

# ECE 220 Computer Systems & Programming

## Lecture 6 – Interrupts & Exceptions



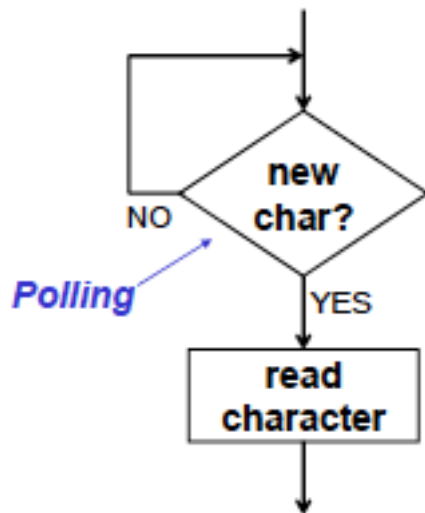
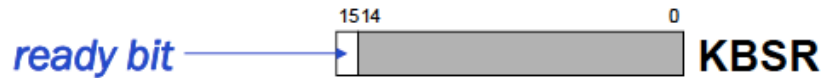
# Outline

- Interrupt driven I/O
- Interrupts and exceptions
- Machinery for processing interrupts

## Concepts

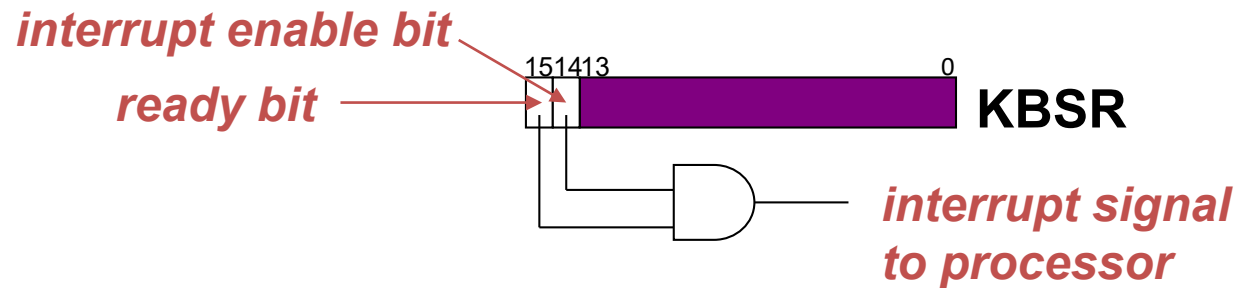
- New hardware to support new feature
- Saving processor state

# Polling V.S. Interrupt Driven I/O



```

POLL  LDI  R0, KBSR
      BRzp POLL
      LDI  R0, KBDR
      ...
KBSR  .FILL xFE00
KBDR  .FILL xFE02
    
```



- Software sets “interrupt enable” bit in device register
- When Ready bit and IE bit are both set, interrupt is signaled

# Interrupt Driven I/O

**An I/O device can:**

- **Force currently executing program to stop**
- **Have the processor carry out the need of the I/O device**
- **Resume the stopped program as if nothing had happened**

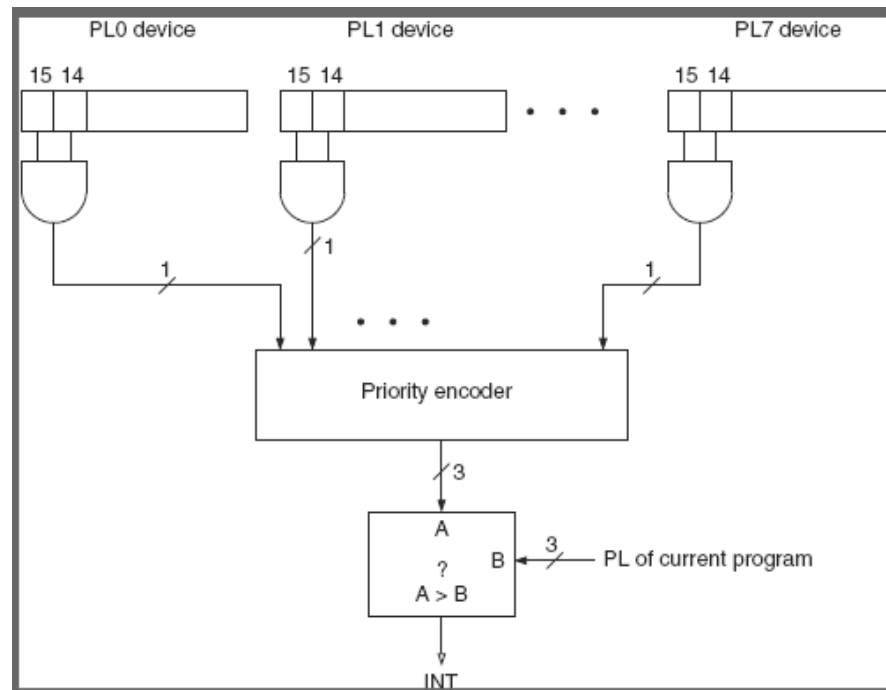
# Interrupt Driven I/O: 4 Steps

1. **Generation of interrupt signal (INT)**
2. **Set-up for servicing interrupt**
3. **Interrupt Service Routine**
4. **Resume execution main program**

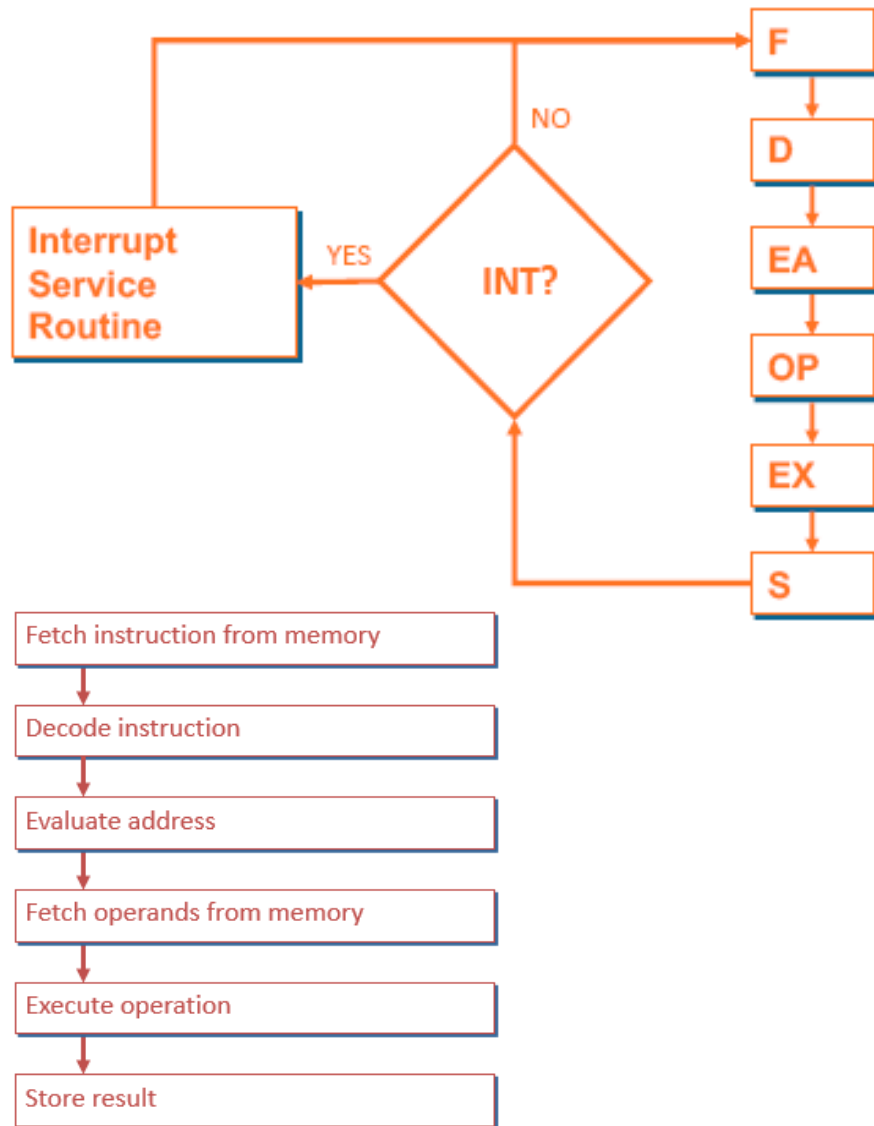
# Step 1. Generation of Interrupt Signal

Three conditions must be true for an I/O device to interrupt the processor:

1. The I/O device must want service (Ready Bit in KBSR)
2. The device must have the right to request service (IE bit in KBSR)
3. The device request must be more urgent than what the processor is currently doing



## Step 2: Instruction Cycle and Interrupt



- If INT is not asserted
  - The usual instruction cycle
  - Ends with check for interrupt
- If INT is asserted, control unit does 2 things before fetching the next instruction:
  - Store state of the program
  - Service the interrupt

# Processor Status Register (PSR)

- Enough state info saved for interrupted program to resume later
- Enough state info loaded for the interrupt service routine to begin service

## State of a Program:

- Contents of all GPRs (should be saved by ISR, callee-saved)
- PC
- PSR (processor status register)
- User Stack Pointer

PSR[15] – privileged (supervisor - 0) or unprivileged (user - 1) mode

PSR[10:8] – priority level, PSR[2:0] – condition code





# Supervisor Stack (where the state is saved)

Suppose the user program uses a stack, and the stack pointer is stored in R6

A special region of memory used as the stack (why?) for interrupt service routines.

- Supervisor Stack Pointer (SSP) stored in Saved.SSP
- Another register for storing User Stack Pointer (USP) Saved.USP

When switching from User mode to Supervisor mode (as result of interrupt), save R6 to Saved.USP

Access both Stacks using R6 as the stack pointer.

## 2. Invoking the Service Routine

I/O device transmits Interrupt Vector (INTV, 8-bit) along with interrupt signal and priority level (PL0..PL7). When an interrupt is taken:

1. If Priv = 1 (user), i.e. PSR[15] = 1  
Saved.USB = R6, then R6 = Saved.SSP.
2. Push PSR and PC to Supervisor Stack.
3. Set PSR[15] = 0 (supervisor mode).
4. Set PSR[10:8] = priority of interrupt being serviced. (keyboard PL4)
5. Set PSR[2:0] = 0.
6. Set MAR = x01vv, where vv = 8-bit interrupt vector provided by interrupting device (e.g., keyboard = x80).
7. Load memory location (M[x01vv]) into MDR.
8. Set PC = MDR; now first instruction of ISR will be fetched.

## Step 3: Interrupt Service Routine

- PC contains the starting address of the ISR.
- Callee-save for general purpose registers.
- Execute code for servicing the interrupt
- There may be other interrupts during this execution (in which case go back to Step 2)

## Step 4: Returning from interrupt

Special instruction (RTI) – restores state

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
<b>RTI</b>	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

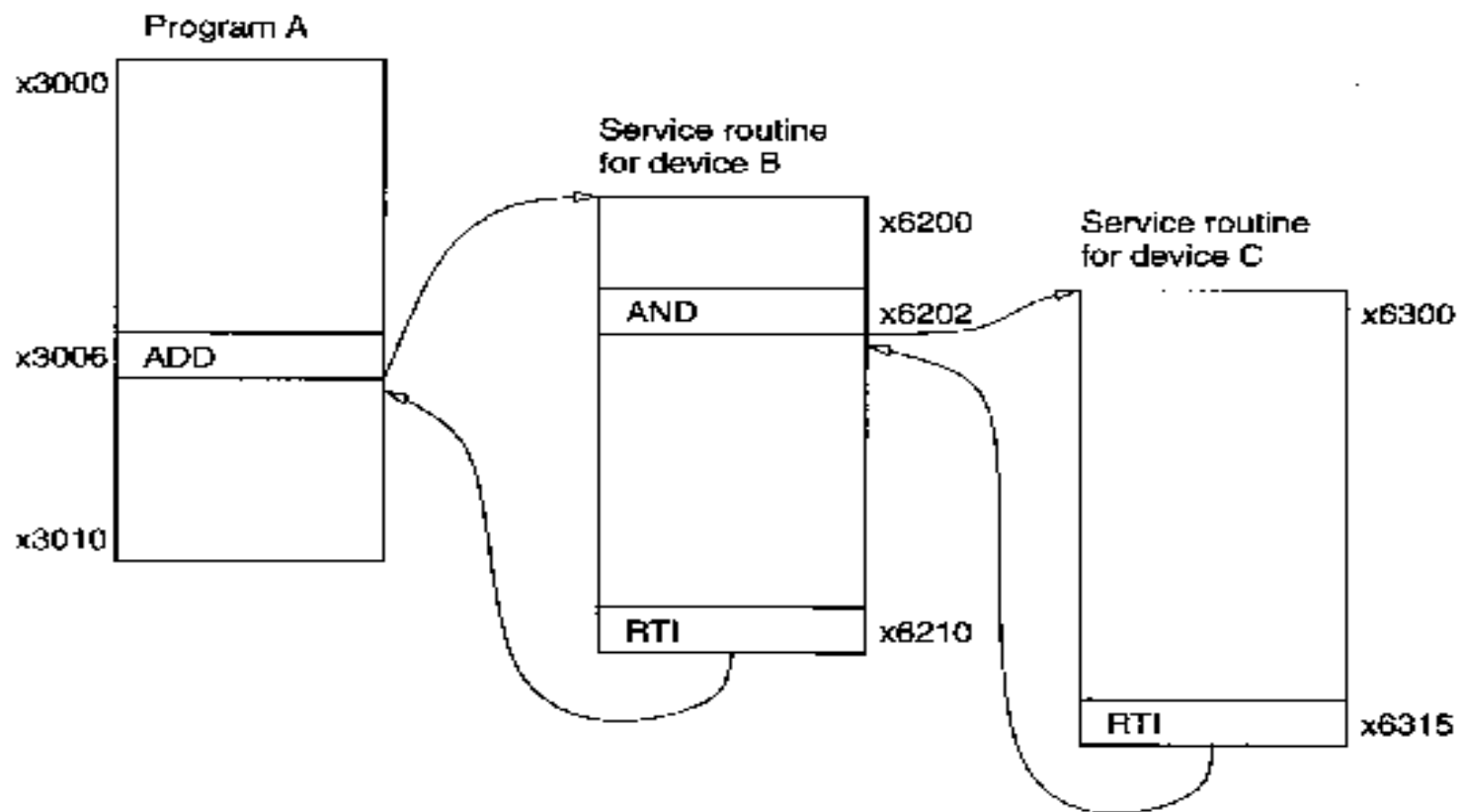
1. Pop PC from supervisor stack. ( $PC = M[R6]; R6 = R6 + 1$ )
2. Pop PSR from supervisor stack. ( $PSR = M[R6]; R6 = R6 + 1$ )
3. If  $PSR[15] = 1$ ,  $R6 = \text{Saved.USP}$ .  
(If going back to user mode, need to restore User Stack Pointer.)

RTI is a privileged instruction.

- Can only be executed in Supervisor Mode.
- If executed in User Mode, causes an exception.

# Nested Interrupts

Example:



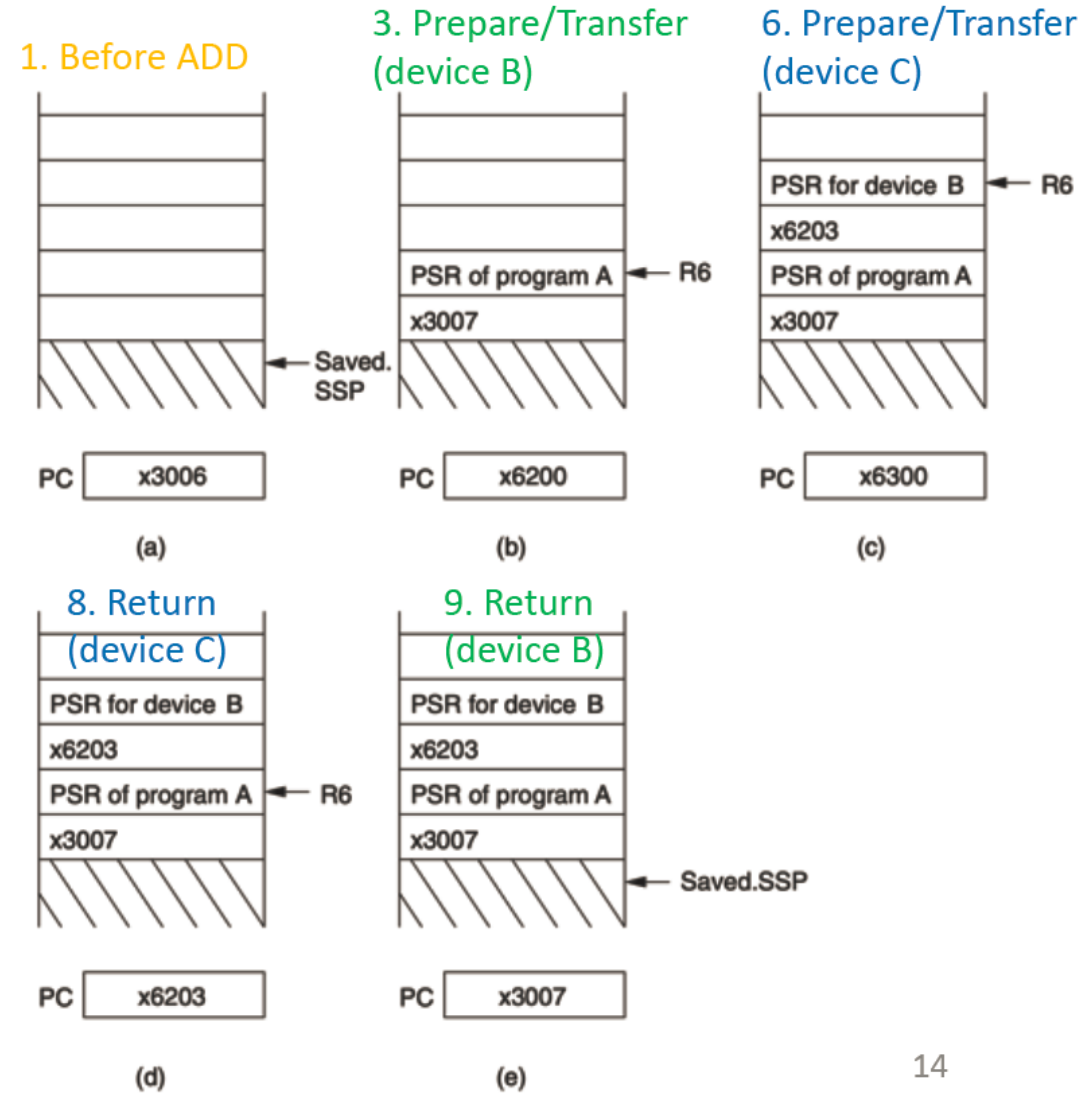
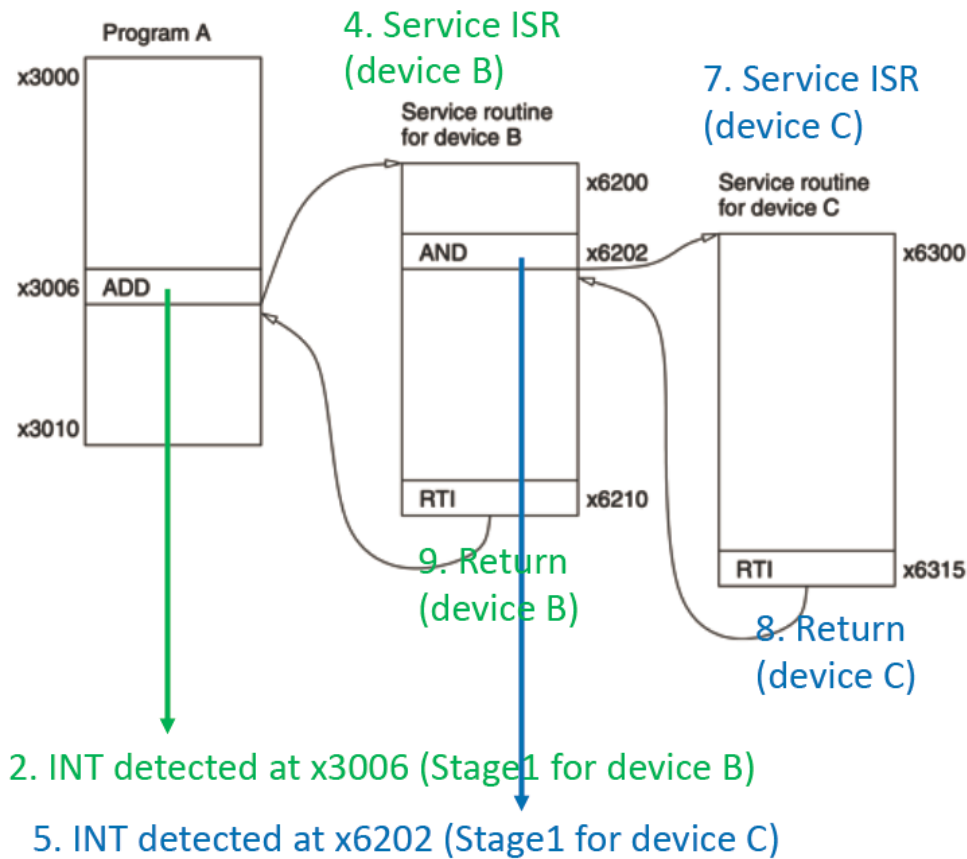
## Interrupt vector table

Addr	Data
x01F1	x6200
x01F2	x6300

## INTV

Device B = xF1  
Device C = xF2

PL: A<B<C



**Example Code: (See Github)**

# Exceptions: Internal Interrupt

When something unexpected happens inside the processor, it may cause an exception.

Examples of Exception in LC-3:

- Privileged operation (e.g., RTI in user mode)
- Executing an illegal opcode (Bits[15:12] = 1101)

Handled just like an interrupt

- Vector is determined internally by type of exception
- Priority is the same as running program

## Interrupt Vector Table

Exception Service Routines – x0100 to x017F

Interrupt Service Routines – x0180 to 01FF

