

# **More on MIPS Assembly and Memory**

**CS 64: Computer Organization and Design Logic**  
**Lecture #9**

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## AN ENGINEER HELPING A DESIGNER



# Lecture Outline

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- Midterm Exam: Sample Qs
- Shifting bits
- Stdin process
- How to implement arrays
- Intro to MIPS Calling Convention
  - Functions

# MIDTERM IS COMING!

- **Thursday, 2/15** in this classroom
- **Starts at 3:30pm **\*\*SHARP\*\*****
  - Please start arriving 5-10 minutes before class
- **I may ask you to change seats**
- Please bring your UCSB IDs with you
- **Closed book: no calculators, no phones, no computers**
- **Only the MIPS Reference Card is allowed**
- **You will write your answers on the exam sheet itself.**



# What's on the Midterm?? 1/2

- Data Representation
  - Convert bin  $\leftrightarrow$  hex  $\leftrightarrow$  decimal  $\leftrightarrow$  bin
  - Signed and unsigned binaries
- Logic and Arithmetic
  - Binary addition, subtraction
    - Carry and Overflow
  - Bitwise AND, OR, NOT, XOR
  - General rules of AND, OR, XOR, using NOR as NOT
- All demos done in class
- Lab assignments 1, 2 , 3 and 4

# What's on the Midterm?? 2/2

## Assembly

- Core components of a CPU
  - How instructions work
- Registers (\$t, \$s, \$a, \$v)
- Arithmetic in assembly (add, subtract, multiply, divide)
  - What's the difference between add, addi, addu, addui, etc...
- Conditionals and loops in assembly
- Conversion to and from Assembly and C/C++
- syscall and its various uses (printing output, taking input, ending program)
- **.data** and **.text** declarations
- Memory in MIPS
- Big Endian vs Little Endian
- R-type and I-type instructions
- Pseudo instructions

# About the Midterm Exam

- Made up of Multiple Choice & Short Answers/Coding.

## EXAMPLES:

**Complete the following MIPS assembly code that is supposed to add the number in register \$t0 to 15:**

```
li $t0, 12
```

---

- A. add \$t2, \$t1, \$t2
- B. addu \$t2, \$t1, \$t2
- C. addi \$t2, \$t0, F
- D. addi \$t2, \$t0, 0x15
- E. addui \$t2, \$t0, 0xF

**What is the 2's complement of 0x5EC?**

- A. 1x5EC
- B. 0x5EC
- C. 0xA13
- D. 0xA14
- E. 0xA15

# Sample Questions

Translate this C-style code into 4 lines of MIPS assembly code:

```
int t1 = 10, t2= 3;  
int t3= t1 + 2 * t2
```

```
li $t1, 10  
li $t2, 3  
sll $t2, $t2, 1  
add $t3, $t2, $t1
```

What is the result of these operations?

$0xF2 \ \& \ \sim(0x55)$       **0xA2**

$0x6701 \ | \ 0x1076$       **0x7777**

$0x102A99D8 \ \wedge \ 0xABA11CAB$

```
0001 0000 0010 1010 1001 1001 1101 1000  
^ 1010 1011 1010 0001 0001 1100 1010 1011  
= 1011 1011 1000 1011 1000 0101 0111 0011  
= 0xBB8B8573
```



# Sample Questions

Translate this MIPS assembly code into pseudo-code (C or C++ accepted):

```
li $s0, 2
li $s1, 6
li $t0, 2
add $s2, $s1, $s0
sll $s2, 3
mult $s2, $t0
mflo $s3
```

Translate this C code into MIPS:

```
int t0 = 3;
if (t0 < 7)
{
    t1 = 1;
    printf (t1);
}
else
    t1 = t0 + t0;
```

# srl vs sra

- srl replaces the “lost” MSBs with 0s
- sra replaces the “lost” MSBs with *either* 0s (if number is +ve) *or* 1s (if number is –ve)

## IMPLICATIONS:

***shiftDemo.asm***

- srl should NOT be used for negative numbers
- sra use for positive numbers is redundant
- When using negative numbers, use sra

# More on I/O Instructions

- Recall: to call an output, we use `syscall` which looks at what's in `$v0`
  - If `$v0 = 1`, `syscall` will print (output) an integer
  - If `$v0 = 4`, `syscall` will print (output) a string
- Recall: we can also execute inputs with the same method
  - If `$v0 = 5`, `syscall` will get (input) an integer
  - The integer is returned in `$v0`

# *StdIn*: Using the $\$v0 = 5$ syscall

- If  $\$v0 = 5$ , `syscall` will get (input) an integer
- The integer is returned in  $\$v0$

## **EXAMPLE:**

```
li $v0, 5
syscall
move $t0, $v0
...
li $v0, 1
move $a0, $t0
syscall
```

***inputDemo.asm***



# Arrays

- Question:

As far as memory is concerned, what is the *major difference* between an **array** and a **global variable**?

- Arrays contain multiple elements

- Let's take a look at:

- print\_array1.asm

- print\_array2.asm

- print\_array3.asm

# print\_array1.asm

```
int myArray[]  
    = {5, 32, 87, 95, 286, 386};  
int myArrayLength = 6;  
int x;  
  
for (x = 0; x < myArrayLength; x++)  
{  
    print(myArray[x]);  
    print("\n");  
}
```

```

# C code:
# int myArray[] =
#     {5, 32, 87, 95, 286, 386}
# int myArrayLength = 6
# for (x = 0; x < myArrayLength; x++) {
#     print(myArray[x])
#     print("\n") }

.data
newline: .asciiz "\n"
myArray: .word 5 32 87 95 286 386
myArrayLength: .word 6

.text
main:
    # t0: x
    # initialize x
    li $t0, 0
loop:
    # get myArrayLength, put result in $t2
    # $t1 = &myArrayLength
    la $t1, myArrayLength
    lw $t2, 0($t1)

    # see if x < myArrayLength
    # put result in $t3
    slt $t3, $t0, $t2
    # jump out if not true
    beq $t3, $zero, end_main

```

```

# get the base of myArray
la $t4, myArray

# figure out where in the array we need
# to read from. This is going to be the array
# address + (index << 2). The shift is a
# multiplication by four to index bytes
# as opposed to words.
# Ultimately, the result is put in $t7
sll $t5, $t0, 2
add $t6, $t5, $t4
lw $t7, 0($t6)

# print it out, with a newline
li $v0, 1
move $a0, $t7
syscall
li $v0, 4
la $a0, newline
syscall

# increment index
addi $t0, $t0, 1

# restart loop
j loop

end_main:
    # exit the program
    li $v0, 10
    syscall

```



# print\_array3.asm

```
int myArray[]  
    = {5, 32, 87, 95, 286, 386};  
int myArrayLength = 6;  
int* p;  
  
for ( p = myArray; p < myArray + myArrayLength; p++)  
{  
    print(*p);  
    print("\n");  
}
```

---

This is the end of the material of  
what is on your midterm exam

# **MIPS Function Calling Convention**

*(not on the midterm exam)*

# Functions

- Up until this point, we have not discussed **functions**
- Why not?
  - Memory management is a must for the call stack ...though we can make some progress without it
- Think of recursion...
  - How many variables are we going to need ahead of time?
  - What memory do we end up using in recursive functions?

# Implementing Functions

## **What capabilities do we need for functions?**

1. Ability to execute code elsewhere
  - Branches and jumps
2. Way to pass arguments
  - There a way (convention) to do that...
3. Way to return values
  - Registers

# Jumping to Code

- We have ways to jump to code (**j** instruction)

```
void foo() {  
    bar();  
    baz();  
}
```

```
void bar() {  
    ...  
}
```

```
void baz() {  
    ...  
}
```

- But what about jumping back?
  - We'll need a way to *save* where we were (so we can “jump” back)
- **Q:** What do need so that we can do this on MIPS?
  - **A:** A way to store the program counter (\$PC) (to tell us where the *next* instruction is so that we know *where* to return!)

# Calling Functions on MIPS

- Two crucial instructions: **jal** and **jr**
- One specialized register: **\$ra**
- **jal** (**jump-and-link**)
  - Simultaneously **jump to an address**, and **store the location of the next instruction** in register **\$ra**
- **jr** (**jump-register**)
  - **Jump to the address stored in a register**, often **\$ra**

# Simple Call Example

- See program: **simple\_call.asm**

**# Calls a function (test) which immediately returns**

**.text**

**test:**

**# return to whoever made the call**

**jr \$ra**

**main:**

**# call test**

**jal test**

**# exit**

**li \$v0, 10**

**syscall**





# Passing and Returning Values

- We want to be able to call arbitrary functions without knowing the implementation details
- How might we achieve this?
  - Designate specific registers for **arguments** and **return values**



# Passing and Returning Values in MIPS

- Registers **\$a0** thru **\$a3**
  - **Argument registers**, for passing function arguments
- Registers **\$v0** and **\$v1**
  - **Return registers**, for passing return values

# Passing and Returning Values in MIPS

## Demos

- ***print\_int.asm***
  - Illustrates the use of a printing sub-routine  
(i.e. like a simple function)
- ***add\_ints.asm***
  - Illustrates the use of an adding sub-routine  
(i.e. like a simple function that returns a value)

# YOUR TO-DOs

- Assignment/Lab #5 is next week
- **No lab this week**
- **Study for your midterm exam!! 😊**
- **IMPORTANT:**  
Read the ***MIPS Calling Convention*** handout for next week (after the exam)
  - Handout is on the class website

**</LECTURE>**