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Real-Time Task Assignment in Hyperlocal Spatial Crowdsourcing under Budget Constraints

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- ✓ Precipitation estimation technologies fail to provide real-time, fine-grained data
- ✓ Smartphones are awesom!!
- ✓ Weather crowdsourcing becomes popular



mPing



SkyMotion



Weewave



WEDDAR



Weathermob

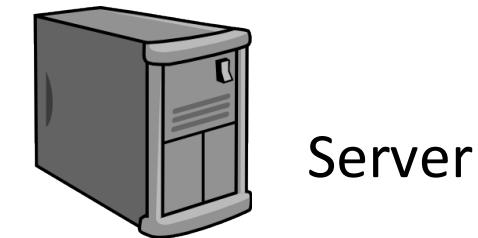
Crowdsourcing is regarded as “the future of weather forecasting”

Problem Setup



Requesters submit a set of tasks to the Server

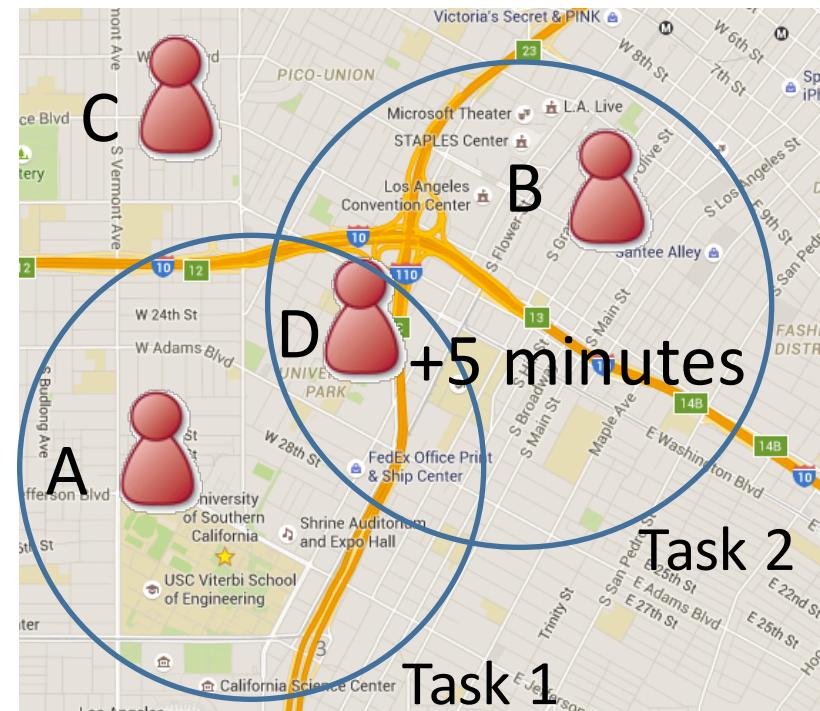
Tasks with spatiotemporal continuity in measurement, e.g., rainfall, temperature



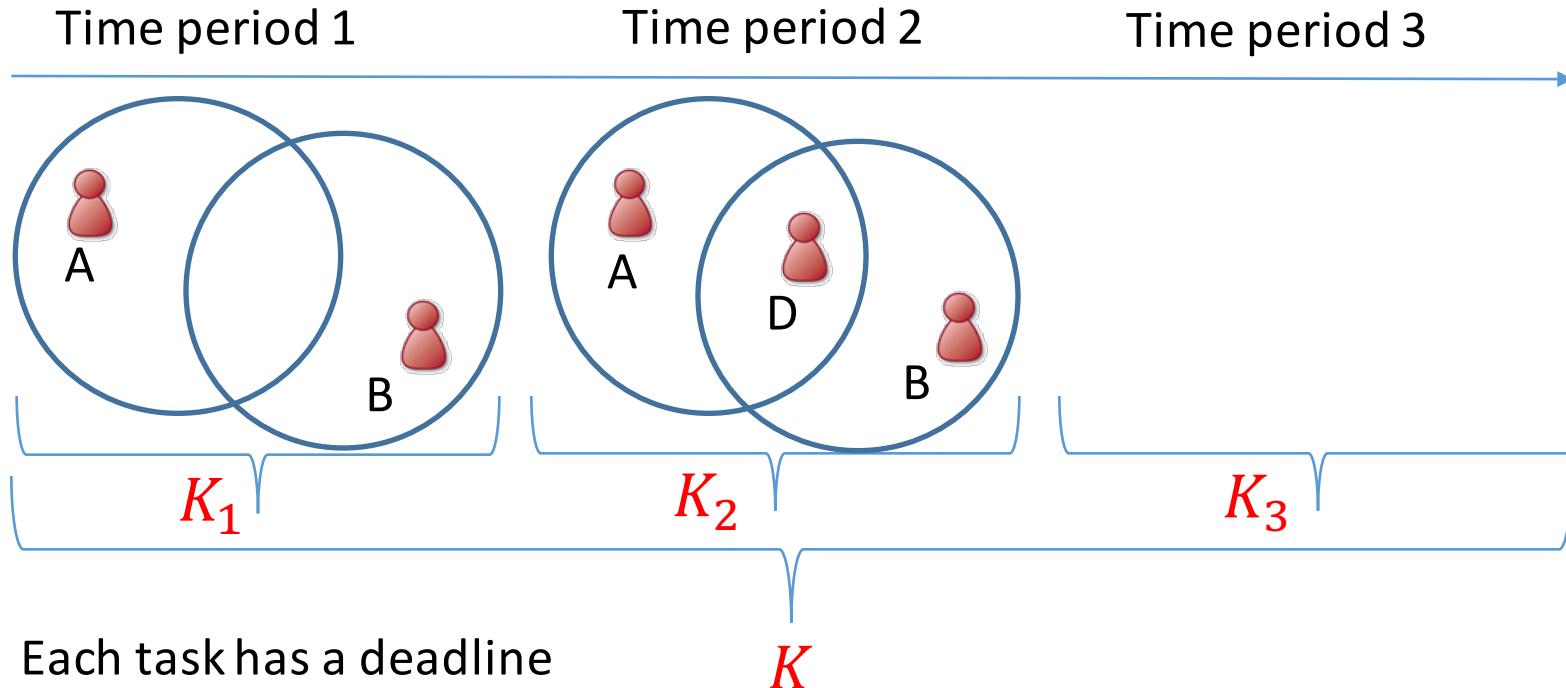
A worker can respond to a task if he is within the task region

Maximize #covered tasks given #workers to activate (budget K)

Dynamic arrivals of workers and tasks



Problem Space



given a total budget K
for all time periods

given a particular budget
 K_i per time period

server does not
know the future

server knows
everything

	Dynamic-budget	Fixed-budget
Online	Online dMTC	Online fMTC
Offline	Offline dMTC	Offline fMTC

Offline MTC Variants

Given Q time period $s_i \in \{s_1, \dots, S_Q\}$,

Dynamic-budget maximum task coverage (dMTC) is to select a set of workers L_i at every s_i s.t. #covered tasks is maximized and $\sum_{i=1}^Q |L_i| \leq K$

Fixed-budget MTC (fMTC) is similar to dMTC, except the budget is given for each time period $|L_i| \leq K_i$

dMTC is more challenging than fMTC

dMTC is NP-hard by reduction from maximum coverage problem (MCP) [Feige, JACM 1998]

fMTC is NP-hard by reduction from maximum coverage with group budgets constraints problem [Chekuri et. al. APPROX 2014]

Online dMTC – Which worker to select?

Difficult to achieve global optimal solution

Optimizes locally at every time period

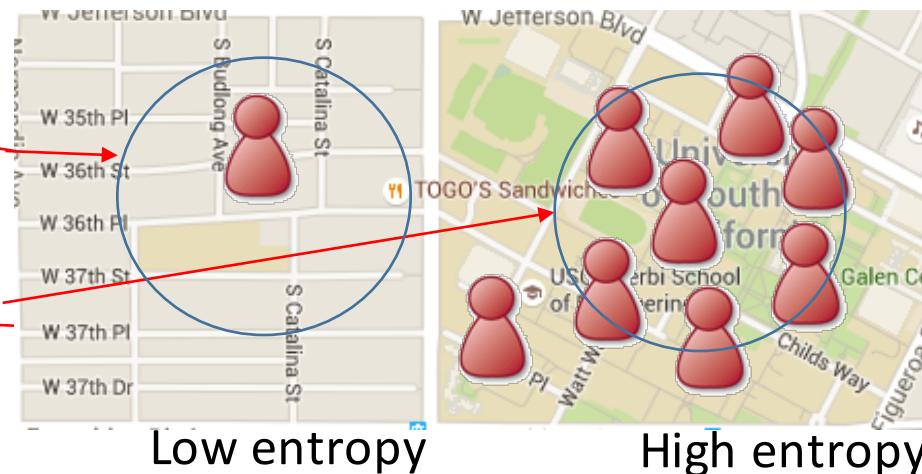
fMTC for one time period is NP-hard by reduction from MCP

Greedy algorithm (*Basic*) to MCP chooses the worker that covers the largest number of uncovered elements [Feige JACM'98]

Basic does not consider spatiotemporal information of each task

- *Temporal heuristic*: expiring tasks have higher priority
- *Spatial heuristic*: tasks in worker-sparse area have higher priority

Expired in 5 time periods
Task in worker-sparse region



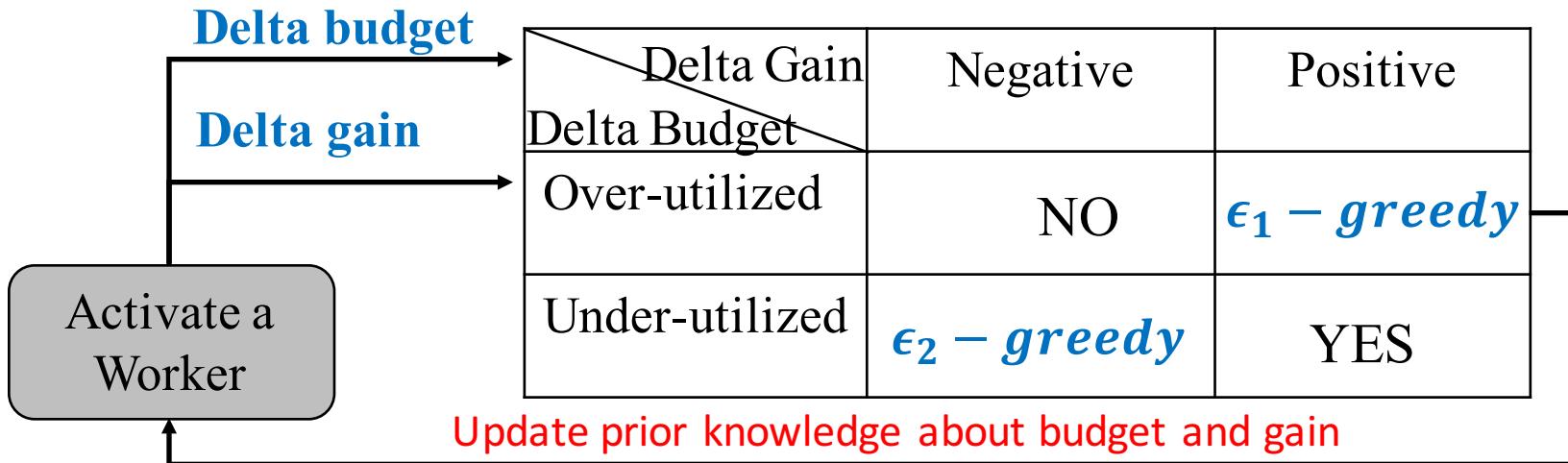
Task in worker-dense region
Expired in 1 time period

Online dMTC – How many to select?

Allocate overall budget K over Q time periods

Equal: strategy equally divides K to Q time periods K/Q

Adapt: decision of how many workers to activate depends on 1) budget utilization 2) distribution of workers and tasks



Balance with explore (YES) and exploit (NO) tradeoff

Use contextual $\epsilon - \text{greedy}$ algorithm

[Li et. al. WWW 2010]

Experimental Setup

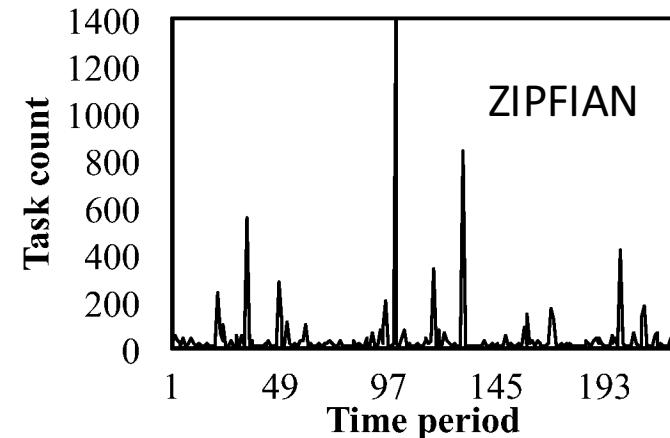
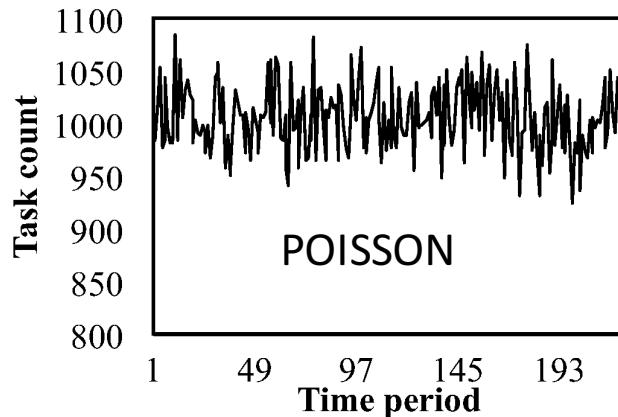
Real-world datasets

Name	#Tasks (Check-ins)	#Workers (Users)	Unit time period	#periods
Gowalla	151,075	6,160	1 day	224
Foursquare	89,968	45,138	1 hour	384

Time period Q = 28. Task duration = 5 time periods. Task radius = 5 km

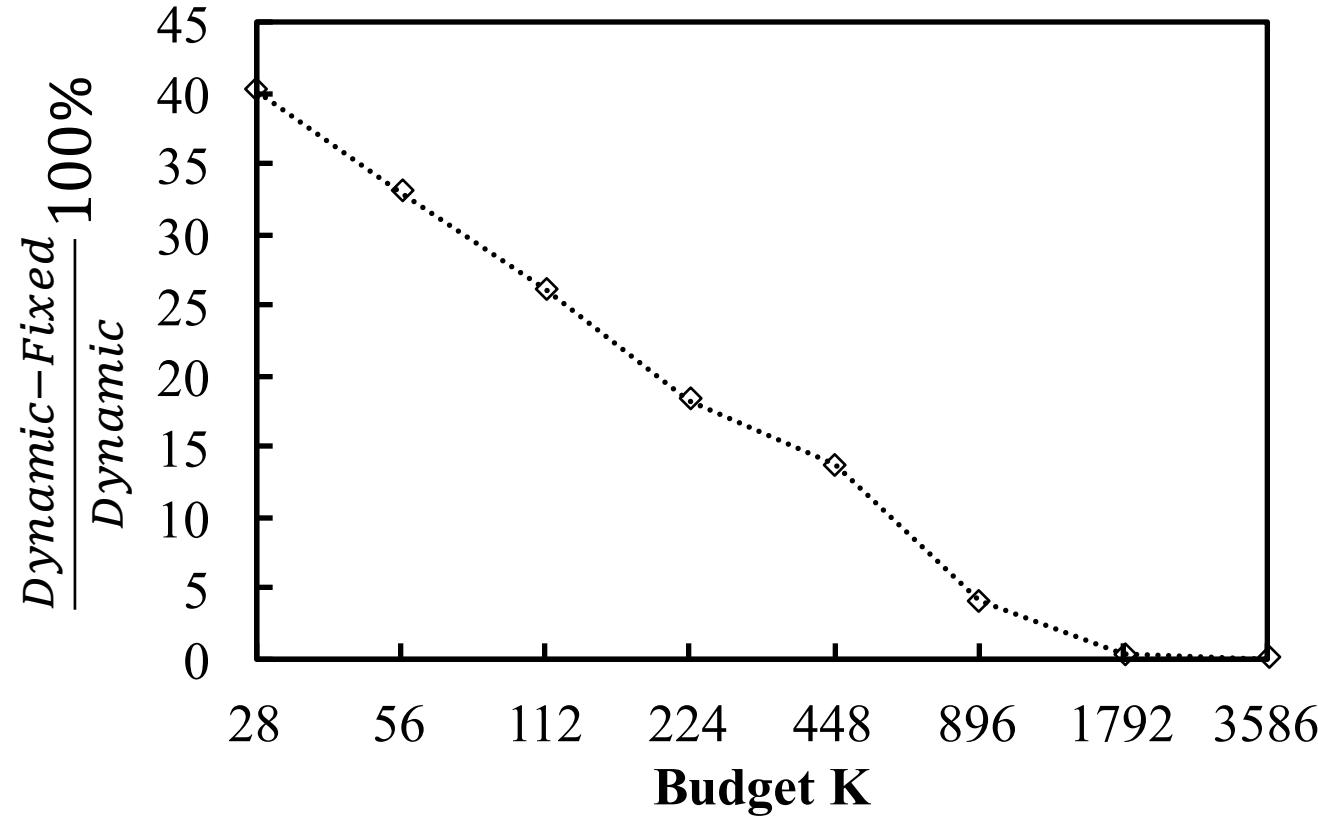
Synthetic datasets

- Synthetically arrival rates of workers and tasks, e.g., GO-ZIPFIAN



Offline Solutions

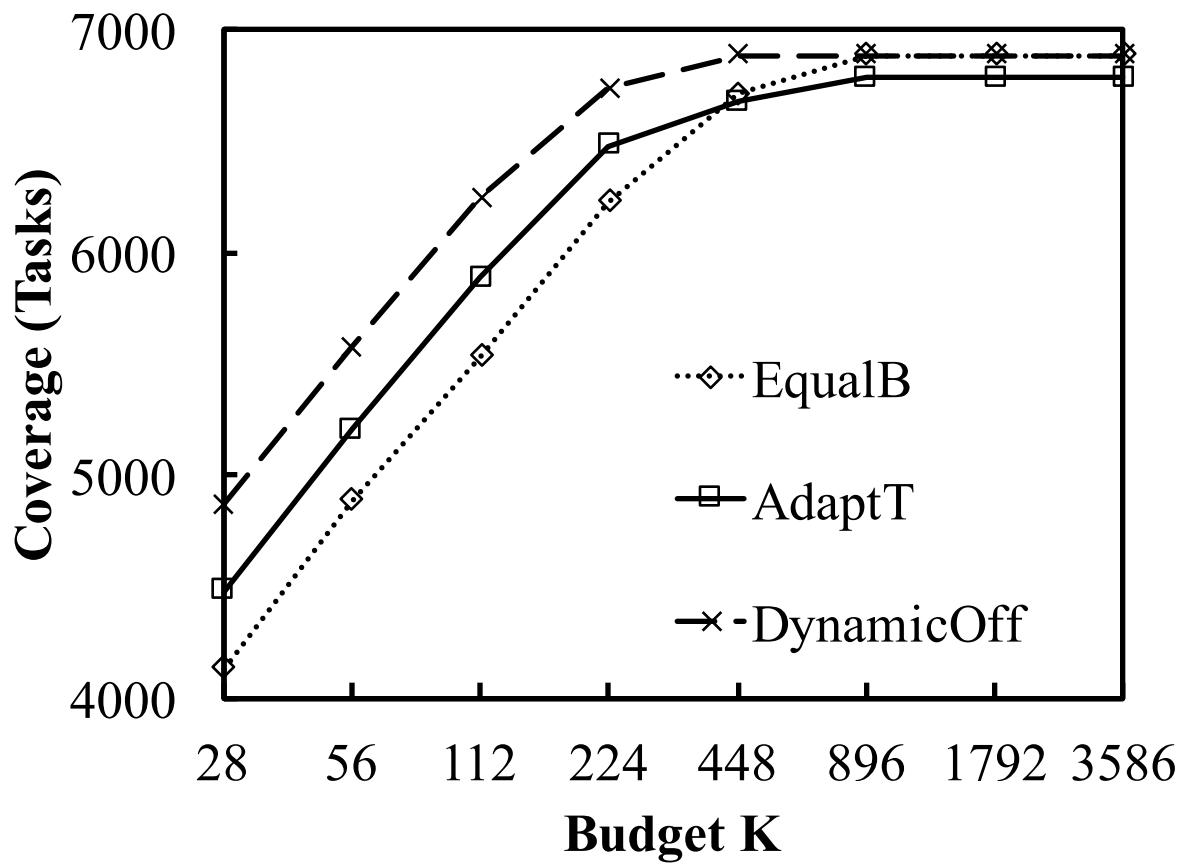
Varying budget K (Go-ZIPFIAN)



- ✓ Dynamic budget allocation is more effective
- ✓ The higher the budget, the smaller the improvement

Dynamic-Budget Online

Varying radius r (Go-POISSON)



- ✓ *AdaptT (Adapt + Temporal) comes closer to offline solution*

Conclusion

- ✓ Proposed a framework to maximize task coverage under budget constraints
- ✓ Developed three online heuristics, *Basic*, *Temporal* and *Spatial* for fMTC
- ✓ Developed online budget allocation to dynamically allocate total budget to time periods for dMTC
- ✓ Empirical results confirm that online solutions increase task coverage by 40% over basic approach

Q/A

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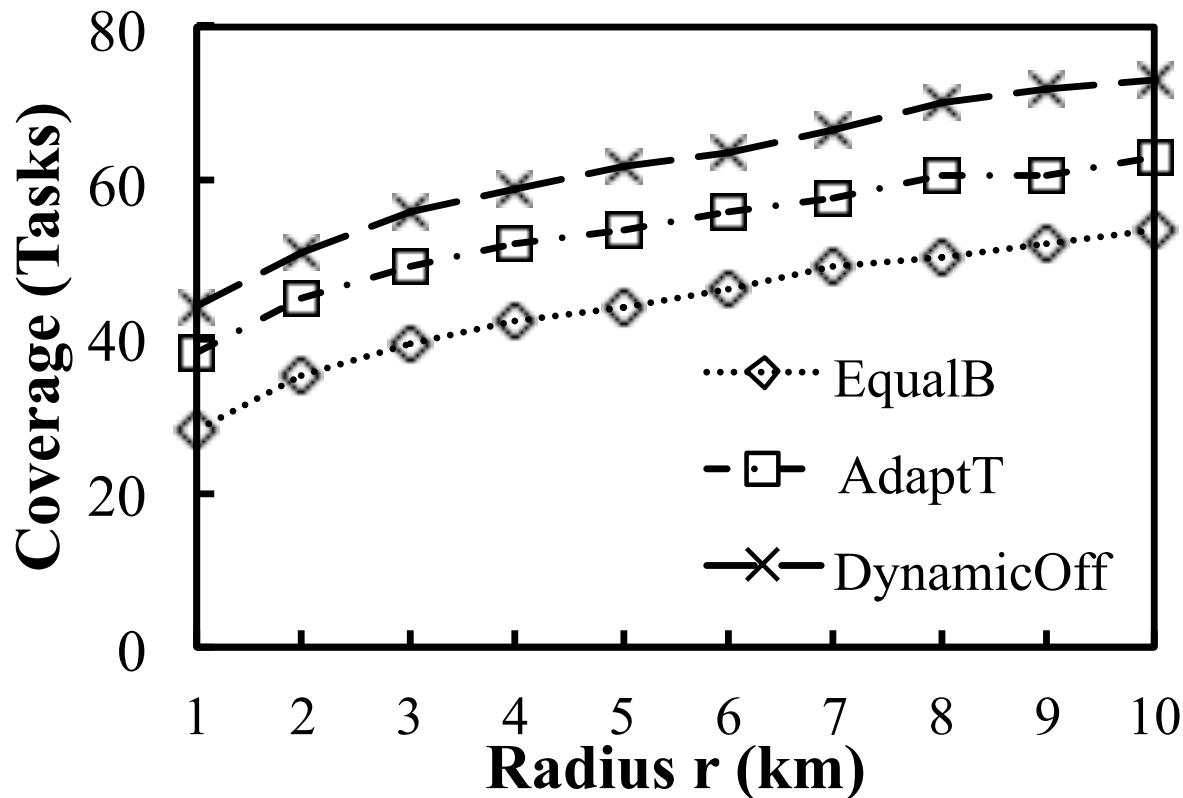
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Dynamic-Budget Online

Varying radius r (Go-POISSON)



- ✓ *AdaptT (Adapt + Temporal) obtains up to 40% improvements (at $r=1$)*