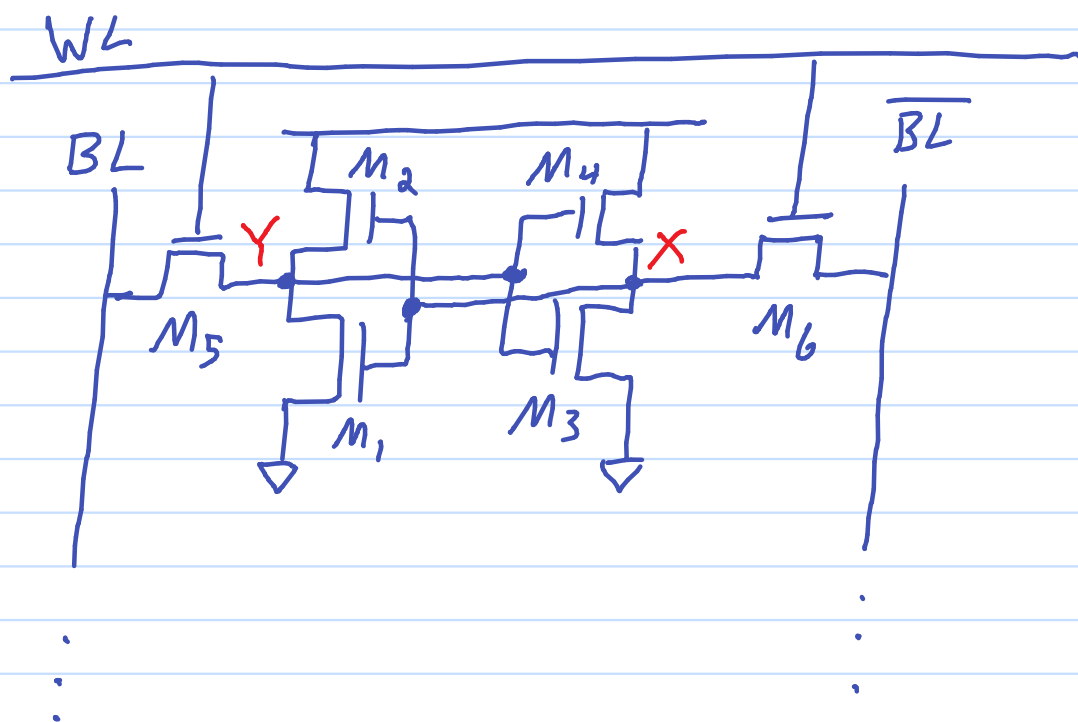


# EECS 2513 HW3 solutions

①



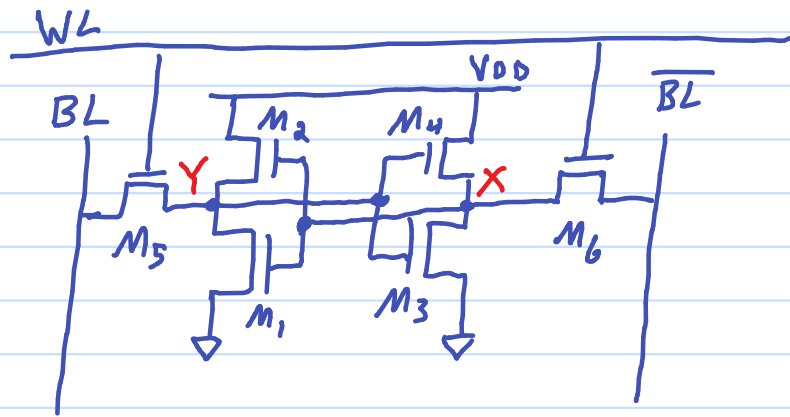
a)  $V_{WL} \downarrow \Rightarrow M_5 \text{ \& } M_6$  effective resistance increases relative to  $M_1 \text{ \& } M_3$ . This increases read stability by lowering the voltage seen across nodes X/Y during a read.

b)  $V_{WL} \downarrow \Rightarrow$  Assuming  $M_1$  and  $M_3$  are not limiting the current from the bitline, then decreasing the wordline voltage will increase the read access time

c)  $V_{WL} \downarrow \Rightarrow M_5 \text{ \& } M_6$   $R_{eff}$  increases relative to  $M_2$  and  $M_4$ . This decreases write stability by increasing the voltage at nodes X/Y during a write.

① cont...

d)



$\uparrow V_{DD} \Rightarrow$  Increased cell supply voltage will decrease  $R_{eff}$  of  $M_2 \parallel M_4$  relative to the access transistors  $M_5 \parallel M_6$ . This decreases write stability as the voltage drop will be larger across the access transistors  $M_5 \parallel M_6$  during a write.

② 2MB

128 x 128 bit arrays

4 columns / sense amp

$$C_{BL} = 1 \text{ pF}$$

$$V_{DD} = 1V$$

$$I_{on} = 100 \text{ mA}$$

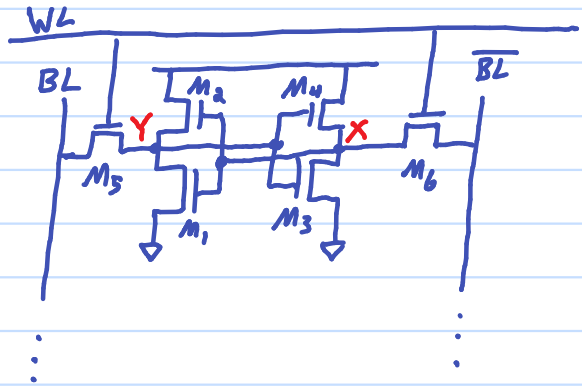
$$I_{OFF} = 400 \text{ nA}$$

$$t_{\text{read}} = 1 \text{ ns}$$

$$M_{offset} = 0 \text{ V}$$

$$\sigma_{offset} = 15 \text{ mV}$$

← sense and offset normally distributed



The  $I_{on}$  and  $I_{off}$  currents are the currents from the bitline through the access transistor and through either  $M_1$  or  $M_3$  during a read (depends on whether node  $X$  or  $Y$  is at  $0V$ ). We want to know the minimum  $\Delta V_{BL}$  and make sure it is greater than the sense amp's offset voltage. This is what determines whether or not the sense amp "yields".

2. cont...

First step is to calculate the minimum (i.e. worst case)  $\Delta V_{BL}$ . This is seen when all cells in the column store 0's while the cell being read stores a 1 (or vice versa).

The  $\Delta V_{BL}$  in this case becomes

*change in BL* *change in  $\overline{BL}$*

*All cells in column except 1*

$$\Delta V_{BL} = \Delta V_{BL-1} - \Delta V_{BL-2} = \frac{I_{on} \cdot t_{read}}{C_{BL}} - \frac{127 \cdot I_{off} \cdot t_{read}}{C_{BL}}$$

Plugging in we get

$$\Delta V_{BL} = \frac{(100 \mu A)(1 \text{ ns})}{1 \text{ pF}} - \frac{(127)(400 \text{ nA})(1 \text{ ns})}{1 \text{ pF}}$$

$$\Delta V_{BL} = 100 \text{ mV} - 51 \text{ mV} = 49 \text{ mV}$$

$$\Delta V_{BL} = 49 \text{ mV}$$

Now calculate the number of columns in the SRAM array.

*MiB =  $2^{20}$  bytes*

$$\#_{\text{cols}} = \frac{\text{Size of Mem}}{\text{Column Length}} = \frac{2 \text{ MiB}}{1286} = 131,072$$

Now get the # of sense amps

$$\#_{SA} = \frac{\#_{\text{cols}}}{4} = \frac{131,072}{4} = 32,768$$

② cont...

Now let's get the probability that a single sense amp will yield. This is done by first computing the z-score of our minimum  $\Delta V_{BL}$

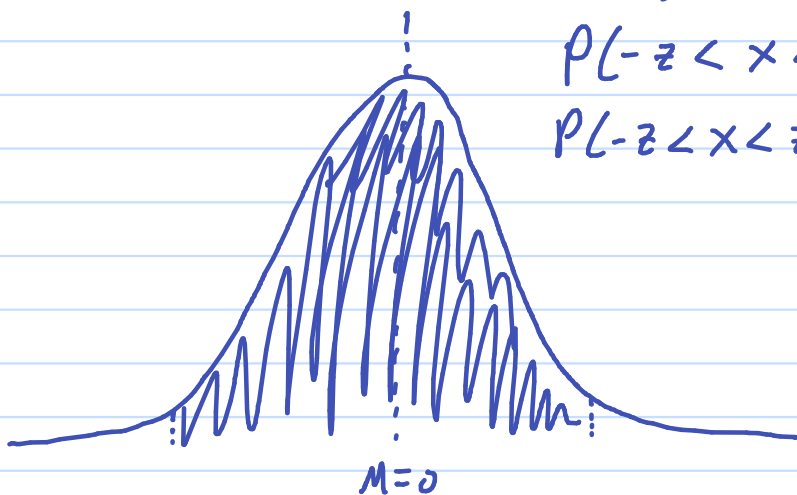
$$z = \frac{\Delta V_{BL} - \mu}{\sigma} = \frac{49 \text{ nV}}{15 \text{ nV}} = 3.27$$

Now we need to get the two sided cumulative probability when using a z-score table because an offset of  $> 49 \text{ nV}$  or  $< -49 \text{ nV}$  will render the sense amp useless in this worst-case scenario.

$$z = 3.27$$

$$P(-z < x < z) = 1 - 2(0.0005)$$

$$P(-z < x < z) = 0.999$$



Thus,  $P_{SA} = 0.999$

Without redundancy our array would have a yield of

$$\text{yield} = (P_{SA})^{\#SA_s} = (0.999)^{32,768} = 5.8e-15$$

This is crazy low... Now we'll use redundancy!

② cont...

Redundant columns can only be used to replace defective columns in their own sub-array. Let's calculate the yield needed for each sub-array to get  $\geq 99\%$  yield in the entire SRAM

$$\#_{\text{sub-arrays}} = \frac{\text{Size of Mem}}{\text{Size of sub}} = \frac{2\text{MiB}}{128 \times 128} = 1024$$

$$P_{\text{sub}} = (0.99)^{\frac{1}{1024}} = 0.99999$$

Now we need to figure out how many additional sense amps we need to get our sub array to have this yield.

$$\#_{\text{SAs/sub}} = \frac{\#_{\text{cols}}}{4} = \frac{128}{4} = 32 \text{ SAs/sub-array}$$

0 additional

$$P_{\text{sub-work}} = P_{\text{SA}}^{32} = 0.968$$

Too low

All 32 work

Only 1 SA is broken



1 additional

$$P_{\text{sub-work}} = P_{33\text{-work}} + P_{32\text{-work}} = P_{\text{SA}}^{33} + \binom{33}{32} P_{\text{SA}}^{32} (1 - P_{\text{SA}})$$

$$P_{\text{sub-work}} = 0.9995$$

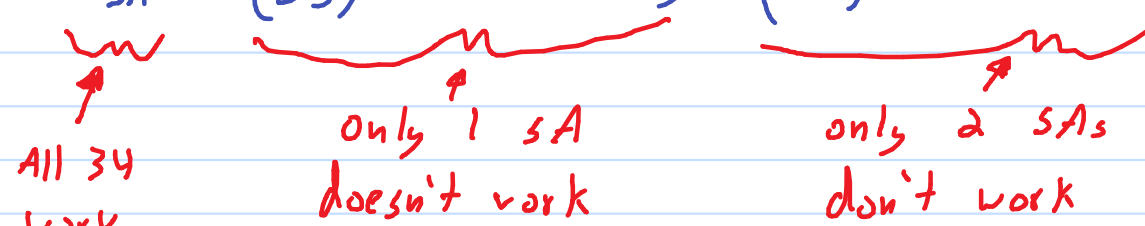
still too low

② cont...

2 additional sense amps

$$P_{\text{sub-work}} = P_{34\text{-work}} + P_{33\text{-work}} + P_{32\text{-work}}$$

$$P_{\text{sub-work}} = P_{SA}^{34} + \binom{34}{33} P_{SA}^{33} (1 - P_{SA}) + \binom{34}{32} P_{SA}^{32} (1 - P_{SA})^2$$



All 34 work      only 1 SA doesn't work      only 2 SAs don't work

$$P_{\text{sub-work}} = 0.999994$$

This is large enough!

We have found that adding 2 additional sense amps gives us a sub-array yield greater than what is needed to get the entire array's yield to 99%.

2 sense amps means we need  $2 \times 4 = 8$  additional columns per sub-array

8 redundant columns  
per sub-array