

CS118:

Computer Network Fundamentals

Lecture-1: introduction



CS118: explains (roughly) **how the Internet works**

- ◆ Internet: a huge, complex network of networks
- ◆ Divide-and-conquer
 - Figure out how many major parts,
 - Learn one piece at a time
- ◆ Your job:
 - Read textbook, think, collect a list of questions
 - review every lecture slide deck after each class
 - Ask questions in class/office hours/via Piazza
 - Practice what you learn through homework and projects

James F. Kurose | Keith W. Ross

COMPUTER NETWORKING

A TOP-DOWN APPROACH

Eighth Edition



Brief Contents

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Course assignment and due schedule

Midterm	In-class, Wednesday Feb 5 (Location TBD)
Final	3:00PM-6:00PM Saturday March 21 (Location TBD)
Homework	Release: on Thursday of week 1, 3, 5, 7; Due: 11:59pm Tuesday of week 3, 5, 7, 9.
Project 0	Release: Monday Jan 6, 2024 (Week 1) Due: 11:59pm Wednesday, Jan 15, 2024 (Week 2) 1.5 weeks
Project 1	Release: Thursday Jan 16, 2024 (Week 2) Due: 11:59pm Sunday, Feb 16, 2024 (Week 6) 4.5 weeks Grading: auto-grading script (sample tests will be provided to let everyone test their code before submission)
Project 2	Release: Monday, Feb 17, 2024 (Week 7) Due: 11:59pm Sunday, March 16, 2024 (Week 10) 4 weeks Grading: auto-grading script (sample tests will be provided to let everyone test their code before submission)

FOR ALL OTHER COURSE INFO, PLEASE SEE [HTTPS://BRUINLEARN.UCLA.EDU/](https://bruinlearn.ucla.edu/)

Course workload and grading

- ◆ Bi-weekly homework assignments
- ◆ 3 programming projects,
 0. UDP socket (individual)
 1. Reliable data delivery (2-3 people team)
 2. Secured reliable data delivery (2-3 people team)
- ◆ Midterm and final exams (cheat sheets allowed, 2 pages double sided)
- ◆ **Strict Grading Policy**
 - Homework: do it yourself; no credit for late submission
 - Project: 20% credit reduction per late day
 - *No make-up exam*

Homework	20%
Programming Projects	25% (5/ 10/ 10)
Midterm	25%
Final exam	30%

2% extra credits based on piazza

1% extra credits course evaluation and TA/LA feedbacks

Class Policy

The following actions are **strictly prohibited**

- ◆ Posting/sharing/selling class material, with or without answers, to anyone outside this class, during or after this quarter.
- ◆ Use of old homework/midterm/finals in doing homework or exams
- ◆ Use ChatGPT in doing assignments
- ◆ Making your project code publicly available either during or *after* this quarter
 - *you must use private repository* on GitHub or GitLab

Hints for Getting Good Grade

- ◆ Review previous lecture slides
- ◆ Read textbook before coming to each lecture
- ◆ Ask *questions*
- ◆ Get your work done early
 - Lecture slides uploaded to BruinLearn *one day before the lecture*
 - Get HWs and projects done **before** the deadline

Let's get started

Today we cover the basic concepts in
Chapter 1 of the textbook

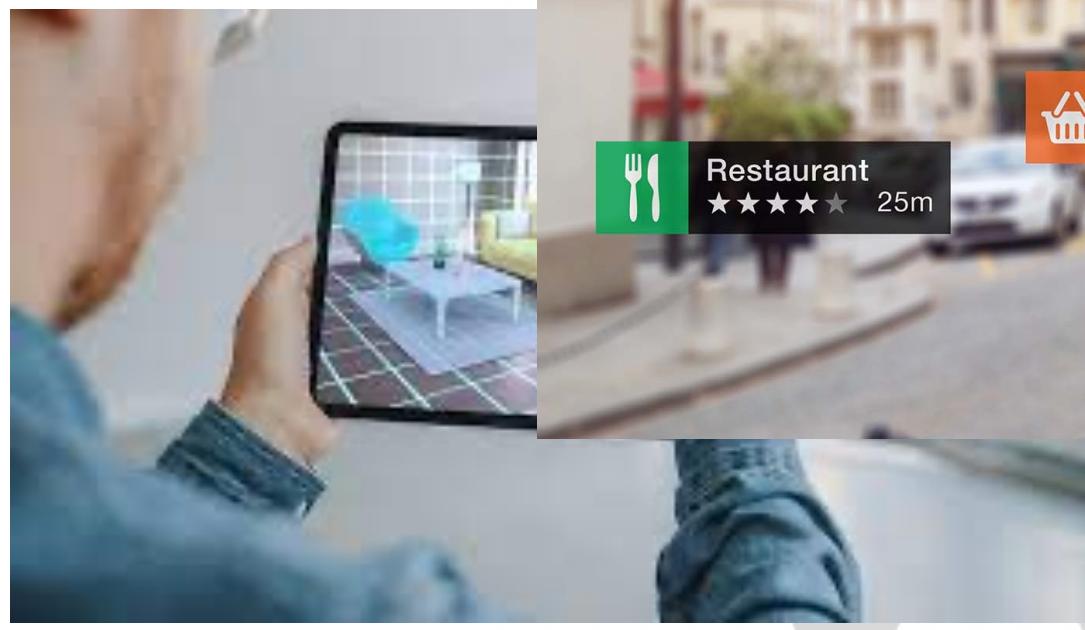
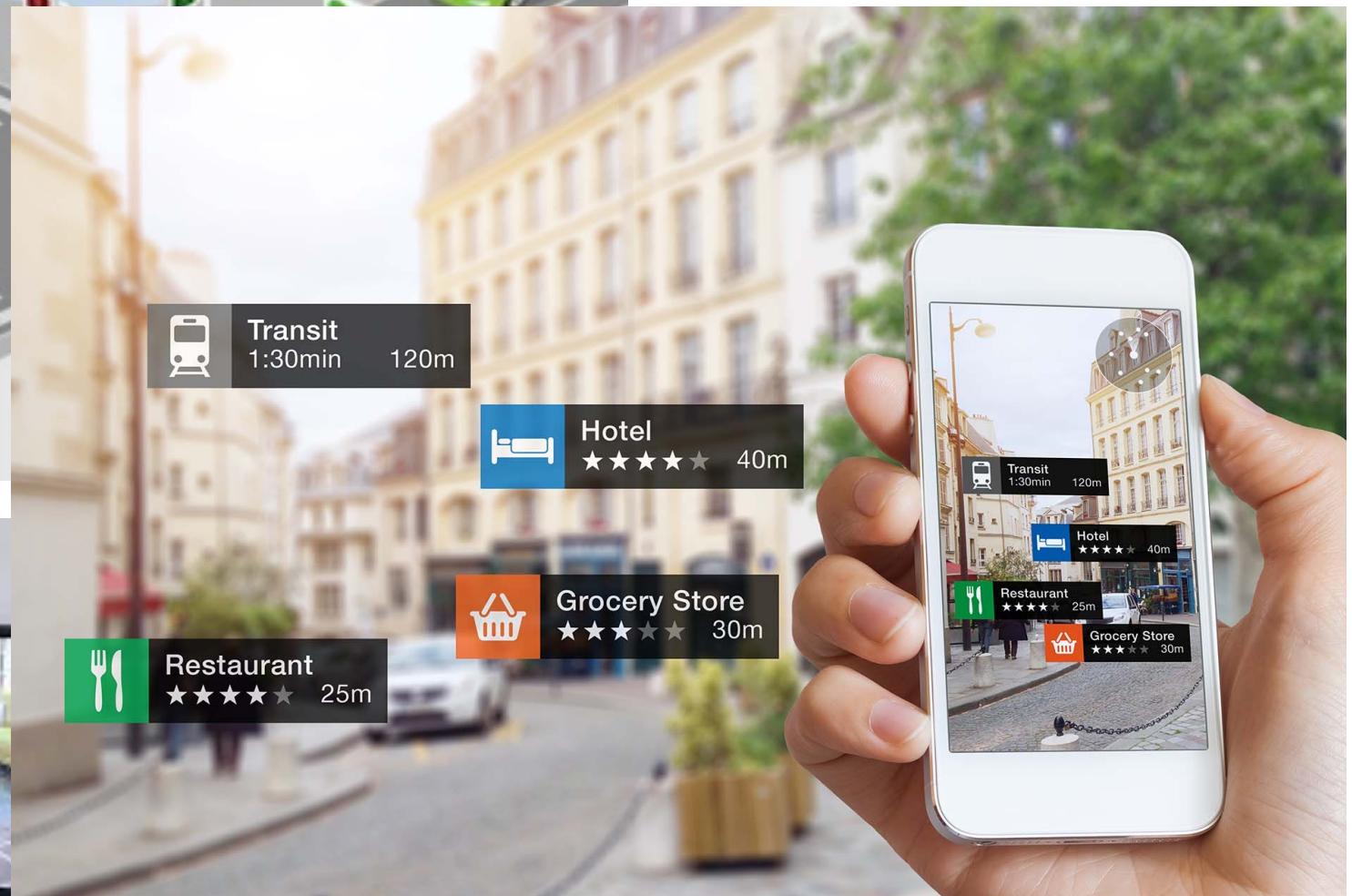
What is a Computer Network



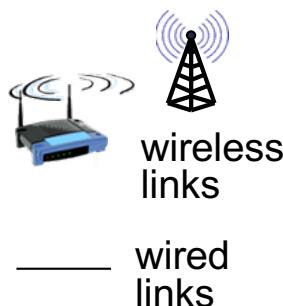
What is a Computer Network



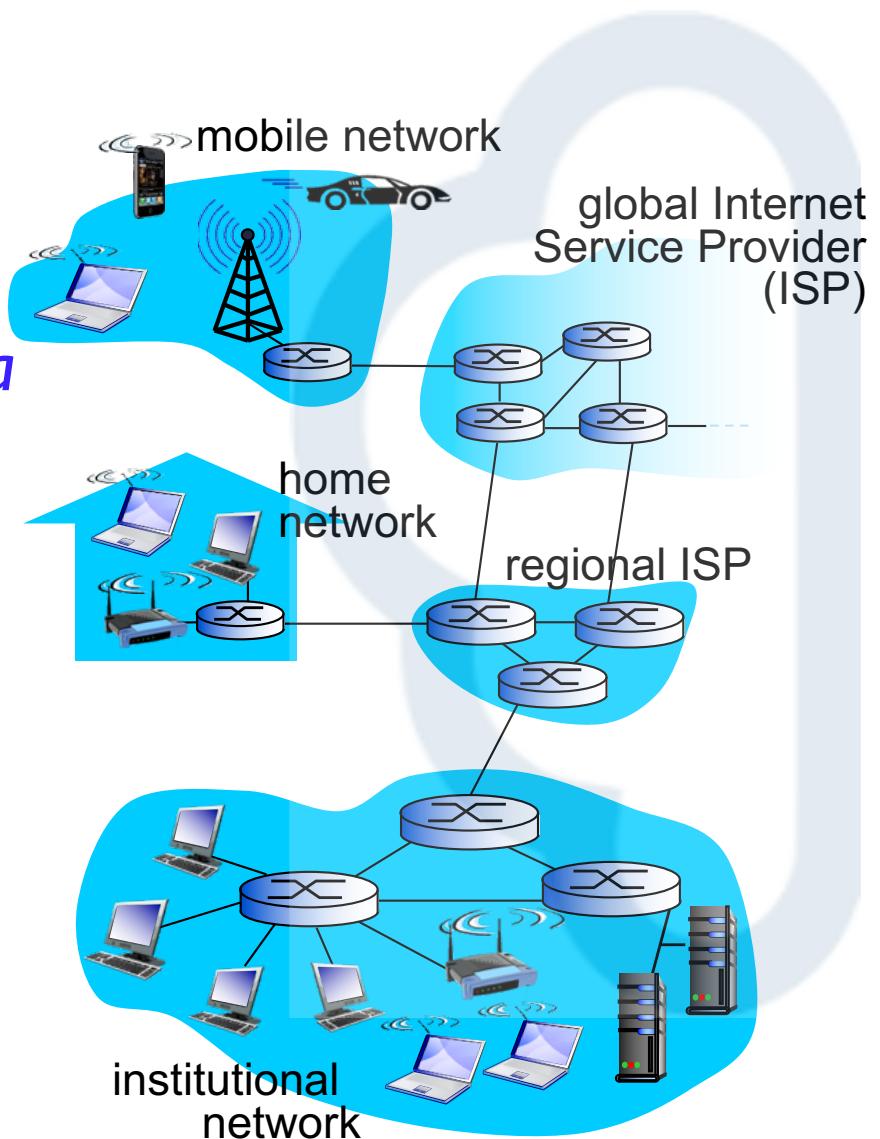
Coming to you



Terminology



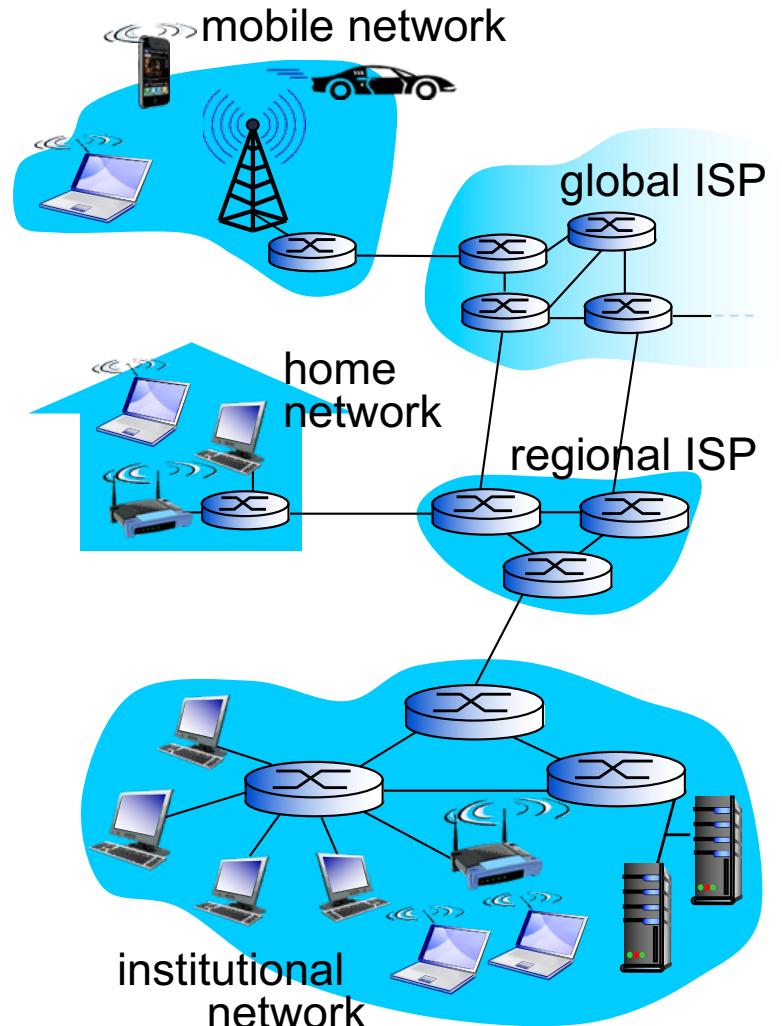
- billions of connected computing devices:
 - *hosts = end systems*
 - running *network apps*
 - Apps send/receive *data packets*
- *Routers = packet switches* inside network
- *communication links*
 - fiber, copper, radio, satellite
 - transmission rate = *bandwidth (BW)*



Recent years witnessed rapid growth of giant cloud service providers

“Nuts and Bolts”

- ◆ *Internet: “network of networks”*
 - Interconnected ISPs, enterprise networks, now also cloud service providers
- ◆ *Protocols: define how to send, receive packets*
 - e.g., HTTP, TCP, IP, 802.11
- ◆ *Internet protocol standards*
 - RFCs: “Request for Comments”
 - <https://www.rfc-editor.org/rfc-index.html>
 - Developed by Internet Engineering Task Force (IETF)
 - IEEE Standards
 - W3C (World Wide Web Consortium), and others



What is a protocol?

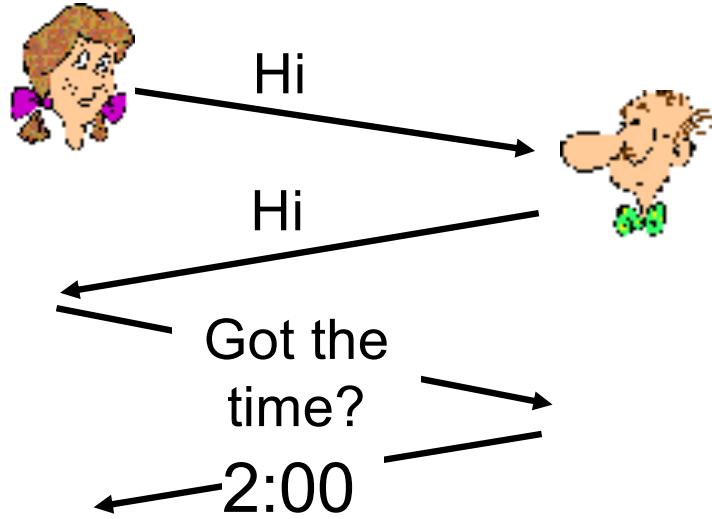
Traffic light protocol

- ◆ Green: go
- ◆ Red: stop
- ◆ Yellow: slow down - stop

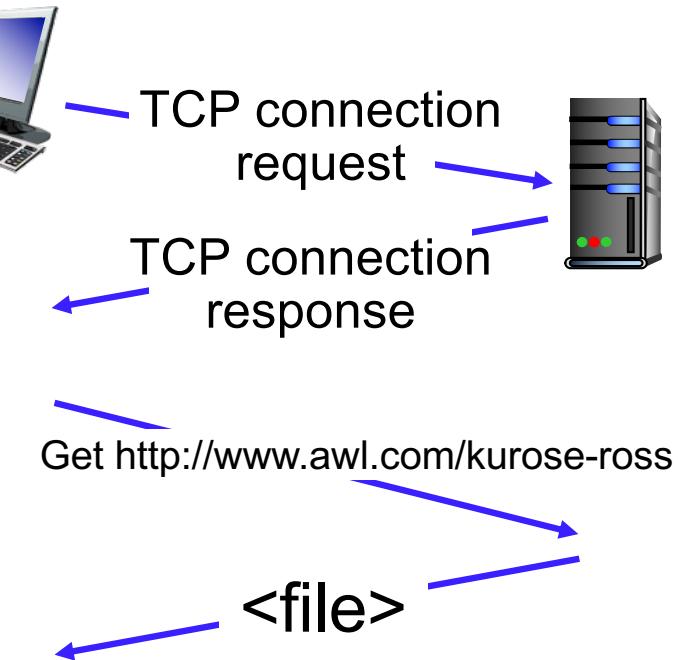
... specific messages sent

... specific actions taken when
the messages received

human protocols:

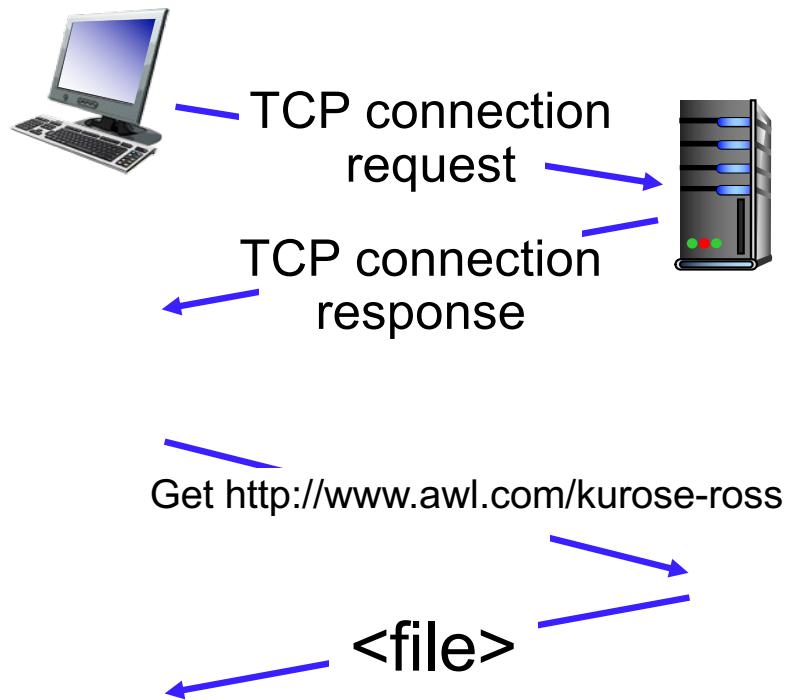


computer protocols:



Internet protocols

computer protocols:

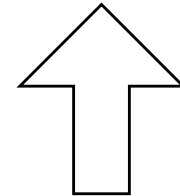


- Communication between machines rather than humans
- all communication activity governed by protocols

protocols define format, order of packets sent and received among network entities, and actions taken on packet transmission, receipt

Delivering data over the global Internet is a complicated process, involving many many steps

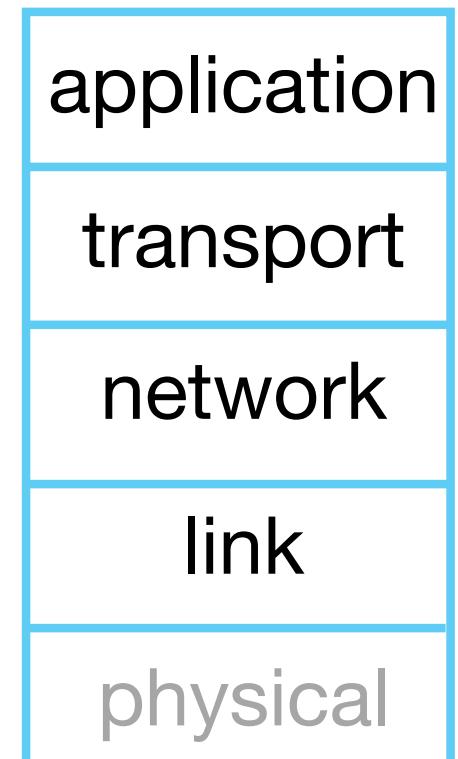
How to get the work done: divide and conquer
Group functions to a few modules



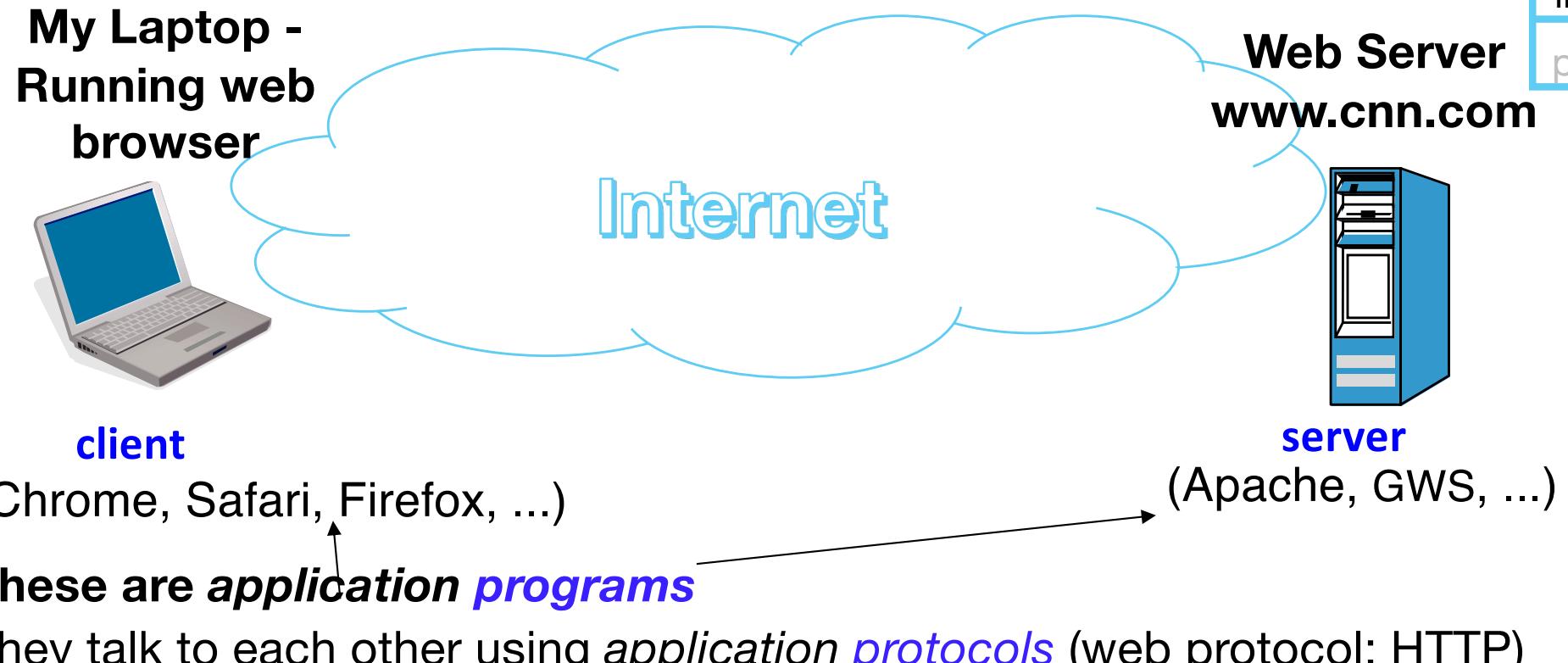
How many?

Internet protocol stack

- ◆ Application layer protocols
 - Support data exchange between application processes
 - Example: SMTP, HTTP, DNS
(Simple Mail Transfer Protocol)
- ◆ Transport layer protocols
 - handling delivery reliability, multiplex within a host
 - Example: TCP, UDP
- ◆ Network layer protocols
 - forward packets from source to destination
 - Example: IP
- ◆ Link layer protocols
 - transfer data between directly connected network elements
 - Example: Ethernet protocol, WiFi
- ◆ Physical layer: bits “on the wire”



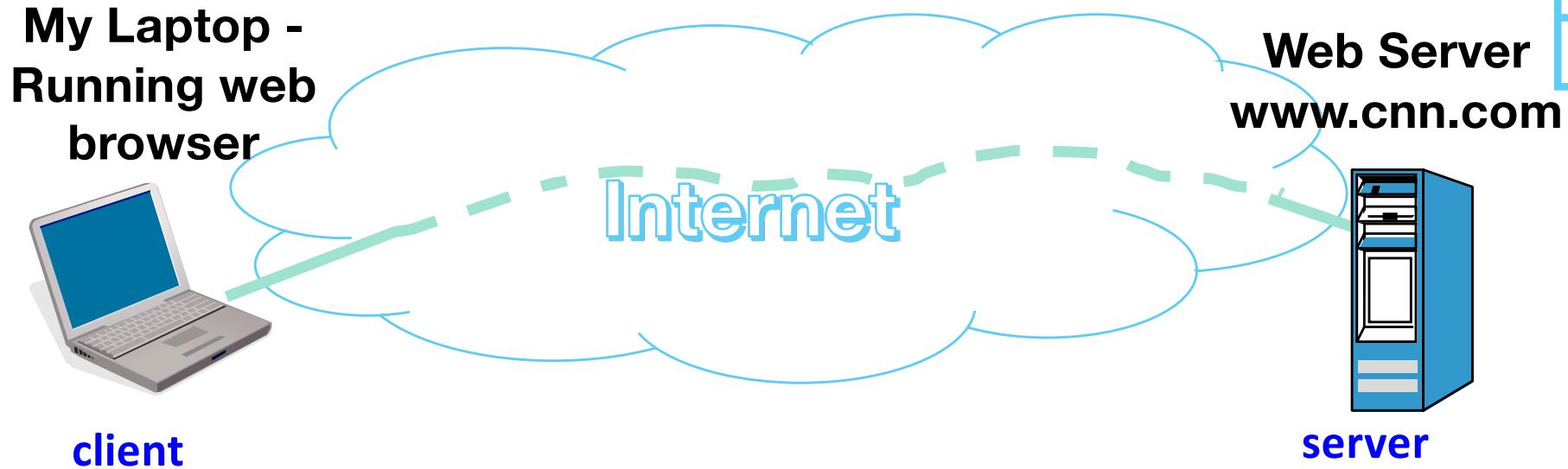
Application View



Application protocols

- Assume network can send data to any hosts on the Internet
- Don't know/care how data is sent, and assume all data delivered reliably
- Runs on top of a transport protocol

Transport View

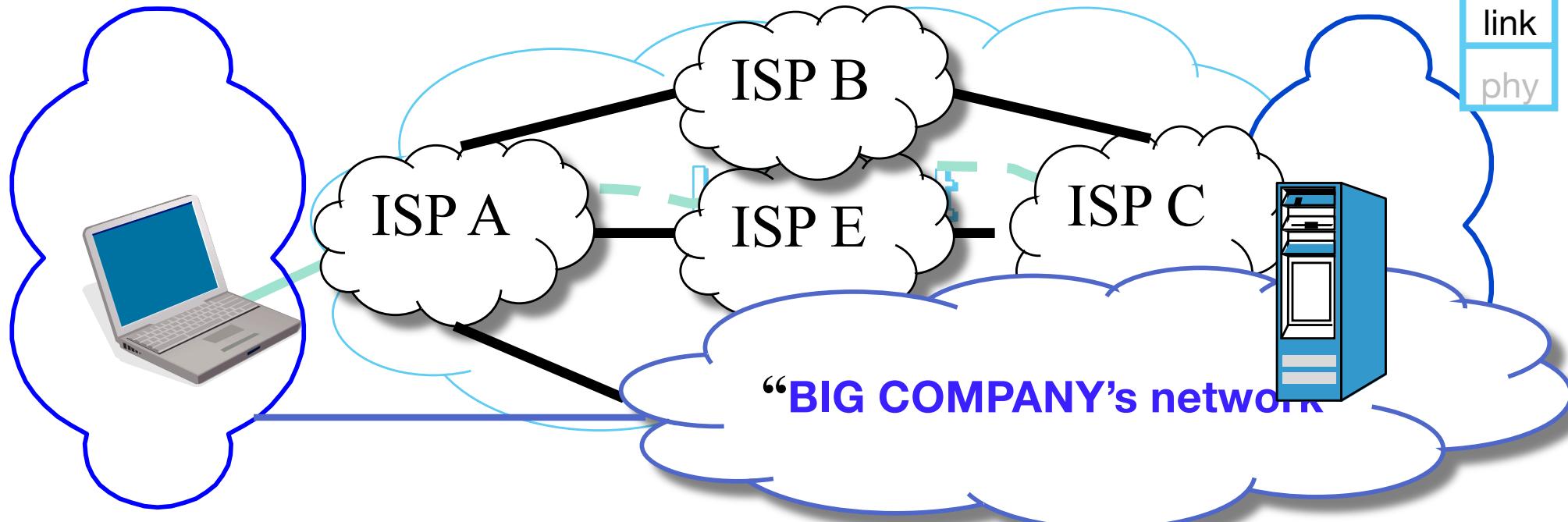


- ◆ **A transport protocol's job: delivering data between the two communicating ends**
 - *Don't know or care about which paths data may traverse through the network*
- ◆ Multiple transport protocols exist, each offers somewhat different functions (e.g. reliability, congestion control)

Actually, transport protocols don't do delivery → network protocol's job

Network Layer View

apps
trans
net
link
phy



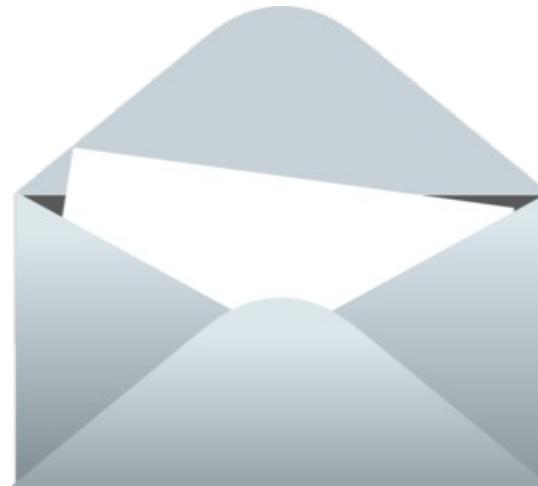
- ◆ **network protocol's job:** forward packets from source to destination host
- ◆ A really hard problem: the Internet is large, run by many different parties
 - connection from laptop to CNN.com:
WiFi → campus backbone → local ISP → other ISP → CNN website

Link Layer View



- ◆ **Link layer's job:** Get a packet transmitted across some communication medium to **next hop**
- ◆ Different medium → different link layer protocol

What protocol “layer” really means



application

transport

network

link

physical

Link
layer
protocol

Network
protocol

Transport
protocol

Application
protocol

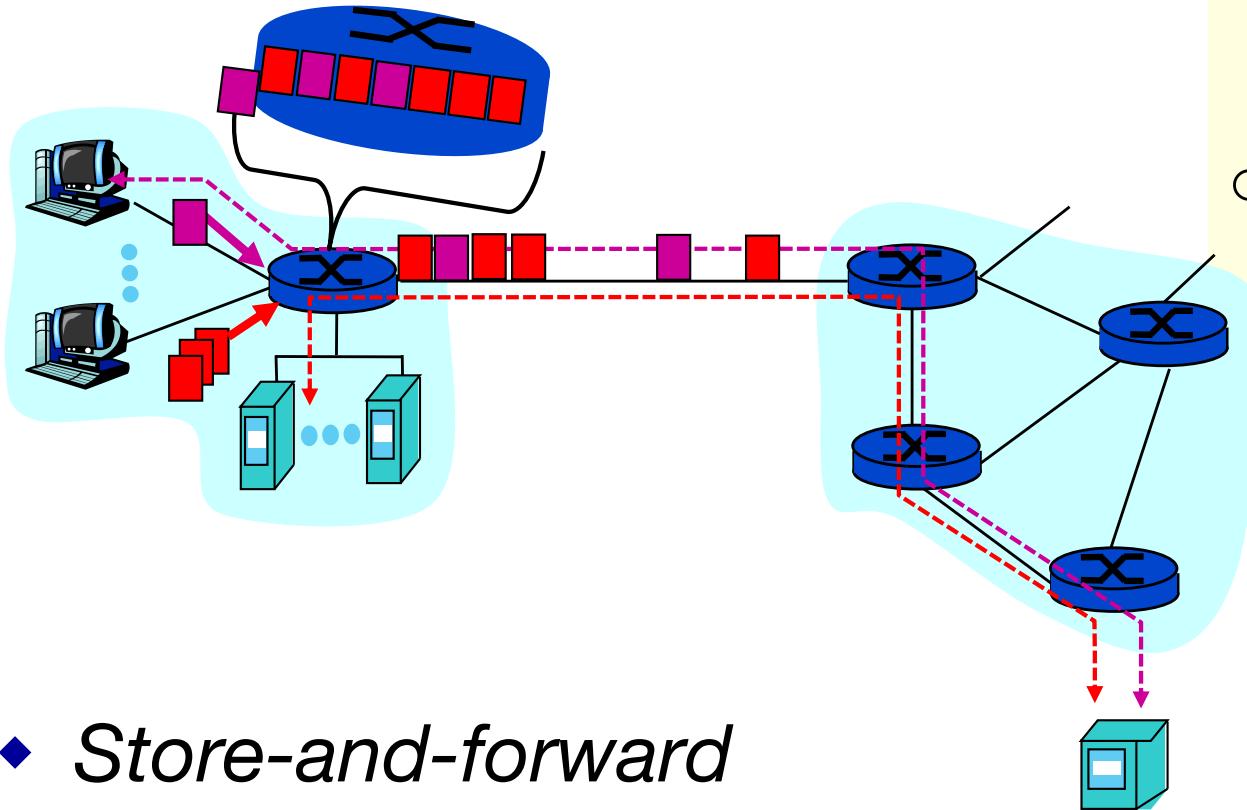
Application
data

(Tentative) Schedule of the Quarter

Week:	1	2	3	4	5
Mon	1/6 Course intro BW& delay	1/13 HTTP	1/20 Martin Luther King Jr. Day	1/27 Transport protocols	2/3 Congestion Control
Wed	1/8 Socket programming, Web & HTTP	1/15 DNS	1/22 DNS	1/29 TCP	2/5 Midterm
	6	7	8	9	10
Mon	2/10 Security 101	2/17 Presidents' Day	2/24 Routing algorithms & protocols	3/3 Routing in the Internet	3/10 Hubs and switches
Wed	2/12 Internet Protocol (IP)	2/19 Addressing, NAT, IPv6	2/26 Routing algorithms & protocols	3/5 Link layer (Ethernet)	3/12 Course review
					3/21: Final Exam

- The big yellow numbers indicate the chapter numbers in the textbook.

Packet Switching: Statistical Multiplexing



- Each node sends packets as soon as link available
- Receiver gets a full packet first, then forwards it towards the destination

- ◆ *Store-and-forward*
- ◆ Packet switch can temporarily buffer up packets
 - Introduce *delay*
 - Packets get *dropped* when the queue is full

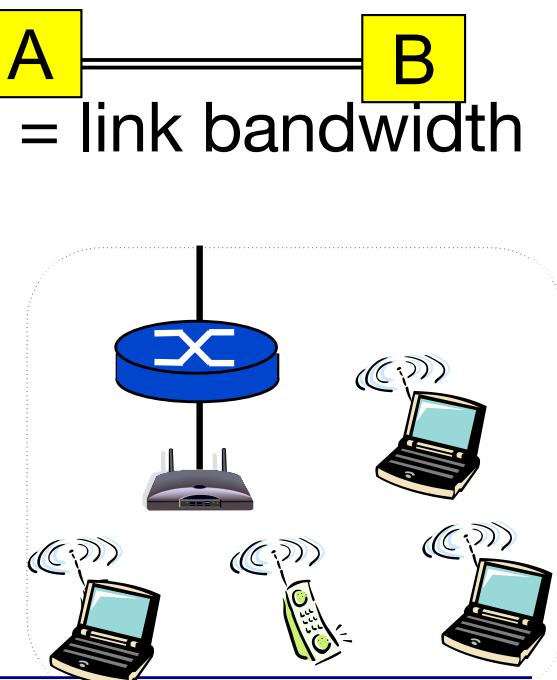
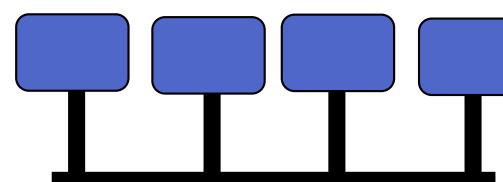
Network Performance

- ◆ 3 basic measurements

- Throughput (bits/sec, Kbps=1000 bits/sec, Mbps)
- Loss rate (% of packets lost)
- Delay (sec, msec)

Throughput

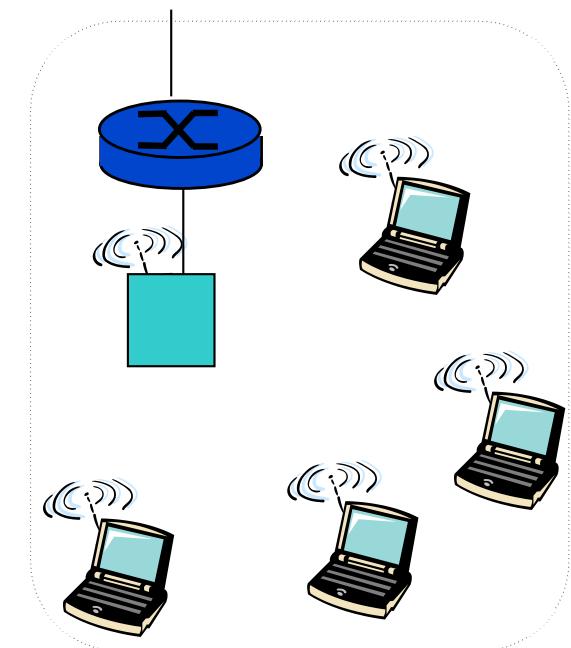
- over a single link: point-to-point
 - Pumping data into the pipe: throughput = link bandwidth
- Multi-access:
 - a lot more difficult to measure, *why?*



Packet Losses

- ◆ Wired links
 - Loss due to transmission errors
 - Loss due to congestion
- ◆ Wireless links
 - Limited transmission rate
 - Higher (than wire) bit error rate
 - Host mobility: high variance in the number of hosts sharing the same wireless channel

Do users know there are packet losses?
Do users' performance get affected by losses?



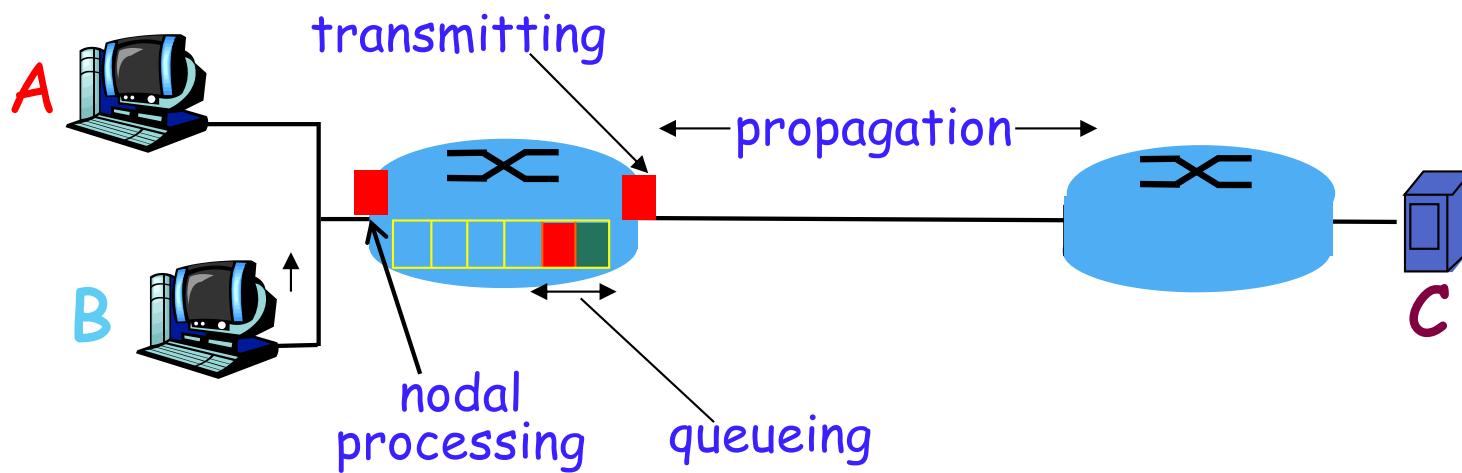
Delay in packet-switched networks

4 sources of delay at each hop

- ◆ node processing:
 - check bit errors
 - determine output link
- ◆ Queuing = #packets in queue X transmission time of each packet

Transmission = Length / rate
 R = link bandwidth (bps)
 L = packet length (bits)

Propagation = distance/sec
 d = length of physical link
 s = propagation speed in medium ($\sim 2 \times 10^8$ m/sec)

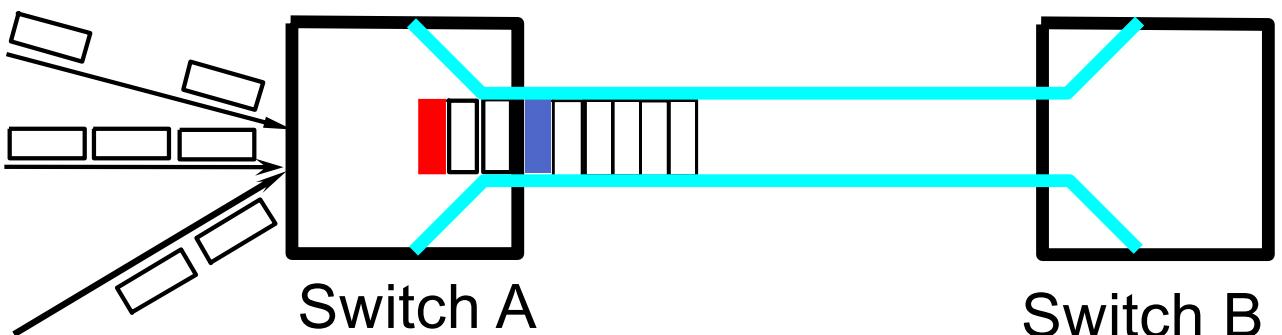


Example: calculating one hop delay

total delay (A → B) = ?

- ❖ Queuing delay = ?
- ❖ transmission delay = ?
- ❖ Propagation delay = ?

link length = 100 km
Bandwidth= 1 Mbps
packet size= 1000 bits
(all pkts equal length)



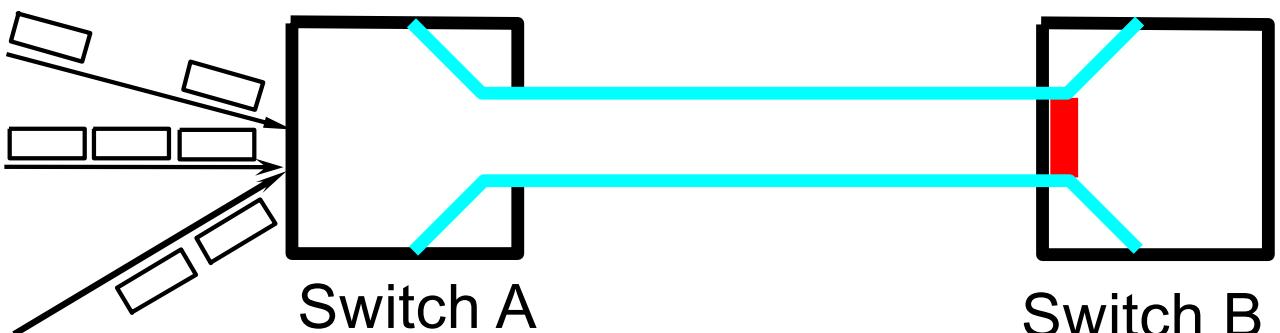
(2.0×10^8 meters/sec in a fiber)

Example: calculating one hop delay

total delay (A → B) = $1ms \times 2 + 1ms + 0.5ms = 3.5ms$

- ❖ Queuing delay = **Waiting time for 2 pkts**
- ❖ transmission delay = $\frac{1000bits}{1000000bits/sec} = 1 \text{ msec}$
- ❖ Propagation delay = $\frac{100,000m}{2 \times 10^8 m/sec} = 0.5 \text{ msec}$

link length = 100 km
Bandwidth= 1 Mbps
packet size= 1000 bits
(all pkts equal length)



(2.0×10^8 meters/sec in a fiber)

Transmission vs. propagation delay

Transmission delay: L / R

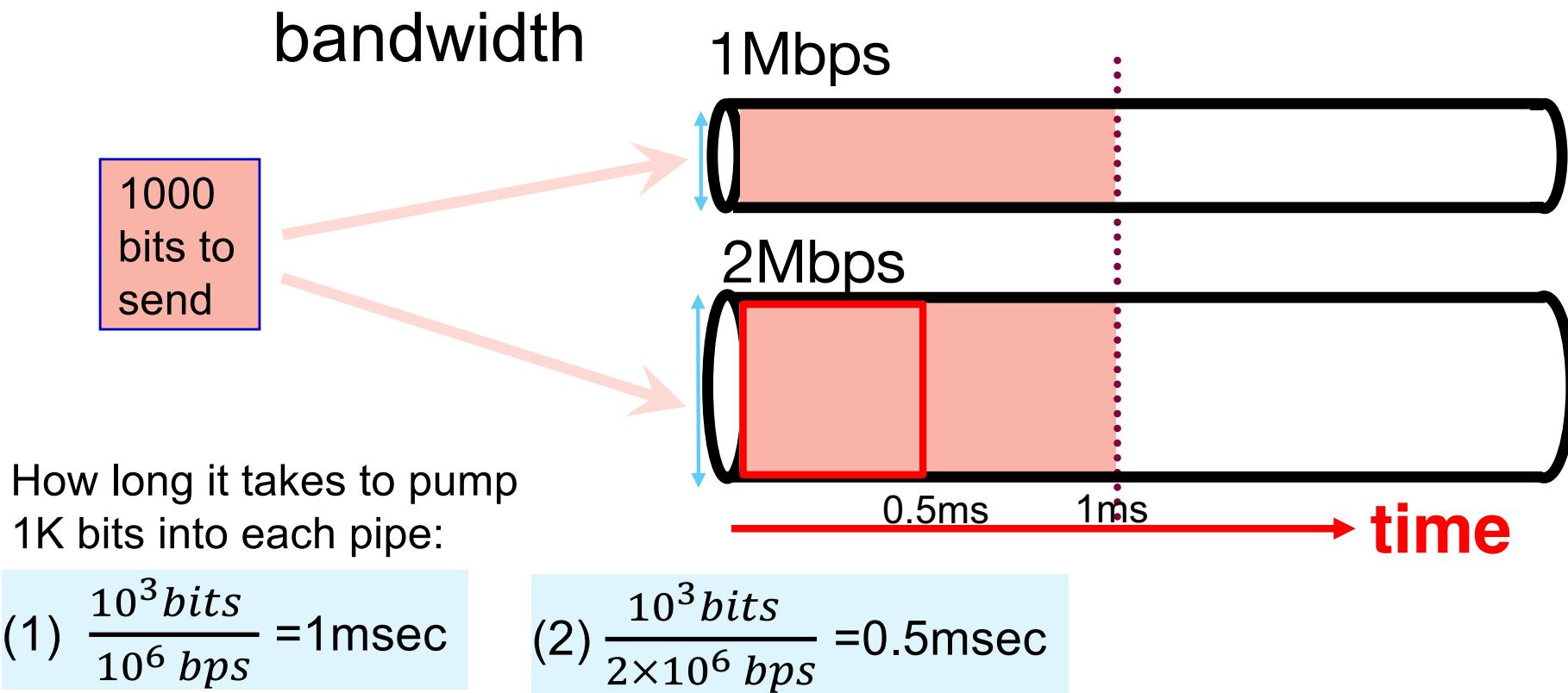
R = link bandwidth (bit-per-second, bps)

L = packet length (bits)

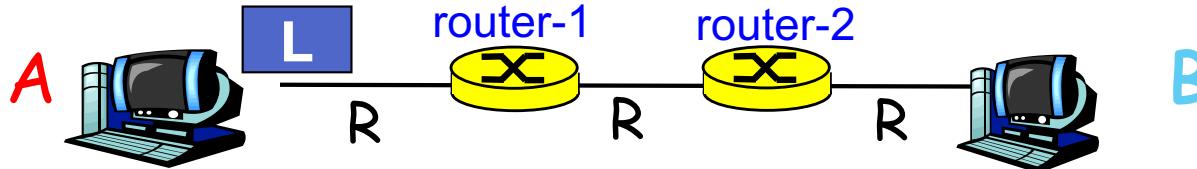
Propagation: d / s

d = length of a physical link

s = signal's propagation speed in the medium ($\sim 2 \times 10^8$ meter/sec)



Packet-switching: store-and-forward

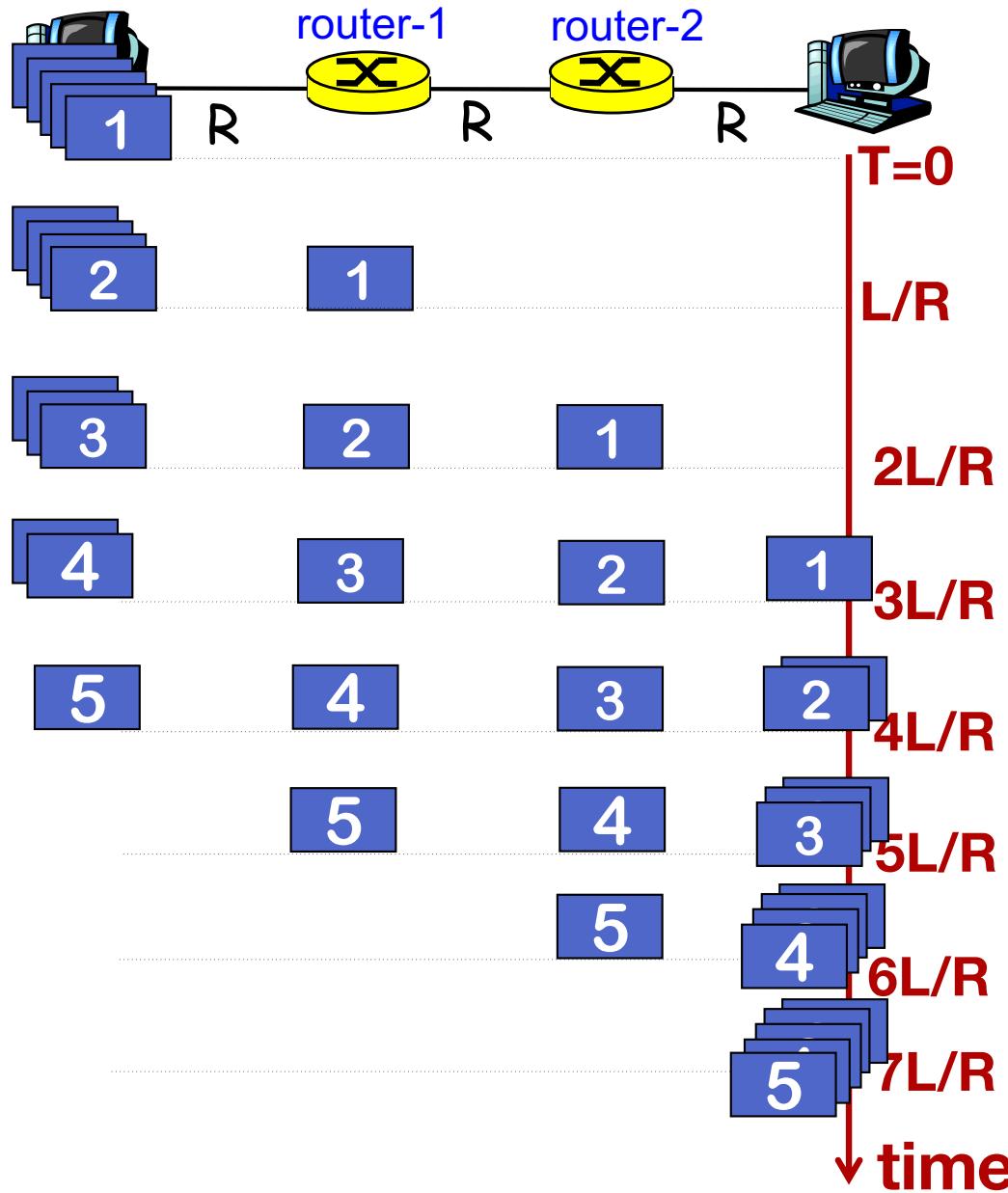


- ◆ Takes L/R seconds to transmit (push out) packet of L bits on to link of R bps
- ◆ Entire packet must arrive at router before it can be transmitted on next link: *store and forward*

Example 1: send L A \rightarrow B

- ◆ $L = 8000$ bits (1000 bytes)
- ◆ Bandwidth $R = 2$ Mbps
- ◆ *If ignore propagation delay:* i.e. when last bit of a packet left A, it arrives at router-1 instantly
$$\text{delay} = 3 \times L/R = 12 \text{ msec}$$

Packet-switching: store-and-forward

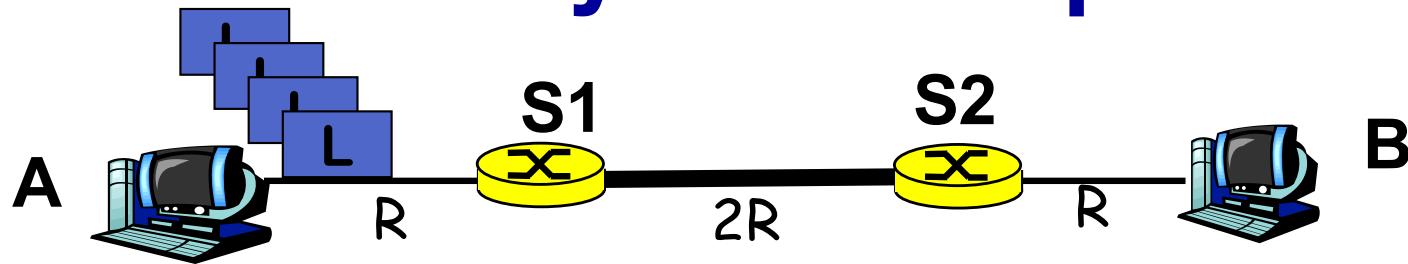


Example 2:

- ◆ A sends 5 packets to B
- ◆ $L = 8000 \text{ bits}$, $R = 2 \text{ Mbps}$
 - *Ignore propagation delay*
- ◆ How long does it take starting from A sending the first bit of first packet till B receives the last bit of the last packet?

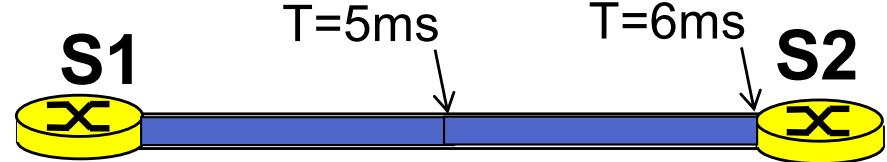
What if one takes into account the propagation delay?

Lets try an example

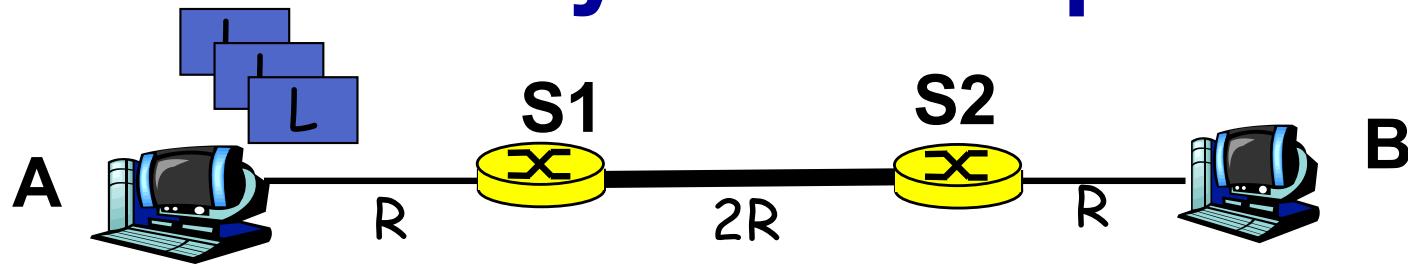


- ◆ $L = 4000 \text{ bits}$, $R = 2 \text{ Mbps}$, A sends 4 packets
 - $L/R = 2\text{msec}$, $L/2R = 1\text{ms}$
- ◆ Propagation delay: 2msec for each link
- ◆ When will first packet get to S2?
 - $A \Rightarrow S1: L/R + D_{\text{propagation}} = 2\text{ms} + 2\text{ms}$
 - $S1 \Rightarrow S2: L/2R + D_{\text{propagation}} = 1\text{ms} + 2\text{ms}$

When 1st packet arrives at S2,
where is the 2nd packet?



Lets try an example



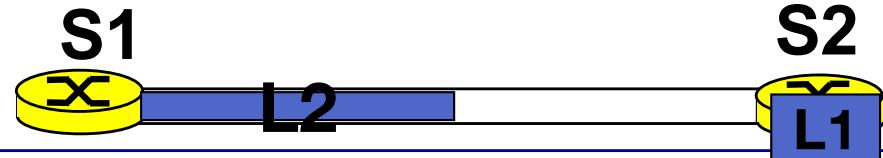
- ◆ $L/R = 2\text{msec}$, $L/2R = 1\text{ms}$
- ◆ Propagation delay: 2msec for each link
- ◆ When 1st packet arrives at S2, where exactly is the 2nd one?

→ $T=7\text{ms}$

$T=4\text{ms}$: last bit leaves A

$T=6\text{ms}$: last bit arrives S1

$T=7\text{ms}$: last bit leaves S1



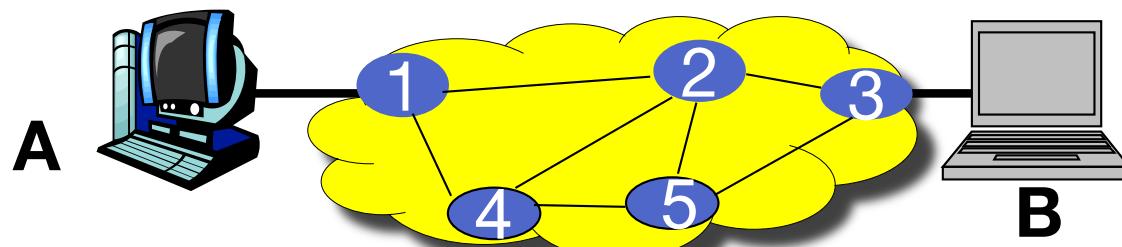
Network latency

- ◆ The time to send **1** packet from host A to B
 - sum of delays across each hop along the path

$$Delay_{A-B} = Delay_{A-1} + Delay_{1-2} + Delay_{2-3} + Delay_{3-B}$$

- ◆ RTT: round-trip-time

$$RTT_{AB} = Delay_{A-B} + Delay_{B-A}$$



What we covered today

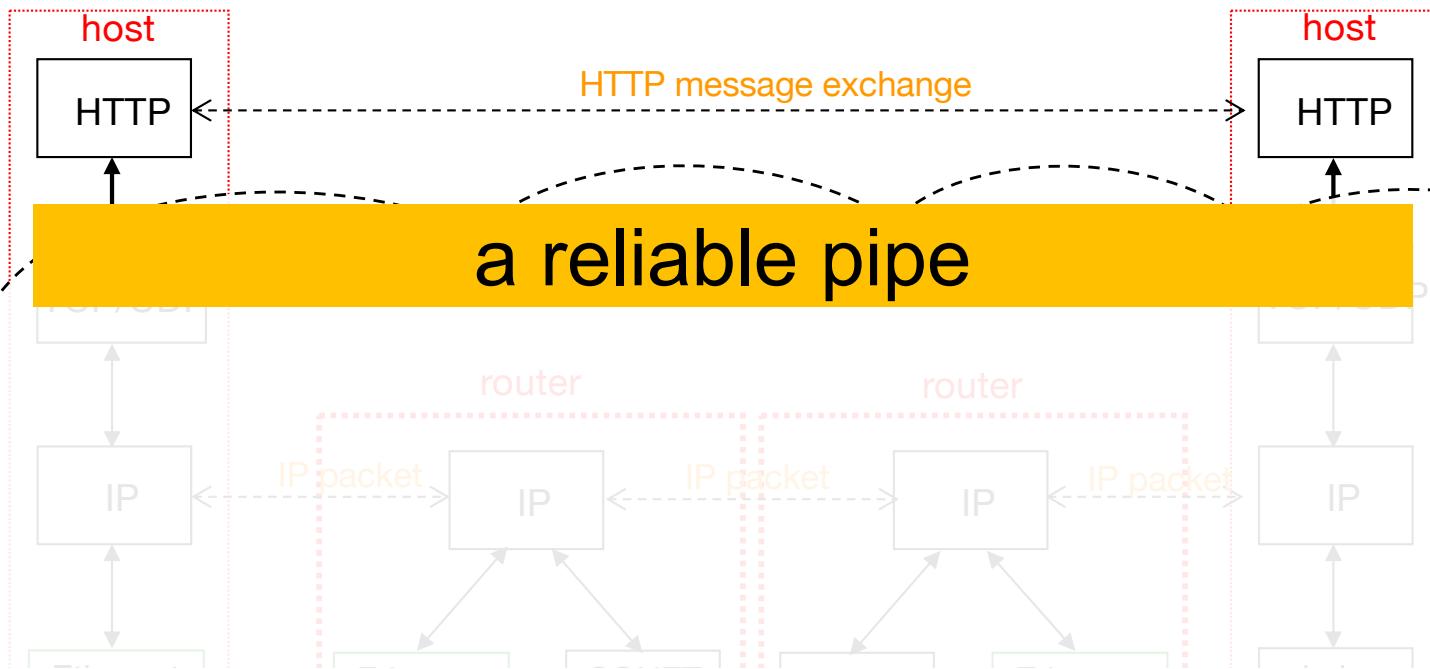
- ◆ Internet: made of a huge number of hosts, routers, wired and wireless links
- ◆ Hosts: run application protocols to exchange data packets with each other
- ◆ Routers: run bunch of protocols to move all packets towards their destinations
- ◆ Why protocols are layered
- ◆ How to calculate packet delays as they move across a packet-switched network

Lecture 1 Review: layered protocol architecture



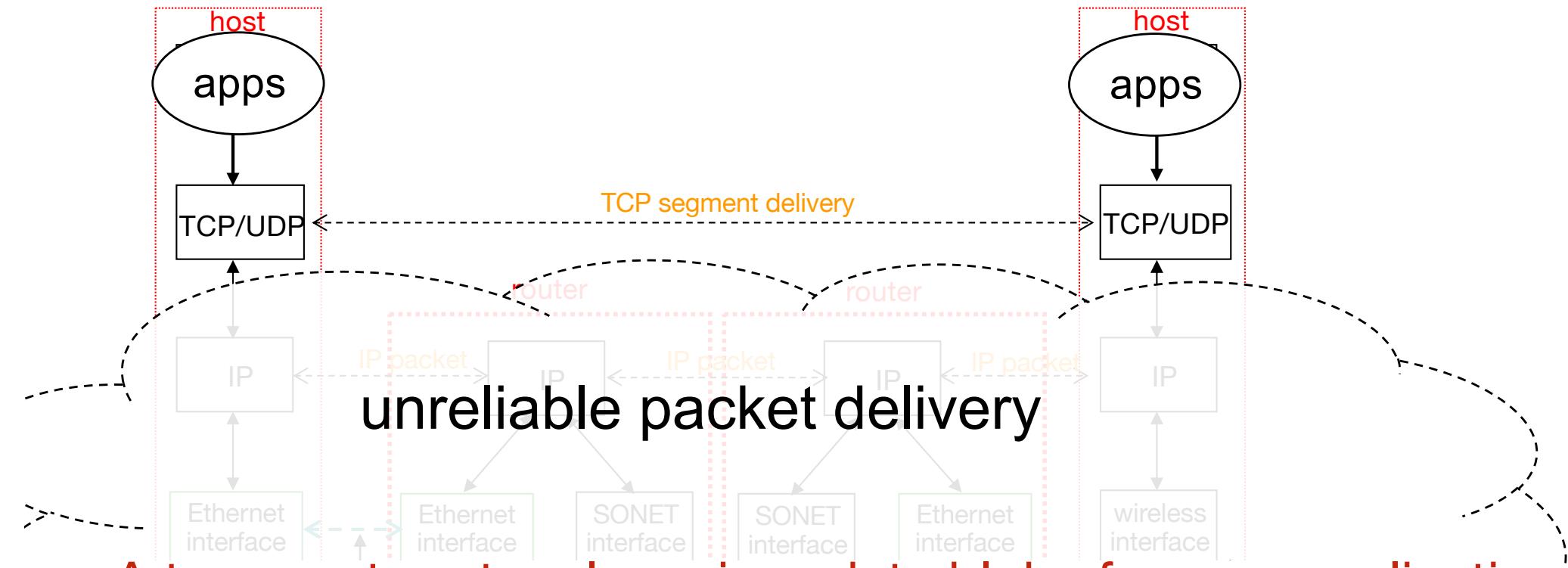
- ◆ Concepts:
 - **Internet**: made of a huge number of hosts and routers, interconnected by physical and wireless links
 - **Host**: a computer running applications and bunch of protocols to let apps exchange data with each other
 - **Router**: a packet switch running bunch of protocols to move packets toward their destinations
- ◆ Protocols are organized in layers:
 - Application protocols
 - Transport protocols
 - Network protocols
 - Link layer protocols
 - Physical layer
- ◆ How to calculate packet delays as they move across one hop

Application protocol's view of the world



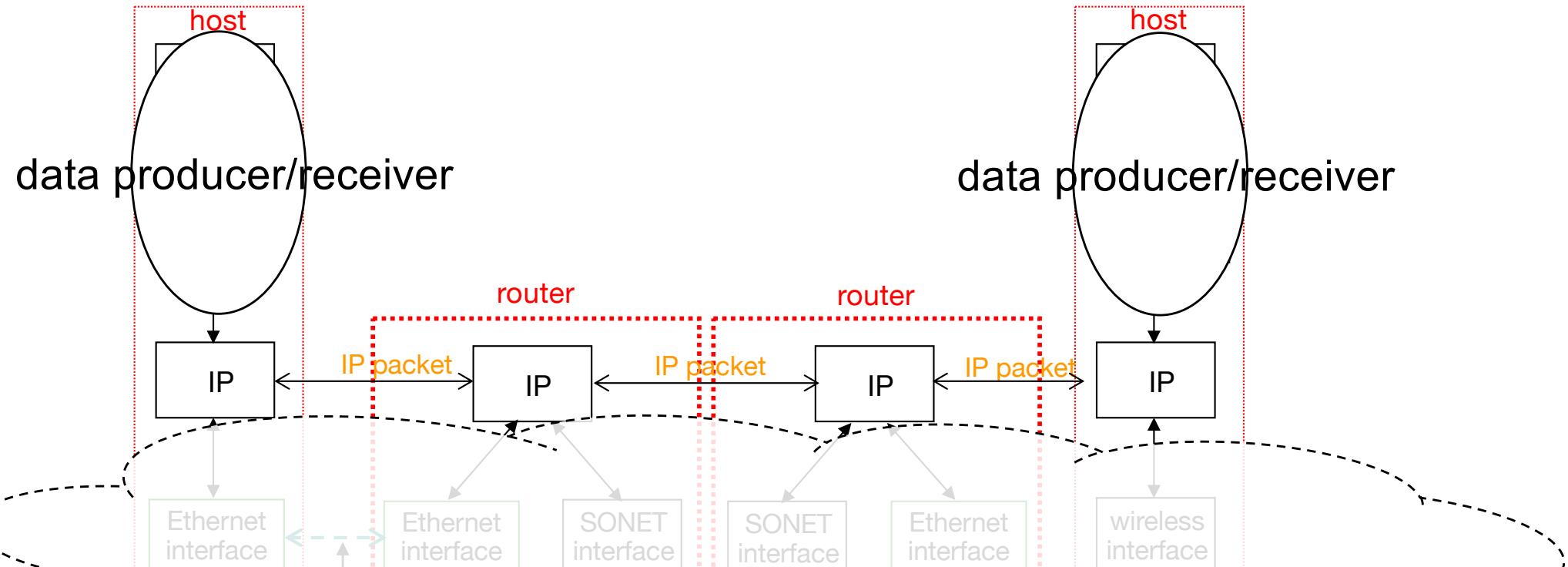
- ◆ Web browser and server exchange formatted HTTP messages over a reliable pipe
- ◆ As an application protocol, HTTP only concerns with the message's presentation format
- ◆ Application decides where msgs should be delivered to
 - The receiving end is identified by its name, which gets translated to IP address

Transport protocol's view of the world



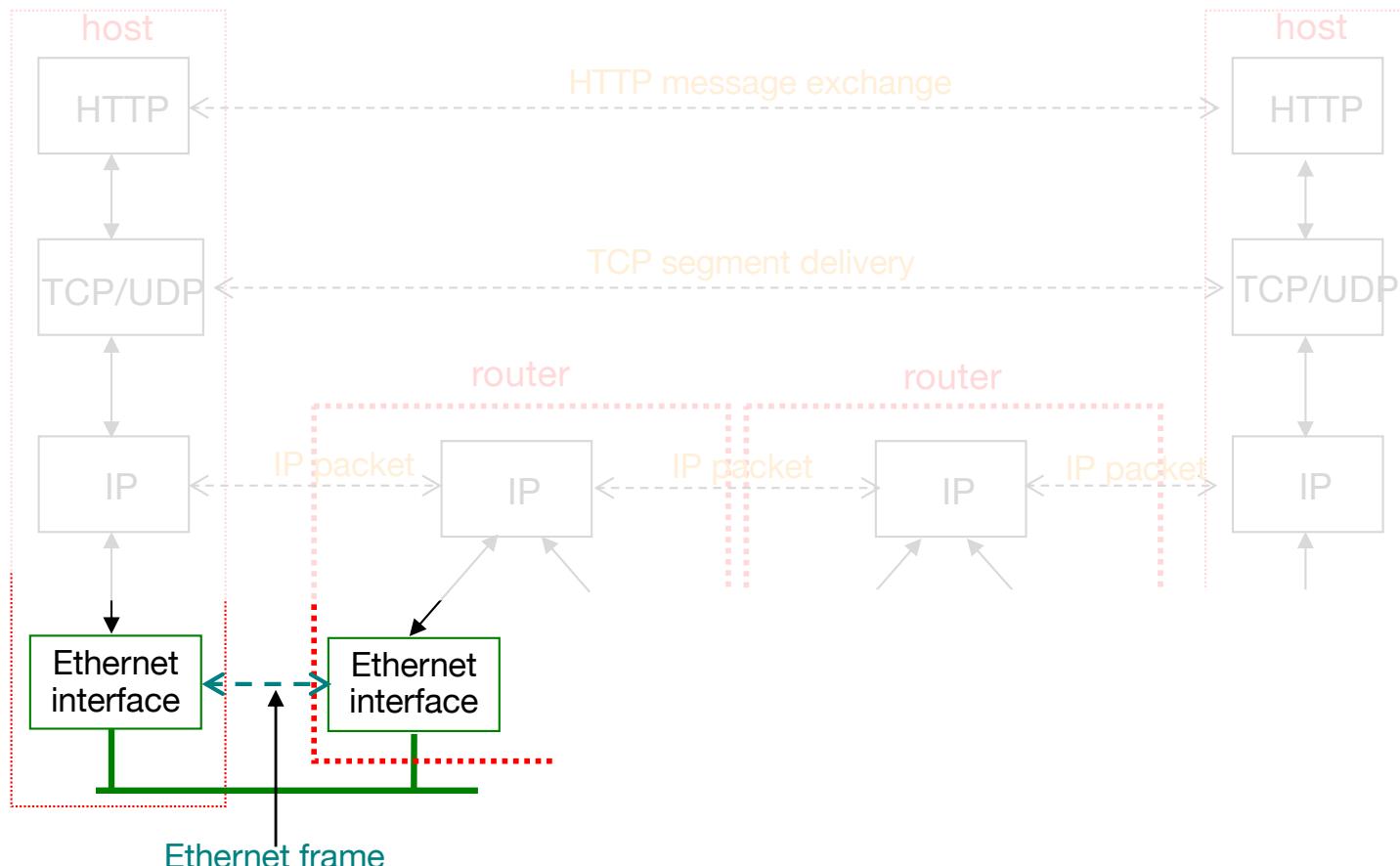
- ◆ A transport protocol receives data blobs from an application process, delivers them to the destination process (reliably)
 - Dest. Process is identified by IP address + (trans)port number
- ◆ It runs between two processes over an unreliable network (where packets can be garbled, lost, or reordered)

Network protocol's view of the world



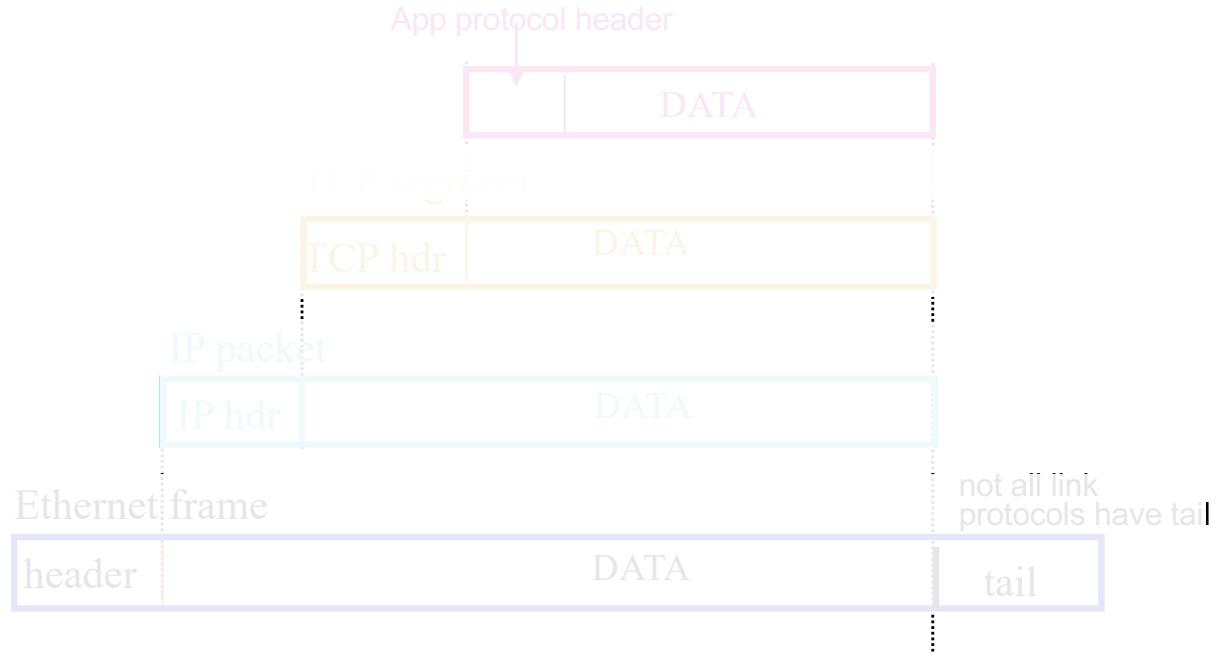
- ◆ Network protocol, IP, sees all IP-speaking nodes
- ◆ It receives data segments, delivers each of them to its destination IP address (with its best effort)
 - A router forwards packets, without looking inside IP envelope

Link layer protocol's view of the world



- ◆ A link layer protocol delivers data frames between two physically connected nodes
 - ◆ A link-layer header is added at sending node, removed by the receiving node
 - ◆ When a packet moves through the network across multiple hops: link-layer header is added and removed multiple times

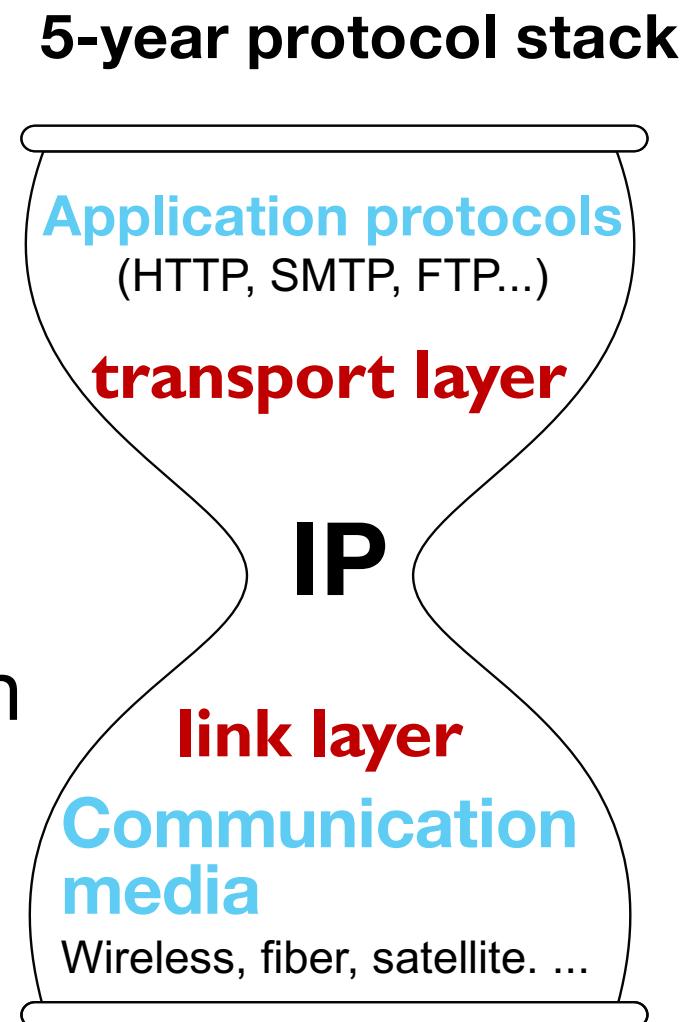
Layered protocol implementation



- ◆ protocol header: contains the information one writes on the “envelope”
- ◆ all the information, and *only* the information, that’s needed to carry out the protocol’s functionality

One more question: why 5 layers?

- ◆ Two layers are taken as given
 - Multiple different **application protocols**
 - Multiple different **physical communication media types**
- ◆ **IP**: the span layer
 - Connecting up all nodes
- ◆ **Link layer**: adaptation between IP and physical media
- ◆ **Transport**: adaptation between what apps want and what IP offers



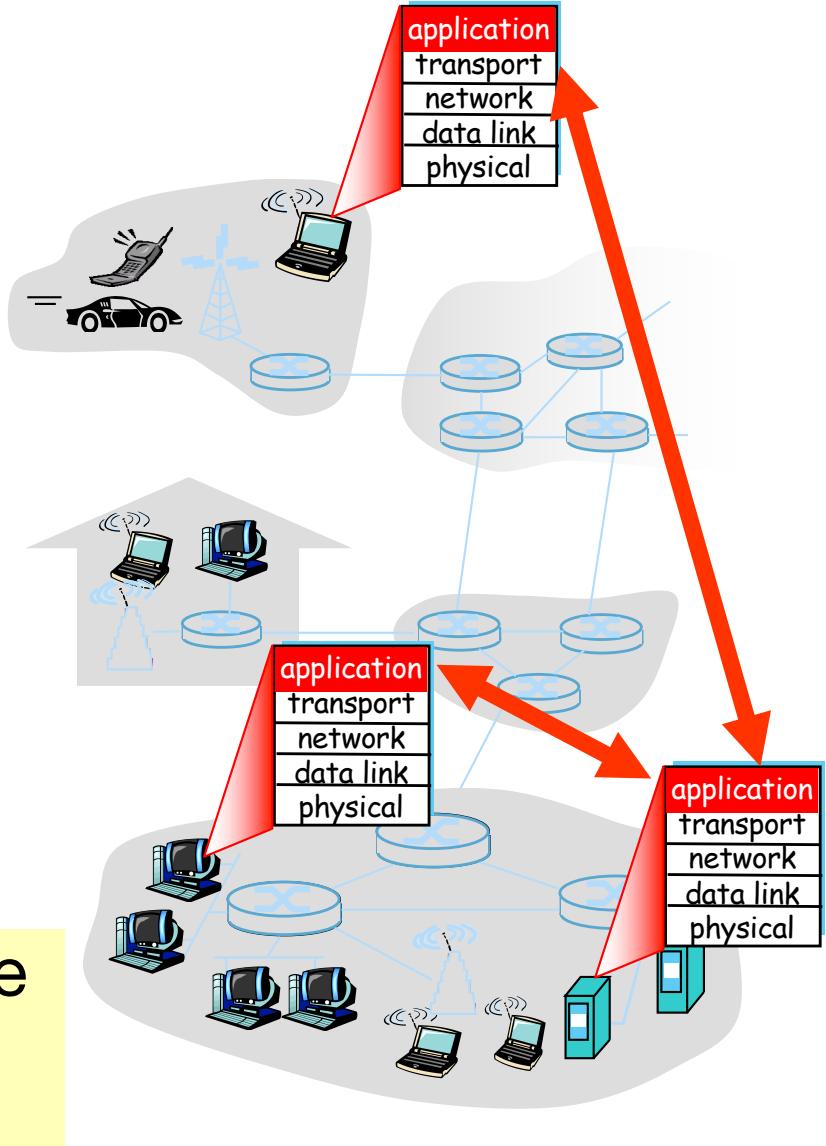
Network applications: how different parties reach each other



Some popular network applications

- ◆ Email
- ◆ Web
- ◆ WhatsApp
- ◆ BitTorrent (P2P file sharing)
- ◆ Online Games
- ◆ YouTube
- ◆ Virtual Conferencing

Application processes communicate with each other using application protocols



Client-server application communication model

servers:

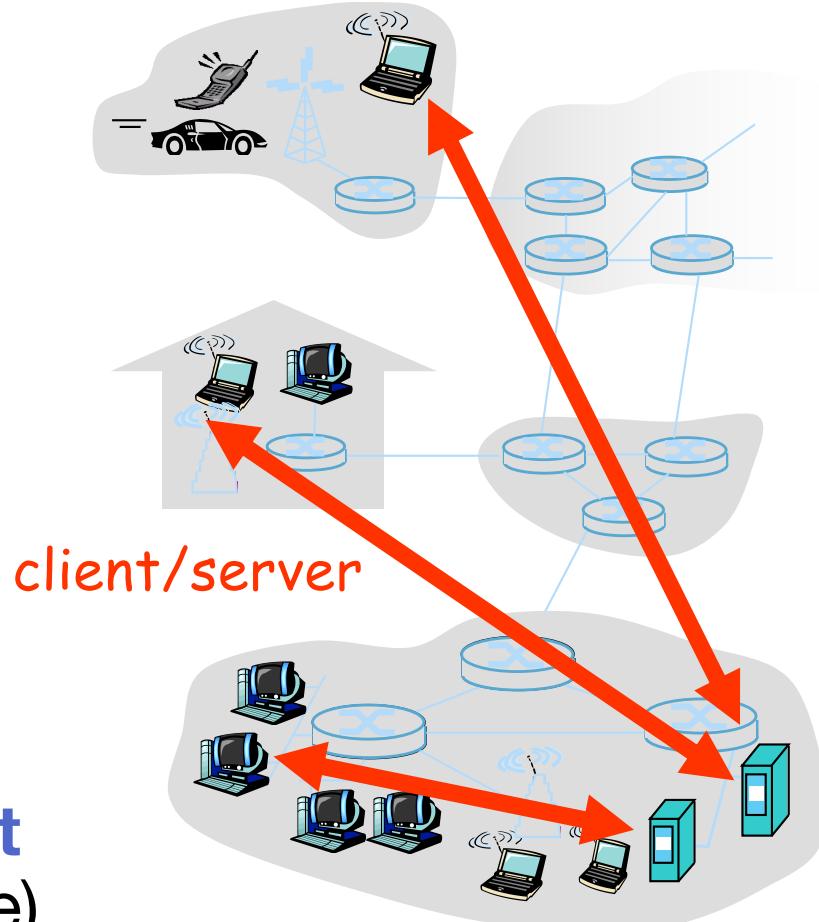
- ◆ Reachable by IP address
- ◆ **always-on**, waiting for incoming requests from clients

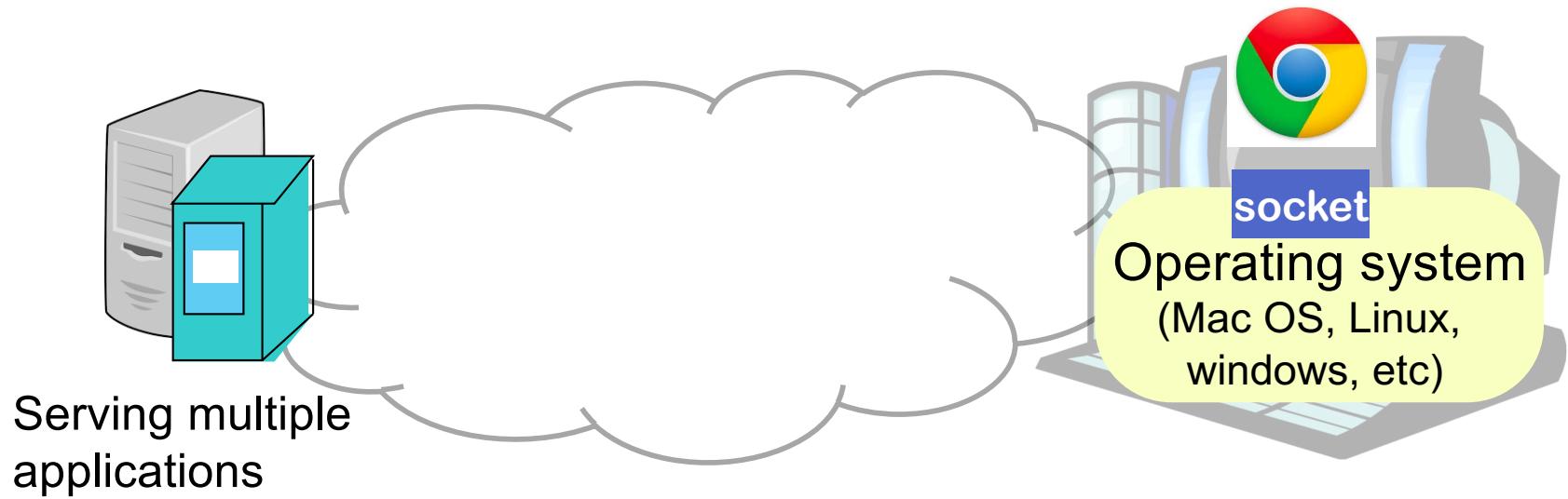
clients:

- ◆ Initiate communication with server

Q: How does a client process **identify** the server process with which it wants to communicate?

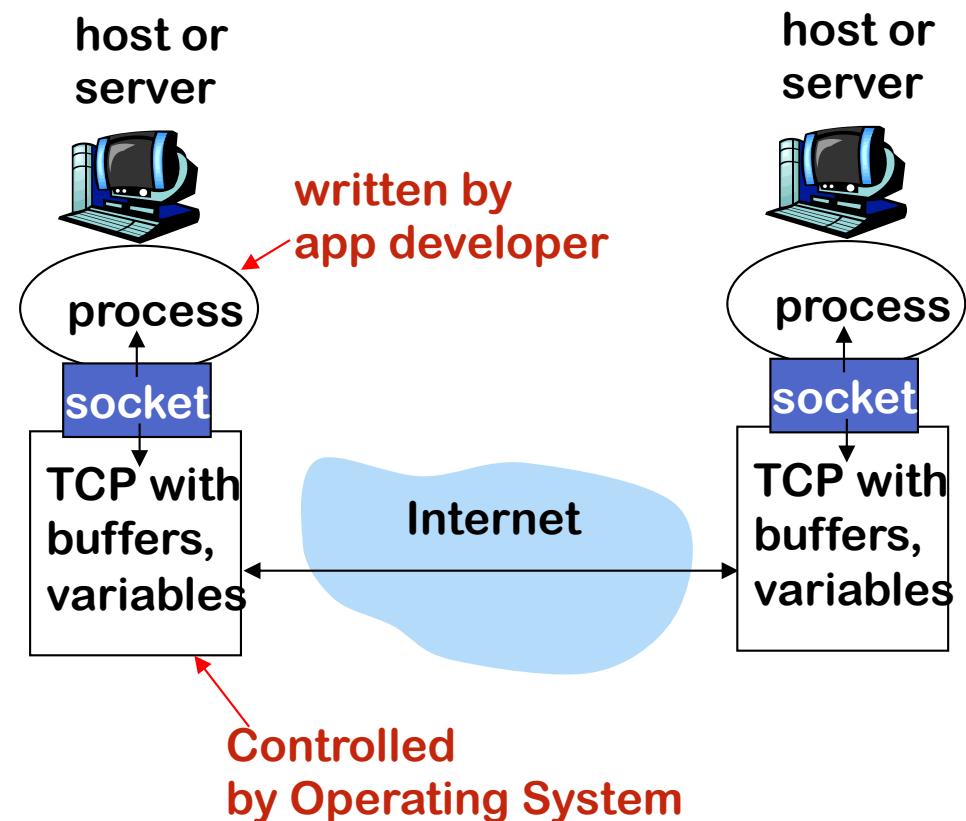
A: Using port numbers via the **socket API** (Application Program Interface)





Socket

- ◆ **Process:** program running on a host
- ◆ Between different hosts: Processes communicate through an **application-layer protocol**
- ◆ A process sends/receives messages to/from its **socket**
- ◆ A socket analogous to a door:
 - sending process shoves message out of the door
 - transport protocol brings message up to the socket at receiving process



What is "socket"

- ◆ A set of system function calls

socket(): Create a socket

bind(): bind a socket to a local IP address and port #

connect(): initiating connection to another socket

listen(): passively waiting for connections

accept(): accept a new connection

Write(): write data to a socket

Read(): read data from a socket

Close()

host or
server

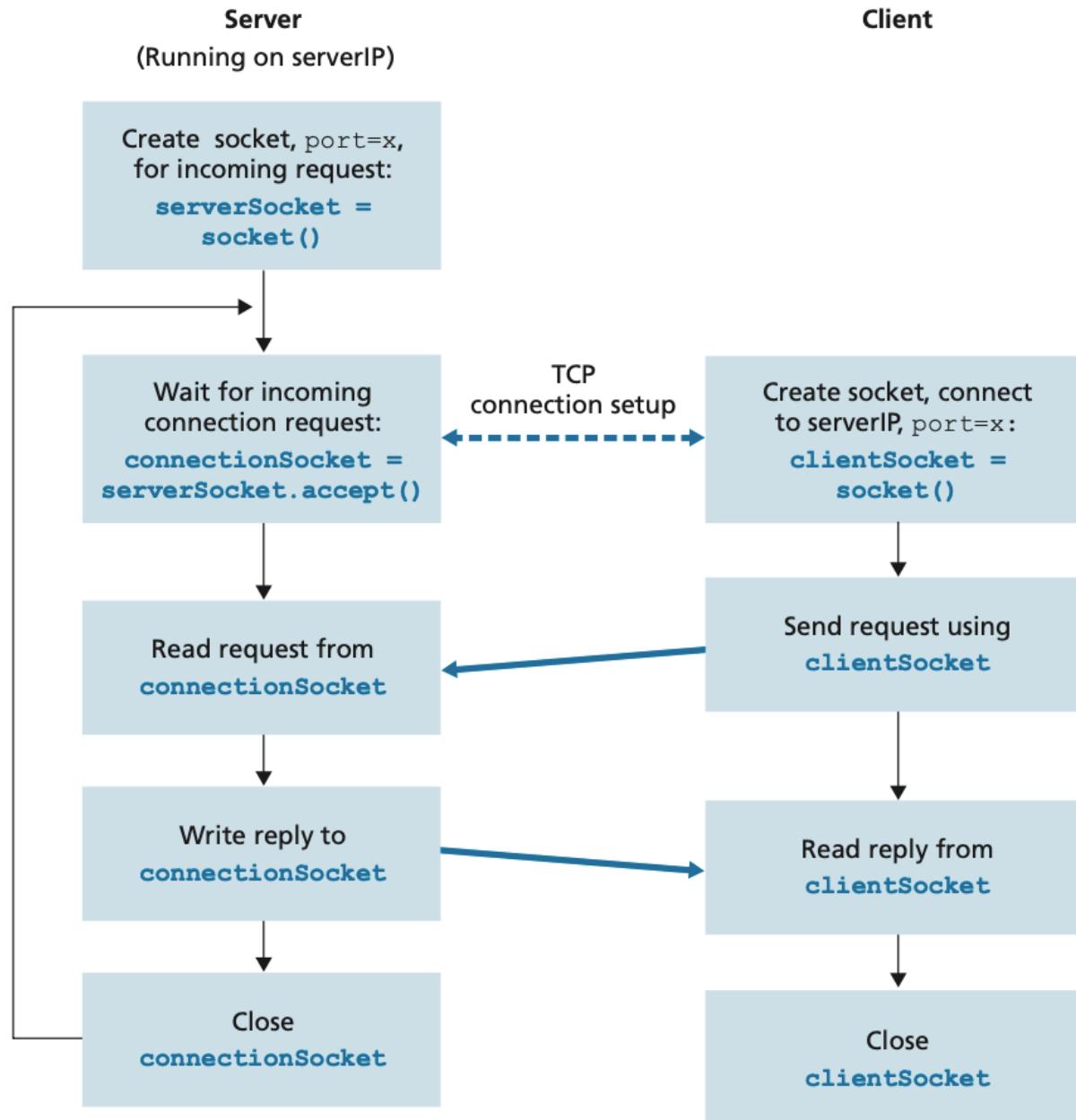


process

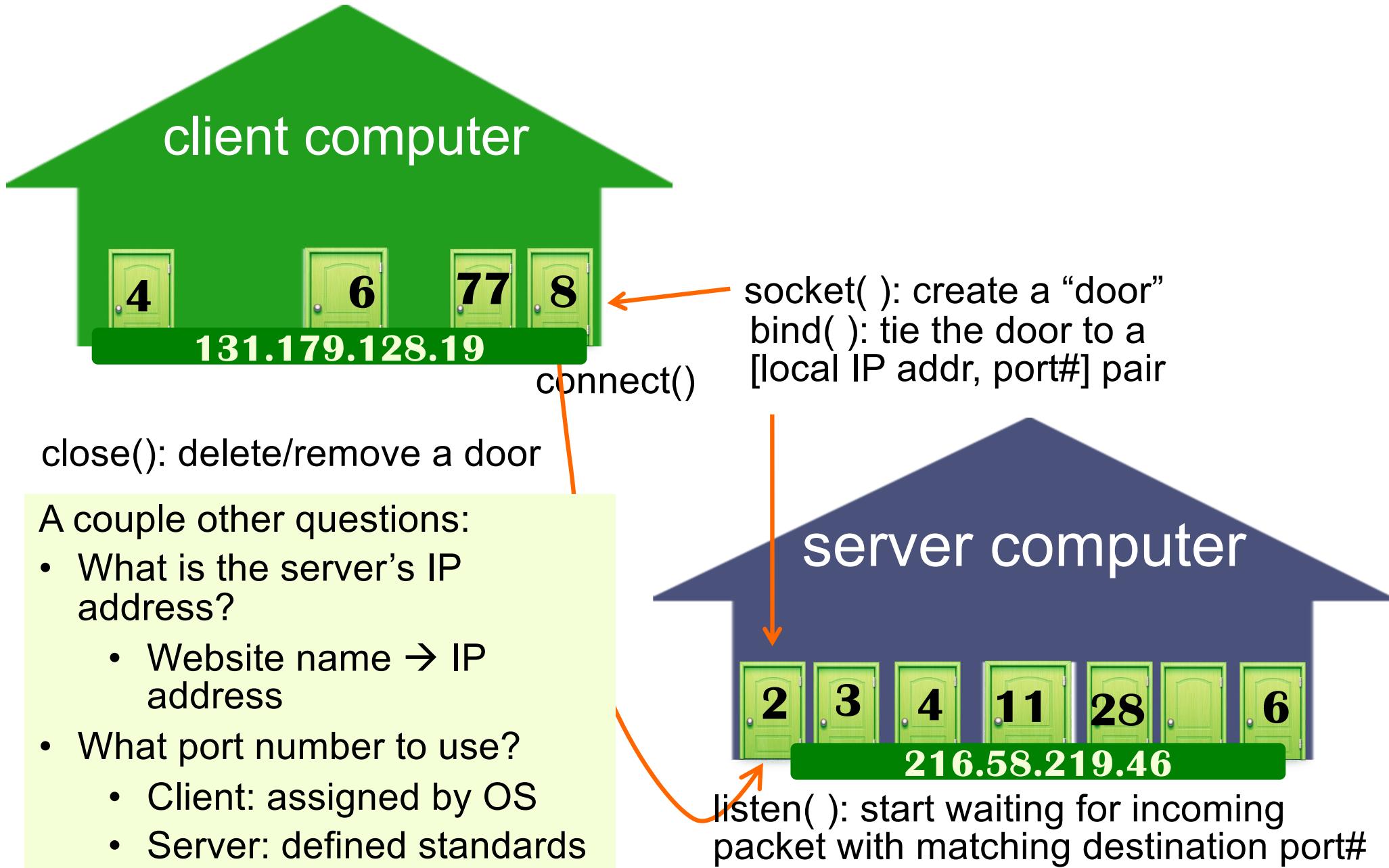
socket

TCP with
buffers,
variables

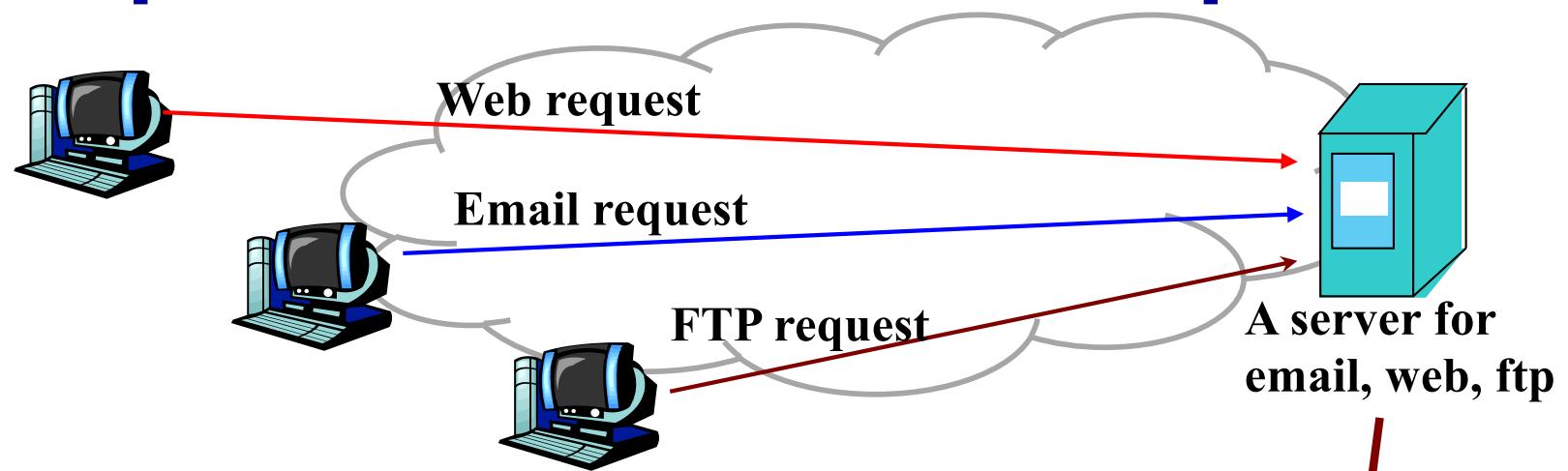
What is "socket"



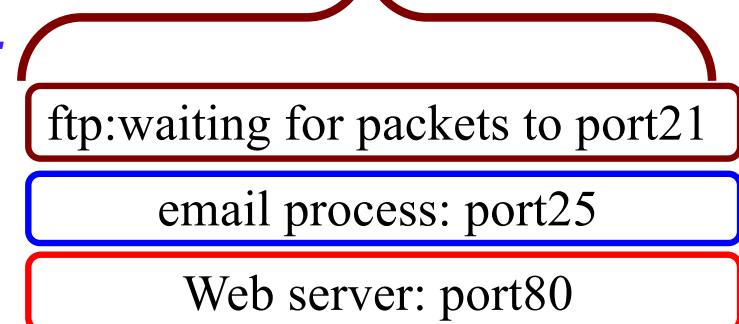
Socket: analogous to a door



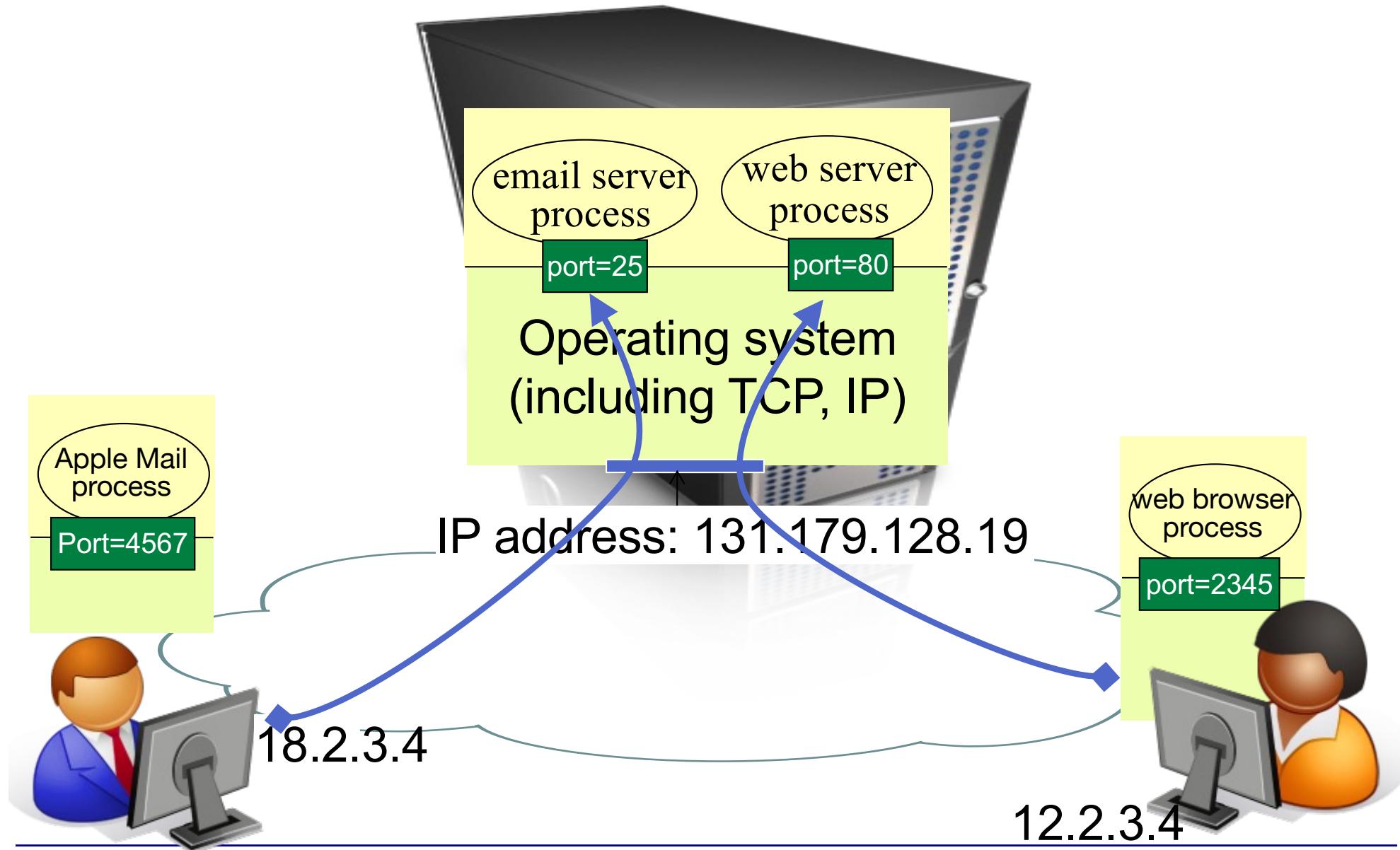
A quick comment about "port"



- ◆ Web, Email, FTP all connect client to server by TCP
- ◆ The server tells which client wants which service by *well-defined port number*:
 - Web: port 80, ftp 21, mail 25



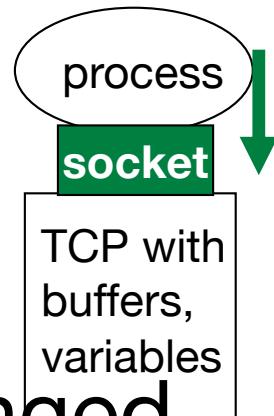
IP address, TCP connection, port number, processes, and sockets



Applications

So far we've talked

- ◆ Application process (executing application program)
- ◆ Application protocol (used by application processes to exchange data)
- ◆ Exactly how data is exchanged
 - Socket
 - Transport protocol
- ◆ Lets look at exactly what data is exchanged



Web and HTTP

- ◆ **Web page:** normally consists of
 - base **HTML-file**, which includes
 - several referenced **objects**
- ◆ An object can be another HTML file, JPEG image, Java applet, audio file,...
- ◆ Each object is addressable by a **URL** (Universal Resource Locator)



What protocol “layer” really means



application

transport

network

link

physical

Link
layer
protocol

Network
protocol

Transport
protocol

Application
protocol

Application
data

Acknowledgment

- ◆ Slides adapted from S24 CS118 instructed by Prof. Lixia Zhang