

COSC 2021: Computer Organization

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Handout # 6

Unsigned integer and floating point instructions

Topics:

1. Unsigned instructions
2. Floating Point: IEEE 754 single and double precision formats
3. Floating Point Registers and Instructions

Patterson: Sections 3.1 – 3.2, 3.5

Unsigned and Signed Arithmetic

MIPS has a separate format for unsigned and signed integers

1. Unsigned integers

- are saved as 32-bit words

Example: Smallest unsigned integer is $00000000_{\text{hex}} = 0_{\text{ten}}$

Largest unsigned integer is $\text{ffffffff}_{\text{hex}} = 4,294,967,295_{\text{ten}}$

2. Signed integers

- are saved as 32-bit words in 2's complement with the MSB reserved for sign
- If MSB = 1, then the number is negative
- If MSB = 0, then the number is positive

Example:

Smallest signed integer: $1000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000_{\text{two}}$
 $= -(2^{31})_{10} = -2,147,483,648_{10}$

Largest signed integer: $0111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111_{\text{two}}$
 $= (2^{31} - 1)_{10} = 2,147,483,647_{10}$

MIPS Commands for Unsigned Numbers

Category	Instruction	Example	Meaning	Comments
Arithmetic	add(u)	add \$s1,\$s2,\$s3	\$s1 ← \$s2+\$s3	Most arithmetic
	Subtract(u)	sub \$s1,\$s2,\$s3	\$s1 ← \$s2-\$s3	instruction have
	addi(u)	add \$s1,\$s2,100	\$s1 ← \$s2+100	unsigned format
Data Transfer	load word	lw \$s1,100(\$s2)	\$s1 ← Mem[\$s2+100]	
	store word	sw \$s1,100(\$s2)	Mem[\$s2+100] ← \$s1	
	load byte unsigned	lbu \$s1,100(\$s2)	\$s1 ← Mem[\$s2+100]	Unsigned
	store byte	sb \$s1,100(\$s2)	Mem[\$s2+100] ← \$s1	1 byte only
Conditional branch	branch on equal	beq \$s1,\$s2,L	if (\$s1==\$s2) go to L	
	branch not equal	bne \$s1,\$s2,L	if (\$s1!=\$s2) go to L	
	set on less than	slt \$s1,\$s2,\$s3	if (\$s2<\$s3) \$s1 = 1 else \$s1 = 0	
	set on less than immediate	slti \$s1,\$s2,10	if (\$s2<10) \$s1 = 1 else \$s1 = 0	
	set less than unsigned	sltu \$s1,\$s2,\$s3	- same as slt -	Unsigned
	slt unsign immediate	sltui \$s1,\$s2,100	- same as slti -	Unsigned
Unconditional jump	jump	j 2500	go to (4 x 2500)	
	jump register	jr \$ra	go to \$t1	
	jump and link	jal fact	go to fact; set \$ra = PC + 4	

Addition and Subtraction

In MIPS, addition and subtraction for signed numbers use 2's complement arithmetic

Example 1: Add 10_{ten} and 15_{ten}

Step 1: Represent the operands in 2's complement

$10_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$

$15_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1111_{\text{two}}$

Step 2: Perform bit by bit addition using table 1.

$10_{\text{ten}} + 15_{\text{ten}}$
 $= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1001_{\text{two}}$
 $= 25_{\text{ten}}$

Example 2: Subtract 15_{ten} from 10_{ten}

The problem is reduced to $(10_{\text{ten}} + (-15_{\text{ten}}))$

$10_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$

$-15_{\text{ten}} = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 0001_{\text{two}}$

$10_{\text{ten}} - 15_{\text{ten}}$
 $= 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1011_{\text{two}}$
 $= -5_{\text{ten}}$

bit 1	bit 2	Prev. Carry	Sum	Next Carry
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Table 1: Truth Table for addition

Overflow (1)

Recall that:

Smallest signed integer: $1000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000_{\text{two}}$
 $= -(2^{31})_{\text{ten}} = -2,147,483,648_{\text{ten}}$

Largest signed integer: $0111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111_{\text{two}}$
 $= (2^{31} - 1)_{\text{ten}} = 2,147,483,647_{\text{ten}}$

What happens if the result of an operation is more than the largest signed integer or less than the smallest signed integer?

Example: Add $2,147,483,640_{\text{ten}}$ and 28_{ten}

$$\begin{array}{rcl} 28_{\text{ten}} & = & 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1100_{\text{two}} \\ 2,147,483,640_{\text{ten}} & = & 0111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1000_{\text{two}} \\ \hline 28_{\text{ten}} + 2,147,483,640_{\text{ten}} & = & 1000\ 0000\ 0000\ 0000\ 0000\ 0001\ 0100_{\text{two}} \\ & = & -2,147,483,628_{\text{ten}} \end{array}$$

Overflow caused the value to be perceived as a negative integer

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Overflow (2)

When can overflow occur?

Operation	Operand A	Operand B	Result indicating overflow
A + B	A \geq 0	B \geq 0	< 0
A + B	A < 0	B < 0	\geq 0
A - B	A \geq 0	B < 0	< 0
A - B	A < 0	B \geq 0	\geq 0

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Floating Point: Single Precision

- In MIPS, decimal numbers are represented with the **IEEE 754 binary representation** that uses the **normalized** standard scientific binary notation defined as

$$(-1)^S \times (1 + \text{fraction})_{\text{two}} \times 2^{\text{exponent} - \text{bias}}$$
- A number in **normalized scientific notation** has a mantissa that has no leading 0's and must be of the form $(1 + \text{fraction})$. For example, the binary representations 2.0×2^{-5} , 0.5×2^{-3} , 4.0×2^{-6} , and 1.0×2^{-4} are all equivalent but only 1.0×2^{-4} is the normalized scientific binary notation.
- MIPS allows for two floating point representations: Single precision and double precision.
- Single precision** has a bias of 127 while double precision has a bias of 1023.
- In single precision, the floating point representation is 32 bit long and has the following form

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
S	exponent								fraction																						
	(8 bits)								(23 bits)																						

where S represents the sign bit, which is 1 for negative numbers and 0 for positive numbers.

Activity 2:

Represent -0.75_{ten} in single precision of IEEE 754 binary representation.

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Floating Point: Double Precision

- In **double precision**, the value of bias in

$$(-1)^S \times (1 + \text{fraction})_{\text{two}} \times 2^{\text{exponent} - \text{bias}}$$
is 1023.
- In single precision, the floating point representation is 64 bit long and has the following form

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
S	exponent											fraction																			
	(11 bits)											(Total of 52 bits)																			
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
fraction (continued)																															

Activity 3:

Represent -0.75_{ten} in double precision of IEEE 754 binary representation.

Activity 4:

Show that the largest magnitude that can be represented using single precision is $\pm 2.0_{\text{ten}} \times 10^{38}$, while the smallest fraction that can be represented is $\pm 2.0_{\text{ten}} \times 10^{-38}$.

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Floating Point Registers

Name	Example	Comments
32 floating point registers each is 32 bits long	<code>\$f0, \$f1, \$f2, \$f3, \$f4, ..., \$f31</code>	MIPS floating point registers are used in pairs for double precision numbers
Memory w/ 2^{30} words	Memory[0], Memory[4], ... Memory[4294967292]	Memory is accessed one floating point (single or double precision) at a time

The following is the established register usage convention for the floating point registers:

<code>\$f0, \$f1, \$f2, \$f3:</code>	Function-returned values
<code>\$f4, \$f5, ..., \$f11:</code>	Temporary values
<code>\$f12, \$f13, \$f14, \$f15:</code>	Arguments passed into a function
<code>\$f16, \$f17, \$f18, \$f19:</code>	More Temporary values
<code>\$f20, \$f21, ..., \$f31:</code>	Saved values

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Example

```
# calculate area of a circle
.data
Ans:    .asciiz    "The area of the circle is: "
Ans_add: .word     Ans           # Pointer to String (Ans)
Pi:     .double   3.1415926535897924
Rad:    .double   12.345678901234567
Rad_add: .word     Rad           # Pointer to float (Rad)
.text
main:   lw $a0, Ans_add($0)      # load address of Ans into $a0
        addi $v0, $0, 4          # Sys Call 4 (Print String)
        syscall

#----- # load float (Assembler Instruction)
        la $s0, Pi              # load address of Pi into $s0
        ldc1 $f2, 0($s0)        # $f2 = Pi
#----- # load float (MIPS Instruction)
        lw $s0, Rad_add($0)     # load address of Rad into $s0
        ldc1 $f4, 0($s0)        # $f4 = Rad
        mul.d $f12, $f4, $f4
        mul.d $f12, $f12, $f2
        addi $v0, $0, 3         # Sys Call 3 (Print Double)
        syscall
exit:   jr $ra
```

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Floating Point Instructions

Category	Instruction	Example	Meaning	Comments
Arithmetic	FP add single	<code>add.s \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 + \$f6$	Single Prec.
	FP subtract single	<code>sub.s \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 - \$f6$	Single Prec.
	FP multiply single	<code>mul.s \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 \times \$f6$	Single Prec.
	FP divide single	<code>div.s \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 / \$f6$	Single Prec.
	FP add double	<code>add.d \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 + \$f6$	Double Prec.
	FP subtract double	<code>sub.d \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 - \$f6$	Double Prec.
	FP multiply double	<code>mul.d \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 \times \$f6$	Double Prec.
	FP divide double	<code>div.d \$f2, \$f4, \$f6</code>	$\$f2 \leftarrow \$f4 / \$f6$	Double Prec.
Data Transfer	load word copr.1	<code>lwc1 \$f2, 100(\$s2)</code>	$\$f2 \leftarrow \text{Mem}[\$s2 + 100]$	Single Prec.
	store word copr.1	<code>swc1 \$f2, 100(\$s2)</code>	$\text{Mem}[\$s2 + 100] \leftarrow \$f2$	Single Prec.
Conditional branch	FP compare single (eq, ne, lt, le, gt, ge)	<code>c.lt.s \$f2, \$f4</code>	$\text{if } (\$f2 < \$f4) \text{ cond} = 1, \text{ else cond} = 0$	Single Prec.
	FP compare double (eq, ne, lt, le, gt, ge)	<code>c.lt.d \$f2, \$f4</code>	$\text{if } (\$f2 < \$f4) \text{ cond} = 1, \text{ else cond} = 0$	Double Prec.
	Branch on FP true	<code>bc1t 25</code>	$\text{if cond} == 1 \text{ go to PC} + 100 + 4$	Single/Double Prec.
	Branch on FP false	<code>bc1f 25</code>	$\text{if cond} == 0 \text{ go to PC} + 100 + 4$	Single/Double Prec.

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