Root directory



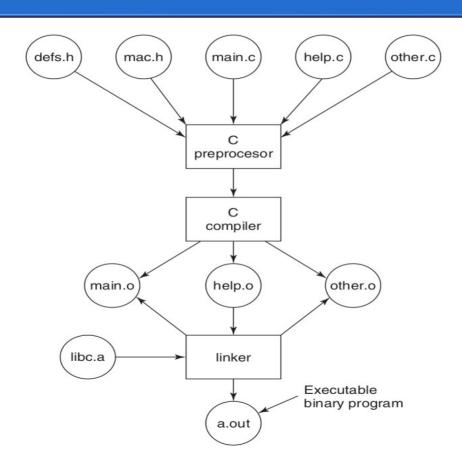
Directories in root (1)

- /bin Binaries.
- /boot Files required for booting.
- /dev Device files.
- <u>letc</u> <u>Et cetera</u>. The name is inherited from the earliest Unixes, which is when it became the spot to put config-files.
- /home Where home directories are kept.
- /lib Where code libraries are kept.
- <u>/media</u> A more modern directory, but where removable media gets mounted.
- /mnt Where temporary file-systems are mounted.
- <u>lopt</u> Where optional add-on software is installed. This is discrete from /usr/local/ for reasons I'll get to later.

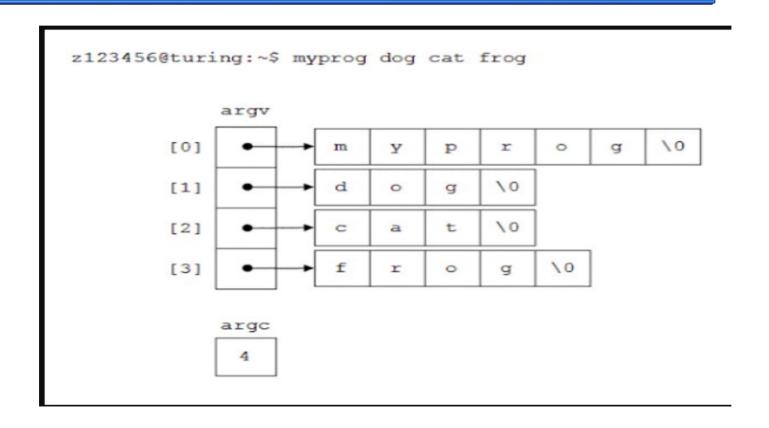
Directories in root (2)

- <u>/run</u> Where runtime variable data is kept.
- <u>/sbin</u> Where super-binaries are stored. These usually only work with root.
- <u>/srv</u> Stands for "serve". This directory is intended for static files that are served out.
 /srv/http would be for static websites, /srv/ftp for an FTP server.
- <u>/tmp</u> Where temporary files may be stored.
- <u>/usr</u> Another directory inherited from the Unixes of old, it stands for "UNIX System
 Resources". It does not stand for "user" (see the <u>Debian Wiki</u>). This directory should be
 sharable between hosts, and can be NFS mounted to multiple hosts safely. It can be mounted
 read-only safely.
- <u>/var</u> Another directory inherited from the Unixes of old, it stands for "variable". This is where
 system data that varies may be stored. Such things as spool and cache directories may be
 located here. If a program needs to write to the local file-system and isn't serving that data to
 someone directly, it'll go here.

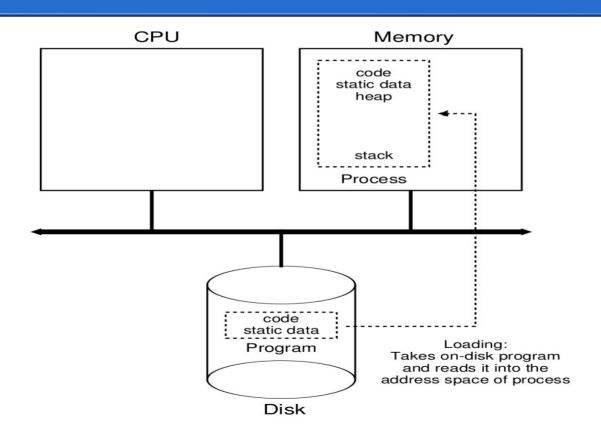
C-preprocess, compile and link



C main arguments



Unix process loading

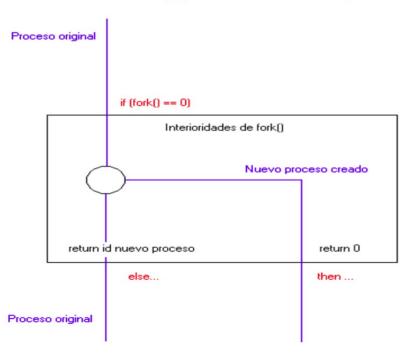


Direct execution no limits

OS	Program
Create entry for process list	
Allocate memory for program	
Load program into memory	
Set up stack with argc/argv	
Clear registers	
Execute call main()	
	Run main()
	Execute return from main
Free memory of process	
Remove from process list	

Fork process creation

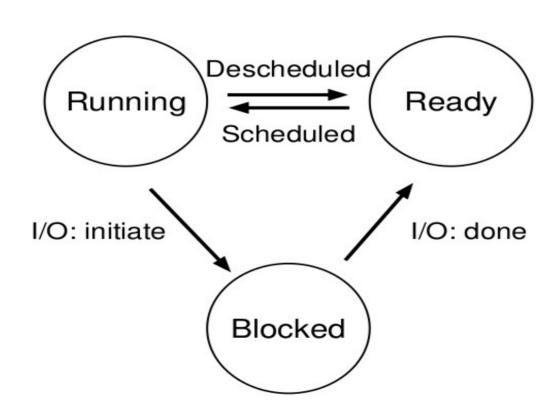
Interioridades de fork() y creación de un nuevo proceso



Fork program layout

```
switch (fork())
    case -1:
        /* Código de error */
        break;
    case 0:
        /* Código del proceso hijo */
        break;
    default:
        . . .
        /* Código del proceso original */
```

Process states



Context switch

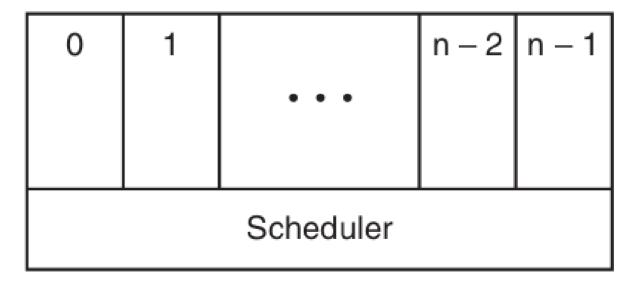
OS @ boot (kernel mode)	Hardware	
initialize trap table	remember addresses of syscall handler	
start interrupt timer	timer handler	
	start timer interrupt CPU in X ms	
OS @ run (kernel mode)	Hardware	Program (user mode)
		Process A
Handle the trap	timer interrupt save regs(A) to k-stack(A) move to kernel mode jump to trap handler	
Call switch() routine save regs(A) to proc-struct(A) restore regs(B) from proc-struct(B) switch to k-stack(B) return-from-trap(into B)	restore regs(B) from k-stack(B)	
	move to user mode jump to B's PC	
		Process B

Kernel interrupt service

- 1. Hardware stacks program counter, etc.
- 2. Hardware loads new program counter from interrupt vector.
- 3. Assembly-language procedure saves registers.
- 4. Assembly-language procedure sets up new stack.
- 5. C interrupt service runs (typically reads and buffers input).
- 6. Scheduler decides which process is to run next.
- 7. C procedure returns to the assembly code.
- 8. Assembly-language procedure starts up new current process.

Scheduler processes





Process table fields

Process management

Registers

Program counter

Program status word

Stack pointer

Process state

Priority

Scheduling parameters

Process ID

Parent process

Process group

Signals

Time when process started

CPU time used

Children's CPU time

Time of next alarm

Memory management

Pointer to text segment info Pointer to data segment info

Pointer to stack segment info

File management

Root directory

Working directory

File descriptors

User ID

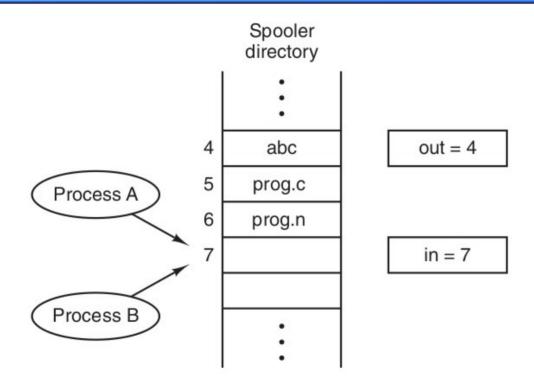
Group ID

Threads items

Per-process items	Per-thread items
Address space	Program counter
Global variables	Registers
Open files	Stack
Child processes	State
Pending alarms	
Signals and signal handlers	
Accounting information	

Figure 2-12. The first column lists some items shared by all threads in a process. The second one lists some items private to each thread.

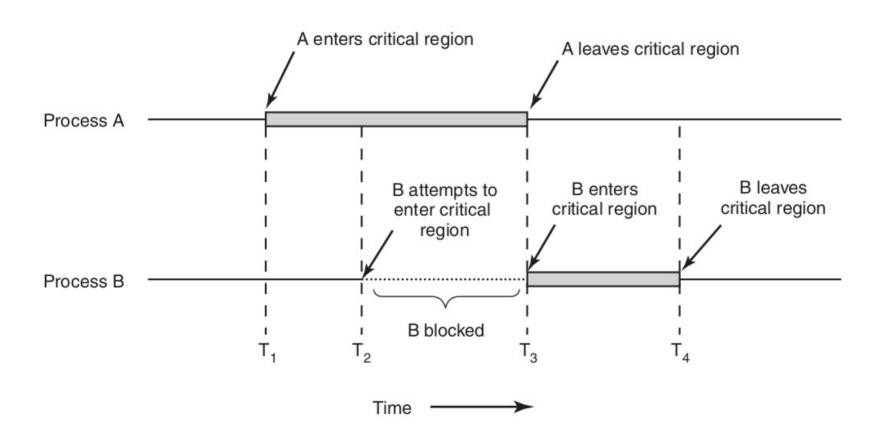
Races



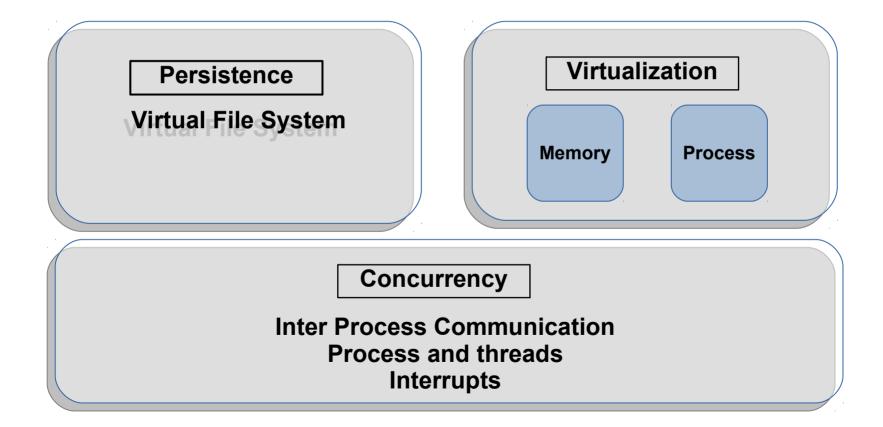
21. Two processes want to access shared memory at the same time.

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Mutual exclusion (critical region)



Operating Systems abstractions



Linux layered structure

