HW#6 (CSC390); Turn in your HW on the Blackboard by 03/13/2018 by 11:30PM

#Q1. Consider the following two C procedures (swap and sort) and their corresponding MIPS assembly codes as shown in the figures 1 and 2, respectively. Using these two C procedures (swap and sort) write a MIPS assembly program that will sort the following array elements in the **ascending order**. What changes would you make to sort the array elements in **descending order**?

Array= [100 50 75 -1 -50 500 20 40 40 17 19 23 5 7 -20]

```
void swap(int v[], int k)
{
  int temp;
  temp = v[k];
  v[k] = v[k+1];
  v[k+1] = temp;
}
```

Procedure body \$t1, \$a1, 2 \$t1, \$a0, \$t1 swap: sll # reg \$t1 = k * 4# reg \$t1 = v + (k * 4)add # reg \$t1 has the address of v[k] \$t0,0(\$t1) # reg \$t0 (temp) = v[k]# reg \$t2 = v[k + 1] \$t2, 4(\$t1) l w #refers to next element of v SW \$t2.0(\$t1) # v[k] = reg \$t2\$t0, 4(\$t1) # v[k+1] = reg \$t0 (temp)SW

Procedure return					
jr	\$ra	# return to calling routine			

Figure 1 MIPS assembly code for swap procedure

```
void sort (int v[], int n)
{
  int i, j;
  for (i = 0; i < n; i += 1) {
    for (j = i - 1; j >= 0 && v[j] > v[j + 1]; j -= 1) {
        swap(v,j);
    }
  }
}
```

Saving registers							
sort:	S W S W	<pre>\$ra, 16(\$sp)# save \$ra on stack \$s3,12(\$sp) # save \$s3 on stack</pre>					
	SW SW SW	\$s2, 8(\$sp)# save \$s2 on stack \$s1, 4(\$sp)# save \$s1 on stack \$s0, 0(\$sp)# save \$s0 on stack					

		Procedure body
Move parameters	mov mov	
Outer loop	mov for1tst:	slt\$t0, \$s0, \$s3 $\#$ reg \$t0 = 0 if \$s0 \S \$s3 (i \S n)
Inner loop	add for2tst: bne sll add lw lw slt beq	slti\$t0, \$s1,0 # reg \$t0 = 1 if \$s1 < 0 (j < 0) \$t0, \$zero, exit2# go to exit2 if \$s1 < 0 (j < 0) \$t1, \$s1, 2# reg \$t1 = j * 4 \$t2, \$s2, \$t1# reg \$t2 = v + (j * 4) \$t3, 0 (\$t2)# reg \$t3 = v[j] \$t4, 4 (\$t2)# reg \$t4 = v[j + 1]
Pass parameters and call	mov mov jal	\$a0, \$s2 #1st parameter of swap is v (old \$a0)
Inner loop	add j	\$s1, \$s1, -1# j -= 1 for2tst # jump to test of inner loop
Outer loop	exit2: add	

Restoring registers							
exit1:	1w	\$s0,0(\$sp)	# restore \$sO from stack				
	1 w	\$s1,4(\$sp)# restore\$s1 from stack					
	1 w	\$s2,8(\$sp)# restore\$s2 from stack					
	1 w	\$s3,12(\$sp) # restore \$s3 from stack					
	1 w	<pre>\$ra,16(\$sp) # restore \$ra from stack</pre>					
	addi	\$sp,\$sp, 20	# restore stack pointer				

		Procedure return
jr	\$ra	# return to calling routine

Figure 2 MIPS assembly version of sort procedure

Results (ascending order):

□ Data Segment								
Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	-50	-20	-1	5	7	17	19	20 📤
0x10010020	23	40	40	50	75	100	500	0

Results (descending order):

Data Segment ਸੰਯੋ								
Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	500	100	75	50	40	40	23	20 📤
0x10010020	19	17	7	5	-1	-20	-50	0

Q2. Given the two object files of procedure A and procedure B, show updated address of the first few instructions of the completed executable file. Ref: pages 127-128 of your text book. Note that the default starting address of Text segment = $0x00400000_{hex}$, Data Segment = $0x100000000_{hex}$ and $$gp = 0x10008000_{hex}$

Object File Header			
	Name	Procedure A	
	Text Size	120 _{hex}	
	Data Size	30 _{hex}	
Text Segment	Address	Instruction	
	0	lw \$a0, O1 (\$gp)	
	4	lw \$a1, O2 (\$gp)	
	8	jal O	
Data Segment	01	(X1)	
	02	(X2)	
Relocation Table	Label	Address	Dependency
	0	lw	X1
	4	lw	X2
	8	jal	В
Symbol Table	Label	Address	
	X1		
	X2		
	В		
Object File Header			
	Name	Procedure B	
	Text Size	250 _{hex}	
	Data Size	25 _{hex}	
Text Segment	Address	Instruction	
	0	lw \$a0, O1 (\$gp)	
	8	jal O	
	4	sw \$a1, O2 (\$gp)	
Data Segment	01	(Y1)	
	02	(Y2)	
Relocation Table	Label	Address	Dependency
	0	lw	Y1
	4	jal	A
	8	SW	Y2
Symbol Table	Label	Address	
	Y1		
	Α		
	Y2		

Table 1 Object files of Procedure A and Procedure B