CSE12 - Lecture 11

Wednesday, April 26, 2023 8:00 AM

PA3 -> due today @ 16pm PAY > released today - due Tuesday Exam 1 = Friday
La BFS/DFS and before

Measuring Runtime

Count how many times each line executes, then say which $\Theta($ statement(s) is(are) true.

```
int maxDifference(int[] arr){
  max = 0;
                                            1 + (N+1)+N
   for (int i=0; i<arr.length; i++) {</pre>
     for (int j=0; j<arr.length; j++) {
   if (arr[i] - arr[j] > max)
                                           N(1+(N+1)+~
          max = arr[i] - arr[j];
                        3+2n+N(2+4n)
```

Assume n = arr.length

A.
$$f(n) = \theta(2^n)$$

 $f(n) = \theta(n^2)$
 $f(n) = \theta(n)$
 $f(n) = \theta(n)$
 $f(n) = \theta(n)$

D. $f(n) = \theta(n^3)$ Other/none/more

$$8^{N^2+3} \qquad g(N) = N^2$$

$$C = 8$$

$$N_0 = 3$$

$$Q(N^2)$$

Count how many times each line executes, then say which $\Theta()$ statement(s) is(are) true.

```
int sumTheMiddle(int[] arr){
                                  int range = 100;
                                    int start = arr.length/2 - range/2;
                              int sum = 0;
for (int i=start; i<start+range; i++) | + | o| + | co| = 
                                                                        sum += arr[i];
                                  return max;
```

Assume n = arr.length

A.
$$f(n) = \theta(2^n)$$

B. $f(n) = \theta(n^2)$
C. $f(n) = \theta(n)$
 $f(n) = \theta(1)$
E. None of these

Big O <u>Upper</u> bound f(n) = O(g(n)), f(n) <= c * g(n)

Big Somega Ower bound $f(n) = \frac{1}{2}(g(n)), f(n) >= c * g(n)$ for all $n \ge n0$

Big 0 theta <u>tight</u> bound f(n) = 0 (g(n)), f(n) = c * g(n)for all $n \ge n0$

For each function in the list below, it is related to the function below it by O, and the reverse is not true. That is, n is $O(n^2)$ but n^2 is not O(n).

 $f(n) = 1/(n^2)$

for all $n \ge n0$

- f(n) = 1/n
- f(n) = 1
- f(n) = log(n)
- N= 1000 f(n) = sqrt(n)1000 000
- f(n) = n
- $f(n) = n^2$
- $f(n) = n^3$
- 12 (0010 $f(n) = n^4$ 100000000
- ... and so on for constant polynomials ...
- $f(n) = 2^{n}$
- f(n) = n!

Stat = 2 -50

start + range = 1 +50

_ PID: _____ Code: <u>40</u>20

```
void printAllItemsTwice(int arr[], int size)
        C=6
N,=4 Q(N)
      What is the tight bound?
         (\mathcal{N}(\mathcal{N}))
      void printFirstItemThenFirstHalfThenSayHil00Times(int arr[], int size)
         printf("First element of array = %d\n",arr[0]);
      for (int i = 0; i < size/2; i++) { 1+(\frac{1}{2}+1)+\frac{3}{2}+2 }

printf("%d\n", arr[i]);
         f(N) = 3N + 304
      What is the tight bound?
           (i)(N)
      void printAllNumbersThenAllPairSums(int arr[], int size)
         for (int i = 0; i < size; i++) {
                                         1+(M1)+N 2N+2 +
N(1+(M1)+N) 3N2 +2N
                                            P(N)= 3~2)+7~+4
いては、What is the tight bound?
```

Selection Sort

```
import java.util.Arrays;
public class Sort {
  public static void sortA(int[] arr) {
     for(int i = 0; i < arr.length; i += 1) {
       System.out.print(Arrays.toString(arr) + " -> ");
       int minIndex = i;
        for(int j = i; j < arr.length; j += 1) {
           if(arr[minIndex] > arr[j]) { minIndex = j; }
 int temp = arr[i];
arr[i] = arr[minIndex];
arr[minIndex] = temp;
       System.out.println(Arrays.toString(arr));
Selection Sort - what does it print out?
Sort.sortA(new int[]{ 53, 83, 15, 45, 49 });
[53, 83, 15, 45, 49] -> 15 8 7 53
                      15 45 49 53 83)
     worse case: reverse sorted away
87 57 47 45 15
       best case:
                                      sorted array
15 49 49 59 83
```

What is the runtime? Consider the shape of the input array.

Worse case: $\Theta(\nu^{\nu})$ Best case: $\Theta(\nu^{\nu}) \longrightarrow ig Sor + eU(ig)$ $\Theta(\nu)$

Insertion Sort

Insertion Sort - what does it print out?

What is the runtime? Consider the shape of the input array.

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