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CS 16: Solving Problems with Computers 1
      Lecture #11 PRE-RECORDED
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Searching and Sorting Algorithms

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Lecture Outline

- Searching Algorithms
 - Sequential Search
 - Binary Search

- Sorting Algorithms
 - Selection Sort
 - Bubble Sort

Note on getline() Function

- Strings can contain whitespace characters
- But when using cin for user inputs, we can't capture these!!!

We will use this function often with File I/O operations (next week)

```
    Introducing: getline() in <string> lib.
```

- Use: getline(cin, string name);
 - Will capture a user input WITH all whitespaces, except newline ('\n')
 - The newline is the default delimiter.
- To use a different delimiter (not the default newline), use it like this:

```
getline(cin, string_name, character_delimiter);
```

Examples:

```
string name;
getline(cin, name);
// if user enters:
// Hello there, friend<return>
// then name will have the value
// "Hello there, friend"
```

```
string name;
getline(cin, name, ',');
// if user enters
// Hello there, friend<return>
// then name will have the value
// "Hello there"
```

Searching and Sorting Data

...Very common algorithms in CS that get used in MANY applications...

- There are many algorithms to do this with
 - Usually there's a correlation between complexity and efficiency

Some Examples

Searching

- Sequential Search
- Binary Search
- Jump Search
- Exponential Search

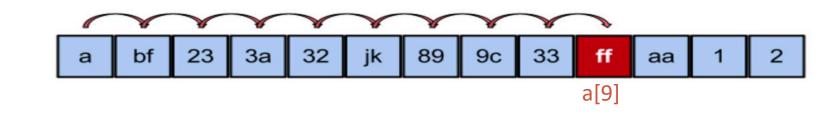
Sorting

- Selection Sort
- Bubble Sort
- Shell Sort
- Merge Sort
- Quick Sort

Sequential Search

aka Linear Search

Task: Search the array for "ff"



Result: in position 9

Sequential Search Algorithm Example

- Given: integer array (and it's size), integer target
- You want to go thru the array
 - Look for the target BUT Quit the search once you found target
- Sounds like a loop where you look thru the array
 - But it could quit at any time
 - What's better: use for or while loop??? (and why)
- While loop: keep searching as long as:
 - the target isn't found AND you haven't reached the end of the array!
 - If you find it, return the index, otherwise, return -1 (since that can't be an index!)

```
int SeqSearch
                                              See Demo File:
(int arr[], int array_size, int target)
                                              seqSearch.cpp
   int index(0);
   bool found(false); // assume not found yet!
   while ((!found) && (index < array size))</pre>
       if (arr[index] == target) {
          found = true;
       } else {
          index++;
       } // end if
   } // end while
   if (found) {
       return index;
   } else {
       return -1;
```

Sequential Search Function Example

Simple

- 1. Look for a target value inside of a given array
- 2. If you find it, return its location (i.e. index) in the array
- 3. If you don't find it, return -1

Pros and Cons of Sequential Search

PROS:

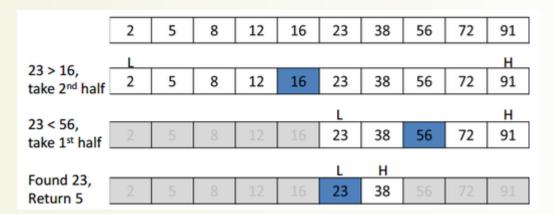
- Simple, easy
- Low complexity

CONS:

- Only finds the 1st occurrence of the target
- Slow
 - Imagine searching a huge array and the target is near the end of it!

Binary Search

- Binary search algorithm assumes the input data is already sorted.
- Example: looking the number "23" in this array:



Binary Search

- Let's call the array indexes: 0 through final_index
- Because the array is sorted, we know:

Like with SeqSearch, if the target is in the list, we want to know where it is
in the list

- Twist: Since the array is sorted, split it in half!
 - Base your initial search on how target compares to median value

Binary Search: Problem Definition

- Let's create a function that will have 5 arguments:
 - Integer array a[] and its size
 - Boolean found
 - Integers location and target
- For this example, let's make the function of type void
 - found and location will be passed-by-reference
 - So we get 2 values that caller can use
 - target only needs to be passed-by-value

Binary Search: Problem Definition

Pre- and Post-conditions for the function:

```
//precondition: a[0] through a[final_index] are
// sorted in increasing order

//postcondition: if target is not in a[0] thru a[final_index]

// found == false; otherwise found == true
```

Our algorithm:



- Start by looking at the item in the middle of the list:
 - If it is the number we are looking for, we are done!
 - If it is greater than the number we are looking for, look in the 1st half of the list
 - If it is less than the number we are looking for, look in the 2nd half of the list

1st attempt at the algorithm:

```
found = false;
mid = approx. midpoint between 0 and final index;
if (target == a[mid]) {  // Hooray! Found it!
   found = true;
   location = mid;
else if (target < a[mid])
        search a[0] through a[mid -1]
else if (target > a[mid])
        search a[mid +1] through a[final index];
```

Has to be repeated, until found or the whole array is searched

Next attempt at the algorithm:

```
found = false;
first index = 0, last index = array size - 1;
mid index = (first index + last index)/2
if (target == a[mid]) {     // Hooray! Found it!
    found = true;
    location = mid;
else if (target > a[mid])
    first_index = mid_index + 1;
```

Has to be repeated, until found or the whole array is searched

Binary Search: A Visualization 1



10	13	66	87	89	92	93	99	101	111	122	129	145
0 first	1	2	3	4	5	6 mid 13 < 93	7 found = 0	8	9	10	11	12 last
Therefore: last = mid - 1 = 5												
10	13	66	87	89	92	93	99	101	111	122	129	145
0	1	2	3	4	5	6	7	8	9	10	11	12
first		mid 13 < 66	found =	0	last							
Therefo	re: last = mi	d-1=1										
10	13	66	87	89	92	93	99	101	111	122	129	145
0 1	1	66 2	87	89	92 5	93	99	101 8	111 9	122 10	129 11	145 12
		2										
0 mid first 13 > 10	1 last	2 = 0										
0 mid first 13 > 10	1 last found	2 = 0										
0 mid first 13 > 1b Therefo	1 last found re: first = mi	2 = 0 id + 1 = 1	3	4	5	6	7	8	9	10	11	12
0 mid first 13 > 10 Therefore	1 last found re: first = m	2 = 0 id + 1 = 1	3 87 3	4 89	5 92	93	7	8	9	10	11	12

Binary Search: A Visualization 2



10	13	66	87	89	92	93	99	101	111	122	129	145
0 first	1	2	3	4	5 9	6 mid 99 > 93	7 found = 0	8	9	10	11	12 last
Therefore: first = mid + 1 = 7												
10	13	66	87	89	92	93	99	101	111	122	129	145
0	1	2	3	4	5	6	7	8	9 ♣	10	11	12
							first		mid 99 < 111	found = 0		last
Therefore: last = mid - 1 = 8												
10	13	66	87	89	92	93	99	101	111	122	129	145
0	1	2	3	4	5	6	7	8	9	10	11	12
							first mid	last 99 = 99	found = 0			

- We must ensure that our algorithm eventually ends
 - No infinite loops!
- We can terminate the loop if target is found in the array
- BUT what if target is not found in the array?
 - If first ever becomes larger than (or equal to) last, we know that there are no more indices to check and key is not in the array

See Demo File: binSearch.cpp

Final Code

```
void binSearch(int a[], int size, int target, bool &found, int &loc){
    int first = 0, last = size - 1, mid;
    found = false;
    while((first <= last) && !(found)){</pre>
        mid = (first + last)/2;
        if (target == a[mid]){
            found = true;
            loc = mid;
        else if (target < a[mid]){</pre>
            last = mid - 1;
        } else { //i.e. target > a[mid]
           first = mid + 1;
    } // while
```

Binary Search Efficiency

- The binary search algorithm is extremely fast compared to an algorithm that checks each
 item in order
- The binary search eliminates about half the elements between a[first] and a[last] from consideration at each recursive call
- For an array of 100 items, a simple serial search will average 50 comparisons and may do as many as 100!
 - N items, N max. comparisons
- For an array of 100 items, the binary search algorithm never compares more than 7 elements to the key!
 - N items, log₂N max. comparisons worst case!
 - As opposed to linear search: N items, N comparisons worst case!

END OF PART 1

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Part 2 of 2

Searching and Sorting Algorithms

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Sorting an Array

- Sorting a list of values is another very common task
 - Create an alphabetical listing
 - Create a list of values in ascending order
 - Create a list of values in descending order
- Many sorting algorithms exist
 - Some are very efficient
 - Some are easier to understand



https://www.toptal.com/developers/sorting-algorithms

Program Example: The **Selection Sort** Algorithm

 When the sort is complete, the elements of the array are ordered in ascending order, such that:

This leads to an outline of an algorithm:

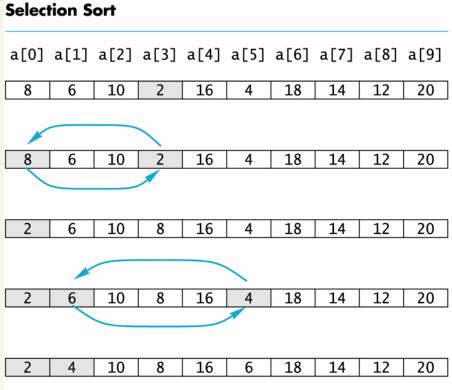
for (int index = 0; index < number_used; index++)
place the indexth smallest element in a[index]

Program Example: Selection Sort Algorithm Development

(Also, see Display 7.10 in the textbook)

- One array is sufficient to do our sorting
 - i.e. you don't really need 2 arrays
- Search for the smallest value in the array
- Place this value in a[0], and place the value that was in a[0] in the location where the smallest was found
 - i.e. swap them
- Starting at a[1], find the smallest remaining value swap it with the value currently in a[1]
- Starting at a[2], continue the process until the array is sorted

Sort from smallest to largest



Final Code

See Demo File: selectionSort.cpp

```
void swap_values(int& v1, int& v2){
                                           void sort(int a[], int number used){
                                               int index of next smallest;
    int temp;
    temp = v1;
                                               for(int index = 0; index < number used - 1; index++){</pre>
    v1 = v2;
                                                   index_of_next_smallest = index_of_smallest(a, index, number_used);
    v2 = temp;
                                                   swap_values(a[index], a[index_of_next_smallest]);
int index of smallest
(int a[], int start_index, int number_used){
    int min = a[start_index];
    int index of min = start index;
    for(int index = start index + 1; index < number used; index++){</pre>
        if (a[index] < min){</pre>
            min = a[index];
            index_of_min = index;
```

return index of min;

Bubble Sort

- Similar to Selection Sort in terms of efficiency
 - A little easier to understand

- It keeps exchanging (sorting) 2 adjacent elements at a time, if they need to be.
 - When no more exchanges are required, the list/array is sorted

Final Code

```
void bubbleSort(int array[], int size)
 int temp;
 for (int i = size-1; i >= 0; i--) {
    for (int j = 1; j <= i; j++) {
      if (array[j-1] > array[j]) {
       temp = array[j-1];
        array[j-1] = array[j];
        array[j] = temp;
     } // if
   } // for j
 } // for i
```

