Sum

- Problem: Sum a given set of numbers
- Solution: Sum the numbers...

Gates

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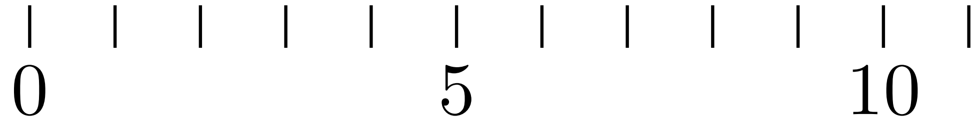

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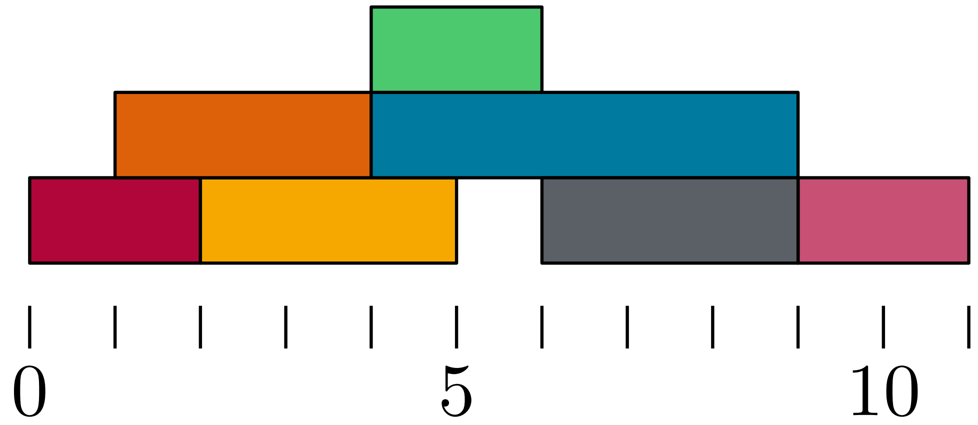

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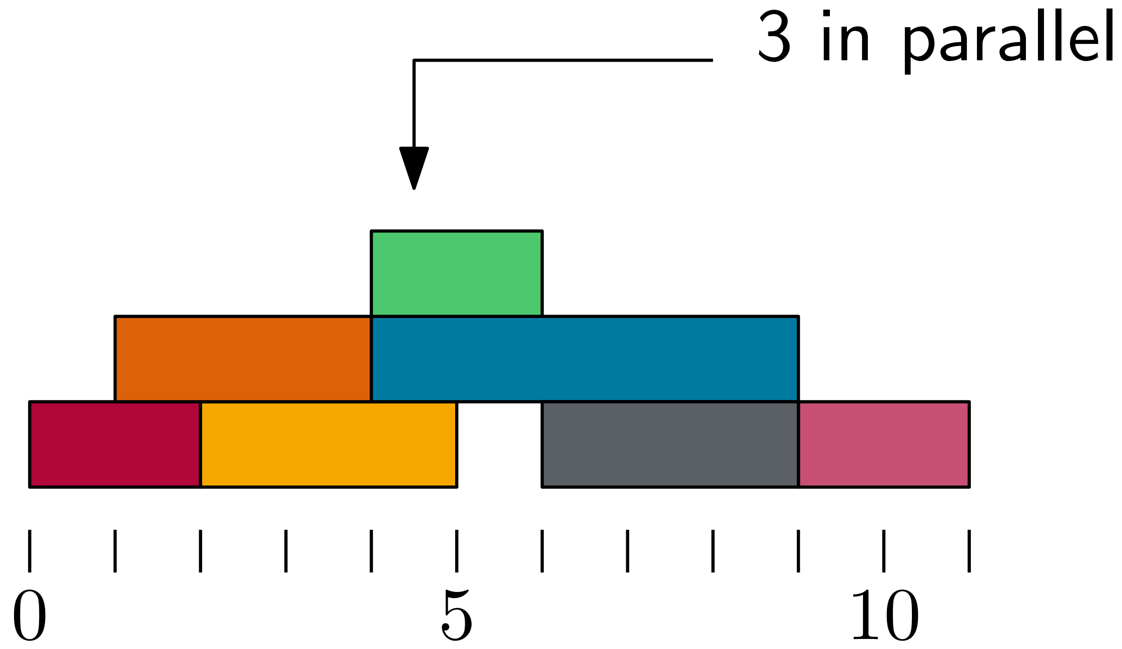

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- dominated by sorting: $\mathcal{O}(n \log n)$

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- small n allows for n^2 or even n^3 brute-force

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- the minimum time from your friends is the sum of the k easiest papers minus a

Never trust a greedy algorithm!

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- Testcase `anti-greedy-3.in`:

```
3
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- Runtime: $\mathcal{O}(n \log n)$

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- For every object keep the interval where it can be at time t and intersect all intervals. If the resulting interval is nonempty, we can choose a meeting point.
- For general k, we have to partition the objects into $k + 1$ groups and place a portal in k of them.

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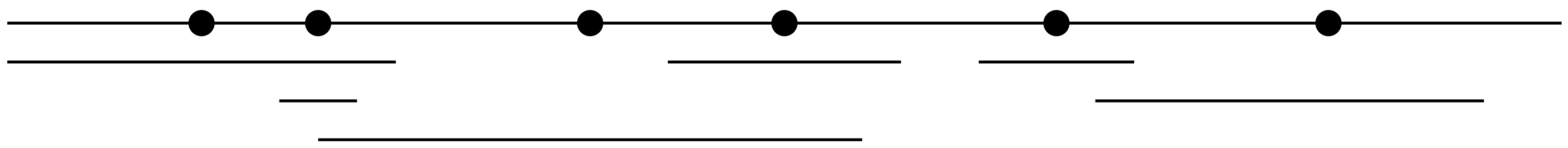

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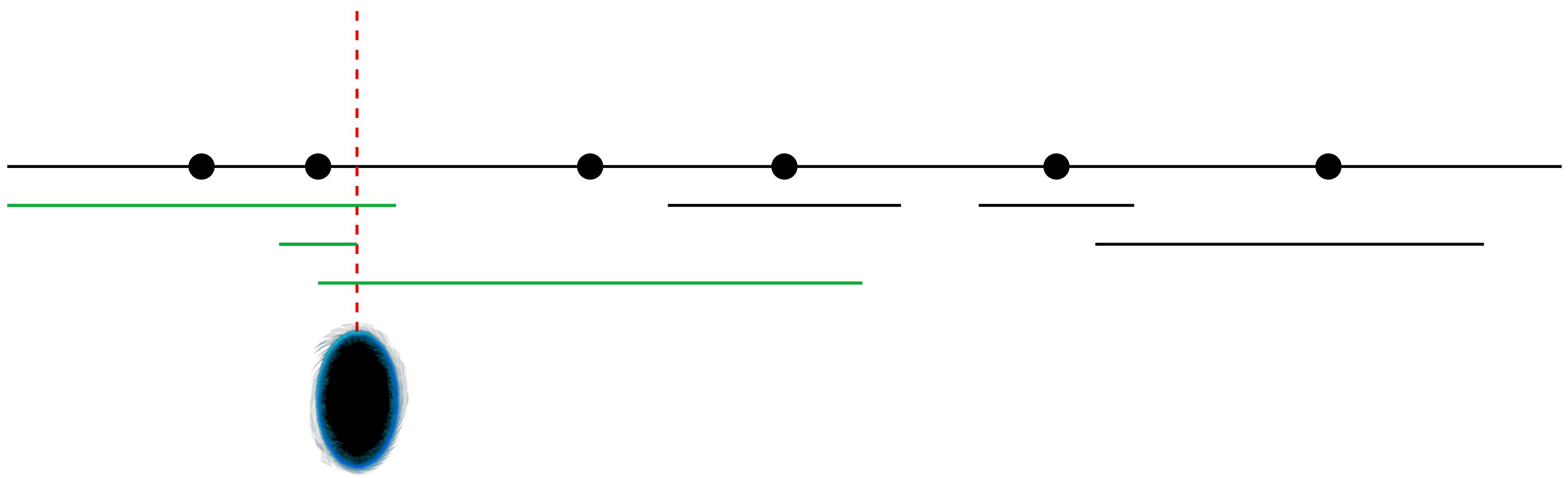

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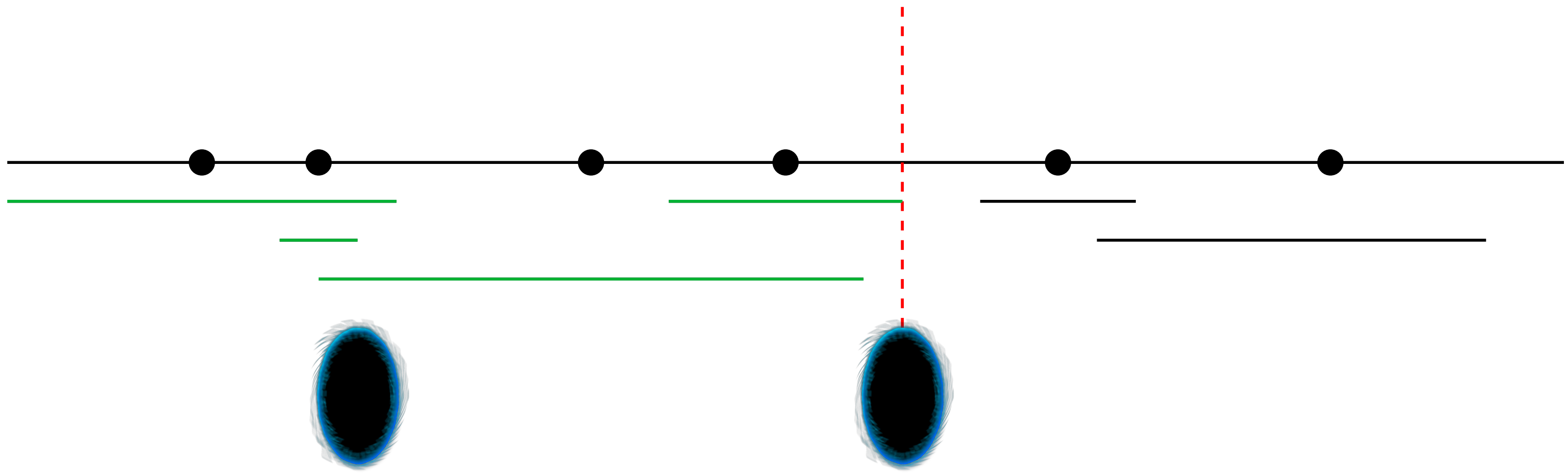

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