

University of Illinois at Urbana-Champaign
Dept. of Electrical and Computer Engineering

ECE 120: Introduction to Computing

The Design of the Lab FSM

What Problem Must be Solved?

In the lab, your task is

- to build a small FSM
- to **control a coin-operated vending machine.**

Inputs are produced by coins.

Outputs specify

- whether a coin should be accepted, and
- whether a product should be released.

The **design is extremely simple** so as to minimize the number of chips needed.

What Purpose Does the Lab Serve?

Help students to make the connection between lines and boxes on paper and wires and chips in a real system.

Help students to realize that the knowledge they have gained in ECE120 enables them to build real systems with sensors and actuators; in other words, to interface with the real world.

Let's Take a Look at the Physical System

These produce
CLK and T.



Gate is
controlled by
output A.

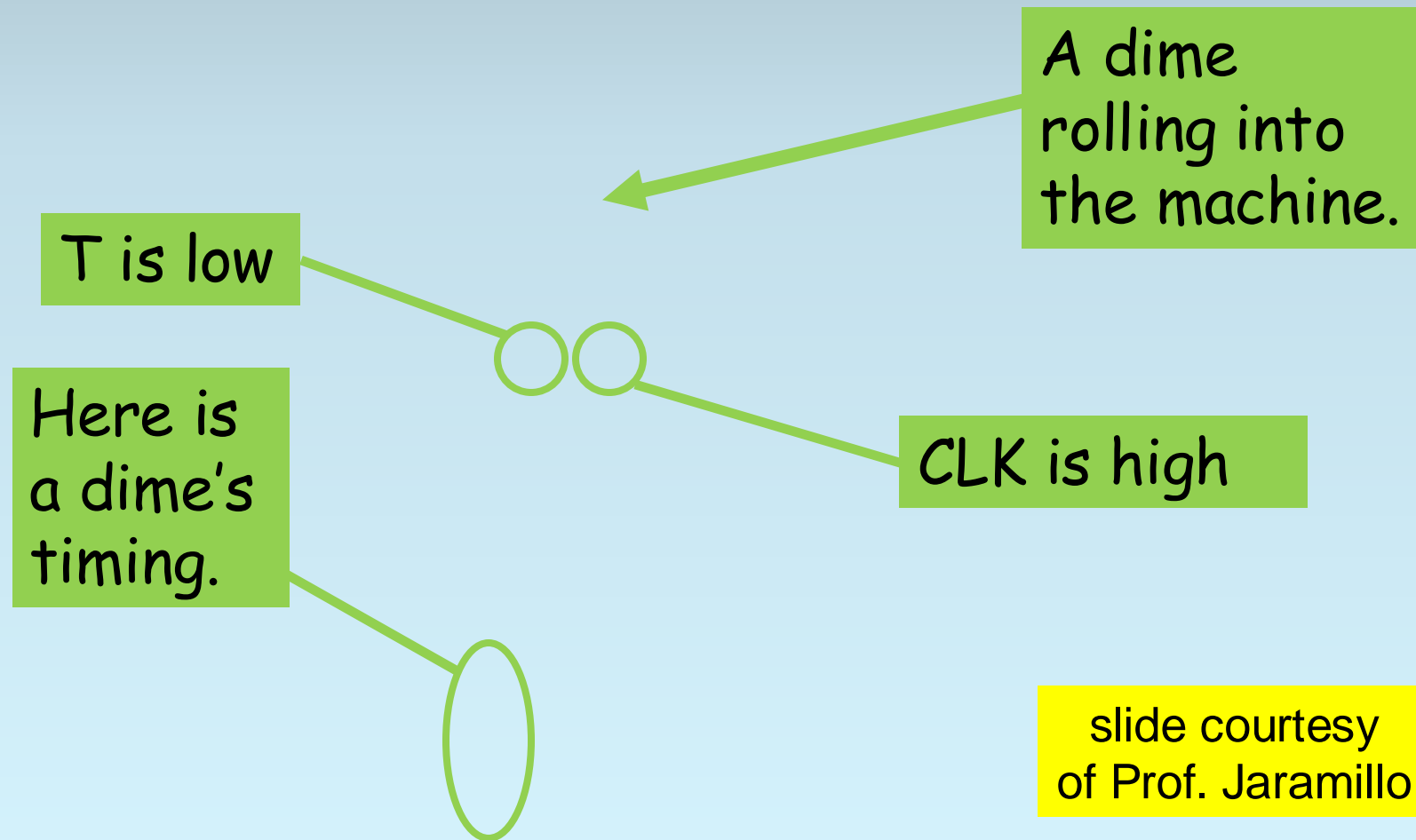
slide courtesy
of Prof. Jaramillo

A Closer Look at Sensors and LED Feedback

CLK

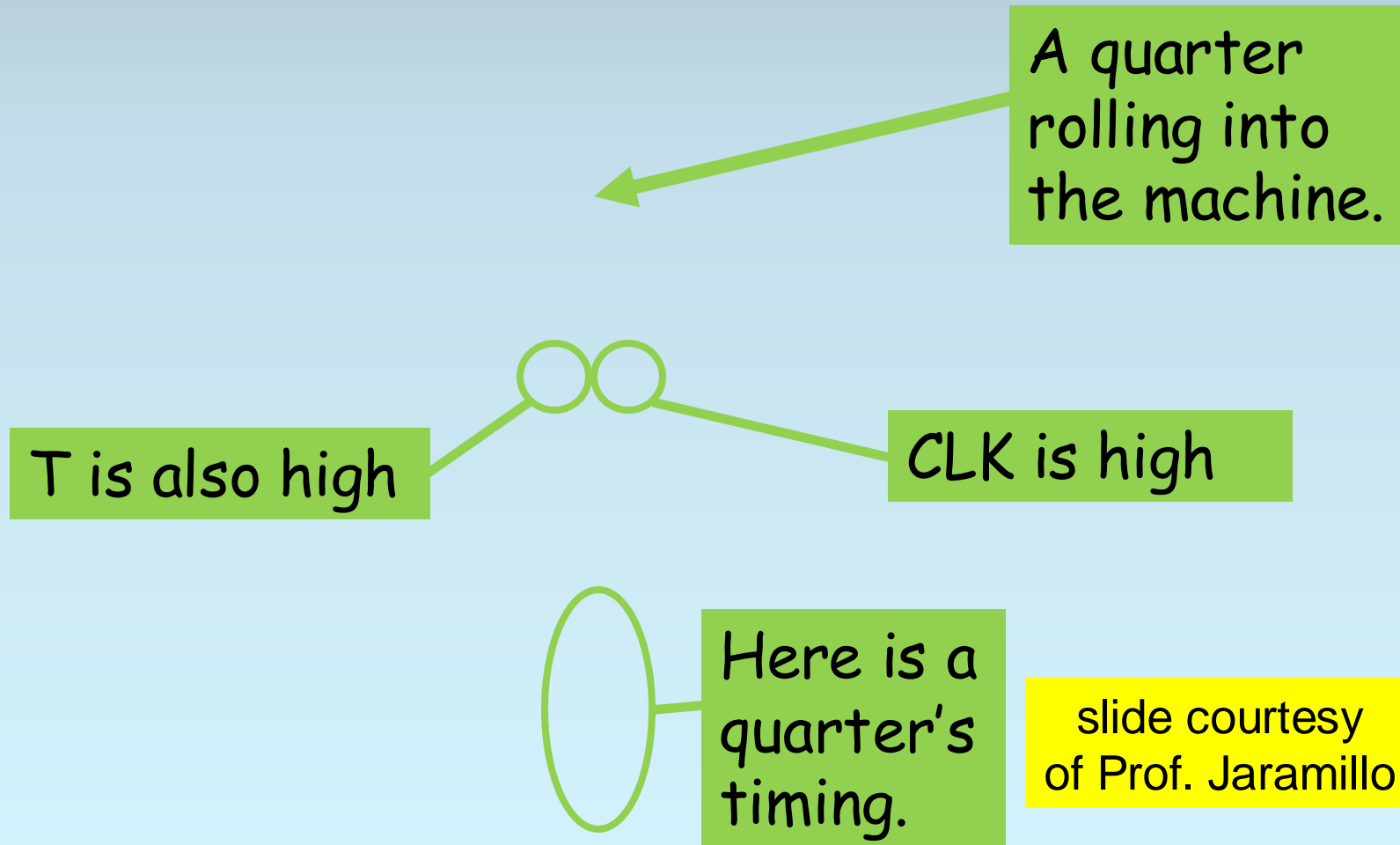
slide courtesy
of Prof. Jaramillo

What Happens When a Dime Rolls In?



slide courtesy
of Prof. Jaramillo

What Happens When a Quarter Rolls In?



The Clock Signal is Unusual

The clock signal **CLK** is

- produced by an optical sensor
- when a coin rolls in front of it.

As a result, **CLK** is

- not a square wave, and
- not even periodic!
- The high pulses are coins.
- The pulse width depends on the coin's speed.
- The cycle time is the time between coins.

CLK Signal is Sufficient for Our FSM's Needs

However, the **CLK** signal is **sufficient for our needs**.

You build with positive edge-triggered D flip-flops.

Because of the positioning of optical sensors, **T** is stable (0 for a dime, 1 for a quarter) when **CLK** rises.

Lab Machine as a Sequence Recognizer

A **sequence recognizer** looks for bit patterns in a serial input stream.

Previously, we developed a 01 sequence recognizer as an example of the difference between Mealy and Moore machines.

For the lab, we can

- treat the sequence of coins as a serial input (0 for dimes—\$0.10; 1 for quarters—\$0.25).
- then look for patterns that total \$0.35.

A Process for Developing a 01/10 Recognizer

If **T** is the serial input of coin types, we must produce product release output **P = 1** whenever we see **01** or **10**.

We use the following process to develop the sequence recognizer:

1. Create a state for each bit in each sequence.
2. Complete the transitions not used in the desired sequences.
3. Minimize by merging redundant states.

Inputs, Output, and Notation

Inputs:

T 0 for dime, **1 for quarter**

(transitions use black/**red** to denote T)

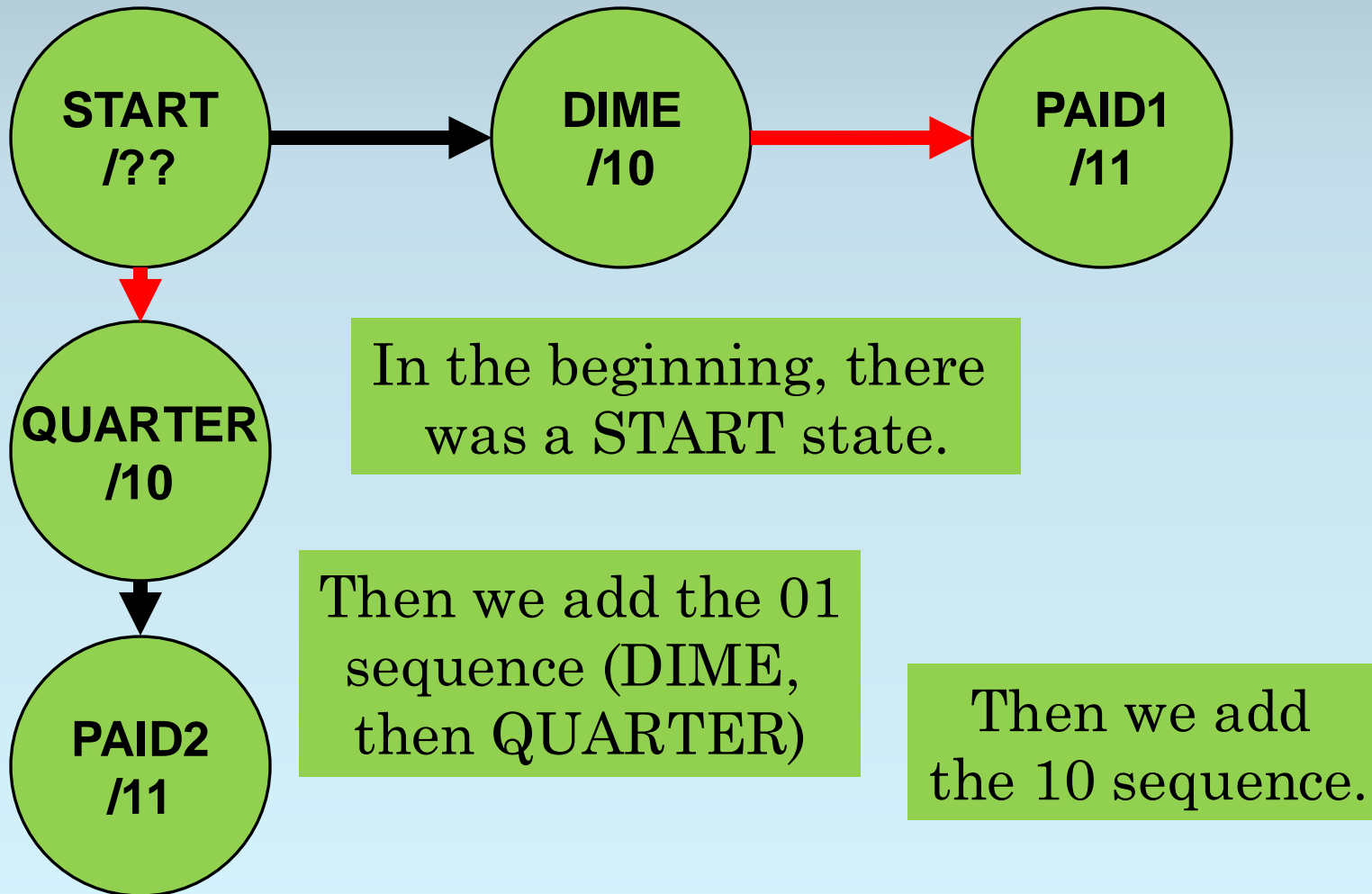
Outputs:

A 1 to accept the coin just inserted
(0 to reject it, returning to the user)

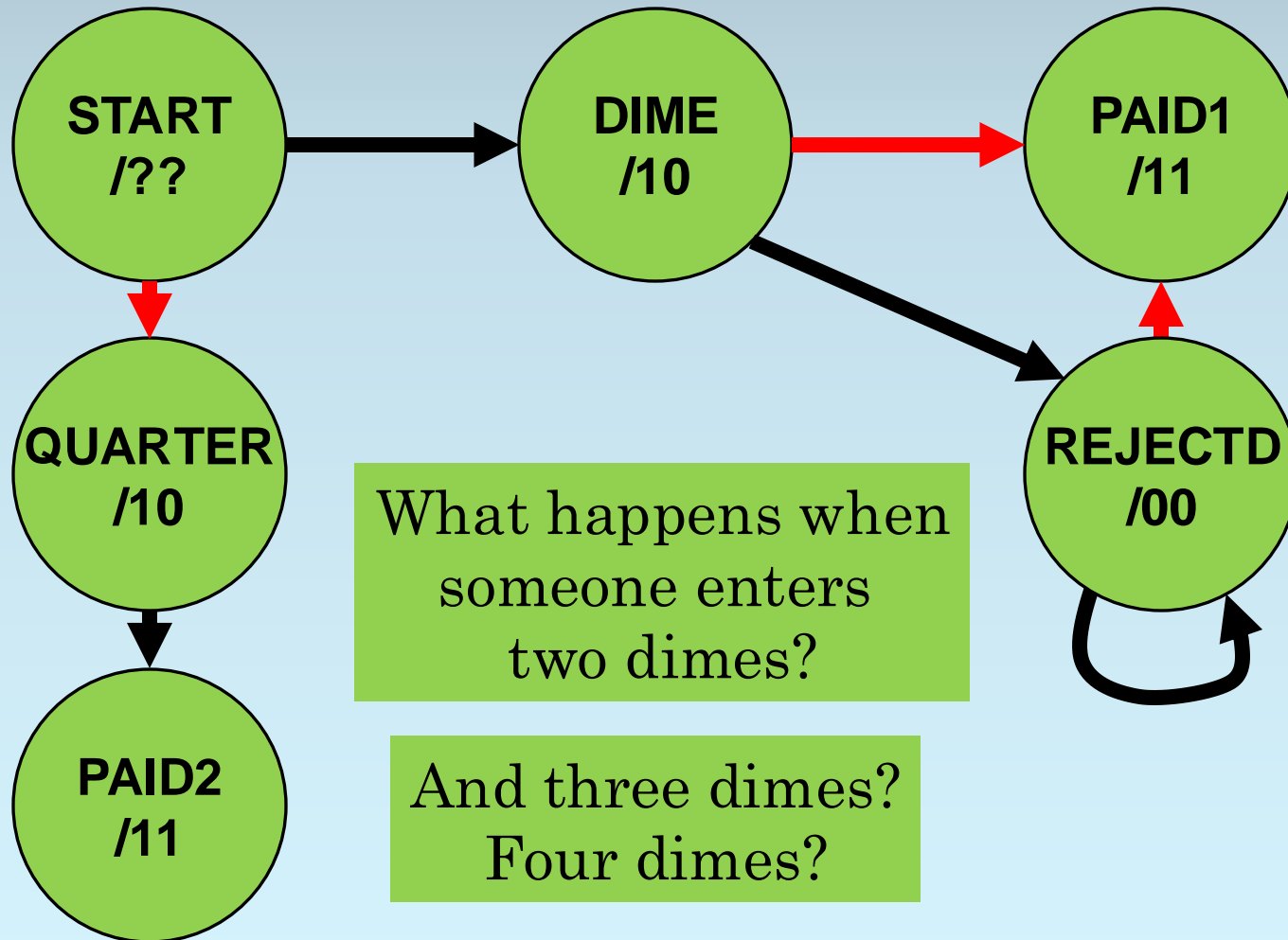
P 1 to release the product

States are marked with /AP.

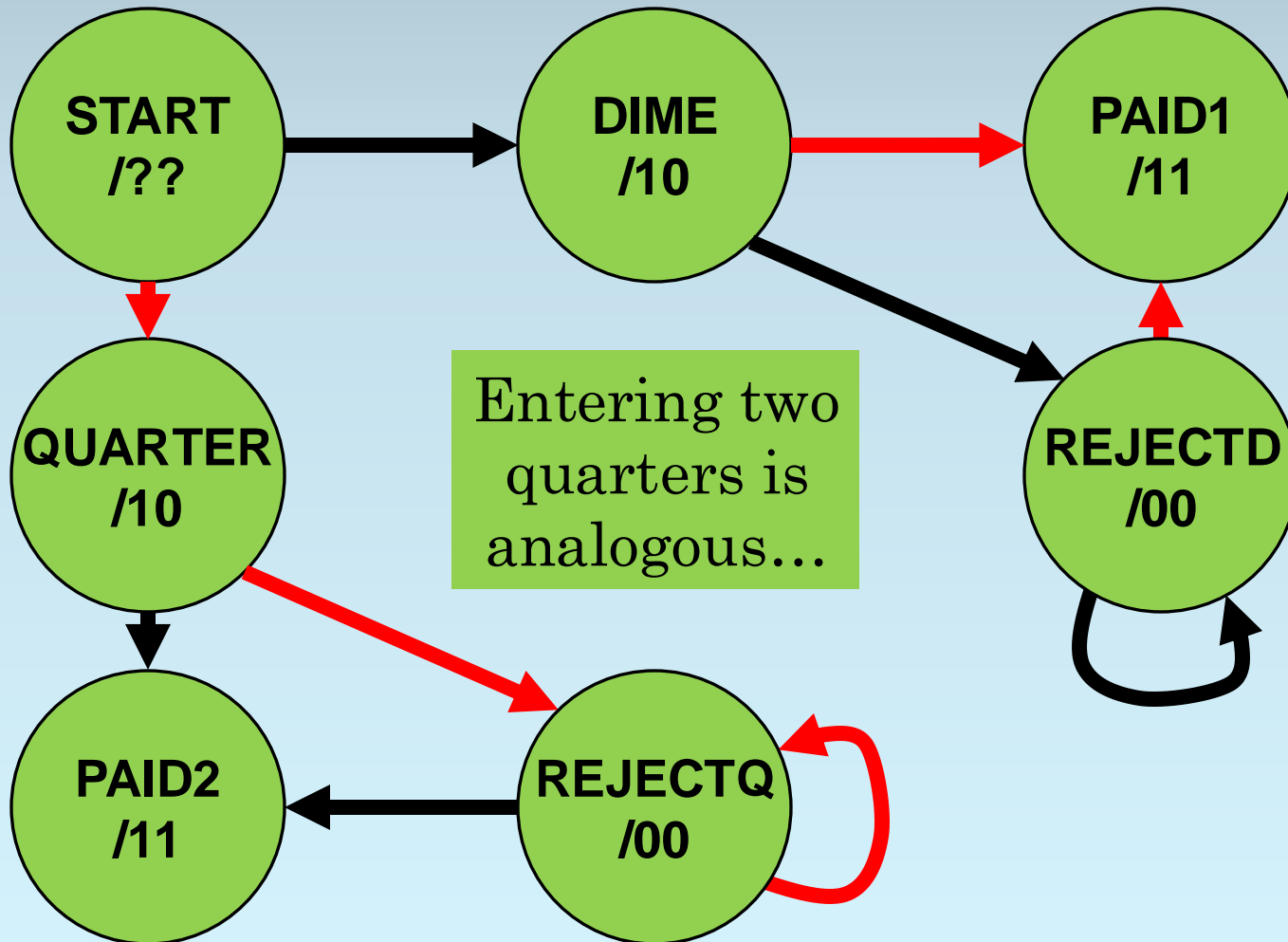
Step 1: Draw States for Both Sequences



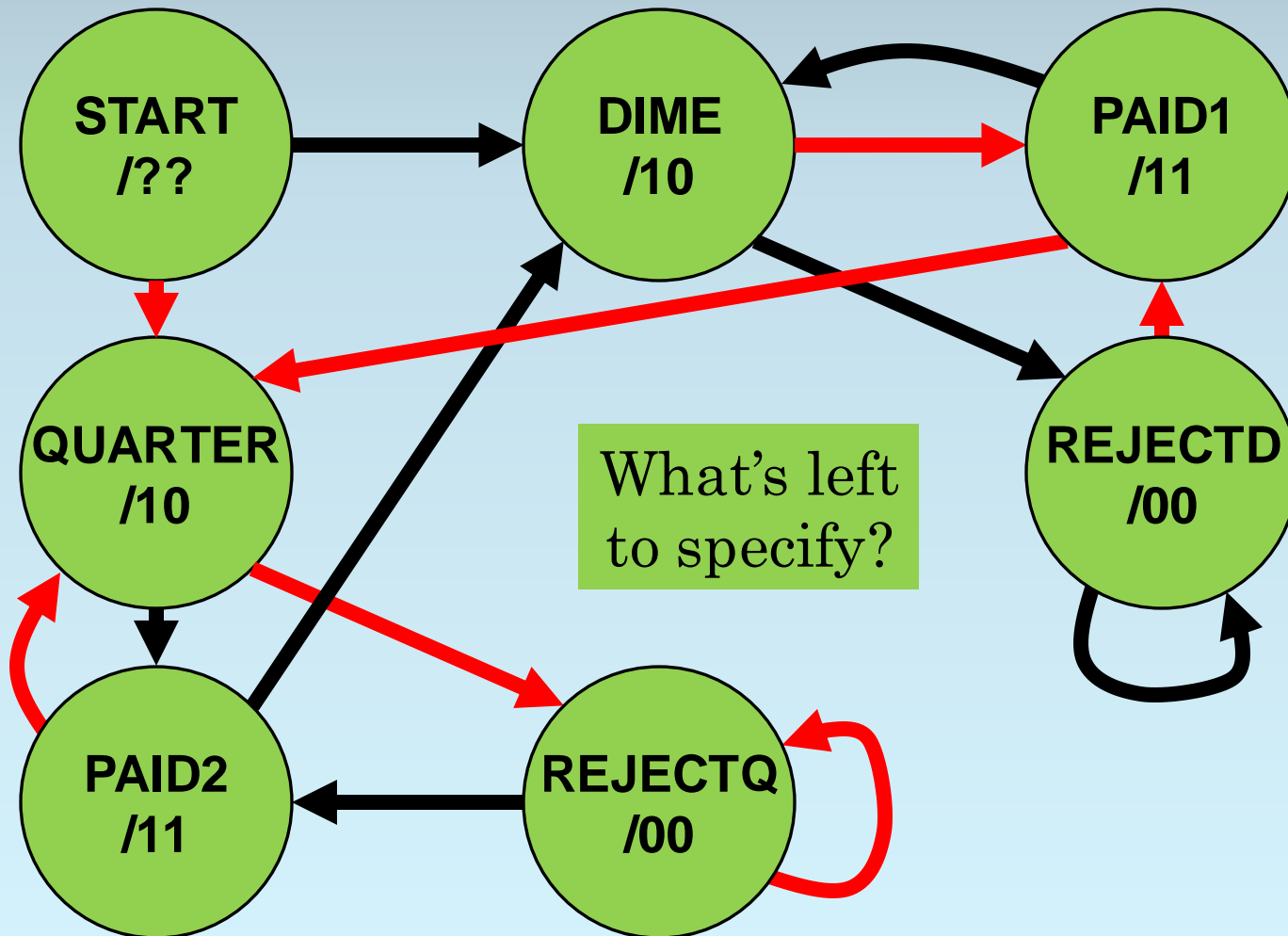
Step 2: Complete the Transitions



Step 2: Complete the Transitions



Step 2: Complete the Transitions



Step 3: Merge Redundant States

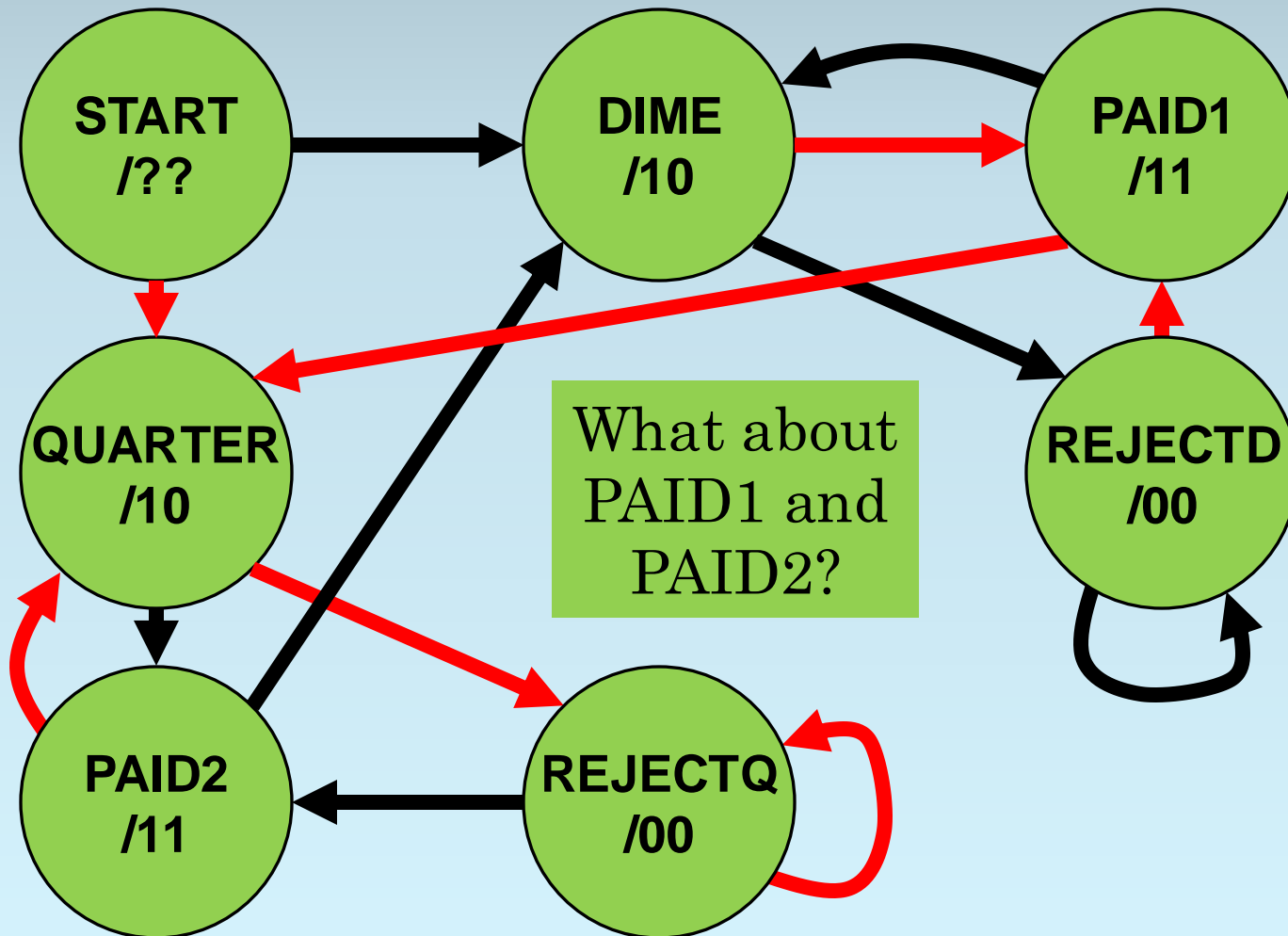
Now we can merge redundant states.

To merge states, it suffices to

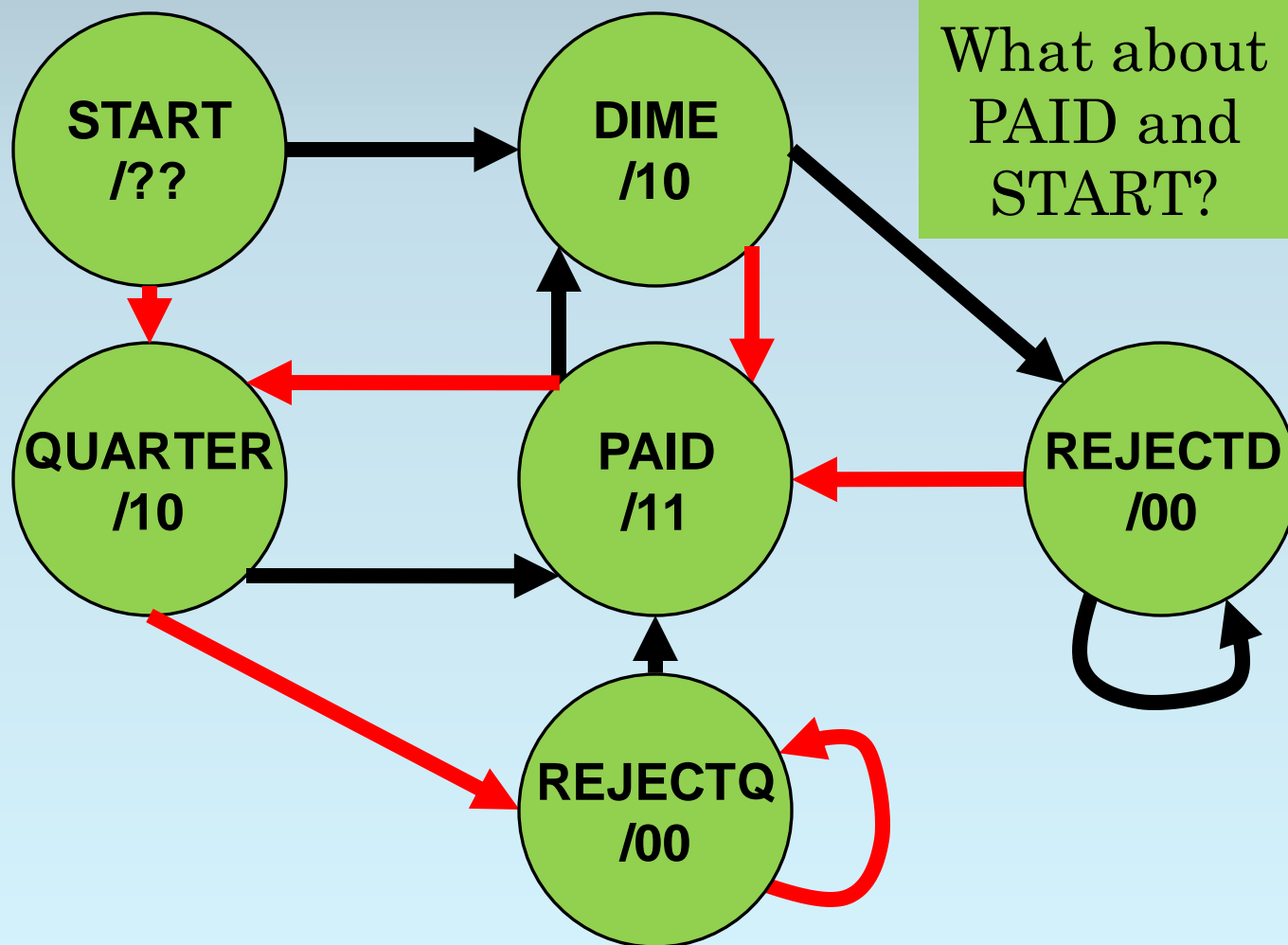
- find two states
- with identical outputs
- and identical next states.

Let's take a look.

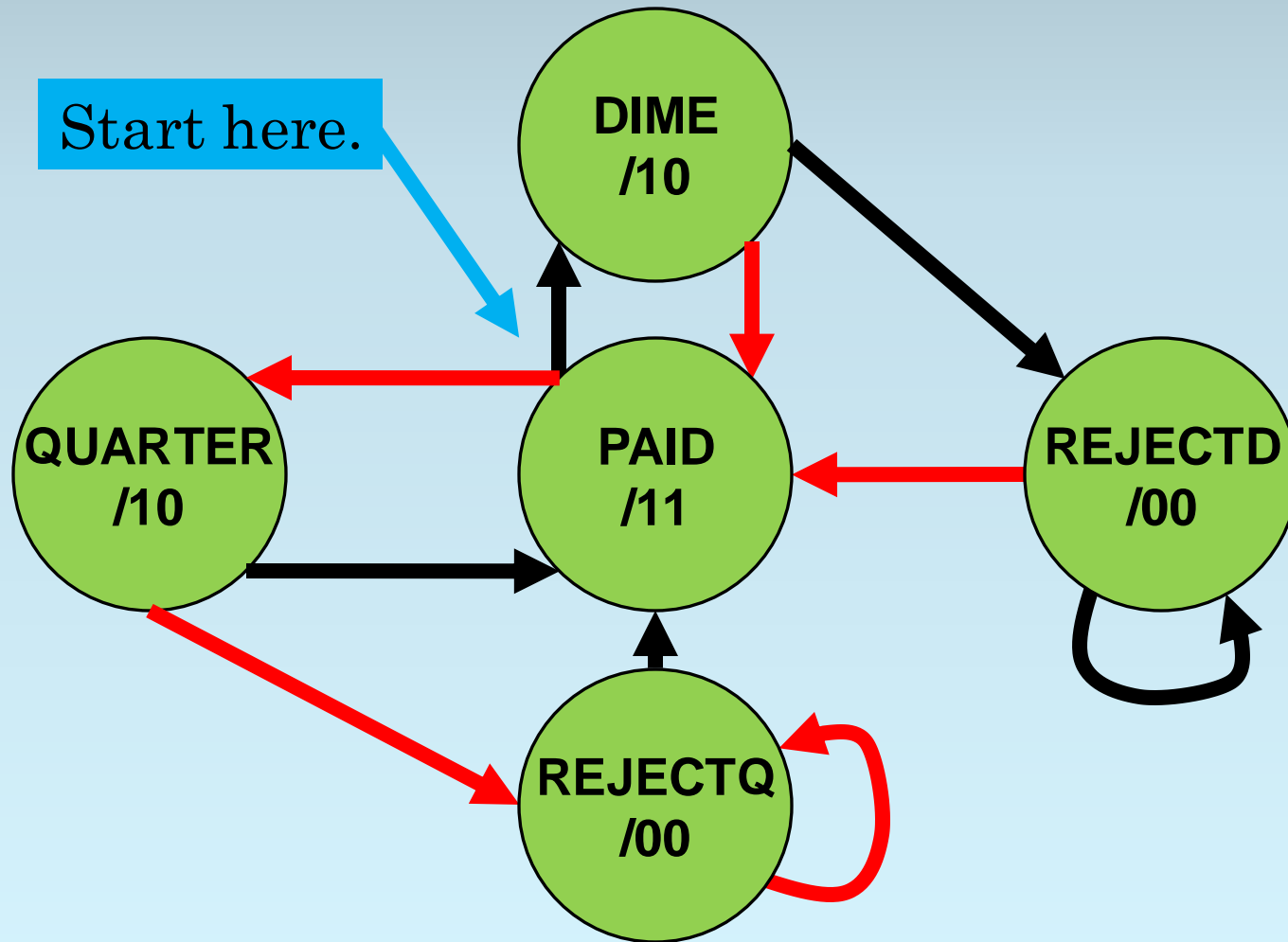
Step 3: Merge Redundant States



Any More Redundant States?



Our Final Abstract Model



Use Human Information to Define the Representation

We need 3 bits for 5 states.

Let's **use human information** to define the representation.

Think about the sequence of coins that has been inserted into the machine.

Let's call the last coin T_0 .

And the one before that T_{-1} .

And so forth.

State Bit Definition

Define the state bits as follows:

S_2 is T_0 , the last coin type.

S_1 is 1 iff one or more **quarters** were inserted

- before the last coin (T_{-1} , T_{-2} , and so on)
- but after the last product release.

S_0 is 1 iff one or more **dimes** were inserted

- before the last coin (T_{-1} , T_{-2} , and so on)
- but after the last product release.

Does that Approach Work? Check the State IDs

For **DIME**, the only coin since the last payment is a single dime.

So we have...

$S_2 = 0$ (the dime)

$S_1 = 0$ (no quarters at all)

$S_0 = 0$ (the dime is unique)

DIME **000**
QUARTER
REJECTD
REJECTQ
PAID

Calculate the State ID for QUARTER

For **QUARTER**, the only coin since the last payment is a single quarter.

So we have...

$S_2 = 1$ (the quarter)

$S_1 = 0$ (the quarter is unique)

$S_0 = 0$ (no dimes at all)

DIME 000

QUARTER 100

REJECTD

REJECTQ

PAID

Calculate the State ID for REJECTD

For **REJECTD**, we have seen two or more dimes but no quarters.

So we have...

$S_2 = 0$ (the last dime)

$S_1 = 0$ (no quarters at all)

$S_0 = 1$ (extra dimes)

DIME 000

QUARTER 100

REJECTD 001

REJECTQ

PAID

Calculate the State ID for REJECTQ

For **REJECTQ**, we have seen two or more quarters but no dimes.

So we have...

$S_2 = 1$ (the last quarter)

$S_1 = 1$ (extra quarters)

$S_0 = 0$ (no dimes at all)

DIME	000
QUARTER	100
REJECTD	001
REJECTQ	110
PAID	

What is the State ID for PAID?

For **PAID**, we could have gotten either coin last! Before the last coin, we got one or more of the other kind.

DIME **000**

For a last quarter, we have... **QUARTER** **100**

$S_2 = 1$ (the last quarter) **REJECTD** **001**

$S_1 = 0$ (no extra quarters) **REJECTQ** **110**

$S_0 = 1$ (at least one dime) **PAID** **101**

What is the State ID for PAID?

But what about a last dime?

$S_2 = 0$ (the last dime)

$S_1 = 1$ (at least one quarter)

$S_0 = 0$ (no extra dimes)

So we have two bit patterns!

DIME	000
QUARTER	100
REJECTD	001
REJECTQ	110
PAID	101
PAID	010

That's All!

You will finish the rest of the design and implement it in the lab...

In Lab 7, you will implement and simulate the FSM in Altera Quartus.

- You already have the output logic done.

In Lab 8, you will assemble your FSM on the prototyping board.

- You already have the output logic done.