Lab of Computer Network: TCP-Stack-2 Fall 2024

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14.1 实验内容

1. 丢包恢复

- 执行 create_randfile.sh,生成待传输数据文件 client-input.dat。
- •运行给定网络拓扑(tcp_topo_loss.py)。
- 在节点 h1 上执行 TCP 程序:
 - 执行脚本 (disable_offloading.sh 和 disable_tcp_rst.sh), 禁止协议栈的相应功能。
 - 在 h1 上运行 TCP 协议栈的服务器模式: ./tcp_stack server 10001。
- 在节点 h2 上执行 TCP 程序:
 - 执行脚本 (disable_offloading.sh 和 disable_tcp_rst.sh), 禁止协议栈的相应功能。
 - 在 h2 上运行 TCP 协议栈的客户端模式: ./tcp_stack client 10.0.0.1 10001。
- Client 发送文件 client-input.dat 给 server, server 将收到的数据存储到文件 server-output.dat。
- 使用 md5sum 比较两个文件是否完全相同。
- 使用 tcp_stack.py 替换两端任意一方,对端都能正确处理数据收发。

2. 拥塞控制

- 执行 create_randfile.sh,生成待传输数据文件 client-input.dat。
- 运行给定网络拓扑(tcp_topo_loss.py)。
- 在节点 h1 上执行 TCP 程序:
 - 执行脚本 (disable_offloading.sh 和 disable_tcp_rst.sh), 禁止协议栈的相应功能。
 - 在 h1 上运行 TCP 协议栈的服务器模式: ./tcp_stack server 10001。
- 在节点 h2 上执行 TCP 程序:
 - 执行脚本 (disable_offloading.sh 和 disable_tcp_rst.sh), 禁止协议栈的相应功能。
 - 在 h2 上运行 TCP 协议栈的客户端模式: ./tcp_stack client 10.0.0.1 10001。

- Client 发送文件 client-input.dat 给 server, server 将收到的数据存储到文件 server-output.dat。
- 使用 md5 sum 比较两个文件是否完全相同。
- 记录 h2 中每次 cwnd 调整的时间和相应值,呈现到二维坐标图中。

14.2 实验过程

14.2.1 丢包恢复

1. 重传定时器操作

本次实验中,我们需要实现定时器的设置、更新、关闭、扫描操作。

设置操作较为简单,先检查是否有定时器,如果没有则直接设置定时器类型为重传定时器,启用并设置超时时间,设置重传次数为0;如果已有定时器,则更新超时时间。最后把定时器加入定时器链表。

更新操作也较为简单,若已建立的连接发送队列为空,则关闭并删除定时器,唤醒发送数据进程。

关闭操作也较为简单,与更新操作类似,只是判断条件改为定时器队列不为空。

扫描操作则需要遍历定时器链表,减少每个定时器的剩余时间,若剩余时间小于等于 0,即超时,则进行处理:重传次数未达上限则重传数据包,重传次数达到上限则直接断开连接,释放资源。扫描操作通过一个线程每 10 秒进行一次扫描。

代码如下:

```
// set the restrans timer of a tcp sock, by adding the timer into
   timer_list
void tcp_set_retrans_timer(struct tcp_sock *tsk)
   if (tsk->retrans timer.enable) {
      tsk->retrans_timer.timeout = TCP_RETRANS_INTERVAL_INITIAL;
      return;
   }
   tsk->retrans_timer.type = TIMER_TYPE_RETRANS;
   tsk->retrans timer.enable = 1;
   tsk->retrans_timer.timeout = TCP_RETRANS_INTERVAL_INITIAL;
   tsk->retrans_timer.retrans_time = 0;
   pthread mutex lock(&retrans timer list lock);
   list_add_tail(&tsk->retrans_timer.list, &retrans_timer_list);
   pthread_mutex_unlock(&retrans_timer_list_lock);
}
void tcp_update_retrans_timer(struct tcp_sock *tsk)
   if (list empty(&tsk->send buf) && tsk->retrans timer.enable) {
```

```
tsk->retrans_timer.enable = 0;
      list_delete_entry(&tsk->retrans_timer.list);
      wake_up(tsk->wait_send);
   }
}
void tcp_unset_retrans_timer(struct tcp_sock *tsk)
   if (!list_empty(&tsk->retrans_timer.list)) {
      tsk->retrans_timer.enable = 0;
      list_delete_entry(&tsk->retrans_timer.list);
      wake_up(tsk->wait_send);
   }
   else
      log(ERROR, "unset an empty retrans timer\n");
}
void tcp_scan_retrans_timer_list(void)
{
   struct tcp_sock *tsk;
   struct tcp_timer *time_entry, *time_q;
   pthread_mutex_lock(&retrans_timer_list_lock);
   list_for_each_entry_safe(time_entry, time_q, &retrans_timer_list,
      list) {
      time_entry->timeout -= TCP_RETRANS_SCAN_INTERVAL;
      tsk = retranstimer_to_tcp_sock(time_entry);
      if (time_entry->timeout <= 0) {</pre>
         if(time_entry->retrans_time >= MAX_RETRANS_NUM && tsk->state !=
            TCP_CLOSED) {
            list_delete_entry(&time_entry->list);
            if (!tsk->parent)
               tcp_unhash(tsk);
            wait_exit(tsk->wait_connect);
            wait_exit(tsk->wait_accept);
            wait_exit(tsk->wait_recv);
            wait_exit(tsk->wait_send);
            tcp_set_state(tsk, TCP_CLOSED);
            tcp_send_control_packet(tsk, TCP_RST);
```

```
else if (tsk->state != TCP_CLOSED) {
            time entry->retrans time += 1;
            log(DEBUG, "retrans time: %d\n", time_entry->retrans_time);
            time_entry->timeout = TCP_RETRANS_INTERVAL_INITIAL;
            tcp_retrans_send_buffer(tsk);
         }
      }
   }
   pthread_mutex_unlock(&retrans_timer_list_lock);
}
void *tcp_retrans_timer_thread(void *arg)
   init_list_head(&retrans_timer_list);
   while(1){
      usleep(TCP_RETRANS_SCAN_INTERVAL);
      tcp_scan_retrans_timer_list();
   }
   return NULL;
}
```

2. 发送队列

对于发送队列,我们需要实现数据包的添加、更新、重传操作,都比较简单。

添加操作只需将数据包加入发送队列即可。注意这里需要使用互斥锁保护发送队列。

更新操作是指遍历队列,将队列中序列号小于收到的 ACK 的数据包删除。这里同样需要使用互斥锁保护发送队列。

重传操作是指超时未收到 ACK 时,将队列中第一个的数据包重传。我将队列中第一个数据包的 TCP 序列号、确认号等信息更新之后,再计算出数据长度和发送窗口大小,将数据包发送出去。

代码如下:

```
// Add a packet to the TCP send buffer
void tcp_send_buffer_add_packet(struct tcp_sock *tsk, char *packet, int
    len) {
    send_buffer_entry_t *send_buffer_entry = (send_buffer_entry_t
        *)malloc(sizeof(send_buffer_entry_t));
    memset(send_buffer_entry, 0, sizeof(send_buffer_entry_t));

send_buffer_entry->packet = (char *)malloc(len);
send_buffer_entry->len = len;
```

```
memcpy(send_buffer_entry->packet, packet, len);
   init_list_head(&send_buffer_entry->list);
  list_add_tail(&send_buffer_entry->list, &tsk->send_buf);
}
//Update the TCP send buffer based on the acknowledgment number
void tcp_update_send_buffer(struct tcp_sock *tsk, u32 ack) {
   send_buffer_entry_t *send_buffer_entry, *send_buffer_entry_q;
   list_for_each_entry_safe(send_buffer_entry, send_buffer_entry_q,
      &tsk->send_buf, list) {
     struct tcphdr *tcp = packet_to_tcp_hdr(send_buffer_entry->packet);
     u32 seq = ntohl(tcp->seq);
     // If the sequence number is less than the acknowledgment number,
         delete the entry
     if (less_than_32b(seq, ack)) {
         list_delete_entry(&send_buffer_entry->list);
         free(send_buffer_entry->packet);
         free(send_buffer_entry);
     }
   }
}
// Retransmit the first packet in the TCP send buffer when ack time exceed
int tcp_retrans_send_buffer(struct tcp_sock *tsk) {
   if (list_empty(&tsk->send_buf)) {
      log(ERROR, "no packet to retrans\n");
     pthread_mutex_unlock(&tsk->send_buf_lock);
      return 0;
  }
  // Retrieve the first send buffer entry
   send_buffer_entry_t *first_send_buffer_entry =
      list_entry(tsk->send_buf.next, send_buffer_entry_t, list);
   char *packet = (char *)malloc(first_send_buffer_entry->len);
  // Copy the packet data and update TCP sequence and acknowledgment
      numbers
```

```
memcpy(packet, first_send_buffer_entry->packet,
    first_send_buffer_entry->len);
struct iphdr *ip = packet_to_ip_hdr(packet);
struct tcphdr *tcp = packet_to_tcp_hdr(packet);
tcp->ack = htonl(tsk->rcv_nxt);
tcp->checksum = tcp_checksum(ip, tcp);
ip->checksum = ip_checksum(ip);

// Calculate TCP data length and update TCP send window
int tcp_data_len = ntohs(ip->tot_len) - IP_BASE_HDR_SIZE -
    TCP_BASE_HDR_SIZE;
tsk->snd_wnd -= tcp_data_len;

log(DEBUG, "retrans seq: %u\n", ntohl(tcp->seq));

// Send the packet
ip_send_packet(packet, first_send_buffer_entry->len);
return 1;
}
```

3. 接收队列

对于接收队列,我们需要实现数据包的添加、移动操作,稍微繁琐。

添加操作要将接收到的数据包按照 seq 顺序插入接收队列中, 若出现重复数据包则直接丢弃, 然后将有效的数据包插入队列中。

移动操作首先遍历接收队列,找到与当前 rcv_nxt 匹配的数据包,然后写入环形缓冲区,唤醒接收进程, 更新 rcv_nxt。注意环形缓冲区需要使用互斥锁保护。

代码如下:

```
// Add an packet to the TCP receive buffer
int tcp_recv_ofo_buffer_add_packet(struct tcp_sock *tsk, struct tcp_cb
    *cb) {
    if (cb->pl_len <= 0)
        return 0;
    recv_ofo_buf_entry_t *recv_ofo_entry = (recv_ofo_buf_entry_t
            *)malloc(sizeof(recv_ofo_buf_entry_t));
    recv_ofo_entry->seq = cb->seq;
    recv_ofo_entry->seq_end = cb->seq_end;
    recv_ofo_entry->len = cb->pl_len;
    recv_ofo_entry->data = (char *)malloc(cb->pl_len);
    memcpy(recv_ofo_entry->data, cb->payload, cb->pl_len);
    init_list_head(&recv_ofo_entry->list);
```

```
// insert the new entry at the correct position
   recv_ofo_buf_entry_t *entry, *entry_q;
   list_for_each_entry_safe (entry, entry_q, &tsk->rcv_ofo_buf, list) {
      if (recv_ofo_entry->seq == entry->seq)
         return 1; // same seq, do not add
      if (less_than_32b(recv_ofo_entry->seq, entry->seq)) {
         list_add_tail(&recv_ofo_entry->list, &entry->list);
         return 1;
      }
   }
   list_add_tail(&recv_ofo_entry->list, &tsk->rcv_ofo_buf);
   return 1;
}
// Move packets from TCP receive buffer to ring buffer
int tcp_move_recv_ofo_buffer(struct tcp_sock *tsk) {
   recv_ofo_buf_entry_t *entry, *entry_q;
   list_for_each_entry_safe(entry, entry_q, &tsk->rcv_ofo_buf, list) {
      if (tsk->rcv_nxt == entry->seq) {
         // Wait until there is enough space in the receive buffer
         while (ring_buffer_free(tsk->rcv_buf) < entry->len)
            sleep_on(tsk->wait_recv);
         pthread_mutex_lock(&tsk->rcv_buf_lock);
         write_ring_buffer(tsk->rcv_buf, entry->data, entry->len);
         tsk->rcv_wnd -= entry->len;
         pthread_mutex_unlock(&tsk->rcv_buf_lock);
         wake_up(tsk->wait_recv);
         // Update seq and free memory
         tsk->rcv_nxt = entry->seq_end;
         list_delete_entry(&entry->list);
         free(entry->data);
         free(entry);
      }
      else if (less_than_32b(tsk->rcv_nxt, entry->seq))
         continue; //the next expected sequence number is not reached yet
      else {
         log(ERROR, "rcv_nxt is more than seq, rcv_nxt: %d, seq: %d\n",
            tsk->rcv_nxt, entry->seq);
         return 0;
      }
```

```
return 1;
}
```

4. TCP 核心函数更新

为了实现可靠传输,tcp_send_packet 函数需要在发送数据包时将数据包加入发送队列,设置重传定时器:

```
// send a tcp packet
//
// Given that the payload of the tcp packet has been filled, initialize
   the tcp
// header and ip header (remember to set the checksum in both header),
   and emit
// the packet by calling ip_send_packet.
void tcp_send_packet(struct tcp_sock *tsk, char *packet, int len)
{
   struct iphdr *ip = packet_to_ip_hdr(packet);
   struct tcphdr *tcp = (struct tcphdr *)((char *)ip + IP_BASE_HDR_SIZE);
   int ip_tot_len = len - ETHER_HDR_SIZE;
   int tcp_data_len = ip_tot_len - IP_BASE_HDR_SIZE - TCP_BASE_HDR_SIZE;
   u32 saddr = tsk->sk_sip;
   u32 daddr = tsk->sk_dip;
   u16 sport = tsk->sk sport;
   u16 dport = tsk->sk_dport;
   u32 seq = tsk->snd_nxt;
   u32 ack = tsk->rcv nxt;
   u16 rwnd = tsk->rcv_wnd;
   tcp_init_hdr(tcp, sport, dport, seq, ack, TCP_PSH|TCP_ACK, rwnd);
   ip_init_hdr(ip, saddr, daddr, ip_tot_len, IPPROTO_TCP);
   tcp->checksum = tcp checksum(ip, tcp);
   ip->checksum = ip_checksum(ip);
   tsk->snd_nxt += tcp_data_len;
   tsk->snd_wnd -= tcp_data_len;
   tcp_send_buffer_add_packet(tsk, packet, len);
   tcp_set_retrans_timer(tsk);
   ip_send_packet(packet, len);
}
```

同样,为了实现可靠传输,tcp_send_control_packet 函数需要在发送控制包时设置重传定时器:

```
// send a tcp control packet
//
// The control packet is like TCP_ACK, TCP_SYN, TCP_FIN (excluding
   TCP_RST).
// All these packets do not have payload and the only difference among
   these is
// the flags.
void tcp_send_control_packet(struct tcp_sock *tsk, u8 flags)
   int pkt_size = ETHER_HDR_SIZE + IP_BASE_HDR_SIZE + TCP_BASE_HDR_SIZE;
   char *packet = malloc(pkt_size);
   if (!packet) {
      log(ERROR, "malloc tcp control packet failed.");
      return ;
   }
   struct iphdr *ip = packet_to_ip_hdr(packet);
   struct tcphdr *tcp = (struct tcphdr *)((char *)ip + IP_BASE_HDR_SIZE);
   u16 tot_len = IP_BASE_HDR_SIZE + TCP_BASE_HDR_SIZE;
   ip_init_hdr(ip, tsk->sk_sip, tsk->sk_dip, tot_len, IPPROTO_TCP);
   tcp_init_hdr(tcp, tsk->sk_sport, tsk->sk_dport, tsk->snd_nxt, \
         tsk->rcv_nxt, flags, tsk->rcv_wnd);
   tcp->checksum = tcp_checksum(ip, tcp);
   if (flags & (TCP_SYN|TCP_FIN))
      tsk->snd nxt += 1;
   if ((flags != TCP_ACK) && !(flags & TCP_RST)) {
      tcp_send_buffer_add_packet(tsk, packet, pkt_size);
      tcp_set_retrans_timer(tsk);
   }
   ip_send_packet(packet, pkt_size);
}
```

TCP 连接的状态机需要做较多的改动。连接建立过程中,上一个状态发送的包可能会丢失,故超时重传需要在 SYN_SENT 和 SYN_RECV 状态下进行,若收到回应则清空发送队列。此外,纯 ACK 包不需要重传:

```
case TCP_SYN_SENT:
   if (cb->flags == (TCP_SYN | TCP_ACK)) {
     tsk->rcv_nxt = cb->seq_end;
     tcp_update_window_safe(tsk, cb);
```

```
tsk->snd_una = greater_than_32b(cb->ack, tsk->snd_una) ? cb->ack :
         tsk->snd_una;;
      tcp_unset_retrans_timer(tsk);
      tcp_update_send_buffer(tsk, cb->ack);
      tcp_set_state(tsk, TCP_ESTABLISHED);
      tcp_send_control_packet(tsk, TCP_ACK);
      wake_up(tsk->wait_connect);
   }
   else if (cb->flags == TCP_SYN) {
      tsk->rcv_nxt = cb->seq_end;
      tcp_set_state(tsk, TCP_SYN_RECV);
      tcp_send_control_packet(tsk, TCP_SYN | TCP_ACK);
   }
   else
      log(DEBUG, "Current state is TCP_SYN_SENT but recv not SYN or
         SYN | ACK");
   break;
case TCP_SYN_RECV:
   if (cb->flags == TCP_ACK) {
      if (!is_tcp_seq_valid(tsk, cb))
         return;
      tsk->rcv_nxt = cb->seq_end;
      tcp_update_window_safe(tsk, cb);
      tsk->snd_una = greater_than_32b(cb->ack, tsk->snd_una) ? cb->ack :
         tsk->snd_una;;
      if (tsk->parent) {
         if (tcp sock accept queue full(tsk->parent)) {
            tcp_set_state(tsk, TCP_CLOSED);
            tcp_send_control_packet(tsk, TCP_RST);
            tcp_unhash(tsk);
            tcp_bind_unhash(tsk);
            list_delete_entry(&tsk->list);
            free_tcp_sock(tsk);
            log(DEBUG, "tcp_sock accept queue is full, so the tsk should
               be freed.");
         }
         else {
            tcp_set_state(tsk, TCP_ESTABLISHED);
```

```
tcp_sock_accept_enqueue(tsk);

tcp_unset_retrans_timer(tsk);

tcp_update_send_buffer(tsk, cb->ack);

wake_up(tsk->parent->wait_accept);
}
else
log(ERROR, "tsk->parent is NULL\n");
}
else
log(DEBUG, "Current state is TCP_SYN_RECV but recv not ACK");
break;
```

关闭连接时,需要在 FIN_WAIT_1 和 LAST_ACK 状态下进行超时重传,若收到回应则清空发送队列,关闭定时器,更新状态:

```
case TCP_FIN_WAIT_1:
   if (!is_tcp_seq_valid(tsk, cb))
      return;
   tsk->rcv_nxt = cb->seq_end;
   if (cb->flags & TCP_ACK) {
      tcp_update_send_buffer(tsk, cb->ack);
      tcp_unset_retrans_timer(tsk);
      tcp_update_window_safe(tsk, cb);
      tsk->snd una = cb->ack;
   }
   if ((cb->flags & TCP_FIN) && (cb->flags & TCP_ACK) && tsk->snd_nxt ==
      tsk->snd_una) {
      tcp_set_state(tsk, TCP_TIME_WAIT);
      tcp_set_timewait_timer(tsk);
      tcp_send_control_packet(tsk, TCP_ACK);
   else if ((cb->flags & TCP_ACK) && tsk->snd_nxt == tsk->snd_una)
      tcp_set_state(tsk, TCP_FIN_WAIT_2);
   else if (cb->flags & TCP_FIN) {
```

```
tcp_set_state(tsk, TCP_CLOSING);
      tcp_send_control_packet(tsk, TCP_ACK);
   }
  break;
case TCP LAST ACK:
if (!is_tcp_seq_valid(tsk, cb))
   return;
tsk->rcv_nxt = cb->seq_end;
if (cb->flags & TCP_ACK) {
  tcp_update_window_safe(tsk, cb);
  tsk->snd_una = cb->ack;
}
if ((cb->flags & TCP_ACK) && tsk->snd_nxt == tsk->snd_una) {
  tcp update send buffer(tsk, cb->ack);
  tcp_unset_retrans_timer(tsk);
  tcp_set_state(tsk, TCP_CLOSED);
  tsk->rcv_nxt = cb->seq;
  tsk->snd_una = cb->ack;
  tcp_set_state(tsk, TCP_CLOSED);
  tcp_unhash(tsk);
  tcp_bind_unhash(tsk);
  free_tcp_sock(tsk);
}
break;
```

还有 ESTABLISH 状态, 先判断收到的序列号是否是预期收到之前的, 若是则丢弃。然后判断数据包是否带数据, 若带数据则交给处理函数进行处理: 先判断数据包长度是否合法, 再看缓冲区是否有足够空间, 满则等待, 否则将数据包加入接收队列, 扫描接收队列, 若有序列号与 rcv_nxt 匹配的数据包则移动到环形缓冲区, 更新接收窗口、发送队列、定时器; 若不带数据则验证收到的序列号与期望的是否一致, 若收到的更新,则重置定时器, 再根据 ACK 更新发送队列和定时器:

```
case TCP_ESTABLISHED:
   if (less_than_32b(cb->seq, tsk->rcv_nxt)) {
     tcp_send_control_packet(tsk, TCP_ACK);
     return;
}
```

```
if (!is_tcp_seq_valid(tsk, cb))
      return;
   if (cb->flags & TCP_ACK) {
      tcp_update_window_safe(tsk, cb);
      tsk->snd_una = greater_than_32b(cb->ack, tsk->snd_una) ? cb->ack :
         tsk->snd_una;;
   }
   if (cb->flags & TCP_FIN) {
      tcp_update_send_buffer(tsk, cb->ack);
      tcp_update_retrans_timer(tsk);
      if (tsk->retrans_timer.enable)
         log(ERROR, "still have no ack packet before close wait\n");
      tcp_set_state(tsk, TCP_CLOSE_WAIT);
      handle_tcp_recv_data(tsk, cb);
      tsk->rcv_nxt = cb->seq_end;
      tcp_send_control_packet(tsk, TCP_ACK);
      wake_up(tsk->wait_recv);
   }
   else {
      if (cb->pl_len != 0)
         handle_tcp_recv_data(tsk, cb);
      else{
         tsk->rcv_nxt = cb->seq_end;
         if (cb->ack > tsk->snd_una) {
            tsk->retrans_timer.retrans_time = 0;
            tsk->retrans_timer.timeout = TCP_RETRANS_INTERVAL_INITIAL;
         }
         tsk->snd_una = cb->ack;
         tcp_update_window_safe(tsk, cb);
         tcp_update_send_buffer(tsk, cb->ack);
         tcp_update_retrans_timer(tsk);
      }
   }
   break;
// handle the recv data from TCP packet
int handle_tcp_recv_data(struct tcp_sock *tsk, struct tcp_cb * cb) {
   if (cb->pl_len <= 0)</pre>
      return 0;
   pthread_mutex_lock(&tsk->rcv_buf_lock);
   while (ring_buffer_full(tsk->rcv_buf)) {
      pthread_mutex_unlock(&tsk->rcv_buf_lock);
```

```
sleep_on(tsk->wait_recv);
}
tcp_recv_ofo_buffer_add_packet(tsk, cb);
pthread_mutex_unlock(&tsk->rcv_buf_lock);

tcp_move_recv_ofo_buffer(tsk);
pthread_mutex_lock(&tsk->rcv_buf_lock);

tsk->rcv_wnd = ring_buffer_free(tsk->rcv_buf);
tcp_update_send_buffer(tsk, cb->ack);
tcp_update_retrans_timer(tsk);
tcp_send_control_packet(tsk, TCP_ACK);
wake_up(tsk->wait_recv);
pthread_mutex_unlock(&tsk->rcv_buf_lock);
return 1;
}
```

14.2.2 拥塞控制

14.3 实验结果

14.3.1 丢包恢复

1. 本实验 server 与本实验 client 进行数据传输,用 md5sum 比较两个文件是否完全相同,结果如下:

```
Routing table of 1 entries has been loaded.

DEBUG: open file server-output.dat

DEBUG: alloc a new top sock, ref_cnt = 1

DEBUG: listening port 10001.

DEBUG: listening port 10001.

DEBUG: listen to port 10001.

DEBUG: listen to port 10001.

DEBUG: listen to port 10001.

DEBUG: alloc a new top sock, ref_cnt = 1

DEBUG: 10.0.0.1;10001 switch state, from CLOSED to SYN_RECV.

DEBUG: Pass 10.0.0.1;10001 <-> 10.0.0.2;12345 from process to listen_queue

DEBUG: 10.0.0.1;10001 switch state, from SYN_RECV to ESTABLISHED.

DEBUG: accept a connection.

DEBUG: 10.0.0.1;10001 switch state, from ESTABLISHED to CLOSE_WAIT.

DEBUG: peer closed.

used time; 10 s

DEBUG: close this connection.

DEBUG: close sock 10.0.0.1;10001 <-> 10.0.0.2;12345, state CLOSE_WAIT

DEBUG: 10.0.0.1;10001 switch state, from CLOSE_WAIT to LAST_ACK.

DEBUG: 10.0.0.1;10001 switch state, from LAST_ACK to CLOSED.

DEBUG: 10.0.0.1;10001 switch state, from CLOSED to CLOSED.

DEBUG: Free 10.0.0.1;10001 <-> 10.0.0.2;12345.
```

图 14.1. my server

```
DEBUG: sent 3944728 Bytes
DEBUG: retrans time: 1

DEBUG: retrans seq: 3911981

DEBUG: sent 3964752 Bytes
DEBUG: sent 3984776 Bytes
DEBUG: sent 4004800 Bytes
DEBUG: sent 4024824 Bytes
DEBUG: sent 4024824 Bytes
DEBUG: retrans time: 1

DEBUG: retrans seq: 3994997

DEBUG: sent 4044848 Bytes
DEBUG: sent 4052632 Bytes
DEBUG: sent 4052632 Bytes
DEBUG: the file has been sent completely.
DEBUG: close sock 10.0.0.2:12345 <-> 10.0.0.1:10001, state ESTABLISHED
DEBUG: 10.0.0.2:12345 switch state, from ESTABLISHED to FIN_WAIT-1.
DEBUG: 10.0.0.2:12345 switch state, from FIN_WAIT-1 to FIN_WAIT-2.
DEBUG: insert 10.0.0.2:12345 <-> 10.0.0.1:10001 to timewait, ref_cnt += 1
DEBUG: 10.0.0.2:12345 switch state, from TIME_WAIT to CLOSED.
DEBUG: Free 10.0.0.2:12345 <-> 10.0.0.1:10001.
```

图 14.2. my client

```
> md5sum client-input.dat server-output.dat c2133d46b8363075aefc0318786630a4 client-input.dat c2133d46b8363075aefc0318786630a4 server-output.dat
```

图 14.3. md5sum

可以看出,出现了丢包,但是通过重传机制,最终文件完全相同。

2. 本实验 server 与标准 client 进行数据传输,用 md5sum 比较两个文件是否完全相同,结果如下:

```
DEBUG: find the following interfaces: h1-eth0.
Routing table of 1 entries has been loaded.

DEBUG: open file server-output.dat

DEBUG: alloc a new top sock, ref_cnt = 1

DEBUG: listening port 10001.

DEBUG: listen to port 10001.

DEBUG: listen to port 10001.

DEBUG: alloc a new top sock, ref_cnt = 1

DEBUG: listen to port 10001 switch state, from CLOSED to SYN_RECV.

DEBUG: pass 10.0.0.1:10001 switch state, from CLOSED to SYN_RECV.

DEBUG: pass 10.0.0.1:10001 <-> 10.0.0.2:54428 from process to listen_queue

DEBUG: 10.0.0.1:10001 switch state, from SYN_RECV to ESTABLISHED.

DEBUG: accept a connection.

DEBUG: peer closed.

used time: 14 s

DEBUG: close this connection.

DEBUG: close this connection.

DEBUG: 10.0.0.1:10001 switch state, from CLOSE_WAIT to LAST_ACK.

DEBUG: 10.0.0.1:10001 switch state, from LAST_ACK to CLOSED.

DEBUG: 10.0.0.1:10001 switch state, from CLOSED to CLOSED.

DEBUG: Free 10.0.0.1:10001 switch state, from CLOSED to CLOSED.
```

图 14.4. my server

```
send:2000000, remain:2052632, total: 2000000/40526
send:2100000, remain:1952632, total: 2100000/40526
send:2200000, remain:1852632, total: 2200000/40526
send:2300000, remain:1752632, total: 2300000/4052632
send:2400000, remain:1652632, total: 2400000/40526
send:2500000, remain:1552632, total: 2500000/405
send:2600000, remain:1452632, total: 2600000/405
send:2700000, remain:1352632, total: 2700000/40526
send:2800000, remain:1252632, total: 2800000/40526
send:2900000, remain:1152632, total: 2900000/4052
send:3000000, remain:1052632, total: 3000000/405263
send:3100000, remain:952632, total: 3100000/405263
send:3200000, remain:852632, total: 3200000/405263
send:3300000, remain:752632, total: 3300000/4052632
send:3400000, remain:652632, total: 3400000/405263%
send:3500000, remain:552632, total: 3500000/4052632
send:3600000, remain:452632, total: 3600000/4052632
send:3700000, remain:352632, total: 3700000/4052632
send:3800000, remain:252632, total: 3800000/4052632
send:3900000, remain:152632, total: 3900000/405
send:4000000. remain:52632. total: 4000000/4052632
```

图 14.5. std client

```
> md5sum client-input.dat server-output.dat
4bb8a034cd44c20105bd3252d9c6495f client-input.dat
4bb8a034cd44c20105bd3252d9c6495f server-output.dat
```

图 14.6. md5sum

可以看出,出现了丢包,但是通过重传机制,最终文件完全相同。

3. 标准 server 与本实验 client 进行数据传输,用 md5sum 比较两个文件是否完全相同,结果如下:

```
('10,0,0,2', 12345)
```

图 14.7. std server

```
DEBUG: sent 4004800 Bytes
DEBUG: sent 4024824 Bytes
DEBUG: sent 4044848 Bytes
DEBUG: retrans time: 1

DEBUG: retrans seq: 4004801

DEBUG: sent 4052632 Bytes
DEBUG: retrans time: 1

DEBUG: retrans seq: 4030665

DEBUG: retrans time: 2

DEBUG: retrans seq: 4042345

DEBUG: retrans seq: 4042345

DEBUG: the file has been sent completely.
DEBUG: close sock 10.0.0.2:12345 <-> 10.0.0.1:10001, state ESTABLISHED
DEBUG: 10.0.0.2:12345 switch state, from ESTABLISHED to FIN_WAIT-1.
DEBUG: 10.0.0.2:12345 switch state, from FIN_WAIT-1 to TIME_WAIT.
DEBUG: insert 10.0.0.2:12345 <-> 10.0.0.1:10001 to timewait, ref_cnt += 1
DEBUG: free 10.0.0.2:12345 switch state, from TIME_WAIT to CLOSED.
DEBUG: Free 10.0.0.2:12345 <-> 10.0.0.1:10001.
```

图 14.8. my_client

```
> md5sum client-input.dat server-output.dat f1a72ea92a9d8735d3a20debdface99e client-input.dat f1a72ea92a9d8735d3a20debdface99e server-output.dat
```

图 14.9. md5sum

可以看出,出现了丢包,但是通过重传机制,最终文件完全相同。

14.3.2 拥塞控制

14.4 实验总结

本次实验中,我们实现了丢包恢复和拥塞控制的功能。通过实验,我们再次深入了解了 TCP 协议栈的工作原理。实现可靠传输使得数据传输更加稳定,实现拥塞控制使得网络更加稳定。通过本次实验,我们对 TCP 协议栈的实现有了更深入的了解,更加接近真实的 TCP 协议栈。