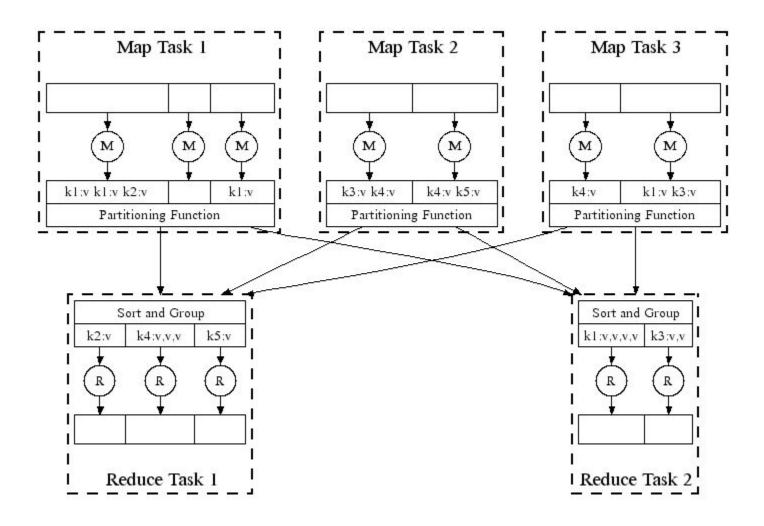
Collect all pairs with same key (Hash merge, Shuffle, Sort, Partition)

Reduce:

Collect all values belonging to the key and output



Map-Reduce: In Parallel



Data Flow

- Input and final output are stored on a distributed file system (DFS):
 - Scheduler tries to schedule map tasks "close" to physical storage location of input data
- Intermediate results are stored on local FS of Map and Reduce workers
- Output is often input to another MapReduce task

Coordination: Master

- Master node takes care of coordination:
 - Task status: (idle, in-progress, completed)
 - Idle tasks get scheduled as workers become available
 - When a map task completes, it sends the master the location and sizes of its R intermediate files, one for each reducer
 - Master pushes this info to reducers
- Master pings workers periodically to detect failures

Dealing with Failures

Map worker failure

- Map tasks completed or in-progress at worker are reset to idle
- Reduce workers are notified when task is rescheduled on another worker

Reduce worker failure

- Only in-progress tasks are reset to idle
- Reduce task is restarted

Master failure

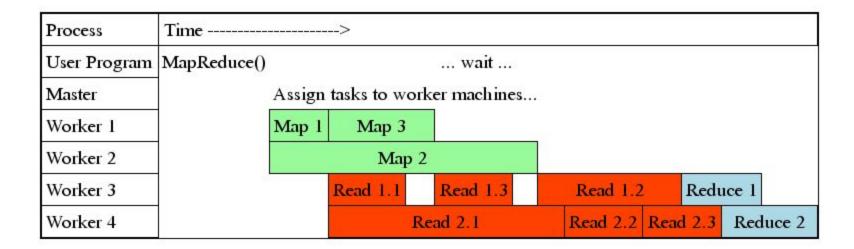
MapReduce task is aborted and client is notified

How many Map and Reduce jobs?

- M map tasks, R reduce tasks
- Rule of a thumb:
 - Make M much larger than the number of nodes in the cluster
 - One DFS chunk per map is common
 - Improves dynamic load balancing and speeds up recovery from worker failures
- Usually R is smaller than M
 - Because output is spread across R files

Task Granularity & Pipelining

- Fine granularity tasks: map tasks >> machines
 - Minimizes time for fault recovery
 - Can do pipeline shuffling with map execution
 - Better dynamic load balancing



Refinement: Backup Tasks

Problem

- Slow workers significantly lengthen the job completion time:
 - Other jobs on the machine
 - Bad disks
 - Weird things

Solution

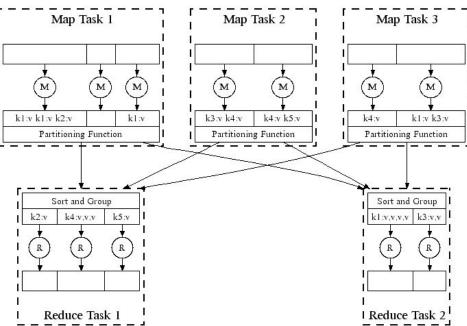
- Near end of phase, spawn backup copies of tasks
 - Whichever one finishes first "wins"

Effect

Dramatically shortens job completion time

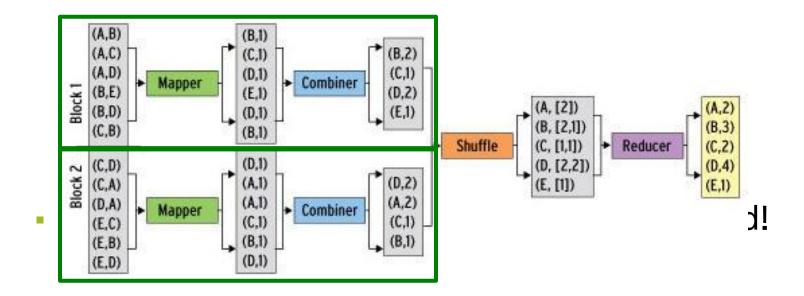
Refinement: Combiners

- Often a Map task will produce many pairs of the form (k,v_1) , (k,v_2) , ... for the same key k
 - E.g., popular words in the word count example
- Can save network time by pre-aggregating values in the mapper:
 - combine(k, list(v₁)) → v₂
 - Combiner is usually same as the reduce function



Refinement: Combiners

- Back to our word counting example:
 - Combiner combines the values of all keys of a single mapper (single machine):



Refinement: Combiners

- Combiner trick works only if reduce function is commutative and associative
- Sum (o)
 - a+b=b+a, (a+b)+c=a+(b+c)
- Average (△)
 - avg(a,b) = avg(b,a), avg(avg(a,b), c) != avg(a, avg(b,c))
 avg(b,c))
 - Combination of sum and count
- Median (x)

Refinement: Partition Function

- Want to control how keys get partitioned
 - The set of keys that go to a single reduce worker
- System uses a default partition function:
 - hash(key) mod R
- Sometimes useful to override the hash function:
 - E.g., hash(hostname(URL)) mod R ensures URLs from a host end up in the same output file

Implementations

- Google MapReduce
 - Uses Google File System (GFS) for stable storage
 - Not available outside Google
- Hadoop
 - Open-source implementation in Java
 - Uses HDFS for stable storage
 - Download: http://hadoop.apache.org/

Questions?