

# RUNNING TIME ANALYSIS

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Problem Solving with Computers-II

C++

```
#include <iostream>
using namespace std;

int main(){
    cout<<"Hola Facebook\n";
    return 0;
}
```



# Problem: Fibonacci Numbers

## Definition:

The Fibonacci numbers are the sequence

1, 1, 2, 3, 5, 8, 13, 21, 34, 55,...

Defined by

$$F_0 = F_1 = 1$$

$$F_n = F_{n-1} + F_{n-2} \text{ for } n \geq 2$$

Problem: Given  $n$ , compute  $F_n$ .

# Which implementation is significantly faster ?

A.

```
F(int n){  
    if(n <= 1) return 1  
    return F(n-1) + F(n-2)  
}
```

B.

```
F(int n){  
    Initialize fib[0 . . . n]  
    fib[0] = fib[1] = 1  
  
    for i = 2 : n  
        fib[i] = fib[i-1] + fib[i-2]  
  
    return fib[n]  
}
```

C. *Both are almost equally fast*

The “right” question is: How does the running time grow?

E.g. How long does it take to compute  $F(200)$  recursively?

....let's say on....a supercomputer that can compute 40 trillion operations per sec

How long does it take to compute  $\text{Fib}(200)$  recursively?

....let's say on.... a supercomputer that runs 40 trillion operations per second

It will take approximately  $2^{92}$  seconds to compute  $F_{200}$ .

Time in seconds

Interpretation

$2^{10}$

17 minutes

$2^{20}$

12 days

$2^{30}$

32 years

$2^{40}$

35000 years  
(cave paintings)

$2^{50}$

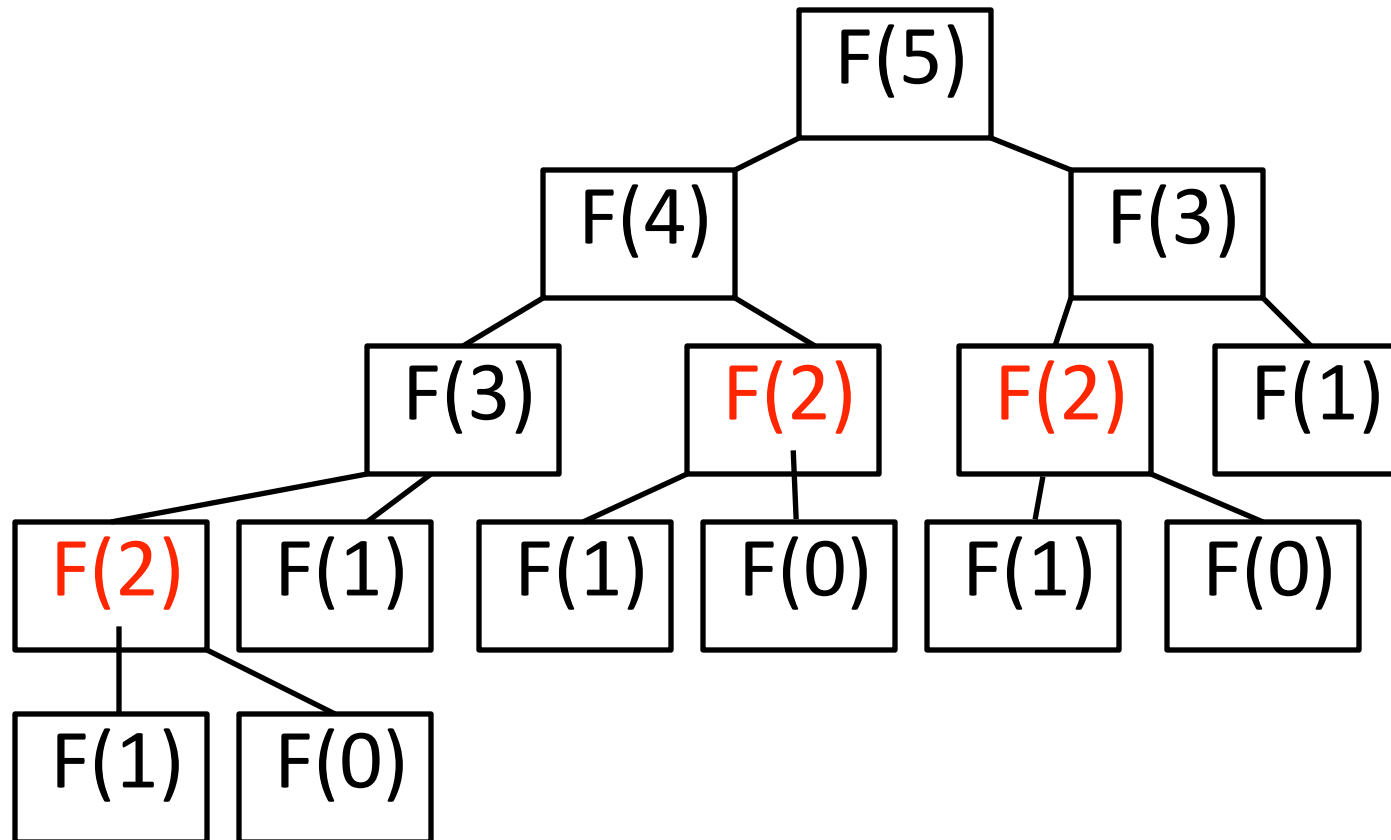
35 million years ago

$2^{70}$

Big Bang

# Why So Slow?

Too many recursive calls.



# Improved Algorithm

Lets compute  $T(n)$  = number of lines of code `Fib(n)` needs to execute.

```
F(int n){  
    Initialize fib[0 . . . n]  
    fib[0] = fib[1] = 1  
  
    for i = 2 : n  
        fib[i] = fib[i-1] + fib[i-2]  
  
    return fib[n]  
}
```

2 lines

$2(n-1)$  lines

1 line

$$T(n) = 2n + 1$$

# Question: Runtime

Is  $T(n) = 2n + 1$  an accurate description of this algorithm? A. Yes. B. No

```
F(int n){  
    Initialize fib[0 . . . n]  
    fib[0] = fib[1] = 1  
  
    for i = 2 : n  
        fib[i] = fib[i-1] + fib[i-2]  
  
    return fib[n]  
}
```

2 lines

$2(n-1)$  lines

1 line

$T(n) = 2n + 1$

# Bottom Line

What we really care about is how long it takes program to run on a real machine.

Unfortunately, this depends on:

- CPU speed
- Memory architecture
- Compiler optimizations
- Background processes

Too much to consider for every analysis



# Run Time Analysis Approach

## **Goal 1: Focus on the impact of the algorithm:**

Simplify the analysis of running time by ignoring “details” which may be an artifact of the underlying implementation

# Run Time Analysis Approach

**Goal 1: Focus on the impact of the algorithm:**

Count operations instead of absolute time!

- Every computer can do some primitive operations in constant time:
  - Data movement (assignment)
  - Control statements (branch, function call, return)
  - Arithmetic and logical operations
- By inspecting the pseudo-code, we can count the number of primitive operations executed by an algorithm

# Run Time Analysis Approach

## **Goal 1: Focus on the impact of the algorithm:**

Simplify the analysis of running time by ignoring “details” which may be an artifact of the underlying implementation

## **Goal 2: Focus on trends as input size increases (asymptotic behavior):**

How does the running time of an algorithm increase with the size of the input in the limit (for large input sizes)

# Run Time Analysis Approach

## **Goal 1: Focus on the impact of the algorithm:**

Simplify the analysis of running time by ignoring “details” which may be an artifact of the underlying implementation

Count operations instead of absolute time!

## **Goal 2: Focus on trends as input size increases:**

How does the running time of an algorithm increase with the size of the input in the limit (for large input sizes)

Describe asymptotic running time using well known functions

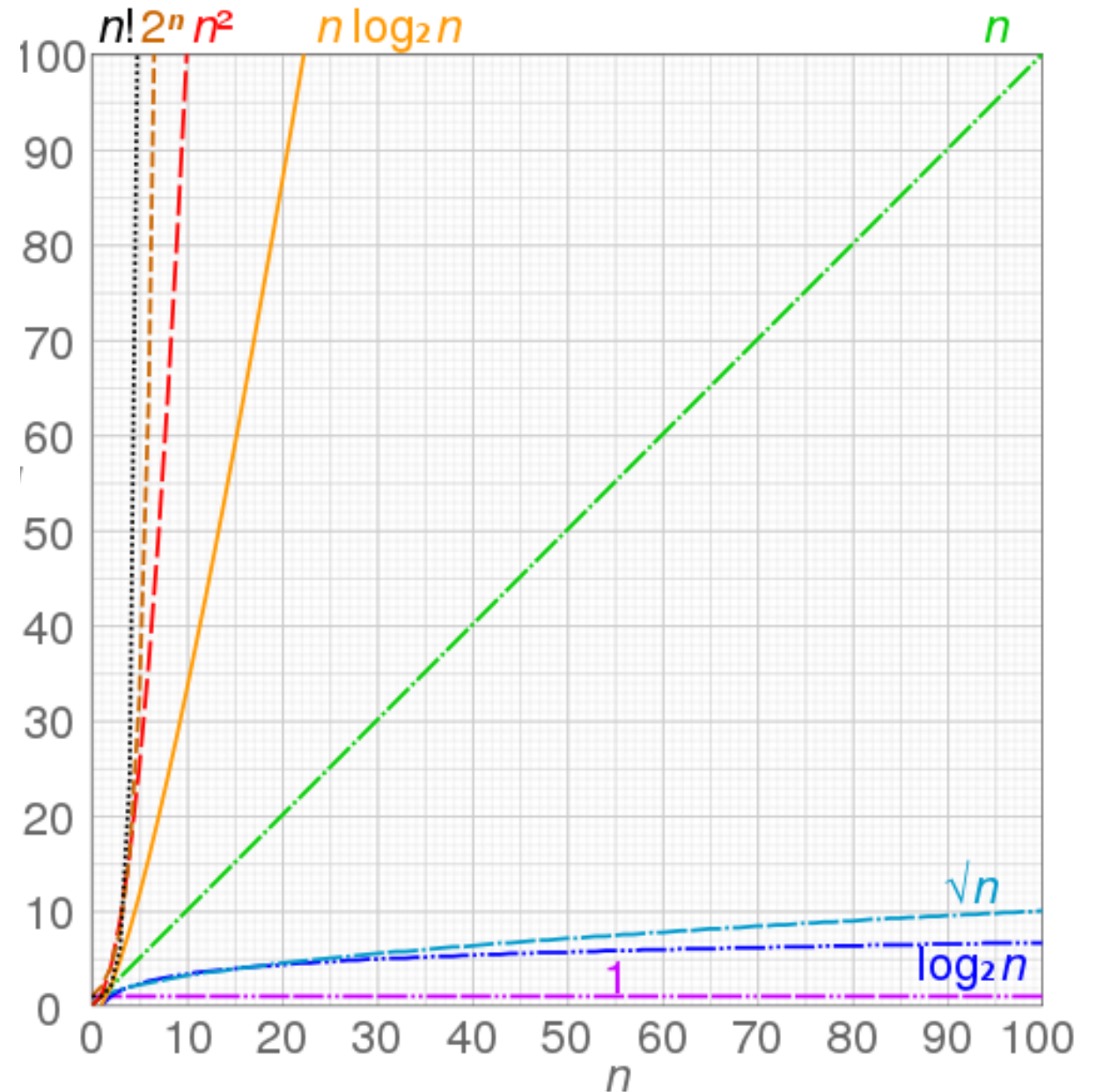
# Orders of growth

An **order of growth** is a set of functions whose asymptotic growth behavior is considered equivalent. For example,  $2n$ ,  $100n$  and  $n+1$  belong to the same order of growth

Which of the following functions has a higher order of growth?

A.  $50n$

B.  $2n^2$



# Big-O notation

- Big-O notation provides an upper bound on the order of growth of a function

# Definition of Big-O

$f(n)$  and  $g(n)$  map positive integer inputs to positive reals.

We say  $f = O(g)$  if there is a constant  $c > 0$  and  $k > 0$  such that  $f(n) \leq c \cdot g(n)$  for all  $n \geq k$ .

$f = O(g)$

means that “ $f$  grows no faster than  $g$ ”

# Express in Big-O notation

1. 100000000
2.  $3*n$
3.  $6*n-2$
4.  $15*n + 44$
5.  $50*n*\log(n)$
6.  $n^2$
7.  $n^2-6n+9$
8.  $3n^2+4*\log(n)+1000$
9.  $3^n + n^3 + \log(3*n)$

## Common sense rules

1. Multiplicative constants can be omitted:  
 $14n^2$  becomes  $n^2$ .
2.  $n^a$  dominates  $n^b$  if  $a > b$ : for instance,  $n^2$  dominates  $n$ .
3. Any exponential dominates any polynomial:  
 $3^n$  dominates  $n^5$  (it even dominates  $2^n$ ).

**For polynomials, use only leading term, ignore coefficients: linear, quadratic**



```
procedure max( $a_1, a_2, \dots, a_n$ : integers)
  max :=  $a_1$ 
  for  $i := 2$  to  $n$ 
    if max <  $a_i$ 
      max :=  $a_i$ 
  return max{max is the greatest element}
```

What is the Big-O running time of *max*?

- A.  $O(n^2)$
- B.  $O(n)$
- C.  $O(n/2)$
- D.  $O(\log n)$
- E. None of the above

# What is the Big O running time of sum()?

```
/* n is the length of the array*/  
int sum(int arr[], int n)  
{  
    int result = 0;  
    for(int i = 0; i < n; i+=2)  
        result+=arr[i];  
    return result;  
}
```

- A.  $O(n^2)$
- B.  $O(n)$
- C.  $O(n/2)$
- D.  $O(\log n)$
- E. None of the above

# What is the Big O running time of sum()?

```
/* n is the length of the array*/  
int sum(int arr[], int n)  
{  
    int result = 0;  
    for(int i = 1; i < n; i=i*2)  
        result+=2*arr[i];  
    return result;  
}
```

- A.  $O(n^2)$
- B.  $O(n)$
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- D.  $O(\log n)$
- E. None of the above

What is the Big O running time of sum()?

```
/* n is the length of the array*/
int sum(int arr[], int n)
{
    int result = 0;
    for(int i = 0; i < n; i = i+2)
        result+=arr[i];
    for(int i = 1; i < n; i =i*2)
        result+=2*arr[i];
    return result;
}
```

- A.  $O(n^2)$
- B.  $O(n)$
- C.  $O(n/2)$
- D.  $O(\log n)$
- E. None of the above

# Next time

- Running time analysis : best case and worst case
- Running time analysis of Binary Search Trees

Credits and references:

Slides by Professors Sanjoy Das Gupta and Daniel Kane at UCSD  
<http://algorithmics.lsi.upc.edu/docs/Dasgupta-Papadimitriou-Vazirani.pdf>