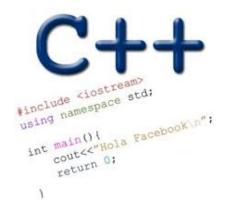
HEAPS

Problem Solving with Computers-II





How is PA02 going?

- A. Done
- B. On track to finish
- C. Having some difficulties
- D. Just started
- E. Haven't started

Heaps

- Clarification
 - heap, the data structure is not related to heap, the region of memory
- What are the operations supported?
- What are the running times?

Heaps

Min-Heaps

Max-Heap

BST

- Insert :
- Min:
- Delete Min:
- Max
- Delete Max

Choose heap if you are doing repeated insert/delete/(min OR max) operations

Applications:

- Efficient sort
- Finding the median of a sequence of numbers
- Compression codes

std::priority_queue (STL's version of heap)

A C++ priority_queue is a generic container, and can store any data type on which an ordering can be defined: for example ints, structs (Card), pointers etc.

```
#include <queue>
priority_queue<int> pq;
```

Methods:

```
*push() //insert
*pop() //delete max priority item
*top() //get max priority item
*empty() //returns true if the priority queue is empty
```

- You can extract object of highest priority in O(log N)
- To determine priority: objects in a priority queue must be comparable to each other

STL Heap implementation: Priority Queues in C++

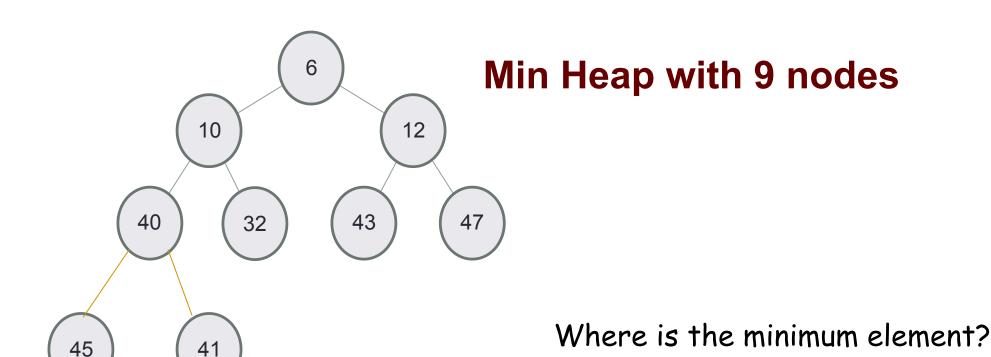
```
What is the output of this code?
```

```
priority queue<int> pq;
pq.push(10);
pq.push(2);
pq.push(80);
cout<<pre>cout<<pre>cout<</pre>
pq.pop();
cout<<pq.top();
pq.pop();
cout<<pq.top();
pq.pop();
```

```
A.10 2 80
B.2 10 80
C.80 10 2
D.80 2 10
E. None of the above
```

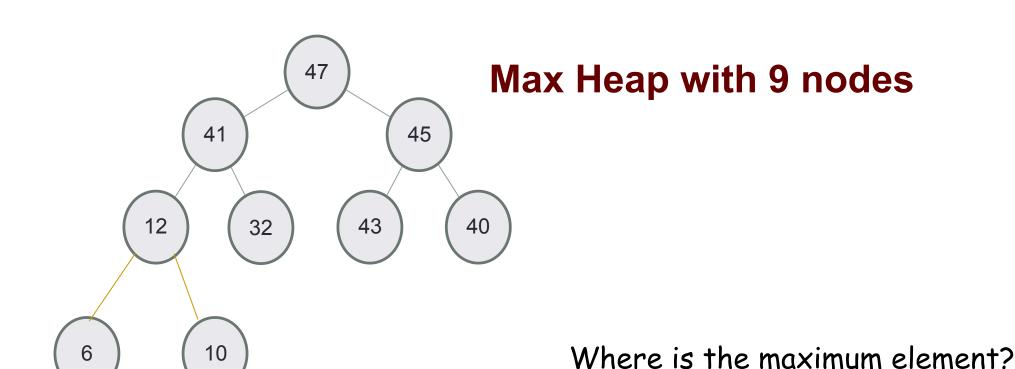
Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a min-Heap, each node satisfies the following heap property:
 key(x)<= key(children of x)



Heaps as binary trees

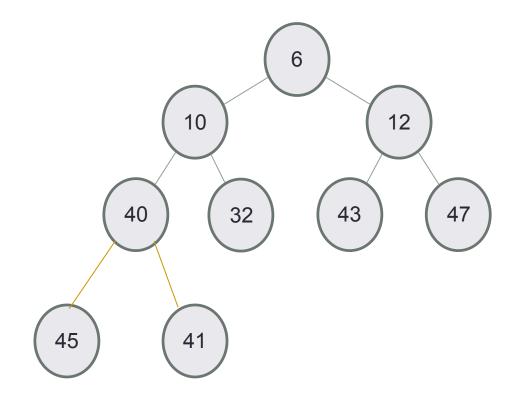
- Rooted binary tree that is as complete as possible
- In a max-Heap, each node satisfies the following heap property:
 key(x)>= key(children of x)



Identifying heaps

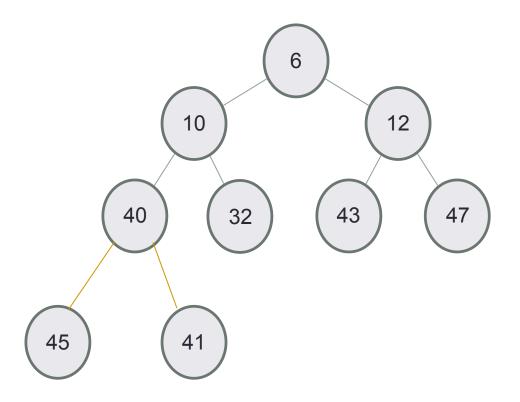
Starting with the following min Heap which of the following operations will result in something that is NOT a min Heap

- A. Swap the nodes 40 and 32
- B. Swap the nodes 32 and 43
- C. Swap the nodes 43 and 40
- D. Insert 50 as the left child of 45
- E. C&D



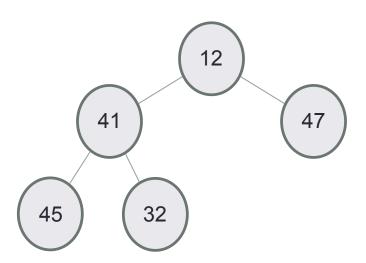
Structure: Complete binary tree

A heap is a complete binary tree: Each level is as full as possible. Nodes on the bottom level are placed as far left as possible

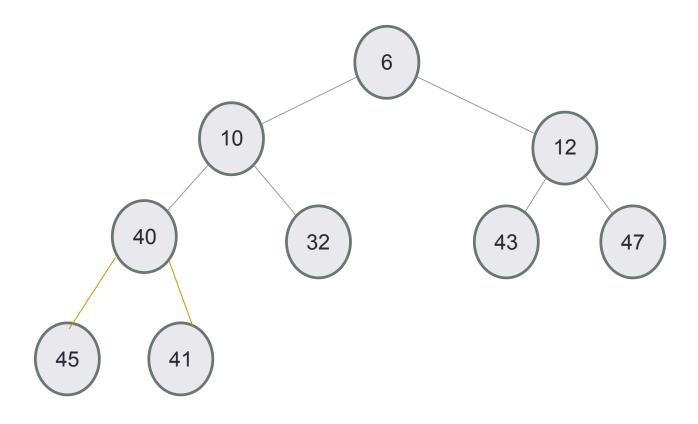


Insert 32 into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else: while(key(parent(x))>key(x)) swap the key(x) with key(parent(x))

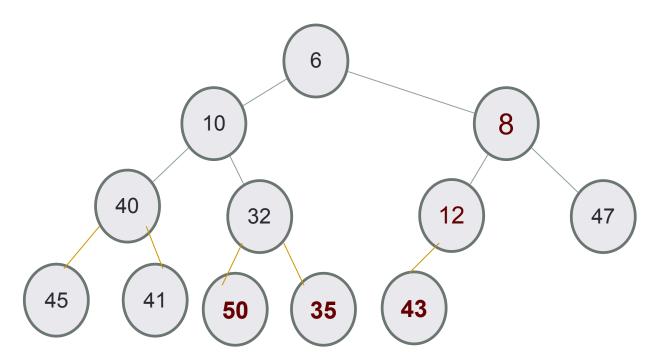


Insert 50, then 35, then 8



Delete min

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored

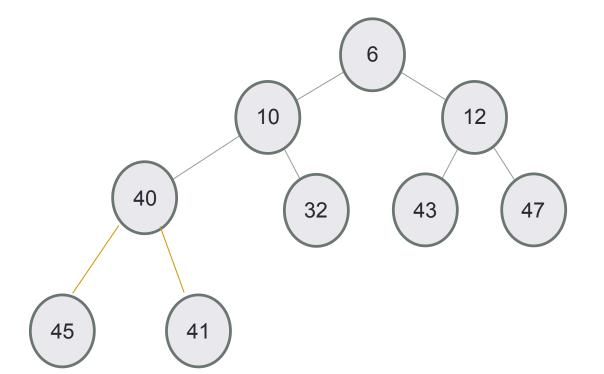


Under the hood of heaps

- An efficient way of implementing heaps is using vectors
- Although we think of heaps as trees, the entire tree can be efficiently represented as a vector!!

Implementing heaps using an array or vector

Value										
Index	0	1	2	3	4	5	6	7	8	



Using vector as the internal data structure of the heap has some advantages:

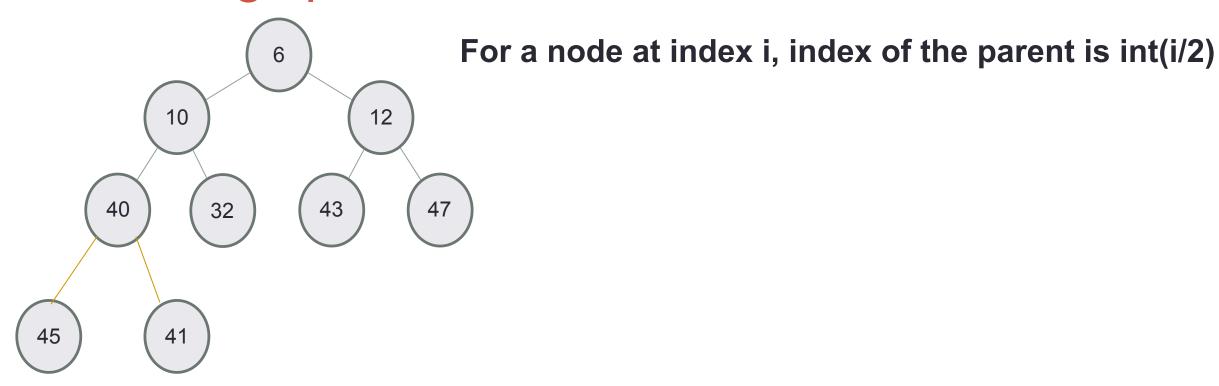
- More space efficient than trees
- Easier to insert nodes into the heap

Insert into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else....

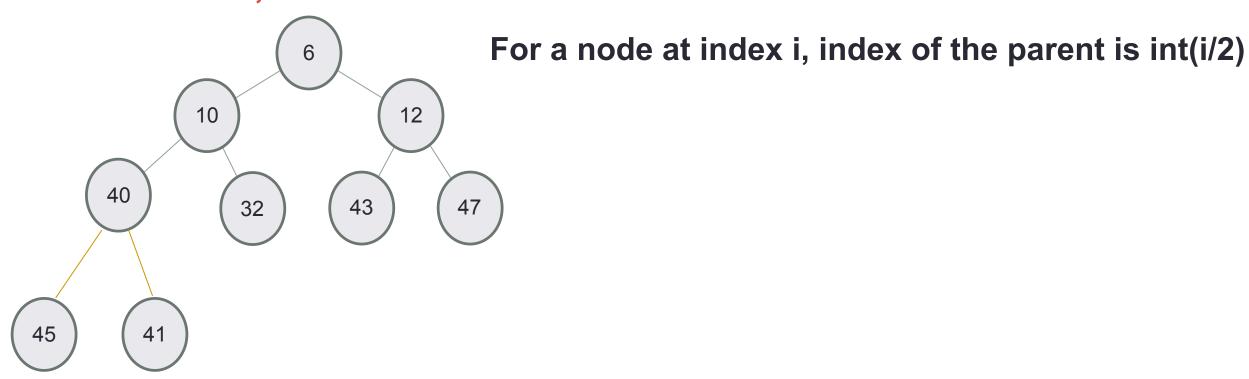
Insert the elements {12, 41, 47, 45, 32} in a min-Heap using the vector representation of the heap

Traversing up the "tree"



Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	

Insert 50, then 35



Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	

Insert 8 into a heap

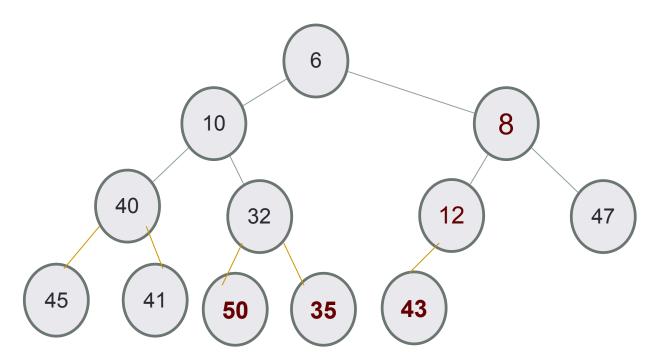
Value	6	10	12	40	32	43	47	45	41	50	35
Index	0	1	2	3	4	5	6	7	8	9	10

After inserting 8, which node is the parent of 8?

- A. Node 6
- **B. Node 12**
- **C.** None 43
- D. None Node 8 will be the root

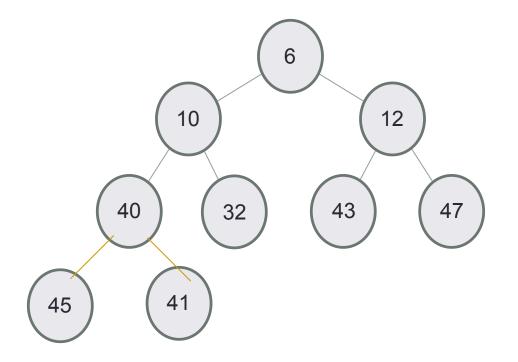
Delete min

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored



Traversing down the tree

Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	



For a node at index i, what is the index of the left and right children?

- A. (2*i, 2*i+1)
- B. (2*i+1, 2*i+2)
- C. (log(i), log(i)+1)
- D. None of the above

Next lecture

- More on STL implementation of heaps (priority queues)
- Queues