

PA4  
PA1 Late / Resubmit > due tonight  
PAS released → due Tuesday

## Sorting Quickly

$O(1)$   $\left\{ \begin{array}{l} \text{public class SortQuickly} \{ \\ \text{public static void swap(String[] array, int i1, int i2) } \\ \text{String temp = array[i1];} \\ \text{array[i1] = array[i2];} \\ \text{array[i2] = temp;} \\ \} \end{array} \right.$

$O(N)$   $\left\{ \begin{array}{l} \text{public static int partition(String[] array, int low, int high) } \\ \text{int pivotStartIndex = high - 1;} \\ \text{String pivot = array[pivotStartIndex];} \\ \text{int smallerBefore = low, largerAfter = high - 2;} \\ \text{while (smallerBefore <= largerAfter) } \\ \text{\{ if (array[smallerBefore].compareTo(pivot) < 0) } \\ \text{\{ smallerBefore += 1;} \\ \text{\} } \\ \text{\{ else } \\ \text{\{ swap(array, smallerBefore, largerAfter);} \\ \text{\{ largerAfter -- 1;} \\ \text{\} } \\ \text{\} } \\ \text{swap(array, smallerBefore, pivotStartIndex);} \\ \text{return smallerBefore;} \\ \} \end{array} \right.$

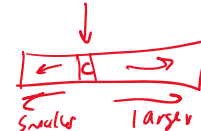
$\rightarrow$  ignored

$\left\{ \begin{array}{l} \text{public static void qsort(String[] array, int low, int high) } \\ \text{if (high - low <= 1) } \{ \text{return;} \} \\ \text{int splitAt = partition(array, low, high);} \\ \text{qsort(array, low, splitAt);} \\ \text{qsort(array, splitAt + 1, high);} \\ \} \end{array} \right.$

$\left\{ \begin{array}{l} \text{public static void sortD(String[] array) } \\ \text{qsort(array, 0, array.length);} \\ \} \end{array} \right.$

$\left\{ \begin{array}{l} \text{public static void main(String[] args) } \\ \text{String[] str = {"f", "b", "a", "e", "d", "c"};} \\ \text{int[] result = SortQuickly.sortD(str);} \\ \text{System.out.println(Arrays.deepToString(result));} \\ \} \end{array} \right.$

$["a", "b", "c", "d", "e", "f"]$



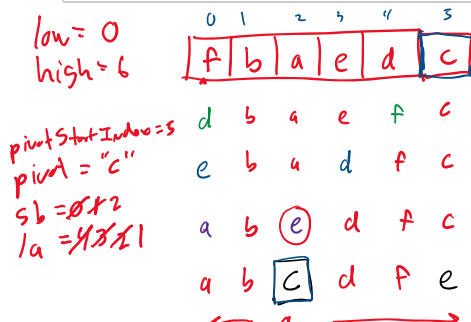
Draw the picture of sortD()

What is the tight bound of sortD:

Best case: median value  $O(N \log_2 N)$   
Worst case: sorted array  $O(N^2)$

Name: \_\_\_\_\_ PID: \_\_\_\_\_ Code: 2970

Size = 6



$N = 6$   
height = 3  
 $\rightarrow \log_2(N)$

1a = 1111

a b c d f e

← 2 →

return 2 → splitAt

0 1 2  
a b  
return 1

3 4 5 6  
d f e

N-1

a

d e f

3 4 5 6  
d f

Partition rules

← p v →  
return index

log<sub>2</sub>(n)

$\Theta(N * \log_2(n))$

worst case → sorted

1 2 3 4 5  
1 2 3 4  
1 2 3  
1 2  
1

$N + N \rightarrow \Theta(n^2)$